

ICT 409

13 Dec 2020

Data Representation

Outline

- ❑ Floating Point Number Representation
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- ❑ Fixed Point Number Representation
- ❑ Data Representation
 - Basic Formats
 - Word lengths
 - Storage Order
 - Big-endian
 - Little-endian

Floating Point Number Representation

Consider a binary number **110101**₂ which represents the value:

$$\begin{aligned} &1 * 2^5 + 1 * 2^4 + 0 * 2^3 + 1 * 2^2 + 0 * 2^1 + 1 * 2^0 \\ &= 32 + 16 + 4 + 1 \\ &= \mathbf{53}_{10} \end{aligned}$$

Divide the number 53 by 2, the result would be **26.5**.

How do we represent it if we only had integer representations?

Floating Point Number Representation

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How do we represent it if we only had integer representations?

The **key to represent fractional numbers**, like 26.5 above, is the concept of *binary point*. A binary point is like the decimal point in a decimal system. It acts as a **divider** between the integer and the fractional part of a number.

Floating Point Number Representation

In a **decimal system**, a **decimal point** denotes the **position in a numeral** that the **coefficient should multiply** by $10^0 = 1$. For example, in the numeral 26.5, the coefficient 6 has a weight of $10^0 = 1$.

$$2 * 10^1 + 6 * 10^0 + 5 * 10^{-1} = 26.5$$

As in the decimal system, a **binary point** represents the coefficient of the term $2^0 = 1$. All digits (or bits) to the **left of the binary point** carries a weight of $2^0, 2^1, 2^2$, and so on. Digits (or bits) on the **right of binary point** carries a weight of $2^{-1}, 2^{-2}, 2^{-3}$, and so on. So **11010.1₂** represents the value

$$1 * 2^4 + 1 * 2^3 + 0 * 2^2 + 1 * 2^1 + 0 * 2^0 + 1 * 2^{-1} \\ = 16 + 8 + 2 + 0.5 = 26.5$$

Shifting Operation

Used Instead of Binary Point

In the case of 53_{10} , there is "no" binary point.

Alternatively, we can say the binary point is located at the far right, at position 0. (Think in decimal, 53 and 53.0 represents the same number.)

2^5	2^4	2^3	2^2	2^1	2^0	Binary Point	2^{-1}	2^{-2}	2^{-3}
1	1	0	1	0	1	.	0	0	0

In the case of 26.5_{10} , binary point is located one position to the *left* of 53_{10} . **Shifting an integer to the right by 1 bit position** is equivalent to dividing the number by 2.

2^5	2^4	2^3	2^2	2^1	2^0	Binary Point	2^{-1}	2^{-2}	2^{-3}
0	1	1	0	1	0	.	1	0	0

Fixed Point Number Representation

To represent a real number in computers (or any hardware in general), we can define a fixed point number type simply by **implicitly *fixing* the binary point to be at some position** of a numeral. We need are two parameters:

- width (w) of the number representation, and
- binary point position (b) within the number

Hence, the notation is **fixed<w,b>**

For example, **fixed<8,3>** denotes a 8-bit fixed point number, of which 3 right most bits are fractional. Therefore, the bit pattern:

0	0	0	1	0	1	1	0
---	---	---	---	---	---	---	---

represents a real number:

$$\begin{aligned} 00010.110_2 &= 1 * 2^1 + 1 * 2^{-1} + 1 * 2^{-2} \\ &= 2 + 0.5 + 0.25 = 2.75 \end{aligned}$$

Fixed Point Number Representation

0	0	0	1	0	1	1	0
---	---	---	---	---	---	---	---

A bit pattern can represent anything. Therefore the same bit pattern, if we "cast" it to another type, such as a **fixed<8,5>** type, will represent the number:

$$\begin{aligned} 000.10110_2 &= 1 * 2^{-1} + 1 * 2^{-3} + 1 * 2^{-4} \\ &= 0.5 + 0.125 + 0.0625 = 0.6875 \end{aligned}$$

If we treat this bit pattern as **integer**, it represents the number:

$$\begin{aligned} 00010110_2 &= 1 * 2^4 + 1 * 2^2 + 1 * 2^1 \\ &= 16 + 4 + 2 = 22 \end{aligned}$$

H.T. Write a program in C++ to represent fixed point number.

Data Representation

□ Basic Formats

Figure 3.15 shows the fundamental division of **information into instructions** (operation or control words) and **data** (operands).

Data can be further subdivided into **numerical and nonnumerical**. **Two main number formats** have evolved: fixed-point and floating-point.

The **binary fixed-point format** takes the form $b_A b_B b_C \dots b_K$, where each b_i is 0 or 1 and a binary point is present in some fixed but implicit position.

A floating-point number, on the other hand, consists of a pair of fixed-point numbers M , E , which denote the number $M \times B^E$, where B is a predetermined base.

Data Representation

Nonnumerical data usually take the form of **variable-length character strings** encoded in one of several standard codes, such as *ASCII* code.

Non-numerical data represents characteristics such as a person's gender, marital status, hometown, ethnicity or the types of movies people like. An example is non-numerical data representing the colors of flowers in a yard: yellow, blue, white, red, etc.

Data Representation

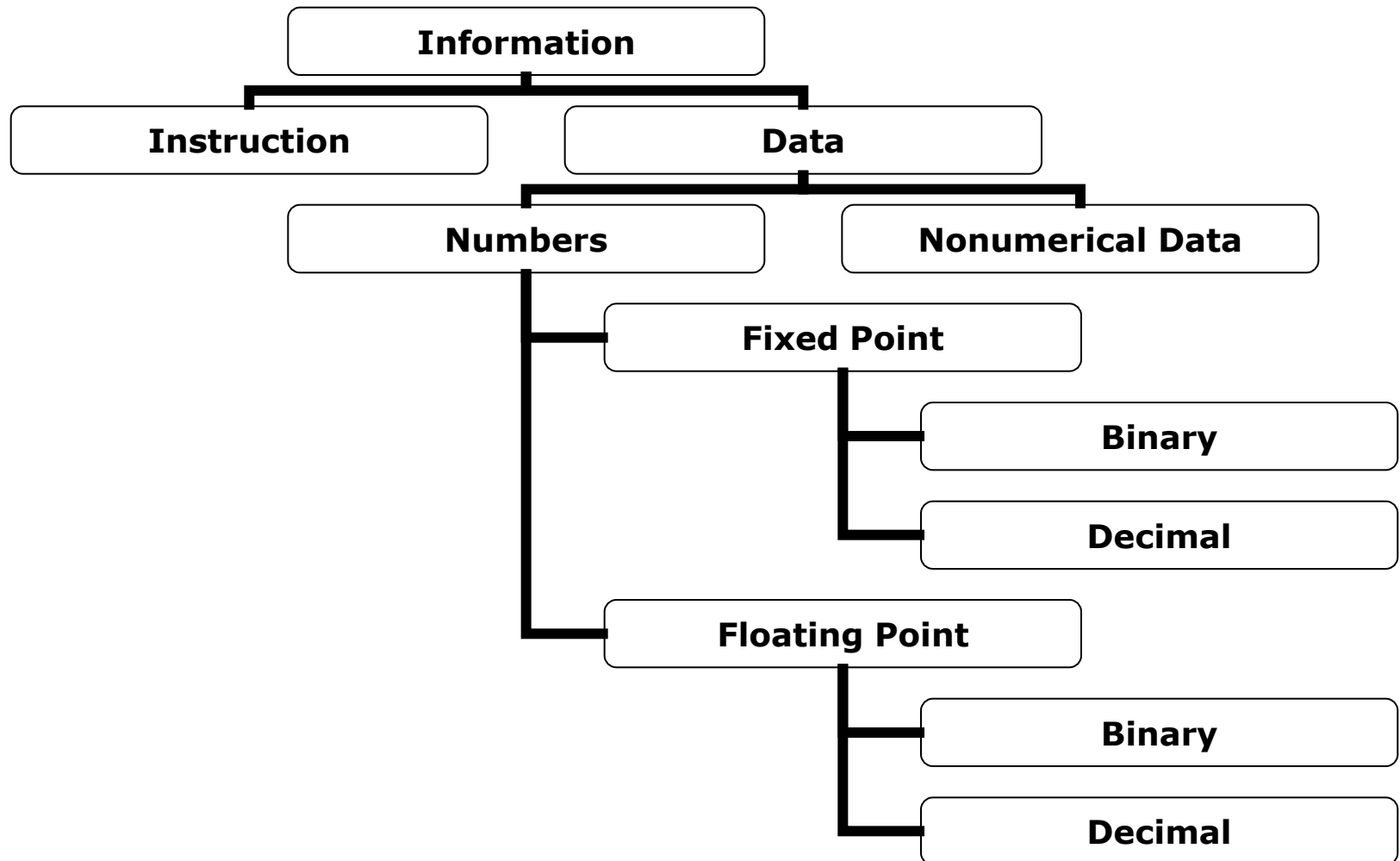


Figure 3.15 The basic information types

Data Representation

□ Word Length

Information is represented in a digital computer by means of **binary words**, where a **word** is a **unit of information** of some fixed length n . An **n -bit word** allows up to 2^n different items to be represented. For example, with $n = 4$, we can encode the 10 decimal digits as follows:

0 = 0000 1 = 0001 2 = 0010 3 = 0011 4 = 0100

5 = 0101 6 = 0110 7 = 0111 8 = 1000 9 = 1001

Number of bits AC can store at a time. Word size is typically a multiple of 8, common CPU word sizes being 8, 16, 32, and 64 bits.

Data Representation

❑ Storage Order

A small but important aspect of data representation is the **way in which the** bits of a word are indexed as shown in Fig.1, where the **right-most bit** is assigned the index 0 and the bits are labeled in increasing order from right to left.

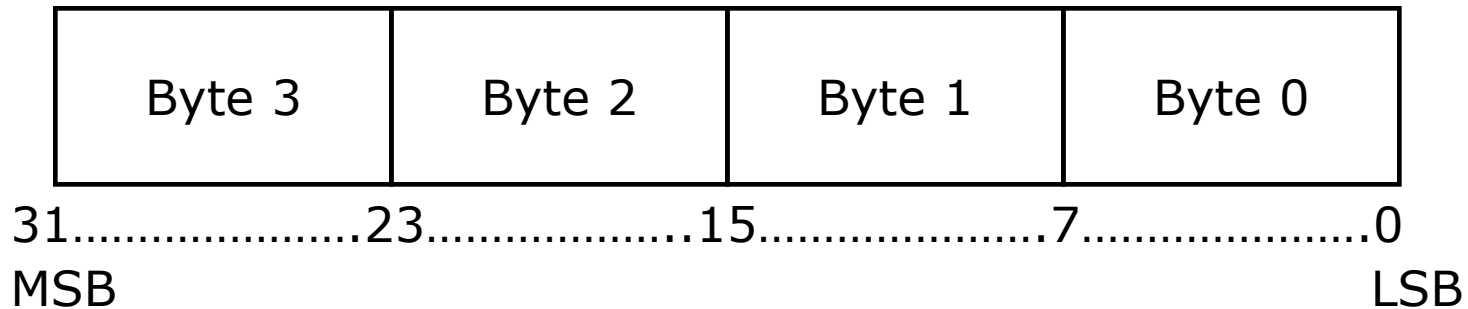


Fig 1. Indexing convention for the bits and bytes of a word.

The advantage of this convention is that when the word is interpreted as an unsigned binary integer, the **low order indexes** correspond to the numerically **less significant bits** and the **high-order indexes more significant bits**.

Data Representation

Suppose a sequence W_0, W_1, \dots, W_m of **m 4-byte number words** is to be stored. Suppose further that W_i as $B_{i,3}, B_{i,2}, B_{i,1}, B_{i,0}$. Hence the entire sequence can be written as

$$W_0, W_1, \dots, W_m = B_{0,3}, B_{0,2}, B_{0,1}, B_{0,0}, B_{1,3}, B_{1,2}, B_{1,1}, B_{1,0}, \dots, B_{m,3}, B_{m,2}, B_{m,1}, B_{m,0}$$

This storage sequence is called **big-endian** and shown in Fig 3.18 a. It is so named because the **most significant byte** $B_{i,3}$ of word W_i is assigned the lowest address and the **least significant byte** $B_{i,0}$ is assigned the highest address. In other words, the big-endian scheme assigns the highest address to byte 0.

Data Representation

The alternative byte-storage scheme called **little-endian** assigns the lowest address to byte 0. This corresponds to

$$W_0, W_1, \dots, W_m = B_{0,0}, B_{0,1}, B_{0,2}, B_{0,3}, B_{1,0}, B_{1,1}, B_{1,2}, B_{1,3}, \dots, B_{m,0}, B_{m,1}, B_{m,2}, B_{m,3}$$

and is illustrated by Fig. 3.18 b.

For example, the Motorola 680X0 uses the big-endian method, whereas the Intel 80X86 series is little-endian.

Data Representation

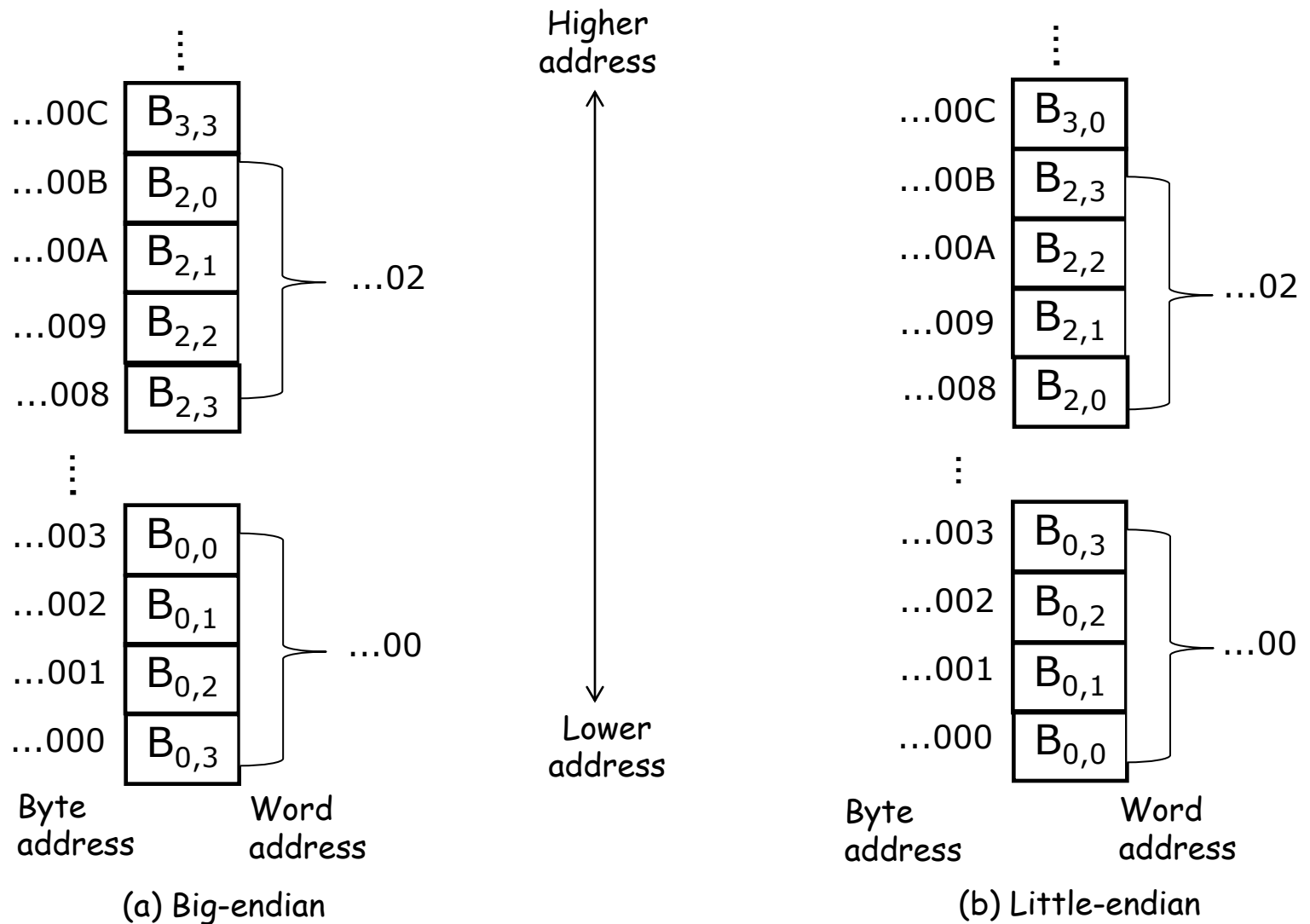


Fig. 3.18 Basic byte storage methods