

# Game Theory, Fall 2022

## Problem Set 6

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1. ST 8.4.

- (a) **Soln:** There are infinite proper subgames because every quantity choice of player 1 results in a proper subgame. Specifically, the set of all the subgames has the same cardinality as  $\mathbb{R}$ .
- (b) **Soln:** This is a game of imperfect information because players 2 and 3 make their choices without observing each other's choice first.
- (c) **Soln:** First we solve for the Nash equilibrium of the simultaneous move stage in which players 2 and 3 make their choices as a function of the choice after player 1. Fix player 1's quantity choice and player 2 maximizes

$$\max_{q_2} \{ \max \{ a - q_1 - q_2 - q_3 - c, 0 \} q_2 \},$$

which yields the best response

$$q_2 = \begin{cases} \frac{a - q_1 - q_3 - c}{2}, & \text{if } q_3 \leq a - c - q_1, \\ 0, & \text{if } q_3 > a - c - q_1. \end{cases}$$

Symmetrically, the best response function of player 3 is

$$q_3 = \begin{cases} \frac{a - q_1 - q_2 - c}{2}, & \text{if } q_2 \leq a - c - q_1, \\ 0, & \text{if } q_2 > a - c - q_1. \end{cases}$$

Hence, from the best response of player 2 and 3, it is apparent that there exists a unique Nash equilibrium of the subgame where player 1 choose  $q_1$ , which is

$$q_2^*(q_1) = q_3^*(q_1) = \begin{cases} \frac{a - q_1 - c}{3}, & \text{if } q_1 \leq a - c, \\ 0, & \text{if } q_1 > a - c. \end{cases}$$

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Back to the first node where player 1 chooses quantity, player 1 would maximize

$$\max_{q_1} \{ \max \{ a - q_1 - q_2^*(q_1) - q_3^*(q_1) - c, 0 \} q_1 \},$$

which leads to the FOC

$$\frac{a - 2q_1 - c}{3} = 0,$$

resulting in a unique solution  $q_1 = \frac{a-c}{2}$ . To sum up, the unique subgame-perfect equilibrium dictates that  $q_1^* = \frac{a-c}{2}$  and

$$q_2^*(q_1) = q_3^*(q_1) = \begin{cases} \frac{a-q_1-c}{3}, & \text{if } q_1 \leq a - c, \\ 0, & \text{if } q_1 > a - c. \end{cases}$$

- (d) **Soln:** There are many Nash equilibria in this game. We now prove that  $q_1^* = \frac{a-c}{2}$  and

$$q_2^*(q_1) = q_3^*(q_1) = \begin{cases} \frac{a-q_1-c}{3}, & \text{if } q_1 = \frac{a-c}{2}, \\ a, & \text{otherwise.} \end{cases}$$

is a Nash equilibrium. When player 1 chooses  $\frac{a-c}{2}$ , player 2 and 3 would not deviate from their best responses. Meanwhile, player 1 would not deviate because any other choice would cause his opponents to choose  $a$  and bring player 1 negative profits. Hence the strategy profile above is a Nash equilibrium.

Apparently this strategy profile is not a subgame-perfect equilibrium, because the strategies of player 2 and 3 are not best responses in the subgames induced by a slightly different action of player 1.

## 2. ST 8.6.

- (a) **Soln:** There are two subgames after firm 1 decides whether to invest or not. If firm 1 decides not to invest, each firm  $i$  would solve

$$\max_{q_i} \{ \max \{ 100 - 10 - q_1 - q_2, 0 \} q_i \}.$$

The solutions to the problems yield the unique Nash equilibrium  $q_1 = q_2 = 30$ , and the profit of firm 1 in this subgame Nash equilibrium is  $(100 - 30 - 30 - 10) \times 30 = 900$ .

If firm 1 decides to invest, firm 1 would solve a new problem, which is

$$\max_{q_1} \{ \max \{ 100 - 5 - q_1 - q_2, 0 \} q_1 - F \},$$

while firm 2 would still solve the same problem above. The solutions to the problems yield another unique Nash equilibrium  $q_1 = \frac{100}{3}$  and  $q_2 = \frac{85}{3}$ . And the profit of firm 1 in this subgame Nash equilibrium is  $(100 - \frac{100}{3} - \frac{85}{3} - 5) \times \frac{100}{3} - F = \frac{10000}{9} - F$ .

Hence, if there is a unique subgame-perfect equilibrium involving firm 1 investing, it must be strictly better to invest than not to invest, which means  $\frac{10000}{9} - F > 900$ . Therefore, all  $F^* \in (0, \frac{1900}{9})$  lead to a unique subgame-perfect equilibrium where firm 1 decides to invest.

- (b) **Soln:** If  $F > \frac{1900}{9}$ ,  $F$  would be larger than all the possible  $F^*$ , and we can construct a Nash equilibrium in which firm 1 decides not to invest, and will choose  $q_1 = 30$  when not investing and  $q_1 = \frac{100}{3}$  when investing, while firm 2 will choose  $q_2 = 30$  when firm 1 does not invest, and  $q_2 = 100$  otherwise. The outcome would be firm 1 doesn't invest and both firms produce 30. It's obvious that both firms would not deviate. Hence the strategy profile is indeed a Nash equilibrium. Moreover, this strategy profile is not a subgame-perfect equilibrium, because  $q_2 = 100$  is not firm 2's best response when firm 1 invests.

### 3. ST 8.9.

- (a) **Soln:** In the normal form representation of the game, the set of firms is  $N = \{1, 2\}$ . Each firm  $i \in \mathcal{N}$  have a strategy space  $S_i = [0, +\infty)$ . The payoff function of firm  $i$  is

$$u_i(q_i, q_{-i}) = \max\{100 - q_1 - q_2, 0\}q_i - k\chi_{\{q_i > 0\}}.$$

in which  $\chi_{\{q_i > 0\}}$  is the indicator function for the set  $\{q_i > 0\}$ .

- (b) **Soln:** The best response of firm  $i$  is

$$q_i^*(q_j) = \begin{cases} \frac{100 - q_j}{2}, & \text{if } q_j < 100 - 20\sqrt{10}, \\ \{10\sqrt{10}, 0\}, & \text{if } q_j = 100 - 20\sqrt{10}, \\ 0, & \text{if } q_j > 100 - 20\sqrt{10}. \end{cases}$$

From the best responses of the firms, we can find three pure Nash equilibria in the game, which are  $(\frac{100}{3}, \frac{100}{3})$ ,  $(50, 0)$ ,  $(0, 50)$ .

- (c) **Soln:** The game tree for question 3c is depicted in Figure 1. The payoff is the same as the payoff in 3a.

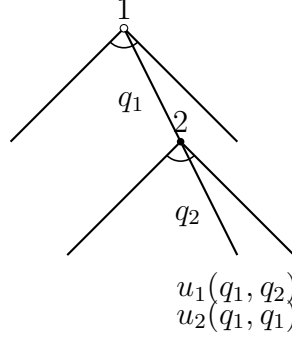


Figure 1: The game tree for Question 3c

When  $k = 25$ , the best response of firm 2 becomes

$$q_2^*(q_1) = \begin{cases} \frac{100-q_1}{2}, & \text{if } q_1 < 90, \\ 0, & \text{if } q_1 \geq 90. \end{cases}$$

Hence firm 1 might solve

$$\max_{q_1 < 90} \left\{ (100 - q_1 - \frac{100 - q_1}{2}) q_1 \right\}$$

if the solution leads to positive profits for both firms, considering the fixed cost. The solution to the problem above is  $q_1 = 50$ , yielding that  $q_2 = 25$ ,  $u_1 = (100 - 50 - 25) \times 50 - 25 > 0$ , and  $u_2 = 25 \times 25 - 25 > 0$ . Moreover, it's easy to verify that choosing  $q_1 \geq 90$  is less profitable than this outcome, no matter what firm 2's strategy is. Hence there is a unique backward induction solution, in which  $q_1 = 50$  and

$$q_2 = \begin{cases} \frac{100-q_1}{2}, & \text{if } q_1 < 90, \\ 0, & \text{if } q_1 \geq 90. \end{cases}$$

(d) **Soln:** When  $k = 725$ , the best response of firm 2 becomes

$$q_2^*(q_1) = \begin{cases} \frac{100-q_1}{2}, & \text{if } q_1 < 100 - 10\sqrt{29}, \\ 0, & \text{if } q_1 \geq 100 - 10\sqrt{29}. \end{cases}$$

For the initial node, first consider the case where  $q_1 \geq 100 - 10\sqrt{29}$ , firm 2 quits the market and firm 1 solves

$$\max_{q_1 \geq 100 - 10\sqrt{29}} \{ (100 - q_1) q_1 \},$$

which yields  $q_1 = 50$ . The corresponding payoff is  $u_1 = 2500 - 725 = 1775 > 0$  and  $u_2 = 0$ .

Otherwise, firm 1 might solve

$$\max_{q_1 < 100 - 10\sqrt{29}} \left\{ (100 - q_1 - \frac{100 - q_1}{2})q_1 \right\}.$$

If the solution leads to positive profits for both firms, considering the fixed cost. However, it's obvious that

$$\begin{aligned} \max_{100 - 10\sqrt{29} \leq q_1 \leq 100} \{(100 - q_1)q_1\} &= \max_{q_1 \leq 100} \{(100 - q_1)q_1\} \\ &\geq \max_{q_1 \leq 100} \left\{ (100 - q_1 - \frac{100 - q_1}{2})q_1 \right\} \\ &> \max_{q_1 < 100 - 10\sqrt{29}} \left\{ (100 - q_1 - \frac{100 - q_1}{2})q_1 \right\}. \end{aligned}$$

Thus, there's a unique backward induction solution, in which  $q_1 = 50$  and

$$q_2 = \begin{cases} \frac{100 - q_1}{2}, & \text{if } q_1 < 100 - 10\sqrt{29}, \\ 0, & \text{if } q_1 \geq 100 - 10\sqrt{29}. \end{cases}$$

4. ST 8.12.

- (a) **Soln:** This is a game of perfect information. There are two players,  $i \in \{1, 2\}$ . The strategy sets are  $S_1 = X = [0, 5]$  and  $S_2 = \{A, R\}$ , where  $A$  denotes accepting the proposal  $x \in X$ , and  $R$  means rejecting it and adopting the status quo  $q = 4$ . The payoffs are given by

$$v_1(s_1, s_2) = \begin{cases} 10 - |s_1 - 1|, & \text{if } s_2 = A, \\ 7, & \text{if } s_2 = R, \end{cases}$$

and

$$v_2(s_1, s_2) = \begin{cases} 10 - |s_1 - 3|, & \text{if } s_2 = A, \\ 9, & \text{if } s_2 = R. \end{cases}$$

- (b) **Soln:** Player 2 can guarantee himself a payoff of 9 by choosing  $R$ , implying that his unique best response is to choose  $A$  if and only if  $10 - |s_1 - 3| > 9$ , which holds for any  $s_1 \in (2, 4)$ , and  $R$  for  $s_1 > 4$  or  $s_1 < 2$ . He would be indifferent between  $A$  and  $R$  if  $s_1 = 2$  or  $s_1 = 4$ . On the other hand, given that player 2 always accepts for any  $s_1 \in (2, 4)$ , player 1 would like to have an alternative proposal that is accepted by player 2 and as close to 1 as possible. If player 2 rejects  $s_1 = 2$ , there would be no subgame-perfect equilibria since for any  $s_1 \in (2, 4]$ , there always exists  $2 < s'_1 = \frac{2 + s_1}{2} < s_1$ , such that  $s'_1$  is closer to 1 than  $s_1$ , which means players

1 can always deviate to get a higher payoff. Thus, we only need to consider the case where player 2 accepts  $s_1 = 2$ . When  $s_1 = 4$ , he is indifferent between  $A$  and  $R$ , and both strategies can be equilibria. Given player 2's strategy, player 1's best response is to choose  $s_1 = 2$ , which is the closest to 1 for any  $s_1 \in [2, 4]$  (or  $s_1 \in [2, 4)$ ), and let player 2 accept. Thus, there are two subgame-perfect equilibria which result in the same payoffs  $(v_1, v_2) = (9, 9)$ .

- (c) **Soln:** One Nash equilibrium is that player 2 adopts the strategy “I will reject anything except  $s_1 = 3$ ” and that player 1 chooses 3. By doing so player 1 gets 8, while any other choice of  $s_1$  is expected to yield a payoff of 7. Hence, player 1's best response to player 2's proposed strategy is indeed to choose  $s_1 = 3$ , and the payoffs from this Nash equilibrium are  $(v_1, v_2) = (8, 10)$ .

Since player 2 can guarantee himself a payoff of 9, there are actually infinite Nash equilibria that are not subgame-perfect following a similar logic: player 2 adopts the strategy “I will reject anything except  $s_1 = x$ ” for some value  $x \in [2, 4)$ . Player 1 would strictly prefer the adoption of  $x$  over 4, and hence would indeed propose  $x$ , and player 2 would accept the proposal. For  $x = 4$  both players are indifferent so it would also be supported as a Nash equilibrium.

5. ST 8.13.

- (a) **Soln:** The game tree for question 5a is depicted in Figure 2.

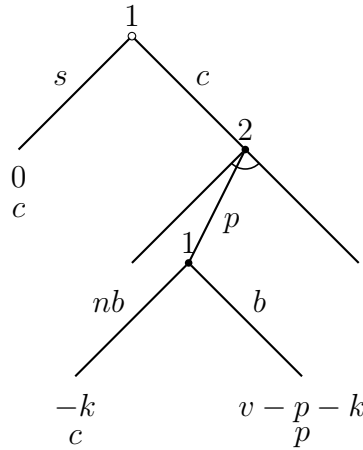


Figure 2: The game tree for Question 5a

Buyer's best response to seller's offer is

$$BR_1(c, p) = \begin{cases} b, & \text{if } p < v, \\ \{b, nb\}, & \text{if } p = v, \\ nb, & \text{if } p > v. \end{cases}$$

- (b) **Soln:** Backward induction implies that buyer should reject the offer when  $p < v$ , and accept the offer when  $p > v$ . It is indifferent for the buyer between accepting  $p = v$  or not. If the buyer reject the offer when  $p = v$ , there would be no best response for the seller because for any  $p < v$  the seller could deviate to  $p' = \frac{p+v}{2} > p$  to get a higher payoff than  $p$ . Hence in a subgame-perfect equilibrium, the buyer should accept  $p = v$ , and the seller would offer  $p = v$ .

In this case, the buyer could only get  $-k$  if he commutes to the store, so he would stay home. This is the unique backward induction solution, in which the buyer decides to stay home, the seller would propose  $p = v$  if the buyer commutes, and the buyer would accept any  $p \geq v$  and reject others.

The outcome of the equilibrium is  $(0, c)$ , which is Pareto dominated by  $(0, v - k)$  from the path that the buyer commutes and accepts the propose  $p = v - k$  from the seller. Hence it is not Pareto optimal.

- (c) **Soln:** Consider the strategy profile in which the buyer commutes to the store, the seller proposes  $p = \frac{v-k+c}{2} > c$ , and the buyer only accepts  $p = \frac{v-k+c}{2}$  and rejects other cases. It's evident that nobody has the incentive to deviate. The payoff of this strategy profile is  $(\frac{v-k-c}{2}, \frac{v-k+c}{2})$ , which is strictly higher than  $(0, c)$ .
- (d) **Soln:** If the seller chooses to send the postcard, he actually sets the price before the buyer decides whether to leave home, rather than after the buyer arrives. Thus, the game has a different timing. Backward induction implies that buyer should leave home and buy the good when  $p < v - k$ , stay home when  $p > v - k$ , and be indifferent when  $p = v - k$ . Similar to 5b, in a subgame-perfect equilibrium, the buyer should accept  $p = v - k$ , and the seller would offer  $p = v - k$ . Finally, the payoff of the buyer remains 0, but the seller gets  $p - \epsilon = v - k - \epsilon > c$  and hence is strictly better off than in the subgame-perfect equilibrium of 5b, which is exactly the case where the seller doesn't send the postcard. Actually, when the seller commits  $c + \epsilon < p < v - k$ , the buyer would choose to buy and then both of them would be better off. Therefore, the seller would choose to send the postcard.