

# Problem Set 1

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(a)

Answer:

There are  $n$  bidders,  $N = [n]$ . Every bidder bids a price  $b_i \geq 0$ . Thus the strategy is  $S_i = [0, +\infty)$   
Then the payoff is

$$v_i(b_i, b_{-i}) = \begin{cases} \frac{v_i - b_i}{|\{j | b_i = b_j\}|}, & \forall j, b_i \geq b_j \\ 0, & o.w. \end{cases}$$

(b)

For the second price auction, we have,

$$v_i(b_i, b_{-i}) = \begin{cases} \frac{v_i - \max_{j < i} b_j}{|\{j | b_i = b_j\}|}, & \forall j, b_i \geq b_j \\ 0, & o.w. \end{cases}$$

## 2. ST Exercise 4.5

Answer:

In the first round,  $U$  for player 1 is strictly dominated by  $M$ .  
In the first round,  $C$  for player 2 is strictly dominated by  $H$ .  
Hence  $\{M, D\} \times \{L, R\}$  survives.

1\2	L	C	R
U	6,8	2,6	8,2
M	8,2	4,4	9,5
D	8,10	4,6	6,7

### 3. ST Exercise 4.6

Answer:

(a)

For  $i$ , when  $t_j \geq 10$ , his best reponse is  $t_i = 0$ . When  $t_j < 10$ , the first order condition is  $10 - t_j - 2t_i = 0$ . Then we find,

$$t_i = \max\{\frac{10 - t_j}{2}, 0\}$$

(b)

- $\forall t_i \in [0, 5]$  is not strictly dominated because it is a best response to  $t_j = 10 - 2t_i$ .
- For  $t_i > 5$ , assume  $k = t_i - 5 > 0$ , then we have

$$\begin{aligned} v_i(5 + k, t_j) &= (10 - t_j)(5 + k) - (5 + k)^2 = 25 - 5t_i - k^2 - t_j k \\ &< 25 - 5t_j = v_i(5, t_j) \end{aligned}$$

It means strategy  $t_i > 5$  is strictly dominated by  $t_i = 5$ .

- We have  $S_1^1 = S_2^2 = [0, 5]$

(c)

- Define  $f(x) = 5 - x/2$ .
- Assume  $S_1^* = S_2^* \subset [l, u]$ 
  - First,  $S_1^* = S_2^* \subset [0, 5]$
  - Second we find,  $[f(u), f(l)] \subset [l, u], \forall l, u > 0, l + u \leq 5$
  - Similar to (b), we can find  $t_i \in [f(u), f(l)]$  is not strictly dominated
  - For  $t_i < f(u)$ , assume  $k = f(u) - t_i > 0$ , similarly,

$$\begin{aligned} v(f(u), t_j) - v(f(u) - k, t_j) &= k(10 - t_j - 2f(u) + k) \\ &\geq k(10 - 5 - 2f(5) + k) > 0 \end{aligned}$$

Thus  $t_i$  is strictly dominated by  $t_i = f(u)$ .

- For  $t_i > f(l)$ , similarly, is strictly dominated by  $t_i = f(l)$ .
- By induction, we can get closed interval  $T_n \subset T_{n-1} \subset \dots \subset T_1 = [0, 5]$ , s.t.

$$T_i = [l_i, u_i] \Rightarrow T_{i+1} = [f(u_i), f(l_i)], i \in \mathbb{N}^+$$

and

$$S_1^* = S_2^* \subset T_n$$

- Given that the measure  $|f(u) - f(l)| = |u - l|/2$ , we can get

$$\lim_{n \rightarrow +\infty} |u_n - l_n| = 0$$

By the closed interval theorem,  $T_{+\infty} = \{c\}$ . And from the method of generating  $T_n$ , we can know  $f(c) = c$ , which means  $c = 10/3$ .

- In conclusion,  $S_1^* = S_2^* = \bigcap_{n=1}^{+\infty} T_n = T_{+\infty} = \{10/3\}$ . The unique pair of strategies that survive IESDS for this game are  $t_1 = t_2 = 10/3$

## 4. ST Exercise 4.7

**Answer:**

**(a)**

The set of the player is  $\{1, 2\}$ . The strategy spaces are  $S_1 = S_2 = \{P, B, N\}$ . The payoffs to player 1 are:

$$\begin{aligned} v_1(P, P) &= 0.5, & v_1(P, B) &= 0, & v_1(P, N) &= 0.3, \\ v_1(B, P) &= 1, & v_1(B, B) &= 0.5, & v_1(B, N) &= 0.4, \\ v_1(N, P) &= 0.7, & v_1(N, B) &= 0.6, & v_1(N, N) &= 0.5 \end{aligned}$$

And,  $\forall s \in \{P, B, N\} \times \{P, B, N\}, v_2(s) = 1 - v_1(s)$ .

**(b)**

The matrix is following.

		Player 2		
		P	B	N
Player 1	P	0.5, 0.5	0, 1	0.3, 0.7
	B	1, 0	0.5, 0.5	0.4, 0.6

<b>N</b>	0.7, 0.3	0.6, 0.4	0.5, 0.5
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**(c)**

In the first round,  $P$  for both players is strictly dominated by  $B$ .

In the second round,  $B$  for both players is strictly dominated by  $N$ .

Thus  $\{N\} \times \{N\}$  survives, which leads to a clear prediction.

## 5. ST Exercise 4.8

**Answer:**

**(a)**

The average is less than 20 regardless of the number of players, and  $3/4$  of the average is less than 15, This means 18 also makes  $i$  the only winner.

**(b)**

- $x = 20$  is not one of the best responses because there will be many winners.
- For  $x < 20$ , the average is  $\frac{20(n-1)+x}{n}$ , thus player  $i$  is the only winner  $\Leftrightarrow$

$$\left| 20 - \frac{20(n-1)+x}{n} \right| > \left| \frac{20(n-1)+x}{n} \right|$$

When  $x < 20$ , it is equivalent to  $x > 10 - \frac{30}{2n-3}$

- Therefore the set of best responses is

$$\left\{ x \in \mathbb{Z} \left| 10 - \frac{30}{2n-3} < x < 20 \right. \right\}, n \geq 2,$$

which depends on the number of players .

**6.**

**Answer:**

**(a)**

$$v_i(x_i, x_j) = \arctan x_i < \arctan(x_i + 1) = v_i(x_i + 1, x_j)$$

Thus  $x_i$  is strictly dominated.

**(b)**

$$\forall x_i > 0, v_i(0, 1) = 2 > \arctan x_i = v_i(x_i, 1)$$

Thus 0 for player i is a best response to  $x_j = 1$ , which is not strictly dominated.

**(c)**

$$v_i(0, 0) = \arctan 0 < \arctan 1 = v_i(1, 0)$$

They are obviously not mutual best responses.

**7.**

**Answer:**

**(a)**

$n$  firms,  $N = [n]$ .  $S_i = \mathbb{R}_+$ . The payoff is:

$$v_i(q_1, q_2, \dots, q_n) = (\max\{0, 100 - \sum_{j \in N} q_j\} - 10)q_i$$

**(b)**

Maximize the above payoff function, we get,

$$q_i = \max \left\{ 0, \frac{90 - \sum_{j \neq i} q_j}{2} \right\}$$

- Given that  $q_j > 0$ , when  $90 - \sum_{j \neq i} q_j = c > 0$ ,  $q_i = c$  is the best response, which traverses  $[0, 45]$  when  $n \geq 3$

- On the contrary, when  $n \geq 3$ , we can always construct a  $q_{-1}$  such that  $90 - \sum_{j \neq i} q_j$  can traverse  $[0, 45]$ , that means  $[0, 45]$  survives IESDS.
- for  $q_j > 45$ ,
  - if  $90 - \sum_{j \neq i} q_j \geq q_i$ ,

$$v_i(0, q_{-i}) - v_i(q_i, q_{-i}) = 0 - (-10q_i) > 0$$

- if  $90 - \sum_{j \neq i} q_j < q_i$ ,

$$v_i(45, q_{-i}) - v_i(q_i, q_{-i}) = (q_i - 45)(q_i - 45 + \sum_{j \neq i} q_j) > 0$$

- That means  $(45, +\infty)$  is strictly dominated. And  $[0, 45]$  survives IESDS when  $n \geq 3$
- When  $n = 2$ , similar to 4.6(c),  $S_1^* = S_2^* = \{\text{Fixed point of function } f(x) = 45 - x/2\} = \{30\}$
- In conclusion,

$$S_i^* = \begin{cases} \{30\}, & n = 2, \\ [0, 45], & n \geq 3 \end{cases}$$