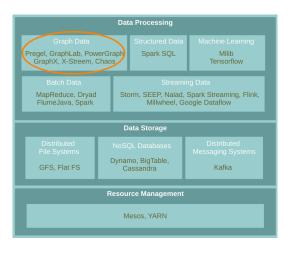


Large Scale Graph Processing - GraphX, Giraph++, and Pegasus

Amir H. Payberah payberah@kth.se 09/10/2018



https://id2221kth.github.io













- ▶ Difficult to extract parallelism based on partitioning of the data.
- ▶ Difficult to express parallelism based on partitioning of computation.



Different Approached to Process Large Scale Graphs

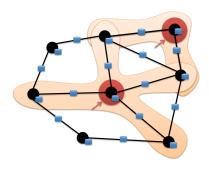
- ► Think like a vertex
- ► Think like an edge
- ► Think like a table
- ► Think like a graph
- ► Think like a matrix



Think Like a Table

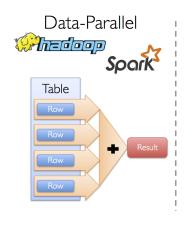


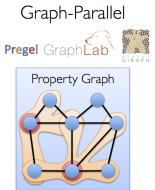
Graph-Parallel Processing Model





Data-Parallel vs. Graph-Parallel Computation





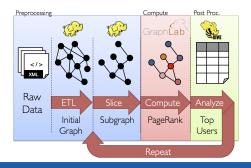
Motivation (2/3)

► Graph-parallel computation: restricting the types of computation to achieve performance.

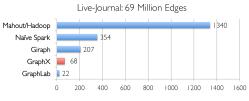


Motivation (2/3)

- Graph-parallel computation: restricting the types of computation to achieve performance.
- ► The same restrictions make it difficult and inefficient to express many stages in a typical graph-analytics pipeline.

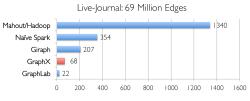




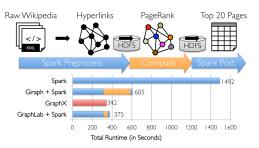


Runtime (in seconds, PageRank for 10 iterations)





Runtime (in seconds, PageRank for 10 iterations)



- ▶ Unifies data-parallel and graph-parallel systems.
- ▶ Tables and Graphs are composable views of the same physical data.





$\mathsf{Graph}\mathsf{X}$



- ► GraphX is the library to perform graph-parallel processing in Spark.
- ► In-memory caching.
- ► Lineage-based fault tolerance.





The Property Graph Data Model

- ► Spark represent graph structured data as a property graph.
- ▶ It is logically represented as a pair of vertex and edge property collections.





| Vertex Table | |
|--------------------|--|
| Property (V) | |
| (rxin, student) | |
| (jgonzal, postdoc) | |
| | |

| | (rxin, student) |
|---|-----------------------|
| | (jgonzal, postdoc) |
| | (franklin, professor) |
| | (istoica, professor) |
| E | dge Table |

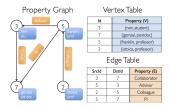
Srcld Dstld Property (E) 3 7 Collaborator 5 3 Advisor 2 5 Colleague



▶ VertexRDD: contains the vertex properties keyed by the vertex ID.

```
class Graph[VD, ED] {
  val vertices: VertexRDD[VD]
  val edges: EdgeRDD[ED]
}

// VD: the type of the vertex attribute
abstract class VertexRDD[VD] extends RDD[(VertexId, VD)]
```





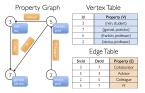


The Edge Collection

► EdgeRDD: contains the edge properties keyed by the source and destination vertex IDs.

```
class Graph[VD, ED] {
  val vertices: VertexRDD[VD]
  val edges: EdgeRDD[ED]
}

// ED: the type of the edge attribute
case class Edge[ED](srcId: VertexId, dstId: VertexId, attr: ED)
abstract class EdgeRDD[ED] extends RDD[Edge[ED]]
```



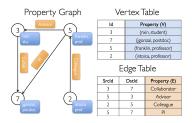


- ► The triplets collection consists of each edge and its corresponding source and destination vertex properties.
- ▶ It logically joins the vertex and edge properties: RDD [EdgeTriplet[VD, ED]].
- ► The EdgeTriplet class extends the Edge class by adding the srcAttr and dstAttr members, which contain the source and destination properties respectively.





Building a Property Graph



Graph Operators

- ► Information about the graph
- Property operators
- Structural operators
- ► Joins
- ► Aggregation
- ► Iterative computation
- **...**



Information About The Graph (1/2)

► Information about the graph

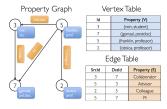
```
val numEdges: Long
val numVertices: Long
val inDegrees: VertexRDD[Int]
val outDegrees: VertexRDD[Int]
val degrees: VertexRDD[Int]
```

Views of the graph as collections

```
val vertices: VertexRDD[VD]
val edges: EdgeRDD[ED]
val triplets: RDD[EdgeTriplet[VD, ED]]
```



Information About The Graph (2/2)



```
// Constructed from above
val graph: Graph[(String, String), String]

// Count all users which are postdocs
graph.vertices.filter { case (id, (name, pos)) => pos == "postdoc" }.count

// Count all the edges where src > dst
graph.edges.filter(e => e.srcId > e.dstId).count
```



Property Operators

- ► Transform vertex and edge attributes
- ► Each of these operators yields a new graph with the vertex or edge properties modified by the user defined map function.

```
def mapVertices[VD2](map: (VertexId, VD) => VD2): Graph[VD2, ED]
def mapEdges[ED2](map: Edge[ED] => ED2): Graph[VD, ED2]
def mapTriplets[ED2](map: EdgeTriplet[VD, ED] => ED2): Graph[VD, ED2]
```



Property Operators

- ► Transform vertex and edge attributes
- ► Each of these operators yields a new graph with the vertex or edge properties modified by the user defined map function.



Property Operators

- ► Transform vertex and edge attributes
- ► Each of these operators yields a new graph with the vertex or edge properties modified by the user defined map function.

► reverse returns a new graph with all the edge directions reversed.

```
def reverse: Graph[VD, ED]
def subgraph(epred: EdgeTriplet[VD, ED] => Boolean, vpred: (VertexId, VD) => Boolean):
    Graph[VD, ED]
def mask[VD2, ED2](other: Graph[VD2, ED2]): Graph[VD, ED]
```



Structural Operators

- reverse returns a new graph with all the edge directions reversed.
- subgraph takes vertex/edge predicates and returns the graph containing only the vertices/edges that satisfy the given predicate.

```
def reverse: Graph[VD, ED]
def subgraph(epred: EdgeTriplet[VD, ED] => Boolean, vpred: (VertexId, VD) => Boolean):
        Graph[VD, ED]
def mask[VD2, ED2](other: Graph[VD2, ED2]): Graph[VD, ED]

// Remove missing vertices as well as the edges to connected to them
val validGraph = graph.subgraph(vpred = (id, attr) => attr._2 != "Missing")
graph.vertices.collect.foreach(println)
validGraph.vertices.collect.foreach(println)

// Restrict the answer to the valid subgraph
val validUserGraph = graph.mask(validGraph)
```



Structural Operators

- ▶ reverse returns a new graph with all the edge directions reversed.
- ▶ subgraph takes vertex/edge predicates and returns the graph containing only the vertices/edges that satisfy the given predicate.
- mask constructs a subgraph of the input graph.

```
def reverse: Graph[VD, ED]
def subgraph(epred: EdgeTriplet[VD, ED] => Boolean, vpred: (VertexId, VD) => Boolean):
        Graph[VD, ED]
def mask[VD2, ED2](other: Graph[VD2, ED2]): Graph[VD, ED]

// Remove missing vertices as well as the edges to connected to them
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validGraph.vertices.collect.foreach(println)

// Restrict the answer to the valid subgraph
val validUserGraph = graph.mask(validGraph)
```

- joinVertices joins the vertices with the input RDD.
 - Returns a new graph with the vertex properties obtained by applying the user defined map function to the result of the joined vertices.
 - Vertices without a matching value in the RDD retain their original value.

```
def joinVertices[U](table: RDD[(VertexId, U)])(map: (VertexId, VD, U) => VD): Graph[VD, ED]
```

- joinVertices joins the vertices with the input RDD.
 - Returns a new graph with the vertex properties obtained by applying the user defined map function to the result of the joined vertices.
 - Vertices without a matching value in the RDD retain their original value.

```
def joinVertices[U](table: RDD[(VertexId, U)])(map: (VertexId, VD, U) => VD): Graph[VD, ED]
```

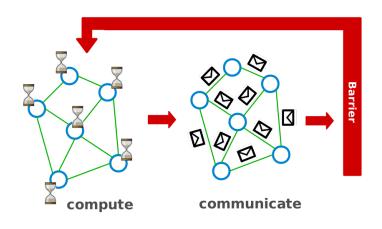
```
val rdd: RDD[(VertexId, String)] = sc.parallelize(Array((3L, "phd")))
val joinedGraph = graph.joinVertices(rdd)((id, user, role) => (user._1, role + " " + user._2))
joinedGraph.vertices.collect.foreach(println)
```

aggregateMessages applies a user defined sendMsg function to each edge triplet in the graph and then uses the mergeMsg function to aggregate those messages at their destination vertex.

```
def aggregateMessages[Msg: ClassTag](
   sendMsg: EdgeContext[VD, ED, Msg] => Unit, // map
   mergeMsg: (Msg, Msg) => Msg, // reduce
   tripletFields: TripletFields = TripletFields.All):
   VertexRDD[Msg]
```

```
// count and list the name of friends of each user
val profs: VertexRDD[(Int, String)] = validUserGraph.aggregateMessages[(Int, String)](
    // map
    triplet => {
        triplet.sendToDst((1, triplet.srcAttr._1))
    },
    // reduce
    (a, b) => (a._1 + b._1, a._2 + " " + b._2)
)
profs.collect.foreach(println)
```





```
Iterative Computation (2/9)
```

```
i_val := val
for each message m
  if m > val then val := m
if i_val == val then
  vote_to_halt
else
  for each neighbor v
    send_message(v, val)
```

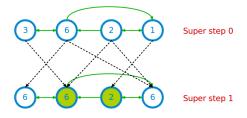


Super step 0

```
i_val := val

for each message m
   if m > val then val := m

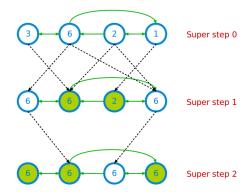
if i_val == val then
   vote_to_halt
else
   for each neighbor v
        send_message(v, val)
```



```
i_val := val

for each message m
  if m > val then val := m

if i_val == val then
  vote_to_halt
else
  for each neighbor v
    send_message(v, val)
```

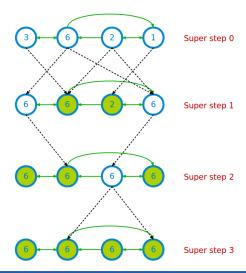


Iterative Computation (5/9)

```
i_val := val

for each message m
   if m > val then val := m

if i_val == val then
   vote_to_halt
else
   for each neighbor v
        send_message(v, val)
```



▶ pregel takes two argument lists: graph.pregel(list1)(list2).

```
def pregel[A]
  (initialMsg: A, maxIter: Int = Int.MaxValue, activeDir: EdgeDirection = EdgeDirection.Out)
  (vprog: (VertexId, VD, A) => VD, sendMsg: EdgeTriplet[VD, ED] => Iterator[(VertexId, A)],
    mergeMsg: (A, A) => A):
  Graph[VD, ED]
```



Iterative Computation (6/9)

- ▶ pregel takes two argument lists: graph.pregel(list1)(list2).
- ► The first list contains configuration parameters
 - The initial message, the maximum number of iterations, and the edge direction in which to send messages (by default along out edges).

```
def pregel[A]
  (initialMsg: A, maxIter: Int = Int.MaxValue, activeDir: EdgeDirection = EdgeDirection.Out)
  (vprog: (VertexId, VD, A) => VD, sendMsg: EdgeTriplet[VD, ED] => Iterator[(VertexId, A)],
    mergeMsg: (A, A) => A):
  Graph[VD, ED]
```



Iterative Computation (6/9)

- ▶ pregel takes two argument lists: graph.pregel(list1)(list2).
- ► The first list contains configuration parameters
 - The initial message, the maximum number of iterations, and the edge direction in which to send messages (by default along out edges).
- ▶ The second list contains the user defined functions.
 - For receiving messages (the vertex program vprog), computing messages (sendMsg), and combining messages (mergeMsg).

```
def pregel[A]
  (initialMsg: A, maxIter: Int = Int.MaxValue, activeDir: EdgeDirection = EdgeDirection.Out)
  (vprog: (VertexId, VD, A) => VD, sendMsg: EdgeTriplet[VD, ED] => Iterator[(VertexId, A)],
    mergeMsg: (A, A) => A):
  Graph[VD, ED]
```

Iterative Computation (7/9)

```
import org.apache.spark._
import org.apache.spark.graphx._
import org.apache.spark.rdd.RDD

val initialMsg = -9999

val vertices: RDD[(VertexId, (Int, Int))] = sc.parallelize(Array((1L, (1, -1)), (2L, (2, -1)), (3L, (3, -1)), (6L, (6, -1))))

val relationships: RDD[Edge[Boolean]] = sc.parallelize(Array(Edge(1L, 2L, true), Edge(2L, 1L, true), Edge(2L, 6L, true), Edge(3L, 6L, true), Edge(6L, 3L, true)))

val graph = Graph(vertices, relationships)
```





Iterative Computation (8/9)

```
// the function for receiving messages
def vprog(vertexId: VertexId, value: (Int, Int), message: Int): (Int, Int) = {
   if (message == initialMsg)
     value
   else
      (math.max(message, value._1), value._1)
}
```



Iterative Computation (8/9)

```
// the function for receiving messages
def vprog(vertexId: VertexId, value: (Int, Int), message: Int): (Int, Int) = {
 if (message == initialMsg)
   value
  else
    (math.max(message, value._1), value._1)
// the function for computing messages
def sendMsg(triplet: EdgeTriplet[(Int, Int), Boolean]): Iterator[(VertexId, Int)] = {
 val sourceVertex = triplet.srcAttr
  if (sourceVertex. 1 == sourceVertex. 2)
    Iterator.empty
  else
    Iterator((triplet.dstId, sourceVertex._1))
```



Iterative Computation (8/9)

```
// the function for receiving messages
def vprog(vertexId: VertexId, value: (Int, Int), message: Int): (Int, Int) = {
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    Iterator.empty
  else
    Iterator((triplet.dstId, sourceVertex._1))
```

```
// the function for combining messages
def mergeMsg(msg1: Int, msg2: Int): Int = math.max(msg1, msg2)
```

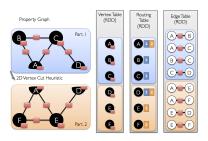
- GraphFrames extends GraphX to provide a DataFrame API.
- ▶ To build a GraphFrame we need to define the vertices and edges as DataFrames.
- spark-shell --packages graphframes:graphframes:0.6.0-spark2.3-s_2.11
 - You may need to delete .ivy2 from your home folder.

```
import org.graphframes._
import org.apache.spark.sql.SQLContext
val sqlContext = new org.apache.spark.sql.SQLContext(sc)
val userDF = sqlContext.createDataFrame(Array(("rxin", "student"), ("jgonzal", "postdoc"),
  ("franklin", "prof"), ("istoica", "prof"))).toDF("id", "role")
val relationshipsDF = sqlContext.createDataFrame(Array(("rxin", "jgonzal", "collab"),
  ("franklin", "rxin", "advisor"), ("istoica", "franklin", "colleague"),
  ("franklin", "franklin", "pi"))).toDF("src", "dst", "relationship")
val graphDF = GraphFrame(userDF, relationshipsDF)
graphDF.edges.where("src = 'franklin'").groupBy("src", "dst").count().show
```



Graph Representation

- Vertex-cut partitioning
- ► Representing graphs using two RDDs: edge-collection and vertex-collection
- ▶ Routing table: a logical map from a vertex id to the set of edge partitions that contains adjacent edges.

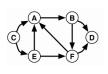




Think Like a Graph



- ▶ Vertex-centric programming model.
 - Operate on a vertex and its edges.
 - Communication to other vertices, via message passing (Pregel), or shared memory (GraphLab).
- ▶ Divide input graphs into partitions.



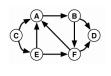
| <u>Partition</u> | <u>Vertex</u> | Edge List |
|------------------|---------------|-----------|
| P1 | A | В |
| | B | DF |
| P2 | © (D) | A E |
| P3 | E F | A F |

Motivation (2/2)

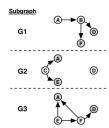
- ▶ In the vertex-centric model, a vertex is very short sighted.
 - A vertex has information about its immediate neighbors.
 - Information is propagated through graphs slowly, one hop at a time.
- ► Graph-centric programming paradigm is proposed to overcome this limitation.

| | Think Like a Vertex | Think Like a Graph |
|---------------|---|---|
| Partition | A collection of vertices | A proper subgraph |
| Computaion | A vertex and its edges | A subgraph |
| Communication | 1-hop at a time, e.g., $\mathtt{A} 	o \mathtt{B} 	o \mathtt{D}$ | Multiple-hops at a time, e.g., $\mathtt{A} 	o \mathtt{D}$ |

Edge List



| P1 | (A) (B) | B D F |
|----|------------|----------|
| P2 | © (D) | A E |
| Р3 | E F | A F |

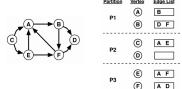


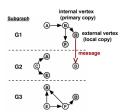


Giraph++



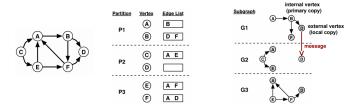
- ► Expose subgraphs to programmers.
- ▶ Internal vertices vs. boundary vertices.





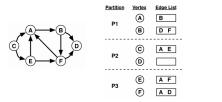


- Expose subgraphs to programmers.
- ► Internal vertices vs. boundary vertices.
 - Information exchange between internal vertices of a partition is immediate.
 - Messages are only sent from boundary vertices of a partition to internal vertices of a different partition.





- Expose subgraphs to programmers.
- Internal vertices vs. boundary vertices.
 - Information exchange between internal vertices of a partition is immediate.
 - Messages are only sent from boundary vertices of a partition to internal vertices of a different partition.
- ▶ A vertex is an internal vertex in exactly one subgraph, but it can be a boundary vertex in zero or more subgraphs.





Execution Model (1/2)

- ▶ A program is executed in sequence of supersteps.
- ► Supersteps are separated by global synchronization barriers.

Execution Model (1/2)

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- ▶ In each superstep, the computation is performed on the whole subgraph in a partition.

Execution Model (1/2)

- ▶ A program is executed in sequence of supersteps.
- ► Supersteps are separated by global synchronization barriers.
- ▶ In each superstep, the computation is performed on the whole subgraph in a partition.
- ▶ Like in Pregel, each internal vertex of a partition has two states: active or inactive.
- ► A boundary vertex does not have any state.

Execution Model (2/2)

▶ Differentiate internal messages and external messages.

Execution Model (2/2)

- ▶ Differentiate internal messages and external messages.
- ▶ What messages can be used in local computation?
 - External messages from previous superstep (global synchronous computation).
 - Internal messages from previous + current superstep (local asynchronous computation).

Execution Model (2/2)

- Differentiate internal messages and external messages.
- ▶ What messages can be used in local computation?
 - External messages from previous superstep (global synchronous computation).
 - Internal messages from previous + current superstep (local asynchronous computation).
- ► This is called hybrid execution model.



Think Like a Matrix

- ► A graph can be represented by an adjacency matrix.
- ▶ Operations on graphs can be performed by algebraic operations on matrices.
- ▶ Linear algebra and matrix theory can be applied to solve graph problems.

Graphs and Matrices (2/2)

- ► Given a graph G = (V, E)
- ► Adjacency matrix A(G), a |V| × |V| matrix

$$\mathtt{A[i][j]} = \begin{cases} 1 & \text{if } i \neq j \text{ and } (v_i, v_j) \in \mathtt{E} \\ 0 & \text{if } i \neq j \text{ and } (v_i, v_j) \not \in \mathtt{E} \\ 0 & \text{if } i = j \end{cases}$$

Adjacency Matrix Example

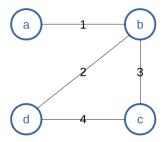
ightharpoonup Produce a vector representing the neighbors of a vertex v_i .



Adjacency Matrix Example

- ▶ Produce a vector representing the neighbors of a vertex v_i.
- ▶ By computing A · x_{vi}
 - $\mathbf{x}_{v_i}[i] = 1$ and all other elements of \mathbf{x}_{v_i} are 0.
- ► For example, to find the neighbors of vertex b

$$\begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 1 & 1 \end{bmatrix}$$





Pegasus



Generalized Iterated Matrix-Vector (GIM-V)

- ► Targets at iterative graph algorithms.
- ► Generalized Iterated Matrix-Vector multiplication (GIM-V)
 - Matrix-vector multiplication
 - Assume M is a $n \times n$ matrix, v is a vector of size n, and $m_{i,j}$ denotes the (i,j) element of M.
 - $\textbf{v} \leftarrow \textbf{M} \cdot \textbf{v}$, where $\textbf{v}_i \leftarrow \sum_{j=1}^n \textbf{m}_{i,j} \textbf{v}_j$.

Generalized Iterated Matrix-Vector (GIM-V)

- ► Targets at iterative graph algorithms.
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 - Matrix-vector multiplication
 - Assume M is a n × n matrix, v is a vector of size n, and m_{i,j} denotes the (i, j) element of M.
 - $\textbf{v} \leftarrow \textbf{M} \cdot \textbf{v}$, where $\textbf{v}_i \leftarrow \sum_{j=1}^n \textbf{m}_{i,j} \textbf{v}_j$.
- ▶ Pegasus models each iteration of the graph computation by a GIM-V operation
 - It is repeated until the vertex values in the vector converge.

GIM-V Operators

 $\blacktriangleright~\textbf{v} \leftarrow \textbf{M} \cdot \textbf{v}, \text{ where } \textbf{v}_i \leftarrow \sum_{j=1}^n \textbf{m}_{i,j} \textbf{v}_j.$

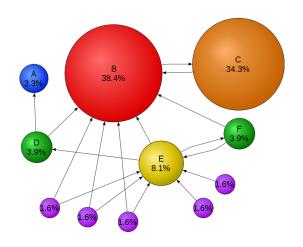
- $\blacktriangleright \ \textbf{v} \leftarrow \textbf{M} \cdot \textbf{v} \text{, where } \textbf{v}_i \leftarrow \textstyle \sum_{j=1}^n \textbf{m}_{i,j} \textbf{v}_j.$
- ▶ combine2(i, j): to combine $m_{i,j}$ and v_j into a value.

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- ► combineAll(i): for each v_i, to combine all the n intermediate results produced by combine2 into a single value.
- ▶ assign: to overwrite the old value of v_i with the new value.

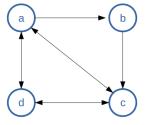


GIM-V Example: PageRank (1/3)



GIM-V Example: PageRank (2/3)

- ▶ PageRank formula: $\mathbf{v} \leftarrow (0.85\mathbf{E}^{\mathrm{T}} + 0.15\mathbf{U}) \cdot \mathbf{v}$.
 - v is a column vector with n elements.
 - E is a is the row-normalized adjacency matrix.
 - U is a n × n matrix, with all elements set to $\frac{1}{n}$.



$$\mathbf{A} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix} \mathbf{E} = \begin{bmatrix} 0 & 1/3 & 1/3 & 1/3 \\ 0 & 0 & 0 & 1 \\ 1/2 & 0 & 0 & 1/2 \\ 1/2 & 0 & 1/2 & 0 \end{bmatrix} \mathbf{U} = \begin{bmatrix} 1/4 & 1/4 & 1/4 & 1/4 \\ 1/4 & 1/4 & 1/4 & 1/4 \\ 1/4 & 1/4 & 1/4 & 1/4 \end{bmatrix} \mathbf{v}_{\text{init}} = \begin{bmatrix} 1/4 \\ 1/4 \\ 1/4 \\ 1/4 \end{bmatrix}$$

▶ If $M = 0.85E^T + 0.15U$, then we can write the PageRank as $\mathbf{v} \leftarrow \mathbf{M} \cdot \mathbf{v}$.

GIM-V Example: PageRank (3/3)

- ▶ PageRank formula: $\mathbf{v} \leftarrow (0.85\mathbf{E}^{\mathrm{T}} + 0.15\mathbf{U}) \cdot \mathbf{v}$.
- ▶ combine2(i,j) = $0.85 \times m_{i,j} \times v_j$
- combineAll(i) = $\frac{0.15}{n} + \sum_{j=1}^{n} combine2(i, j)$
- ▶ assign: $v_i \leftarrow combineAll(i)$



Summary



- ► Think like a table
 - Graphx: unifies data-parallel and graph-parallel systems.
- ► Think like a graph
 - Giraph++: exposes subgraphs to programmers
- ► Think like a matrix
 - Pegasus: linear algebra and matrix theory to solve graph problems.

References

- ▶ J. Gonzalez et al., "GraphX: Graph Processing in a Distributed Dataflow Framework", OSDI 2014
- ▶ Y. Tian et al., "From think like a vertex to think like a graph", VLDB 2013
- ▶ U. Kang et al., "PEGASUS: mining peta-scale graphs", Knowledge and information systems 2011



Questions?