

Large Scale Graph Processing - X-Stream and GraphX

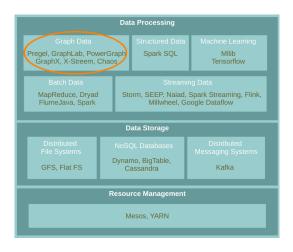
Amir H. Payberah payberah@kth.se 03/10/2018



https://id2221kth.github.io



Where Are We?













- ▶ Difficult to extract parallelism based on partitioning of the data.
- ▶ Difficult to express parallelism based on partitioning of computation.



Think Like an Edge

Could we compute big graphs on a single machine?



► Vertex-centric gather-scatter: iterates over vertices

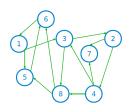
```
Until convergence {
    // the scatter phase
    for all vertices v that need to scatter updates
        send updates over outgoing edges of v

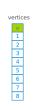
// the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v

}
```



Vertex-Centric Breadth First Search (1/5)





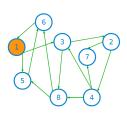


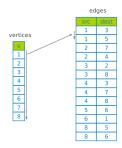
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    // the gather phase
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}
```



Vertex-Centric Breadth First Search (2/5)



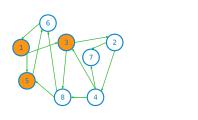


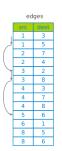
```
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   // the gather phase
   for all vertices v that have updates
      apply updates from inbound edges of v
}
```



Vertex-Centric Breadth First Search (3/5)





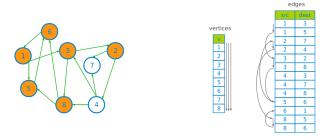
vertices

```
Until convergence {
    // the scatter phase
    for all vertices v that need to scatter updates
        send updates over outgoing edges of v

    // the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v
}
```



Vertex-Centric Breadth First Search (4/5)

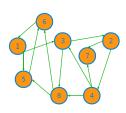


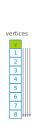
```
Until convergence {
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    for all vertices v that need to scatter updates
        send updates over outgoing edges of v

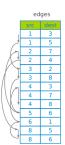
    // the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v
}
```



Vertex-Centric Breadth First Search (5/5)







```
Until convergence {
    // the scatter phase
    for all vertices v that need to scatter updates
        send updates over outgoing edges of v

    // the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v
}
```



X-Stream



- ► Could we process massive graphs on a single machine?
- ► X-Stream makes graph edges accesses sequential.
- ► Edge-centric scatter-gather model.



- Disk-based processing
 - Graph traversal = random access
 - Random access is inefficient for storage

Medium	Read (MB/s)		Write (MB/s)	
	Random	Sequential	Random	Sequential
RAM	567	2605	1057	2248
SSD	22.64	355	49.16	298
Disk	0.61	174	1.27	170

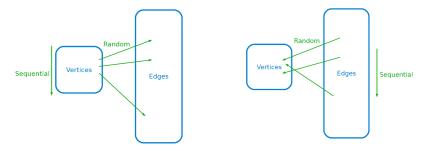
Note: 64 byte cachelines, 4K blocks (disk random), 16M chunks (disk sequential)

Eiko Y., and Roy A., "Scale-up Graph Processing: A Storage-centric View", 2013.



Vertex-Centric vs. Edge-Centric Programming Model (1/2)

- ► Vertex-centric gather-scatter: iterates over vertices
- ► Edge-centric gather-scatter: iterates over edges





Vertex-Centric vs. Edge-Centric Programming Model (2/2)

```
Until convergence {
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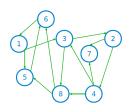
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        apply updates from inbound edges of v
}
```

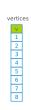
```
Until convergence {
    // the scatter phase
    for all edges e
        send update over e

    // the gather phase
    for all edgaes e that have updates
        apply update to e.destination
}
```



Vertex-Centric Breadth First Search (1/5)





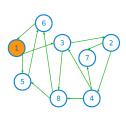


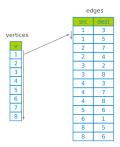
```
Until convergence {
    // the scatter phase
    for all vertices v that need to scatter updates
        send updates over outgoing edges of v

    // the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v
}
```



Vertex-Centric Breadth First Search (2/5)



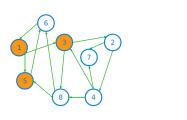


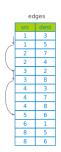
```
Until convergence {
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    for all vertices v that need to scatter updates
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    // the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v
}
```



Vertex-Centric Breadth First Search (3/5)





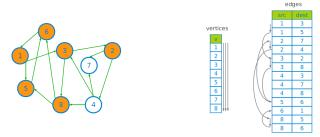
vertices

```
Until convergence {
    // the scatter phase
    for all vertices v that need to scatter updates
        send updates over outgoing edges of v

    // the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v
}
```



Vertex-Centric Breadth First Search (4/5)

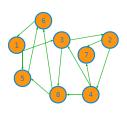


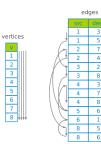
```
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    // the scatter phase
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        send updates over outgoing edges of v

    // the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v
}
```



Vertex-Centric Breadth First Search (5/5)



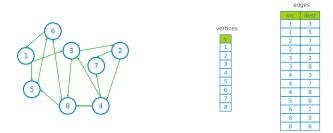


```
Until convergence {
    // the scatter phase
    for all vertices v that need to scatter updates
        send updates over outgoing edges of v

    // the gather phase
    for all vertices v that have updates
        apply updates from inbound edges of v
}
```



Edge-Centric Breadth First Search (1/5)

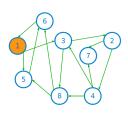


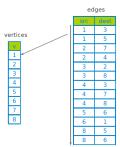
```
Until convergence {
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    // the gather phase
    for all edgaes e that have updates
        apply update to e.destination
}
```



Edge-Centric Breadth First Search (2/5)



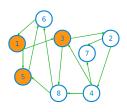


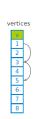
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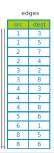
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```



Edge-Centric Breadth First Search (3/5)





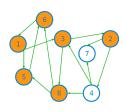


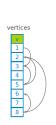
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Until convergence {
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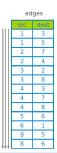
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        apply update to e.destination
}
```



Edge-Centric Breadth First Search (4/5)





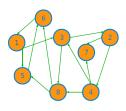


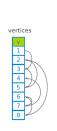
```
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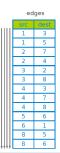
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Edge-Centric Breadth First Search (5/5)







```
Until convergence {
    // the scatter phase
    for all edges e
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    for all edgaes e that have updates
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}
```



Vertex-Centric vs. Edge-Centric Tradeoff

► Vertex-centric scatter-gather: EdgeData
RandomAccessBandwidth

► Edge-centric scatter-gather: Scatters×EdgeData SequentialAccessBandwidth

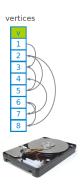
► Sequential Access Bandwidth ≫ Random Access Bandwidth.

► Few scatter gather iterations for real world graphs.



Streaming Partitions (1/4)

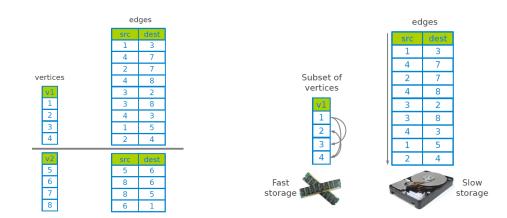
▶ Problem: still have random access to vertex set.



Solution
Partition the graph into streaming partitions.



Streaming Partitions (2/4)



Streaming Partitions (3/4)

- ▶ A streaming partition consists of: a vertex set, an edge list, and an update list.
- ► The vertex set: a subset of the vertex set of the graph that fits into the memory.
 - Vertex sets are mutually disjoint.
 - Their union equals the vertex set of the entire graph.
- ▶ The edge list: all edges whose source vertex is in the partition's vertex set.
- ► The update list: all updates whose destination vertex is in the partition's vertex set.

Streaming Partitions (4/4)

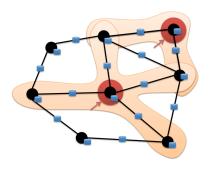
```
// Scatter phase:
for each streaming_partition p
   read in vertex set of p
   for each edge e in edge list of p
        append update to Uout
// shuffle phase:
for each update u in Uout
    p = partition containing target of u
   append u to Uin(p)
destroy Uout
//gather phase:
for each streaming_partition p
   read in vertex set of p
   for each update u in Uin(p)
        edge_gather(u)
   destroy Uin(p)
```



Think Like a Table

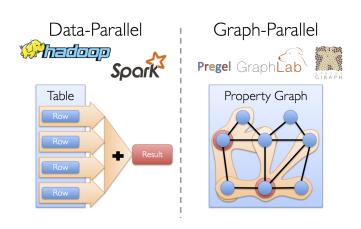


Graph-Parallel Processing Model





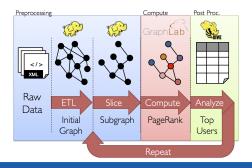
Data-Parallel vs. Graph-Parallel Computation





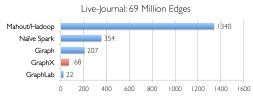
Motivation (2/3)

- Graph-parallel computation: restricting the types of computation to achieve performance.
- ► The same restrictions make it difficult and inefficient to express many stages in a typical graph-analytics pipeline.

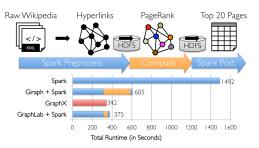




Motivation (3/3)

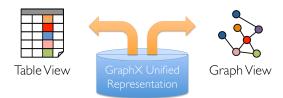








- ▶ Unifies data-parallel and graph-parallel systems.
- ▶ Tables and Graphs are composable views of the same physical data.





GraphX



- ► GraphX is the library to perform graph-parallel processing in Spark.
- ► In-memory caching.
- ► Lineage-based fault tolerance.





The Property Graph Data Model

- ► Spark represent graph structured data as a property graph.
- ▶ It is logically represented as a pair of vertex and edge property collections.





Vertex Table		
ld	Property (V)	
3	(rxin, student)	

3	(rxin, student)
7	(jgonzal, postdoc)
5	(franklin, professor)
2	(istoica, professor)

Edge Table

SrcId	Dstld	Property (E)
3	7	Collaborator
5	3	Advisor
2	5	Colleague
5	7	PI

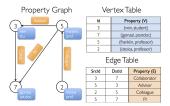


The Vertex Collection

▶ VertexRDD: contains the vertex properties keyed by the vertex ID.

```
class Graph[VD, ED] {
  val vertices: VertexRDD[VD]
  val edges: EdgeRDD[ED]
}

// VD: the type of the vertex attribute
abstract class VertexRDD[VD] extends RDD[(VertexId, VD)]
```





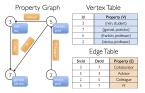


The Edge Collection

► EdgeRDD: contains the edge properties keyed by the source and destination vertex IDs.

```
class Graph[VD, ED] {
  val vertices: VertexRDD[VD]
  val edges: EdgeRDD[ED]
}

// ED: the type of the edge attribute
case class Edge[ED](srcId: VertexId, dstId: VertexId, attr: ED)
abstract class EdgeRDD[ED] extends RDD[Edge[ED]]
```







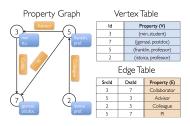
The Triplet Collection

- ► The triplets collection consists of each edge and its corresponding source and destination vertex properties.
- ▶ It logically joins the vertex and edge properties: RDD [EdgeTriplet[VD, ED]].
- ► The EdgeTriplet class extends the Edge class by adding the srcAttr and dstAttr members, which contain the source and destination properties respectively.





Building a Property Graph





Graph Operators

- ► Information about the graph
- Property operators
- Structural operators
- ► Joins
- ► Aggregation
- ► Iterative computation
- **.**..



Information About The Graph (1/2)

► Information about the graph

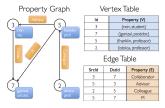
```
val numEdges: Long
val numVertices: Long
val inDegrees: VertexRDD[Int]
val outDegrees: VertexRDD[Int]
val degrees: VertexRDD[Int]
```

Views of the graph as collections

```
val vertices: VertexRDD[VD]
val edges: EdgeRDD[ED]
val triplets: RDD[EdgeTriplet[VD, ED]]
```



Information About The Graph (2/2)



```
// Constructed from above
val graph: Graph[(String, String), String]

// Count all users which are postdocs
graph.vertices.filter { case (id, (name, pos)) => pos == "postdoc" }.count

// Count all the edges where src > dst
graph.edges.filter(e => e.srcId > e.dstId).count
```



Property Operators

- ► Transform vertex and edge attributes
- ► Each of these operators yields a new graph with the vertex or edge properties modified by the user defined map function.

- ▶ reverse returns a new graph with all the edge directions reversed.
- subgraph takes vertex/edge predicates and returns the graph containing only the vertices/edges that satisfy the given predicate.

- ▶ joinVertices joins the vertices with the input RDD.
 - Returns a new graph with the vertex properties obtained by applying the user defined map function to the result of the joined vertices.
 - Vertices without a matching value in the RDD retain their original value.

```
def joinVertices[U](table: RDD[(VertexId, U)])(map: (VertexId, VD, U) => VD): Graph[VD, ED]

val rdd: RDD[(VertexId, String)] = sc.parallelize(Array((3L, "phd")))

val joinedGraph = graph.joinVertices(rdd)((id, user, role) => (user._1, role + " " + user._2))

joinedGraph.vertices.collect.foreach(println)
```

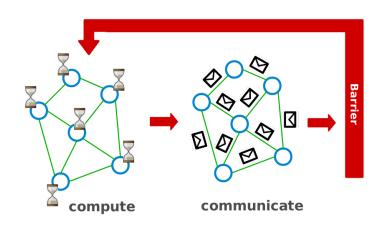
aggregateMessages applies a user defined sendMsg function to each edge triplet in the graph and then uses the mergeMsg function to aggregate those messages at their destination vertex.

```
def aggregateMessages[Msg: ClassTag](
   sendMsg: EdgeContext[VD, ED, Msg] => Unit, // map
   mergeMsg: (Msg, Msg) => Msg, // reduce
   tripletFields: TripletFields = TripletFields.All):
   VertexRDD[Msg]
```

```
// count and list the name of friends of each user
val profs: VertexRDD[(Int, String)] = validUserGraph.aggregateMessages[(Int, String)](
    // map
    triplet => {
        triplet.sendToDst((1, triplet.srcAttr._1))
    },
    // reduce
    (a, b) => (a._1 + b._1, a._2 + " " + b._2)
)
profs.collect.foreach(println)
```



Iterative Computation (1/9)



Iterative Computation (2/9)

```
i_val := val

for each message m
   if m > val then val := m

if i_val == val then
   vote_to_halt
else
   for each neighbor v
        send_message(v, val)
```



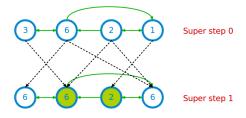
Super step 0

Iterative Computation (3/9)

```
i_val := val

for each message m
  if m > val then val := m

if i_val == val then
  vote_to_halt
else
  for each neighbor v
    send_message(v, val)
```

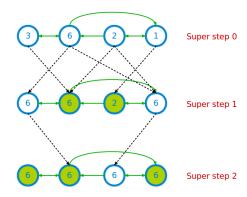


Iterative Computation (4/9)

```
i_val := val

for each message m
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    send_message(v, val)
```



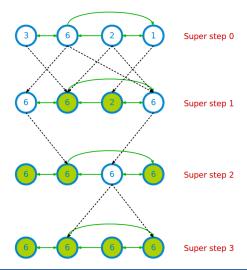


Iterative Computation (5/9)

```
i_val := val

for each message m
   if m > val then val := m

if i_val == val then
   vote_to_halt
else
   for each neighbor v
        send_message(v, val)
```





Iterative Computation (6/9)

- pregel takes two argument lists: graph.pregel(list1)(list2).
- ► The first list contains configuration parameters
 - The initial message, the maximum number of iterations, and the edge direction in which to send messages.
- ▶ The second list contains the user defined functions.
 - Gather: mergeMsg, Apply: vprog, Scatter: sendMsg

```
def pregel[A]
  (initialMsg: A, maxIter: Int = Int.MaxValue, activeDir: EdgeDirection = EdgeDirection.Out)
  (vprog: (VertexId, VD, A) => VD, sendMsg: EdgeTriplet[VD, ED] => Iterator[(VertexId, A)],
    mergeMsg: (A, A) => A):
  Graph[VD, ED]
```

Iterative Computation (7/9)

```
import org.apache.spark._
import org.apache.spark.graphx._
import org.apache.spark.rdd.RDD

val initialMsg = -9999

val vertices: RDD[(VertexId, (Int, Int))] = sc.parallelize(Array((1L, (1, -1)), (2L, (2, -1)), (3L, (3, -1)), (6L, (6, -1))))

val relationships: RDD[Edge[Boolean]] = sc.parallelize(Array(Edge(1L, 2L, true), Edge(2L, 1L, true), Edge(2L, 6L, true), Edge(3L, 6L, true), Edge(6L, 3L, true)))

val graph = Graph(vertices, relationships)
```





Iterative Computation (8/9)

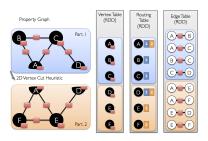
```
// Gather: the function for combining messages
def mergeMsg(msg1: Int, msg2: Int): Int = math.max(msg1, msg2)
// Apply: the function for receiving messages
def vprog(vertexId: VertexId, value: (Int, Int), message: Int): (Int, Int) = {
 if (message == initialMsg)
   value
  else
    (math.max(message, value._1), value._1)
// Scatter: the function for computing messages
def sendMsg(triplet: EdgeTriplet[(Int, Int), Boolean]): Iterator[(VertexId, Int)] = {
 val sourceVertex = triplet.srcAttr
  if (sourceVertex. 1 == sourceVertex. 2)
   Iterator.empty
  else
    Iterator((triplet.dstId, sourceVertex._1))
```

Iterative Computation (9/9)



Graph Representation

- Vertex-cut partitioning
- ► Representing graphs using two RDDs: edge-collection and vertex-collection
- ▶ Routing table: a logical map from a vertex id to the set of edge partitions that contains adjacent edges.





Summary



- ► Think like an edge
 - XStream: edge-centric GAS, streaming partition
- ► Think like a table
 - Graphx: unifies data-parallel and graph-parallel systems.

References

- ► A. Roy et al., "X-stream: Edge-centric graph processing using streaming partitions", ACM SOSP 2013.
- ▶ J. Gonzalez et al., "GraphX: Graph Processing in a Distributed Dataflow Framework", OSDI 2014



Questions?