

# Loop Invariant Generation through Active Learning

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**Abstract.** Loop invariant generation is one of the fundamental problems in program analysis and verification. In this work, we propose an automatic generation method for inductive loop invariants through iterations of runtime sampling, machine learning and constraint solving. In each iteration, our method first collects real data at runtime based on selective sampling. Then, based on their satisfaction of the assumptions and assertions in the program, our method uses support vector machine to learn a loop invariant candidate, i.e., a conjunction of linear and polynomial constraints. Finally, if the candidate can be verified as the inductive loop invariant of the program, our method returns it directly and ends the generation process. Otherwise, the counter-example is used to refine the data sampling in the next iteration. The experiment evaluation shows that our method can be used to learn loop invariants effectively and automatically that cannot be learned by other loop invariant generation tools.

## 1 Introduction

In this work, we propose a method and a tool called ZILU for automatically learning loop invariants. ZILU works through an iterative process combining sampling, active learning and verification.

## 2 Problem Definition and Solution Overview

In the following, we assume that a program contains a finite set of integer variable  $\{x, y, z, \dots\}$  and thus a program state is a valuation of the variables. A predicate on the variables is viewed as the maximum set of program states which satisfies the predicate. We use predicates and sets of program states interchangeably. Without loss of generality, we assume the input to ZILU is a Hoare triple

$$\begin{array}{l} \{Pre\} \\ while\ (Cond)\{ \\ \quad Body \\ \} \\ \{Post\} \end{array}$$

where *Pre* is the pre-condition, which should be satisfied before entering the loop; *Cond* is the loop guard condition, which is the only way to enter or exit the loop; *Body* is

```

void ex1 (int x) {
    int y = 355;
    if (x > 46) x = 46;
    while (x <= 100) {
        if (x >= 46) {
            y = y+1;
        }
        x = x + 1;
    }
    assert (y==409);
}

```

**Fig. 1.** An example adopted from [?]

```

void ex2 () {
    lock=0;new=old+1;
    while (new!=old) {
        lock=1;old=new;
        if (foo(new)) {
            lock=0;new++;
        }
    }
    if (lock==0)
        error();
}

```

**Fig. 2.** An example adopted from [?]

the loop body, in which we assume there is no *break* or *goto* statement which can jump out of the loop without checking *cond*; and *Post* is the post-condition, which should be satisfied after the loop.

For simplicity, we assume that *Body* is a function such that  $Body(s) = s'$  means that starting at a program state  $s$ , executing *Body* would result in a program state  $s'$ . Furthermore, we write  $Body(Pr)$  where  $Pr$  to denote the set  $\{s' \mid \exists s \in Pr : Body(s) = s'\}$ . The goal is thus to automatically obtain a loop invariant such that the following conditions are satisfied.

$$\begin{aligned}
 Pre &\Rightarrow Inv & (1) \\
 Inv &\Rightarrow Body(Inv \wedge Cond) & (2) \\
 Inv \wedge \neg Cond &\Rightarrow Post & (3)
 \end{aligned}$$

*Example 1.* We use the two examples shown in Figure 1 and 2 to illustrate how our approach works. In *ex1*, the precondition of the loop is  $y = 355 \wedge x \leq 46$  and the post-condition is  $y = 409$ . In *ex2*, the precondition is that  $lock = 0 \wedge new = old + 1$  and the post-condition (necessary so that there is no error) is  $lock = 1$ . We remark that  $foo(new)$  is an external function which *deterministically* returns either true or false, i.e., it returns true if *new* is even; otherwise, it returns false. We will discuss in Section how our approach would work if  $foo(new)$  is non-deterministic.

*Problem Definition* In this work, we assume that given the Hoare triple, there is either a counterexample (i.e., a program state  $s$  such that  $s \in Pre$  and executing the program from  $s$  results in failing *post*) or there exists an invariant satisfying (1) and (2) and (3). Furthermore, the invariant *inv* is a boolean formula over a linear inequality constraint of the form  $ax + bx + \dots \geq d$  where  $a, b, d$  are bounded integer constants; and *inv* contains no more than  $k$  such statements. We remark that such invariant is in general not convex and thus existing approaches on learning convex invariants do not work [].

*Overview of Our Approach* Our approach to solve the problem is illustrated in Figure ???. Firstly, we randomly generate a set of program states right before the loop and test the program. Based on the testing results, we obtain program states which must or must not satisfy any invariant satisfying (1), (2) and (3). Secondly, we develop an algorithm for generating candidate invariants based classification techniques from the

machine learning community. Thirdly, to overcome the limitation of the sampled program states, we adopt active learning techniques, in particular, selective sampling, to refine the candidate invariants. Lastly, we rely on constraint solving techniques to check whether the generate invariant satisfies (1) and (2) and (3). If it does, we report that our approach is successful; otherwise, using the counterexamples generated by the constraint solvers, we repeat from the second step. In this following sections, we present details of each step.

### 3 Sampling

In this step, we sample, either randomly or using tools based on the idea of concolic testing [], a set  $T$  of program states and test the program starting with each program state  $s$  in  $T$ . We write  $Body^*(s)$  to denote the set of program states which could be reached after executing zero or more iterations of the loop starting from  $s$ . We write  $Body^*(T)$  to denote  $\{s' \mid \exists s \in T \cdot s' \in Body^*(s)\}$ . Furthermore, we write  $s \Rightarrow s'$  to denote that starting with a program state  $s$  would result in state  $s'$  when the loop terminates.

We categorize program states in  $Body^*(T)$  into four sets:  $C_T$  which stands for counter-example trace,  $P_T$  which stands for traces with positive labels,  $N_T$  which stands for traces with negative labels and  $U_T$  which stands for traces with unknown labels. They can be judged by the following rules:

- Set  $C_T$  is  $\{s \in Body^*(T) \mid s \in Pre \wedge s \Rightarrow s' \wedge s' \notin Post\}$ ;
- Set  $P_T$  is  $\{s \in Body^*(T) \mid s \in Pre \wedge s \Rightarrow s' \wedge s' \in Post\}$ ;
- Set  $N_T$  is  $\{s \in Body^*(T) \mid s \notin Pre \wedge s \Rightarrow s' \wedge s' \notin Post\}$ ;
- Set  $U_T$  is  $\{s \in Body^*(T) \mid s \notin Pre \wedge s \Rightarrow s' \wedge s' \in Post\}$ ;

We remark that anytime a program state in  $C_T$  is identified, a counter-example is found and ZILU reports that verification is failed immediately. Otherwise, because  $Inv$  must satisfy (1),(2) and (3), we know that  $P_T \subseteq Inv$  and  $N_T \cap Inv = \emptyset$ . The program states in  $U_T$  may or not may be in  $Inv$ . If we know that a program state  $s \in U_T$  is in  $Inv$ ,  $Body^*(s) \subseteq Inv$ .

Due to the limited set of samples we have (which is often referred to as labeled samples in the machine learning community), the classifier obtained above might be far from being correct. In fact, without labeled samples which are right on the boundary of the ‘actual’ classifier, it is very unlikely that we would find it. Intuitively, in order to get the ‘actual’ classifier, we would require samples which would distinguish the actual one from any nearby one. This problem has been discussed and addressed in the machine learning community using active learning and selective sampling [?].

The concept of active learning or selective sampling refers to the approaches that aim at reducing the labeling effort by selecting only the most informative samples to be labeled. SVM selective sampling techniques have been proven effective in achieving a high accuracy with fewer examples in many applications [?,?]. The basic idea of selective sampling is that at each round, we select the samples that are the closest to the classification boundary so that they are the most difficult to classify and the most informative to label. Since an SVM classification function is represented by support

**Algorithm 1:** Algorithm *activeLearning***Input:**  $F^+$  and  $F^-$ **Output:** a classifier for  $F^+$  and  $F^-$ 

```

1 let old be null;
2 while true do
3   let  $f = \text{classify}(F^+, F^-)$ ;
4   if  $f$  is identical to old then
5     return  $f$ ;
6   let  $old = f$ ;
7   let sam be a set of samples computed by selective sampling;
8   test the program and update  $F^+$  and  $F^-$  accordingly;
```

vectors which are the samples closest to the boundary, this selective sampling effectively learns an accurate function with fewer labeled data [?]. In our setting, this means that we should sample a program state right by the classifier and test the program with that state to label that feature vector so that the classifier would be improved.

Algorithm 1 presents details on how active learning is implemented in ZILU. At line 2, we obtain a classifier based on Algorithm ???. We compare the newly obtained classifier with the previous one at line 4, if they are identical, we return the classifier; otherwise we apply selective sampling so that we can generate additional labeled samples for improving the classifier. In particular, at line 5, we apply standard techniques [?] to select the most informative sample. Notice that in our setting, the most informative samples are those which are exactly on the lines and therefore can be obtained by solving an equation system. At line 8, we test the program with the newly generated samples so as to label them accordingly.

*Example 2.*

**Proposition 1.** *Algorithm activeLearning always eventually terminates.* □

*Example 3.*

## 4 Classification

After sampling and labeling, we obtain some program states must be in *inv* and some must not. Thus, any candidate invariant must be able to perfectly classify these states. We apply classification techniques from the machine learning community to obtain classifiers as candidate invariants. Due to different technique can results in different forms of classifiers, we apply SVM and some of its derivatives on our training data.

**Algorithm 2: Algorithm overall****Input:** *Pre, Cond, Body, Post***Output:** an invariant which completes the proof or a counterexample

```

1  let  $T$  be a set of random samples;
2  while true do
3    test the program for each sample in  $T$ ;
4    if a state  $s$  in  $CT$  is identified then
5      return  $s$  as a counterexample;
6    let  $P, N$  and  $NP$  be the respective sets accordingly;
7    let  $Inv_u = \text{activeLearning}(P, N \cup NP)$ ;
8    let  $Inv_o = \text{activeLearning}(P \cup NP, N)$ ;
9    let  $Inv_s = \text{activeLearning}(P, N)$ ;
10   for each  $Inv$  in  $\{Inv_u, Inv_o, Inv_s\}$  do
11     if (1) or (2) or (3) is not satisfied then
12       add the counterexample into  $T$ ;
13     else
14       return  $Inv$  as the proof;

```

**4.1 Linear SVM****4.2 Polynomial SVM****4.3 Conjunctive SVM**

In the following, we present how we obtain a classifier automatically using SVM. SVM is a supervised machine learning algorithm for classification and regression analysis. We use its binary classification functionality. Mathematically, the binary classification functionality of (linear) SVM works as follows. Given two sets of feature vectors  $F^+$  and  $F^-$ , it generates, if there is any, a linear constraint in the form of  $ax + by + \dots \geq d$  where  $x$  and  $y$  are feature values and  $a, b, d$  are constants, such that every state  $s \in F^+$  satisfies the constraint and every state  $s' \in F^-$  fails the constraint. In this work, we always choose the *optimal margin classifier* (see the definition in [?]) if possible. This half space could be seen as the strongest witness why the two data states are different. In the following, we write  $svm(F^+, F^-)$  to denote the function which returns a linear classifier

If, however,  $F^+$  and  $F^-$  cannot be perfectly classified by one half space only, a more complicated function  $f$  must be adopted. For instance, if there is a classifier in the form of conjunctive of multiple half spaces, the algorithm presented in [?] can be used to identify such a classifier.

**5 Verification**

Given a learned predicate  $Inv$ , we verify whether constraint (1), (2) and (3) are satisfied using symbolic execution. If all of them are satisfied, we successfully verify the program. Otherwise, if any of them is violated, the counterexample obtained is added to the

set of sample  $X$ , which is then tested, categorized, used for active learning accordingly. The overall algorithm is presented in Figure 2.

We remark that we learn three classifiers as candidates for the loop invariant:  $U$ ,  $OU$ ,  $O$  such that

- $U$  classifies states in  $P$  and those in  $N \cup NP$ .
- $O$  classifies states in  $N$  and those in  $P \cup NP$ .
- $OU$  classifies states in  $P$  and  $N$ ;

Intuitively,  $U$  would be an under-approximation of  $Inv$  (by assuming states in  $NP$  does not satisfy  $Inv$ );  $O$  would be an over-approximation of  $Inv$  (by assuming states in  $NP$  does satisfy  $Inv$ ); and  $OU$  would be an safe-approximation of  $Inv$  (by using states which we are certain whether they are in  $Inv$  or not).

*Example 4.*

**Theorem 1.** *Algorithm overall always eventually terminates and it is correct.* □

## 6 Experiments

## 7 Related Work

## 8 Conclusion

*Limitation and Potential Remedies* We remark that in theory, we could learn non-linear classifier using methods like SVM with kernel methods []. Nonetheless, due to the limitation of proving capability and tools with regards to non-linear constraints, we leave those to our future work. Furthermore, we assume there is a bound  $k$  on the number of clauses in the variant. In practice, we would expect (refer to empirical evidence in Section 6) often  $k$  is of a small value.

## References