

Midsem

(Monsoon 2021-22)

Q1. In ancient times, people used parrots to make predictions. Let us consider a parrot that we are using predict the result of a badminton match between player 1 and player 2. Once we know the outcome of the match, we make a logic function (F) that outputs “1” when the parrot was right and “0” when the parrot was wrong. Make a truth-table for this logic function. Assuming all the possible outcomes in the truth-table are equally likely, what is the probability that our parrot predicted correctly ? Now, let us say we are unhappy with the probability of correct prediction, and want to increase our chances of success. For that, we now employ three parrots and create a function (G) that outputs “1” when one or more of them was right, and “0” only when all of them were wrong simultaneously. Make a truth-table, a K-map and the simplest logic circuit for this function.

A1.

- We go from logic problem -> function I/O -> K-map -> expression -> circuit
- Output of the function is given as F/G. What are the inputs ?
- Match the result and parrot's prediction
- The truth-table will be:

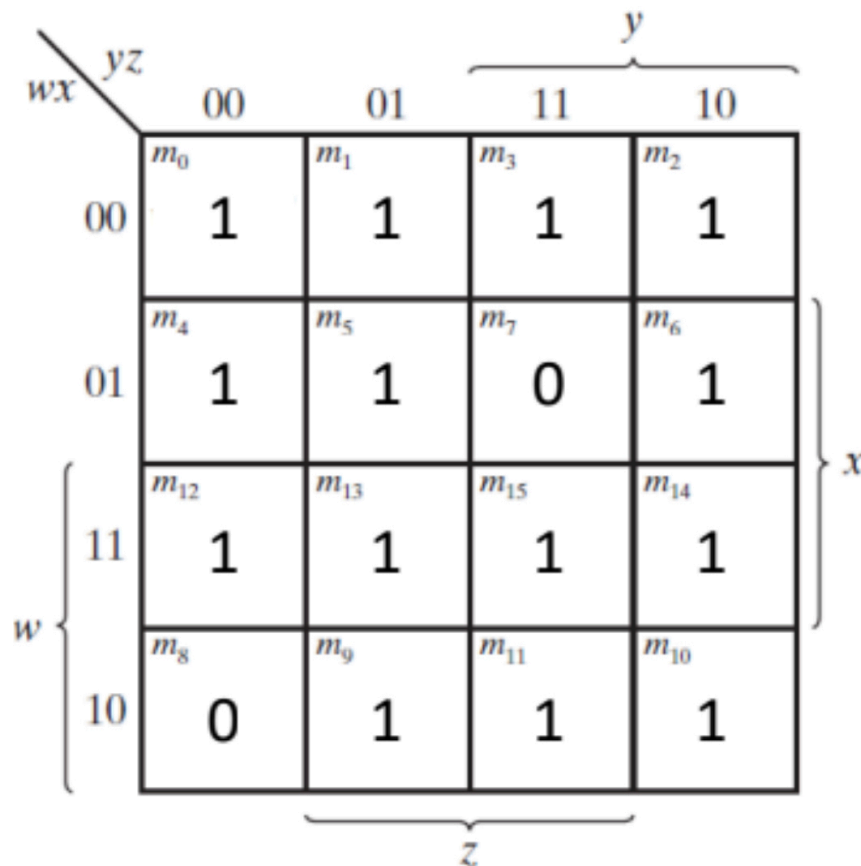
R	P	F
P1	P1	1
P1	P2	0
P2	P1	0
P2	P2	1

- We can take P1 as 0 and P2 as 1 to make it resemble a normal truth-table
- For three parrots:

R	P1	P2	P3	G
0	0	0	0	1
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	1
0	1	0	1	1
0	1	1	0	1
0	1	1	1	0
1	0	0	0	0
1	0	0	1	1
1	0	1	0	1

R	P1	P2	P3	G
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

- “Correctness” of the parrot is a derived variable from the result of the match and the prediction of the parrot
- The K-map for the function G



- No clusters of 16 or 8
- Clusters of 4 ?
- Total 12: 4 vertical/horizontal, 8 squares
- No essential prime implicants
- The function can be written as $G = wz + xy + w'x' + yz'$
- The function can also be written as $G = (w + x' + y' + z')(w' + x + y + z)$

Q2. Let us define a set $B = \{1, p, q, pq\}$, where p and q are distinct prime natural numbers. Let us define two binary operators on the elements of this set as Least Common Multiple denoted by $(a \# b)$, and Greatest Common Factor denoted by $(a * b)$. Are the operations $\#$ and $*$ closed on the set B ? Do p and q need to be prime for the closure to work ? Do the operations $\#$ and $*$ have an identity element, if so, what are they ?

A2.

- We can make a table with all the inputs and see what the outputs are
- P and q should be co-prime for the closure to work
- Identity in this case is 1 and pq for LCM and GCD
- In general case, it is important to define which set you are working on - because operators are defined on sets
- Identity for LCM = 1. Identity for GCD on Integer set is 0; on natural numbers is infinity/not defined

Q3. We can design the 2's complement circuit using adder circuits. But is that the most optimum way? Design an optimal logic circuit for obtaining the 2's complement of a 4-bit binary number (ABCD). Count the number of transistors used and compare the number with the adder implementation.

A3.

- Yes a more optimum circuit exists
- Use K-maps to make 2-level circuit
- The number of transistors can be calculated using the following:
 1. For an AND Gate no of transistors = $2n$ where n is the number of inputs
 2. For a NAND Gate no of transistors = $2n+2$ where n is the number of inputs

Q4. Do the following conversions

1. $(93.25)_{10} = ()_8$
2. $(1011)_{10} = ()_{16}$
3. $(CAD.004)_{16} = ()_8$
4. $(123)_4 = ()_{16}$
4. $(3.3)_{10} = ()_2$

A4.

1.

8	93	5
8	11	3
8	1	1
	0	

Taking the integral part $(93)_{10} = (135)_8$

Taking the decimal part

0.25	2.0	2
0	-	-

$(0.25)_{10} = (0.2)_8$

$\therefore (93.25)_{10} = (135.2)_8$

2.

16	1011	3
16	63	15 (F)
16	3	3
	0	

$\therefore (1011)_{10} = (3F3)_{16}$

3. Path => 16 -> 2 -> 8

	C	A	D	.	0	0	4	
	1100	1010	1101	.	0000	0000	0100	
110	010	101	101	.	000	000	000	100
6	2	5	5	.	0	0	0	4

$$\therefore (\text{CAD.004})_{16} = (6255.0004)_8$$

4. Path => 4 -> 2 -> 16

1	2	3
01	10	11
0001	1011	
1	B	

$$\therefore (123)_4 = (1B)_{16}$$

5. Taking integral part

2	3	1
2	1	1
2	0	

$$\text{Taking integral part } (3)_{10} = (11)_2$$

Taking decimal part

0.3	0.6	0
0.6	1.2	1
0.2	0.4	0
0.4	0.8	0
0.8	1.6	1
0.6	1.2	1
0.2		

Repetition

$$(0.3)_{10} = (0.\overline{1001})_2$$

$$\therefore (3.3)_{10} = (11.0\overline{1001})_2$$

