

# Discrete Mathematics: Lecture 13

- Last time:
  - Chap 4.2: Integer representations and algorithms
  - Chap 4.3: Primes and greatest common divisors
- Today:
  - Chap 4.4: Solving congruences
  - Chap 4.6: Cryptography (a simple introduction)
- Next time:
  - Chap 5.1: Mathematical induction
  - Chap 5.2: Strong induction and well-ordering

# Review of last time

- The fundamental theorem of arithmetic
- There are infinitely many primes
- Greatest common divisors, least common multiples
- Euclidean algorithm
- Base  $b$  representation of integers
- Algorithms for addition and multiplication of integers
- Modular exponentiation algorithm

# Some useful results

- Theorem: If  $a$  and  $b$  are positive integers, then there exist integers  $s$  and  $t$  such that  $\gcd(a, b) = sa + tb$ .
- Example:  $\gcd(252, 198) = 18$
- Lemma 1: If  $a, b$ , and  $c$  are positive integers such that  $\gcd(a, b) = 1$  and  $a \mid bc$ , then  $a \mid c$ .
- Lemma 2: If  $p$  is a prime and  $p \mid a_1 a_2 \dots a_n$ , where each  $a_i$  is an integer, then  $p \mid a_i$  for some  $i$ .
- The uniqueness of prime factorization of integers
- Theorem: Let  $m$  be a positive integer and let  $a, b$ , and  $c$  be integers. If  $ac \equiv bc \pmod{m}$  and  $\gcd(c, m) = 1$ , then  $a \equiv b \pmod{m}$ .

# Linear congruences

- A congruence of the form  $ax \equiv b \pmod{m}$  is called a linear congruence. How to solve linear congruences?
- An integer  $\bar{a}$  such that  $\bar{a}a \equiv 1 \pmod{m}$  is called an inverse of  $a$  modulo  $m$ .
- Theorem: If  $m > 1$  and  $\gcd(a, m) = 1$ , then an inverse of  $a$  modulo  $m$  exists. Moreover, this inverse is unique modulo  $m$ .
- Example: Find an inverse of 3 modulo 7
- Example: Solve  $3x \equiv 4 \pmod{7}$

# Sun-Tsu's puzzle

*There are certain things whose number is unknown.  
When divided by 3, the remainder is 2; when divided by  
5, the remainder is 3; and when divided by 7, the  
remainder is 2. What will be the number of things?*

What are the solutions of the system of congruences:

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

# The Chinese remainder theorem

- Theorem: Let  $m_1, m_2, \dots, m_n$  be pairwise relatively prime positive integers and  $a_1, a_2, \dots, a_n$  arbitrary integers. Then the system  $x \equiv a_i \pmod{m_i}$ ,  $i = 1, 2, \dots, n$  has a unique solution modulo  $m = m_1 m_2 \dots m_n$ .
- Example: solve the system
$$\begin{aligned}x &\equiv 2 \pmod{3} \\x &\equiv 3 \pmod{5} \\x &\equiv 2 \pmod{7}\end{aligned}$$

# Format's little theorem

- An integer  $n$  is prime when it is not divisible by any prime  $p \leq \sqrt{n}$ .
- Are there more efficient ways for primality testing?
- Ancient Chinese mathematicians believed:  
 $n$  is prime iff  $2^{n-1} \equiv 1 \pmod{n}$
- Counterexample:  $2^{340} \equiv 1 \pmod{341}$ , but  $341=11 \cdot 31$  (Ex 27)
- Format's little theorem: If  $p$  is prime and  $\gcd(a, p) = 1$ , then  $a^{p-1} \equiv 1 \pmod{p}$ . Furthermore, for every integer  $a$  we have:  $a^p \equiv a \pmod{p}$ . (Ex 19)

# Applications of number theory: Cryptology (optional)

- Cryptology is the study of secret messages.
- Encryption is the process of making a secret message.
- Decryption is the process of recovering the original message from the secret message.



# Caesar's encryption method

- One of the earliest known uses of cryptology is by Caesar.
- He made messages secret by shifting each letter 3 letters forward in the alphabet.
- $f : \{0, 1, \dots, 25\} \rightarrow \{0, 1, \dots, 25\}$   
 $f(p) = (p + 3) \bmod 26$ , called a shift cipher
- $f^{-1}(p) = (p - 3) \bmod 26$
- Generalization of Caesar's method:  $f(p) = (ap + b) \bmod 26$  such that  $f$  is a bijection, called an affine transformation
- Example: what letter replace the letter K (10) when  $f(p) = (7p + 3) \bmod 26$

# Public key cryptography

- Caesar's encryption method is an example of private key cryptosystems
- Knowing the encryption key lets one quickly find the decryption key
- People wanting to communicate must have a separate key
- Public key cryptosystem: everybody can have a publicly known encryption key, and the encryption key does not let one find the decryption key without an extraordinary amount of time