AES 仿真实现

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实验目的

- 1. 通过实验加深对 AES 加密算法的基本原理的理解;
- 2. 了解 AES 算法的详细步骤

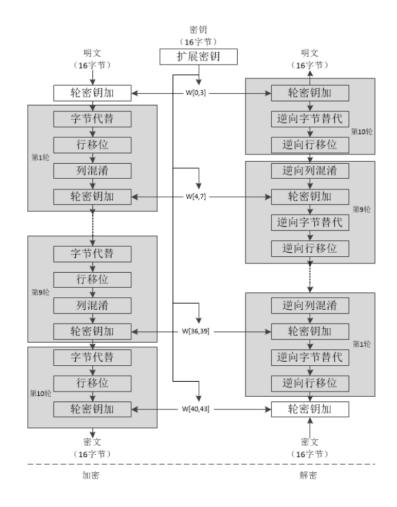
实验原理

AES 的基本结构

AES 为分组密码,分组密码将明文分成多个等长的小组,每次加密一组数据,直到加密 完整个明文。在 AES 标准规范中,分组长度只能为 128 位,那么每个分组为 16 个字节。密 钥的长度可以使用 128 位、192 位或 256 位,密钥的长度不同,加密轮数也不同,如下表:

AES	密钥长度(32 位比特字)	分组长度(32位比特字)	加密轮数
AES-128	4	4	10
AES-192	6	4	12
AES-256	8	4	14

AES 加密算法涉及:字节代替、行位移、列混淆和轮密钥加。下图为 AES 加解密的流程,从图中可以看出:(1)解密算法的每一步分别对应加密算法的逆操作;(2)加解密所有操作的顺序正好是相反的。加解密中每轮的密钥分别由种子密钥经过密钥扩展算法得到。算法中 16 字节的明文、密文和轮子密钥都以一个 4x4 的矩阵表示。



算法原理

1) 字节替代

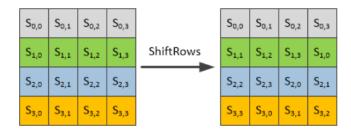
字节替代是通过 S 盒完成一个字节到另外一个字节的映射。这里直接给出构造好的结果, 下图为 S 盒。S 盒用于提供密码算法的混淆性。

									9	y							
		0	1	2	3	4	5	6	7	8	9	A	В	С	D	Е	F
x	0	63	7C	77	7B	F2	6B	6F	C5	30	01	67	2B	FE	D7	AB	76
	1	CA	82	C9	7D	FA	59	47	F0	AD	D4	A2	AF	9C	A4	72	C0
	2	B7	FD	93	26	36	3F	F7	CC	34	A5	E5	F1	71	D8	31	15
	3	04	C7	23	C3	18	96	05	9A	07	12	80	E2	EB	27	B2	75
	4	09	83	2C	1A	1B	6E	5A	A0	52	3B	D6	В3	29	E3	2F	84
	5	53	D1	00	ED	20	FC	B1	5B	6A	CB	BE	39	4A	4C	58	CF
	6	D0	EF	AA	FB	43	4D	33	85	45	F9	02	7F	50	3C	9F	A8
	7	51	A3	40	8F	92	9D	38	F5	BC	B6	DA	21	10	FF	F3	D2
	8	CD	0C	13	EC	5F	97	44	17	C4	A7	7E	3D	64	5D	19	73
	9	60	81	4F	DC	22	2A	90	88	46	EE	B8	14	DE	5E	0B	DB
	Α	E0	32	3A	0A	49	06	24	5C	C2	D3	AC	62	91	95	E4	79
	В	E7	C8	37	6D	8D	D5	4E	A9	6C	56	F4	EA	65	7A	AE	08
	C	BA	78	25	2E	1C	A6	B4	C6	E8	DD	74	1F	4B	BD	8B	8A
	D	70	3E	B5	66	48	03	F6	0E	61	35	57	В9	86	C1	1D	9E
	E	E1	F8	98	11	69	D9	8E	94	9B	1E	87	E9	CE	55	28	DF
	F	8C	A1	89	0D	BF	E6	42	68	41	99	2D	0F	В0	54	BB	16

S 为 16x16 的矩阵,完成一个 8 比特输入到 8 比特输出的映射,输入的高 4-bit 对应的值作为行标,低 4-bit 对应的值作为列标。假设输入字节的值为 $A=a_7a_6a_5a_4a_3a_2a_1a_0$,则输出值为 $S[a_7a_6a_5a_4][a_3a_2a_1a_0]$ 。

2) 行移位

行移位是一个 4x4 的矩阵内部字节之间的置换,用于提供算法的扩散性。正向行移位用于加密,原理:第一行保持不变,第二行循环左移 8 比特,第三行循环左移 16 比特,第四行循环左移 24 比特。假设矩阵名字为 state,用公式表示: state'[i][j] = state[i][(j+1)%4],其中 i、j属于[0,3]。



3) 列混淆

列混淆利用 GF(28)域上算术特性的一个代替,同样用于提供算法的扩散性。 列混合变换是通过矩阵相乘来实现的,经行移位后的状态矩阵与固定的矩阵相乘,得到混淆后的状态矩阵,如下图的公式所示:

状态矩阵中的第j列($0 \le j \le 3$)的列混合可以表示为下图所示:

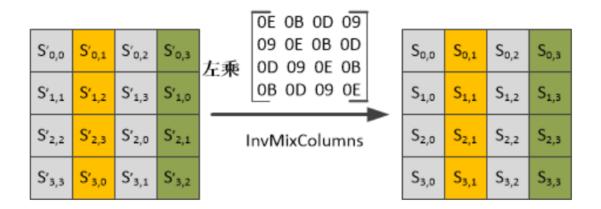
$$s'_{0,j} = (2 * s_{0,j}) \oplus (3 * s_{1,j}) \oplus s_{2,j} \oplus s_{3,j}$$

$$s'_{1,j} = s_{0,j} \oplus (2 * s_{1,j}) \oplus (3 * s_{2,j}) \oplus s_{3,j}$$

$$s'_{2,j} = s_{0,j} \oplus s_{1,j} \oplus (2 * s_{2,j}) \oplus (3 * s_{3,j})$$

$$s'_{3,j} = (3 * s_{0,j}) \oplus s_{1,j} \oplus s_{2,j} \oplus (2 * s_{3,j})$$

逆向列混淆的原理图如下:

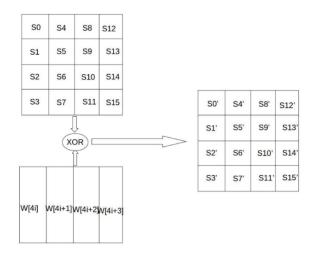


由于:

说明两个矩阵互逆,经过一次逆向列混淆后即可恢复原文,从而实现解密得到明文。

4) 轮密钥加

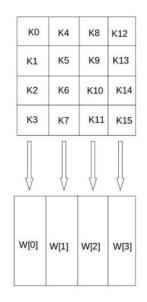
轮密钥加是将 128 位轮密钥 K_i 同状态矩阵中的数据进行逐位异或操作,如下图所示。 其中,密钥 K_i 中每个字 W[4i],W[4i+1],W[4i+2],W[4i+3]为 32 位比特字,包含 4 个字节, 他们的生成算法将在下面介绍。轮密钥加过程可以看成是字逐位异或的结果,也可以看成字 节级别或者位级别的操作。也就是说,可以看成 $S_0S_1S_2S_3$ 组成的 32 位字与 W[4i]的异或运 算。

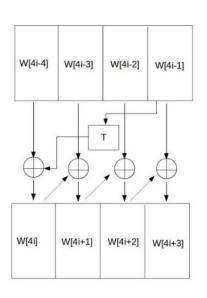


轮密钥加的逆运算同正向的轮密钥加运算完全一致,这是因为异或的逆操作是其自身。 轮密钥加非常简单,但却能够影响 S 数组中的每一位。

5) 密钥扩展

AES 首先将初始密钥输入到一个 4*4 的状态矩阵中, 如下图所示:





这个 4x4 矩阵的每一列的 4 个字节组成一个字,矩阵 4 列的 4 个字依次命名为 W[0]、W[1]、W[2]和 W[3],它们构成一个以字为单位的数组 W。

实验内容

密钥设置为"16337341zhuzhiru", 密钥偏移量为"123":

```
    const char original_key[17] = "16337341zhuzhiru";
    const char original_iv[17] = "123";
```

AES 加密:

```
    string EncryptionAES(const string& strSrc) { //AES 加密
    size_t length = strSrc.length();
    int block_num = length / BLOCK_SIZE + 1;
    //明文
    char* szDataIn = new char[block_num * BLOCK_SIZE + 1];
    memset(szDataIn, 0x00, block_num * BLOCK_SIZE + 1);
```

```
7.
8. #pragma warning(disable:4996)
9.
       strcpy(szDataIn, strSrc.c_str());
10.
       //进行 PKCS7Padding 填充。
11.
12.
       int k = length % BLOCK_SIZE;
        int j = length / BLOCK SIZE;
13.
       int padding = BLOCK_SIZE - k;
14.
15.
        for (int i = 0; i < padding; i++)</pre>
16.
            szDataIn[j * BLOCK_SIZE + k + i] = padding;
17.
        szDataIn[block_num * BLOCK_SIZE] = '\0';
18.
       //加密后的密文
19.
       char *szDataOut = new char[block_num * BLOCK_SIZE + 1];
20.
       memset(szDataOut, 0, block_num * BLOCK_SIZE + 1);
21.
22.
23.
       //进行进行 AES 的 CBC 模式加密
24.
       AES aes;
25.
       aes.MakeKey(original_key, original_iv, 16, 16);
       aes.Encrypt(szDataIn, szDataOut, block_num * BLOCK_SIZE, AES::CBC);
26.
27.
       string str = base64_encode((unsigned char*)szDataOut, block_num * BLOCK_
   SIZE);
28.
       delete[] szDataIn;
       delete[] szDataOut;
29.
30.
        return str;
31. }
```

AES 解密:

```
string DecryptionAES(const string& strSrc) { //AES 解密
2.
       string strData = base64_decode(strSrc);
       size_t length = strData.length();
3.
4.
       //密文
5.
       char *szDataIn = new char[length + 1];
       memcpy(szDataIn, strData.c_str(), length + 1);
7.
       //明文
       char *szDataOut = new char[length + 1];
8.
9.
       memcpy(szDataOut, strData.c_str(), length + 1);
10.
       //进行 AES 的 CBC 模式解密
11.
12.
       AES aes;
13.
       aes.MakeKey(original_key, original_iv, 16, 16);
14.
       aes.Decrypt(szDataIn, szDataOut, length, AES::CBC);
15.
```

```
//去 PKCS7Padding 填充
16.
17.
        if (0x00 < szDataOut[length - 1] <= 0x16) {</pre>
            int tmp = szDataOut[length - 1];
18.
            for (int i = length - 1; i >= (length - tmp); i--) {
19.
                if (szDataOut[i] != tmp) {
20.
21.
                    memset(szDataOut, 0, length);
                    cout << "去填充失败! 解密出错!!" << endl;
22.
23.
                    break;
24.
                }
25.
                else
26.
                    szDataOut[i] = 0;
27.
            }
28.
        string strDest(szDataOut);
29.
       delete[] szDataIn;
30.
31.
       delete[] szDataOut;
32.
       return strDest;
33. }
```

XOR 运算:

```
1. void AES::Xor(char* buff, char const* chain) {
2.    if (false == m_bKeyInit)
3.        return;
4.    for (int i = 0; i < m_blockSize; i++)
5.        *(buff++) ^= *(chain++);
6. }</pre>
```

扩展密钥:

```
    void AES::MakeKey(char const* key, char const* chain, int keylength, int blo

   ckSize) {
2.
       if (NULL == key)
3.
            return;
       if (!(16 == keylength || 24 == keylength || 32 == keylength))
4.
5.
            return;
       if (!(16 == blockSize || 24 == blockSize || 32 == blockSize))
6.
7.
            return;
       m_keylength = keylength;
8.
       m_blockSize = blockSize;
9.
       //初始化
10.
       memcpy(m_chain0, chain, m_blockSize);
11.
12.
       memcpy(m_chain, chain, m_blockSize);
```

```
13.
        //计算轮数
14.
        switch (m keylength) {
15.
            case 16:
                m_iROUNDS = (m_blockSize == 16)?10:(m_blockSize == 24?12:14);
16.
17.
                break;
18.
            case 24:
19.
                m_iROUNDS = (m_blockSize != 32) ? 12 : 14;
20.
21.
                break;
22.
            default: // 32 bytes = 256 bits
23.
                m_iROUNDS = 14;
24.
25.
        }
        int BC = m_blockSize / 4;
26.
        int i, j;
27.
28.
        for (i = 0; i <= m_iROUNDS; i++) {</pre>
29.
            for (j = 0; j < BC; j++)
30.
                m Ke[i][j] = 0;
31.
        }
32.
        for (i = 0; i <= m_iROUNDS; i++) {</pre>
33.
            for (j = 0; j < BC; j++)
34.
                m_Kd[i][j] = 0;
35.
        int ROUND_KEY_COUNT = (m_iROUNDS + 1) * BC;
36.
        int KC = m_keylength / 4;
37.
38.
        //将 bytes 转成 int
39.
        int* pi = tk;
        char const* pc = key;
40.
41.
        for (i = 0; i < KC; i++) {</pre>
42.
            *pi = (unsigned char) *(pc++) << 24;
43.
            *pi |= (unsigned char) *(pc++) << 16;
            *pi |= (unsigned char) *(pc++) << 8;
44.
45.
            *(pi++) |= (unsigned char) *(pc++);
       }
46.
        //轮密钥数组赋值
47.
        int t = 0;
48.
        for (j = 0; (j < KC) && (t < ROUND_KEY_COUNT); j++, t++) {</pre>
49.
50.
            m_{Ke[t / BC][t \% BC] = tk[j];
            m_Kd[m_iROUNDS - (t / BC)][t % BC] = tk[j];
51.
52.
        }
53.
        int tt, rconpointer = 0;
54.
        while (t < ROUND_KEY_COUNT) {</pre>
55.
            //使用 phi 计算
            tt = tk[KC - 1];
56.
```

```
tk[0] = (sm_S[(tt >> 16) \& 0xFF] \& 0xFF) << 24
57.
58.
                ^ (sm_S[(tt >> 8) & 0xFF] & 0xFF) << 16
                ^ (sm_S[tt & 0xFF] & 0xFF) << 8
59.
                ^ (sm_S[(tt >> 24) & 0xFF] & 0xFF)
60.
                ^ (sm_rcon[rconpointer++] & 0xFF) << 24;</pre>
61.
            if (KC != 8)
62.
                for (i = 1, j = 0; i < KC;)
63.
64.
                    tk[i++] ^= tk[j++];
65.
            else {
                for (i = 1, j = 0; i < KC / 2;)
66.
67.
                    tk[i++] ^= tk[j++];
                tt = tk[KC / 2 - 1];
68.
69.
                tk[KC / 2] ^= (sm_S[tt & 0xFF] & 0xFF)
70.
                    ^ (sm_S[(tt >> 8) & 0xFF] & 0xFF) << 8
71.
                    ^ (sm_S[(tt >> 16) & 0xFF] & 0xFF) << 16
72.
                    ^ (sm_S[(tt >> 24) & 0xFF] & 0xFF) << 24;
73.
                for (j = KC / 2, i = j + 1; i < KC;)
74.
                    tk[i++] ^= tk[j++];
75.
            }
76.
            //轮密钥数组赋值
            for (j = 0; (j < KC) && (t < ROUND_KEY_COUNT); j++, t++) {</pre>
77.
78.
                m_Ke[t / BC][t % BC] = tk[j];
79.
                m Kd[m iROUNDS - (t / BC)][t % BC] = tk[j];
80.
       }
81.
        //逆行混淆
82.
        for (int r = 1; r < m_iROUNDS; r++)</pre>
83.
84.
            for (j = 0; j < BC; j++) {
85.
                tt = m_Kd[r][j];
86.
                m_Kd[r][j] = sm_U1[(tt >> 24) & 0xFF] ^ sm_U2[(tt >> 16) & 0xFF]
                    ^ sm_U3[(tt >> 8) & 0xFF] ^ sm_U4[tt & 0xFF];
87.
88.
89.
       m_bKeyInit = true;
90.}
```

定义加密块:

```
    void AES::DefEncryptBlock(char const* in, char* result) {
    if (false == m_bKeyInit)
    return;
    int* Ker = m_Ke[0];
    int t0 = ((unsigned char) *(in++) << 24);</li>
    t0 |= ((unsigned char) *(in++) << 16);</li>
```

```
7.
       t0 |= ((unsigned char) *(in++) << 8);
8.
        (t0 |= (unsigned char) *(in++)) ^= Ker[0];
        int t1 = ((unsigned char) *(in++) << 24);</pre>
9.
       t1 |= ((unsigned char) *(in++) << 16);
10.
11.
       t1 |= ((unsigned char) *(in++) << 8);
        (t1 |= (unsigned char) *(in++)) ^= Ker[1];
12.
13.
        int t2 = ((unsigned char) *(in++) << 24);</pre>
14.
       t2 |= ((unsigned char) *(in++) << 16);
       t2 |= ((unsigned char) *(in++) << 8);
15.
16.
        (t2 |= (unsigned char) *(in++)) ^= Ker[2];
17.
        int t3 = ((unsigned char) *(in++) << 24);</pre>
       t3 |= ((unsigned char) *(in++) << 16);
18.
19.
       t3 |= ((unsigned char) *(in++) << 8);
20.
        (t3 |= (unsigned char) *(in++)) ^= Ker[3];
21.
       int a0, a1, a2, a3;
22.
        //轮加密
        for (int r = 1; r < m_iROUNDS; r++) {</pre>
23.
24.
            Ker = m Ke[r];
25.
            a0 = (sm_T1[(t0 >> 24) \& 0xFF] ^ sm_T2[(t1 >> 16) \& 0xFF]
                ^ sm_T3[(t2 >> 8) & 0xFF] ^ sm_T4[t3 & 0xFF]) ^ Ker[0];
26.
            a1 = (sm_T1[(t1 >> 24) \& 0xFF] \land sm_T2[(t2 >> 16) \& 0xFF]
27.
28.
                ^ sm_T3[(t3 >> 8) & 0xFF] ^ sm_T4[t0 & 0xFF]) ^ Ker[1];
29.
            a2 = (sm T1[(t2 >> 24) \& 0xFF] ^ sm T2[(t3 >> 16) \& 0xFF]
                ^ sm_T3[(t0 >> 8) & 0xFF] ^ sm_T4[t1 & 0xFF]) ^ Ker[2];
30.
31.
            a3 = (sm_T1[(t3 >> 24) \& 0xFF] ^ sm_T2[(t0 >> 16) \& 0xFF]
                ^ sm_T3[(t1 >> 8) & 0xFF] ^ sm_T4[t2 & 0xFF]) ^ Ker[3];
32.
33.
            t0 = a0;
34.
            t1 = a1;
35.
            t2 = a2;
36.
            t3 = a3;
       }
37.
        //最后一轮加密不同
38.
39.
       Ker = m_Ke[m_iROUNDS];
40.
        int tt = Ker[0];
        result[0] = sm_S[(t0 >> 24) & 0xFF] ^ (tt >> 24);
41.
        result[1] = sm_S[(t1 >> 16) & 0xFF] ^ (tt >> 16);
42.
        result[2] = sm_S[(t2 >> 8) \& 0xFF] ^ (tt >> 8);
43.
44.
       result[3] = sm_S[t3 \& 0xFF] ^ tt;
45.
       tt = Ker[1];
        result[4] = sm_S[(t1 >> 24) & 0xFF] ^ (tt >> 24);
46.
47.
        result[5] = sm_S[(t2 >> 16) & 0xFF] ^ (tt >> 16);
       result[6] = sm_S[(t3 >> 8) \& 0xFF] ^ (tt >> 8);
48.
49.
        result[7] = sm_S[t0 \& 0xFF] ^ tt;
50.
        tt = Ker[2];
```

```
51.
        result[8] = sm_S[(t2 >> 24) \& 0xFF] ^ (tt >> 24);
52.
        result[9] = sm_S[(t3 >> 16) & 0xFF] ^ (tt >> 16);
        result[10] = sm_S[(t0 >> 8) \& 0xFF] ^ (tt >> 8);
53.
54.
        result[11] = sm_S[t1 & 0xFF] ^ tt;
55.
        tt = Ker[3];
        result[12] = sm_S[(t3 >> 24) & 0xFF] ^ (tt >> 24);
56.
57.
        result[13] = sm S[(t0 \Rightarrow 16) & 0xFF] ^ (tt \Rightarrow 16);
        result[14] = sm_S[(t1 >> 8) & 0xFF] ^ (tt >> 8);
58.
59.
        result[15] = sm_S[t2 \& 0xFF] ^ tt;
60.}
```

定义解密块:

```
1. void AES::DefDecryptBlock(char const* in, char* result) {
2.
        if (false == m_bKeyInit)
3.
            return;
4.
        int* Kdr = m_Kd[0];
5.
        int t0 = ((unsigned char) *(in++) << 24);</pre>
        t0 = t0 | ((unsigned char) *(in++) << 16);
6.
        t0 |= ((unsigned char) *(in++) << 8);</pre>
7.
8.
        (t0 |= (unsigned char) *(in++)) ^= Kdr[0];
9.
        int t1 = ((unsigned char) *(in++) << 24);</pre>
10.
        t1 |= ((unsigned char) *(in++) << 16);
11.
        t1 |= ((unsigned char) *(in++) << 8);
        (t1 |= (unsigned char) *(in++)) ^= Kdr[1];
12.
13.
        int t2 = ((unsigned char) *(in++) << 24);</pre>
        t2 |= ((unsigned char) *(in++) << 16);
14.
        t2 |= ((unsigned char) *(in++) << 8);
15.
        (t2 |= (unsigned char) *(in++)) ^= Kdr[2];
16.
17.
        int t3 = ((unsigned char) *(in++) << 24);</pre>
        t3 |= ((unsigned char) *(in++) << 16);
18.
19.
        t3 |= ((unsigned char) *(in++) << 8);
20.
        (t3 |= (unsigned char) *(in++)) ^= Kdr[3];
        int a0, a1, a2, a3;
21.
22.
        //轮解密
        for (int r = 1; r < m_iROUNDS; r++) {</pre>
23.
24.
            Kdr = m Kd[r];
            a0 = (sm_T5[(t0 >> 24) \& 0xFF] ^ sm_T6[(t3 >> 16) \& 0xFF]
25.
26.
                ^ sm_T7[(t2 >> 8) & 0xFF] ^ sm_T8[t1 & 0xFF]) ^ Kdr[0];
            a1 = (sm_T5[(t1 >> 24) \& 0xFF] ^ sm_T6[(t0 >> 16) \& 0xFF]
27.
                ^ sm_T7[(t3 >> 8) & 0xFF] ^ sm_T8[t2 & 0xFF]) ^ Kdr[1];
28.
29.
            a2 = (sm_T5[(t2 >> 24) \& 0xFF] ^ sm_T6[(t1 >> 16) \& 0xFF]
30.
                ^ sm_T7[(t0 >> 8) & 0xFF] ^ sm_T8[t3 & 0xFF]) ^ Kdr[2];
            a3 = (sm_T5[(t3 >> 24) \& 0xFF] ^ sm_T6[(t2 >> 16) \& 0xFF]
31.
```

```
32.
                ^ sm_T7[(t1 >> 8) & 0xFF] ^ sm_T8[t0 & 0xFF]) ^ Kdr[3];
33.
            t0 = a0;
34.
           t1 = a1;
35.
            t2 = a2;
36.
            t3 = a3;
37.
       }
       //最后一轮解密不同
38.
39.
       Kdr = m_Kd[m_iROUNDS];
40.
       int tt = Kdr[0];
41.
       result[0] = sm_Si[(t0 >> 24) \& 0xFF] ^ (tt >> 24);
42.
       result[1] = sm_Si[(t3 >> 16) & 0xFF] ^ (tt >> 16);
       result[2] = sm_Si[(t2 >> 8) & 0xFF] ^ (tt >> 8);
43.
44.
       result[3] = sm_Si[t1 \& 0xFF] ^ tt;
45.
       tt = Kdr[1];
       result[4] = sm_Si[(t1 >> 24) & 0xFF] ^ (tt >> 24);
46.
47.
       result[5] = sm_Si[(t0 >> 16) & 0xFF] ^ (tt >> 16);
48.
       result[6] = sm_Si[(t3 >> 8) & 0xFF] ^ (tt >> 8);
       result[7] = sm_Si[t2 & 0xFF] ^ tt;
49.
       tt = Kdr[2];
50.
       result[8] = sm_Si[(t2 >> 24) & 0xFF] ^ (tt >> 24);
51.
       result[9] = sm_Si[(t1 >> 16) & 0xFF] ^ (tt >> 16);
52.
53.
       result[10] = sm_Si[(t0 >> 8) \& 0xFF] ^ (tt >> 8);
54.
       result[11] = sm Si[t3 \& 0xFF] ^ tt;
       tt = Kdr[3];
55.
56.
       result[12] = sm_Si[(t3 >> 24) & 0xFF] ^ (tt >> 24);
       result[13] = sm_Si[(t2 >> 16) & 0xFF] ^ (tt >> 16);
57.
       result[14] = sm_Si[(t1 >> 8) & 0xFF] ^ (tt >> 8);
58.
       result[15] = sm_Si[t0 & 0xFF] ^ tt;
59.
60.}
```

加密:

```
1. void AES::Encrypt(char const* in, char* result, size_t n, int iMode) {
2.
       if (false == m_bKeyInit)
3.
        // n 应该大于 0 且 m_blockSize 不等于 0
4.
        if (0 == n || n % m_blockSize != 0)
5.
6.
            return;
7.
       int i;
8.
        char const* pin;
9.
       char* presult;
        if (CBC == iMode) { //CBC 模式
10.
11.
            for (i = 0, pin = in, presult = result; i < n / m_blockSize; i++) {</pre>
12.
                Xor(m_chain, pin);
```

```
13.
                EncryptBlock(m_chain, presult);
14.
                memcpy(m chain, presult, m blockSize);
15.
                pin += m_blockSize;
                presult += m_blockSize;
16.
17.
            }
18.
        else if (CFB == iMode) { //CFB 模式
19.
            for (i = 0, pin = in, presult = result; i < n / m_blockSize; i++) {</pre>
20.
21.
                EncryptBlock(m_chain, presult);
22.
                Xor(presult, pin);
23.
                memcpy(m chain, presult, m blockSize);
24.
                pin += m_blockSize;
25.
                presult += m_blockSize;
26.
27.
       }
28.
       else { //ECB 模式
            for (i = 0, pin = in, presult = result; i < n / m_blockSize; i++) {</pre>
29.
30.
                EncryptBlock(pin, presult);
31.
                pin += m_blockSize;
32.
                presult += m_blockSize;
33.
            }
34.
35.}
```

解密:

```
void AES::Decrypt(char const* in, char* result, size_t n, int iMode) {
2.
       if (false == m_bKeyInit)
3.
            return;
       // n 应该大于 0 且 m_blockSize 不等于 0
5.
       if (0 == n || n % m_blockSize != 0)
6.
            return;
7.
       int i;
8.
       char const* pin;
        char* presult;
9.
10.
        if (CBC == iMode) { //CBC 模式
11.
            for (i = 0, pin = in, presult = result; i < n / m_blockSize; i++) {</pre>
12.
                DecryptBlock(pin, presult);
13.
                Xor(presult, m_chain);
14.
                memcpy(m_chain, pin, m_blockSize);
15.
                pin += m_blockSize;
16.
                presult += m blockSize;
17.
            }
18.
```

```
19.
        else if (CFB == iMode) {//CFB 模式
20.
            for (i = 0, pin = in, presult = result; i < n / m blockSize; i++) {</pre>
                EncryptBlock(m_chain, presult);
21.
22.
                Xor(presult, pin);
23.
                memcpy(m_chain, pin, m_blockSize);
24.
                pin += m_blockSize;
                presult += m blockSize;
25.
26.
27.
        }
        else { //ECB 模式
28.
            for (i = 0, pin = in, presult = result; i < n / m blockSize; i++) {</pre>
29.
30.
                DecryptBlock(pin, presult);
31.
                pin += m_blockSize;
                presult += m_blockSize;
32.
33.
            }
34.
35.}
```

实验结果

运行程序 AES.exe, 对 "Cryptography and Network Security Principles and Practice, Sixth Edition"加密后得到的密文为 "nHDxvEH0iTh0tXVYkWec99Bk0pNMpeS4Y2sDvieHw6uK 1RmGK1PHfa0PPZf/04b6uWqsHETH1tJRoE/B7Rd7swB1VDMeEGHduKV44lkhYyw=",如图所示:

```
指令: 1.加密明文 2.解密密文
请输入: 1
请输入明文:
Cryptography and Network Security Principles and Practice, Sixth Edition
加密后:nHDxvEHOiThOtXVYkWec99BkOpNMpeS4Y2sDvieHw6uK1RmGK1PHfaOPPZf/O4b6uWqsHETH1tJRoE/B7Rd7swB1VDMeEGHduKV441khYyw=
请问是否继续? 1.继续 2.退出 _
```

对密文 "nHDxvEH0iTh0tXVYkWec99Bk0pNMpeS4Y2sDvieHw6uK1RmGK1PHfa0PPZf /04b6uWqsHETH1tJRoE/B7Rd7swB1VDMeEGHduKV44lkhYyw= "解密可得明文 "Cryptography and Network Security Principles and Practice, Sixth Edition",如图所示:

```
指令: 1. 加密明文 2. 解密密文
请输入: 1
请输入明文;
Cryptography and Network Security Principles and Practice, Sixth Edition
加密后:nHDxvEHOiThOtXVYkWec99BkOpNMpeS4Y2sDvieHw6uK1RmGK1PHfaOPPZf/O4b6uWqsHETH1tJRoE/B7Rd7swB1VDMeEGHduKV44lkhYyw=
请问是否继续? 1. 继续 2. 退出 1
指令: 1. 加密明文 2. 解密密文
请输入: 2
请输入: 2
请输入: 2
时加xvEHOiThOtXVYkWec99BkOpNMpeS4Y2sDvieHw6uK1RmGK1PHfaOPPZf/O4b6uWqsHETH1tJRoE/B7Rd7swB1VDMeEGHduKV44lkhYyw=
解密后:Cryptography and Network Security Principles and Practice, Sixth Edition
请问是否继续? 1. 继续 2. 退出
```

由此可以看出 AES 加密和解密过程正确无误。