

# The x86 PC

assembly language, design, and interfacing

fifth  
edition

Prentice Hall

Dec	Hex	Bin
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ORG ; SIX

## Signed Numbers Strings and Tables

# The x86 PC

assembly language,  
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fifth edition

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# OBJECTIVES

this chapter enables the student to:

- Convert a number to its 2's complement.
- Code signed arithmetic instructions:
  - ADD, SUB, IMUL, and IDIV.
- Demonstrate how arithmetic instructions affect the sign flag.
- Explain the difference between a carry and an overflow.
- Prevent overflow errors by sign-extending data.
- Code signed shift instructions:
  - SAL and SAR.

# OBJECTIVES

(*cont*)

this chapter enables the student to:

- Code logic instruction CMP for signed numbers and explain its effect on the flag register.
- Code conditional jump instructions after CMP of signed data.
- Explain the function of registers SI and DI in string instructions.
- Describe the operation of the direction flag in string instructions.
- Code instructions CLD and STD to control the direction flag.

# OBJECTIVES

(*cont*)

this chapter enables the student to:

- Describe the operation of the REP prefix
- Code string instructions:
  - MOVSB and MOVSW for data transfer.
  - STOS, LODS to store and load the contents of AX.
  - CMPS to compare two strings of data.
  - SCAS to scan a string for data matching in AX.
  - XLAT for table processing .

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

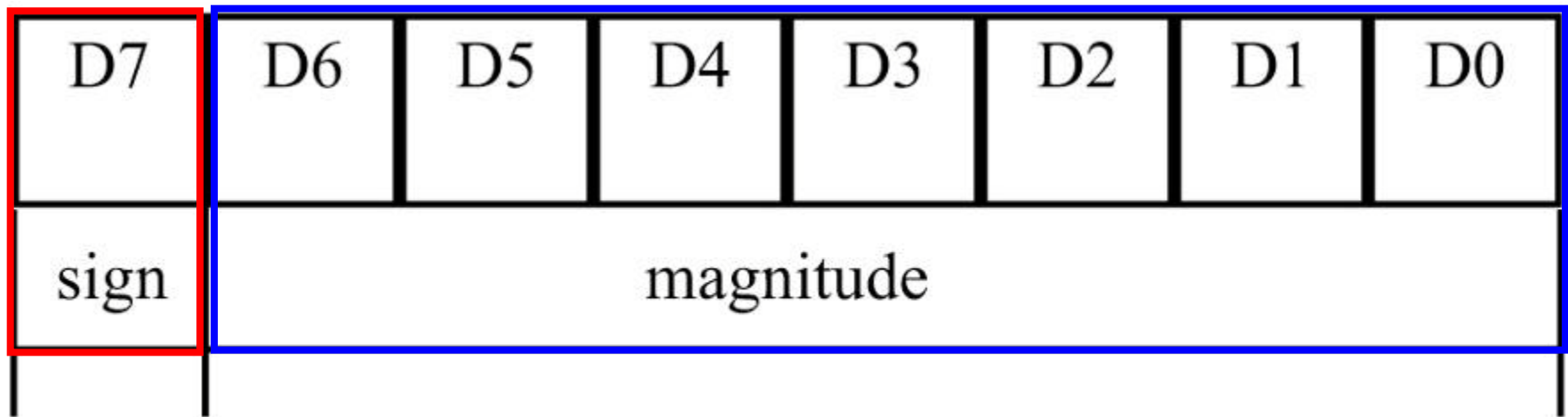
## Concept of signed numbers in computers

- Many applications require signed data, & computers must be able to accommodate such numbers.
  - The most significant bit (MSB) is set aside for the sign (+ or -) & the rest of the bits are used for the magnitude.
  - The sign is represented by 0 for positive (+) numbers and 1 for negative (-) numbers.

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

### signed byte operands

- In signed byte operands, **D7 (MSB)** is the sign and **D0 to D6** are set aside for the *magnitude* of the number.
  - If D7 = 0, the operand is *positive*.
  - If D7 = 1, it is *negative*.



## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

### positive numbers

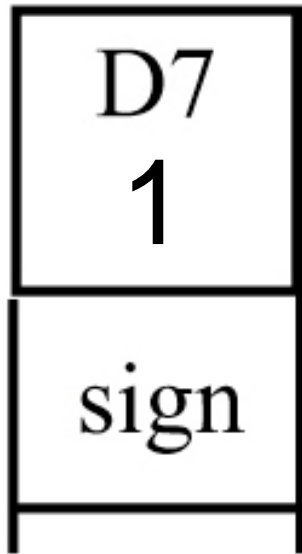
- The range of positive numbers that can be represented by the format above is 0 to +127.
  - If a positive number is larger than +127, a word-sized operand must be used.

0	0000	0000
+1	0000	0001
+5	0000	0101
...	...	...
+127	0111	1111

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## negative numbers

- For negative numbers D7 is 1, but the magnitude is represented in 2's complement.



To convert to negative number representation (2's complement), follow these steps:

1. Write the magnitude of the number in 8-bit binary (no sign).
2. Invert each bit.
3. Add 1 to it.



# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## negative numbers

### Example 6-1

Show how the computer would represent -5.

#### Solution:

1. 0000 0101      5 in 8-bit binary
2. 1111 1010      invert each bit
3. 1111 1011      add 1 (hex = FBH)

The signed number representation in 2's complement for -5.

The range of byte-sized negative numbers is -1 to -128.

### Example 6-2

Show -34H as it is represented internally.

#### Solution:

1. 0011 0100
2. 1100 1011
3. 1100 1100      (which is CCH)

### Example 6-3

Show the representation for -128<sub>10</sub>.

#### Solution:

1. 1000 0000
2. 0111 1111
3. 1000 0000      Notice that this is not negative zero (-0).

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## negative numbers

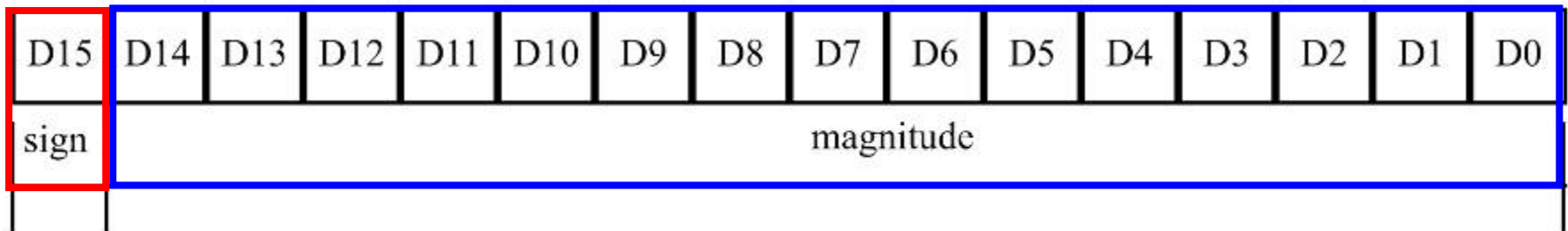
Byte-sized signed number ranges.

<u>Decimal</u>	<u>Binary</u>	<u>Hex</u>
-128	1000 0000	80
-127	1000 0001	81
-126	1000 0010	82
...	.... ...	..
-2	1111 1110	FE
-1	1111 1111	FF
0	0000 0000	00
+1	0000 0001	01
+2	0000 0010	02
...	... ...	...
+127	0111 1111	7F

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## word-sized signed numbers

- A word is 16 bits in length in x86 computers.
  - Setting aside the **MSB (D15)** for the sign leaves a total of 15 bits (**D14 – D0**) for the magnitude.
    - A range of -32768 to +32767.



- Larger numbers must be treated as a multiword operand, processed chunk by chunk.
  - The same way as unsigned numbers.

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## word-sized signed numbers

The range of signed word operands.

<u>Decimal</u>	<u>Binary</u>	<u>Hex</u>
-32,768	1000 0000 0000 0000	8000
-32,767	1000 0000 0000 0001	8001
-32,766	1000 0000 0000 0010	8002
...	...	...
...	...	...
-2	1111 1111 1111 1110	FFFE
-1	1111 1111 1111 1111	FFFF
0	0000 0000 0000 0000	0000
+1	0000 0000 0000 0001	0001
+2	0000 0000 0000 0010	0002
...	...	...
...	...	...
+32,766	0111 1111 1111 1110	7FFE
+32,767	0111 1111 1111 1111	7FFF

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

### word-sized signed numbers

- To convert a negative number to its word operand representation, the same three steps discussed in negative byte operands are used:

1. Write the magnitude of the number in 8-bit binary (no sign).
2. Invert each bit.
3. Add 1 to it.

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS signed number operations overflow problem

- When using signed numbers, a serious issue, the overflow problem, must be dealt with.
  - The CPU understands only 0s & 1s, ignoring the human convention of positive/negative numbers.
    - The CPU indicates the problem with the OF (overflow) flag.
- If the result of an operation on signed numbers is too large for the register, an overflow occurs.
  - Review Example 6-4.

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS signed number operations overflow problem

## Example 6-4

```
DATA1      DB      +96
DATA2      DB      +70
...
MOV        AL, DATA1      ;AL=0110 0000 (AL=60H)
MOV        BL, DATA2      ;BL=0100 0110 (BL=46H)
ADD        AL, BL          ;AL=1010 0110 (AL=A6H= 90 invalid!)

+ 96 0110 0000
+ 70 0100 0110
+166 1010 0110 According to the CPU, this is 90, which is wrong. (OF = 1, SF = 1, CF = 0)
```

**+96** is added to **+70** and the result according to the CPU is **-90**.

*Why?* The result was more than **AL** could handle, because like all other 8-bit registers, **AL** could only contain up to +127.

The CPU designers created the overflow flag to inform the programmer that the result of the signed number operation is erroneous.

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS the overflow flag in 8-bit operations

- In 8-bit signed number operations, OF is set to 1 if:
  - There is a carry from D6 to D7 but no carry out of D7.
    - (CF = 0)
  - There is a carry from D7 out, but no carry from D6 to D7.
    - (CF = 1)
- The overflow flag is set to 1 if there is a carry from D6 to D7 or from D7 out, but *not* both.
  - If there is a carry both from D6 to D7, and from D7 out, then OF = 0.



## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS the overflow flag in 8-bit operations

- In Example 6-4, since there is only a carry from D6 to D7 and no carry from D7 out,  $OF = 1$ .
  - Examples 6-5, 6-6, and 6-7 give further illustrations of the overflow flag in signed arithmetic.

**Example 6-5** Observe the results of the following:

```
MOV    DL,-128      ;DL=1000 0000 (DL=80H)
MOV    CH,-2        ;CH=1111 1110 (CH=FEH)
ADD    DL,CH        ;DL=0111 1110 (DL=7EH=+126 invalid!)
```

```
  -128  1000 0000
+   -2   1111 1110
-----
 -130   0111 1110  OF=1, SF=0 (positive), CF=1
```

According to the CPU, the result is +126, which is wrong. The error is indicated by  $OF = 1$ .

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS the overflow flag in 8-bit operations

**Example 6-6** Observe the results of the following:

```
MOV    AL, -2          ;AL=1111 1110 (AL=FEH)
MOV    CL, -5          ;CL=1111 1011 (CL=FBH)
ADD    CL, AL          ;CL=1111 1001 (CL=F9H=7 which is correct)
```

```
  -2  1111 1110
+ -5  1111 1011
-----
  -7  1111 1001  OF = 0, CF = 0 , and SF = 1 (negative); the result is correct since OF = 0.
```

**Example 6-7** Observe the results of the following:

```
MOV    DH, +7          ;DH=0000 0111 (DH=07H)
MOV    BH, +18         ;BH=0001 0010 (BH=12H)
ADD    BH, DH          ;BH=0001 1001 (BH=19H=+25, correct)
```

```
  +7      0000 0111
+ +18     0001 0010
-----
 +25     0001 1001  OF = 0, CF = 0, and SF = 0 (positive).
```

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

### overflow flag in 16-bit operations

- In a 16-bit operation, OF is set to 1 in two cases:
  - A carry from D14 to D15, but no carry out of D15.
    - (CF = 0).
  - A carry from D15 out but no carry from D14 to D15.
    - (CF = 1)
- Again, the overflow flag is low (*not* set) if there is a carry from both D14 to D15 and from D15 out.
  - See examples 6-8 & 6-9.

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## overflow flag in 16-bit operations

OF is set to 1 only when there is carry from D14 to D15 or from D15 out, but *not* from both.

### Example 6-8

```
MOV    AX, 6E2FH    ; 28,207
MOV    CX, 13D4H    ; + 5,076
ADD    AX, CX        ; = 33,283 – expected answer
```

6E2F	0110	1110	0010	1111	
+13D4	0001	0011	1101	0100	
8203	1000	0010	0000	0011	= - 32,253 incorrect!
	OF = 1,	CF = 0,	SF = 1		

### Example 6-9

```
MOV    DX, 542FH    ; 21,551
MOV    BX, 12E0H    ; +4,832
ADD    DX, BX        ; =26,383
```

543F	0101	0100	0010	1111	
+12E0	0001	0010	1110	0000	
670F	0110	0111	0000	1111	= 26,383 (correct answer); OF = 0, CF = 0, SF = 0

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

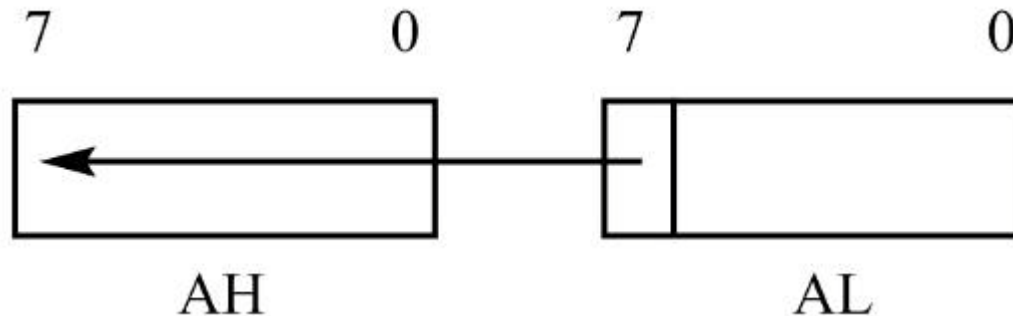
### avoiding erroneous results

- Sign-extending the operand can avoid problems associated with signed number operations.
  - This copies the sign bit (D7) of the lower byte of a register into the upper bits of the register.
    - Or copies the sign bit of a 16-bit register into another register.
- Two directives are used to perform sign extension:
  - CBW (convert signed byte to signed word)
  - CWD (convert signed word to signed double word)

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

### avoiding erroneous results

- CBW will copy D7 (the sign flag) to all bits of AH.
  - The operand is assumed to be AL, and the previous contents of AH are destroyed.



```
MOV    AL, +96          ;AL=0110 0000
CBW                                ;now AH=0000 0000 and AL=0110 0000
```

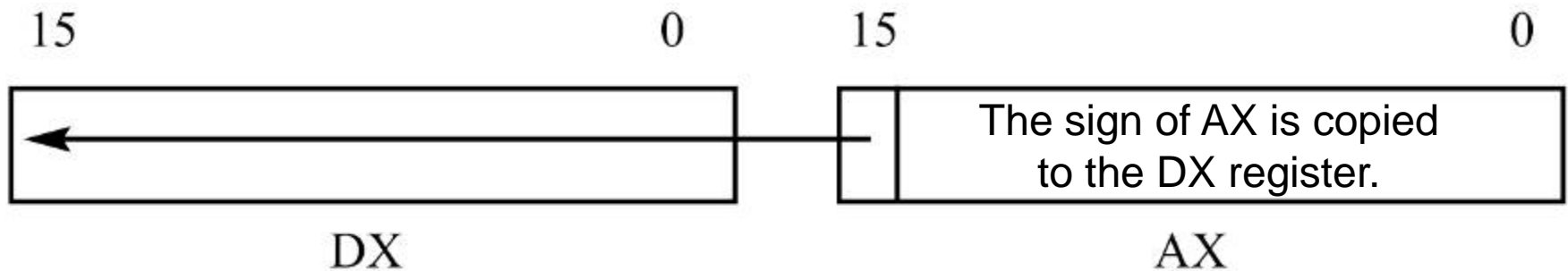
OR:

```
MOV    AL, -2           ;AL=1111 1110
CBW                                ;AH=1111 1111 and AL=1111 1110
```

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

### avoiding erroneous results

- CWD sign-extends AX, copying D15 of AX to all bits of the DX register.
  - This is used for signed word operands.



```
MOV    AX, +260      ;AX=0000 0001 0000 0100 or AX=0104H
CWD                      ;DX=0000H and AX=0104H
```

Another example :

```
MOV    AX, -32766     ;AX=1000 0000 0000 0010B or AX=8002H
CWD                      ;DX=FFFF and AX=8002
```



# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## avoiding erroneous results

Example 6-10 shows 6-4 rewritten to correct for overflow.

If the overflow flag is not raised, the result of the signed number is correct & **JNO** (jump if no overflow) will jump to **OVER**.

### Example 6-10

Rewrite Example 6-4 to provide for handling the overflow problem.

#### Solution:

```
DATA1      DB      +96
DATA2      DB      +70
RESULT     DW      ?

        . . . . .
SUB       AH, AH           ; AH=0
MOV       AL, DATA1       ; GET OPERAND 1
MOV       BL, DATA2       ; GET OPERAND 2
ADD       AL, BL           ; ADD THEM
JNO       OVER             ; IF OF=0 THEN GO TO OVER
MOV       AL, DATA2       ; OTHERWISE GET OPERAND 2 TO
CBW                           ; SIGN EXTEND IT
MOV       BX, AX           ; SAVE IT IN BX
MOV       AL, DATA1       ; GET BACK OPERAND 1 TO
CBW                           ; SIGN EXTEND IT
ADD       AX, BX           ; ADD THEM AND
OVER:     MOV       RESULT, AX ; SAVE IT
```



# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## avoiding erroneous results

Example 6-10 shows 6-4 rewritten to correct for overflow.

If OF = 1, result is erroneous & each operand must be sign-extended, then added.

This works for addition of any two signed bytes.

### Example 6-10

Rewrite Example 6-4 to provide for handling the overflow problem.

#### Solution:

```
DATA1      DB      +96
DATA2      DB      +70
RESULT     DW      ?

        . . . . .
SUB       AH,AH          ; AH=0
MOV       AL,DATA1       ; GET OPERAND 1
MOV       BL,DATA2       ; GET OPERAND 2
ADD       AL,BL          ; ADD THEM
JNO       OVER           ; IF OF=0 THEN GO TO OVER
MOV       AL,DATA2       ; OTHERWISE GET OPERAND 2 TO
CBW                          ; SIGN EXTEND IT
MOV       BX,AX          ; SAVE IT IN BX
MOV       AL,DATA1       ; GET BACK OPERAND 1 TO
CBW                          ; SIGN EXTEND IT
ADD       AX,BX          ; ADD THEM AND
OVER:     MOV       RESULT,AX ; SAVE IT
```

## 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

### avoiding erroneous results

- Analysis of the values in Example 6-10.
  - Each is sign-extended and added as follows:

<u>S</u>	<u>AH</u>	<u>AL</u>	
0	000 0000	0110 0000	+96 after sign extension
0	000 0000	0100 0110	+70 after sign extension
0	000 0000	1010 0110	+166

- If the possibility of overflow exists, byte-sized signed numbers should be sign-extended into a word,
- Word-sized signed operands should be sign-extended before they are processed

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## IDIV signed number division

- IDIV means "*integer division*", used for signed number division.
  - All 8088/86 arithmetic instructions are for integer numbers, regardless of whether the operands are signed or unsigned.
    - Real numbers operations are done by the 8087 coprocessor.

**Table 6-1: Signed Division Summary**

Division	Numerator	Denominator	Quotient	Rem.
byte/byte	AL = byte CBW	register or memory	AL	AH
word/word	AX = word CWD	register or memory	AX	DX
word/byte	AX = word	register or memory	AL <sup>1</sup>	AH
doubleword/word	DXAX = doubleword	register or memory	AX <sup>2</sup>	DX

Notes: 1. Divide error interrupt if  $-127 > AL > +127$ . 2. Divide error interrupt if  $-32,767 > AL > +32,767$ .

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## IMUL signed number multiplication

- Similar in operation to the unsigned multiplication.
  - Operands in signed number operations can be positive or negative, and the result must indicate the sign.

**Table 6-2: Signed Multiplication Summary**

Multiplication	Operand 1	Operand 2	Result
byte $\times$ byte	AL	register or memory	AX <sup>1</sup>
word $\times$ word	AX	register or memory	DX AX <sup>2</sup>
word $\times$ byte	AL = byte CBW	register or memory	DX AX <sup>2</sup>

*Notes:* 1. CF = 1 and OF = 1 if AH has part of the result, but if the result is not large enough to need the AH, the sign bit is copied to the unused bits and the CPU makes CF = 0 and OF = 0 to indicate that.  
2. CF = 1 and OF = 1 if DX has part of the result, but if the result is not large enough to need the DX, the sign bit is copied to the unused bits and the CPU makes CF = 0 and OF = 0 to indicate that. One can use the J condition to find out which of the conditions above has occurred. The rest of the flags are undefined.

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## IMUL signed number multiplication

Program 6-1 computes the average of temperatures:  
+13, -10, +19, +14, -18, -9, +12, -19, +16, Celsius.

```
TITLE      PROG 6-1 FIND THE AVERAGE TEMPERATURE
PAGE      60,132
          .MODEL SMALL
          .STACK 64

;-----
          .DATA
SIGN_DAT  DB +13,-10,+19,+14,-18,-9,+12,-19,+16
          ORG 0010H
AVERAGE  DW ?
REMAINDER DW ?
;-----
          .CODE
MAIN PROC FAR
    MOV    AX,@DATA
    MOV    DS,AX
    MOV    CX,9           ;LOAD COUNTER
    SUB    BX,BX          ;CLEAR BX, USED AS SUM
    MOV    SI,OFFSET SIGN_DAT
    ;-----
```

Each data byte was sign-extended and added to BX, computing the sum, which is a signed word. The sum & count were sign-extended, and by dividing the total sum by the count (9), the average was calculated.

**See the entire program listing on page 182 of your textbook.**

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## SAR/SAL



- SAR destination, count
  - As the bits of the destination are shifted to the right into CF, the empty bits are filled with the sign bit.

```
MOV    AL, -10        ;AL=-10=F6H=1111 0110
SAR     AL, 1          ;AL is arithmetic shifted right once
                        ;AL=1111 1011=FDH=-5
```

**Example 6-11** Using DEBUG, evaluate the results of the following:

```
MOV     AX, -9
MOV     BL, 2
IDIV    BL             ;divide -9 by 2 results in FCH
MOV     AX, -9
SAR     AX, 1          ;divide -9 by 2 with arithmetic shift
                        ;results in FBH
```

### Solution:

The DEBUG trace demonstrates that an IDIV of  $-9$  by  $2$  gives FCH ( $-4$ ), whereas SAR  $-9$  gives FBH ( $-5$ ). This is because SAR rounds negative numbers down but IDIV rounds up.



# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## SAR/SAL



- SAL (shift arithmetic left) and SHL (shift left) do exactly the same thing.
  - Basically the same instruction with two mnemonics.
  - As far as signed numbers are concerned, there is no need for SAL.

# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## CMP signed number comparison

- CMP dest, source
  - The same for both signed and unsigned numbers.
    - The J condition instruction used to make a decision is different from that used for the unsigned numbers.
- In unsigned number comparisons CF and ZF are checked for conditions of larger, equal, and smaller.
  - In signed number comparison, OF, ZF, SF are checked.

destination > source	OF=SF or ZF=0
destination = source	ZF=1
destination < source	OF=negation of SF



# 6.1: SIGNED NUMBER ARITHMETIC OPERATIONS

## CMP signed number comparison

Mnemonics used to detect the conditions above:

JG	Jump Greater	jump if OF=SF or ZF=0
JGE	Jump Greater or Equal	jump if OF=SF
JL	Jump Less	jump if OF=inverse of SF
JLE	Jump Less or Equal	jump if OF=inverse of SF or ZF=1
JE	Jump if Equal	jump if ZF = 1

- Example 6-12 on page 185 should help clarify how the condition flags are affected by the compare instruction.
- Program 6-2 on page 186 is an example of the application of the signed number comparison.
  - It uses the data in Program 6-1, and finds the lowest temperature.

## 6.2: STRING AND TABLE OPERATIONS

- String instructions in the x86 family are capable of operations on a series of operands located in consecutive memory locations.
- While CMP can compare only 2 bytes (or words) of data, CMPS (compare string) compares two arrays of data located in memory locations.
  - Pointed at by the SI and DI registers.

## 6.2: STRING AND TABLE OPERATIONS

### SI & DI, DS & ES in string instructions

- For string operations to work, CPU designers set aside certain registers for specific functions, to permanently provide source/destination operands.
  - In 8088/86 processors, SI & DI registers always point to the source and destination operands, respectively.
- To generate the physical address, 8088/86 always uses SI as the offset of the DS (data segment) register. and DI as the offset of ES (extra segment).
  - The ES register must be initialized for the string operation to work.

## 6.2: STRING AND TABLE OPERATIONS

### byte/word operands in string instructions

The operand can be a byte or a word, distinguished by letters B (byte) & W (word) in the mnemonic.

**Table 6-3: String Operation Summary**

Instruction	Mnemonic	Destination	Source	Prefix
move string byte	MOVSB	ES:DI	DS:SI	REP
move string word	MOVSW	ES:DI	DS:SI	REP
store string byte	STOSB	ES:DI	AL	
store string word	STOSW	ES:DI	AX	REP
load string byte	LODSB	AL	DS:SI	none
load string word	LODSW	AX	DS:SI	none
compare string byte	CMPSB	ES:DI	DS:SI	REPE/REPNE
compare string word	CMPSW	ES:DI	DS:SI	REPE/REPNE
scan string byte	SCASB	ES:DI	AL	REPE/REPNE
scan string word	SCASW	ES:DI	AX	REPE/REPNE

## 6.2: STRING AND TABLE OPERATIONS

### DF, the direction flag

- To process operands in consecutive locations requires the pointer to be incremented/decremented.
  - In string operations, achieved by the direction flag.
    - Flag register bit 11 (D10) is set aside for the direction flag (DF).
- The programmer specifies the choice of increment or decrement by setting the direction flag high or low.
  - CLD (clear direction flag) will reset (zeroes) DF, telling the string instruction to increment the pointers automatically.
    - Sometimes is referred to as *autoincrement*.
  - STD (set the direction flag) performs the opposite function.
    - Sets DF to 1, indicating that the pointers SI and DI should be decremented automatically.

## 6.2: STRING AND TABLE OPERATIONS

### REP prefix

- The REP (repeat) prefix allows a string instruction to perform the operation repeatedly.
  - REP assumes that CX holds the number of times the instruction should be repeated. (until CX becomes zero)
- In Example 6-13, after transfer of every byte by the MOVSB instruction, both the SI and DI registers are incremented automatically once only (notice CLD).
  - The REP prefix causes the CX counter to decrement, and MOVSB is repeated until CX becomes zero.
    - Both DS and ES are set to the same value.

## 6.2: STRING AND TABLE OPERATIONS

### REP prefix

**Example 6-13** Using string instructions, write a program to transfer a block of 20 bytes of data.

**Solution:**

;in the data segment:

```
DATA1 DB      'ABCDEFGHJKLMNOPQRST'
```

```
        ORG    30H
```

```
DATA2 DB      20 DUP (?)
```

;in the code segment:

```
MOV     AX,@DATA
```

```
MOV     DS,AX           ;INITIALIZE THE DATA SEGMENT
```

```
MOV     ES,AX           ;INITIALIZE THE EXTRA SEGMENT
```

```
CLD                     ;CLEAR DIRECTION FLAG FOR AUTOINCREMENT
```

```
MOV     SI,OFFSET DATA1 ;LOAD THE SOURCE POINTER
```

```
MOV     DI,OFFSET DATA2 ;LOAD THE DESTINATION POINTER
```

```
MOV     CX,20           ;LOAD THE COUNTER
```

```
REP     MOVSB           ;REPEAT UNTIL CX BECOMES ZERO
```



## 6.2: STRING AND TABLE OPERATIONS

### REP prefix

- An alternative solution would change only two lines of code:

```
MOV    CX, 10  
REP    MOVSW
```

- MOVSW will transfer a word (2 bytes) at a time and increment SI & DI registers each twice.
  - REP will repeat that process until CX becomes zero.
  - CX has the value of 10 in it, as 10 words is equal to 20 bytes.



## 6.2: STRING AND TABLE OPERATIONS

### STOS instruction

- STOSB stores the byte in register AL into memory locations pointed at by ES:DI.
  - If DF = 0, then DI is *incremented once*.
  - If DF = 1, then DI is decremented.
- STOSW stores the contents of AX in ES:DI and ES:DI+1 (AL into ES:DI and AH into ES:DI+1)
  - If DF = 0, then DI is *incremented twice*.
  - If DF = 1, then DI is decremented twice

## 6.2: STRING AND TABLE OPERATIONS

### LODS instructions

- LODSB loads the contents of memory locations pointed at by DS:SI into AL.
  - Increments SI once, if DF = 0.
  - Decrements SI once, if DF = 1.
- LODSW loads the contents of memory locations pointed at by DS:SI into AL, and DS:SI+1 into AH.
  - SI is incremented *twice* if DF = 0, else it is decremented twice.
- LODS is *never* used with a REP prefix.

## 6.2: STRING AND TABLE OPERATIONS

### testing memory using STOSB and LODSB

- Example 6-14 on page 189 uses string instructions STOSB & LODSB to test an area of RAM memory.
  - First AAH is written into 100 locations by using word-sized operand AAAAH, and a count of 50.
  - In the test, LODSB brings the contents of memory locations into AL, one by one, and each is eXclusive-ORed with AAH. (register AH has hex value AA).
    - If they are the same, ZF = 1 and the process is continued.
    - Otherwise, the pattern written there by the previous routine is not there and the program will exit.

## 6.2: STRING AND TABLE OPERATIONS

### testing memory using STOSB and LODSB

#### Example 6-14

##### Solution:

Assuming that ES and DS have been assigned in the ASSUME directive, the following is from the code segment:

```
;PUT PATTERN AAAAH IN TO 50 WORD LOCATIONS
MOV    AX,DTSEG                ;INITIALIZE
MOV    DS,AX                  ;DS REG
MOV    ES,AX                  ;AND ES REG
CLD                            ;CLEAR DF FOR INCREMENT
MOV    CX,50                  ;LOAD THE COUNTER (50 WORDS)
MOV    DI,OFFSET MEM_AREA     ;LOAD THE POINTER FOR DESTINATION
MOV    AX,0AAAAH              ;LOAD THE PATTERN
REP    STOSW                  ;REPEAT UNTIL CX=0
;BRING IN THE PATTERN AND TEST IT ONE BY ONE
MOV    SI,OFFSET MEM_AREA     ;LOAD THE POINTER FOR SOURCE
MOV    CX,100                 ;LOAD THE COUNT (COUNT 100 BYTES)
```

**See the entire program listing on page 189 of your textbook.**

## 6.2: STRING AND TABLE OPERATIONS

### REPZ and REPNZ prefixes

- Used with CMPS & SCAS for testing purposes.
  - **REPZ (repeat zero)** - the same as REPE (repeat equal), will repeat the string operation as long as the source and destination operands are equal ( $ZF = 1$ ) or until CX becomes zero.
  - **REP NZ (repeat not zero)** - the same as REPNE (repeat not equal), will repeat the string operation as long as the source and destination operands are not equal ( $ZF = 0$ ) or until CX becomes zero.
  - **CMPS (compare string)** - allows the comparison of two arrays of data pointed at by registers SI & DI.
    - CMPS can test inequality of two arrays using "REPNE CMPSB".

## 6.2: STRING AND TABLE OPERATIONS

### SCASB (scan string)

- SCASB compares each byte of the array pointed at by ES:DI with the contents of the AL register.
  - Depending on which prefix, REPE or REPNE, is used, a decision is made for equality or inequality.
- In Example 6-16, on page 191, the letter "G" is compared with "M".
  - Since they are not equal, DI increments, CX decrements
    - Scanning is repeated until letter "G" is found or the CX register is zero.
- SCASB can search for a character in an array & if found, it will be replaced with the desired character.

## 6.2: STRING AND TABLE OPERATIONS

### XLAT instruction and look-up tables

- A table is commonly referred to as a *look-up table*.
  - To access elements of a table, 8088/86 processors provide the XLAT (translate) instruction.
- Assume a need for a table for the values of  $x^2$ , where  $x$  is between 0 and 9.
  - First the table is generated and stored in memory:

```
SQUR_TABLE  DB      0,1,4,9,16,25,36,49,64,81
```

- To access the square of any number from 0 to 9, by use of XLAT, register BX must have the offset address of the look-up table, and the number whose square is sought must be in the register AL.
  - After XLAT execution, AL will have the square of the number.

## 6.2: STRING AND TABLE OPERATIONS

### XLAT instruction and look-up tables

- To get the square of 5 from the table:

```
MOV    BX,OFFSET SQUR_TABLE ;load the offset address of table
MOV    AL,05                ;AL=05 will retrieve 6th element
XLAT                    ;pull the element out of table
                        ;and put in AL
```

– After execution AL will have 25 (19H), the square of 5.

- XLAT is one instruction, equivalent to the following:

```
SUB    AH,AH                ;AH=0
MOV    SI,AX                ;SI=000X
MOV    AL,[BX+SI]           ;GET THE SIth ENTRY FROM BEGINNING
                        ;OF THE TABLE POINTED AT BY BX
```



## 6.2: STRING AND TABLE OPERATIONS

### code conversion using XLAT

- XLAT can translate the hex keys of non-ASCII keyboards to ASCII.
  - Assuming keys are 0–F, the following is the program to convert the hex digits of 0–F to their ASCII equivalents.

```
;data segment:
ASC_TABL      DB      '0','1','2','3','4','5','6','7','8'
               DB      '9','A','B','C','D','E','F'
HEX_VALU      DB      ?
ASC_VALU      DB      ?
;code segment:
              MOV      BX,OFFSET ASC_TABL ;BX= TABLE OFFSET
              MOV      AL,HEX_VALU        ;AL=THE HEX DATA
              XLAT                     ;GET THE ASCII EQUIVALENT
              MOV      ASC_VALU,AL        ;MOVE IT TO MEMORY
```

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