

# Chapter 4 Boolean Algebra and Logic Simplification

## 4.1 Boolean Operations and Expressions

- Variable: a symbol used to represent a logical quantity;
- Complement: the inverse of a variable
  - Be indicated by a bar over the variable: NOT operation
- Boolean Addition: OR operation
- Boolean Multiplication: equivalent to the AND operation.

## 4.2 Laws and Rules of Boolean Algebra

- Laws of Boolean Algebra

- **Commutative Laws**

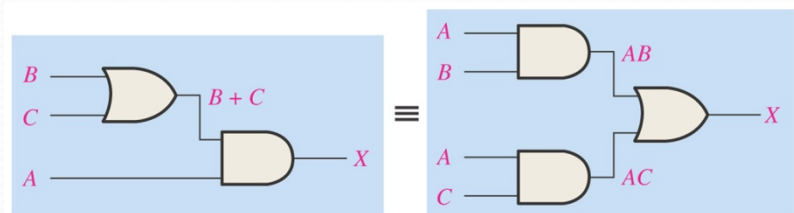
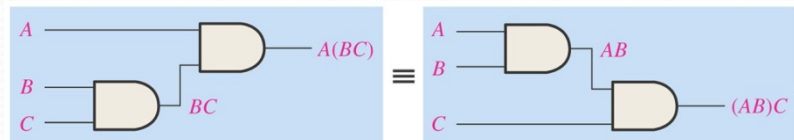
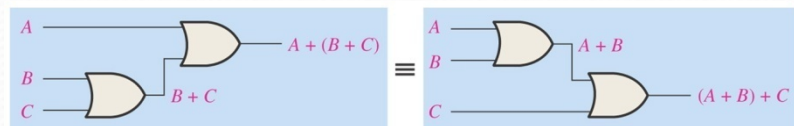
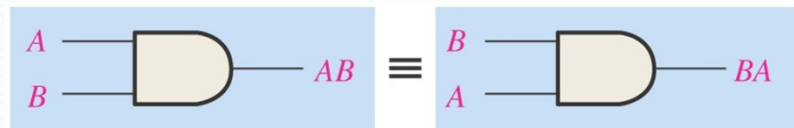
- $A+B=B+A$
- $AB=BA$

- **Associative Laws**

- $A+(B+C)=(A+B)+C$
- $A(BC)=(AB)C$

- **Distributive Law**

- $A(B+C)=AB+AC$



$$X = A(B + C)$$

$$X = AB + AC$$

# Rules of Boolean Algebra

$$0 \bullet A = 0$$

$$1 \bullet A = A$$

$$A \bullet A = A$$

$$A \bullet \bar{A} = 0$$

$$A \bullet B = B \bullet A$$

$$A \bullet (B \bullet C) = (A \bullet B) \bullet C$$

$$A \bullet (B + C) = A \bullet B + A \bullet C$$

$$\overline{A \bullet B} = \bar{A} + \bar{B}$$

$$\overline{\bar{A}} = A$$

$$\bar{1} = 0; \bar{0} = 1$$

$$1 + A = 1$$

$$0 + A = A$$

$$A + A = A$$

$$A + \bar{A} = 1$$

$$A + B = B + A$$

$$A + (B + C) = (A + B) + C$$

$$A + B \bullet C = (A + B) \bullet (A + C)$$

$$\overline{A + B} = \bar{A} \bullet \bar{B}$$



$$A + A \bullet B = A$$

$$A + \bar{A} \bullet B = A + B$$

$$A \bullet B + A \bullet \bar{B} = A$$

$$A \bullet (A + B) = A$$

$$A \bullet B + \bar{A} \bullet C + B \bullet C = A \bullet B + \bar{A} \bullet C$$

$$A \bullet B + \bar{A} \bullet C + B \bullet C \bullet D = A \bullet B + \bar{A} \bullet C$$

$$A \bullet \overline{A \bullet B} = A \bullet \bar{B}$$

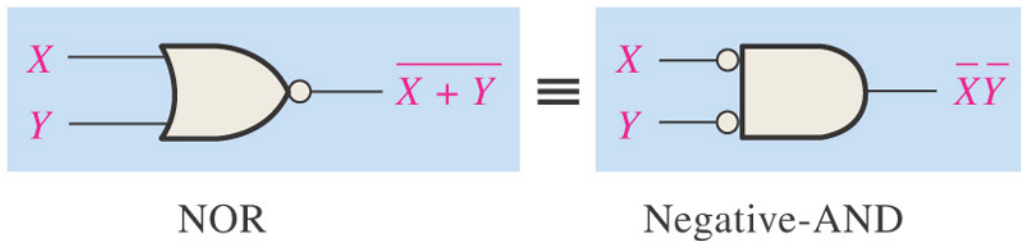
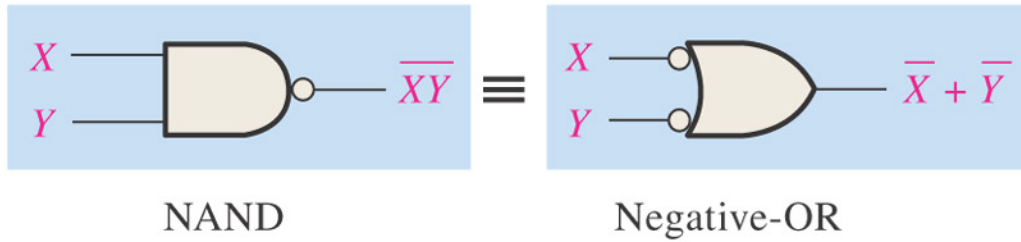
## 4.3 Demorgan's Theorems

- The complement of a product of variables is equal to the sum of the complements of the variables
- The complement of a sum of variables is equal to the product of the complements of the variables

$$\overline{A \bullet B} = \bar{A} + \bar{B}$$

$$\overline{A + B} = \bar{A} \bullet \bar{B}$$

f  
t



Inputs		Output	
X	Y	$\overline{XY}$	$\overline{X} + \overline{Y}$
0	0	1	1
0	1	1	1
1	0	1	1
1	1	0	0

Inputs		Output	
X	Y	$\overline{X + Y}$	$\overline{X}\overline{Y}$
0	0	1	1
0	1	0	0
1	0	0	0
1	1	0	0

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Example: Apply DeMorgan's theorems to the following expressions

$$\overline{XYZ} = ?$$

$$\overline{X + Y + Z} = ?$$

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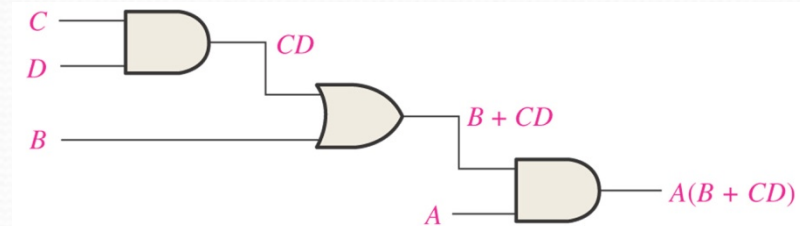
$$\overline{XYZ} = \overline{X} + \overline{Y} + \overline{Z}$$

$$\overline{X + Y + Z} = \overline{X} \overline{Y} \overline{Z}$$



## 4.4 Boolean Analysis of Logic Circuits

- Boolean Expression for a Logic Circuit
- Constructing a Truth Table for a Logic Circuit



A	B	C	D	$A(B+CD)$
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	0
0	1	0	0	0
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	0
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

## 4.5 Simplification Using Boolean Algebra

*Examples :*

$$Y_1 = \overline{\overline{A}BCD} + \overline{A\overline{B}CD}$$

$$Y_2 = \overline{A}\overline{B} + ACD + \overline{\overline{A}B} + \overline{A}CD$$

$$Y_3 = \overline{A}\overline{B}\overline{C} + A\overline{C} + \overline{\overline{B}C}$$

$$Y_4 = \overline{B}\overline{C}D + BC\overline{D} + \overline{B}CD + BCD$$

$$A + \overline{A} = 1$$

*Solutions :*

$$Y_1 = \overline{\overline{A}BCD} + \overline{A}BCD = A(\overline{\overline{BCD}} + \overline{BCD}) = A$$

$$Y_2 = \overline{A}B + ACD + \overline{\overline{A}B} + \overline{A}CD = A(\overline{B} + CD) + \overline{A}(\overline{B} + CD) = \overline{B} + CD$$

$$Y_3 = \overline{A}B\overline{C} + A\overline{C} + \overline{\overline{B}C} = \overline{A}B\overline{C} + (A + \overline{B})\overline{C} = (\overline{A}B)\overline{C} + (\overline{\overline{A}B})\overline{C} = \overline{C}$$

$$\begin{aligned} Y_4 &= \overline{B}CD + B\overline{C}\overline{D} + \overline{\overline{B}CD} + BCD = B(\overline{C}D + C\overline{D}) + B(\overline{\overline{C}\overline{D}} + CD) \\ &= B(C \oplus D) + B(\overline{C \oplus D}) = B \end{aligned}$$



*Examples :*

$$A + 1 = 1$$

$$Y_1 = (\overline{\overline{AB}} + C) ABD + AD$$

$$Y_2 = AB + AB\overline{C} + ABD + AB(\overline{C} + \overline{D})$$

$$Y_3 = A + \overline{\overline{A}} \bullet \overline{\overline{BC}} (\overline{A} + \overline{\overline{BC}} + D) + BC$$



*Solutions :*

$$Y_1 = (\overline{\overline{AB}} + C) ABD + AD = \left[ (\overline{\overline{AB}} + C) B \right] AD + AD = AD$$

$$Y_2 = AB + ABC\overline{C} + ABD + AB(\overline{C} + \overline{D}) = AB + AB \left[ \overline{C} + D + (\overline{C} + \overline{D}) \right] = AB$$

$$\begin{aligned} Y_3 &= A + \overline{\overline{A}} \bullet \overline{\overline{BC}} (\overline{A} + \overline{\overline{BC}} + D) + BC \\ &= (A + BC) + (A + BC) \left( \overline{A} + \overline{\overline{\overline{BC}}} + D \right) = A + BC \end{aligned}$$

*Examples :*

$$Y_1 = \overline{B} + ABC$$

$$Y_2 = A\overline{B} + B + \overline{A}B$$

$$Y_3 = AC + \overline{A}D + \overline{C}D$$

$$A + \overline{A}B = A + B$$

*Solution :*

$$Y_1 = \overline{B} + ABC = \overline{B} + AC$$

$$Y_2 = A\overline{B} + B + \overline{A}B = A + B + \overline{A}B = A + B$$

$$Y_3 = AC + \overline{A}D + \overline{C}D = AC + (\overline{A} + \overline{C}) D = AC + \overline{ACD} \\ = AC + D$$



*Example :*

$$A + A = A$$

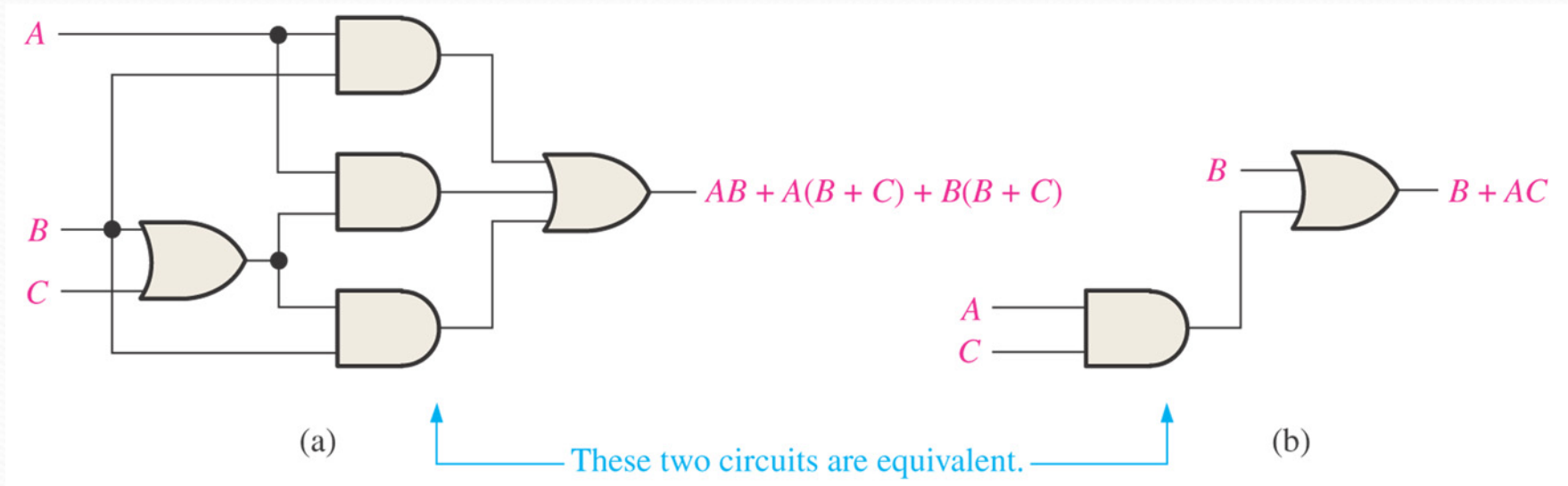
$$Y = A\overline{B} + \overline{A}B + B\overline{C} + \overline{B}C$$



*Solution :*

$$\begin{aligned} Y &= \overline{A}\overline{B} + \overline{A}B(C + \overline{C}) + B\overline{C} + (A + \overline{A})\overline{B}C \\ &= \overline{A}\overline{B} + \overline{A}BC + \overline{A}B\overline{C} + B\overline{C} + A\overline{B}C + \overline{A}\overline{B}C \\ &= (\overline{A}\overline{B} + \overline{A}BC) + (B\overline{C} + \overline{A}B\overline{C}) + (\overline{A}\overline{B}C + \overline{A}\overline{B}C) \\ &= \overline{A}\overline{B} + B\overline{C} + \overline{A}C \end{aligned}$$

**Figure 4–17** Gate circuits for Example 4–8



$$\begin{aligned}
 AB + A(B + C) + B(B + C) &= A(B + B + C) + B(B + C) \\
 &= A(B + C) + B(B + C) = AB + AC + B + BC \\
 &= B(A + 1) + AC + B(1 + C) \\
 &= B + AC
 \end{aligned}$$

## 4.6 Standard Forms of Boolean Expressions

- The sum-of-products (SOP) form
  - A single overbar cannot extend over more than one variable;
  - More than one variable in a term can have an overbar
- The product-of-sums (POS) form
  - A single overbar cannot extend over more than one variable
  - More than one variable in a term can have an overbar

Example:

$$\overline{\overline{A}}\overline{\overline{B}}\overline{C}D + \overline{A}CD + AC$$

$$(A + B)(B + C + D)(A + C).$$



- Conversion of a General Expression to SOP Form
  - Applying the distributive law
- The standard SOP Form
  - **All the variables in the domain appear in each product term in the expression**
- Converting Product Terms to Standard SOP
  - Using  $A + \overline{A} = 1$



*Example :* Convert the following expression to the standard SOP form

$$Y = \overline{A}\overline{B}\overline{C}D + \overline{A}CD + AC$$

*Solution :*

$$Y = \overline{A}\overline{B}\overline{C}D + \overline{A}(B + \overline{B})CD + A(B + \overline{B})C$$

$$= \overline{A}\overline{B}\overline{C}D + \overline{A}BCD + \overline{A}\overline{B}CD + ABC(D + \overline{D}) + ABC(D + \overline{D})$$

$$= \overline{A}\overline{B}\overline{C}D + \overline{A}BCD + \overline{A}\overline{B}CD + ABCD + ABC\overline{D} + \overline{A}BCD + \overline{A}\overline{B}C\overline{D}$$

$$= \sum_i m_i (i = 3, 7, 9, 10, 11, 14, 15)$$

*Example*  $\overline{A}\overline{B}C + \overline{A}\overline{B} + A\overline{B}\overline{C}D$

$$\overline{A}\overline{B}C = \overline{A}\overline{B}C(D + \overline{D}) = \overline{A}\overline{B}CD + \overline{A}\overline{B}C\overline{D}$$

$$\overline{A}\overline{B} = \overline{A}\overline{B}(C + \overline{C})(D + \overline{D}) = \overline{A}\overline{B}CD + \overline{A}\overline{B}C\overline{D} + \overline{A}\overline{B}\overline{C}D + \overline{A}\overline{B}\overline{C}\overline{D}$$

$$\overline{A}\overline{B}C + \overline{A}\overline{B} + A\overline{B}\overline{C}D$$

$$= \overline{A}\overline{B}CD + \overline{A}\overline{B}C\overline{D} + \overline{A}\overline{B}\overline{C}D + \overline{A}\overline{B}\overline{C}\overline{D} + A\overline{B}\overline{C}D + \overline{A}\overline{B}\overline{C}D + \overline{A}\overline{B}\overline{C}\overline{D}$$

- The Standard POS Form

- All the variables in the domain appear in each sum term in the expression.

- Converting a Sum Term to Standard POS

- Use

$$A\overline{A} = 0 \quad A + BC = (A + B)(A + C)$$

*Example*

$$(A + \overline{B} + C)(\overline{B} + C + \overline{D})$$

$$= (A + \overline{B} + C + D\overline{D})(A\overline{A} + \overline{B} + C + \overline{D})$$

$$= (A + \overline{B} + C + D)(A + \overline{B} + C + \overline{D})(A + \overline{B} + C + \overline{D})(\overline{A} + \overline{B} + C + \overline{D})$$



# Converting standard SOP to standard POS

$$Y = \overline{\overline{A}\overline{B}\overline{C}} + \overline{\overline{A}\overline{B}C} + \overline{\overline{A}B\overline{C}} + \overline{\overline{A}BC} + \overline{ABC} \quad \begin{array}{l} \text{(standard SOP form)} \\ \text{(standard POS form)} \end{array}$$
$$= ?$$

$$\because \overline{Y} + Y = 1$$

$$\overline{\overline{A}\overline{B}\overline{C}} + \overline{\overline{A}\overline{B}C} + \overline{\overline{A}B\overline{C}} + \overline{\overline{A}BC} + \overline{ABC} + \overline{\overline{A}B\overline{C}} + \overline{\overline{A}BC} + \overline{ABC} = 1$$

$$Y = \overline{\overline{A}\overline{B}\overline{C}} + \overline{\overline{A}\overline{B}C} + \overline{\overline{A}B\overline{C}} + \overline{\overline{A}BC} + \overline{ABC}$$

$$\therefore \overline{Y} = \overline{\overline{A}\overline{B}\overline{C}} + \overline{\overline{A}\overline{B}C} + \overline{\overline{A}B\overline{C}}$$

$$Y = \overline{\overline{\overline{A}\overline{B}\overline{C}} + \overline{\overline{A}\overline{B}C} + \overline{\overline{A}B\overline{C}}} = \left( A + B + \overline{C} \right) \left( \overline{A} + B + C \right) \left( \overline{A} + \overline{B} + C \right)$$





## **4.7 Boolean Expressions and Truth Tables**

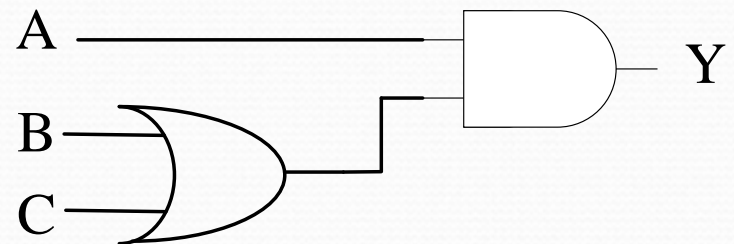
## Truth Table

Input			Output Y
A	B	C	
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

## Boolean Expression

$$Y = A(B + C)$$

## Implementation



## From Truth Table to Boolean Expression

Input			Output Y
A	B	C	
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

→  $A\bar{B}C$

→  $AB\bar{C}$

→  $ABC$

$$\begin{aligned}Y &= A\bar{B}C + AB\bar{C} + ABC \\&= A\bar{B}C + ABC + AB\bar{C} + ABC \\&= AC(\bar{B} + B) + AB(\bar{C} + C) \\&= AC + AB \\&= A(B + C)\end{aligned}$$

# SOP, POS expressions and truth table

Input			Output	P-Term	S-Term
A	B	C	X		
0	0	0	0		$(A + B + C)$
0	0	1	1	$\overline{A}\overline{B}C$	
0	1	0	0		$(A + \overline{B} + C)$
0	1	1	0		$(A + \overline{B} + \overline{C})$
1	0	0	1	$A\overline{B}\overline{C}$	
1	0	1	0		$(\overline{A} + B + \overline{C})$
1	1	0	0		$(\overline{A} + \overline{B} + C)$
1	1	1	1	$ABC$	

$$X = ?$$





## 4.8 The Karnaugh Map

## Example

Simplify the following Boolean expression

$$Y = AC + \overline{B}C + B\overline{D} + C\overline{D} + A(B + \overline{C}) + \overline{A}BC\overline{D} + \overline{A}BDE$$

$$Y = AC + \overline{B}C + B\overline{D} + C\overline{D} + A(B + \overline{C}) + \overline{A}BC\overline{D} + \overline{A}BDE$$

$$= AC + \overline{B}C + B\overline{D} + C\overline{D} + A(B + \overline{C}) + \overline{A}BDE$$

$$= AC + \overline{B}C + B\overline{D} + C\overline{D} + \overline{A}BC + \overline{A}BDE$$

$$= AC + \overline{B}C + B\overline{D} + C\overline{D} + A + \overline{A}BDE$$

$$= A + \overline{B}C + B\overline{D} + C\overline{D}$$

$$= A + \overline{B}C + B\overline{D}$$

$$C\overline{D} + \overline{A}BC\overline{D} = C\overline{D}(1 + \overline{A}B) = C\overline{D}$$

$$A(B + \overline{C}) = \overline{\overline{A}BC}$$

$$\overline{B}C + \overline{\overline{A}BC} = \overline{B}C + A$$

$$AC + A + \overline{A}BDE = A$$

$$\overline{B}C + B\overline{D} + C\overline{D} = \overline{B}C + B\overline{D}$$

# How do you think about this method?

- Many Boolean rules are used
- A little hard
- ...
- Are there any other methods?

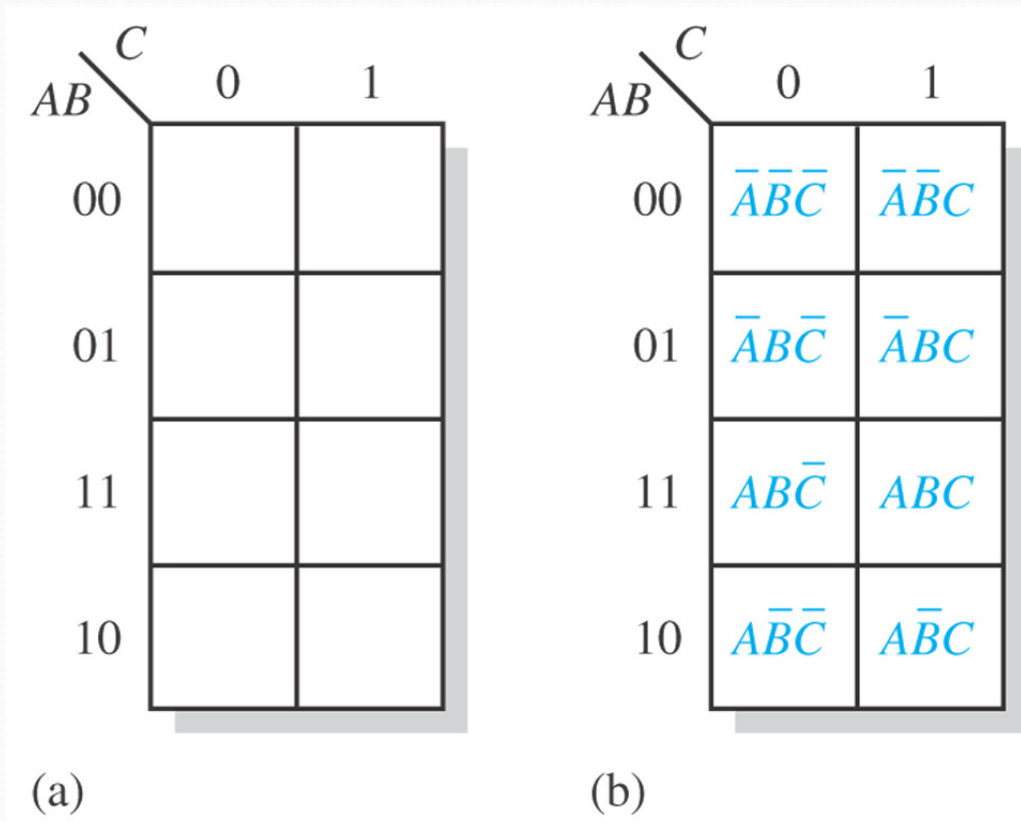


# Karnaugh Map

- A systematic method for simplifying Boolean expressions
- Produce the simplest SOP or POS expression
- Presents all of the possible values of input variables
  - An array of cells
  - Each cell represents a binary value of the input variables
  - Adjacency in position equivalents to adjacency in Boolean algebra



# The Construction of Karnaugh Map



**Figure 4-21** A 3-variable Karnaugh map showing product terms.

# The Construction of Karnaugh Map

AB \ CD				
	00	01	11	10
00				
01				
11				
10				

(a)

AB \ CD				
	00	01	11	10
00	$\bar{A}\bar{B}\bar{C}\bar{D}$	$\bar{A}\bar{B}\bar{C}D$	$\bar{A}\bar{B}C\bar{D}$	$\bar{A}\bar{B}CD$
01	$\bar{A}B\bar{C}\bar{D}$	$\bar{A}B\bar{C}D$	$\bar{A}BC\bar{D}$	$\bar{A}BCD$
11	$AB\bar{C}\bar{D}$	$AB\bar{C}D$	$ABC\bar{D}$	$ABCD$
10	$A\bar{B}\bar{C}\bar{D}$	$A\bar{B}\bar{C}D$	$A\bar{B}C\bar{D}$	$A\bar{B}CD$

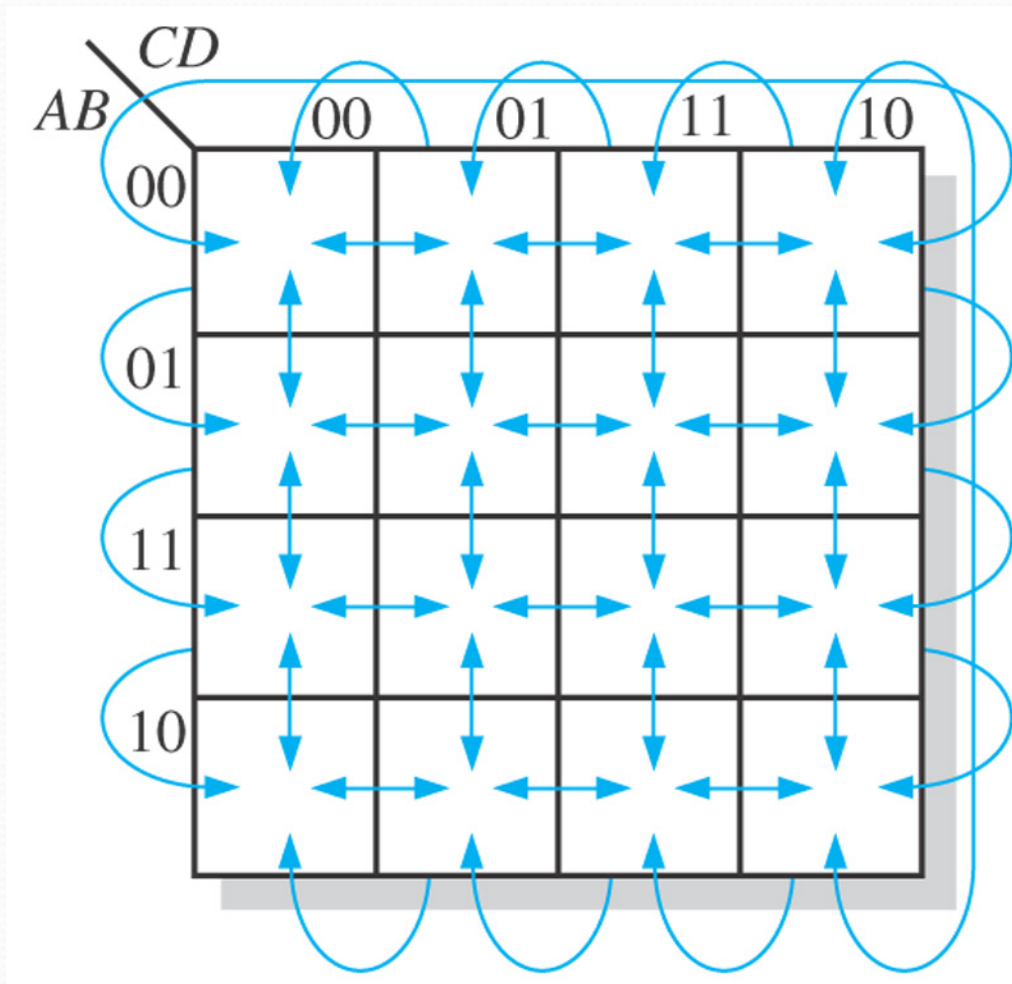
(b)

**Figure 4–22** A 4-variable Karnaugh map.

# Cell Adjacency

- Adjacency (in logic): a single-variable change
  - Cells that differ by only one variable are adjacency
  - Cells that differ by more than one variable are not adjacency
- Adjacency (in location)
  - Cells locate next to others

# Cell Adjacency



**Figure 4–23** Adjacent cells on a Karnaugh map are those that differ by only one variable. Arrows point between adjacent cells.



		<i>B</i>	
		0	1
<i>A</i>	0	$A'B'$ $m_0$	$A'B$ $m_1$
	1	$AB'$ $m_2$	$AB$ $m_3$

(a)

		<i>BC</i>			
		00	01	11	10
<i>A</i>	0	$m_0$	$m_1$	$m_3$	$m_2$
	1	$m_4$	$m_5$	$m_7$	$m_6$

(b)

		<i>CD</i>			
		00	01	11	10
<i>AB</i>	00	$m_0$	$m_1$	$m_3$	$m_2$
	01	$m_4$	$m_5$	$m_7$	$m_6$
	11	$m_{12}$	$m_{13}$	$m_{15}$	$m_{14}$
	10	$m_8$	$m_9$	$m_{11}$	$m_{10}$

(c)

		<i>CDE</i>							
		000	001	011	010	110	111	101	100
<i>AB</i>	00	$m_0$	$m_1$	$m_3$	$m_2$	$m_6$	$m_7$	$m_5$	$m_4$
	01	$m_8$	$m_9$	$m_{11}$	$m_{10}$	$m_{14}$	$m_{15}$	$m_{13}$	$m_{12}$
	11	$m_{24}$	$m_{25}$	$m_{27}$	$m_{26}$	$m_{30}$	$m_{31}$	$m_{29}$	$m_{28}$
	10	$m_{16}$	$m_{17}$	$m_{19}$	$m_{18}$	$m_{22}$	$m_{23}$	$m_{21}$	$m_{20}$

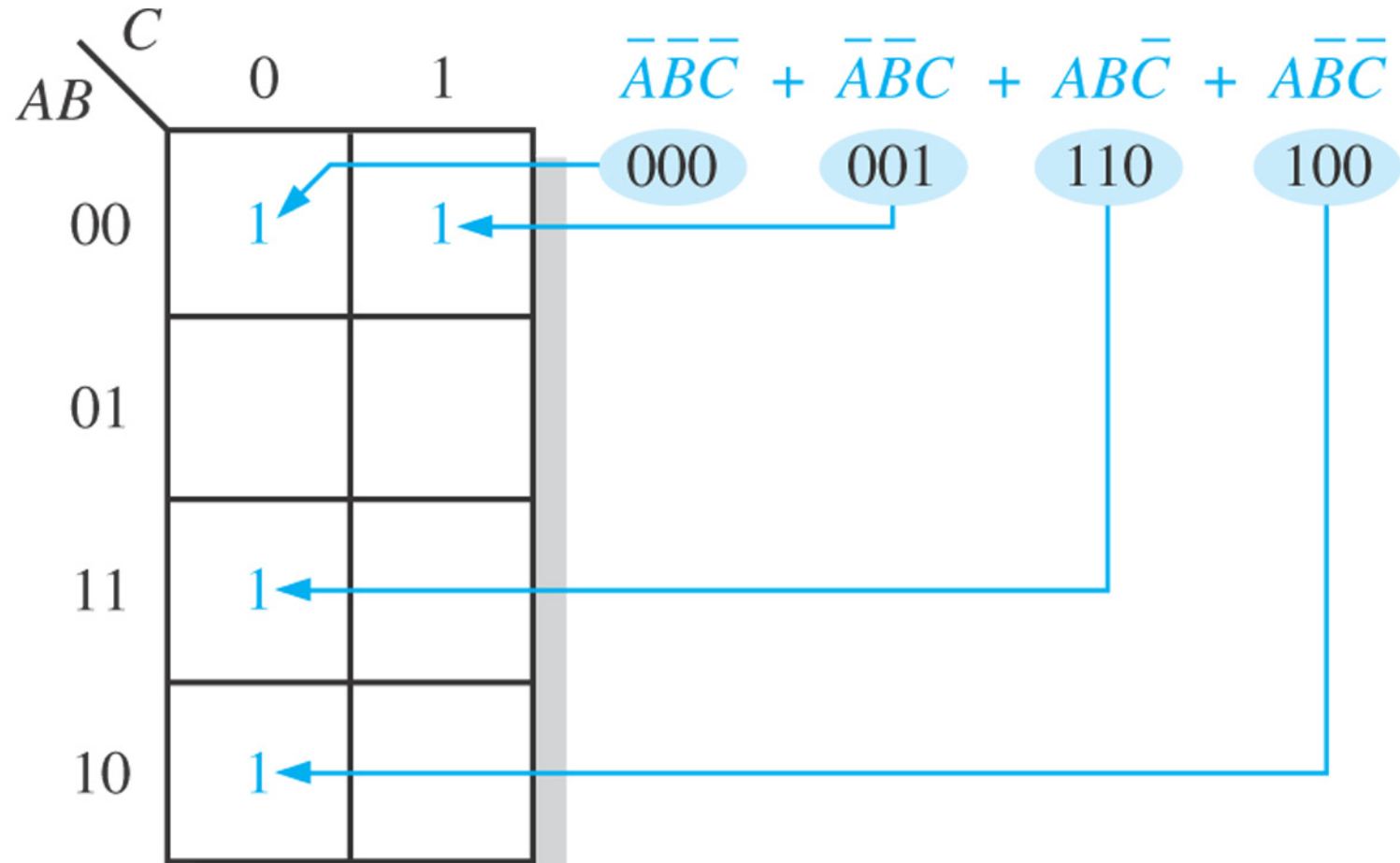
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## 4.9 Karnaugh Map SOP Minimization

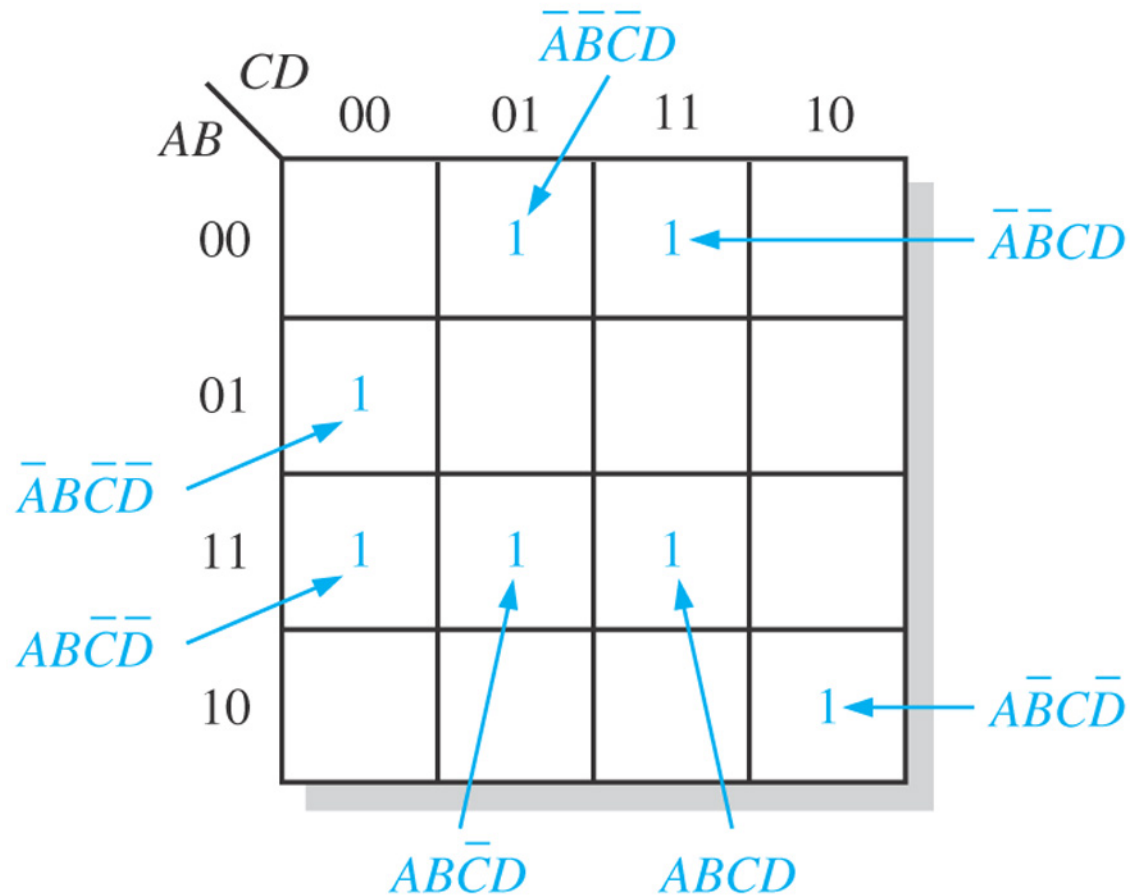
- Mapping a SOP Expression
- Karnaugh Map Simplification of SOP Expressions

## Example: Mapping a standard SOP expression





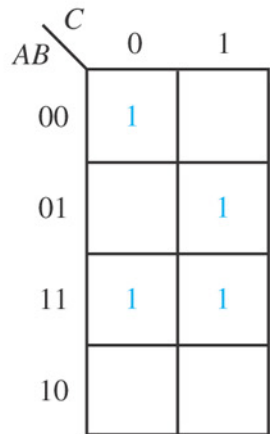
## Another mapping example



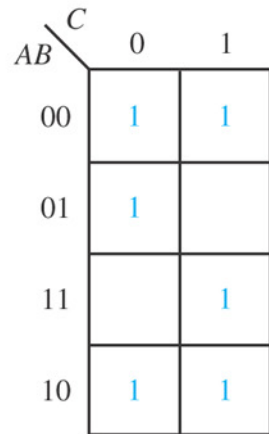
# Karnaugh Map Simplification of SOP Expressions

- Group the 1s
  - Maximize the size of the groups
  - Minimize the number of groups
- Rules
  - A group must contain  $2^n$  cells
  - Each cell must be adjacent to one or more cells in that group
  - Include 1s as much as possible
  - Each 1 on the map must be included at least one group
  - Cell with 1 can be included into more than one group

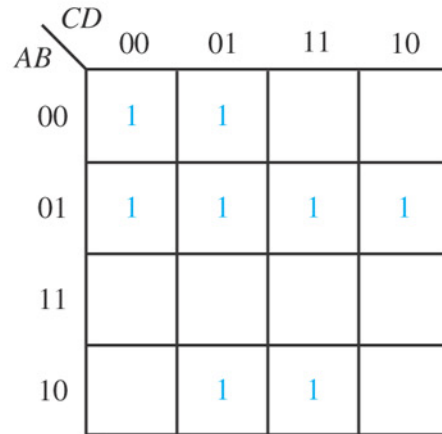
## Example: Group the 1s in each of the Karnaugh maps in the following



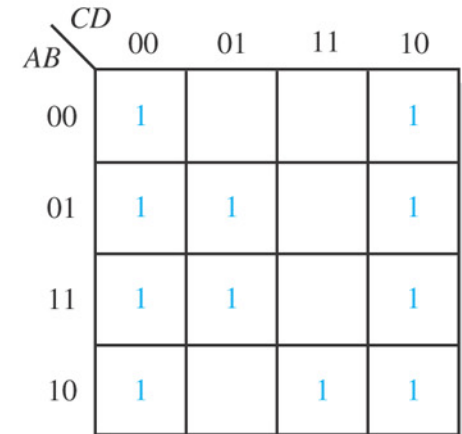
(a)



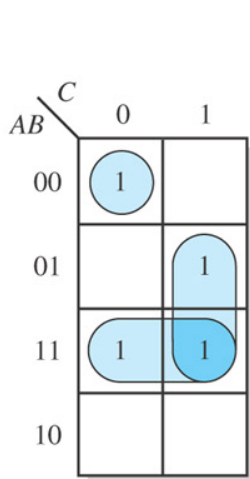
(b)



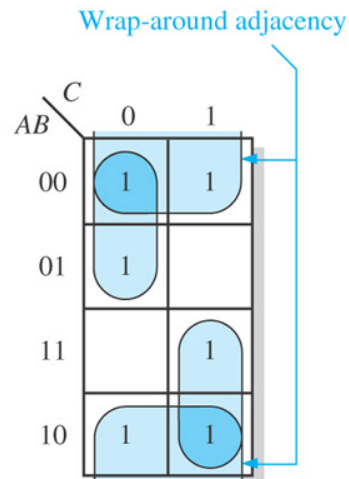
(c)



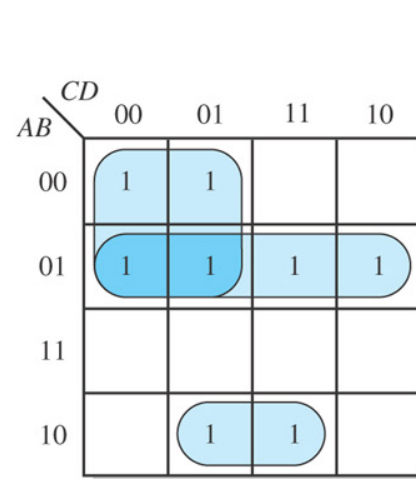
(d)



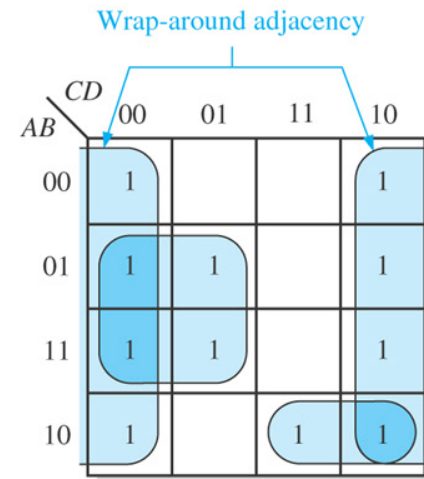
(a)



(b)

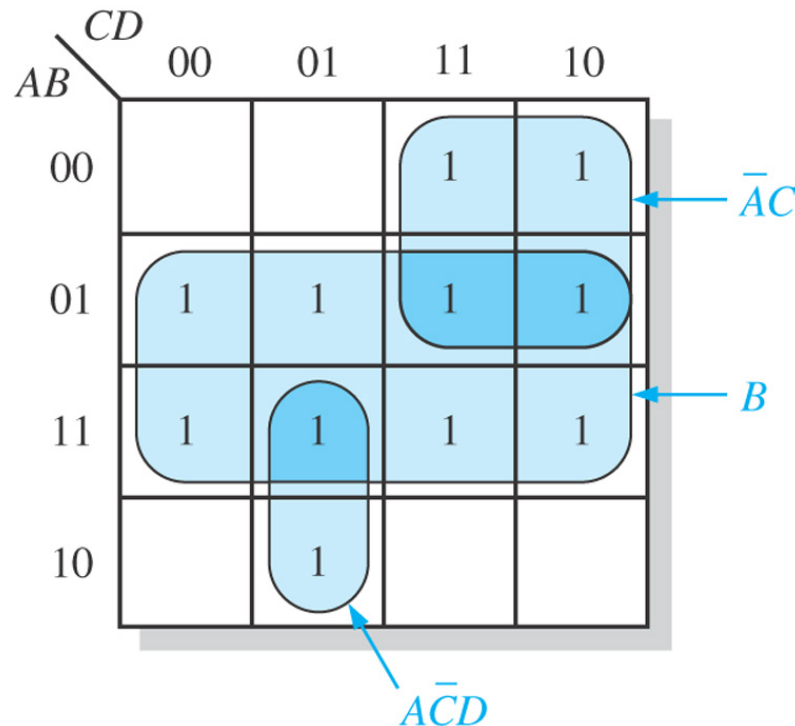


(c)



(d)

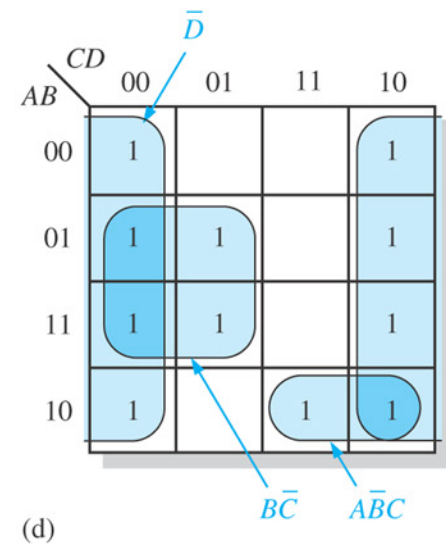
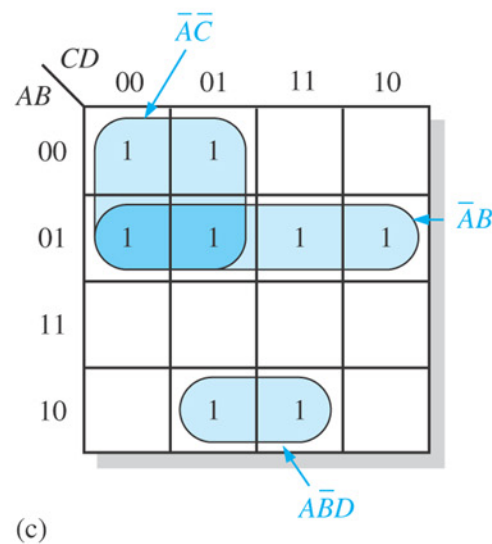
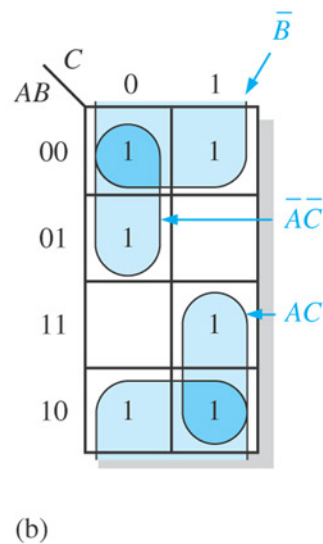
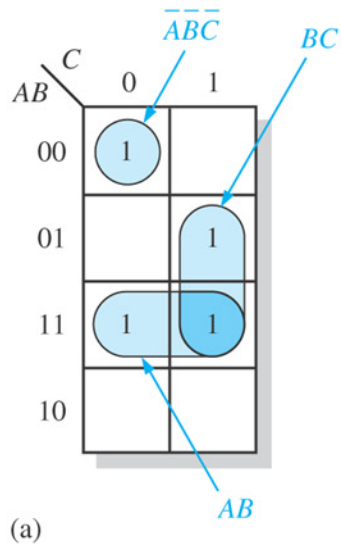
**Example: Determine the product terms for the Karnaugh map in the following figure and write the resulting minimum SOP expression.**



$$B + \overline{A}C + A\overline{C}D$$



**Example: Determine the product terms for the Karnaugh map in the following figure and write the resulting minimum SOP expression.**



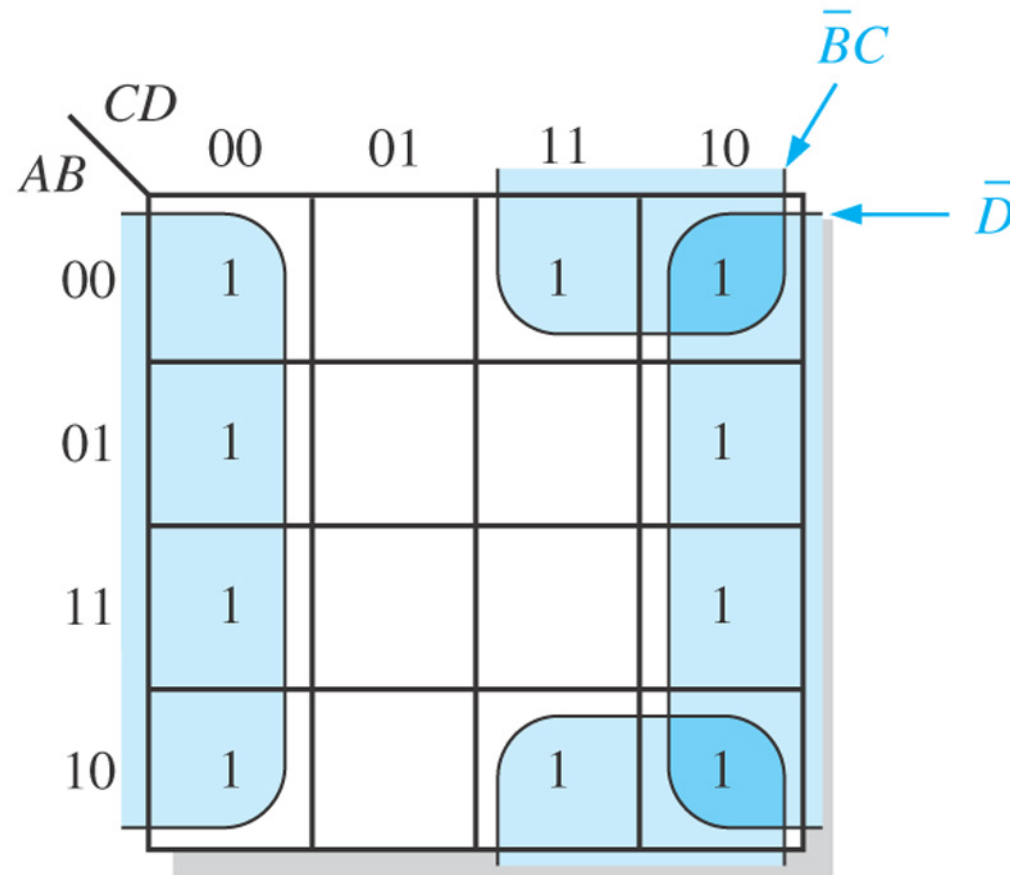
$$(a) AB + BC + \overline{A}\overline{B}\overline{C}$$

$$(b) \overline{B} + \overline{A}\overline{C} + AC$$

$$(c) \overline{A}B + \overline{A}\overline{C} + \overline{A}BD$$

$$(d) \overline{D} + \overline{A}\overline{B}C + B\overline{C}$$

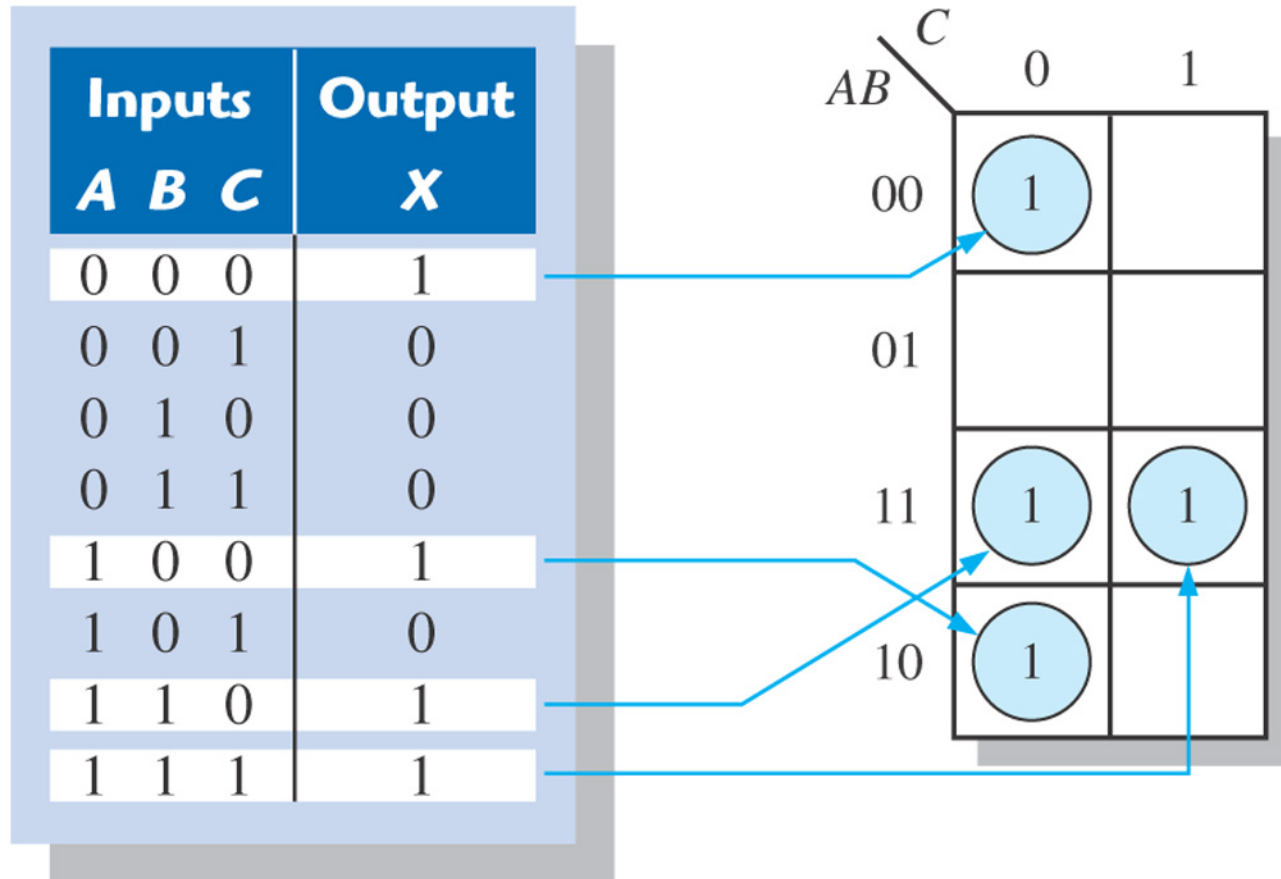
**Example:** Determine the product terms for the Karnaugh map in the following figure and write the resulting minimum SOP expression.



$$\bar{D} + \bar{B}C$$

**Figure 4–35** Example of mapping directly from a truth table to a Karnaugh map.

$$X = \bar{A}\bar{B}\bar{C} + A\bar{B}\bar{C} + ABC\bar{C} + ABC$$



- Don't care conditions
  - Some input variable combinations are not allowed
    - These unallowed states will never occur in application
    - They can be treated as “don't care” terms
  - “don't care” terms either a 1 or 0 may be assigned to the output

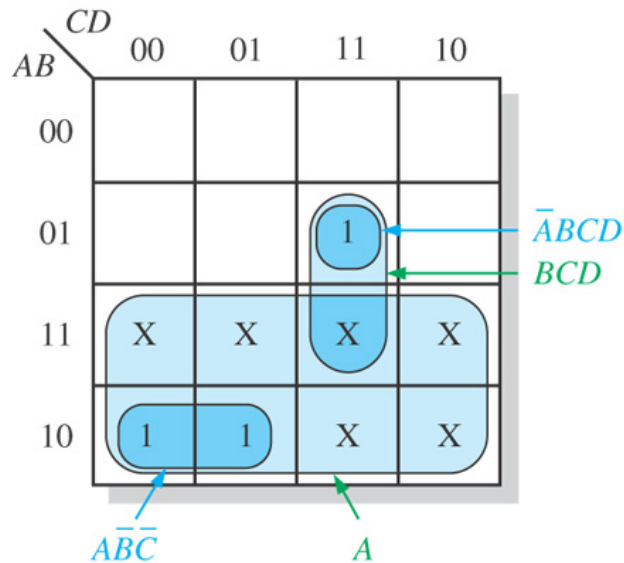


**Figure 4–36** Example of the use of “*don’t care*” conditions to simplify an expression.

Inputs <i>ABCD</i>	Output <i>Y</i>
0 0 0 0	0
0 0 0 1	0
0 0 1 0	0
0 0 1 1	0
0 1 0 0	0
0 1 0 1	0
0 1 1 0	0
0 1 1 1	1
1 0 0 0	1
1 0 0 1	1
1 0 1 0	X
1 0 1 1	X
1 1 0 0	X
1 1 0 1	X
1 1 1 0	X
1 1 1 1	X

(a) Truth table

Don't cares



- (b) Without “don’t cares”  $Y = A\bar{B}\bar{C} + \bar{A}BCD$   
 With “don’t cares”  $Y = A + BCD$

# Summary

- Boolean Algebra
  - Variable
  - Operation
  - Laws and rules
- Simplification of Boolean expression
- SOP form and POS form
- Karnaugh maps

# HW

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