

# APPENDIX R

# PROOF OF THE RSA

# ALGORITHM

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The basic elements of the RSA algorithm can be summarized as follows. Given two prime numbers  $p$  and  $q$ , with  $n = pq$  and a message block  $M < n$ , two integers  $e$  and  $d$  are chosen such that

$$M^{ed} \bmod n = M$$

We state in Section 9.2, that the preceding relationship holds if  $e$  and  $d$  are multiplicative inverses modulo  $\phi(n)$ , where  $\phi(n)$  is the Euler totient function. It is shown in Chapter 8 that for  $p, q$  prime,  $\phi(pq) = (p - 1)(q - 1)$ . The relationship between  $e$  and  $d$  can be expressed as

$$ed \bmod \phi(n) = 1$$

Another way to state this is that there is an integer  $k$  such that  $ed = k\phi(n) + 1$ . Thus, we must show that

$$M^{k\phi(n) + 1} \bmod n = M^{k(p-1)(q-1) + 1} \bmod n = M \quad \textbf{(R.1)}$$

## R.1 BASIC RESULTS

Before proving Equation (R.1), we summarize some basic results. In Chapter 4, we showed that a property of modular arithmetic is the following

$$[(a \bmod n) \times (b \bmod n)] \bmod n = (a \times b) \bmod n$$

From this, it should be easy to see that if we have  $x \bmod n = 1$ , then  $x^2 \bmod n = 1$  and, for any integer  $y$ , we have  $x^y \bmod n = 1$ . Similarly, if we have  $x \bmod n = 0$  for any integer  $y$ , we have  $x^y \bmod n = 0$ .

Another property of modular arithmetic is

$$[(a \bmod n) - (b \bmod n)] \bmod n = (a - b) \bmod n$$

The other result we need is Euler's theorem, which was developed in Chapter 8. If integers  $a$  and  $n$  are relatively prime, then  $a^{\phi(n)} \bmod n = 1$ .

## R.2 PROOF

First we show that  $M^{k(p-1)(q-1)+1} \bmod p = M \bmod p$ . There are two cases to consider.

**Case 1:**  $M$  and  $p$  are not relatively prime; that is,  $p$  divides  $M$ . In this case  $M \bmod p = 0$  and therefore  $M^{k(p-1)(q-1)+1} \bmod p = 0$ . Thus,  
 $M^{k(p-1)(q-1)+1} \bmod p = M \bmod p$ .

**Case 2:** If  $M$  and  $p$  are relatively prime, by Euler's theorem,  
 $M^{\phi(p)} \bmod p = 1$ . We proceed as

$$\begin{aligned} M^{k(p-1)(q-1)+1} \bmod p &= [(M)M^{k(p-1)(q-1)}] \bmod p \\ &= [(M)(M^{(p-1)})^{k(q-1)}] \bmod p \\ &= [(M)(M^{\phi(p)})^{k(q-1)}] \bmod p \\ &= (M \bmod p) \times [(M^{\phi(p)} \bmod p)^{k(q-1)}] \\ &= (M \bmod p) \times (1)^{k(q-1)} \quad (\text{by Euler's theorem}) \\ &= M \bmod p \end{aligned}$$

We now observe that

$$[M^{k(p-1)(q-1)+1} - M] \bmod p = [M^{k(p-1)(q-1)+1} \bmod p] - [M \bmod p] = 0$$

Thus,  $p$  divides  $[M^{k(p-1)(q-1)+1} - M]$ . By the same reasoning, we can show that  $q$  divides  $[M^{k(p-1)(q-1)+1} - M]$ . Because  $p$  and  $q$  are distinct primes, there must exist an integer  $r$  that satisfies

$$[M^{k(p-1)(q-1)+1} - M] = (pq)r = nr$$

Therefore,  $n$  divides  $[M^{k(p-1)(q-1)+1} - M]$ , and so  $M^{k\phi(n)+1} \bmod n = M^{k(p-1)(q-1)+1} \bmod n = M$ .