Chapter 13: I/O Systems







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- I/O Hardware
- Application I/O Interface
- Kernel I/O Subsystem
- Transforming I/O Requests to Hardware Operations
- Streams
- Performance





Objectives

- Explore the structure of an operating system's I/O subsystem
- Discuss the principles of I/O hardware and its complexity
- Provide details of the performance aspects of I/O hardware and software



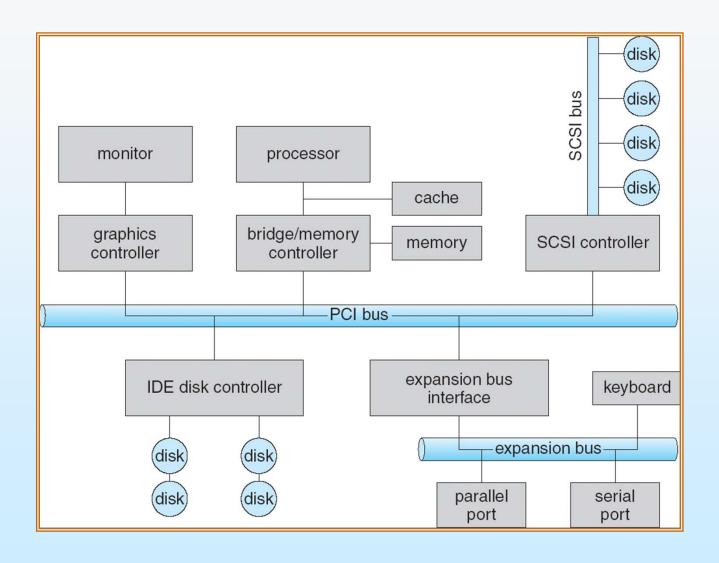
I/O Hardware

- Incredible variety of I/O devices
- Common concepts
 - Port
 - Bus (daisy chain or shared direct access)
 - Controller (host adapter)
- I/O instructions control devices
- Devices have addresses, used by
 - Direct I/O instructions
 - Memory-mapped I/O





A Typical PC Bus Structure







Device I/O Port Locations on PCs (partial)

I/O address range (hexadecimal)	device	
000-00F	DMA controller	
020–021	interrupt controller	
040–043	timer	
200–20F	game controller	
2F8–2FF	serial port (secondary)	
320-32F	hard-disk controller	
378–37F	parallel port	
3D0-3DF	graphics controller	
3F0-3F7	diskette-drive controller	
3F8-3FF	serial port (primary)	



Polling

- Determines state of device
 - command-ready
 - busy
 - Error
- Busy-wait cycle to wait for I/O from device





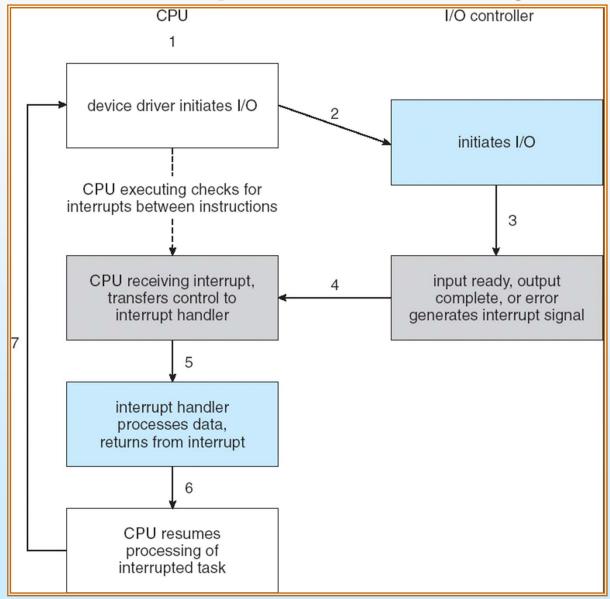
Interrupts

- CPU Interrupt-request line triggered by I/O device
- Interrupt handler receives interrupts
- Maskable to ignore or delay some interrupts
- Interrupt vector to dispatch interrupt to correct handler
 - Based on priority
 - Some nonmaskable
- Interrupt mechanism also used for exceptions





Interrupt-Driven I/O Cycle







Intel Pentium Processor Event-Vector Table

vector number	description	
0	divide error	
1	debug exception	
2	null interrupt	
3	breakpoint	
4	INTO-detected overflow	
5	bound range exception	
6	invalid opcode	
7	device not available	
8	double fault	
9	coprocessor segment overrun (reserved)	
10	invalid task state segment	
11	segment not present	
12	stack fault	
13	general protection	
14	page fault	
15	(Intel reserved, do not use)	
16	floating-point error	
17	alignment check	
18	machine check	
19–31	(Intel reserved, do not use)	
32–255	maskable interrupts	



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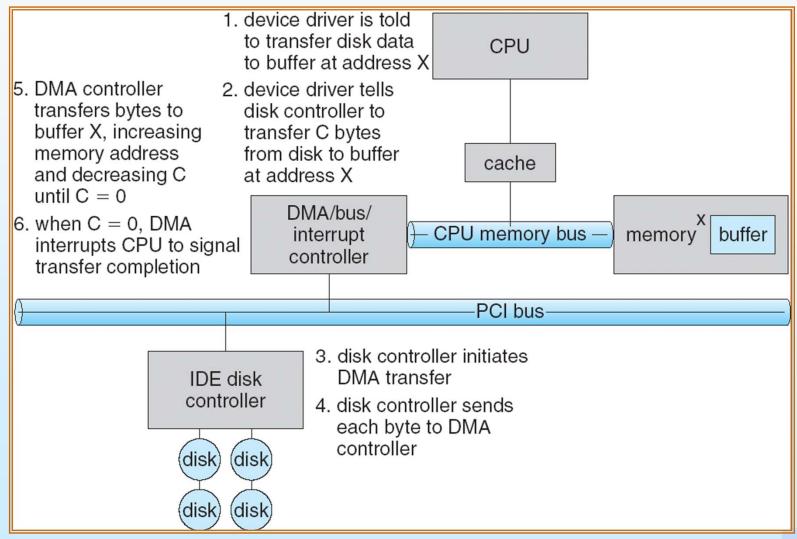
Direct Memory Access

- Used to avoid programmed I/O for large data movement
- Requires DMA controller
- Bypasses CPU to transfer data directly between I/O device and memory





Six Step Process to Perform DMA Transfer





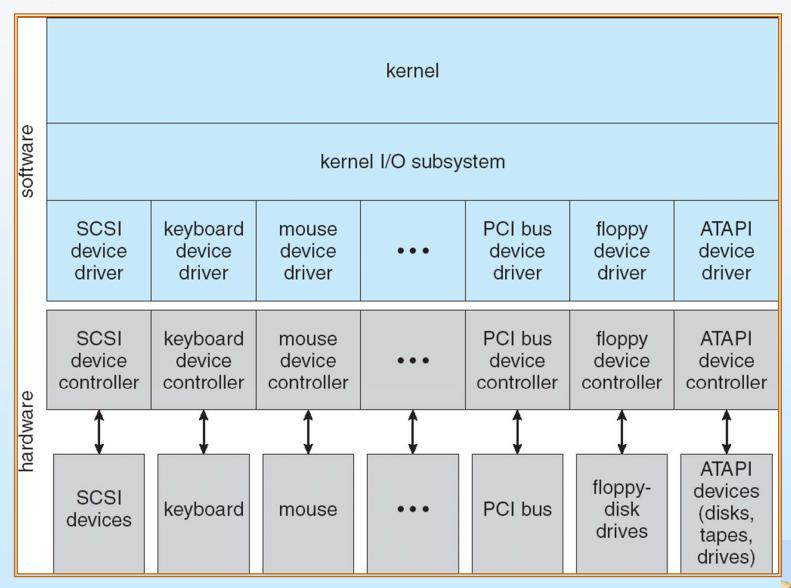
Application I/O Interface

- I/O system calls encapsulate device behaviors in generic classes
- Device-driver layer hides differences among I/O controllers from kernel
- Devices vary in many dimensions
 - Character-stream or block
 - Sequential or random-access
 - Sharable or dedicated
 - Speed of operation
 - read-write, read only, or write only





A Kernel I/O Structure





Characteristics of I/O Devices

aspect	variation	example
data-transfer mode	character block	terminal disk
access method	sequential random	modem CD-ROM
transfer schedule	synchronous asynchronous	tape keyboard
sharing	dedicated sharable	tape keyboard
device speed	latency seek time transfer rate delay between operations	
I/O direction	read only write only read–write	CD-ROM graphics controller disk





Block and Character Devices

- Block devices include disk drives
 - Commands include read, write, seek
 - Raw I/O or file-system access
 - Memory-mapped file access possible
- Character devices include keyboards, mice, serial ports
 - Commands include get, put
 - Libraries layered on top allow line editing





Network Devices

- Varying enough from block and character to have own interface
- Unix and Windows NT/9x/2000 include socket interface
 - Separates network protocol from network operation
 - Includes select functionality
- Approaches vary widely (pipes, FIFOs, streams, queues, mailboxes)





Clocks and Timers

- Provide current time, elapsed time, timer
- Programmable interval timer used for timings, periodic interrupts
- ioctl (on UNIX) covers odd aspects of I/O such as clocks and timers



Blocking and Nonblocking I/O

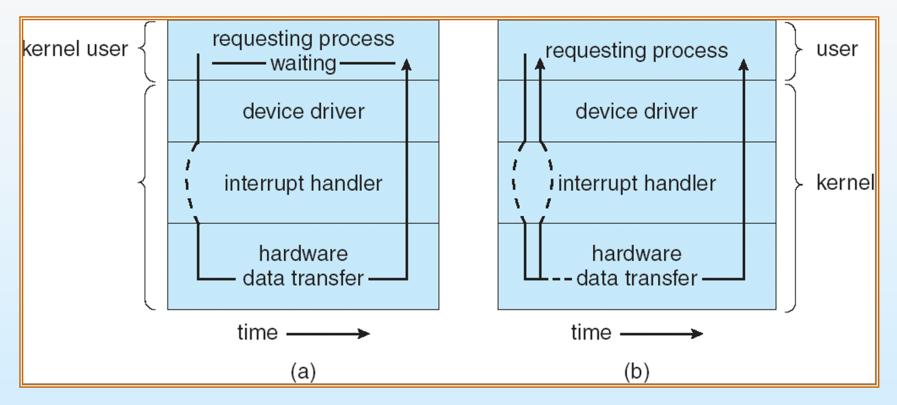
- Blocking process suspended until I/O completed
 - Easy to use and understand
 - Insufficient for some needs
- Nonblocking I/O call returns as much as available
 - User interface, data copy (buffered I/O)
 - Implemented via multi-threading
 - Returns quickly with count of bytes read or written
- Asynchronous process runs while I/O executes
 - Difficult to use
 - I/O subsystem signals process when I/O completed



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Two I/O Methods



Synchronous

Asynchronous





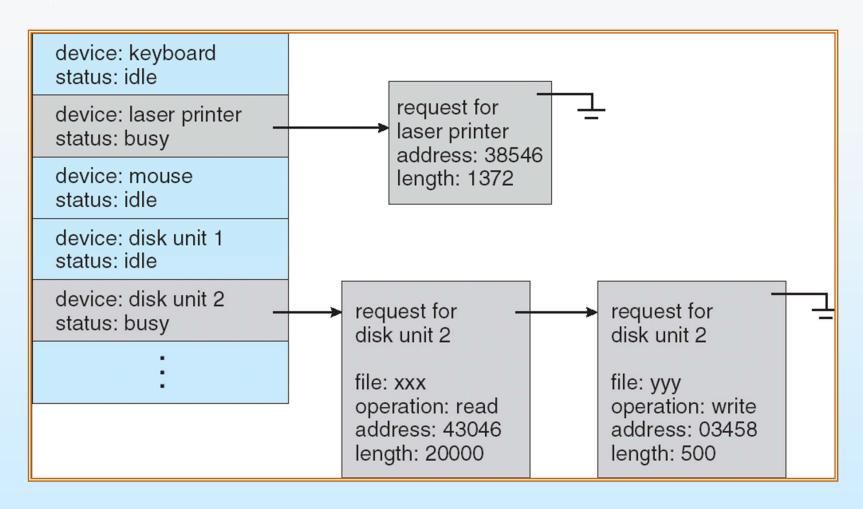
Kernel I/O Subsystem

- Scheduling
 - Some I/O request ordering via per-device queue
 - Some OSs try fairness
- Buffering store data in memory while transferring between devices
 - To cope with device speed mismatch
 - To cope with device transfer size mismatch
 - To maintain "copy semantics"



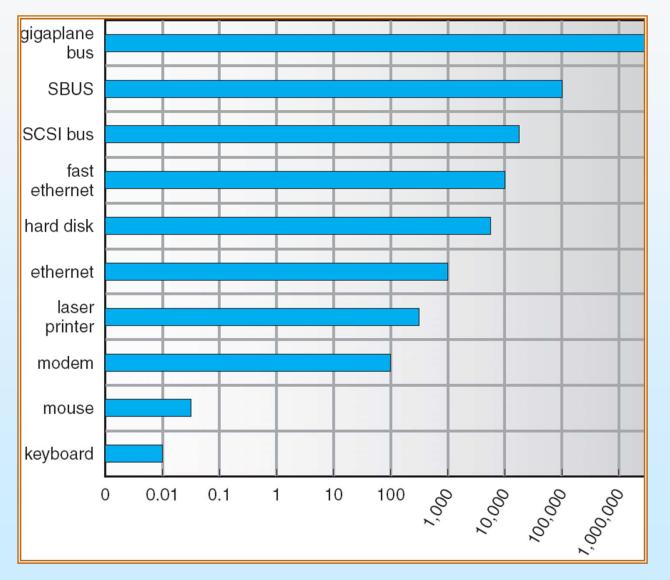


Device-status Table





Sun Enterprise 6000 Device-Transfer Rates







Kernel I/O Subsystem

- Caching fast memory holding copy of data
 - Always just a copy
 - Key to performance
- Spooling hold output for a device
 - If device can serve only one request at a time
 - i.e., Printing
- **Device reservation** provides exclusive access to a device
 - System calls for allocation and deallocation
 - Watch out for deadlock





Error Handling

- OS can recover from disk read, device unavailable, transient write failures
- Most return an error number or code when I/O request fails
- System error logs hold problem reports





Improving Performance

- Reduce number of context switches
- Reduce data copying
- Reduce interrupts by using large transfers, smart controllers, polling
- Use DMA
- Balance CPU, memory, bus, and I/O performance for highest throughput



End of Chapter 13



