

Building Whiteboxes: attacks and defenses

—
Matthieu Daumas (me), Adrien Guinet, Philippe Teuwen

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- ▶ Wait, protect what? the keys? the data?

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We need to better define our context. . .
...and our **attacker model**

Attacker Model

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To sum up: think of the **worst possible situation!**

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In a nutshell...

Example: standard AES

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1 C=AES.new("this is some key")  
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1 C=WBAES.new() # extracting a key from WBAES is **hard enough**  
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Whiteboxes: introduction

Introduction (2/2)

- ▶ Idea: generate cryptographic algorithms that process encoded ² keys

²read this as “somewhat concealed from such kind of omniscient attacker”

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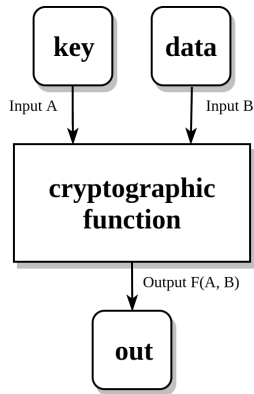
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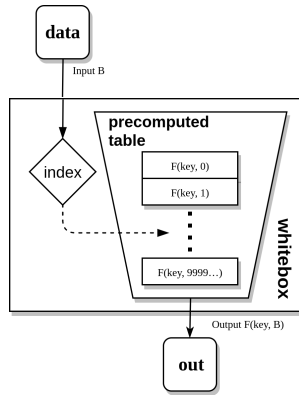
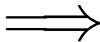
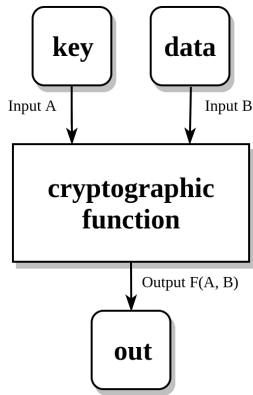
Making the attacker's work *hard enough*

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Building Whiteboxes with Tables

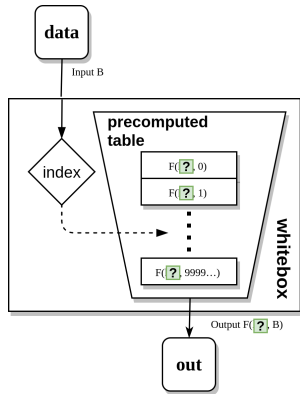
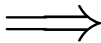
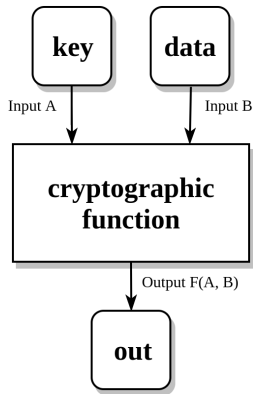


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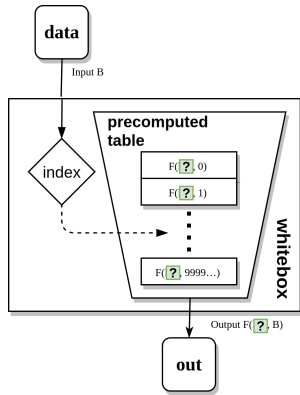
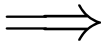
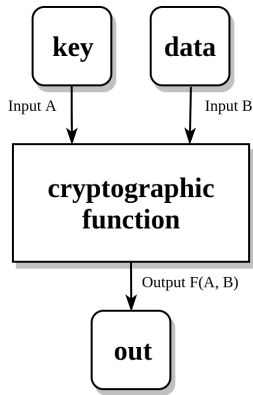
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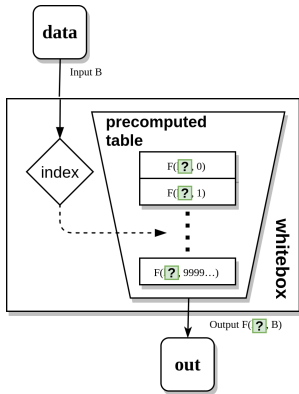
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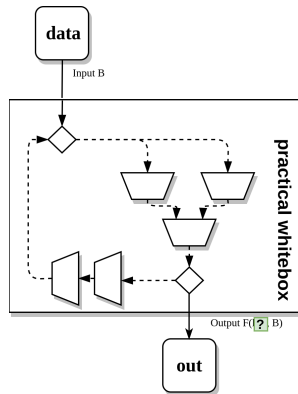
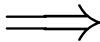
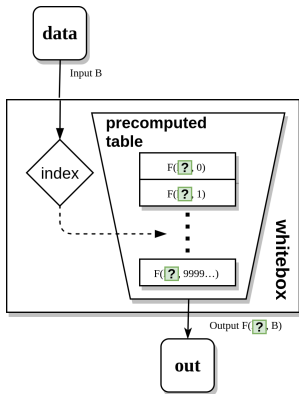
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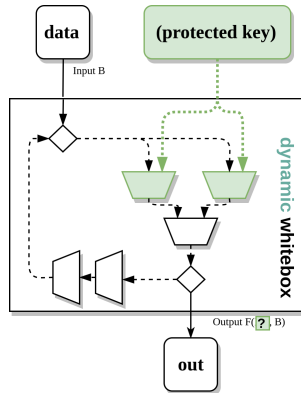
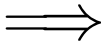
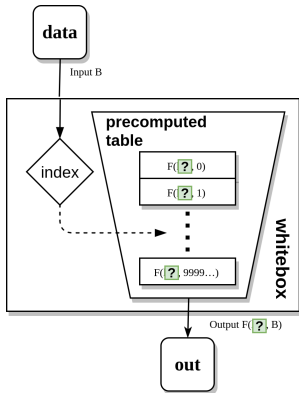
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Building Whiteboxes with Tables



Tabulate **smaller operations** on **internal states**

Building Whiteboxes with Tables



Tabulate **smaller operations** on **internal states**
(you can even make a **dynamic** whitebox)

Building Whiteboxes: Tradeoffs

Practical whiteboxes

- ▶ Work on *encoded* internal states of the algorithm

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The **unprotected** OpenSSL AES implementation (`aes_core.c`)

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The goal is to give the user the choice of a **tradeoff**!

Whiteboxes: Usage

Digital Right Managements

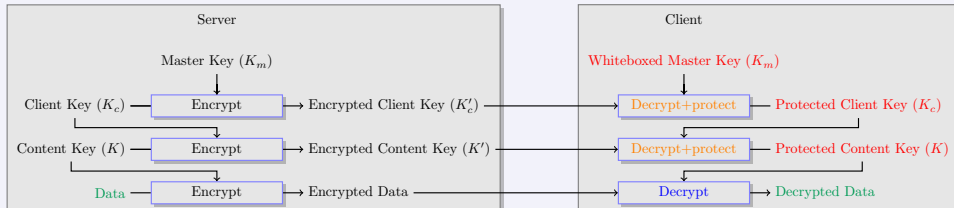


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- Code Lifting

- Dataflow (for complex schemes)

- Case Study: AES

- Attacking ECDSA

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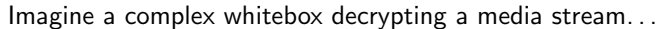
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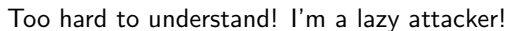
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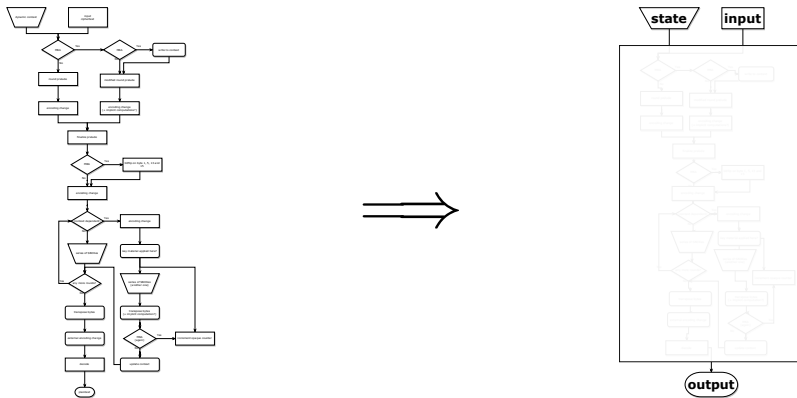
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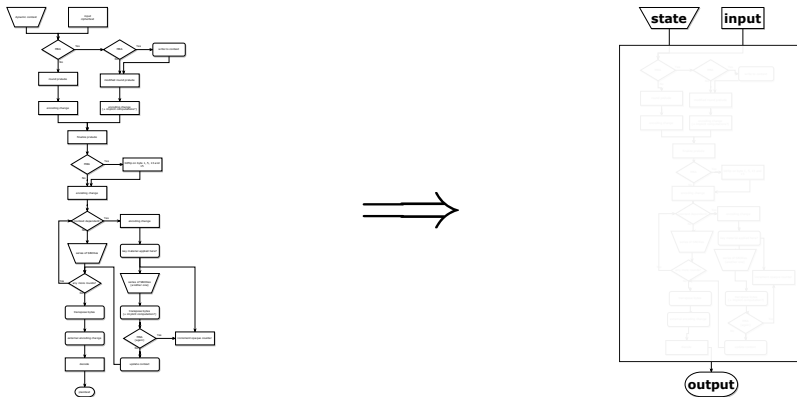


Code Lifting: Example



Too hard to understand! I'm a lazy attacker! No problem!

Code Lifting: Example



Me attacker can steal movie now,



Code Lifting: In Practice

Code Lifting

- ▶ Idea: extract useful high-level primitives from the whitebox

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Ask yourself: is my whitebox protected against code lifting?
...and protections are not trivial to implement ⁴

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⁴especially when integrating with large projects or third-party code

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The dataflow problem

In practice: whiteboxing “end-to-end”

- ▶ Cryptographic schemes involve more **than one primitive**,

The dataflow problem

In practice: whiteboxing “end-to-end”

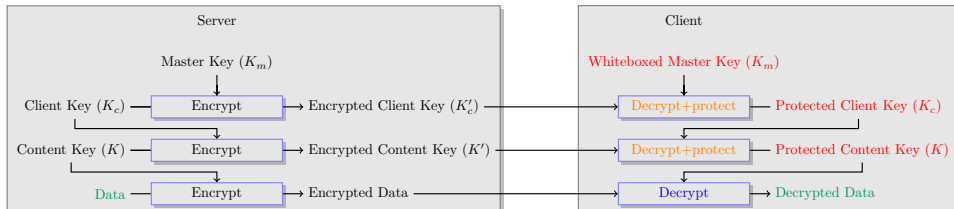
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The dataflow problem

In practice: whiteboxing “end-to-end”

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- ▶ and complex key derivation processes. . .

The dataflow problem: Example

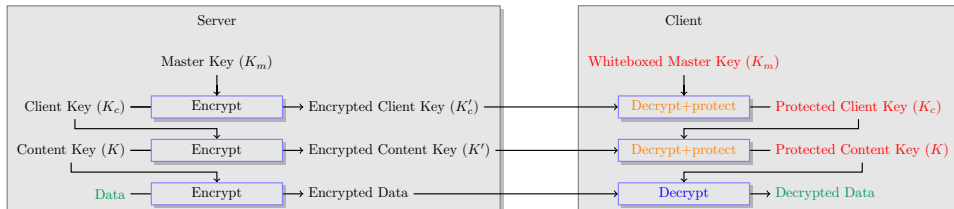


Example

In the example above:

- ▶ K_m is used to **decrypt+protect** K_c (also called *unwrapping*)
- ▶ K_c is used to **decrypt+protect** K
- ▶ K is used to **decrypt** the encrypted data to **plaintext data**

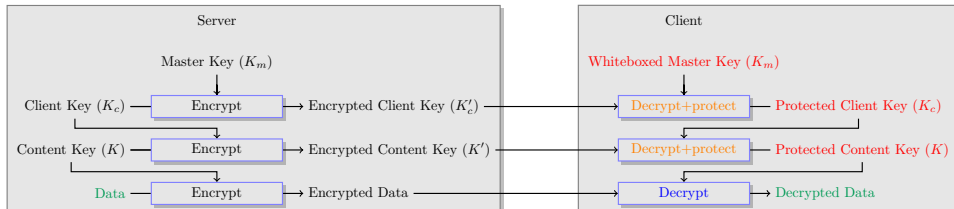
The dataflow problem: Example



In the whitebox code, we thus have **two functions**:

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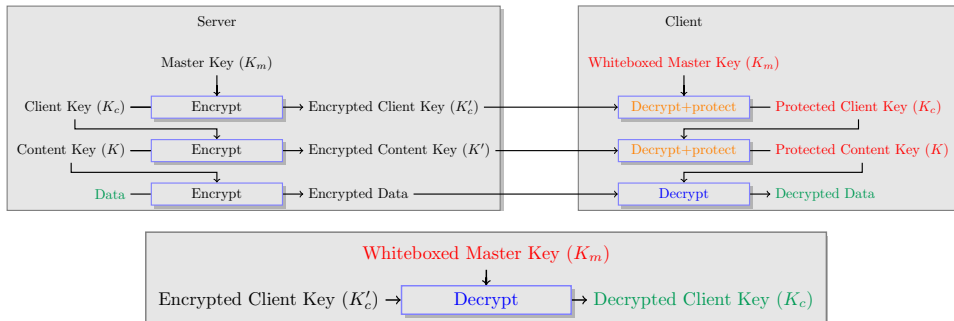
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The attacker **modifies the dataflow** to extract an **unprotected** copy of K_c :



The dataflow problem: Example



The dataflow problem: Solutions

- ▶ Code lifting protections only help to somewhat **mitigate** this problem
- ▶ Tweaking how each key is unwrapped can **enforce** where it is supposed to be used

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Ask yourself: is my whitebox dataflow properly enforced?

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- Design-Specific

- Generic Attacks – Tracing

- Generic Attacks – Fault Injection

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Can we do **cheaper**?

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Greybox Attacks

Bringing back the **Greybox Attack Model** from the hardware world:

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- ▶ Attack Campaign **Automation** 🙄

Tracing



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script-kiddie p0w3r 

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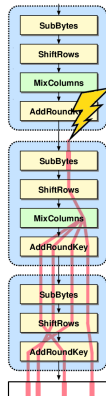
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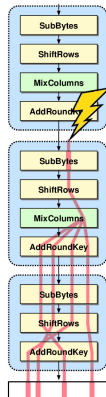


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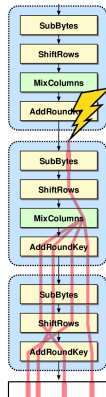


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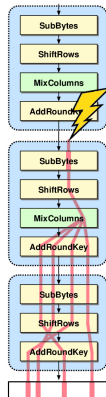
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script-kiddie p0w3r (again)



Counter-Measures

“Ask yourself” checklist:

Counter-Measures

“**Ask yourself**” checklist:

- ▶ is my whitebox protected against **code lifting**?

Counter-Measures

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If not protected **explicitly**, my whitebox will likely be an easy break!

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- Case Study: AES

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Enter The Void

What about whiteboxing more exotic cryptography. . .

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Elliptic Curve Digital Signature Algorithm

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How to attack an “ECDSA whitebox”?

Just as AES: **Inspire yourself from hardware attacks**

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How to attack an “ECDSA whitebox”?

(but first, a crash-course on ECDSA)

Signing with ECDSA

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How to attack an “ECDSA whitebox”?

Let's read literature on **ECDSA hardware attacks!**

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ECDSA vs Fault Injection

ECDSA Textbook Fault Injection

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Bonus: works on actual designs commercially available!

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Spoilers: ECDSA tricks you don't know about
(the last one will surprise you)

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protecting the thing that protects your secrets

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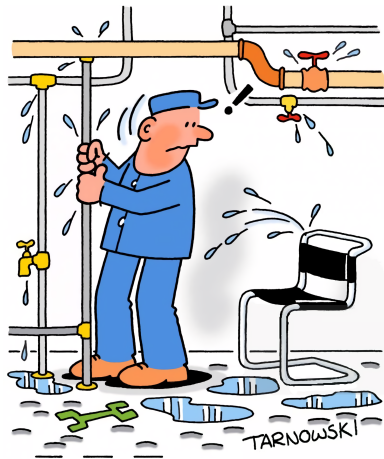
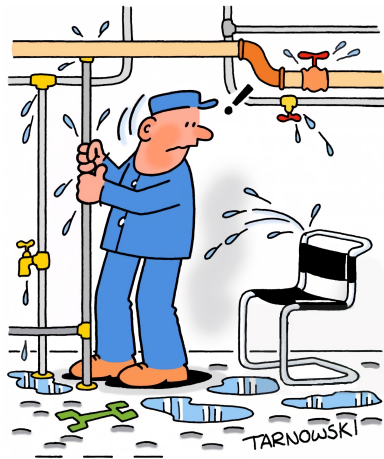


Fig. 1: whitebox cryptography

Conclusion



*Fig. 1: whitebox cryptography
(without proper software protections)*

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Questions?

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