Appendix. Online Supplement

A. Detailed Results of IRMS under the Time Limit $\hat{t} = 7200$ Seconds

Table 1 summarizes the detailed results of IRMS under a long time limit $\hat{t}=3600$ seconds. In Table 1, columns 1-3 describe for each instance its name (Instance), K value, and best known value (BKV) reported in the literature, respectively. Columns 4-6 report the results of IRMS, including the best result (i.e., \hat{f}) found during 30 runs, the average result (i.e., \bar{f}), and average computation time (i.e., \bar{t}) at each run. From it, we observe that IRMS also performs excellent performance. In particular, it finds new upper bounds for nine instances, and matches previous best-known upper bounds on 23 instances.

Table 1 Results of IRMS on Synthetic and Real-world Benchmarks under $\hat{t}=7200$ Seconds

| | | | IRMS | | | | |
|--------------------|--------------------|------------------------|-----------------------|-------------------------|------------------|--|--|
| Instance | K | BKV | \hat{f} | $ar{f}$ | \bar{t} | | |
| BA500 | 50 | 195* | 195 | 195.0 | 0.0 | | |
| BA1000 | 75 | 558* | 558 | 558.0 | 1.5 | | |
| BA2500 | 100 | 3704* | 3704 | 3704.0 | 4.5 | | |
| BA5000 | 150 | 10196* | 10196 | 10196.0 | 18.8 | | |
| ER235 | 50 | 295 | 295 | 295.0 | 7.9 | | |
| ER466 | 80 | 1524 | $\bf 1524$ | 1524.0 | 79.2 | | |
| ER941 | 140 | 5012 | $\boldsymbol{5012}$ | 5023.3 | 1817.8 | | |
| ER2344 | 200 | 902498 | 918952 | 941170.2 | 5673.8 | | |
| FF250 | 50 | 194* | 194 | 194.0 | 0.0 | | |
| FF500 | 110 | 257* | 257 | 257.0 | 0.7 | | |
| FF1000 | 150 | 1260* | 1260 | 1260.0 | 18.5 | | |
| FF2000 | 200 | 4545* | 4545 | 4545.0 | 264.2 | | |
| WS250 | 70 | 3083 | 3083 | 3132.7 | 4127.2 | | |
| WS500 | 125 | 2072 | 2072 | 2078.1 | 813.2 | | |
| WS1000 | 200 | 109677 | 137766 | 144237.3 | 3149.4 | | |
| WS1500 | 265 | 13098 | 13098 | 13103.4 | 2461.6 | | |
| Bovine | 3 | 268 | 268 | 268.0 | 0.0 | | |
| Circuit | 25 | 2099 | 2099 | 2099.0 | 1.3 | | |
| Ecoli | 15 | 806 | 806 | 806.0 | 0.0 | | |
| USAir97 | 33 | 4336 | 4336 | 4596.0 | 1452.8 | | |
| HumanDi | 52 | 1115 | 1115 | 1115.0 | 1.5 | | |
| TreniR | 26 | 918 | 918 | 918.0 | 1.3 | | |
| EU_fli | 119 | 348268 | 348268 | 348268.0 | 749.4 | | |
| openfli | 186 | 26783 | 26875 | 28363.8 | 4400.1 | | |
| yeast1 | 202 | 1412 | 1412 | 1412.0 | 34.0 | | |
| H1000 | 100 | 306349 | 306349 | 308345.0 | 3561.5 | | |
| H2000 | 200 | 1242739 | 1236887* | 1250761.8 | 5352.7 | | |
| H3000a | 300 | 2840690 | 2799868* | 2840491.0 | 6055.2 | | |
| H3000b | 300 | 2837584 | 2794262* | 2821455.9 | 6184.0 | | |
| H3000c | 300 | 2835369 | 2783248* | 2825113.4 | 6682.2 | | |
| H3000d | 300 | 2828492 | 2802615* | 2838392.3 | 5612.8 | | |
| H3000e | 300 | 2843000 | 2798688* | 2841275.9 | 5836.7 | | |
| H4000 | 400 | 5038611 | 4977344* | 5074711.6 | 5894.7 | | |
| H5000 | 500 | 7964765 | 7956481* | 8058861.6 | 6062.2 | | |
| powergr Oclinks | $\frac{494}{190}$ | 15862 611253 | 15863 | 15870.8 | 5764.8 1494.1 | | |
| faceboo | 404 | 420334 | 614467 705403 | 614467.0 723233.5 | 5915.9 | | |
| | - | | | | | | |
| grqc | 524 988 | $\frac{13591}{106276}$ | 13590 * 110352 | 13598.0 114505.1 | 5945.4 6367.1 | | |
| hepth | $\frac{988}{1201}$ | 6155877 | 9309386 | 9543464.4 | 5009.6 | | |
| hepph | 1877 | 53963375 | 9309386 56849750 | 9543464.4 57688316.6 | 5149.1 | | |
| astroph condmat | 2313 | 2298596 | 9359852 | 10415923.6 | 5149.1 4155.8 | | |
| Condinat | 2313 | 2290090 | 9509602 | 10410923.0 | 4100.8 | | |

 $^{^{\}ast}$ Optimal results obtained by branch-and-cut algorithm within 5 days.

* Improved best upper bounds.

B. Detailed Results of CEMCNP under the Time Limit $\hat{t} = 3600$ Seconds

Since the source code of CEMCNP is not available to us, we have re-implemented it according its pseudo code. Detailed comparison between our implemented CEMCNP algorithm and reported CEMCNP algorithm are summarized in Table 2. From it, we observe that our implemented CEMCNP algorithm performs slightly worse performance than the reported results on most of instances. However, some reported results are obtained under a longer computation time than the time limit. For example, at least 4000 seconds are required to achieve the results of H3000e, hepth, hepph and condmat. It is worthy noting that for some instances (e.g, grqc, hepth and condmat), our implemented CEMCNP algorithm achieves better performance than the original version.

Table 2 Detailed Results of CEMCNP on Synthetic and Real-world Benchmarks under $\hat{t}=3600$ Seconds

| | | | Reported 1 | Results | Implemented Results | | |
|----------|------|----------|---------------------|------------------|---------------------|----------------|----------------|
| Instance | K | BKV | \hat{f} | \overline{t} | \hat{f} | \overline{f} | \overline{t} |
| BA500 | 50 | 195 | 195 | 0.0 | 195 | 195.0 | 0.0 |
| BA1000 | 75 | 558 | 558 | 0.2 | 558 | 558.0 | 54.0 |
| BA2500 | 100 | 3704 | 3704 | 0.2 | 3704 | 3704.0 | 347.3 |
| BA5000 | 150 | 10196 | 10196 | 1.8 | 10196 | 10196.0 | 9.6 |
| ER235 | 50 | 295 | 295 | 19.8 | 297 | 302.8 | 0.0 |
| ER466 | 80 | 1524 | $\bf 1524$ | 98.4 | 1569 | 1630.6 | 1.7 |
| ER941 | 140 | 5012 | $\boldsymbol{5012}$ | 694.6 | 5363 | 5635.3 | 2.7 |
| ER2344 | 200 | 902498 | 912346 | 3069.0 | 1012527 | 1060618.6 | 847.9 |
| FF250 | 50 | 194 | 194 | 0.3 | 194 | 194.0 | 178.2 |
| FF500 | 110 | 257 | 257 | 0.2 | 257 | 258.6 | 0.3 |
| FF1000 | 150 | 1260 | 1260 | 120.6 | 1260 | 1260.0 | 1080.9 |
| FF2000 | 200 | 4545 | 4545 | 486.4 | 4546 | 4552.5 | 1458.1 |
| WS250 | 70 | 3083 | 3083 | 623.1 | 4203° | 5447.9 | 0.5 |
| WS500 | 125 | 2072 | 2072 | 105.8 | 2085 | 2193.0 | 153.7 |
| WS1000 | 200 | 109677 | 109935 | 1256.2 | 154899 | 169877.2 | 19.3 |
| WS1500 | 265 | 13098 | 13098 | 658.2 | 13664 | 27810.6 | 108.0 |
| Bovine | 3 | 268 | 268 | 0.0 | 268 | 268.0 | 0.0 |
| Circuit | 25 | 2099 | 2099 | 0.1 | 2101 | 2188.7 | 0.1 |
| Ecoli | 15 | 806 | 806 | 0.0 | 806 | 808.8 | 0.1 |
| USAir97 | 33 | 4336 | 4336 | 856.4 | 4336 | 5149.5 | 0.1 |
| HumanDi | 52 | 1115 | 1115 | 0.5 | 1115 | 1115.0 | 319.9 |
| TreniR | 26 | 918 | 918 | 0.2 | 918 | 918.0 | 10.9 |
| EU_fli | 119 | 348268 | 348325 | 431.6 | 350762 | 357502.7 | 89.5 |
| openfli | 186 | 26783 | 26796 | 3165.8 | 29481 | 31377.7 | 86.4 |
| yeast1 | 202 | 1412 | 1412 | 9.8 | 1412 | 1414.3 | 1085.8 |
| H1000 | 100 | 306349 | 307113 | 1068.4 | 330493 | 337217.0 | 211.4 |
| H2000 | 200 | 1242739 | 1245637 | 1527.2 | 1324988 | 1344914.2 | 126.4 |
| H3000a | 300 | 2840690 | 2842695 | 1204.3 | 2962661 | 3042695.4 | 165.2 |
| H3000b | 300 | 2837584 | 2840867 | 2040.0 | 2968601 | 3035272.7 | 438.0 |
| H3000c | 300 | 2835369 | 2831643 | 2159.3 | 2956916 | 3008793.3 | 903.6 |
| H3000d | 300 | 2828492 | 2830284 | 3895.2° | 2967747 | 3030236.3 | 51.6 |
| H3000e | 300 | 2843000 | 2846536 | 4035.6° | 3012046 | 3043889.5 | 25.9 |
| H4000 | 400 | 5038611 | 5096437 | 1563.7 | 5261850 | 5383301.3 | 347.5 |
| H5000 | 500 | 7964765 | 8007638 | 2024.1 | 8206499 | 8348589.5 | 696.5 |
| powergr | 494 | 15862 | 15906 | 856.7 | 15965 | 16084.9 | 3393.5 |
| Oclinks | 190 | 611253 | 615467 | 267.8 | 622237 | 626701.5 | 836.1 |
| faceboo | 404 | 420334 | 589763 | 3526.9 | 794938 | 857245.1 | 685.4 |
| grqc | 524 | 13591 | 13743 | 1025.0 | 13666 | 13701.7 | 3391.1 |
| hepth | 988 | 106276 | 115309 | 4521.4° | 109504 | 111495.6 | 1602.3 |
| hepph | 1201 | 6155877 | 7556094 | 4025.1° | 8464199 | 10530358.9 | 3513.9 |
| astroph | 1877 | 53963375 | 57895042 | 3105.2 | 59476077 | 60449655.3 | 3512.9 |
| condmat | 2313 | 2298596 | 7658643 | 4203.5° | 3756761 | 4632078.6 | 3555.9 |

 $^{^{\}circ}$ indicates a longer computation time than 3600 seconds.

C. TarjanInComponent Procedure

Algorithm 1 realizes the TarjanInComponent procedure used in the articulation point impact strategy. It starts the search from a root node, and recursively builds a depth first search (DFS) tree. During DFS phase, all articulation points are identified. Once all nodes are visited, the evaluation of all nodes is finished.

Algorithm 1 Pseudo Code of TarjanInComponent Procedure

```
Input: A large connected component C, root node x, time stamp Count, \gamma and \eta
1: dfn[x] \leftarrow + + Count
2: low[x] \leftarrow dfn[x]
```

```
3: for each neighbour w of x in C do
       if w has not been visited then
          TarjanInComponent(C, w, Count, \gamma, \eta)
 5:
          low[x] \leftarrow min\{low[x], low[w]\}
 6:
          //w is not the parent of v
          if dfn[x] < dfn[w] then
 7:
             \gamma[x] + = \gamma[w]
 8:
          end if
9:
10:
          if dfn[x] < low[w] then
              The number of x's subtrees increases one
11:
              if x is not the root node then
12:
                 x is marked as an articulation point
13:
                 \eta[x] + = \gamma[w]
14:
                 \psi[x] + = \frac{\gamma[w](\gamma[w] - 1)}{2}
15:
16:
17:
                 if x is the root \wedge x has more than one subtree then
18:
                    x is marked as an articulation point
                 end if
19:
              end if
20:
          end if
21:
22:
       else
          low[x] \leftarrow min\{low[x], dfn[w]\}
23:
24:
       end if
```

25: end for