

Asymptotics

The asymptotics of Guava's utilities are entirely predictable, but are listed here for completeness.

List

Implementation	<code>add</code>	<code>add(i, elem)</code>	<code>remove(i)</code>	<code>contains</code>	Iteration	size
<code>ArrayList</code> (JDK)	$O(1)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(1)$
<code>LinkedList</code> (JDK)	$O(1)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(1)$
<code>CopyOnWriteArrayList</code> (JDK)	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(1)$
<code>ImmutableList</code>	N/A	N/A	N/A	$O(n)$	$O(n)$	$O(1)$
<code>ImmutableSet.asList()</code>	N/A	N/A	N/A	$O(1)$	$O(n)$	$O(1)$

Set

Implementation	<code>add</code>	<code>remove</code>	<code>contains</code>	Iteration	size
<code>HashSet</code> (JDK)	$O(1)$	$O(1)$	$O(1)$	$O(\max n) *$	$O(1)$
<code>LinkedHashSet</code> (JDK)	$O(1)$	$O(1)$	$O(1)$	$O(n)$	$O(1)$
<code>TreeSet</code> (JDK)	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(n)$	$O(1) **$
<code>CopyOnWriteArraySet</code> (JDK)	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(1)$
<code>ImmutableSet</code>	N/A	N/A	$O(1)$	$O(n)$	$O(1)$
<code>ImmutableSortedSet</code>	N/A	N/A	$O(\log n)$	$O(n)$	$O(1)$

* `HashSet` iteration takes time proportional to the maximum number of elements the `HashSet` has ever had, not proportional to the current number of elements.

** `TreeSet.subSet(...).size()` takes time proportional to the size of the subset.

Multiset

Note: n is the number of **distinct** elements in the multiset.

Implementation	Performs like a...	<code>size()</code>	<code>count(E)</code>	<code>add(E, int)</code>	<code>remove(E, int)</code>	<code>setCount(int)</code>
<code>HashMultiset</code>	<code>HashMap<E, Integer></code>	$O(1)$	$O(1)$	$O(1)$	$O(1)$	$O(1)$
<code>LinkedHashMultiset</code>	<code>LinkedHashMap<E, Integer></code>	$O(1)$	$O(1)$	$O(1)$	$O(1)$	$O(1)$

TreeMultiset	TreeMap<E, Integer>	$O(1)$ **	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(\log n)$
ConcurrentHashMultiset	ConcurrentHashMap<E, AtomicInteger>	$O(n)$	$O(1)$	$O(1)$	$O(1)$	$O(1)$
ImmutableMultiset	ImmutableMap<E, Integer>	$O(1)$	$O(1)$	$O(1)$	$O(1)$	$O(1)$
ImmutableSortedMultiset	ImmutableSortedMap<E, Integer>	$O(1)$	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(\log n)$

* Like `HashMap`, the iteration cost through the `entrySet` is linear in the maximum number of elements the `HashMultiset` has ever had, not the number it has now.

** `TreeMultiset.subMultiset().size()` takes time $O(\log n)$.

Multimap

k is the number of distinct keys; n is the number of distinct entries; $\#(key)$ is the number of entries associated with `key`. Where not specified, the asymptotics are equivalent to the "obvious" implementation based on the "Performs like a..." column.

Implementation	Performs like a...	<code>size()</code>	<code>get(K)</code>	<code>put(K, v)</code>	<code>containsEntry(K, v)</code>	Iteration through entries
ArrayListMultimap	HashMap<K, ArrayList<V>>	$O(1)$	$O(1)$	$O(1)$	$O(\#(key))$	$O(max + n)$
LinkedListMultimap	LinkedHashMap<K, LinkedList<V>>	$O(1)$	$O(1)$	$O(1)$	$O(\#(key))$	$O(n)$
HashMultimap	HashMap<K, HashSet<V>>	$O(1)$	$O(1)$	$O(1)$	$O(1)$	$O(max + n)$
LinkedHashMultimap	LinkedHashMap<K, LinkedHashSet<V>>	$O(1)$	$O(1)$	$O(1)$	$O(1)$	$O(n)$
TreeMultimap	TreeMap<K, TreeSet<V>>	$O(1)$	$O(\log k)$	$O(\log k + \log \#(key))$	$O(\log k + \log \#(key))$	$O(n)$
ImmutableListMultimap	ImmutableMap<K, ImmutableList<V>>	$O(1)$	$O(1)$	N/A	$O(\#(key))$	$O(n)$
ImmutableSetMultimap	ImmutableMap<K, ImmutableSet<V>>	$O(1)$	$O(1)$	N/A	$O(1)$	$O(n)$