## Asymptotics

The asymptotics of Guava's utilities are entirely predictable, but are listed here for completeness.

List

Implementation	add	add(i, elem)	remove(	(i)contai	ns Iteratic	n size
ArrayList (JDK)	O(1)	O(n)	O(n)	O(n)	O(n)	O(1)
${ t LinkedList} \ ({ t JDK})$	O(1)	O(n)	O(n)	O(n)	O(n)	O(1)
$ \begin{array}{c} {\tt CopyOnWriteArrayLi} \\ ({\tt JDK}) \end{array} $	stO(n)	O(n)	O(n)	O(n)	O(n)	O(1)
ImmutableList	N/A	N/A	N/A	O(n)	O(n)	O(1)
ImmutableSet.asLis	t(N/A	N/A	N/A	O(1)	O(n)	O(1)

## Set

Implementation	add	remove	contains	Iteration	size
HashSet (JDK) LinkedHashSet (JDK)	O(1) O(1)	O(1) O(1)	O(1) O(1)	$O(max \ n) *$ $O(n)$	O(1) $O(1)$
TreeSet (JDK) CopyOnWriteArraySet	$O(\log n)$ $O(n)$	$O(\log n)$ $O(n)$	$O(\log n)$ $O(n)$	O(n) $O(n)$	O(1) ** O(1)
(JDK) ImmutableSet ImmutableSortedSet	N/A N/A	N/A N/A	$O(1)$ $O(\log n)$	$O(n) \ O(n)$	O(1) O(1)

<sup>\*</sup> HashSet iteration takes time proportional to the maximum number of elements the HashSet has ever had, not proportional to the current number of elements.

## Multiset

Note: n is the number of **distinct** elements in the multiset.

<sup>\*\*</sup> TreeSet.subSet(...).size() takes time proportional to the size of the subset.

D 6			(-		<i>'</i>	Iterate through
Performs			add(E	,remove	( <b>£</b> etCou	ntenftrySet() or
Implementaltikma	size()	count	(Et)	int)	int)	<pre>elementSet()</pre>
HashMultisHæntshMap <e,< td=""><td>O(1) (</td><td>O(1)</td><td>O(1)</td><td>O(1)</td><td>O(1)</td><td><math>O(max \ n) *</math></td></e,<>	O(1) (	O(1)	O(1)	O(1)	O(1)	$O(max \ n) *$
Integer>						
LinkedHashMunkteidsHatshMa	p <b>©</b> Ę[]) (	O(1)	O(1)	O(1)	O(1)	O(n)
Integer>						
TreeMultiscreeMap <e,< td=""><td>O(1)</td><td>O(log</td><td>O(log</td><td>O(log</td><td>O(log</td><td>O(n)</td></e,<>	O(1)	O(log	O(log	O(log	O(log	O(n)
Integer>	** T	$_{l})$	n)	n)	n)	
Concurrent Chbansh Murletnits Hen	s6Map <e< td=""><td><math>\dot{\Omega}(1)</math></td><td>O(1)</td><td>O(1)</td><td>O(1)</td><td>O(n)</td></e<>	$\dot{\Omega}(1)$	O(1)	O(1)	O(1)	O(n)
AtomicIntege	r>					
ImmutableMunhtuitsabdeMap	< <b>©</b> (1) (	O(1)	O(1)	O(1)	O(1)	O(n)
Integer>						
Immutable Sommutueto and built Sioco	±1 <b>⊝</b> ¢Màp∢	<b>E</b> (log	O(log	O(log	O(log	O(n)
Integer>	r	n)	n)	n)	n)	

<sup>\*</sup> Like HashMap, the iteration cost through the entrySet is linear in the maximum number of elements the HashMultiset has ever had, not the number it has now.

## Multimap

 ${\tt k}$  is the number of distinct keys;  ${\tt n}$  is the number of distinct entries;  ${\tt \#(key)}$  is the number of entries associated with key. Where not specified, the asymptotics are equivalent to the "obvious" implementation based on the "Performs like a. . ." column.

Performs Implementlakteon	size@et(]	put(K,	contains V)	, ,	<pre>Iterate through asMap().entrySet()</pre>
ArrayLis <b>tHasht</b> Ma <b>rp</b> eopK,	$O(1) \ O(1)$	O(1)	O(#(key)	)O(max	$O(\max k)$
ArrayList <v< td=""><td>&gt;&gt;</td><td></td><td></td><td>k + n</td><td></td></v<>	>>			k + n	
LinkedListiMuketdHaasphM	ap(K), O(1)	O(1)	O(#(key))	O(n)	O(k)
LinkedList<	V>>				
HashMultHaasphMap <k,< td=""><td><math>O(1) \ O(1)</math></td><td>O(1)</td><td>O(1)</td><td>O(max</td><td><math>O(max \ k)</math></td></k,<>	$O(1) \ O(1)$	O(1)	O(1)	O(max	$O(max \ k)$
HashSet <v>&gt;</v>				n)	
LinkedHashiMuketdHaasphM	ap(K), O(1)	O(1)	O(1)	O(n)	O(k)
LinkedHashS	et <v>&gt;</v>				
TreeMultimæpeMap <k,< td=""><td><math>O(1) O(\log</math></td><td><math>O(\log k)</math></td><td><math>O(\log k)</math></td><td>O(n)</td><td>O(k)</td></k,<>	$O(1) O(\log$	$O(\log k)$	$O(\log k)$	O(n)	O(k)
TreeSet <v>&gt;</v>	k)	+ log	+ log		
	ŕ	#(key))	#(key))		
ImmutableTminsuttMublteMax	app O(1)	N/A	O(#(key))	O(n)	O(k)
ImmutableLi	st <v>&gt;</v>	•	. , .,		

<sup>\*\*</sup> TreeMultiset.subMultiset().size() takes time  $O(\log n)$ .

Performs put(K Implementiation size@et(K))	, contair V)	Iterate nsE <b>htnoy</b> l@h, entries	Iterate through () asMap().en	ntrySet()
ImmutableShentiMaabdieMaqpe@K(1) O(1) N/A ImmutableSet <v>&gt;</v>	O(1)	O(n)	O(k)	