

# RCU and lockdep checking

All flavors of RCU have lockdep checking available, so that lockdep is aware of when each task enters and leaves any flavor of RCU read-side critical section. Each flavor of RCU is tracked separately (but note that this is not the case in 2.6.32 and earlier). This allows lockdep's tracking to include RCU state, which can sometimes help when debugging deadlocks and the like.

In addition, RCU provides the following primitives that check lockdep's state:

```
rcu_read_lock_held() for normal RCU.  
rcu_read_lock_bh_held() for RCU-bh.  
rcu_read_lock_sched_held() for RCU-sched.  
srcu_read_lock_held() for SRCU.
```

These functions are conservative, and will therefore return 1 if they aren't certain (for example, if `CONFIG_DEBUG_LOCK_ALLOC` is not set). This prevents things like `WARN_ON(!rcu_read_lock_held())` from giving false positives when lockdep is disabled.

In addition, a separate kernel config parameter `CONFIG_PROVE_RCU` enables checking of `rcu_dereference()` primitives:

```
rcu_dereference(p):  
    Check for RCU read-side critical section.  
rcu_dereference_bh(p):  
    Check for RCU-bh read-side critical section.  
rcu_dereference_sched(p):  
    Check for RCU-sched read-side critical section.  
srcu_dereference(p, sp):  
    Check for SRCU read-side critical section.  
rcu_dereference_check(p, c):  
    Use explicit check expression "c" along with rcu_read_lock_held(). This is useful in code that is invoked by both  
    RCU readers and updaters.  
rcu_dereference_bh_check(p, c):  
    Use explicit check expression "c" along with rcu_read_lock_bh_held(). This is useful in code that is invoked by  
    both RCU-bh readers and updaters.  
rcu_dereference_sched_check(p, c):  
    Use explicit check expression "c" along with rcu_read_lock_sched_held(). This is useful in code that is invoked  
    by both RCU-sched readers and updaters.  
srcu_dereference_check(p, c):  
    Use explicit check expression "c" along with srcu_read_lock_held(). This is useful in code that is invoked by both  
    SRCU readers and updaters.  
rcu_dereference_raw(p):  
    Don't check. (Use sparingly, if at all.)  
rcu_dereference_protected(p, c):  
    Use explicit check expression "c", and omit all barriers and compiler constraints. This is useful when the data  
    structure cannot change, for example, in code that is invoked only by updaters.  
rcu_access_pointer(p):  
    Return the value of the pointer and omit all barriers, but retain the compiler constraints that prevent duplicating or  
    coalescing. This is useful when when testing the value of the pointer itself, for example, against NULL.
```

The `rcu_dereference_check()` check expression can be any boolean expression, but would normally include a lockdep expression. However, any boolean expression can be used. For a moderately ornate example, consider the following:

```
file = rcu_dereference_check(fdt->fd[fd],  
                             lockdep_is_held(&files->file_lock) ||  
                             atomic_read(&files->count) == 1);
```

This expression picks up the pointer "`fdt->fd[fd]`" in an RCU-safe manner, and, if `CONFIG_PROVE_RCU` is configured, verifies that this expression is used in:

1. An RCU read-side critical section (implicit), or
2. with `files->file_lock` held, or
3. on an unshared `files_struct`.

In case (1), the pointer is picked up in an RCU-safe manner for vanilla RCU read-side critical sections, in case (2) the `->file_lock` prevents any change from taking place, and finally, in case (3) the current task is the only task accessing the `file_struct`, again preventing any change from taking place. If the above statement was invoked only from updater code, it could instead be written as follows:

```
file = rcu_dereference_protected(fdt->fd[fd],  
                                lockdep_is_held(&files->file_lock) ||  
                                atomic_read(&files->count) == 1);
```

This would verify cases #2 and #3 above, and furthermore lockdep would complain if this was used in an RCU read-side critical section unless one of these two cases held. Because `rcu_dereference_protected()` omits all barriers and compiler constraints, it generates better code than do the other flavors of `rcu_dereference()`. On the other hand, it is illegal to use `rcu_dereference_protected()` if either the RCU-protected pointer or the RCU-protected data that it points to can change concurrently.

Like `rcu_dereference()`, when lockdep is enabled, RCU list and hlist traversal primitives check for being called from within an RCU read-side critical section. However, a lockdep expression can be passed to them as a additional optional argument. With this lockdep expression, these traversal primitives will complain only if the lockdep expression is false and they are called from outside any RCU read-side critical section.

For example, the workqueue `for_each_pwq()` macro is intended to be used either within an RCU read-side critical section or with `wq->mutex` held. It is thus implemented as follows:

```
#define for_each_pwq(pwq, wq)
    list_for_each_entry_rcu((pwq), &(wq)->pwqs, pwqs_node,
        lock_is_held(&(wq->mutex).dep_map))
```