



Theorembeweiserpraktikum

Introduction & Propositional Logic

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THEOREM PROVER

Outline

- What is interactive theorem proving?
- Terms and types
- The Curry-Howard correspondence
- The structure of a Lean file
- Some basic types in Lean
- Propositional logic in Lean

What is Interactive Theorem Proving (ITP)?

- Language and IDE supporting the user in writing *formal proofs* about certain structures which are then *verified* by the computer.
- ITP is *not* automated theorem proving (ATP)! User input is usually required to guide the system.
- Formal foundation of Lean: Type theory, inductive types (Calculus of Inductive Constructions, CIC).

Interactive Theorem Proving – How is it Useful?

- Verify proofs independently of an academic review system.
- Write programs and verify them – in the same language!

A Cherry-Picked List of Theorem Proving Projects

2005 *Gonthier & Werner*: four color theorem [Coq]

2014 *Hales et al*: Kepler conjecture [HOL Light/Isabelle]

2014+ *Leroy et al*: CompCert, a verified C compiler [Coq]

2020 *Han & van Doorn*: independence of the continuum hypothesis [Lean 3]

On Lean Versions

- This course is about *Lean 4*, a rewritten (in Lean!) and greatly extended implementation of Lean
- Lean 4 is **not** backwards-compatible with Lean 3, the latest stable release
- Most information you can find online will be about Lean 3
- Lean's sizable *mathlib* is not available for Lean 4 yet, but we won't need it
- If you find bugs or room for improvement, do tell us!

Terms and Types

- Martin-Löf type theory (MLTT) is a very strong, static type system.
- It is based on *type judgments*:

$$t : \alpha$$

means “ t is an instance of α ” for a term t and a type α .

- Everything has a fixed type, even types themselves: Universes are types that contain types:

$$\begin{array}{lcl}
 t & : & \alpha \\
 \alpha & : & \text{Sort}_u \\
 \text{Sort}_u & : & \text{Sort}_{u+1} \\
 \text{Sort}_{u+1} & : & \dots
 \end{array}$$

The Curry-Howard Correspondence



- Named after Haskell Curry and William Alvin Howard, 1969
- Paraphrasing:
 - “**Proving** in intuitionistic (constructive) logic
is the same as
programming in typed lambda calculus.”
- We prove statements by constructing instances of certain types.

The Curry-Howard Correspondence

	Proving	Programming
$\alpha : \text{Sort}_u$	α is a proposition	α is a data type
$t : \alpha$	t is a proof of α	t is an element of α
$\alpha \rightarrow \beta$	α implies β	type of functions from α to β
$\alpha \times \beta$	α and β	cartesian product of α and β
$\alpha + \beta$	α or β	disjoint sum of α and β

The Difference Between Proof and Data

- Some provers do not distinguish between proof and data
- The biggest difference in Lean is:

We can treat instances $h : p$ and $h' : p$ for a proposition p as equal,

but

we need to distinguish instances $x : \alpha$ and $y : \alpha$ of a data type.

- Lean keeps propositions and data types in different universes, has aliases

$$\begin{aligned}\text{Prop} &::= \text{Sort}_0 \\ \text{Type} &::= \text{Type}_0 ::= \text{Sort}_1 \\ \text{Type}_1 &::= \text{Sort}_2 \\ &\dots \quad \dots\end{aligned}$$

- Definitions of data are (conventionally) preceded by the keyword **def**, propositional definitions by **theorem**

The Structure of a Lean File

```
-- single-line comments start with "--"
import MyProject.Utils -- import file `MyProject/Utils.lean` from current project/dependency

def add_three_mul (x y : Nat) : Nat := -- define a function with two arguments x and y
  x + 3 * y

theorem add_three_mul_gt {x y : Nat} : add_three_mul x y >= x :=
  ...

example {y : Nat} : 5 + 3 * y >= 5 := -- an example is a theorem without a name
  add_three_mul_gt

#check add_three_mul_gt -- check the type of our theorem
```

Propositional Logic in Lean: True and False

The canonical way to prove the proposition `True` is called `True.intro` :

```
example : True := True.intro
```

There is no way of proving false (hopefully!) but we can use the principle of “ex falso quodlibet” by using `False.elim` :

```
example (p : Prop) (fa : False) : p := False.elim fa
```

Propositional Logic in Lean: Implication

Implication is modelled by function types. So, applying an implication is function application:

```
section
variable (p q r : Prop) -- automatically become arguments when referenced in the section

theorem modus_ponens (hpq : p → q) (hp : p) : q :=
  hpq hp -- function application is written with a space instead of parentheses!

example (hpqr : p → q → r) (hp : p) (hq : q) : r := -- → has implicit parentheses on the right
  hpqr hp hq -- function application has implicit parentheses on the left

example (htp : True → p) : p :=
  htp True.intro
end
```

We can input \rightarrow writing `\to`.

Convention: We write *curried* functions: $\alpha \rightarrow \beta \rightarrow \gamma$ instead of $\alpha \times \beta \rightarrow \gamma$ or $\alpha \wedge \beta \rightarrow \gamma$!

Propositional Logic in Lean: Implication

Proving an implication is done by lambda abstraction:

```
section
variable (p q r : Prop)

example : p → p → p :=
  fun hp hp' => hp

example (hq : q) : p → q :=
  fun hp => hq

example (hpq : p → q) (hqr : q → r) : p → r :=
  fun hp => hqr (hpq a)
end
```

The expression `fun a => b` is what Mathematicians write as $a \mapsto b$!

Arguments vs. Implications

The following two Lean snippets are the same:

```
example (p q r : Prop) (hpq : p → q) (hqr : q → r) : p → r :=  
  fun hp => hqr (hpq hp)
```

and

```
example (p q r : Prop) : (p → q) → (q → r) → p → r :=  
  fun hpq hqr hp => hqr (hpq hp)
```

Arguments of a definition are just shorthand for iterated implications (functions)!

Propositional Logic in Lean: Disjunction

- The disjunction of two propositions is p and q is written $p \vee q$.
- A case distinction on $p \vee q$ can be done by *matching* on the proof.

```
example (p q : Prop) (hp : p) : p ∨ q := Or.inl hp

example (p q r s : Prop) (hp : p) (hs : s) : (p ∨ q) ∧ (r ∨ s) :=
  And.intro (Or.inl hp) (Or.inr hs)

example (p q r : Prop) (hpq : p ∨ q) (hpr : p → r) (hqr : q → r) : r :=
  match hpq with
  | Or.inl hp => hpr hp
  | Or.inr hq => hqr hq
```

Write \vee as `\or`.

Propositional Logic in Lean: Conjunction

- The conjunction of two propositions p and q is written $p \wedge q$.
- A proof of $p \wedge q$ is created using `And.intro`.
- A proof of p can be recovered using matching:

```
example (p q : Prop) (hpq : p ∧ q) : p :=  
  match hpq with  
  | And.intro hp hq => hp
```

```
example (p q : Prop) (hp : p) (hq : q) : p ∧ q := And.intro hp hq
```

We can input \wedge using `\and`.

Algebraic Data Types and Propositions

Where does the strange `match` syntax come from? `Or` is defined as an algebraic data type:

```
inductive Or (p q : Prop) : Prop where
| inl (hp : p) : Or p q
| inr (hq : q) : Or p q
```

Each line starting with `|` is a *constructor* to `Or p q`. It states one way to construct an instance of the proposition.

`And` can be defined by the same mechanism, but only has *one* constructor with *multiple* parameters:

```
inductive And (p q : Prop) : Prop where
| intro (hp : p) (hq : q) : And p q
```

Algebraic Data Types and Propositions

`True` and `False` are algebraic data types as well, with one and zero constructors, respectively!

```
inductive True : Prop where
| intro : True

inductive False : Prop where
```

The term `nomatch` lets us use the fact that `False` does not have any constructors, and `False.elim` is defined by using `nomatch`.

Algebraic Data Types and Propositions

And and Or are given more convenient *infix notations* as we've seen.

```
infixr:35 " ^ " => And
infixr:30 " v " => Or
```

Both notations associate to the right, with And binding more tightly, i.e.

$$a \wedge b \wedge c \vee d \equiv (a \wedge (b \wedge c)) \vee d \equiv \text{Or } (\text{And } a \ (\text{And } b \ c)) \ d$$

Algebraic Data Types and Propositions

Data types can be defined the same way, like this type of days of the week:

```
inductive Weekday : Type where
| monday : Weekday
| tuesday : Weekday
| wednesday : Weekday
| thursday : Weekday
| friday : Weekday
| saturday : Weekday
| sunday : Weekday
```

This definition then enables us to *match* on variables of this data type:

```
def dayNumber (w : Weekday) : Nat :=
  match w with
  | Weekday.monday => 0
  | Weekday.tuesday => 1
  | Weekday.wednesday => 2
  | Weekday.thursday => 3
  | Weekday.friday => 4
  | Weekday.saturday => 5
  | Weekday.sunday => 6
```

Lean will throw an error whenever we forget one of the constructors!

Placeholders

Often, you will not write a proof in one go, but instead compose the proof term bit by bit. In this situation we can use the underscore `_` as a placeholder. Lean will try to infer its type and display it when hovering over the underscore as well as in the *Infoview*:

```
example (f :  $\alpha \rightarrow \beta$ ) (a :  $\alpha$ ) :  $\beta$  :=  
  f _ -- shows that we need to provide a term of type  $\alpha$ 
```

have Expressions

To create a local variable which abbreviates another expressions, we can use **have** :

```
example (f :  $\beta \rightarrow \beta \rightarrow \gamma$ ) (g :  $\alpha \rightarrow \beta$ ) (a :  $\alpha$ ) :  $\gamma$  :=  
  have b :  $\beta$  := g a  
  f b bha
```

For longer proofs, you will find that using **have** is a way to make them more readable and concise!

Demo

Summary

<i>type</i> p	<i>introduction</i> $\dots \vdash p$ “how do I prove this”	<i>elimination</i> $h : p, \dots \vdash \dots$ “how do I use this”
<i>implication/function type</i> $q \rightarrow r$	<i>abstraction</i> <code>fun hq => (_ : r)</code>	<i>application</i> $h \ hq$
<i>inductive type</i>	<i>constructor app.</i>	<i>matching</i>
$q \vee r$	<code>Or.inl hq , Or.inr hr</code>	<code>match h with</code> <code> Or.inl hq => _</code> <code> Or.inr hr => _</code>
$q \wedge r$	<code>And.intro hq hr</code>	<code>match h with</code> <code> And.intro hq hr => _</code>
<code>True</code>	<code>True.intro</code>	<code>—</code>
<code>False</code>	<code>—</code>	<code>False.elim h / nomatch h</code>