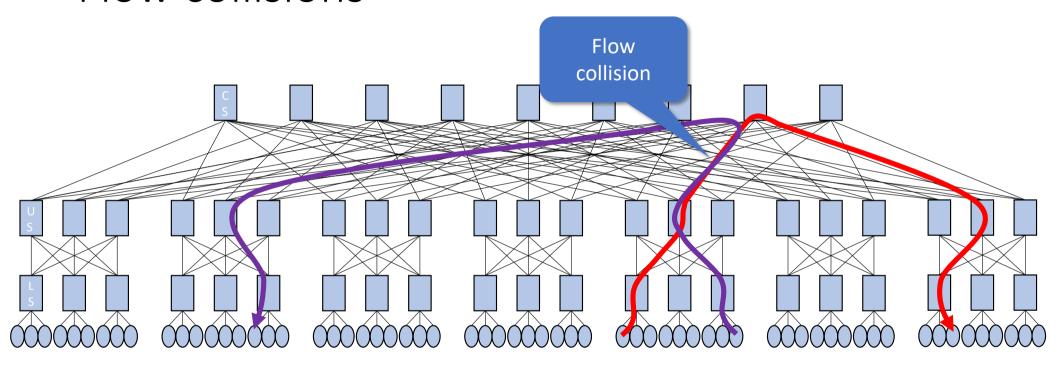
Load balancing strategies in AIML networks

Costin Raiciu

Broadcom and Politehnica of Bucharest

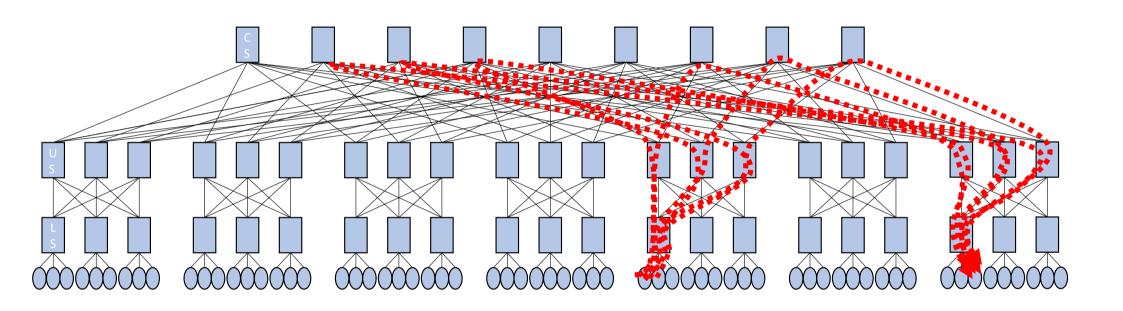
Flow collisions



Why not use MPTCP for AIML networks?

- Need to use many paths for the common case of
 - Symmetric highly loaded networks + short flows.
 - Best way to load balance short flows is to use many paths.
 - Load balancing works well for long flows, not so well for shorter flows.
- But with many paths, minimum MPTCP total window is #paths.
 - E.g. 256 paths means min 256 packet window. This equals BDP at 800Gbps.
 - Congestion collapse in incast.
- Path state for MPTCP is quite costly.
 - CWND, flight_size, sequence numbers, etc. (tens of bytes).
 - 256 * 20 = 5KB per connection!

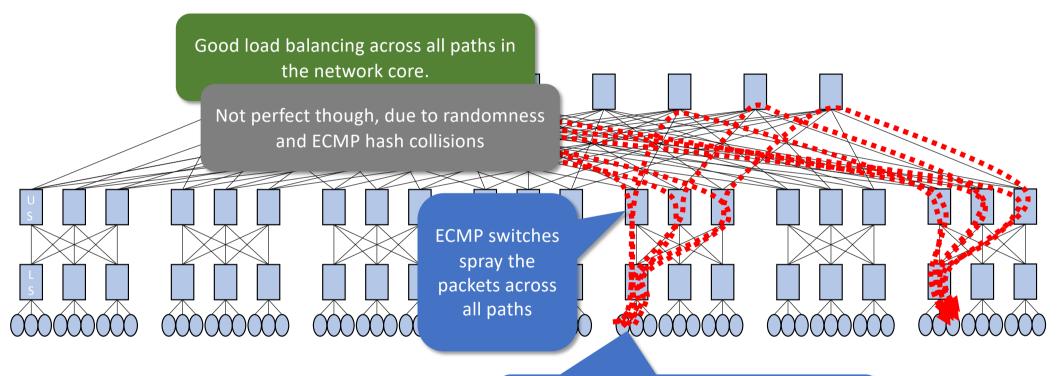
Packet spraying in AIML networks



Congestion control with packet spraying

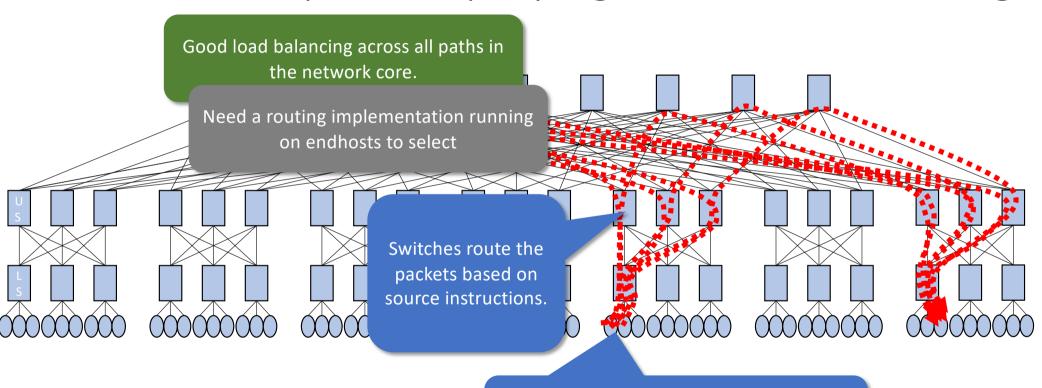
- Maintain a single congestion window that upper bounds flight size.
- Sender-driven congestion control (e.g. UET NSCC)
 - Targets sub-BDP standing queue at the bottleneck.
 - Use ECN and delay simultaneously.
 - Aggressive increase when queue ~ 0. Linear increase otherwise.
 - Multiplicative decrease when ECN mark & average delay above threshold.
 - Average delay across all paths.
- But how to load balance packets across paths?
 - Bad load balancing results in reducing CWND (across all paths).

End-to-end packet spraying via ECMP

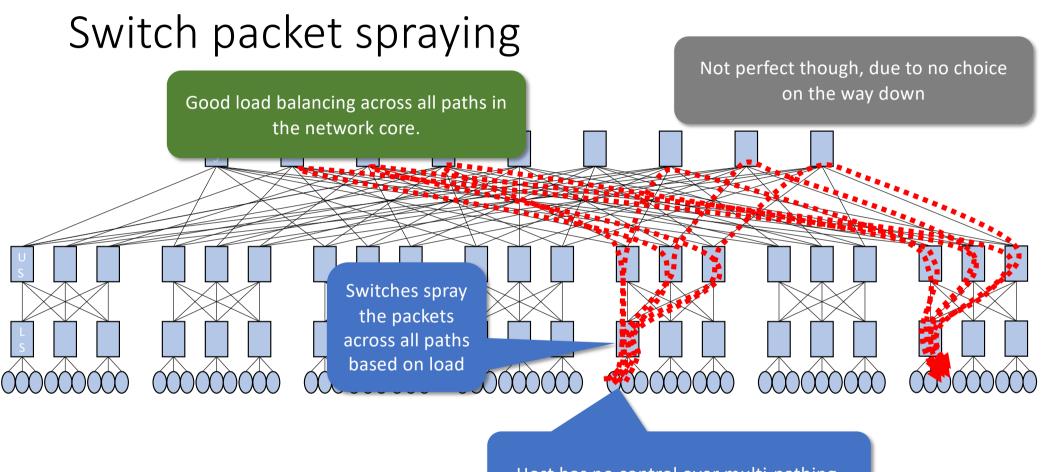


Change entropy value (e.g. UDP src port) with **every** packet sent

End-to-end packet spraying with source routing.



Change path for **every** packet sent



Two basic approaches for spraying

- Host-based spraying
 - ECMP + standard routing protocol (e.g. BGP).
 - Source routing requires an SDN routing protocol to compute paths and deliver them to hosts.
- Switch spraying adaptive routing / dynamic load balancing.

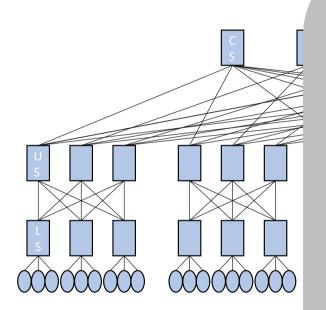
What are the pros and cons of each?

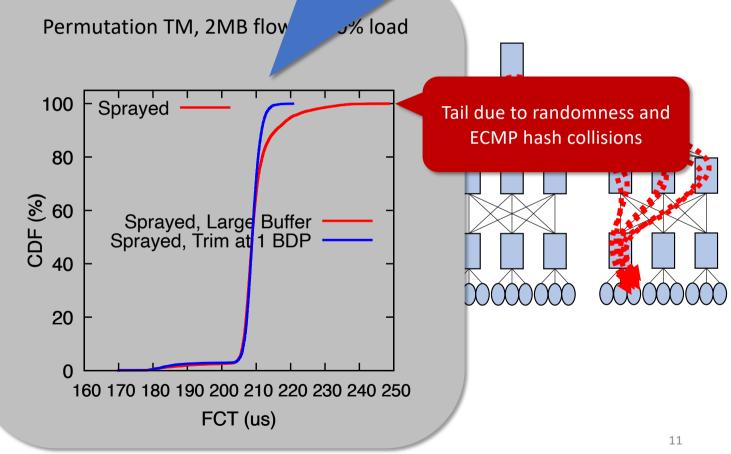
Host based spraying What is the best way to spray packets?

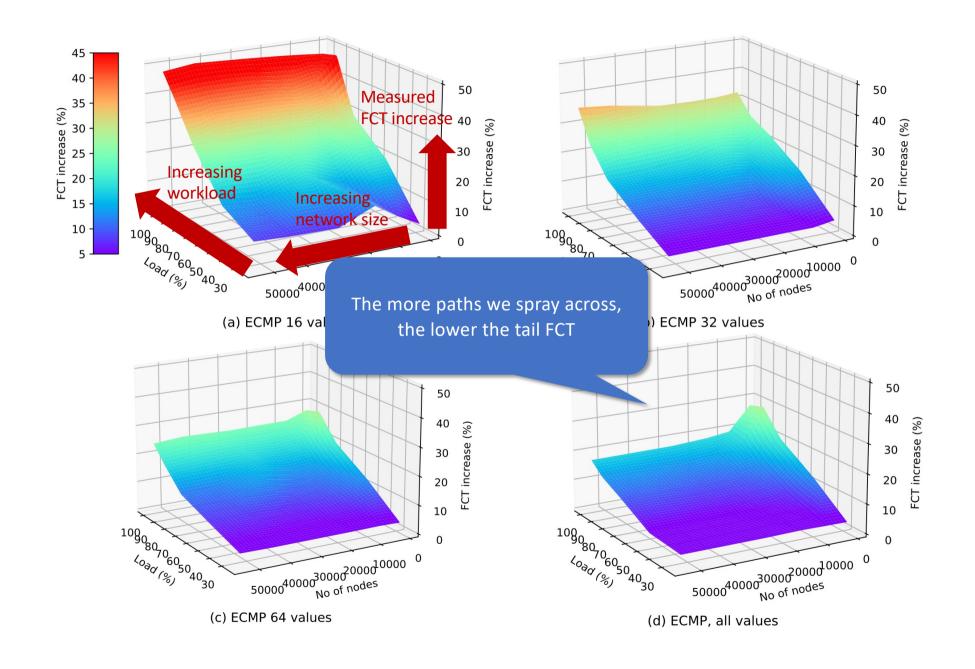
- Simplest: oblivious load balancing
 - Pick a random EV for each packet.
 - Works very well if network capacity is uniform.
- Bitmap load balancing (e.g. UET bitmap algorithm)
 - Per EV state one or a few bits.
 - When ACK indicates ECN mark, increment EV state.
 - When EV is next to be picked but non-zero state, decrement state, skip.
- Recycled entropies (REPS):
 - Keep EV cache for which we got an ACK without ECN set.
 - Path selection: pick EV from cache if non-empty. Otherwise pick random EV.

Oblivious packe

When sprayed load balancing is imperfect, queues can still build.
Trimming prevents queue building.
Packet gets trimmed, NACKed,
RTX on a different less loaded path.



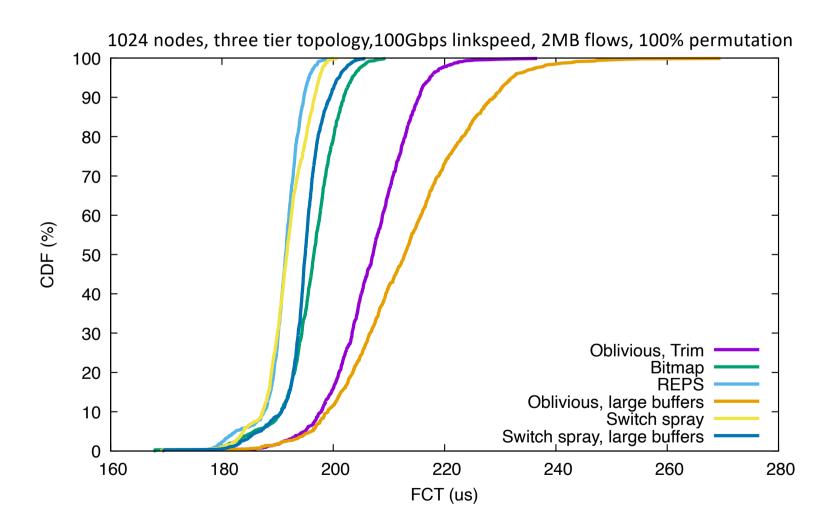




How does switch spray work?

- ECMP group => DLB group.
- During route lookup, routes in DLB group are consulted.
 - Contains all available paths towards destination.
 - Switch uses local information to decide which route (and associated egress port) to pick
 - Example metrics:
 - Queue length
 - Bandwidth utilization
 - PFC Port state.
 - Combination of the above possible.
- Works very well when path choice exists (e.g. going up the tree).
- Less well on the downward path / with asymmetries.
 - At the limit, behaves like oblivious endhost spraying.

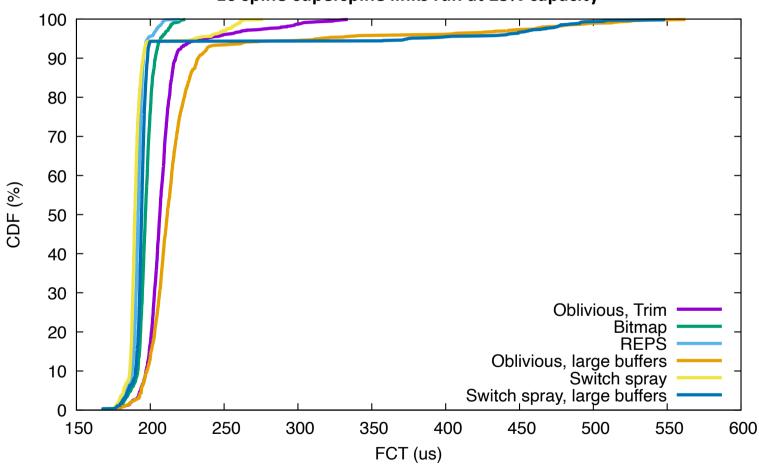
Load balancing algorithms comparison



Load balancing with asymmetric capacity

1024 nodes, three tier topology,100Gbps linkspeed, 2MB flows, 100% permutation

10 spine-superspine links run at 25% capacity



Summary

- Packet spraying enables exploring many/all paths.
- Endhost or switch spraying possible.
- All schemes work almost perfect when network is perfectly symmetric.
 - Switch spraying
 - Lower average buffer utilization than endhost load balancing.
 - Works very well when path choice available (e.g. link bundles).
 - Endhost spraying:
 - State-based schemes can achieve similar FCT to switch spraying.
 - But lead to higher queue utilization.
 - Even oblivious works quite well.
- When network is asymmetric, poor load balancing leads to large FCT.
 - Switch spray on its own struggle need additional mechanisms.
 - State-based endhost load balancing copes fairly well.