

PESU Center for Information Security, Forensics and Cyber Resilience



PESU Center for Information Security, Forensics and Cyber Resilience

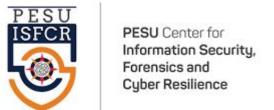
Cybsec Workshop

Introduction to reversing



Contents

- 1) What is Reverse Engineering
- 2)Introduction to decompilers and disassemblers
- 3) Basic x86_64 architecture
- 4) Analyzing and Reversing Simple Hello World Program
- 5) Solving A basic Reverse Engineering Challlenge
- 6) Ghidra and Other reverse engineering tools and References

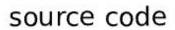


What is Reverse Engineering

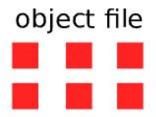
 Generally To understand what Is reversing engineering.one must first understand how any higher level compiled languages is converted into machine readable / byte code, Ever Wondered what happens on the inside, how The computer reads and gives output magically from source code

 The Compiler which is a special program converts Source Code into machine language module called object file. And Another specialized program called the Linker combines this object file with other previously compiled object files (in particular run-time modules) to create an executable file





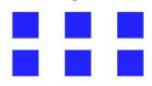
compiler





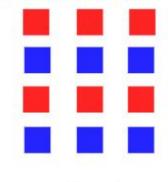
PESU Center for Information Security, Forensics and Cyber Resilience

library files

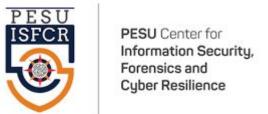


linker

executable file



output



 To see the compiler and linker in action you can use "-c" option in gcc to compile a source without linking

```
PES2UG19CS350)-[~/rev]
cc -c helloworld.c
      PES2UG19CS350)-[~/rev]
oworld.c helloworld.o peda-session-helo.txt
   PES2UG19CS350)-[~/rev]
hmod +x ./helloworld.o
   pessug19cs350)-[~/rev]
/helloworld.o
exec format error: ./helloworld.o
   PES2UG19CS350)-[~/rev]
ile helloworld.o
oworld.o: ELF 64-bit LSB relocatable, x86-64, ve
```



 However you Cant run the executable file Yet because The Linker has not yet linked the Run Time Modules To the Object file (hence the exec Format error)

You Can Link the object file with the Linker Using the Standard gcc command

```
PES2UG19CS350)-[~/rev]
    gcc helloworld.c -o hello
          PES2UG19CS350)-[~/rev]
zsh: no such file or directory: ./hel
        PES2UG19CS350)-[~/rev]
Hello world LMAO
```



- Since Now you Have a Basic understanding Of how pseudo-readable code is converted into machine language by the compiler and Interpreter
- Reverse Engineering is typically the process of taking a compiled (machine code, bytecode) program and trying to understand what the program does and converting it back into a more human readable format.



Decompilers and Disassemblers

 A Diassembler is a Program which converts Machine / byte code into intermediate assembly instructions

```
ump of assembler code for function main:
 0x0000000000040052d <+0>:
                                     %гьр
                                     %rsp,%rbp
                                     $0x10,%rsp
                              movq $0x4005e4,-0x10(%rbp)
 0x000000000040053d <+16>:
                             movq $0x0,-0x8(%rbp)
                                     -0x10(%rbp),%rax
                                    -0x10(%rbp),%rcx
                                     $0x0,%edx
                                     %rcx,%rsi
                                     %rax.%rdi
                              callq 0x400420 <execve@plt>
 0x00000000000400558 <+43>:
                              leaveg
 0x0000000000040055e <+49>:
nd of assembler dump
```

 A Decompiler is program which "converts back" machine code into human readable format and pseudo-code



X86_64 architecture

 x86 Assembly is the assembly instruction code used by the non ARM (Intel/AMD) processors

- There are 8 general purpose registers in both 32bit and 64 bit types
- And a special instruction pointer

- You can think of registers as special variables of the assembly world
- Which can be used to hold any type of data to which some have acquired specific use which are used in programs



- 32 bit arch has 6 general purpose registers and 2 special registers
- The 8 registered are named as EAX, EBX, ECX, EDX, ESI, EDI, ESP, EBP
- In 64 bit they are named as RAX, RBX, RCX, RDX, RSI, RDI, RSP, RBP
- RSP is the stack pointer always points to the top of the process stack
- You Can perform basic assembly operations on these Registers Such as
- Mov
- Cmp
- Test
- lea
- Push
- Pop

Analyzing and Reversing Simple progr

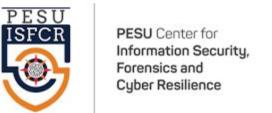
- Lets see a Simple compiled program Aka executable File which just prints text to output
- This is our source code for Reference

```
#include<stdio.h>
int main(){
printf("Hello world LMAO\n");
return 0; }
```

```
(root PES2UG19CS350)-[~/rev]

# ./hello
Hello world LMAO
```

 Now Lets use gdb(disassembler) to take a look inside and Try to understand the Assembly!



 Pretty Small disassembly as expected let us go step by step to understand Whats happening

```
disas main
Dump of assembler code for function main:
   0×0000000000001135 <+0>:
                                push
                                       rbp
   0×00000000000001136 <+1>:
                                       rbp,rsp
   0×0000000000001139 <+4>:
                                lea
                                       rdi,[rip+0×ec4]
                                                               # 0×2004
   0×0000000000001140 <+11>:
                                call
                                       0×1030 <puts@plt>
   0×0000000000001145 <+16>:
                                       eax,0×0
                                mov
   0×000000000000114a <+21>:
                                       rbp
                                pop
   0×000000000000114b <+22>:
                                ret
End of assembler dump.
```

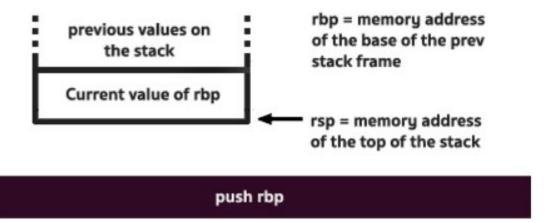
rbp



- The first Two instructions are Used for Setting up the Stack frame Of the process.so that it creates Space for the local variables and function parameters of the program
- Push Rbp: pushing the Current value present in RBP to the stack
- Because it "pushes" onto the stack, now the value of RSP contains memory address on top of new stack.



- The next instruction mov rbp,rsp copies the value of rsp into rbp
- And both point to top of the stack now.this is Done so that the process can subtract rsp by offset to make space for local variables
- while keeping rbp intact





• The third instruction Uses Lea (load effective address). In simple Terms it takes value present in that specific address and loads it into rdi general purpose register

lea

- The [] in assembly is used for dereferencing or in simple terms finds value present in that address .Similar to The * operator in c and c++
- so the third instruction.takes the value present in the address rip+0xec4 and puts it in the rdi register in our case it's the hello world string in memory



Lets see It live on gdb!

 As You can see the Register Rdi contains our string "hello World"

 Alternatively we can also check the RIP offset address just to make sure

```
RAX: 0×401122 (<main>: push rbp)

RBX: 0×0

RCX: 0×7fffff7fac718 → 0×7fffff7faeb00 → 0×0

RDX: 0×7fffffffe178 → 0×7fffffffe482 ("COLORFGBG=15;0")

RSI: 0×7fffffffe168 → 0×7fffffffe473 ("/root/rev/test")

RDI: 0×402004 ("Hello world LMAO")
```





PESU Center for Information Security, Forensics and Cyber Resilience

```
0×000000000001140 <+11>: call 0×1030 <puts@plt>
0×0000000000001145 <+16>: mov eax,0×0
0×00000000000114a <+21>: pop rbp
0×00000000000114b <+22>: ret
```

- In x64 assembly the function parameters are passed through the registers. In certain order. Visit https://docs.microsoft.com/en-us/cpp/build/x64-calling-convention?view=msvc-160 to know more about this
- The 4th instruction calls the function Puts. But wait A minute you must have a doubt! Even though we used printf It is shown as puts in the disassembly.why?
- The compiler replaces printf with puts if there are no format specifiers such as %d for optimization reasons
- And after 4th instruction ends it prints out string to output
- The rest two instructions just clear the rax register and pop rbp from the stack
- And Voila we have just analyzed our first disassembly of a simple program



Lets solve a simple crackme. Using Reverse Engineering

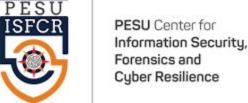
```
root⊕ PES2UG19CS350)-[~/rev]

# ./crackMe
input PassCode! :3123
3123
wrong passcode!
```

So The executable asks for a passcode and checks based according to that. But Since we Do not have the Source code we do not know the Passcode.But don't worry lets try to crack this using Reversing engineering

The disassembly

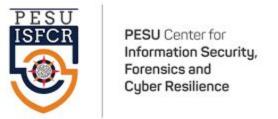
- First lets take a look at the disassembly
- Seems like a lot of disassembly.But don't
 Worry lets go through only important ones I
 have underlined



```
push
                                       rbp
  0×0000555555555556 <+1>:
                                        rbp, rsp
                                       rsp,0×10
                                       DWORD PTR [rbp-0×4],0×130104
                                       rdi,[rip+0×e99]
                                                               # 0×55555556004
                                lea
                                        eax,0×0
                      <+27>:
                                       0×5555555555040 <printfaplt>
                                       rax,[rbp-0×8]
                                        rsi, rax
                                                               # 0×55555556016
                                       rdi,[rip+0×e93]
                      <+46>:
                                       eax,0×0
                                       0×5555555555050 <__isoc99_scanf@plt>
  0×0000555555555518d <+56>:
                                        eax, DWORD PTR [rbp-0×8]
                                                               # 0×55555556016
                                lea
                                       rdi,[rip+0×e7d]
                                       eax,0×0
  0×0000555555555519e <+73>:
                                       0×5555555555040 <printf@plt>
                                call
                                       eax, DWORD PTR [rbp-0×8]
  0×000055555555551a3 <+78>:
                                       DWORD PTR [rbn-0×4].eax
                                       0×55555555551d9 <main+132>
                                       rdi,[rip+0×e67]
                                                               # 0×55555556019
                                lea
                                       0×5555555555040 <printf@plt>
  0×000055555555551b7 <+98>:
                                call
                                       edi,0×a
                                       0×5555555555030 <putchar@plt>
                                call
                                       rdi,[rip+0×e54]
                                                               # 0×55555556021
                                       eax,0×0
                                call
                                       0×5555555555040 <printf@plt>
                                       0×5555555551f4 <main+159>
                                       0×5555555555030 <putchar@plt>
                                       rdi,[rip+0×e41]
                                        eax,0×0
                                       0×5555555555040 <printf@plt>
                                       edi,0×a
                                       0×5555555555030 <putchar@plt>
                                call
                                       eax.0×0
  0×00005555555555204 <+175>:
End of assembler dump.
```



- The first important instruction calls printf with parameters. So it prints "input passcode" string
- •call 0x55555555555040 <printf@plt>
- Second Important instruction <u>mov eax,DWORD PTR [rbp-0x8]</u> seems to take Our Input from [rbp-0x8] and store it in eax



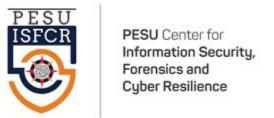
- The next most Important Instruction performs a cmp operation. Basically a if condition in Assembly
- So cmp checks if compares two registers by subtracting them
- Its similar to the SUB intstruction but IT does not affect operands
- so the assembly takes two paths one is JNE instruction fails(indicated by blue arrow) and IF JNE succeeds(indicated by red arrow)
- cmp DWORD PTR [rbp-0x4],eax
- The cmp instruction is comparing our value present in eax to value present in the address [rbp-0x4]
- So IF we can find the value present in [rbp-0x4] we have cracked the program!!



 Lets try to find the Value in Gdb By setting breakpoint just before the compare instruction

```
gdb-peda$ b *main+81
Breakpoint 2 at 0×5555555551a6
gdb-peda$ r
Starting program: /root/rev/crackMe
input PassCode! :1234
```

- We can examine a memory region in gdb by x/wd command 'w' is for examining in word length or 4 bytes 'd' is for viewing values in integer type
- We have to examine [rbp-0x4] so we can type x/wd \$rbp-0x4 to that memory region
- Lets do that in gdb



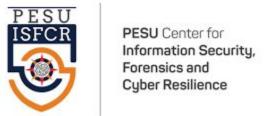
Looks like we Found our passcode present in the Binary

```
gdb-peda$ x/wd $rbp-0×4
0×7fffffffe05c: 1245444
```

Lets Input this value to the program and confirm whether this is correct.

```
root PES2UG19CS350)-[~/rev]
// ./crackMe
input PassCode! :1245444
1245444right !
Good job!
```

 Hooray!! Seems like the Passcode worked. We Have successfully cracked a simple binary using reverse engineering



- Sometimes Reading a Lot of disassembly on gdb can be quite difficult hence there are abundant of reversing engineering tools that make it easier to reverse
- 1) Ghidra Ghidra is a free and open source reverse engineering tool developed by the National Security Agency of the United States.you can get it From their official github page https://github.com/NationalSecurityAgency/ghidra
- 2) Ida interactive disassembler is also quite popular and used world wide by a lot of professionals You can visit their page to download or know more about ida https://hex-rays.com/ida-pro/

THANK YOU



PESU Center for Information Security, Forensics and Cyber Resilience