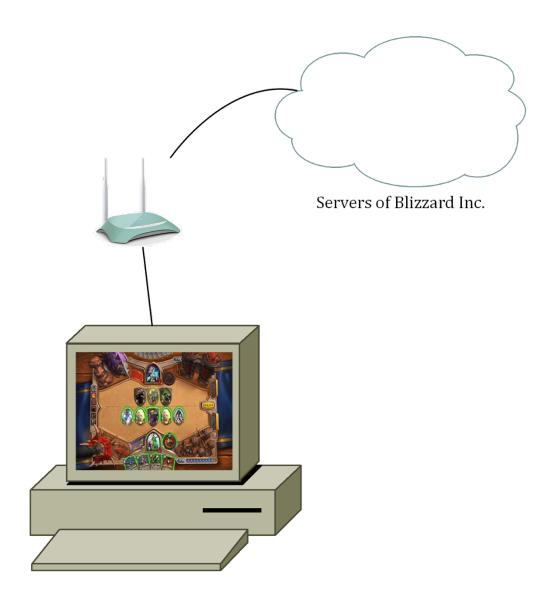
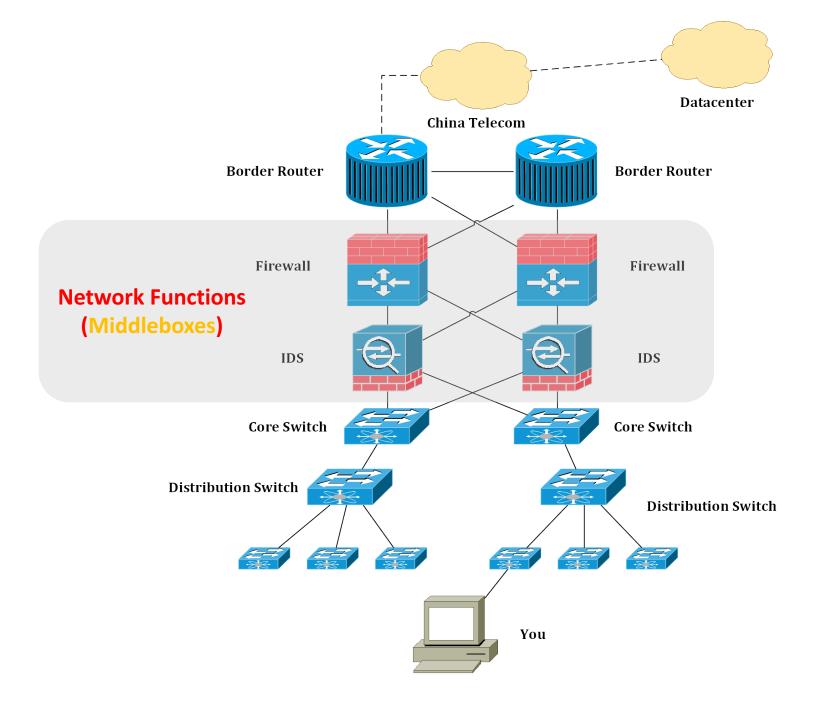
Network Function Virtualization

Haibo Chen
Institute of Parallel and Distributed Systems (IPADS)
Shanghai Jiao Tong University

http://ipads.se.sjtu.edu.cn/haibo_chen

How do you access servers?

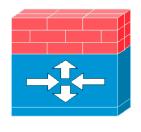








\$235,000



Firewall



\$**4,923**



Intrusion Detection System



\$43,700





Intel E5-2650 v3 with 20 cores, 40 Gbps Ethernet NIC, 64 GB RAM

¥45,000

Goal of NFV

- Save money
 - Commodity server hardware
 - Workload consolidation
 - Lower power consumption
 - Simplified network maintenance
- Make money
 - Accelerated service deployment
 - Network infrastructure as a service

What is NFV?

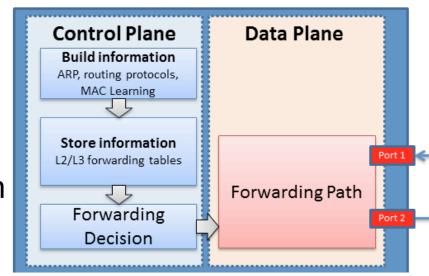
- Network Function
 - Middleboxes (firewall, IDS, WAN optimizer)

- Virtualization
 - Commodity hardware
 - Consolidation

 Network Services running on Cloud Infrastructure

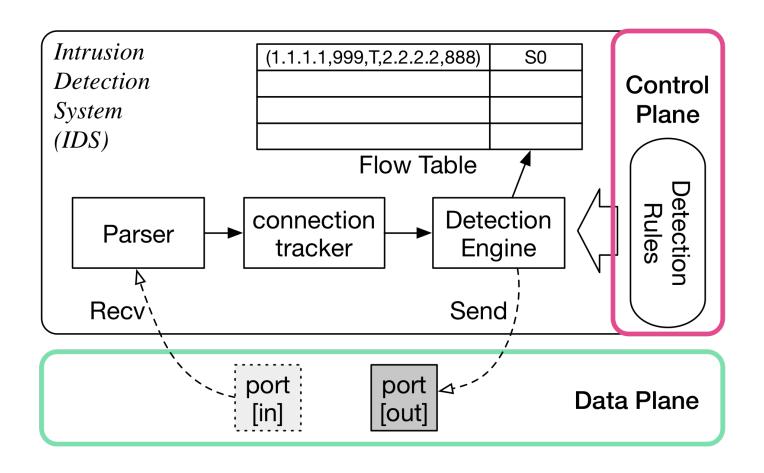
Relationship with SDN

- SDN: Software-defined networking
 - Focus on control plane innovation/virtualization
- NFV: network function virtualization
 - Focus on virtualizing data plane
- SDN & NFV
 - Can innovate separately
 - Or together
 - Like policy vs. mechanism

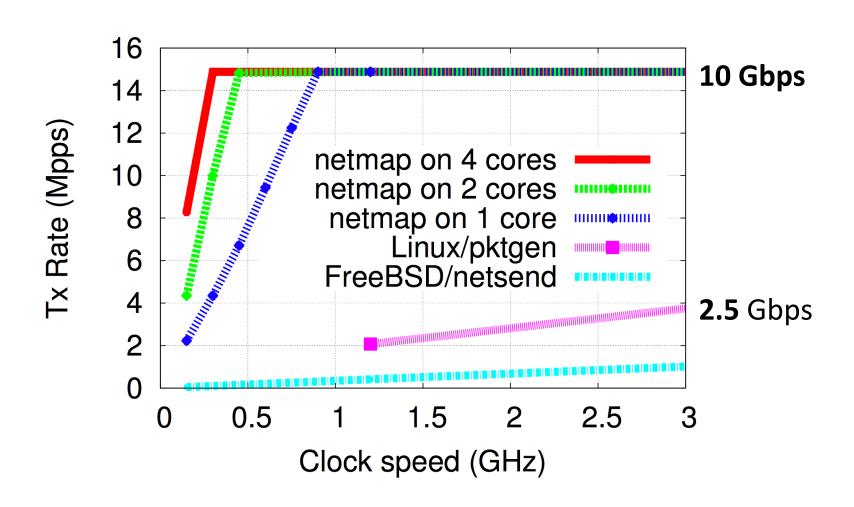


NFV: CHALLENGES

Challenges: Data Plane Performance



Challenges: Data Plane Performance



Middlebox	Functionality	Examples
NIDS	Monitor network activities for malicious behaviors or policy violations and report to operations staff	Bro IDS, Snort
Firewall	A firewall based on packet filter drops packets matching specific patterns and forwards the rest to next middlebox	iptables, pf
Load Balancer	Load balancers distribute network flows across several down- stream middleboxes	HAProxy, LVS
NAT	Remap one IP address space to another	iptables

Process

- No performance penalty
- No performance isolation

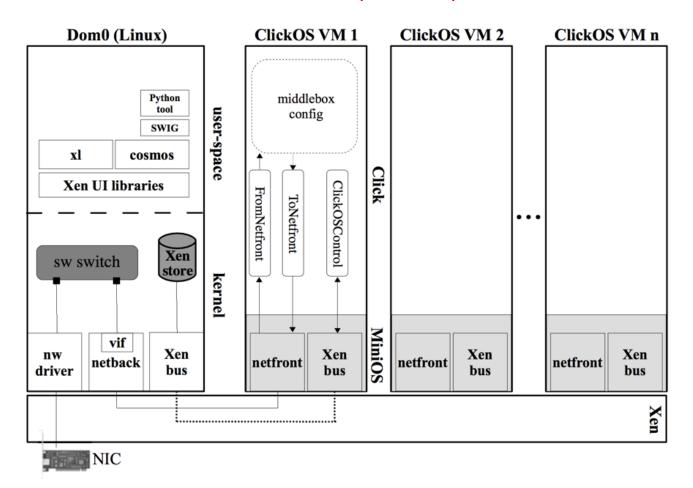
Virtual machine

- Virtualization overhead
- Performance isolation
- Easy deployment
- Mature technology

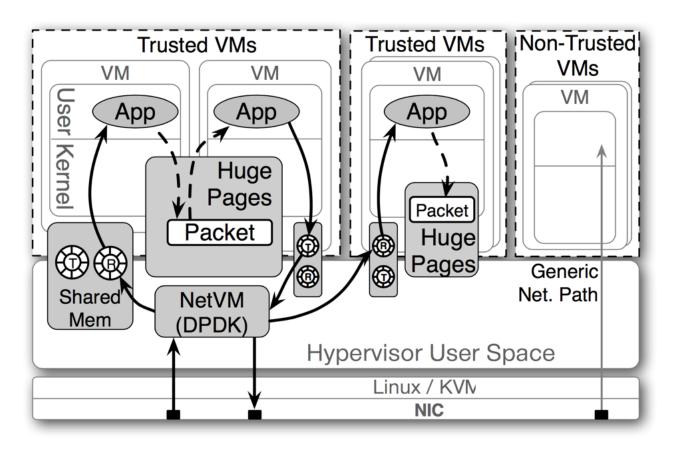
Container

- No performance penalty
- Performance isolation
- Easy deployment

Virtual machine clickos (NSDI'14)



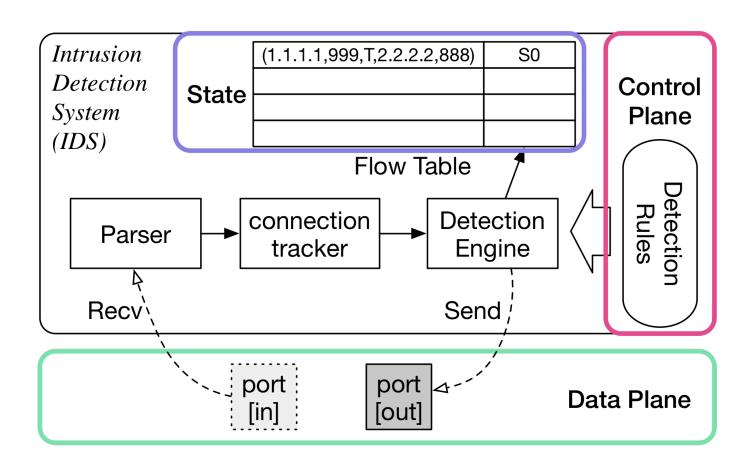
Virtual machine NetVM (NSDI'14)



Challenges: Scalability

- Dynamic network
 - Workload
 - Reconfiguration
 - Resource allocation
- Scalability
 - Scale up
 - Multicore hardware
 - NUMA aware
 - Scale out
 - Replication
 - Migration

Challenges: State Management



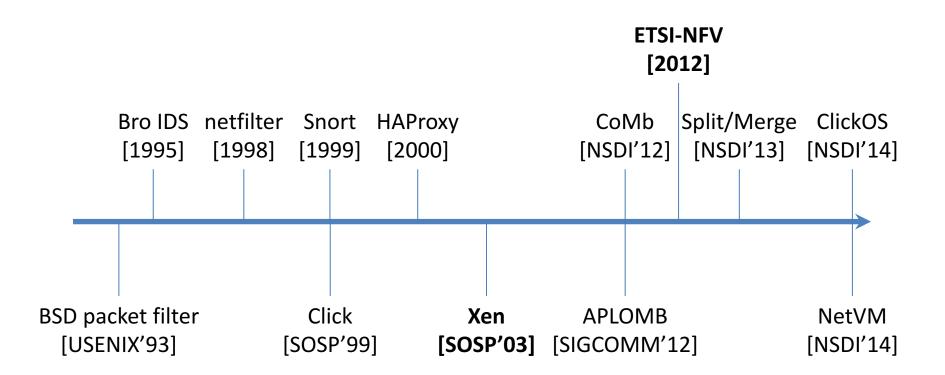
Challenges: State Management

- State consistency
 - Distributed instances
 - Recovery and migration
- Requirements
 - Classified states: internal & external
 - Transactional boundary
 - Move states between replicas
 - Correctly route traffic

NFV: History

- Software middlebox
 - IDS: Bro, Snort
 - Firewall: iptables (1998), BSD pf (USENIX'93)
 - Load Balancer: HAProxy, LVS
 - NAT: iptables
- Click (SOSP'99)
- Middlebox consolidation
 - NIDS Cluster (RAID'07)
 - APLOMB (SIGCOMM'12)
 - CoMb (NSDI'12)
 - ClickOS (NSDI'14)

NFV: History



Software Packet Processing

Network Function Virtualization

NFV: SUMMARY

NFV: Summary

















NFV: Summary

Virtualization of network appliances

Performance and flexibility

- System design
 - Multicore and NUMA hardware
 - Distributed processing
 - I/O optimization

The Click Modular Router

Robert Morris, Eddie Kohler, John Jannotti, M. Frans Kaashoek

Background

- New network functions
 - NIDS
 - Random early detection dropping policy
 - Deep packet inspection
 - GFW
- Commercial off-the-shelf (COTS) hardware
 - Expensive
 - Closed
 - Static
 - Inflexible

Motivation

- Flexible
 - New functionality
- Modular
 - Combine elements
- Open
 - Allow developers to add elements
- Efficient
 - Comparable with hardware

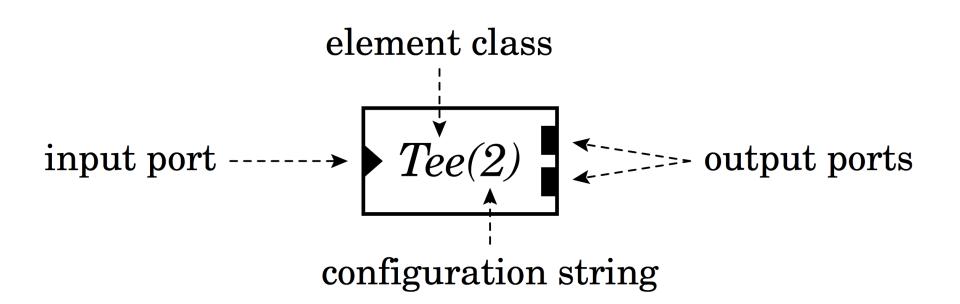
Click: Overview

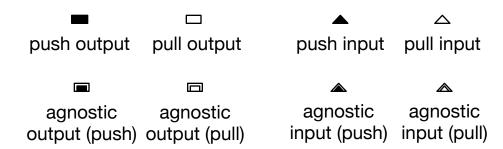
- A modular design to allow easy construction of routers and other middleboxes
 - Elements
 - Push
 - Pull
 - A domain-specific language (DSL)
 - Declare elements
 - Connect elements
 - Lightweight framework
 - Orchestrate elements into a router

Elements

- Building block of Click router
 - C++ classes
 - Scheduled by framework
- Ports
 - Endpoints of connections
 - Associated with push or pull function (or agnostic)
- Configuration
 - Parameters to initialize elements
 - Elements specific

Element Example

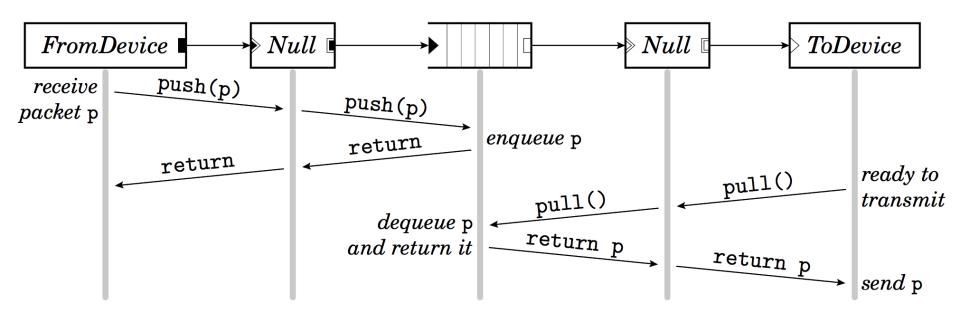




C++ Class of Elements

```
class NullElement : public Element {
   public:
   NullElement()
        { add_input(); add_output(); }
   const char *class_name() const
       { return "Null"; }
   PushOrPull default_processing() const
        { return AGNOSTIC; }
   NullElement *clone() const
        { return new NullElement; }
   void push(int port_number, Packet *p)
        { output(0).push(p); }
   Packet *pull(int port_number)
        { return input(0).pull(); }
```

Push and Pull

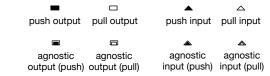


Push

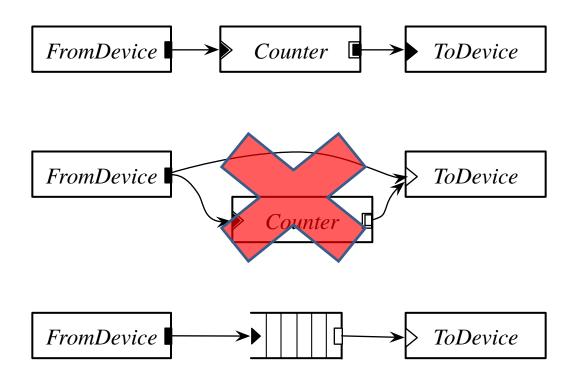
From source element to downstream Triggered by event, e.g. packet arrival Denoted by solid square or triangle

Pull

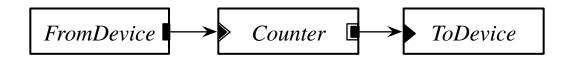
From destination element to upstream Packet transmission
Denoted by empty square or triangle



More Examples



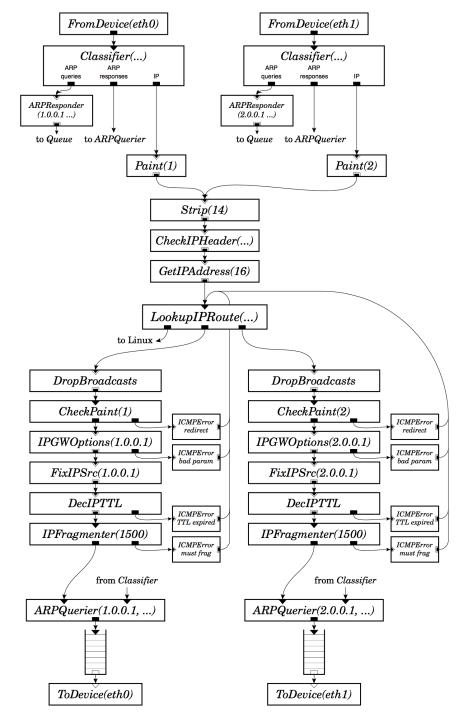
Packet Counter



```
// Declare three elements...
src :: FromDevice(eth0);
ctr :: Counter;
sink :: Discard;
// ... and connect them together
src -> ctr;
ctr -> sink;
```

Ping Responder

```
1 define($IP 192.168.13.235);
 2 define($MAC 00-15-17-15-5d-75);
 4 tap :: KernelTap(192.168.13.235/16, ETHER $MAC);
 6 c :: Classifier(
          12/0806 20/0001, // ARP Requests goes to output 0
       12/0806 20/0002, // ARP Replies to ouput 1
           12/0800, // ICMP Requests to output 2
10
           -);
11
12 arpg :: ARPQuerier($IP, $MAC);
13 arpr :: ARPResponder($IP $MAC);
14
15 tap-> c;
16 c[0] -> ARPPrint -> arpr -> Print("arpr", 60) -> tap;
17 c[1] -> [1]arpq;
18 Idle -> [0]arpq;
19 arpq -> Print("arpq", 60) -> ARPPrint -> tap;
20 c[2] -> CheckIPHeader(14) -> ICMPPingResponder() -> EtherMirror() -> tap;
21 c[3] -> Discard;
```



Scheduling

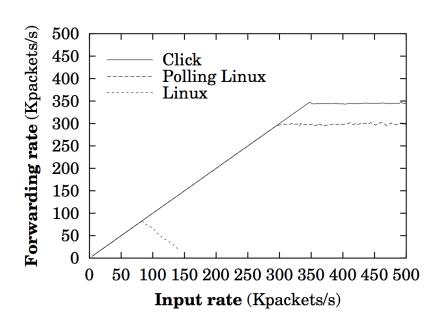
- Single thread in the original paper
- Only some elements need to be scheduled
- Click schedules the CPU with a task queue
 - Flexible and lightweight stride scheduling [Waldspurger'95]
 - Element is the unit of CPU scheduling
 - Most elements are never placed on the task queue
 - Elements frequently initiates push or pull should be scheduled (FromDevice, ToDevice)

Implementation

- Userlevel
 - Run as normal thread
 - FromDevice: TUN/TAP, libpcap
- Kernel module
 - Run as kernel thread
 - Loops over the task queue and run each task
 - Only interrupts can preempt this thread
 - Voluntarily gives up CPU to Linux
- ClickOS [NSDI'14]
 - Run in MiniOS on Xen
 - Use netmap [ATC'12] to get packets

Maximum Loss-Free Forwarding Rate

- Click IP router
 - 8 interfaces
 - 161 elements
 - 64 bytes packets size
- Forwarding rate versus input rate
- MLFFR reflects the router's behavior under load



Adding Network Functions

Simple

Passes packets from one interface to another

Switch

L2 switch

IP

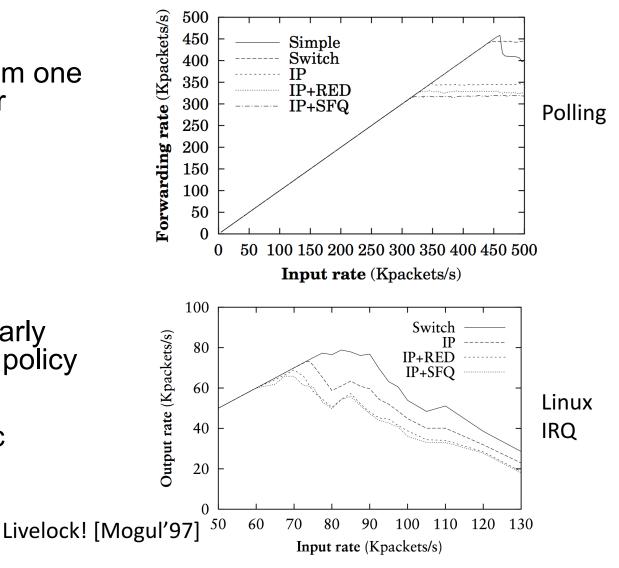
IP router

IP+RED

Router + random early detection dropping policy

IP+SFQ

Router + stochastic



Summary

- NFV: everything old is new again
- Cloud networks need
 - Flexible operational model
 - State-of-the-art infrastructure services
 - Programmable network functions
- NFV is changing the networking industry

