

Task 4

In this case we need to convert a 4-bit number into its complement to two. A truth table is built first with four outputs corresponding to the four input bits of the number complemented, as shown below.

| 4-Bit In | | | | 2-Comp. Out | | | |
|----------|-----|-----|-----|-------------|-------|-------|-------|
| d | c | b | a | f_d | f_c | f_b | f_a |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 |
| 0 | 0 | 1 | 0 | 1 | 1 | 1 | 0 |
| 0 | 0 | 1 | 1 | 1 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 1 | 0 | 0 | 1 |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | 1 |

Table 1: Outputs with complement to two.

The output functions are expressed based on the minterms. They are simplified using Karnaugh's Maps. Starting with f_d function:

$$f_d = \sum (m_1, m_2, m_3, m_4, m_5, m_6, m_7, m_8)$$

| | | | | | |
|----|----|----|----|----|----|
| | | ba | | | |
| | | 00 | 01 | 11 | 10 |
| dc | 00 | 0 | 1 | 1 | 1 |
| | 01 | 1 | 1 | 1 | 1 |
| | 11 | 0 | 0 | 0 | 0 |
| | 10 | 1 | 0 | 0 | 0 |

Solving the map with the indicated groups, the simplified function remains as:

$$f_d = (d \cdot \bar{c} \cdot \bar{b} \cdot \bar{a}) + (\bar{d} \cdot c) + (\bar{d} \cdot \bar{c} \cdot a) + (\bar{d} \cdot b \cdot \bar{a})$$

Now, taking the f_c function, we do the same:

$$f_c = \sum (m_1, m_2, m_3, m_4, m_9, m_{10}, m_{11}, m_{12})$$

| | | ba | | | |
|----|----|----|----|----|----|
| | | 00 | 01 | 11 | 10 |
| dc | 00 | 0 | 1 | 1 | 1 |
| | 01 | 1 | 0 | 0 | 0 |
| | 11 | 1 | 0 | 0 | 0 |
| | 10 | 0 | 1 | 1 | 1 |

With the indicated groups, we get the simplified function:

$$f_c = (c \cdot \bar{b} \cdot \bar{a}) + (\bar{c} \cdot a) + (\bar{c} \cdot b \cdot \bar{a})$$

Next, with the f_b function:

$$f_b = \sum (m_1, m_2, m_5, m_6, m_9, m_{10}, m_{13}, m_{14})$$

| | | ba | | | |
|----|----|----|----|----|----|
| | | 00 | 01 | 11 | 10 |
| dc | 00 | 0 | 1 | 0 | 1 |
| | 01 | 0 | 1 | 0 | 1 |
| | 11 | 0 | 1 | 0 | 1 |
| | 10 | 0 | 1 | 0 | 1 |

Solving the map, we get:

$$f_b = (\bar{b} \cdot a) + (b \cdot \bar{a})$$

For the last function f_a :

$$f_a = \sum m_1, m_3, m_5, m_7, m_9, m_{11}, m_{13}, m_{15}$$

In the table it can be observed that the output depends directly from input a . We can write the simplified function without making the Karnaugh's map:

$$f_a = a$$

Having already the four output functions, the implementation can be carried out in a circuit with AND, OR and NOT logic gates, as shown in the next page.

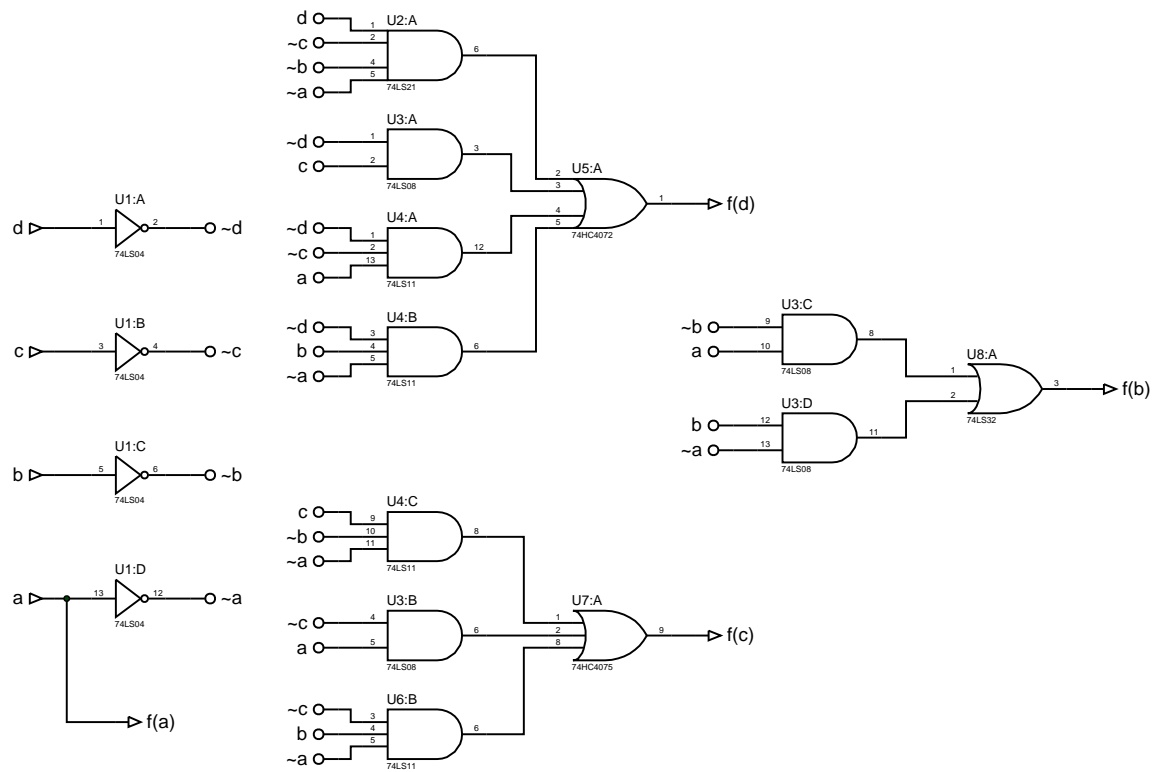


Figure 1: Implementation of 2-complement circuit for a 4-bit input number - Designed in Proteus 7.8