



Instituto Tecnológico de Buenos Aires
18 de Octubre de 2018

Trabajo de Laboratorio N° 2

Electrónica III

Grupo 7

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Task 1

In this section, a state machine will be developed for controlling the switching on and off of two pumps, to fill a tank. The are controlled by two sensors from the upper part of the tank (S) and the lower part of the tank (I). The actions to take are as follows:

- Tank full: $S = I = 1$ - Pumps OFF
- Tank empty: $S = I = 0$ - Pumps ON
- Half full tank: $S = 0$ & $I = 1$ - Pumps alternate

With this in mind, a Moore machine is developed as follows.

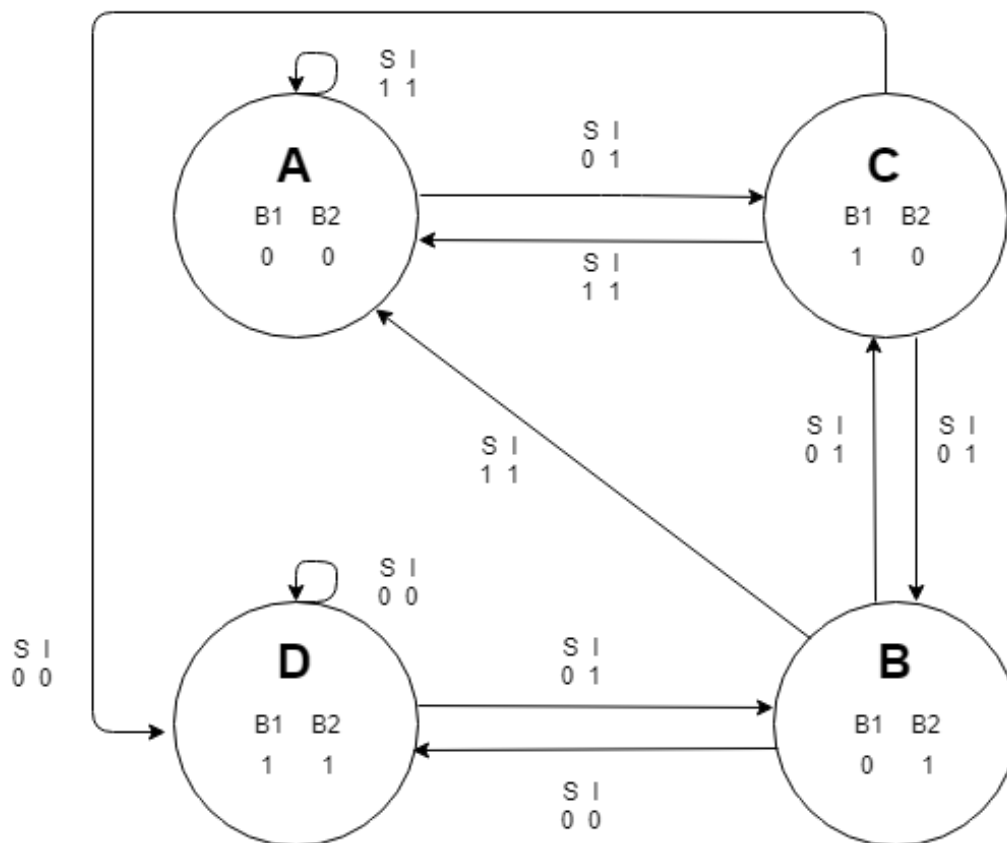


Figure 1: Moore state machine diagram

Using two bits to assign the states, a table of transitions is made, as shown below.

Estado Actual		Estado Siguiente				Salidas	
	y2 - y1	S - I				B1	B2
		00	01	10	11		
A	00	X	C	X	A	0	0
B	01	D	C	X	A	0	1
C	10	D	B	X	A	1	0
D	11	D	B	X	X	1	1

Figure 2: Moore state machine - Transitions

From the table, using Karnaugh's maps the functions for the state variables and the two pumps outputs are made as shown below.

		ba			
		00	01	11	10
dc	00	X	1	1	1
	01	1	1	0	0
	11	0	0	X	0
	10	X	X	X	X

		ba			
		00	01	11	10
dc	00	X	1	1	1
	01	1	1	0	0
	11	0	0	X	0
	10	X	X	X	X

Figure 3: Maps for Y_2 (left) and Y_1 (right) functions.

Where from the left table $Y_2 = \bar{I} + \bar{S} \cdot \bar{y}_2$, and from the right $Y_1 = \bar{I} + \bar{S} \cdot y_2$. And for the outputs, $B_1 = y_2$ and $B_2 = y_1$. Finally, the state machine is implemented using D Flip Flops:

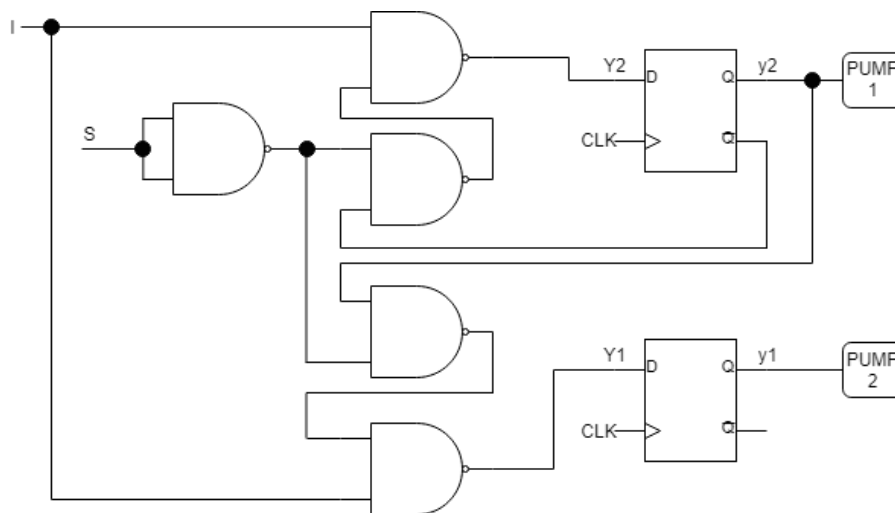


Figure 4: Moore state machine - Circuit implementation

On the other side, the same system is implemented now using a Mealy state machine, as shown below.

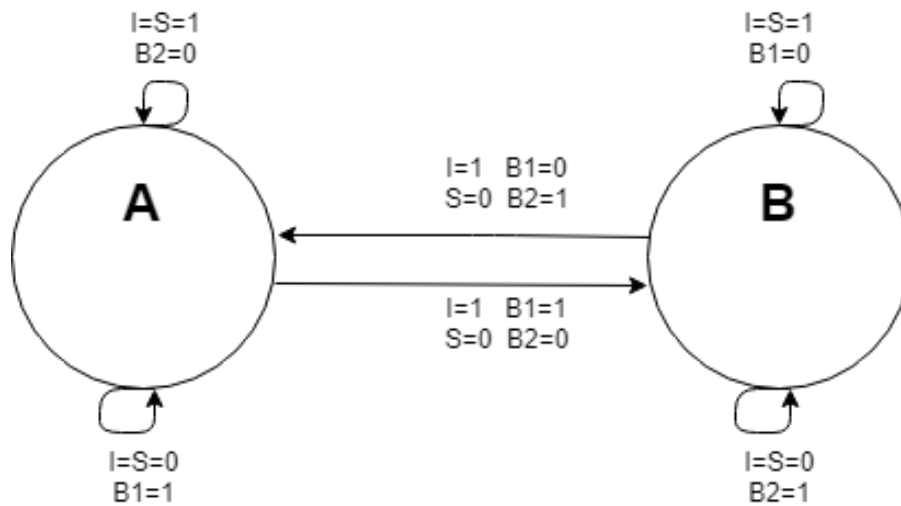


Figure 5: Mealy state machine diagram

Notice that the direct connection between the input and the pumps outputs reduces the number of states from four to two, in comparison with the Moore machine. Using one bit for the states, a table of transitions is made as follows.

Estado Actual		Estado Siguiente				Salidas			
	y	S - I				S - I			
		00	01	10	11	00	01	10	11
		y				B1 - B2	B1 - B2	B1 - B2	B1 - B2
A	0	A	B	A	X	11	10	XX	00
B	1	B	A	B	X	11	01	XX	00

Figure 6: Mealy state machine - Transitions

Using again Karnaugh's maps, the functions for the state variable and the two pumps outputs are made as shown below.

a	bc			
	00	01	11	10
0	0	1	0	X
1	1	0	1	X

a	bc			
	00	01	11	10
0	1	1	0	X
1	1	0	0	X

a	bc			
	00	01	11	10
0	1	0	0	X
1	1	1	0	X

Figure 7: Maps for Y (left), B_1 (center) and B_2 (right) functions.

Where from the left table $Y = \bar{y} \cdot \bar{S} \cdot I + y \cdot \bar{I} + y \cdot S$. From the center table $B_1 = \bar{y} \cdot \bar{S} + \bar{S} \cdot \bar{I}$, and from the right table $B_2 = \bar{S} \cdot \bar{I} + y \cdot \bar{S}$. Finally, the state machine is implemented using one D Flip Flop.

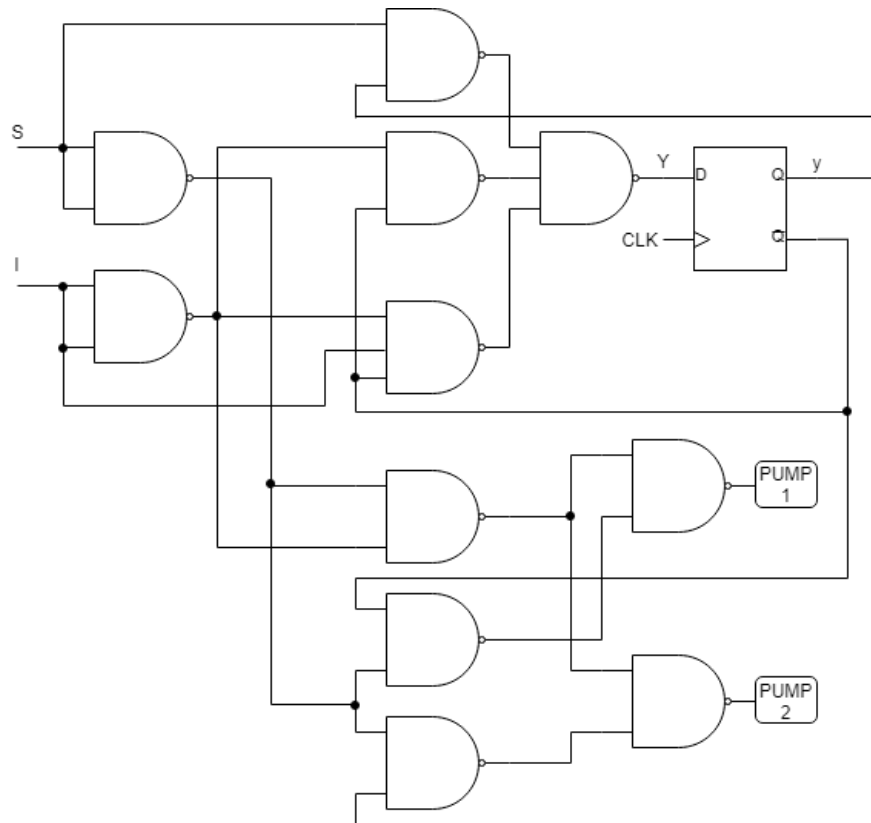


Figure 8: Mealy state machine - Circuit implementation

Verilog tests

This task was tested using verilog tests. Here we show the results of part A and part B.

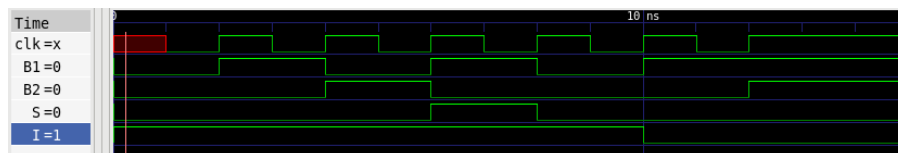


Figure 9: Mealy state machine - Verilog test results

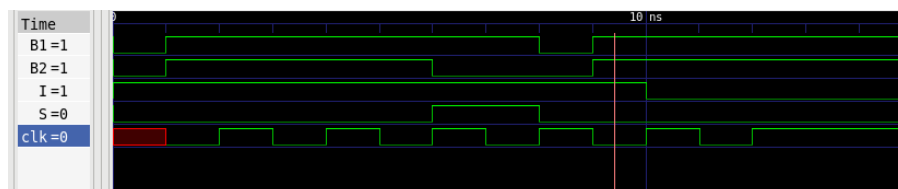


Figure 10: Moore state machine - Verilog test results

Apparently, they seem correct.

Task 2

In this section, the objective is to recognize a sequence of 4 bits that come in a synchronized way. If the sequence is recognized, an output is turned on. Using a Moore's state machine, the resulting diagram is as shown below.

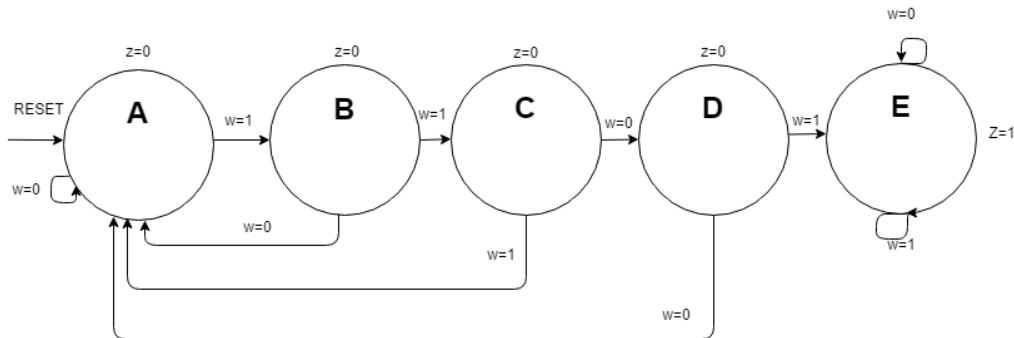


Figure 11: Moore state machine diagram

Notice that when the sequence is recognized, the machine needs to be reseted to detect a new combination. With the diagram, the following transition table is made.

Estado Actual		Estado Siguiente		Salidas
	y3 - y2 - y1	W		Z
		0	1	
A	000	A	B	0
B	001	A	C	0
C	010	D	A	0
D	011	A	E	0
E	100	E	E	1

Figure 12: Moore state machine - Transitions

Using Karnaugh's maps, the functions for the diferent states and the output are made as follows.

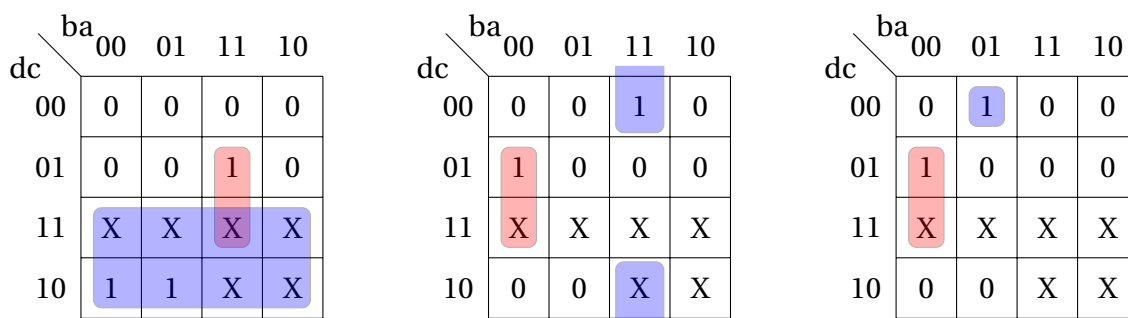


Figure 13: Maps for Y_3 (left), Y_2 (right), and Y_1 (center) functions.

Where $Y_3 = y_3 + y_2 \cdot y_1 \cdot W$, $Y_2 = y_2 \cdot \overline{y_1} \cdot \overline{W} + \overline{y_2} \cdot y_1 \cdot W$, and $Y_1 = \overline{y_3} \cdot \overline{y_2} \cdot \overline{y_1} \cdot W + y_2 \cdot \overline{y_1} \cdot \overline{W}$. From the table of transitions, $Z = y_3$.

With the functions, the state machine is implemented using three D Flip Flops as shown below.

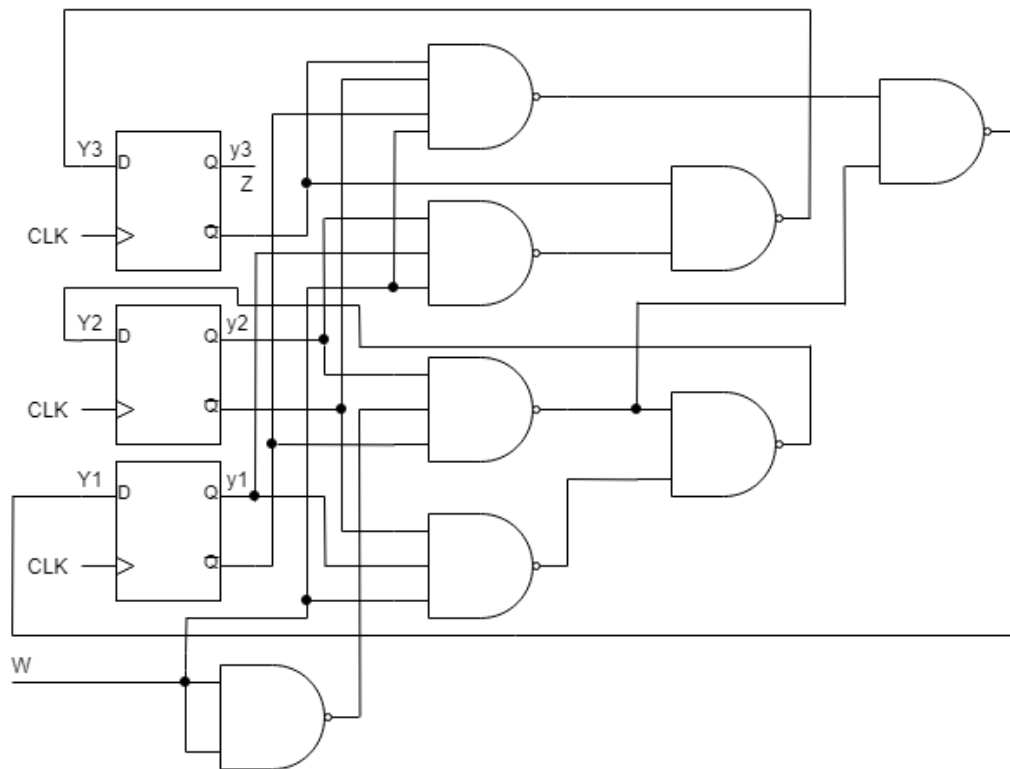


Figure 14: Moore state machine - Circuit implementation

Now, the same system is implemented using a Mealy's state machine, wick resulting diagram is shown below.

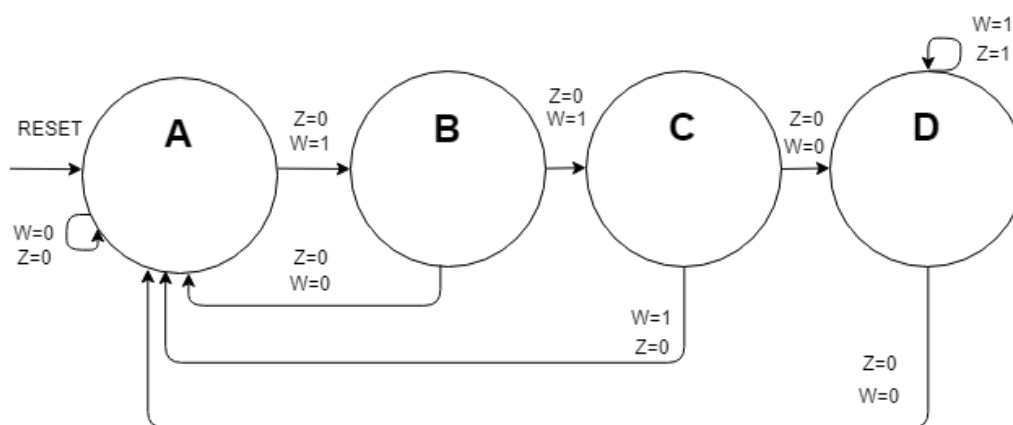


Figure 15: Mealy state machine diagram

Using the diagram, a table with the state transitions is made.

Estado Actual		Estado Siguiente		Salidas	
	y2 - y1	W		Z	
		0	1	W=0	W=1
A	00	A	B	0	0
B	01	A	C	0	0
C	10	D	A	0	0
D	11	A	D	0	1

Figure 16: Mealy state machine - Transitions

Using Karnaugh's maps, the functions for the states and the output are made below.

		bc			
		00	01	11	10
a					
0		0	0	0	1
1		0	1	1	0

		bc			
		00	01	11	10
a					
0		0	0	0	1
1		1	0	1	0

		bc			
		00	01	11	10
a					
0		0	0	0	0
1		0	0	1	0

Figure 17: Maps for Y_2 (left), Y_1 (center) and Z (right) functions.

Where $Y_2 = W \cdot y_1 + \overline{W} \cdot y_2 \cdot \overline{y_1}$, $Y_1 = W \cdot \overline{y_2} \cdot \overline{y_1} + W \cdot y_2 \cdot y_1 + \overline{W} \cdot y_2 \cdot \overline{y_1}$, and $Z = W \cdot y_2 \cdot y_1$. With the defined functions, the state machine is implemented with 2 D Flip Flops. In this case is used one less flip flop, and the machine can used again without reset it.

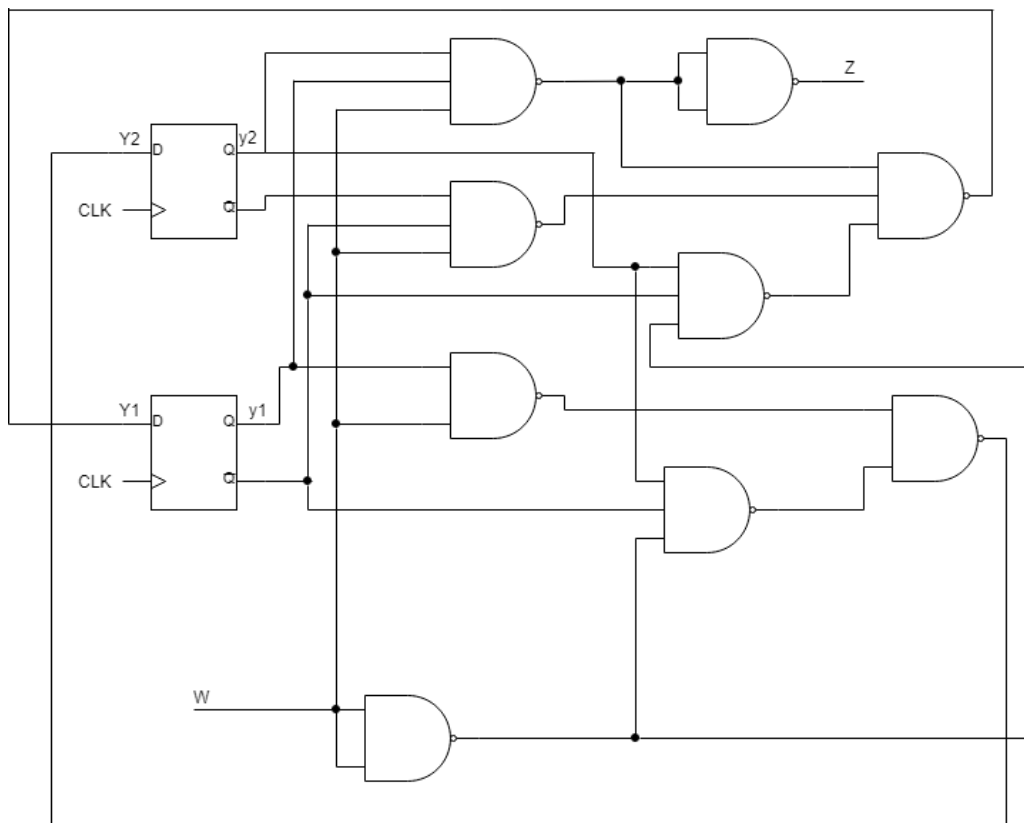


Figure 18: Mealy state machine - Circuit implementation

Task 3

In this section, is implemented a Moore's state machine already defined, as shown below.

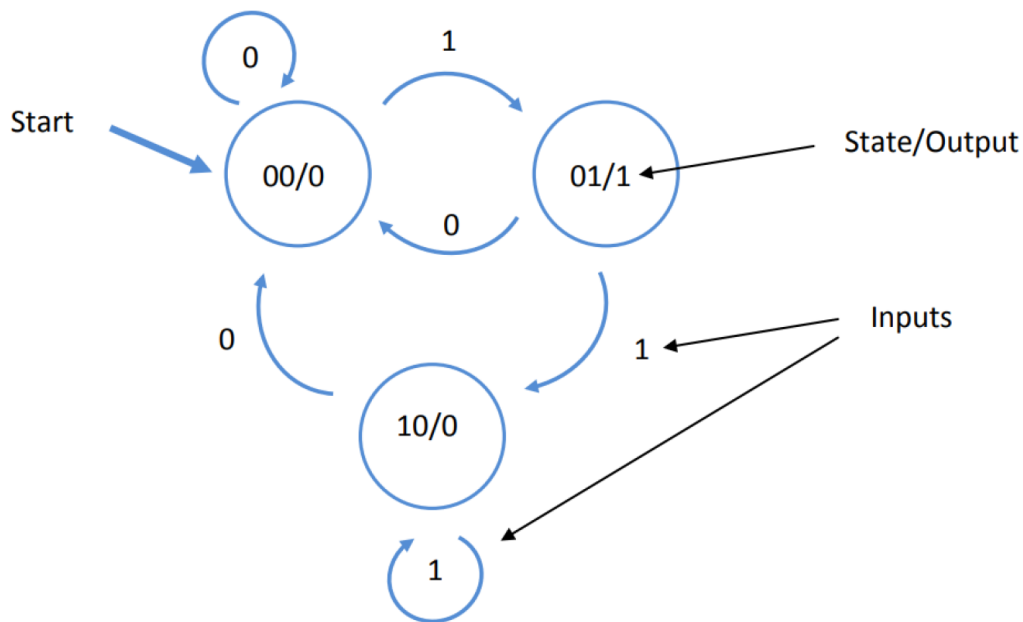


Figure 19: Moore state machine diagram

Using the diagram, the following table of transitions is made.

Estado Actual		Estado Siguiente		Salidas
	y2 - y1	W		Z
		0	1	
A	00	A	B	0
B	01	A	C	1
C	10	A	C	0

Figure 20: Moore state machine - Transitions

With the transitions, using Karnaugh's maps (see resolution in *Annex*), the functions for the states and the output result as follows: $Y_2 = W \cdot y_1 + W \cdot y_2$, and $Y_1 = W \cdot \overline{y_2} \cdot \overline{y_1}$.

From the transitions table, it is simple to see that $Z = y_1$. With the functions, the state machine is implemented using two D Flip Flops as follows.

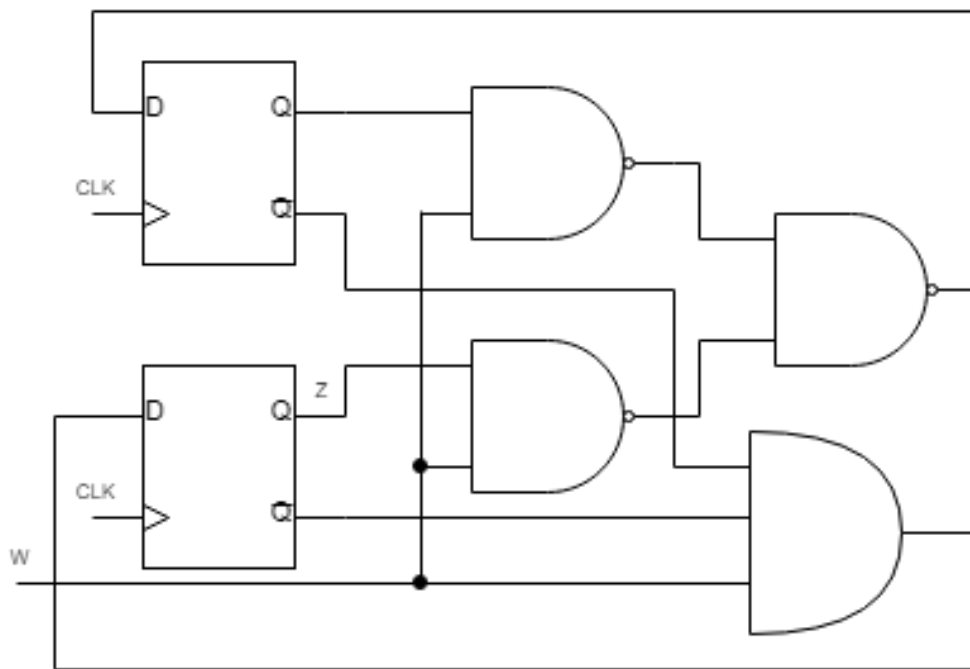


Figure 21: Moore state machine - Circuit implementation

Now the same system is implemented using a Mealy's state machine, wick resulting diagram is shown below.

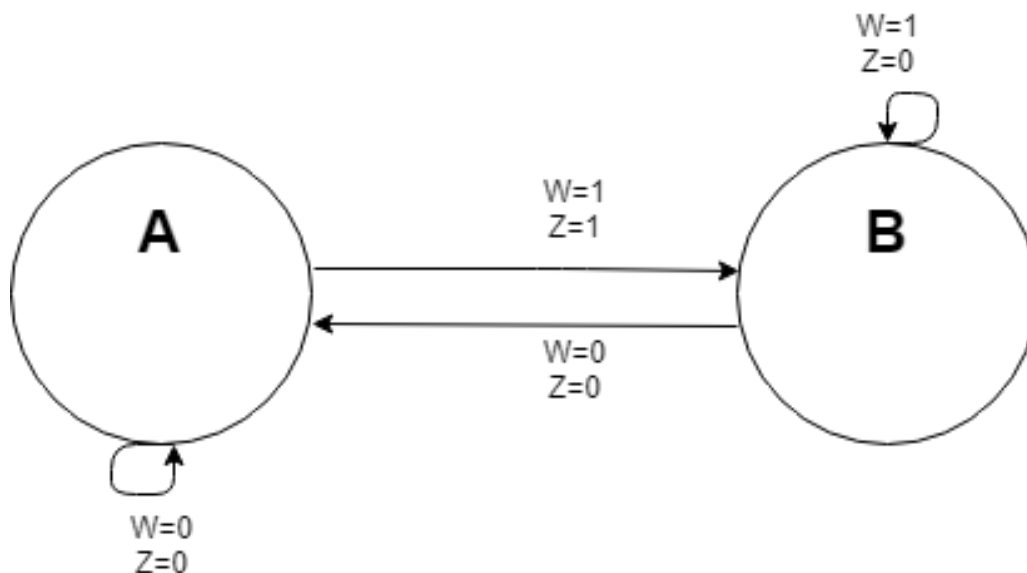


Figure 22: Mealy state machine diagram

Notice that it requires one less state than Moore's machine because of the direct connection of the from the input to the output.

The following transition table is made using the diagram.

Estado Actual		Estado Siguiente		Salidas	
	y	W		Z	
		0	1	W=0	W=1
A	0	A	B	0	1
B	1	A	B	0	0

Figure 23: Mealy state machine - Transitions

With the table, using Karnaugh's maps (see resolution in *Annex*), are made the functions for the states and the output as follows: $Y = W$, and $Z = \bar{y} \cdot W$. With the defined functions, the state machine is implemented using one D Flip Flop as shown below.

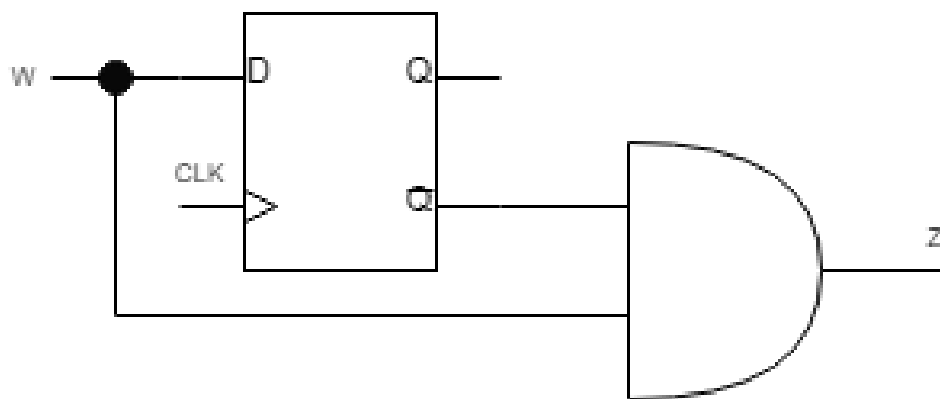


Figure 24: Mealy state machine - Circuit implementation

Since the internal logic works with 3.3V power supply, and the external signals work with 5V, level shifters are implemented using BJT transistors. For adapting the inputs of CLK and W, the circuits are shown below.

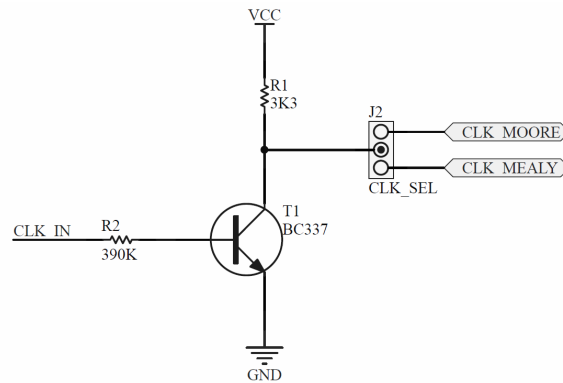


Figure 25: Level shifter for CLK from 5V to 3.3V (VCC)

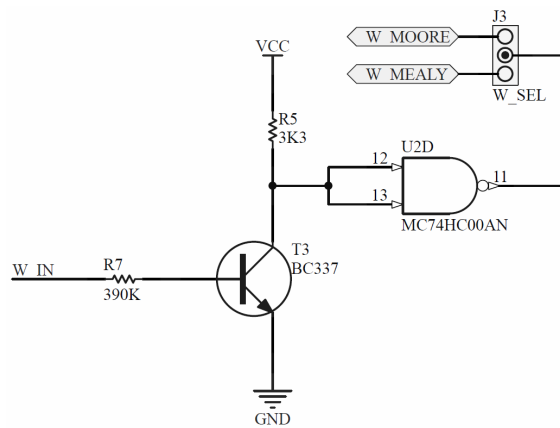


Figure 26: Level shifter for W from 5V to 3.3V (VCC). The inverting gate is to compensate the transistor logic inversion.

And for the outputs (Moore and Mealy machines) the driver circuit is as shown below.

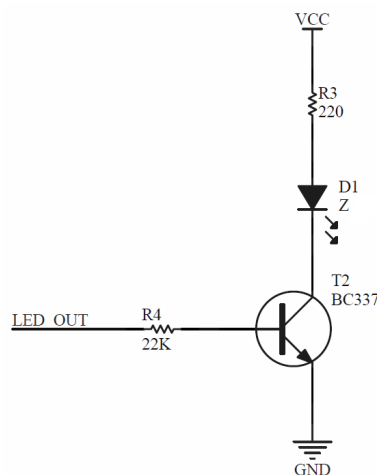


Figure 27: Driver for LED output.

Level shifter for inputs

From the implemented circuit:

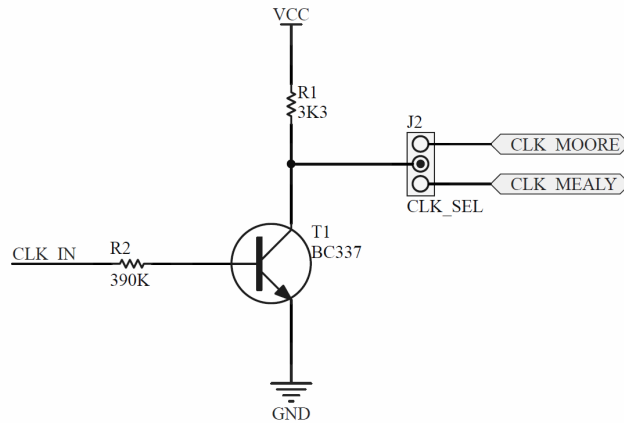


Figure 28: Level shifter for CLK from 5V to 3.3V (VCC)

Usign for $I_{SAT} = 1mA$, considering $VCE_{SAT} = 0.2V$, the equation from the out mesh:

$$3.3V - VCE_{SAT} - I_{SAT}R_1 = 0$$

$$\frac{3.3V - VCE_{SAT}}{I_{SAT}} = R_1 = 3.1K\Omega$$

Normalizing we have $R_1 = 3.3K\Omega$. Considering $HFE_{MIN} = 100$, from the input mesh:

$$5V - VBE_{ON} - \frac{I_C}{HFE_{MIN}}R_2 = 0$$

$$\frac{5V - VBE_{ON}}{I_C}HFE_{MIN} = R_2 = 430K\Omega$$

Normalizing we have $R_2 = 390K\Omega$.

Driver for output led

Taking the implemented circuit:

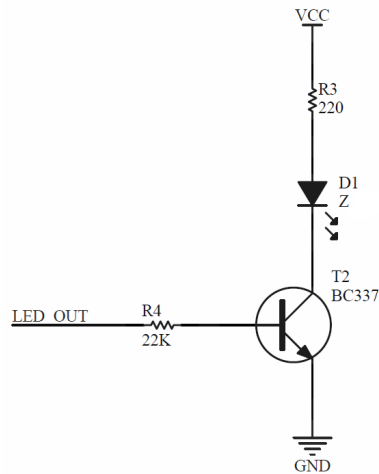


Figure 29: Driver for LED output.

Usign for $I_{LED} = 10mA$, considering $V_{CE_{SAT}} = 0.2V$ and $V_{LED} = 2V$, the equation from the out mesh:

$$3.3V - V_{LED} - V_{CE_{SAT}} - I_{LED}R_3 = 0$$

$$\frac{3.3V - V_{LED} - V_{CE_{SAT}}}{I_{LED}} = R_3 = 110\Omega$$

Normalizing we have $R_3 = 220\Omega$. Considering $HFE_{MIN} = 100$, from the input mesh:

$$3.3V - V_{BE_{ON}} - \frac{I_C}{HFE_{MIN}}R_4 = 0$$

$$\frac{3.3V - V_{BE_{ON}}}{HFE_{MIN}} = R_4 = 26K\Omega$$

Normalizing we have $R_4 = 22K\Omega$, to guarantee saturation.

Verilog tests

This task was tested using verilog tests. Here we show the results of part A and part B.

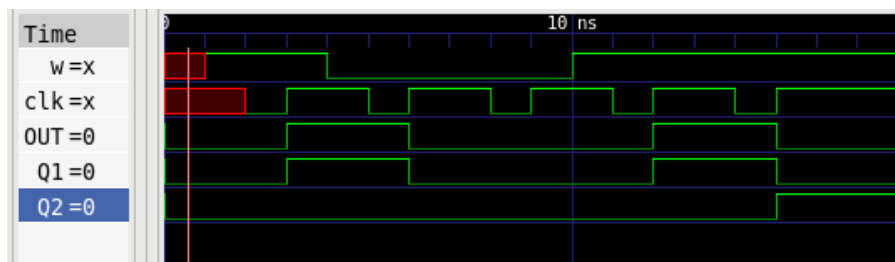


Figure 30: Mealy state machine - Verilog test results

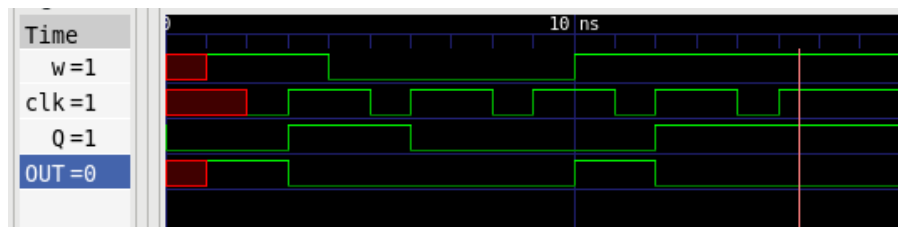


Figure 31: Moore state machine - Verilog test results

They seem correct.