

A Simple Analysis Of CTF Challenge From A CTF Beginner

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0x0 Analysis

First, Let's run the crack.exe in CMD.exe with the test SN "12345":

```
C:\Users\T\Desktop\CTF3r\第一周11月19-24\Re\题目1\CrackMe>CrackMe.exe 12345
registration failed
```

It shows the common string of registration failed, then I try to search the string in IDA, but no result, so it should be encrypted by author, let's extract all strings with string.exe, and I found some suspicious strings:

```
umen
!fac
{{T5|Y
ucce
lly
Y
rastontiai fqed
```

0000000140038050	93 90 8D 96 86 03 BF F7	4A 14 7B 7B	54 35 7C 59	舜·嶂·..亏·J·f{T5 Y
0000000140038060	9B 9B DF 8C 75 63 63 65	8E 8C 99 8A 6C 6C 79 00		凍·邵·ucce峒·模·lly
0000000140038070	8E 9A 98 96 72 61 73 74	6F 6E 74 69 61 69 20 66		韶·標·rastontiai·f
0000000140038080	71 65 64 00 00 00 00 00	34 94 03 40 01 00 00 00		qed.....4..@....
0000000140038090	FF FF FF FF 00 00 00 00	00 00 00 00 00 00 00 00	
00000001400380A0	0A 00 00 00 00 00 00 00	00 00 00 00 00 00 00 00	

successfully

it seems like the encrypted registration successfully string and registration failed string, look them in IDA pseudo-code:

```

if ( v118 >= 0 )
{
LABEL_118:
v321 = -dword_140038070;
v322 = (dword_140038074 >> 16) + (dword_140038074 << 16);
v323 = (dword_140038078 >> 16) + (dword_140038078 << 16);
v324 = (dword_14003807C >> 16) + (dword_14003807C << 16);
v325 = dword_140038080 - 5;
sub_140001A40(&v506, &v321, 19i64);
v180 = &v506;
if ( v508 >= 0x10 )
{
v180 = v506;
v181 = sub_140001340(&qword_140039930, v180, v507); // registration failed
sub_140001800(v181);
sub_140001FD0(&v506);
v226 = -293063854;
v178 = -1773405844;
v179 = -827325493;
}
else
{
v326 = 5 * dword_140038058;
v327 = 7 * dword_14003805C - 217;
v328 = -dword_140038060;
v224 = dword_140038064;
v329 = dword_140038064 + (dword_140038064 & 0xFFFF0000);
v330 = 1 - dword_140038068;
v331 = dword_14003806C;
sub_140001A40(&v503, &v326, 23i64);
v176 = &v503;
if ( v505 >= 0x10 )
{
v176 = v503;
v177 = sub_140001340(&qword_140039930, v176, v504); // registration successfully
sub_140001800(v177);
sub_140001FD0(&v503);
v225 = 0xEE883352;
v178 = -858993452;
v179 = -613566752;
}
}
}

```

We can see they are in the if-else structure statement, if the value of `v118 >= 0`, it shows us the failer string, otherwise the successful string (so the function `sub_1400061A0` is the `SN_Check` function) where does the `v118` come from? look forward for the previous reference, we found lots of Junk Code (judge them by the the var's reference):

```

v124 = -(0xCCCCCCCC
+ (((0x33333333 * v122 + 7) >> 16)
+ ((0x33333333 * v122 + 7 + (((0x33333333 * v122 + 7) >> 16) * ((0x33333333 * v122 + 7) >> 16))) << 16)) >> 16)
+ ((0xCCCC * ((0x33333333 * v122 + 7) >> 16)) << 16)
+ 7;
v125 = -(v122
+ ((7 - v123) << 32)
+ 1
- (((v124
+ ((v124 >> 8) + 7 + (v124 << 8))
- (7
- (0xCCCCCCCC
* (((0x33333333 * v122 + 7) >> 16)
+ ((0x33333333 * v122 + 7 + (((0x33333333 * v122 + 7) >> 16) * ((0x33333333 * v122 + 7) >> 16))) << 16)) >> 16)
+ ((5
- (v124
+ (v124 >> 8)
+ (v124 << 8)
+ (0xCCCCCCCC
* (((0x33333333 * v122 + 7) >> 16)
+ ((0x33333333 * v122 + 7 + (((0x33333333 * v122 + 7) >> 16) * ((0x33333333 * v122 + 7) >> 16))) << 16)) >> 16)
- v122);
v126 = v125
+ 1
- (((v124
+ ((v124 >> 8) + 7 + (v124 << 8))
- (7
- (0xCCCCCCCC
* (((0x33333333 * v122 + 7) >> 16)
+ ((0x33333333 * v122 + 7 + (((0x33333333 * v122 + 7) >> 16) * ((0x33333333 * v122 + 7) >> 16))) << 16)) >> 16)
+ ((5

```

Junk Code

Finally we came to the key point:

```

v107 = ((v94 + *(&v94[*(v94 + 2 * v98 + v95[9])] + v95[7])))(&v701);
qword_140039408 = v107;
for ( j = *(*( __readgsqword(0x30u) + 96) + 24i64) + 16i64); ; j = *j )
{
    v109 = j[6];
    v110 = (v109 + *(v109 + v109[15] + 136));
    if ( v110 != v109 )
    {
        v111 = 0;
        v112 = v110[6];
        if ( v112 )
            break;
    }
}
LABEL_76:
;
}
while ( 1 )
{
    v113 = v109 + *(&v109[v111] + v110[8]);
    v114 = 0x811C9DC5;
    v115 = *v113;
    if ( *v113 )
    {
        do
        {
            v114 = 16777619 * (v114 ^ v115);
            v115 = ***v113;
        }
        while ( *v113 );
        if ( v114 == 0xF8F45725 )
            break;
    }
    if ( ++v111 >= v112 )
        goto LABEL_76;
}
v116 = ((v109 + *(&v109[*(v109 + 2 * v111 + v110[9])] + v110[7])))(v107, v706);
if ( v116 )
{
    v117 = -1i64;
    do
    {
        ++v117;
        while ( byte_140039300[v117] );
        v118 = v116(0i64, v117);
    }
}

```

We can see that v118 equals to the return value of v116() function,

```

loc_1400552B:                                ; CODE XREF: sub_140002E00+271C↑j
mov     ecx, [rdi+1Ch]
add     rcx, r9
mov     eax, [rdi+24h]
add     rax, r9
movzx   eax, word ptr [rax+r10*2]
mov     r8d, [rcx+rax*4]
add     r8, r9
lea     rdx, [rbp+0D40h+var_140]
mov     rcx, r13
call    r8
mov     r8, rax
test    rax, rax
mov     r14d, 5
jz      short loc_14005589
mov     rax, 0FFFFFFFFFFFFFFFFh
lea     rcx, byte_140039300
nop     dword ptr [rax+00h]

loc_14005570:                                ; CODE XREF: sub_140002E00+2777↓j
inc     rax
cmp     byte ptr [rcx+rax], 0
jnz     short loc_14005570
mov     rdx, rax
xor     ecx, ecx
call    r8                                ; v116
movsxd  rdx, eax

```

From the disassemble code above, we can see that it's the function call from export table, so we decided to debug it for clear (we set an bp at statement `call r8`), but we found the control flow can't come here, it exited instead, then we found the following condition must be satisfied for running to our bp by cross-reference analysis:

```

5 if ( sub_140002AC0(byte_140039300) == 0x4F8075587499C0FFi64 )
6 {
7     LODWORD(v46) = 31;
8     LODWORD(v47) = 0;
9 }


```

```

1 signed __int64 __fastcall sub_140002AC0(char *a1)
2 {
3     char v1; // dl
4     signed __int64 result; // rax
5     signed __int64 v3; // rax
6
7     v1 = *a1;
8     for ( result = 0xcBF29CE48422325i64; *a1; result = 0x100000001B3i64 * v3 )
9     {
10         ++a1;
11         v3 = result ^ v1;
12         v1 = *a1;
13     }
14     return result;
15 }

```

By Google the “magic number” `0xcBF29CE48422325` in function `sub_140002AC0()`, It’s the hash calculation algorithm `FNV-1(64bit)`:



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Fowler–Noll–Vo hash function - Wikipedia

https://en.wikipedia.org/wiki/Fowler–Noll–Vo_hash_function ▼

Fowler–Noll–Vo is a non-cryptographic hash function created by Glenn Fowler, Landon Curt ... The FNV_offset_basis is the 64-bit FNV offset basis value: 14695981039346656037 (in hex, **0xcbf29ce48422325**). The FNV_prime is the 64-bit ...

[The hash](#) · [FNV-1 hash](#) · [FNV prime](#) · [FNV hash parameters](#)

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FNV1a c++11 constexpr compile time hash functions, 32 and 64 bit ...

<https://gist.github.com/underscorediscovery/81308642d0325fd386237cfa3b44785c> ▼

uint64_t hash = **0xcbf29ce48422325**; uint64_t prime = 0x100000001b3; for(int i = 0; i < len; ++i) {
uint8_t value = data[i]; hash = hash ^ value; hash *= prime; }.

picoquic/fnv1a.h at master · private-octopus/picoquic · GitHub

<https://github.com/private-octopus/picoquic/blob/master/picoquic/fnv1a.h> ▼

14695981039346656037 (in hex, **0xcbf29ce48422325**). The prime is. 1099511628211 (in hex, 0x100000001b3; or as an expression $2^{40} + 2^8 + 0xb3$).

FNV-1 hash [edit]

The FNV-1 hash algorithm is as follows:[8][9]

```
hash = FNV_offset_basis
for each byte_of_data to be hashed
    hash = hash × FNV_prime
hash = hash XOR byte_of_data
return hash
```

In the above pseudocode, all variables are unsigned integers. All variables, except for *byte_of_data*, have the same number of bits as the FNV hash. The variable, *byte_of_data*, is an 8-bit unsigned integer.

As an example, consider the 64-bit FNV-1 hash:

- All variables, except for *byte_of_data* are 64-bit unsigned integers.
- The variable, *byte_of_data*, is an 8-bit unsigned integer.
- The *FNV_offset_basis* is the 64-bit *FNV offset basis* value: 14695981039346656037 (in hex `0xcbf29ce484222325`).
- The *FNV_prime* is the 64-bit *FNV prime* value: 1099511628211 (in hex `0x100000001b3`).
- The *multiply* returns the lower 64-bits of the product.
- The *XOR* is an 8-bit operation that modifies only the lower 8-bits of the hash value.
- The *hash* value returned is a 64-bit unsigned integer.

And the string var `byte_140039300` is confirmed to be the inputed SN string by debugging:

```
.text:00007FF6CA4A69EE ; -----
.text:00007FF6CA4A69EE
.text:00007FF6CA4A69EE
.text:00007FF6CA4A69EE 41 89 04 01 00 00 loc_7FF6CA4A69EE: ; CODE XREF: sub_7
mov     r9d, 104h ; MaxCount
.text:00007FF6CA4A69F4 4C 8B 42 08 mov     r8, [rdx+8] ; Src
.text:00007FF6CA4A69F8 41 8B 01 mov     edx, r9d ; SizeInBytes
.text:00007FF6CA4A69FB 48 8D 0D FE 28 03 00 lea     rcx, byte_7FF6CA4D9300 ; Dst
.text:00007FF6CA4A6A02 EB 55 7B 00 00 call    strncpy_s
RIP .text:00007FF6CA4A6A07 48 BA 00 00 00 06 00 00+mov     rdx, 600000000h ; SizeInBytes
.text:00007FF6CA4A6A11 48 8D 4D 77 lea     rcx, [rbp+77h] ; Dst
.text:00007FF6CA4A6A15 EB E6 C3 FF FF call    CheckSN
.text:00007FF6CA4A6A1A 33 C0 xor     eax, eax
.text:00007FF6CA4A6A1C
.text:00007FF6CA4A6A1C loc_7FF6CA4A6A1C: ; CODE XREF: sub_7
lea     r11, [rsp+0C8h+var_18]
.text:00007FF6CA4A6A24 49 8B 5B 28 mov     rbx, [r11+28h]
.text:00007FF6CA4A6A28 49 8B 73 38 mov     rsi, [r11+38h]
.text:00007FF6CA4A6A2C 49 8B E3 mov     rsp, r11

UNKNOWN|00007FF6CA4A69FB: sub_7FF6CA4A6A1A0+05B| (Synchronized with RIP)

Hex View-1
00007FF6CA4D9300 69 74 68 34 63 68 65 72 00 00 00 00 00 00 00 00 ith4cker .....
00007FF6CA4D9310 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....
00007FF6CA4D9320 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....
00007FF6CA4D9330 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....
00007FF6CA4D9340 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....
00007FF6CA4D9350 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....
```

Then the code execution flow goes into `sub_140002E00()`, it is confirmed to be the SN check function, all the check logics are in it, as I said before, there are also lots of junkcode in it for disturbing out analysis, it doesn't matter, just follow the inputed SN string and related vars, you will see the real codes, there are several check logics in the function, the first is the length of SN string must equal to 30 and there are at least 3 '9' in the SN string:

```

do
    ++v6;
while ( SN_Input[v6] );
LODWORD(v307) = v6;
HIDWORD(v307) = 5 - HIDWORD(v6);
v529 = (v6 + (HIDWORD(v6) << 32)) >> 12;
v530 = (v6 + (HIDWORD(v6) << 32)) >> 12;
v531 = (v6 + (HIDWORD(v6) << 32)) >> 12;
v532 = (v6 + (HIDWORD(v6) << 32)) >> 12;
v304 = v6 & 0xFFF;
v305 = v6 & 0xFFF;
v234 = v6 & 0xFFF;
v235 = v6 & 0xFFF;
v533 = (v6 + (HIDWORD(v6) << 32)) >> 12;
v534 = (v6 + (HIDWORD(v6) << 32)) >> 12;
v535 = (v6 + (HIDWORD(v6) << 32)) >> 12;
v236 = v6 & 0xFFF;
v237 = v6 & 0xFFF;
v238 = v6 & 0xFFF;
v7 = 4i64;
if ( (v6 & 0xFFF) + (v529 << 12) == 30 ) // length(SN) == 30
{

```

```

    if ( (v6 & 0xFFF) + (v529 << 12) == 30 ) // length(SN) == 30
    {
        v8 = 0;
        v9 = 0;
        if ( v6 )
        {
            v10 = SN_Input;
            do
            {
                v184 = 0;
                v11 = *v10;
                LOBYTE(v184) = v11;
                v12 = &v184;
                for ( i = 0xCBF29CE48422325i64; *v12; v11 = *v12 )
                {
                    i = 0x100000001B3i64 * (i ^ v11);
                    v12 = (v12 + 1);
                }
                count = v8 + 1;
                if ( i != 0xAF63B44C8601A894i64 ) // FNU-1(9) = 0xAF63B44C8601A894
                    count = v8;
                v8 = count;
                ++v9;
                ++v10;
            }
            while ( v9 < v6 );
            v1 = 32i64;
            if ( count >= 3 )
            {

```

The second check is that the first 9 chars are “KXCTF2018”:

```

LOBYTE(v183) = SN_Input[0];
if ( FNV_1_hash(&v183) == 0xAF64064C860233EAi64 )// 'K'
{
LOBYTE(v183) = SN_Input_1_;
if ( FNV_1_hash(&v183) == 0xAF64154C86024D67i64 )// 'X'
{
LOBYTE(v183) = SN_Input_2_;
if ( FNV_1_hash(&v183) == 0xAF63FE4C86022652i64 )// 'C'
{
LOBYTE(v183) = SN_Input_3_;
if ( FNV_1_hash(&v183) == 0xAF64094C86023903i64 )// 'T'
{
LOBYTE(v183) = SN_Input_4_;
if ( FNV_1_hash(&v183) == -5808522788293435079i64 )// 'F'
{
LOBYTE(v183) = SN_Input_5_;
if ( FNV_1_hash(&v183) == -5808606351177179115i64 )// '2'
{
LOBYTE(v183) = SN_Input_6_;
if ( FNV_1_hash(&v183) == -5808608550200435537i64 )// '0'
{
LOBYTE(v183) = SN_Input_7_;
if ( FNV_1_hash(&v183) == -5808609649712063748i64 )// '1'
{
LOBYTE(v183) = SN_Input_8_;
if ( FNV_1_hash(&v183) == -5808599754107409849i64 )// '8'
{
v42 = 0i64;
v37 = 0;
v43 = 0x50700i64;

```

The third is the FNV-1 hash of the whole SN string is **0x4F8075587499C0FF**(As we discussed earlier):

```

5 if ( FNV_1_hash(SN_Input) == 0x4F8075587499C0FFi64 )
7 {
3 LODWORD(v46) = 31;
9 LODWORD(v47) = 0;
9 }
1 else
2 {

```

After going through these checks, it came to the following string handling code block:

```

memset(&v701, 0, 260ui64);
v80 = strstr(SN_Input, "9"); // locate the first '9'
v81 = v80;
v701 = v80[1];
v702 = v80[2];
v703 = v80[3];
v704 = v80[4];
v705 = v80[5];
strncat_s(&v701, 260ui64, ".DLL", 260ui64); // XXXXX.DLL
v82 = strstr(v81 + 1, "9"); // locate the second '9'
v83 = v82;
*strstr(v82 + 1, "9") = 0; // locate the third '9'

```

We can know that the layout of the SN string is like following:

KXCTF20189XXXXX9YYYYYYYYYYYYY9

XXXXX represent the name of some a dll system directory, totally 5 bytes.

YYYYYYYYYYYYY represent the name of some a api in XXXXX.dll, totally 13 bytes.

After knowing the layout of the SN string, now you can debug it in IDA for a further analysis, you will find that it will retrieve the address of API `LoadLibrary` and `GetProcAddress` by PEB, and load the `XXXXX.dll` using `LoadLibrary`, after which call API `YYYYYYYYYYYY`, it using 32 bit `FNv_hash` to retrieve the info it wanted.

```

v93 = *((*(__readgsqword(0x30u) + 96) + 24i64) + 16i64); // get the module address by PEB
if ( (HIDWORD(v92) - abs(v92)) + (v92 << 32) )
{
    while ( 1 )
    {
        v94 = v93[6];
        v95 = (v94 + *(v94 + v94[15] + 136));
        if ( v95 != v94 )
        {
            v96 = 0;
            v97 = v95[6];
            if ( v97 )
                break;
        }
    }
LABEL_61:
    v93 = *v93;
    while ( 1 )
    {
        v98 = v96;
        v99 = v94 + *(v94[v96] + v95[8]);
        v100 = 0x811C9DC5;
        v101 = *v99;
        if ( *v99 )
        {
            do
            {
                v100 = 0x1000193 * (v100 ^ v101);
                v101 = ++v99;
            }
            while ( *v99 );
            if ( v100 == 0x53B2070F ) // LoadLibrary
                break;
        }
        if ( ++v96 >= v97 )
            goto LABEL_61;
    }
}

hModule_XXXXX = ((v94 + *(v94 + *(v94 + 2 * v98 + v95[9]) + v95[7]))(&XXXXX)); // LoadLibraryA("XXXXX.dll")
qword_14089408 = hModule_XXXXX;
for ( j = *((*(__readgsqword(0x30u) + 96) + 24i64) + 16i64); ; j = *j )
{
    v109 = j[6];
    v110 = (v109 + *(v109 + v109[15] + 136));
    if ( v110 != v109 )
    {
        v111 = 0;
        v112 = v110[6];
        if ( v112 )
            break;
    }
}
LABEL_76:
;
while ( 1 )
{
    v113 = v109 + *(v109[v111] + v110[8]);
    v114 = 0x811C9DC5;
    v115 = *v113;
    if ( *v113 )
    {
        do
        {
            v114 = 16777619 * (v114 ^ v115);
            v115 = ++v113;
        }
        while ( *v113 );
        if ( v114 == 0xF0F45725 ) // GetProcAddress
            break;
    }
    if ( ++v111 >= v112 )
        goto LABEL_76;
}
v116 = ((v109 + *(v109 + *(v109 + 2 * v111 + v110[9]) + v110[7]))(hModule_XXXXX, VVVVVVVVVVVVV)); // GetProcAddress(hModule_XXXXX, VVVVVVVVVVVVV)
if ( v116 )
{
    v117 = -1i64;
    do
    {
        ++v117;
        while ( SH_Input[v117] );
        v118 = v116(0i64, v117);
    }
}

```



```

xor     ebx, ebx

loc_7FF6CA4A5468:                ; CODE XREF: sub_7FF6CA4A2E00+2606↑j
                                   ; sub_7FF6CA4A2E00+2612↑j
mov     rsi, [rsi]
jmp     short loc_7FF6CA4A53F0
;-----
loc_7FF6CA4A546D:                ; CODE XREF: sub_7FF6CA4A2E00+25CC↑j
                                   ; sub_7FF6CA4A2E00+265C↑j
mov     eax, [rdi+24h]
add     rax, r9
mov     ecx, [rdi+1Ch]
movzx   eax, word ptr [rax+r10*2]
add     rcx, r9
mov     edx, [rcx+rax*4]
lea     rcx, [rbp+8D40h+var_250]
add     rdx, r9
;-----
; LoadLibraryA("NTDLL.DLL")
mov     r13, rax
mov     cs:qword_7FF6CA4D9408, rax
mov     rax, [rsi+30h]
mov     rcx, [rax+60h]
mov     rax, [rcx+18h]
mov     rsi, [rax+10h]
xor     r14d, r14d
nop

loc_7FF6CA4A54B0:                ; CODE XREF: sub_7FF6CA4A2E00+2729↑j
mov     r9, [rsi+30h]
movssd  rax, dword ptr [r9+3Ch]
mov     edi, [rax+r9+88h]
;-----
; sub_7FF6CA4A2E00+26AF (Synchronized with RIP)
;-----
; Hex View-1
;-----
; Stack view

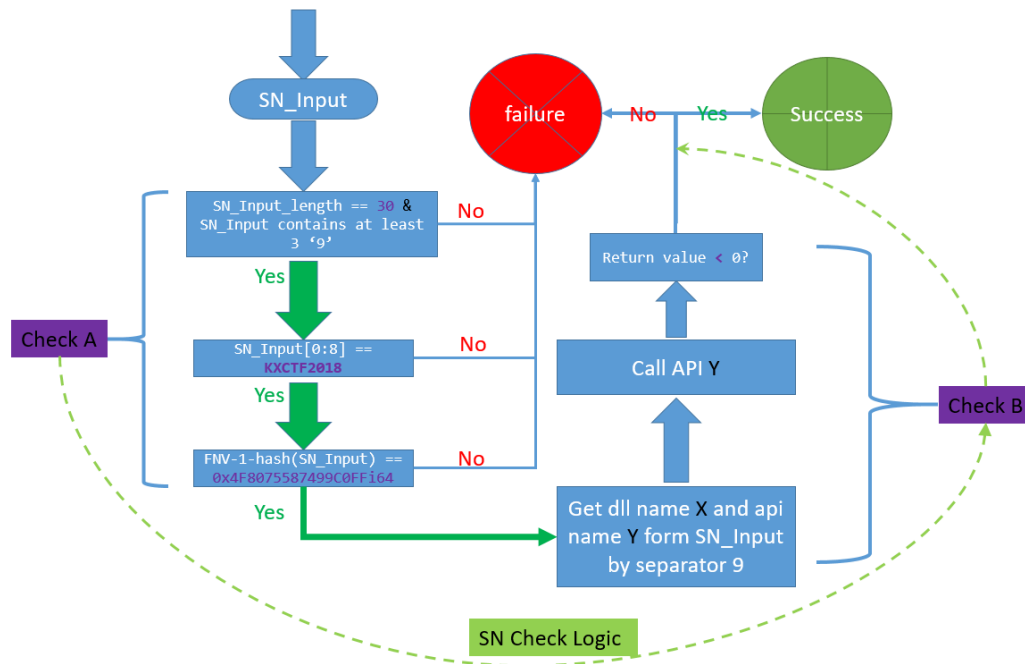
```

```

;-----
; sub_7FF6CA4A2E00+26D2↑j
mov     rsi, [rsi]
jmp     short loc_7FF6CA4A54B0
;-----
loc_7FF6CA4A552B:                ; CODE XREF: sub_7FF6CA4A2E00+271C↑j
mov     ecx, [rdi+1Ch]
add     rcx, r9
mov     eax, [rdi+24h]
add     rax, r9
movzx   eax, word ptr [rax+r10*2]
mov     r8d, [rcx+rax*4]
add     r8, r9
lea     rdx, [rbp+8D40h+var_140]
mov     rcx, r13
;-----
; GetProcAddress(NTDLL.DLL,YYYYYYYYYYY)
mov     r8, rax
test    rax, rax
mov     r14d, 5
;-----
; loc_7FF6CA4A5589
;-----
; loc_7FF6CA4A5570:                ; CODE XREF: sub_7FF6CA4A2E00+2777↑j
;-----
; sub_7FF6CA4A5565: sub_7FF6CA4A2E00+2765 (Synchronized with RIP)
;-----
; Hex View-1
;-----
; Stack view

```

So the whole check logic seems clear now as following flow chart shows:



The above flow chart is the key check logic,I divide the check logics into 2 groups **Check A and B** ,why? as there are 3 additional if-else judgment statements pointing to the **registration failure** code block:

```

- 1164);
v118 = 5 - (v175 & 0xFFFF0000) - v175 + (v175 - 5) - (5 - v175) - 7;
}
if ( v118 >= 0 )
{
  Registration_Failure:
  v321 = -dword_140038070;
  v322 = (dword_140038074 >> 16) + (dword_140038074 << 16);
  v323 = (dword_140038078 >> 16) + (dword_140038078 << 16);
  v324 = (dword_14003807C >> 16) + (dword_14003807C << 16);
  v325 = dword_140038080 - 5;
  sub_140001040(v0506, 60321, 19164);
  v180 = v0506;
  if ( v508 >= 0x10 )
  {
    v180 = v506;
    v181 = sub_140001340(dword_140039230, v180, v507); // registration failed
  }
}

```

Xref	Line	Column	Pseudocode line
o	708	0	goto Registration_Failure;
o	761	0	goto Registration_Failure;
o	698	0	goto Registration_Failure;
o	1447	0	v321 = -dword_140038070;

Line 3 of 4

```

v668 = v39;
v669 = v39;
v670 = v39;
v671 = v39;
v672 = v39;
v673 = v39;
v674 = 5i64;
v675 = (v39 - 5) & 0x7FFFFFFFFFFFFFFF164;
v676 = 7 * v675 & 0x7FFFFFFFFFFFFFFF164;
v677 = 5i64;
v678 = -(v676 - 5) & 0x7FFFFFFFFFFFFFFF164 & 0x7FFFFFFFFFFFFFFF164;
if ( !(5 * v678 & 0x7FFFFFFFFFFFFFFF164) + ((5 * v297 & 0x1FFFF) << 47)) )
goto Registration_Failure;
v679 = v45;
v680 = v45;
v681 = v45;

```

In dynamic debugging,the rip can run here,but I don't know what condition it need?when run here,the RAX and RDX will be zero:

```

C4B06 48 8B C0 07      imul    rax, 7
C4B0A 48 23 C0           and     rax, r8
C4B0D 48 89 85 D8 09 00 00  mov     [rbp+0040h+var_368], rax
C4B14 4C 89 85 E8 09 00 00  mov     [rbp+0040h+var_360], r14
C4B1B 48 83 C0 FB      add     rax, 0FFFFFFFFFFFFFFFBh
C4B1F 48 23 C0           and     rax, r8
C4B22 48 F7 D8      neg     rax
C4B25 48 23 C0           and     rax, r8
C4B28 48 89 85 E8 09 00 00  mov     [rbp+0040h+var_358], rax
C4B2F 48 23 C0           and     rax, r8
C4B32 48 8D 14 80      lea     rdx, [rax+rax*4]
C4B36 8B C1      mov     eax, ecx
C4B38 48 C1 E8 2F      chl     rax, 2Fh
C4B3C 48 23 D8      and     rdx, r8
C4B3F 48 83 C2      add     rax, rdx
RIP 48 8B C0 07      imul    rax, 7
C4B48 48 89 9D F8 09 00 00  mov     [rbp+0040h+var_348], rbx
C4B4F 48 89 9D 08 0A 00 00  mov     [rbp+0040h+var_340], rbx
C4B56 48 89 9D 08 0A 00 00  mov     [rbp+0040h+var_338], rbx
C4B5D 48 89 9D 18 0A 00 00  mov     [rbp+0040h+var_330], rbx
C4B64 48 89 9D 18 0A 00 00  mov     [rbp+0040h+var_328], rbx
C4B6B 48 89 9D 28 0A 00 00  mov     [rbp+0040h+var_320], rbx

```

Only when we input the right flag, the RAX and RDX will not be zero, and the jump will be false, so it seems that all check condition except the last condition (return value < 0) will jump here when the condition isn't satisfied, but I don't have a clear and accurate description for there are too many garbage instructions...

then we need to brute force the dll name (5 bytes) and api name (13 bytes) for the final flag using python script, the final crack output will look like as following:

```

Found SN_Input[0]:K
Found SN_Input[1]:X
Found SN_Input[2]:C
Found SN_Input[3]:T
Found SN_Input[4]:F
Found SN_Input[5]:2
Found SN_Input[6]:0
Found SN_Input[7]:1
Found SN_Input[8]:8
input[0:8]: KXCTF2018
Found SN_Input[9]:9
Found SN_Input[29]:9
Found API(e463da3c): GetModuleHandleA
Found API(f8f45725): GetProcAddress
Found API(53b2070f): LoadLibraryA
Congratulations to ITh4cker, Found flag: KXCTF20189NTDLL9DbgUiContinue9

```

0x1 Reference

0. <https://ctf.pediy.com/game-fight-66.htm>
1. https://en.wikipedia.org/wiki/Fowler%E2%80%93Noll%E2%80%93Vhe_hash_function#FNV-1_hash
2. <http://www.isthe.com/chongo/tech/comp/fnv/index.html#FNV-param>