

Leetcode: Record

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Problem 1

Problem 234 Palindrome Linked List

Given a singly linked list, determine if it is a palindrome.

Follow up:

Could you do it in $O(n)$ time and $O(1)$ space?

<https://leetcode.com/problems/palindrome-linked-list/>

Solution 1

Listing 1 shows solution with $O(n)$ runtime, $O(n)$ space.

The key idea is to use a stack storing elements from List, where popping elements out from a stack is exactly reversing the List.(The stack may use $O(n)$ space.)

Listing 1: Thought 1

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public boolean isPalindrome(ListNode head) {
        int len = getLength(head);
        if (len == 0 || len == 1){
            return true;
        }
        Stack<ListNode> stack = new Stack<ListNode>();
        int idx = 0;
        ListNode cur = head;
        boolean flag = true;
        while (cur != null){
            stack.push(cur);
            if (len % 2 == 0 && idx == len / 2 - 1){
                break;
            }
            if (len % 2 == 1 && idx == len / 2){
                stack.pop();
                break;
            }
            idx++;
            cur = cur.next;
        }
        cur = cur.next;
        while (cur != null){
            ListNode compare = stack.pop();
            if (compare.val != cur.val){
                return false;
            }
        }
    }
}
```

```
        }
        cur = cur.next;
    }
    return true;
40 }

    private int getLength(ListNode head){
        int len = 0;
        ListNode cur = head;
45        while (cur != null){
            len ++;
            cur = cur.next;
        }
        return len;
50    }
}
```

Solution 2

Listing 2 shows Solution 2 with $O(n)$ runtime, $O(n)$ space.

Instead of using stack, this solution finds element in the middle of whole List, reverse the second half and compare the elements with those in the first half.

Some key ideas:

1. The key idea of getting the middle element is to keep two iteration pointers, one with 1 step while another with 2 steps every time they iterate.
2. Take care of how to reverse a Linked-list in-place.

Listing 2: Thought2

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
5  *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
10    public boolean isPalindrome(ListNode head) {
        if (head == null || head.next == null){
            return true;
        }
        ListNode cur = head;
15        // get middle
        ListNode mid = getMid(head);
        // reverse from mid to end
        mid = reverse(mid);
        // check equations
20        while (mid != null){
            if (mid.val != cur.val){
                return false;
            }
        }
    }
}
```

```
        }
        mid = mid.next;
25         cur = cur.next;
    }
    return true;
}

30 public ListNode getMid(ListNode head){
    ListNode slow = head;
    ListNode fast = head;
    while (fast.next != null && fast.next.next != null){
        slow = slow.next;
35         fast = fast.next.next;
    }
    return slow.next;
}

40 public ListNode reverse(ListNode head){
    if (head.next == null){
        return head;
    }
    ListNode p,q,r;
45     p = head;
    q = p.next;
    r = q.next;
    p.next = null;
    while (q.next != null){
50         q.next = p;
        p = q;
        q = r;
        r = q.next;
    }
55     q.next = p;
    return q;
}

}
```

Problem 2

Problem 35 Search Insert Position

Given a sorted array and a target value, return the index if the target is found.

If not, return the index where it would be if it were inserted in order.

You may assume no duplicates in the array.

Here are few examples.

[1,3,5,6], 5 -> 2

[1,3,5,6], 2 -> 1

[1,3,5,6], 7 -> 4

[1,3,5,6], 0 -> 0

<https://leetcode.com/problems/search-insert-position/>

Solution

Basic Binary Search algorithm

Listing 3: Solution

```
public class Solution {
    public int searchInsert(int[] nums, int target) {
        if (target <= nums[0]){
            return 0;
        }
        if (target > nums[nums.length - 1]){
            return nums.length;
        }
        int l = 0 , r = nums.length - 1;
        while (l < r){
            int mid = (l + r) / 2;
            if (target == nums[mid]){
                return mid;
            }
            if (mid == l && target != nums[mid]){
                return mid + 1;
            }
            if (target > nums[mid]){
                l = mid;
            }
            else{
                r = mid;
            }
        }
        return -1;
    }
}
```


Problem 3

Problem 121 Best Time to Buy and Sell Stock

Say you have an array for which the i th element is the price of a given stock on day i .

If you were only permitted to complete at most one transaction (ie, buy one and sell one share of the stock), design an algorithm to find the maximum profit.

<https://leetcode.com/problems/best-time-to-buy-and-sell-stock/>

Solution 1

Brute force algorithm, list all possibilities, $O(n^2)$ runtime, undoubtedly Time Limit Exceed.

Listing 4: Solution 1

```
public class Solution {  
    // TLE  
    public int maxProfit(int[] prices) {  
        int max = 0;  
        for (int i = 0 ; i < prices.length ; i++){  
            for (int j = i + 1 ; j < prices.length ; j++){  
                if (prices[j] - prices[i] > max){  
                    max = prices[j] - prices[i];  
                }  
            }  
        }  
        return max;  
    }  
}
```

Solution 2

Maintain a minimum value until just before current index, so that we may get a better result when selling current stock while minimum price bought, $O(n)$ runtime.

Listing 5: Solution 2

```
public class Solution {  
    public int maxProfit(int[] prices) {  
        int max = 0;  
        if (prices.length == 0){  
            return 0;  
        }  
        int curMin = prices[0];  
        for (int i = 1 ; i < prices.length ; i++){  
            int profit = prices[i] - curMin;  
            if (profit > max){  
                max = profit;  
            }  
            if (prices[i] < curMin){  
                curMin = prices[i];  
            }  
        }  
    }  
}
```

```
        }  
        return max;  
    }  
}
```

Problem 4

Problem 122 Best Time to Buy and Sell Stock II

Say you have an array for which the i th element is the price of a given stock on day i .

Design an algorithm to find the maximum profit.

You may complete as many transactions as you like (ie, buy one and sell one share of the stock multiple times).

However, you may not engage in multiple transactions at the same time (ie, you must sell the stock before you buy again).

<https://leetcode.com/problems/best-time-to-buy-and-sell-stock-ii/>

Solution

Don't hesitate to use Greedy algorithm to solve the problem.

Listing 6: Solution

```
public class Solution {  
    public int maxProfit(int[] prices) {  
        int sum = 0;  
        for (int i = 1 ; i < prices.length ; i++){  
            if (prices[i] > prices[i - 1]){  
                sum += prices[i] - prices[i - 1];  
            }  
        }  
        return sum;  
    }  
}
```

Problem 5

Problem 123 Best Time to Buy and Sell Stock III

Say you have an array for which the i th element is the price of a given stock on day i .

Design an algorithm to find the maximum profit. You may complete at most two transactions.

Note:

You may not engage in multiple transactions at the same time (ie, you must sell the stock before you buy again).

<https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iii/>

Solution

This solution here is a variant of Problem “Best Time to Buy and Sell Stock”. We use each day as split point, respectively calculating maximum profits before and after the i^{th} day, then compare all of them to get the real maximum profit.

BTW, when we look at the 4th problem of this series, the count of transaction is extended to k , so that we can just substitute 2 into k to get the result of this problem. The solution will be demonstrated in next Problem.

However, simply apply such process will get TLE. We need to do some pre-processing. We may calculate maximum profit of $[0,1],[0,2],\dots$ and $[n-1,n],[n-2,n]$ so that we can reduce the complexity to $O(n)$

Listing 7: Solution

```
public class Solution {
    public int maxProfit(int[] prices) {
        if (prices.length == 0){
            return 0;
        }
        int maxSum = 0;
        int[] maxProfit = new int[prices.length];
        int[] maxProfitRev = new int[prices.length];

        int max = 0;
        int min = prices[0];
        for (int i = 1 ; i < prices.length ; i++){
            int diff = prices[i] - min;
            max = Math.max(diff, max);
            if (prices[i] < min){
                min = prices[i];
            }
            maxProfit[i] = max;
        }
        max = 0;
        int high = prices[prices.length - 1];
        for (int i = prices.length - 2 ; i >= 0 ; i--){
            int diff = high - prices[i];
            max = Math.max(max, diff);
            if (prices[i] > high){
                high = prices[i];
            }
        }
    }
}
```

```
        }
        maxProfitRev[i] = max;
    }
30  for (int i = 0 ; i < prices.length ; i++){
        maxSum = Math.max(maxSum, maxProfit[i] + maxProfitRev[i]);
    }
    return maxSum;
35 }
```

Problem 6

Problem 188 Best Time to Buy and Sell Stock IV

Say you have an array for which the i th element is the price of a given stock on day i .

Design an algorithm to find the maximum profit. You may complete at most k transactions.

Note:

You may not engage in multiple transactions at the same time (ie, you must sell the stock before you buy again).

<https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iv/>

Solution

It's not a conventional way of using DP. Detailed Description can be seen below.

题目的关键是要写出下面的动态转移方程:

```
local[i][j]=max(global[i-1][j-1]+max(diff,0),local[i-1][j]+diff),
global[i][j]=max(local[i][j],global[i-1][j]),
```

而这个方程的具体理解是这样的。

首先 global 比较简单，不过是不断地和已经计算出的 local 进行比较，把大的保存在 global 中。

然后看 local 关键是要理解 local 的定义，local[i][j] 表示，前 i 天进行了 j 次交易，并且第 i 天进行了第 j 次交易的最大利润，所以 local[i][j] 中必然有一次交易，也就是当近一次交易，发生在第 i 天。local 由两个部分的比较完成。

第一部分是， $global[i-1][j-1]+max(diff,0)$ ，表示的就是，前面把之前的 $j-1$ 次交易，放在之前的 $i-1$ 天，然后让第 i 天来进行第 j 次交易，那么加入此时 $diff(price[i] - price[i-1])$ 大于零，那么正好可以借助这次交易的机会增长里利润(利润=diff)，否则的话，如果 diff 小于零，那就在第 i 天当天进行一次买卖，凑一次交易的次数，但是产生利润为 0。

第二部分是， $local[i-1][j]+diff$ ，这里的 $local[i-1][j]$ 表示的是，前面 j 次交易在第 $i-1$ 天就已经完成了，可是因为说了 $local[i][j]$ 一定要表达在第 i 天完成了 j 次交易的最大利润，所以就强制使得交易在第 i 天发生，为了实现这一点，只需要在 $local[i-1][j]$ 的基础上，加上 $diff (= price[i] - price[i-1])$ 就可以了。如果 $diff < 0$ 那也没有办法，因为必须满足 local 的定义。接下来算 global 的时候，总会保证取得一个更大的值。

关于正道题目的全文的分析，Code Ganker 的还挺不错。下面全文摘抄作者的分析：

"这道题是 Best Time to Buy and Sell Stock 的扩展，现在我们可以进行两次交易。我们仍然使用动态规划来完成，事实上可以解决非常通用的情况，也就是最多进行 k 次交易的情况。

这里我们先解释最多可以进行 k 次交易的算法，然后最多进行两次我们只需要把 k 取成 2 即可。我们还是使用“局部最优和全局最优解法”。我们维护两种量，一个是当前到达第 i 天可以最多进行 j 次交易，最好的利润是多少 ($global[i][j]$)，另一个是前到达第 i 天，最多可进行 j 次交易，并且最后一次交易在当天卖出的最好的利润是多少 ($local[i][j]$)。下面我们来看递推式。

全局的比较简单，

```
global[i][j]=max(local[i][j],global[i-1][j]),
```

也就是去当前局部最好的，和过往全局最好的中大的那个（因为最后一次交易如果包含当前天一定在局部最好的里面，否则一定在过往全局最优的里面）。

对于局部变量的维护，递推式是

```
local[i][j]=max(global[i-1][j-1]+max(diff,0),local[i-1][j]+diff),
```

也就是看两个量，第一个是全局到 $i-1$ 天进行 $j-1$ 次交易，然后加上今天的交易，如果今天是赚钱的话（也就是前面只要 $j-1$ 次交易，最后一次交易取当前天），第二个量则是取 local 第 $i-1$ 天 j 次交易，然后加上今天的差值（这里因为 $local[i-1][j]$ 比包含第 $i-1$ 天卖出的交易，所以现在变成第 i 天卖出，并不会增加交易次数，而且这里无论 diff 是不是大于 0 都一定要加上，因为否则就不满足 $local[i][j]$ 必须在最后一天卖出的条件了）。

上面的算法中对于天数需要一次扫描，而每次要对交易次数进行递推式求解，所以时间复杂度是 $O(n*k)$ ，如果是最多进行两次交易，那么复杂度还是 $O(n)$ 。空间上只需要维护当天数据皆可，所以是 $O(k)$ ，当 $k=2$ ，则是 $O(1)$ 。”

(Honestly, I don't quite understand the solution here.)

Listing 8: Solution

```
public class Solution {
    public int maxProfit(int k, int[] prices) {
        if (prices.length == 0){
            return 0;
        }
        // for those special input data
        if (k > prices.length){
            int sum = 0;
            for (int i = 1 ; i < prices.length ; i++){
                if (prices[i] > prices[i - 1]){
                    sum += prices[i] - prices[i - 1];
                }
            }
        }
    }
}
```

```
15         return sum;
    }
    else{
        int [][] total = new int[prices.length][k + 1];
        int [][] part = new int[prices.length][k + 1];
        for (int i = 1 ; i < prices.length ; i++){
            int diff = prices[i] - prices[i - 1];
            for (int j = k ; j >= 1 ; j--){
                part[i][j] = Math.max(total[i - 1][j - 1]
                    + Math.max(0, diff), part[i - 1][j] + diff);
                total[i][j] = Math.max(part[i][j], total[i - 1][j]);
            }
        }

        return total[prices.length - 1][k];
    }
}
}
```

Problem 7

Problem 309 Best Time to Buy and Sell Stock with Cooldown

Say you have an array for which the i th element is the price of a given stock on day i .

Design an algorithm to find the maximum profit.

You may complete as many transactions as you like (ie, buy one and sell one share of the stock multiple times) with the following restrictions:

You may not engage in multiple transactions at the same time (ie, you must sell the stock before you buy again).
After you sell your stock, you cannot buy stock on next day. (ie, cooldown 1 day)

Example:

```
prices = [1, 2, 3, 0, 2]
maxProfit = 3
transactions = [buy, sell, cooldown, buy, sell]
```

<https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-cooldown/>

Solution

Another problem I didn't understand.....

分析

因为当前日期买卖股票会受到之前日期买卖股票行为的影响，首先考虑到用DP解决。

这道题比较麻烦的是有个cooldown的限制，其实本质也就是买与卖之间的限制。对于某一天，股票有三种状态: buy, sell, cooldown, sell与cooldown我们可以合并成一种状态，因为手里最终都没股票，最终需要的结果是sell，即手里股票卖了获得最大利润。所以我们可以用两个DP数组分别记录当前持股跟未持股的状态。然后根据题目中的限制条件，理清两个DP数组的表达式。

对于当天最终未持股的状态，最终最大利润有两种可能，一是今天没动作跟昨天未持股状态一样，二是昨天持股了，今天卖了。所以我们只要取这两者之间最大值即可，表达式如下：

```
sellDp[i] = Math.max(sellDp[i - 1], buyDp[i - 1] + prices[i]);
```

对于当天最终持股的状态，最终最大利润有两种可能，一是今天没动作跟昨天持股状态一样，二是前天还没持股，今天买了股票，这里是因为cooldown的原因，所以今天买股票要追溯到前天的状态。我们只要取这两者之间最大值即可，表达式如下：

```
buyDp[i] = Math.max(buyDp[i - 1], sellDp[i - 2] - prices[i]);
```

最终我们要求的结果是

```
sellDp[n - 1] 表示最后一天结束时手里没股票时的累积最大利润
```

当然，这里空间复杂度是可以降到 $O(1)$ 的，具体见第二种代码实现。

Listing 9: Solution

```
public class Solution {
    public int maxProfit(int[] prices) {
        if (prices.length == 0){
            return 0;
        }
    }
}
```



```
5      }
      int sell[] = new int[prices.length];
      int buy[] = new int[prices.length];

      sell[0] = 0;
10     buy[0] = -prices[0];

      for (int i = 1 ; i < prices.length ; i++){
          sell[i] = Math.max(sell[i - 1], buy[i - 1] + prices[i]);
          if (i == 1){
15             buy[i] = Math.max(buy[i - 1], -prices[i]);
          }
          else{
              buy[i] = Math.max(buy[i - 1], sell[i - 2] - prices[i]);
          }
20     }

    return sell[prices.length - 1];
}
}
```

Problem 8

Problem 9 Palindrome Number

Given a string, determine if it is a palindrome, considering only alphanumeric characters and ignoring cases.

For example,

"A man, a plan, a canal: Panama" is a palindrome.

"race a car" is not a palindrome.

Note:

Have you consider that the string might be empty? This is a good question to ask during an interview.

For the purpose of this problem, we define empty string as valid palindrome.

<https://leetcode.com/problems/valid-palindrome/>

Solution

Though it's just a easy string-handling problem with not many edge cases, I made several mistakes:

- Ignoring the case problems of letters, should lowercase characters in sentence.
- Forgot to check boundaries, especially when all characters are non-alphanumeric, there will be an `IndexOutOfBoundsException`.

Listing 10: Solution

```
public class Solution {
    public boolean isPalindrome(String s) {
        if (s == null || s.length() == 0){
            return true;
        }
        int i = 0, j = s.length() - 1;
        while (i < j){
            while (i < s.length() && !Character.isLetter(s.charAt(i))
                && !Character.isDigit(s.charAt(i))){
                i ++;
            }
            while (j > 0 && !Character.isLetter(s.charAt(j))
                && !Character.isDigit(s.charAt(j))){
                j --;
            }
            if (i >= j){
                return true;
            }

            if (Character.toLowerCase(s.charAt(i)) !=
                Character.toLowerCase(s.charAt(j))){
                return false;
            }
            i ++;
            j --;
        }
    }
}
```

```
        return true;  
    }  
}
```

Problem 9

Problem 206 Reverse Linked List

Reverse a singly linked list.

Hint:

A linked list can be reversed either iteratively or recursively.
Could you implement both?

<https://leetcode.com/problems/reverse-linked-list/>

Solution 1

Non-recursive solution, basically maintains 3 pointers, 'cur', 'next', next node of 'next', see details in Listing 11

Listing 11: Solution 1

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode reverseList(ListNode head) {
        if (head == null || head.next == null){
            return head;
        }
        ListNode cur = head;
        ListNode next = cur.next;
        cur.next = null;

        while (next.next != null){
            ListNode tmp = next.next;
            next.next = cur;
            cur = next;
            next = tmp;
        }
        next.next = cur;
        return next;
    }
}
```

Solution 2

Recursive solution.

Listing 12: Solution 2

```
/**
```

```

5  * Definition for singly-linked list.
   * public class ListNode {
   *     int val;
   *     ListNode next;
   *     ListNode(int x) { val = x; }
   * }
   */
10 public class Solution {
   public ListNode reverseList(ListNode head) {
       if (head == null || head.next == null){
           return head;
       }
       ListNode nextNode = head.next;
15  head.next = null;
       ListNode retNode = reverseList(nextNode);
       nextNode.next = head;
       return retNode;
   }
20 }
```

Problem 10

92 Reverse Linked List II

Reverse a linked list from position m to n .
Do it in-place and in one-pass.

For example:

Given $1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow \text{NULL}$, $m = 2$ and $n = 4$,

return $1 \rightarrow 4 \rightarrow 3 \rightarrow 2 \rightarrow 5 \rightarrow \text{NULL}$.

Note:

Given m , n satisfy the following condition:

$1 \leq m \leq n \leq \text{length of list}$.

<https://leetcode.com/problems/reverse-linked-list-ii/>

Solution

Dummy node is a excellent aid when dealing with problems involving Linked-list.

The main idea of my solution is to firstly use dummy node to avoid talking about conditions when the reversion starts from the 1st element, which may need a lot of code to deal with empty head nodes; secondly iterate to the element where the reversion starts, record its previous node as the tail of 1st part of the whole 3 parts; reverse until the index reaches n ; reconnect the tails of each parts by assigning `.next` attribute.

Listing 13: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode reverseBetween(ListNode head, int m, int n) {
        if (n == m){
            return head;
        }
        int idx = 0;
        ListNode dummy = new ListNode(-1);
        dummy.next = head;
        ListNode cur = dummy;
        ListNode prev = cur;
        while (cur.next != null && idx < m) {
            prev = cur;
            cur = cur.next;
            idx++;
        }
        ListNode tail1 = prev;
        ListNode tail2 = cur;
        ListNode next = cur.next;
```

```
        while (cur.next != null && idx < n){
            ListNode tmp = next.next;
            next.next = cur;
30         cur = next;
            next = tmp;
            idx++;
        }
        tail2.next = next;
35         tail1.next = cur;
        return dummy.next;
    }
}
```

Problem 11

19 Remove Nth Node From End of List

Given a linked list, remove the nth node from the end of list and return its head.

For example,

Given linked list: 1->2->3->4->5, and n = 2.

After removing the second node from the end, the linked list becomes 1->2->3->5.

Note:

Given n will always be valid.

Try to do this in one pass.

<https://leetcode.com/problems/remove-nth-node-from-end-of-list/>

Solution

We can solve this problem by calculating its length first and get index of deleted node by length-n. However, this is not a solution in one-pass;

A tactful idea is that we first iterate one pointer n times, then iterate pointer describe above together with a pointer from head of the list, when the 1st point reaches the end of list, the 2nd point stays just at the node we want to remove. See 14

Listing 14: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode removeNthFromEnd(ListNode head, int n) {
        ListNode p1 = head;
        ListNode p2 = head;
        while (n-- > 0) {
            p1 = p1.next;
        }
        if (p1 == null) {
            return head.next;
        }
        ListNode prev = p2;
        while (p1 != null) {
            p1 = p1.next;
            prev = p2;
            p2 = p2.next;
        }
        prev.next = p2.next;
        return head;
    }
}
```


Problem 12

203 Remove Linked List Elements

Remove all elements from a linked list of integers that have value val.

Example

Given: 1 --> 2 --> 6 --> 3 --> 4 --> 5 --> 6, val = 6

Return: 1 --> 2 --> 3 --> 4 --> 5

<https://leetcode.com/problems/remove-linked-list-elements/>

Solution

An easy problem, just be aware of some edge cases such as inputting an empty list or a list full of values to be removed.

Listing 15: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode removeElements(ListNode head, int val) {
        while (head != null && head.val == val){
            head = head.next;
        }
        if (head == null || head.next == null)
            return head;
        ListNode cur = head.next;
        ListNode prev = head;
        while (cur != null){
            while (cur != null && cur.val == val){
                prev.next = cur.next;
                cur = cur.next;
            }
            if (cur == null){
                return head;
            }
            prev = cur;
            cur = cur.next;
        }
        return head;
    }
}
```

Problem 13

83 Remove Duplicates from Sorted List

Given a sorted linked list,
delete all duplicates such that each element appear only once.

For example,

Given 1->1->2, return 1->2.

Given 1->1->2->3->3, return 1->2->3.

<https://leetcode.com/problems/remove-duplicates-from-sorted-list/>

Solution

An easy problem, since list is sorted, just check if value of neighbor nodes are equal.

Listing 16: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
5
public class Solution {
10
    public ListNode deleteDuplicates(ListNode head) {
        if (head == null || head.next == null){
            return head;
        }
        ListNode cur = head;
        ListNode next = cur.next;
15
        while (next != null){
            while (next != null && cur.val == next.val){
                cur.next = next.next;
                next = next.next;
20
            }
            if (next == null)
                return head;
            cur = next;
            next = next.next;
25
        }
        return head;
    }
}
```

Problem 14

82 Remove Duplicates from Sorted List II

Given a sorted linked list, delete all nodes that have duplicate numbers, leaving only distinct numbers from the original list.

For example,

Given 1->2->3->3->4->4->5, return 1->2->5.

Given 1->1->1->2->3, return 2->3.

<https://leetcode.com/problems/remove-duplicates-from-sorted-list-ii/>

Solution

My solution is to use 'prev' and 'cur' pointer to deal with such problem. 'prev' maintains lists which already be processed to have only distinct values; 'cur' is used to skip all adjacent elements whose values are equal.

There is another interesting train of thought, where distinct value means the value of element is not equal to its neighbors(both left side and right side.).

Listing 17: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode deleteDuplicates(ListNode head) {
        if (head == null || head.next == null){
            return head;
        }
        ListNode dummy = new ListNode(-1);
        dummy.next = head;
        ListNode former = dummy;
        ListNode cur = dummy.next;
        ListNode next = cur.next;
        while(next != null){
            if (cur.val == next.val){
                int val = cur.val;

                while (cur != null && cur.val == val){
                    cur = cur.next;
                }
                if (cur == null){
                    former.next = null;
                    return dummy.next;
                }
                else {
                    next = cur.next;
                }
            }
        }
    }
}
```

```
35         else{
           former.next = cur;
           former = cur;

           cur = cur.next;
           next = cur.next;
40     }
    }
    former.next = cur;
    return dummy.next;
45 }
```

Problem 15

328 Odd Even Linked List

Given a singly linked list,
group all odd nodes together followed by the even nodes.
Please note here we are talking about
the node number and not the value in the nodes.

You should try to do it in place. The program should run in
 $O(1)$ space complexity and $O(\text{nodes})$ time complexity.

Example:

Given 1->2->3->4->5->NULL,
return 1->3->5->2->4->NULL.

Note:

The relative order inside both the even and odd groups should remain as it was
in the input.

The first node is considered odd, the second node even and so on ...

<https://leetcode.com/problems/odd-even-linked-list/>

Solution

An easy problem, just to be aware of conditions when length of list is odd and even.

Listing 18: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
5  *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
10  public ListNode oddEvenList(ListNode head) {
    if (head == null || head.next == null){
        return head;
    }
    ListNode cur = head;
15  ListNode next = cur.next;
    ListNode evenHead = next;
    while (next != null){
        cur.next = next.next;
        if (cur.next == null){
20  cur.next = evenHead;
            return head;
        }
        next.next = cur.next.next;
        cur = cur.next;
25  next = next.next;
    }
    cur.next = evenHead;
    return head;
}
```

30

```
}  
}
```

Problem 16

237 Delete Node in a Linked List

Write a function to delete a node (except the tail) in a singly linked list, given only access to that node.

Supposed the linked list is 1 -> 2 -> 3 -> 4 and you are given the third node with value 3, the linked list should become 1 -> 2 -> 4 after calling your function.

<https://leetcode.com/problems/delete-node-in-a-linked-list/>

Solution

An easy but weird problem, must use deep copy to finish the task, since OJ System will check the value of node with current index.

Listing 19: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public void deleteNode(ListNode node) {
        //input check
        if(node==null) return;
        node.val = node.next.val;
        node.next = node.next.next;
    }
}
```

Problem 17

160 Intersection of Two Linked Lists

Write a program to find the node at which the intersection of two singly linked lists begins.

For example, the following two linked lists:

```
A:          a1 -> a2
              ->
              c1 -> c2 -> c3
              ->
B:    b1 -> b2 -> b3
begin to intersect at node c1.
```

Notes:

If the two linked lists have no intersection at all, return null.
The linked lists must retain their original structure after the function returns.
You may assume there are no cycles anywhere in the entire linked structure.
Your code should preferably run in $O(n)$ time and use only $O(1)$ memory.

<https://leetcode.com/problems/intersection-of-two-linked-lists/>

Solution

Solution of this problem needs tactics.

Brute force solution is definitely not good enough. The idea of $O(n)$ runtime and $O(1)$ space solution is similar to the idea in *Problem 19 Remove Nth Node From End of List*, which needs to iterate the longer list several times (equals to the difference between two length), then iterate both lists since they are of same length currently.

Listing 20: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) {
 *         val = x;
 *         next = null;
 *     }
 * }
 */
public class Solution {
    public ListNode getIntersectionNode(ListNode headA, ListNode headB) {
        if (headA == null || headB == null){
            return null;
        }
        int lenA = 1 ; int lenB = 1;
        ListNode curA = headA;
```



```
20     ListNode curB = headB;
    while (curA.next != null){
        lenA ++;
        curA = curA.next;
    }
25     while (curB.next != null){
        lenB ++;
        curB = curB.next;
    }
    if (curA.val != curB.val){
        return null;
30    }
    ListNode longer, shorter;
    int diff = Math.abs(lenA - lenB);
    if (lenA > lenB){
        longer = headA;
35        shorter = headB;
    }
    else{
        longer = headB;
        shorter = headA;
40    }
    boolean found = false;
    ListNode inter = null;
    while (diff -- > 0){
        longer = longer.next;
45    }
    while(longer != null){
        if (!found && longer.val == shorter.val){
            found = true;
            inter = longer;
50        }
        if (longer.val != shorter.val){
            found = false;
            inter = null;
        }
55        longer = longer.next;
        shorter = shorter.next;
    }
    return inter;
60 }
```

Problem 18

21 Merge Two Sorted Lists

Merge two sorted linked lists and return it as a new list.

The new list should be made by splicing together the nodes of the first two lists.

<https://leetcode.com/problems/merge-two-sorted-lists/>

Solution

This problem indicates the importance of dummy node in solving linked list problems.

Listing 21: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
5
public class Solution {
10
    public ListNode mergeTwoLists(ListNode l1, ListNode l2) {
        ListNode dummy = new ListNode(-1);
        ListNode cur = dummy;
        if (l1 == null){
            return l2;
15
        }
        if (l2 == null){
            return l1;
        }
        while (l1 != null && l2 != null){
20
            if (l1.val <= l2.val){
                cur.next = l1;
                l1 = l1.next;
            }
            else{
25
                cur.next = l2;
                l2 = l2.next;
            }
            cur = cur.next;
        }
        if (l1 == null){
30
            cur.next = l2;
        }
        if (l2 == null){
            cur.next = l1;
35
        }
        return dummy.next;
    }
}
```

Problem 19

23 Merge k Sorted Lists

Merge k sorted linked lists and return it as one sorted list.
Analyze and describe its complexity.

<https://leetcode.com/problems/merge-k-sorted-lists/>

Solution 1 & 2

Solution 1 is Brute force. With $O(nk^2)$ runtime and $O(1)$ space. (nk^2 comes from $2n + 3n + \dots + kn$). We may skip Solution 1 here.

Solution 2 makes use of heap. Keep a heap with min root, pop one, link it into result, push its next node.

Since Heap has a complexity of $\log(k)$ runtime with nk elements, the whole process has a $O(kn \log k)$ runtime and $O(k)$ space complexity.

Listing 22: Solution

```

/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    private static final Comparator<ListNode> listComparator
    = new Comparator<ListNode>(){
        @Override
        public int compare(ListNode x, ListNode y){
            return x.val - y.val;
        }
    };

    public ListNode mergeKLists(ListNode[] lists) {
        if (lists.length == 0){
            return null;
        }
        Queue<ListNode> queue
    = new PriorityQueue<ListNode>(lists.length, listComparator);
        for (ListNode list : lists){
            if (list != null){
                queue.offer(list);
            }
        }

        ListNode dummy = new ListNode(-1);
        ListNode head = dummy;
        while (!queue.isEmpty()){
            ListNode curNode = queue.poll();
            head.next = curNode;
            head = head.next;
        }
    }
}

```

```

        if (curNode.next != null){
            queue.offer(curNode.next);
        }
    }
    return dummy.next;
}
}

```

Solution 3

Solution 3 is still of $O(kn \log k)$ runtime but with $O(1)$ space.

The number of lists reduces from:

$$k \rightarrow k/2 \rightarrow k/4 \rightarrow \dots \rightarrow 2 \rightarrow 1$$

When the size of lists increases from :

$$n \rightarrow 2n \rightarrow 4n \rightarrow \dots \rightarrow 2^{\log k} n$$

So the runtime complexity is:

$$kn + k/2 * 2n + \dots + 2^{\log k} n = nk \log k$$

Listing 23: Solution

```

/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode mergeTwoLists(ListNode list1, ListNode list2){
        ListNode dummy = new ListNode(-1);
        ListNode head = dummy;
        ListNode cur1 = list1;
        ListNode cur2 = list2;

        while (cur1 != null && cur2 != null){
            if (cur1.val < cur2.val){
                head.next = cur1;
                cur1 = cur1.next;
                head = head.next;
            }
            else{
                head.next = cur2;
                cur2 = cur2.next;
                head = head.next;
            }
        }

        while (cur1 != null){
            head.next = cur1;
            cur1 = cur1.next;
            head = head.next;
        }
    }
}

```

```
35         while (cur2 != null){
            head.next = cur2;
            cur2 = cur2.next;
            head = head.next;
        }

40     return dummy.next;
}

public ListNode mergeKLists(ListNode[] lists) {
    if (lists.length == 0){
45         return null;
    }
    int begin = 0 , end = lists.length - 1;

    while (end > 0){
50         begin = 0;
        while (begin < end){
            // two pointers
            lists[begin] = mergeTwoLists(lists[begin], lists[end]);
            begin ++;
55             end --;
        }
    }

    return lists[0];
60 }
}
```

Problem 20

143 Reorder List

Given a singly linked list L: L₀->L₁->...->L_{n-1}->L_n,
reorder it to: L₀->L_n->L₁->L_{n-1}->L₂->L_{n-2}->...

You must do this in-place without altering the nodes' values.

For example,

Given {1,2,3,4}, reorder it to {1,4,2,3}.

<https://leetcode.com/problems/reorder-list/>

Solution

The train of thought is somewhat straightforward, just find the middle element, reverse the latter half, and reconnect elements in first half and second half. But not coding in IDE to solve a problem of Medium difficulty is impossible to me currently...

Listing 24: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
10 public void reorderList(ListNode head) {
    if (head == null || head.next == null || head.next.next == null){
        return;
    }
    ListNode slow = head;
    ListNode fast = head;
15     ListNode prev = slow;
    while (fast != null && fast.next != null){
        prev = slow;
        slow = slow.next;
20         fast = fast.next.next;
    }
    prev.next = null;
    // reverse latter half
    ListNode cur = slow;
    ListNode next = cur.next;
25     cur.next = null;
    while (next != null){
        ListNode tmp = next.next;
        next.next = cur;
30         cur = next;
        next = tmp;
    }
    // cur is the reversed list
    ListNode p = head;
```

```
35     ListNode q = cur;
    while (p.next != null){
        ListNode r = p.next;
        ListNode s = q.next;
        p.next = q;
40     q.next = r;
        p = r;
        q = s;
    }
    p.next = q;
45 }
}
```

Problem 21

326 Power of Three

Given an integer, write a function to determine if it is a power of three.

Follow up:

Could you do it without using any loop / recursion?

<https://leetcode.com/problems/power-of-three/>

Solution

Using recursion or loop is extremely easy, but not using both is a challenge. Note that $\log_3(n) = \frac{\log(n)}{\log(3)}$ and it will return a integer if n is power of 3. Then we should just check if *div* and *Math.round(div)* has a difference less than 0.000000001 (If we just check equation we may find 243 returns false because of precision problem in Java).

Listing 25: Solution

```
public class Solution {  
    public boolean isPowerOfThree(int n) {  
        double div = Math.log(n) / Math.log(3);  
        double diff = Math.abs(div - Math.round(div));  
5         if (diff <= 0.000000001){  
            return true;  
        }else{  
            return false;  
        }  
10    }  
}
```


Problem 22

231 Power of Two

Given an integer, write a function to determine if it is a power of two.

Follow up:

Could you do it without using any loop / recursion?

<https://leetcode.com/problems/power-of-two/>

Solution 1

Solution 1 is nearly the same as solution in Pro 326, *Power of Two*, just change 3 into 2.

Listing 26: Solution 1

```
public class Solution {
    public boolean isPowerOfTwo(int n) {
        double div = Math.log(n) / Math.log(2);
        double diff = Math.abs(div - Math.round(div));
5      if (diff <= 0.0000000001)
            return true;
        else
            return false;
    }
10 }
```

Solution 2

This solution is interesting. Speaking of 2(or 4,16, etc, we can do sqrt calculation first to make problem familiar), we may use bit operators to solve it. Both checking n-1 is all 1 in Binary or in form of 1000000 in Binary are fine.

Listing 27: Solution 2

```
public class Solution {
    public boolean isPowerOfTwo(int n) {
        if (n <= 0)
            return false;
5      int cur = (int)n - 1;
        while (cur != 0){
            if ((cur & 1) != 1)
                return false;
            cur = (cur >> 1);
10     }
        return true;
    }
}
// OR JUST!!!
15 public class Solution {
    public boolean isPowerOfTwo(int n) {
        return (n != 0
                && n != Integer.MIN_VALUE
                && (n & (n - 1)) == 0);
20     }
}
```

Problem 23

342 Power of Four

Given an integer (signed 32 bits), write a function to check whether it is a power of 4.

Example:

Given num = 16, return true. Given num = 5, return false.

Follow up: Could you solve it without loops/recursion?

<https://leetcode.com/problems/power-of-four/>

Solution 1

First identify whether it is power of 2 (simply using $n \& (n-1) == 0$), then if it is, the form of the number must be like 10...0, check whether there are even number of 0 behind the heading 1.

However, I'm still using a loop.

Listing 28: Solution 1

```
public class Solution {  
    public boolean isPowerOfFour(int num) {  
        if (((num & (num - 1)) != 0) || num == 0){  
            return false;  
        }  
        while ((num & 3) == 0){  
            num >>= 2;  
        }  
        return ((num & 3) == 1);  
    }  
}
```

Solution 2

Still, check if it is power of 2. Then use 0x55555555 (10101010101010101010101010101010 in binary) to check if the single 1 only appears at odd position.

Listing 29: Solution 2

```
public class Solution {  
    public boolean isPowerOfFour(int num) {  
        return num > 0  
            && (num & (num - 1)) == 0 && (num & 0x55555555) != 0;  
    }  
}
```

Problem 24

191 Number of 1 Bits

Write a function that takes an unsigned integer and returns the number of '1' bits it has (also known as the Hamming weight).

For example, the 32-bit integer '11' has binary representation 00000000000000000000000001011, so the function should return 3.

<https://leetcode.com/problems/number-of-1-bits/>

Solution 1

To be accurate, the 2 solutions are 2 ways of thinking rather than ways of solving problems. The first one uses \gg (unsigned bit operator) and check if n is still larger than 0. *DO NOT USE ' \ggg '! It will overflow for 2^{32} !*

Listing 30: Solution 1

```
public class Solution {  
    // you need to treat n as an unsigned value  
    public int hammingWeight(int n) {  
        int count = 0;  
        while(n != 0){  
            count += (n & 1);  
            n >>= 1;  
        }  
        return count;  
    }  
}
```

Solution 2

The 2nd solution make use of the fact that the integer is at most 32 bits. It iterates the number for 32 bits to see whether 1 exists.

Listing 31: Solution 2

```
public class Solution {  
    // you need to treat n as an unsigned value  
    public int hammingWeight(int n) {  
        int count = 0;  
        for(int i = 0; i < 32; i++){  
            count += (n & 1);  
            n >>= 1;  
        }  
        return count;  
    }  
}
```

What's more, using $n \& (n-1)$ also does excellent work. Mind those boundaries.

Problem 25

202 Happy Number

Write an algorithm to determine if a number is "happy".

A happy number is a number defined by the following process:

Starting with any positive integer,

replace the number by the sum of the squares of its digits,

and repeat the process until the number equals 1 (where it will stay),

or it loops endlessly in a cycle which does not include 1.

Those numbers for which this process ends in 1 are happy numbers.

Example: 19 is a happy number

$$1*1 + 9*9 = 82$$

$$8*8 + 2*2 = 68$$

$$6*6 + 8*8 = 100$$

$$1*1 + 0*0 + 0*0 = 1$$

<https://leetcode.com/problems/happy-number/>

Solution

A set is enough for this problem.

Listing 32: Solution

```
public class Solution {
    public boolean isHappy(int n){
        Set<Integer> set = new HashSet<Integer>();
        int sum = 0;
5       while (true){
            while (n != 0){
                sum += (n % 10) * (n % 10);
                n /= 10;
            }
10          if (sum == 1){
                break;
            }
            if (set.contains(sum)){
15              return false;
            }
            else{
                set.add(sum);
            }
            n = sum;
            sum = 0;
20        }
        return true;
    }
}
```

Problem 26

61 Rotate List

Given a list, rotate the list to the right by k places, where k is non-negative.

For example:

Given $1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow \text{NULL}$ and $k = 2$,
return $4 \rightarrow 5 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow \text{NULL}$.

<https://leetcode.com/problems/rotate-list/>

Solution

Though this problem seems not that difficult, there are still some edge cases which needs to be dealt with. The main idea of my solution is to use 2 pointers to get the 'Nth element from the end'(sounds familiar), one goes faster for k steps.

Here comes the **TRAP!** When k is (much) larger than length of the list, we will get Time Limit Exceed error. So mod operator is in need here.

Also we need to be aware of the condition when $k == 0$ (or $k == n$ where $n \% n == 0$)

Listing 33: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode rotateRight(ListNode head, int k) {
        if (head == null || head.next == null || k == 0) {
            return head;
        }
        ListNode dummy = new ListNode(-1);
        dummy.next = head;
        int len = 0;
        ListNode cur = dummy;
        while (cur.next != null) {
            len++;
            cur = cur.next;
        }
        k = k % len;
        ListNode fast = head;
        while (k > 0) {
            fast = fast.next;
            if (k > 0 && fast == null) {
                fast = head;
            }
        }
        ListNode slow = head;
        if (fast == null) {
```

```
        return head;
    }
    while (fast.next != null){
35         fast = fast.next;
        slow = slow.next;
    }
    fast.next = dummy.next;
    dummy.next = slow.next;
40     slow.next = null;
    return dummy.next;
}
}
```

Problem 27

24 Swap Nodes in Pairs

Given a linked list, swap every two adjacent nodes and return its head.

For example,

Given 1->2->3->4, you should return the list as 2->1->4->3.

Your algorithm should use only constant space.

You may not modify the values in the list, only nodes itself can be changed.

<https://leetcode.com/problems/swap-nodes-in-pairs/>

Solution

Not many tactics, just maintain 2 pointers to reverse cur and next.

Listing 34: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode swapPairs(ListNode head) {
        if (head == null || head.next == null) {
            return head;
        }
        ListNode dummy = new ListNode(-1);
        ListNode tail = head;
        dummy.next = head.next;
        ListNode cur = head;
        ListNode next = cur.next;
        while (cur != null && next != null){
            ListNode tmp = next.next;
            tail.next = next;
            next.next = cur;
            cur.next = null;
            tail = cur;
            cur = tmp;
            if (cur != null){
                next = cur.next;
            }
            if (next == null){
                tail.next = cur;
            }
        }
        return dummy.next;
    }
}
```

Problem 28

66 Plus One

Given a non-negative number represented as an array of digits, plus one to the number.

The digits are stored such that the most significant digit is at the head of the list.

<https://leetcode.com/problems/plus-one/>

Solution

Though it is correct to use carry like we usually do when doing sum action, plus one is not as complicated to use carry digit. See listing 35, it's a pretty lovely solution.

Listing 35: Solution

```
public class Solution {  
    public int[] plusOne(int[] digits){  
        for (int idx = digits.length - 1 ; idx >= 0 ; idx--){  
            int digit = digits[idx];  
            if (digit < 9){  
                digits[idx] = digit + 1;  
                return digits;  
            }  
            else{  
                digits[idx] = 0;  
            }  
        }  
        int[] newDigits = new int[digits.length + 1];  
        newDigits[0] = 1;  
        for (int idx = 0 ; idx < digits.length ; idx++){  
            newDigits[idx + 1] = 0;  
        }  
        return newDigits;  
    }  
}
```


Problem 29

2 Add Two Numbers

You are given two linked lists representing two non-negative numbers.
The digits are stored in reverse order and each of their nodes contain
a single digit.
Add the two numbers and return it as a linked list.

Input: (2 -> 4 -> 3) + (5 -> 6 -> 4)

Output: 7 -> 0 -> 8

<https://leetcode.com/problems/add-two-numbers/>

Solution

Some sort of addition calculator with high precision. Some edge cases to be considered.
However, my solution is far from 'Clean Code'.

Listing 36: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode addTwoNumbers(ListNode head1, ListNode head2) {
        ListNode l1 = head1, l2 = head2;
        ListNode dummy = new ListNode(-1);
        ListNode cur = dummy;
        int carry = 0;
        while (l1 != null && l2 != null){
            int digit1 = l1.val;
            int digit2 = l2.val;
            int sum = digit1 + digit2 + carry;
            int digit = sum % 10;
            carry = sum / 10;
            ListNode newNode = new ListNode(digit);
            cur.next = newNode;
            cur = cur.next;
            l1 = l1.next;
            l2 = l2.next;
        }
        if (l1 == null){
            while (l2 != null){
                int digit = (l2.val + carry) % 10;
                carry = (l2.val + carry) / 10;
                ListNode newNode = new ListNode(digit);
                cur.next = newNode;
                cur = cur.next;
                l2 = l2.next;
            }
        }
        if (l2 == null){
            while (l1 != null){
                int digit = (l1.val + carry) % 10;
                carry = (l1.val + carry) / 10;
                ListNode newNode = new ListNode(digit);
                cur.next = newNode;
                cur = cur.next;
                l1 = l1.next;
            }
        }
        if (carry > 0){
            ListNode newNode = new ListNode(carry);
            cur.next = newNode;
        }
        return dummy.next;
    }
}
```

```
35         }
        }
        else if (l2 == null){
            while (l1 != null){
                int digit = (l1.val + carry) % 10;
                carry = (l1.val + carry) / 10;
40                ListNode newNode = new ListNode(digit);
                cur.next = newNode;
                cur = cur.next;
                l1 = l1.next;
45            }
        }
        if (l1 == null && l2 == null){
            if (carry > 0){
50                ListNode newNode = new ListNode(carry);
                cur.next = newNode;
            }
        }
55        return dummy.next;
    }
}
```

Problem 30

223 Rectangle Area

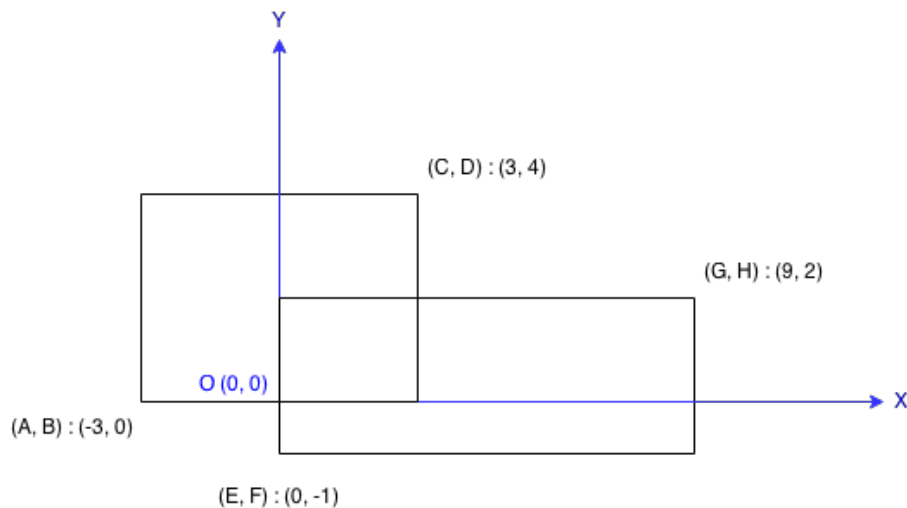
Find the total area covered by two rectilinear rectangles in a 2D plane.

Each rectangle is defined by its bottom left corner and top right corner as shown in the figure.

For Rectangle Area, see image below.

Assume that the total area is never beyond the maximum possible value of int.

<https://leetcode.com/problems/rectangle-area/>



Solution

Quite an easy problem, just read the description of problem more carefully.

Listing 37: Solution

```
public class Solution {
    public int computeArea(int A, int B, int C, int D,
        int E, int F, int G, int H) {
        int newXMin = Math.max(A, E);
        int newYMin = Math.max(B, F);
        int newXMax = Math.min(C, G);
        int newYMax = Math.min(D, H);
        int area1 = (D - B) * (C - A);
        int area2 = (H - F) * (G - E);
        if (newXMin < newXMax && newYMin < newYMax){
            int common = (newYMax - newYMin) * (newXMax - newXMin);
            return area1 + area2 - common;
        }
        else{
            return area1 + area2;
        }
    }
}
```

Problem 31

7 Reverse Integer

Reverse digits of an integer.

Example1: x = 123, return 321

Example2: x = -123, return -321

click to show spoilers.

Have you thought about this?

Here are some good questions to ask before coding.

Bonus points for you if you have already thought through this!

If the integer's last digit is 0, what should the output be?
ie, cases such as 10, 100.

Did you notice that the reversed integer might overflow?

Assume the input is a 32-bit integer, then the reverse of 1000000003 overflows.
How should you handle such cases?

For the purpose of this problem,

assume that your function returns 0 when the reversed integer overflows.

<https://leetcode.com/problems/reverse-integer/>

Solution

Two points which may attract interviewers if you have mentioned,

- What will happen if I input minus number?
- What if the reversed number overflows?

And solution in 'Clean Code Handbook' seems mistaken since it only consider the condition when $\text{sum} > \text{MAX_VALUE}/10$ and there still remains unresolved number. What about 2147483649?

Honestly, I don't think it is a problem as easy as I imagined at first sight.

Listing 38: Solution

```
public class Solution {  
    public int reverse(int x) {  
        int maxHead = Integer.MAX_VALUE / 10;  
        int maxTail = Integer.MAX_VALUE % 10;  
5       int minTail = Integer.MIN_VALUE % 10;  
        int origin = x;  
        int sum = 0;  
        while (origin != 0){  
            int digit = origin % 10;  
10       origin /= 10;  
            if (Math.abs(sum) > maxHead && digit != 0){  
                return 0;  
            }  
        }  
    }  
}
```

```
15         if (Math.abs(sum) == maxHead){
            if ((digit > 0 && digit > maxTail)
                || (digit < 0 && digit < minTail))
                return 0;
            }
            sum = sum * 10 + digit;
20     }
    return sum;
}
}
```

Problem 32

190 Reverse Bits

Reverse bits of a given 32 bits unsigned integer.

For example, given input 43261596

(represented in binary as 00000010100101000001111010011100),

return 964176192 (represented in binary as 00111001011110000010100101000000).

Follow up:

If this function is called many times, how would you optimize it?

<https://leetcode.com/problems/reverse-bits/>

Solution

Basically it is an easy bit manipulation problem. But when we need to call it several times, the method of optimization can be found <http://articles.leetcode.com/reverse-bits>.

Listing 39: Solution

```
public class Solution {  
    // you need treat n as an unsigned value  
    public int reverseBits(int n) {  
        int newNum = 0;  
5         for (int i = 0 ; i < 32 ; i++){  
            newNum = (newNum << 1);  
            newNum += (n & 1);  
            n = n >>> 1;  
        }  
10        return newNum;  
    }  
}
```

Problem 33

141 Linked List Cycle

Given a linked list, determine if it has a cycle in it.

Follow up:

Can you solve it without using extra space?

<https://leetcode.com/problems/linked-list-cycle/>

Solution

Still, the slow and fast pointer, if the list has cycle inside, the slow pointer will be caught sooner or later.

Listing 40: Solution

```
/**
 * Definition for singly-linked list.
 * class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) {
 *         val = x;
 *         next = null;
 *     }
 * }
 */
public class Solution {
    public boolean hasCycle(ListNode head) {
        ListNode slow = head;
        ListNode fast = head;
        while (fast != null && fast.next != null){
            slow = slow.next;
            fast = fast.next.next;

            if (slow == fast){
                return true;
            }
        }
        return false;
    }
}
```

Problem 34

142 Linked List Cycle II

Given a linked list, return the node where the cycle begins.
If there is no cycle, return null.

Note: Do not modify the linked list.

Follow up:
Can you solve it without using extra space?

<https://leetcode.com/problems/linked-list-cycle-ii/>

Solution

If no restrictions in that 'Note', we can easily handle this by HashSet.

Now we can only use mathematical method to solve this problem.

Forget about this problem for a moment. How can we get the length of cycle? When slow and fast pointers continues iterating so as to meet for the 2nd time, the slow pointer passes exactly the length of the whole cycle(or an integral multiple of cycle length)!

Coming back to this problem, we assign 'a' as length from head of list to where the cycle begins, 'b' as length from where cycle begins to where 2 pointers met, 'c' as length passed by fast pointer after it first passed the node where they met.

Then?

$2(a + b) = a + b + c$, here c may be $l \cdot R$, then we get $a = c - b!!!$

Which means, if we re-assign 2 pointers, let one starts from head of list, the other starts from where they met, they will meet at where the cycle begins!

For more details, refer to http://blog.sina.com.cn/s/blog_6f611c300101fs11.html

Listing 41: Solution

```
/**
 * Definition for singly-linked list.
 * class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) {
 *         val = x;
 *         next = null;
 *     }
 * }
 */
public class Solution {
    public ListNode detectCycle(ListNode head) {
        ListNode slow = head;
        ListNode fast = head;
        ListNode l1 = null, l2 = null;
        ListNode cycleBegins = null;
        while (fast != null && fast.next != null){
            slow = slow.next;
            fast = fast.next.next;
        }
```



```
        if (slow == fast){
            l1 = head;
            l2 = slow;

            break;
        }
    }
    if (fast != null && fast.next != null && l1 != null && l2 != null){
        while (l1 != l2){
            l1 = l1.next;
            l2 = l2.next;
        }

        cycleBegins = l1;
    }
    return cycleBegins;
}
```

Problem 35

13 Roman to Integer

Given a roman numeral, convert it to an integer.

Input is guaranteed to be within the range from 1 to 3999.

<https://leetcode.com/problems/roman-to-integer/>

Solution

Main idea is simple, just to check each character and sum up those values.

However, those IX(for 9, 11 is XI) and XL are special(they are designed for the purpose that the same notation won't appear for more than 3 times in a row.), honestly we can deal with all kinds of patterns, we still have a simplified algorithm: when cur digit is larger than prev(IX, X is larger), then we should add cur to sum and remove prev for 2 times. $sum = sum + cur - 2 * prev$.

Listing 42: Solution

```
public class Solution {
    private static final HashMap<Character, Integer> values
        = new HashMap<Character, Integer>(){
        {
5           put('I', 1);
            put('V', 5);
            put('X', 10);
            put('L', 50);
            put('C', 100);
10          put('D', 500);
            put('M', 1000);
        }
    };
    public int romanToInt(String s) {
15        int sum = 0;
        int prev = -1;
        char[] chars = s.toCharArray();
        for (char c : chars){
            int cur = values.get(c);
20            if (prev != -1 && prev < cur){
                sum += cur - 2 * prev;
            }
            else{
                sum += cur;
25            }
            prev = cur;
        }
        return sum;
    }
30 }
```

Problem 36

12 Integer to Roman

Given an integer, convert it to a roman numeral.

Input is guaranteed to be within the range from 1 to 3999.

<https://leetcode.com/problems/integer-to-roman/>

Solution

Find biggest value, remove it, and append corresponding Roman notation to String.

Difficult thing is that there exists IX or XV to indicate reduction. Luckily, there is a given information that integer is 1~3999, so that we can list all possible notation patterns since IX always appear after X(9 < 10, 10 comes first).

Listing 43: Solution

```
public class Solution {
    private static final String[] romans = { "M", "CM", "D", "CD", "C", "XC",
        "L", "XL", "X", "IX", "V", "IV", "I" };
    private static final int[] values = { 1000, 900, 500, 400, 100, 90, 50, 40,
5      10, 9, 5, 4, 1 };
    public String intToRoman(int num) {
        StringBuilder sb = new StringBuilder();
        while (num != 0){
            int idx = 0;
10         while (true){
                while (num < values[idx]){
                    idx ++;
                }
                sb.append(romans[idx]);
15         num -= values[idx];
                if (num == 0){
                    return sb.toString();
                }
            }
20     }
    return sb.toString();
}
```

Problem 37

168 Excel Sheet Column Title

Given a positive integer,
return its corresponding column title as appear in an Excel sheet.

For example:

```
1 -> A
2 -> B
3 -> C
...
26 -> Z
27 -> AA
28 -> AB
```

<https://leetcode.com/problems/excel-sheet-column-title/>

Solution

We can regard the integer as an number with base-26(binary is base-2 and decimal is base-10), but there is another thing we need to be aware of - base-2 consists only 0 and 1, where A-Z usually represents 1-26. So we should normalize our number, by reducing 1 to the integer we can normalize A-Z to 0-25.

Listing 44: Solution

```
public class Solution {
    public String convertToTitle(int n) {
        char[] titleDigits = new char[26];
        char a = 'A';
5      for (int idx = 0 ; idx < 26 ; idx++){
            titleDigits[idx] = (char)(a + idx);
        }
        String result = "";
        int cur = n;
10     int rem = -1;

        while (cur != 0){
            rem = (cur - 1) % 26;
            result = titleDigits[rem] + result;
15         cur = (cur - 1) / 26;
        }
        return result;
    }
}
```

Problem 38

171 Excel Sheet Column Number

Related to question Excel Sheet Column Title

Given a column title as appear in an Excel sheet, return its corresponding column number.

For example:

```
A -> 1
B -> 2
C -> 3
...
Z -> 26
AA -> 27
AB -> 28
```

<https://leetcode.com/problems/excel-sheet-column-number/>

Solution

It's just a reversion of its related problem.

Listing 45: Solution

```
public class Solution {
    public int titleToNumber(String s) {
        char[] title = s.toCharArray();
        int sum = 0;

        for (int idx = 0 ; idx < title.length ; idx++){
            char c = title[idx];
            int digit = c - 'A' + 1;

            sum = sum * 26 + digit;
        }

        return sum;
    }
}
```

Problem 39

60 Permutation Sequence

The set $[1, 2, 3, \dots, n]$ contains a total of $n!$ unique permutations.

By listing and labeling all of the permutations in order,
We get the following sequence (ie, for $n = 3$):

"123"

"132"

"213"

"231"

"312"

"321"

Given n and k , return the k th permutation sequence.

Note: Given n will be between 1 and 9 inclusive.

<https://leetcode.com/problems/permutation-sequence/>

Solution 1

This is maybe the 1st backtracking problem solved by myself!!

In fact, the problem can be solved without recursion in Listing 47

The idea of solving this problem is to justify each digit from the most significant digit. Using $(k - 1) / (n - 1)!$ we can get the index of current digit and remaining numbers.

For example, for $n = 3$ and $k = 4$, the 1st digit is decided by $(4 - 1) / (3 - 1)! = 1$ remains 1, so the 1st digit is the (1+1), 2nd biggest element in remaining set, which is 2; recursively, the whole number is 231.

Listing 46: Solution 1

```

public class Solution {
    public String getPermutation(int n, int k) {
        Set<Integer> numbers = new HashSet<Integer>();
        for (int i = 1; i <= n; i++){
5           numbers.add(i);
        }
        return getPermutation(numbers, n, k);
    }
    private String getPermutation(Set<Integer> numbers, int n, int k) {
10        if (n == 0){
            return "";
        }
        int qou = (k - 1) / getFact(n - 1);
        int rem = (k - 1) % getFact(n - 1);
15        Iterator<Integer> ite = numbers.iterator();
        int num = ite.next();
        int count = qou;

        while (ite.hasNext() && count-- > 0){
20            num = ite.next();
        }
        numbers.remove(num);
    }
}

```

```

        return String.valueOf(num) + getPermutation(numbers, n - 1, rem + 1);
    }
25
    private int getFact(int n){
        for (int i = 1 ; i <= n ; i++){
            power *= i;
        }
30    return power;
    }
}

```

Solution 2

Listing 47: Solution 2

```

class Solution {
    int getPermutationNumber(int n) {
        int result = 1;
        for(int i=1;i<=n;++i) {
5            result *=i;
        }

        return result;
    }
10 public:
    string getPermutation(int n, int k) {
        vector<int> result;
        vector<int> nums;
        for(int i=1;i<=n;++i) {
15            nums.push_back(i);
        }
        k=k-1;
        for(int i=0;i<n;++i) {
            int perms = getPermutationNumber(n-1-i);
            int index = k/perms;
20            result.push_back(nums[index]);
            k%=perms;
            nums.erase(nums.begin()+index);
        }
25        string s="";
        for(int i=0;i<result.size();++i) {
            s += std::to_string(result[i]);
        }
        return s;
30    }
};

```

Solution 3 There is a more interesting solution, given the fact that the numbers only contains 1-9.

Listing 48: Solution 3

```

class Solution {
public:
    string getPermutation(int n, int k) {

```

```
5      string res;
      string nums = "123456789";
      int f[10] = {1, 1, 2, 6, 24, 120, 720, 5040, 40320, 362880};
      --k;
      for (int i = n; i >= 1; --i) {
10         int j = k / f[i - 1];
         k %= f[i - 1];
         res.push_back(nums[j]);
         nums.erase(nums.begin() + j);
      }
15     return res;
};
```


Problem 40

46 Permutations

Given a collection of distinct numbers, return all possible permutations.

For example,

[1,2,3] have the following permutations:

[1,2,3], [1,3,2], [2,1,3], [2,3,1], [3,1,2], and [3,2,1].

<https://leetcode.com/problems/permutations/>

Solution

It is a representative problem of backtracking or recursion.

In fact, there are all kinds of algorithms to generate all-permutations.

See <http://blog.csdn.net/ljiabin/article/details/40151393>

Listing 49: Solution

```
public class Solution {
    Set<Integer> used = new HashSet<Integer>();
    List<List<Integer>> result = new ArrayList<List<Integer>>();
    List<Integer> array = new ArrayList<Integer>();
5   public List<List<Integer>> permute(int[] nums) {
        for (int idx = 0; idx < nums.length; idx++) {
            int num = nums[idx];
            if (used.contains(num)) {
                continue;
10            } else {
                used.add(num);
                array.add(num);
                if (array.size() == nums.length) {
                    ArrayList<Integer> newList = new ArrayList<Integer>();
15                    for (int ele : array) {
                        newList.add(ele);
                    }
                    result.add(newList);
                }

                permute(nums);
                array.remove(array.size() - 1);
                used.remove(num);
20            }
        }
25    }
    return result;
}
```

Problem 41

31 Next Permutation

Implement next permutation, which rearranges numbers into the lexicographically next greater permutation of numbers.

If such arrangement is not possible, it must rearrange it as the lowest possible order (ie, sorted in ascending order).

The replacement must be in-place, do not allocate extra memory.

Here are some examples. Inputs are in the left-hand column and its corresponding outputs are in the right-hand column.

1,2,3 -> 1,3,2

3,2,1 -> 1,2,3

1,1,5 -> 1,5,1

<https://leetcode.com/problems/next-permutation/>

Solution

Utilizes a lexicographical order Algorithm.

The algorithm is explained as below.

例2.3 设有排列(p) = 276**3**541, 按照字典式排序, 它的下一个排列是谁?

(q) = 276**4**135.

(1) 276**3**541 [找最后一个**正序**35]

(2) 276**3**541 [找**3**后面比**3**大的最后一个数]

(3) 276**4**531 [交换**3,4**的位置]

(4) 276**4**1**3**5 [把4后面的**531**反序排列为

135即得到最后的排列(q)]

● 求 $(p) = p_1 \dots p_{i-1} p_i \dots p_n$ 的下一个排列 (q) :

(1) 求 $i = \max\{j \mid p_{j-1} < p_j\}$ (找最后一个正序)

(2) 求 $j = \max\{k \mid p_{i-1} < p_k\}$ (找最后大于 p_{i-1} 者)

(3) 互换 p_{i-1} 与 p_j 得

$p_1 \dots p_{i-2} p_j p_i p_{i+1} \dots p_{j-1} p_{i-1} p_{j+1} \dots p_n$

(4) 反排 p_j 后面的数得到 (q) :

$p_1 \dots p_{i-2} p_j p_{i-1} p_{j+1} p_i p_{j+2} \dots p_{j-1} p_{j+1} p_i$

Listing 50: Solution

```

public class Solution {
    public void nextPermutation(int[] nums) {
        int preIdx = -1, postIdx = -1;
        for (int i = 0; i < nums.length - 1; i++) {
            if (nums[i] < nums[i + 1])
                preIdx = i;
        }
        if (preIdx < 0) {
            reverse(nums, 0, nums.length - 1);
            return;
        }
        for (int i = preIdx + 1; i < nums.length; i++) {
            if (nums[preIdx] < nums[i])
                postIdx = i;
        }
        int tmp = nums[preIdx];
        nums[preIdx] = nums[postIdx];
        nums[postIdx] = tmp;
        reverse(nums, preIdx + 1, nums.length - 1);
    }
    private void reverse(int[] nums, int i, int j) {
        while (i < j) {
            int tmp = nums[i];
            nums[i] = nums[j];
            nums[j] = tmp;
            i++;
            j--;
        }
    }
}

```

Problem 42

47 Permutations II

Given a collection of numbers that might contain duplicates, return all possible unique permutations.

For example,
[1,1,2] have the following unique permutations:
[1,1,2], [1,2,1], and [2,1,1].

<https://leetcode.com/problems/permutations-ii/>

Solution

We can use the same algorithm from Pro-31 to get the next permutation, but remember to sort the array first, or it may lose some results.

It can also be solved by backtracking method.

Listing 51: Solution

```
public class Solution {
    public List<List<Integer>> permuteUnique(int[] nums) {
        List<List<Integer>> list = new ArrayList<List<Integer>>();
        Arrays.sort(nums);
5       List<Integer> newPerm = new ArrayList<Integer>();
        for (int idx = 0 ; idx < nums.length ; idx++){
            newPerm.add(nums[idx]);
        }
        list.add(newPerm);
10      while (nextPermutation(nums)){
            newPerm = new ArrayList<Integer>();
            for (int idx = 0 ; idx < nums.length ; idx++){
                newPerm.add(nums[idx]);
            }
15         list.add(newPerm);
        }
        return list;
    }

    public boolean nextPermutation(int[] nums) {
20        int preIdx = -1 , postIdx = -1;
        for (int i = 0 ; i < nums.length - 1 ; i++){
            if (nums[i] < nums[i + 1]){
                preIdx = i;
            }
25        }
        if (preIdx < 0){
            return false;
        }
        for (int i = preIdx + 1; i < nums.length ; i++){
30            if (nums[preIdx] < nums[i]){
                postIdx = i;
            }
        }
    }
}
```

```
35     int tmp = nums[preIdx];
    nums[preIdx] = nums[postIdx];
    nums[postIdx] = tmp;
    reverse(nums, preIdx + 1, nums.length - 1);
    return true;
}
40
private void reverse(int[] nums, int i, int j) {
    while (i < j){
        int tmp = nums[i];
        nums[i] = nums[j];
45     nums[j] = tmp;

        i++;
        j--;
    }
50 }
}
```

Problem 43

77 Combinations

Given a collection of numbers that might contain duplicates,
return all possible unique permutations.

For example,

[1,1,2] have the following unique permutations:

[1,1,2], [1,2,1], and [2,1,1]. Given two integers n and k,
return all possible combinations of k numbers out of 1 ... n.

For example,

If n = 4 and k = 2, a solution is:

```
[
  [2,4],
  [3,4],
  [2,3],
  [1,2],
  [1,3],
  [1,4],
]
```

<https://leetcode.com/problems/combinations/>

Solution

A dfs problem, similar but not quite same as Problem “Permutations”, Permutations aim at swapping elements, this one is more like “N-Queens”.

Listing 52: Solution

```
public class Solution {
    List<List<Integer>> result = new ArrayList<List<Integer>>();

    public void dfsCombine(List<Integer> curList, int totalDigit, int curDigit, int totalNum) {
5         if (curDigit == totalDigit) {
            result.add(new ArrayList<Integer>(curList));
            return;
        }
        for (int number = curNumber + 1; number <= totalNumber; number++) {
10             if (curList.size() <= curDigit) {
                curList.add(number);
            }
            else {
                curList.set(curDigit, number);
15             }
            dfsCombine(curList, totalDigit, curDigit + 1, totalNumber, number);
        }
    }

20     public List<List<Integer>> combine(int n, int k) {
        List<Integer> list = new ArrayList<Integer>();
    }
```

```
25         dfsCombine(list , k , 0, n, 0);  
           return result;  
       }  
   }
```

Problem 44

172 Factorial Trailing Zeroes

Given an integer n , return the number of trailing zeroes in $n!$.

Note: Your solution should be in logarithmic time complexity.

<https://leetcode.com/problems/factorial-trailing-zeroes/>

Solution

What decides the number of trailing 0 is how many numbers has the factor of 5 (since $5 \times 2 = 10$ and count of 2 is much more than 5). Note that not only 5 should be considered, 25 provides two 5s, which also needs to be taken into consideration.

Listing 53: Solution

```
public class Solution {  
    public int trailingZeroes(int n) {  
        int sum = 0;  
        while (n != 0){  
            n = n / 5;  
            sum += n;  
        }  
        return sum;  
    }  
}
```


Problem 45

8 String to Integer (atoi)

Implement atoi to convert a string to an integer.

Hint: Carefully consider all possible input cases. If you want a challenge, please do not see below and ask yourself what are the possible input cases.

Notes: It is intended for this problem to be specified vaguely (ie, no given input specs). You are responsible to gather all the input requirements up front.

Requirements for atoi:

The function first discards as many whitespace characters as necessary until the first non-whitespace character is found.

Then, starting from this character, takes an optional initial plus or minus sign followed by as many numerical digits as possible, and interprets them as a numerical value.

The string can contain additional characters after those that form the integral number, which are ignored and have no effect on the behavior of this function.

If the first sequence of non-whitespace characters in str is not a valid integral number, or if no such sequence exists because either str is empty or it contains only whitespace characters, no conversion is performed.

If no valid conversion could be performed, a zero value is returned.

If the correct value is out of the range of representable values, INT_MAX (2147483647) or INT_MIN (-2147483648) is returned.

<https://leetcode.com/problems/string-to-integer-atoi/>

Solution

The biggest difficulty of this problem is that there are so many edge cases(though less than those in valid number).

Listing 54: Solution

```
public class Solution {
    public int myAtoi(String str) {
        if (str == null || str.length() == 0){
            return 0;
        }
        int sign = 1;
        char[] ch = str.toCharArray();
        int idx = 0;
        while (ch[idx] == ' '){
            idx ++;
        }
        if (ch[idx] == '-'){
```

```
        sign = -1;
        idx++;
15    }
    else if (ch[idx] == '+'){
        idx++;
    }
    int sum = 0;
20    while (idx < str.length()){
        int temp = ch[idx] - '0';
        if (temp > 9 || temp < 0){
            return sum * sign;
        }
25    if (sum > Integer.MAX_VALUE / 10 ||
        (sum == Integer.MAX_VALUE / 10 && temp > Integer.MAX_VALUE % 10)){
        return (sign == 1) ? Integer.MAX_VALUE : Integer.MIN_VALUE;
    }
    sum = sum * 10 + temp;
30    idx ++;
    }
    return sum * sign;
}
}
```

Problem 46

28 Implement strStr()

Implement strStr().

Returns the index of the first occurrence of needle in haystack,
or -1 if needle is not part of haystack.

<https://leetcode.com/problems/implement-strstr/>

Solution

An easy string manipulating problem. Since I'm using a $O(mn)$ solution, and there's another famous algorithm named KMP which can solve this problem with lower complexity, the solution has the space for improvement.

And also don't forget the solution on "Clean Code Handbook".

Listing 55: Solution

```
public class Solution {  
    public int strStr(String haystack, String needle) {  
        if (needle.length() == 0){  
            return 0;  
        }  
        for (int hIdx = 0 ; hIdx < haystack.length() ; hIdx++){  
            if (haystack.length() - hIdx < needle.length()){  
                return -1;  
            }  
  
            if (haystack.charAt(hIdx) == needle.charAt(0)){  
                int hStart = hIdx;  
                int nStart = 0;  
                while (nStart < needle.length()  
                    && haystack.charAt(hStart) == needle.charAt(nStart)){  
                    hStart++;  
                    nStart++;  
                }  
                if (nStart == needle.length()){  
                    return hIdx;  
                }  
            }  
        }  
  
        return -1;  
    }  
}
```

Problem 47

169 Majority Element

Given an array of size n , find the majority element.
The majority element is the element that appears more than $\text{floor}(n/2)$ times.

You may assume that the array is non-empty and the majority element always exist in the array.

<https://leetcode.com/problems/majority-element/>

Solution 1

This is my solution, luckily write out **Moore voting algorithm**, though not 'Clean Code':

We maintain a current candidate and a counter initialized to 0. As we iterate the array, we look at the current element x :

If the counter is 0, we set the current candidate to x and the counter to 1.

If the counter is not 0, we increment or decrement the counter based on whether x is the current candidate.

After one pass, the current candidate is the majority element. Runtime complexity = $O(n)$.

Listing 56: Solution 1

```
public class Solution {
    public int majorityElement(int[] nums) {
        int curNum = -1;
        int curCount = 0;
        int idx1 = 0, idx2 = 1;
        while (idx1 < nums.length && idx2 < nums.length) {
            if (nums[idx1] == nums[idx2]) {
                curCount++;
                curNum = nums[idx1];
                idx2++;
            } else {
                if (nums[idx1] == curNum) {
                    if (curCount > 0) {
                        curCount--;
                        idx2++;
                    }
                    continue;
                }
                else {
                    curNum = -1;
                }
            }
            idx1 = idx2 + 1;
            idx2 = idx2 + 2;
        }
        if (idx1 == nums.length - 1 && curNum == -1) {
            curNum = nums[idx1];
        }
    }
}
```

```
        return curNum;
    }
}
```

Solution 2

The reason why I record Solution 2 in such an excitement is that the algorithm is FXXXING excellent!!

We would need 32 iterations, each calculating the number of 1's for the *i*th bit of all *n* numbers. Since a majority must exist, therefore, either count of 1's > count of 0's or vice versa (but can never be equal). The majority numbers *i*th bit must be the one bit that has the greater count. (Words below are copied from Solutions of this problem in Leetcode)

Listing 57: Solution 2

```
public int majorityElement(int[] num) {
    int ret = 0;
    for (int i = 0; i < 32; i++) {
        int ones = 0, zeros = 0;
5         for (int j = 0; j < num.length; j++) {
            if ((num[j] & (1 << i)) != 0) {
                ++ones;
            }
            else
10                ++zeros;
        }

        if (ones > zeros)
            ret |= (1 << i);
15    }
    return ret;
}
```

Problem 48

229 Majority Element II

Given an integer array of size n ,
find all elements that appear more than $\text{floor}(n/3)$ times.
The algorithm should run in linear time and in $O(1)$ space.

Hint:

How many majority elements could it possibly have?

<https://leetcode.com/problems/majority-element-ii/>

Solution

Also, use Moore's Voting Algorithm, this time we have 2 candidates comparing to other nodes. Remember to check finally whether the two candidates are all majority elements we want since there may be only one majority element.

Listing 58: Solution 1

```
public class Solution {
    public List<Integer> majorityElement(int[] nums) {
        int cand1 = -1, cand2 = -1;
        int count1 = 0, count2 = 0;

        for (int num : nums){
            if (num == cand1){
                count1++;
            } else if (num == cand2){
                count2++;
            }
            else {
                // num <> cand1 & num <> cand2
                if (count1 == 0){
                    cand1 = num;
                    count1 = 1;
                }
                else if (count2 == 0){
                    cand2 = num;
                    count2 = 1;
                }
                else{
                    count1--;
                    count2--;
                }
            }
        }
        count1 = 0;
        count2 = 0;
        // check if there is only one majority element
        for (int num : nums){
            if (num == cand1){
                count1++;
            }
        }
    }
}
```

```
    }
    35     if (num == cand2){
        count2 ++;
    }
}
List<Integer> result = new ArrayList<Integer>();
40     if (count1 > nums.length / 3){
        result.add(cand1);
    }
    if (count2 > nums.length / 3){
        result.add(cand2);
45     }

    return result;
}
}
```

Another interesting solution

From 'Clean Code Handbook', with comments served as description.

Listing 59: Solution 2

```
class Solution {
public:
    5     int singleNumber(int A[], int n) {
        int one,two,three;
        one=two=three=0;
        for(int i=0;i<n;i++){
            // Already appeared twice
            three = two & A[i];
10         // Appeard twice add circumstances
            // when appeared once and in A[i] there is another
            two = two | one & A[i];
            // Appeared once
            one = one | A[i];
15         // When already appeared for the 3rd time,
            // clear one and two
            one = one & ~three;
            two = two & ~three;
        }
20         return one;
    }
};
```

Problem 49

67 Add Binary

Given two binary strings, return their sum (also a binary string).

For example,
a = "11"
b = "1"
Return "100".

<https://leetcode.com/problems/add-binary/>

Solution

An easy problem just like “Add Two Number”, easier.

Listing 60: Solution

```
public class Solution {  
    public String addBinary(String a, String b) {  
        StringBuilder sb = new StringBuilder();  
        int lenA = a.length();  
        int lenB = b.length();  
        if (a.length() < b.length()){  
            return addBinary(b, a);  
        }  
        char[] chA = a.toCharArray();  
        char[] chB = b.toCharArray();  
        int carry = 0;  
        for (int idxA = lenA - 1 ; idxA >= 0 ; idxA--){  
            int idxB = idxA - (lenA - lenB);  
            int digitA = chA[idxA] - '0';  
            int digitB = (idxB < 0) ? 0 : chB[idxB] - '0';  
            int rem = (digitA + digitB + carry) % 2;  
            carry = (digitA + digitB + carry) / 2;  
            sb.append(String.valueOf(rem));  
        }  
        if (carry == 1){  
            sb.append("1");  
        }  
        return sb.reverse().toString();  
    }  
}
```


Problem 50

258 Add Digits

Given a non-negative integer `num`, repeatedly add all its digits until the result has only one digit.

For example:

Given `num = 38`, the process is like: $3 + 8 = 11$, $1 + 1 = 2$.
Since 2 has only one digit, return it.

Follow up:

Could you do it without any loop/recursion in $O(1)$ runtime?

Hint:

A naive implementation of the above process is trivial.
Could you come up with other methods?

<https://leetcode.com/problems/add-digits/>

Solution 1

First come up with the “naive” solution.

Listing 61: Solution 1

```
public class Solution {  
    public int addDigits(int num) {  
        int sum = 0;  
        while (num != 0){  
            int digit = num % 10;  
            num /= 10;  
            sum += digit;  
            if (num == 0){  
                if (sum >= 10){  
                    num = sum;  
                    sum = 0;  
                }  
            }  
        }  
        return sum;  
    }  
}
```

Solution 2

Solution 2 is interesting and simple! See figure below.

另一个方法比较简单，可以举例说明一下。假设输入的数字是一个5位数字num，则num的各位分别为a、b、c、d、e。

有如下关系： $\text{num} = a * 10000 + b * 1000 + c * 100 + d * 10 + e$

即： $\text{num} = (a + b + c + d + e) + (a * 9999 + b * 999 + c * 99 + d * 9)$

因为 $a * 9999 + b * 999 + c * 99 + d * 9$ 一定可以被9整除，因此num模除9的结果与 $a + b + c + d + e$ 模除9的结果是一样的。

对数字 $a + b + c + d + e$ 反复执行同类操作，最后的结果就是一个1-9的数字加上一串数字，最左边的数字是1-9之间的，右侧的数字永远都是可以被9整除的。

这道题最后的目标，就是不断将各位相加，相加到最后，当结果小于10时返回。因为最后结果在1-9之间，得到9之后将不会再对各位进行相加，因此不会出现结果为0的情况。因为 $(x + y) \% z = (x \% z + y \% z) \% z$ ，又因为 $x \% z \% z = x \% z$ ，因此结果为 $(\text{num} - 1) \% 9 + 1$ ，只模除9一次，并将模除后的结果加一返回。

Listing 62: Solution 2

```
public class Solution {  
    public int addDigits(int num) {  
        return (num - 1) % 9 + 1;  
    }  
}
```

5

Problem 51

233 Number of Digit One

Given an integer n , count the total number of digit 1 appearing in all non-negative integers less than or equal to n .

For example:

Given $n = 13$,

Return 6, because digit 1 occurred in the following numbers: 1, 10, 11, 12, 13.

Hint:

Beware of overflow.

<https://leetcode.com/problems/number-of-digit-one/>

Solution

Honestly, I don't think I totally understand the solution, so that I'll list my current train of thought in case of forgetting.

例子更容易理解，所以以例子辅助说明。
主要思路是基于每十个数（0-9）就会有1个个位数为1，每100个数就有10个十位数为1，.....
这样的话做循环看每一位就行了。
但是需要注意的是，以十位数为例，对于205, 215, 255，所得到的值是不一样的（215比205多210~215这6个；255多所有210开头的十位数为1的。）
因此这几种情况要分别考虑。
0: 就是高位数乘以现在的base，对于205就是2（2是0的高位，我们现在正在处理十位数）*10；
1: 高位数乘以base还要加上低位数（例子中是5, 0的低位），还要加1，因为0~5是6个而不是5个。
2: 是（高位数+1）乘以base，1在例子中就表示21开头的数字们
然后累加即可

Listing 63: Solution

```

public class Solution {
    public int countDigitOne(int n) {
        if (n < 0){
            return 0;
        }
        int base = 1;
        int num = n;
        int count = 0;
        while (num != 0){
            int lower = n - base * num;
            int quo = num / 10;
            int rem = num % 10;
            switch(rem){
                case 0:
                    count += quo * base;
                    break;
                case 1:
                    count += quo * base + lower + 1;
                    break;
                default:
                    count += (quo + 1) * base;
            }
            base *= 10;
            num = quo;
        }
    }
}

```

```
25         num /= 10;  
        }  
        return count;  
    }  
}
```

Problem 52

204 Count Primes

Count the number of prime numbers less than a non-negative number, n .

<https://leetcode.com/problems/count-primes/>

Solution 1

First solution is straightforward but a bit inefficient.

Listing 64: Solution 1

```
public class Solution {  
    public int countPrimes(int n) {  
        int count = 0;  
        for (int i = 1; i < n; i++) {  
            if (isPrime(i))  
                count++;  
        }  
        return count;  
    }  
  
    private boolean isPrime(int num) {  
        if (num <= 1)  
            return false;  
        // Loop's ending condition is  $i * i \leq \text{num}$  instead of  $i \leq \text{sqrt}(\text{num})$   
        // to avoid repeatedly calling an expensive function sqrt().  
        for (int i = 2; i * i <= num; i++) {  
            if (num % i == 0)  
                return false;  
        }  
        return true;  
    }  
}
```

Solution 2

Solution 2 is actually extracted from the hint spoilers from problem descriptions, since it is such a litany, I put it in solution description rather than in problem.

1. Let's start with a `isPrime` function. To determine if a number is prime, we need to check if it is not divisible by any number less than n . The runtime complexity of `isPrime` function would be $O(n)$ and hence counting the total prime numbers up to n would be $O(n^2)$. Could we do better?
2. As we know the number must not be divisible by any number $> n / 2$, we can immediately cut the total iterations half by dividing only up to $n / 2$. Could we still do better?
3. Let's write down all of 12's factors:
 $2 * 6 = 12$
 $3 * 4 = 12$
 $4 * 3 = 12$
 $6 * 2 = 12$

4. As you can see, calculations of $4 * 3$ and $6 * 2$ are not necessary. Therefore, we only need to consider factors up to \sqrt{n} because, if n is divisible by some number p , then $n = p * q$ and since $p \leq q$, we could derive that $p \leq \sqrt{n}$.

Our total runtime has now improved to $O(n^{1.5})$, which is slightly better. Is there a faster approach? Code see solution 1.

5. The Sieve of Eratosthenes is one of the most efficient ways to find all prime numbers up to n . But don't let that name scare you, I promise that the concept is surprisingly simple.

For figure of the algorithm,

see https://leetcode.com/static/images/solutions/Sieve_of_Eratosthenes_animation.gif.

We start off with a table of n numbers. Let's look at the first number, 2. We know all multiples of 2 must not be primes, so we mark them off as non-primes. Then we look at the next number, 3. Similarly, all multiples of 3 such as $3 * 2 = 6$, $3 * 3 = 9$, ... must not be primes, so we mark them off as well. Now we look at the next number, 4, which was already marked off. What does this tell you? Should you mark off all multiples of 4 as well?

6. 4 is not a prime because it is divisible by 2, which means all multiples of 4 must also be divisible by 2 and were already marked off. So we can skip 4 immediately and go to the next number, 5. Now, all multiples of 5 such as $5 * 2 = 10$, $5 * 3 = 15$, $5 * 4 = 20$, $5 * 5 = 25$, ... can be marked off. There is a slight optimization here, we do not need to start from $5 * 2 = 10$. Where should we start marking off?
7. In fact, we can mark off multiples of 5 starting at $5 * 5 = 25$, because $5 * 2 = 10$ was already marked off by multiple of 2, similarly $5 * 3 = 15$ was already marked off by multiple of 3. Therefore, if the current number is p , we can always mark off multiples of p starting at p^2 , then in increments of p : $p^2 + p$, $p^2 + 2p$, ... Now what should be the terminating loop condition?
8. It is easy to say that the terminating loop condition is $p < n$, which is certainly correct but not efficient. Do you still remember Hint #3?
9. Yes, the terminating loop condition can be $p < \sqrt{n}$, as all non-primes $\geq \sqrt{n}$ must have already been marked off. When the loop terminates, all the numbers in the table that are non-marked are prime.

The Sieve of Eratosthenes uses an extra $O(n)$ memory and its runtime complexity is $O(n \log \log n)$. For the more mathematically inclined readers, you can read more about its algorithm complexity on Wikipedia. Code snippets are in Solution listing.

Listing 65: Solution 2

```

public class Solution {
    public int countPrimes(int n) {
        boolean[] isPrime = new boolean[n];
        for (int i = 2; i < n; i++) {
5           isPrime[i] = true;
        }
        // Loop's ending condition is i * i < n instead of i < sqrt(n)
        // to avoid repeatedly calling an expensive function sqrt().
        for (int i = 2; i * i < n; i++) {
10          if (!isPrime[i])
              continue;
          for (int j = i * i; j < n; j += i) {
              isPrime[j] = false;
          }
        }
    }
}

```

```
15     }  
    int count = 0;  
    for (int i = 2; i < n; i++) {  
        if (isPrime[i])  
            count++;  
20     }  
    return count;  
    }  
}
```

Problem 53

20 Valid Parentheses

Given a string containing just the characters '(', ')', '{', '}', '[' and ']', determine if the input string is valid.

The brackets must close in the correct order, "`()`" and "`()[]{}`" are all valid but "`()`" and "`()[]`" are not.

<https://leetcode.com/problems/valid-parentheses/>

Solution

A representative stack problem, use stack to hold left parentheses, pop when meet right hands, check if match.

Listing 66: Solution

```
public class Solution {
    public boolean isValid(String s) {
        HashMap<Character, Character> pairs
        = new HashMap<Character, Character>(){
5          {
              put('(', ')');
              put('[', ']');
              put('{', '}');
          }
        };
10      Stack<Character> stack = new Stack<Character>();
        for (int idx = 0; idx < s.length(); idx++){
            char ch = s.charAt(idx);
            if (pairs.containsKey(ch)){
15                stack.push(ch);
            }
            else{
                if (stack.isEmpty()){
20                    return false;
                }
                else{
                    char left = stack.pop();
                    char expRight = pairs.get(left);
                    if (expRight != ch){
25                        return false;
                    }
                }
            }
        }
30      if (stack.isEmpty()){
          return true;
        }
        else {
            return false;
35      }
    }
}
```



```
}
```

Problem 54

32 Longest Valid Parentheses

Given a string containing just the characters '(' and ')',
find the length of the longest valid (well-formed) parentheses substring.

For "()", the longest valid parentheses substring is "()",
which has length = 2.

Another example is "()()()", where
the longest valid parentheses substring is "()()", which has length = 4.

<https://leetcode.com/problems/longest-valid-parentheses/>

Solution

It is a dp problem, but we can first think about how to solve it in Brute force Method, which is of $O(n^2)$ runtime and will get Time Limit Exceed error.

The dp solution is as listing. Interestingly, the array iterates from the end of string, $dp[i]$ indicates the length of longest valid parentheses starts from index i . Of course when $s[i]$ is ')' it cannot be the start point, so we only need to consider the '(' condition. if $dp[i+1] > 0$, which means there are valid ()s starts from $i+1$, then we should check if $s[i+dp[i+1]+1]$ is ')', if is the valid ()s ends at $i+dp[i+1]+1$ then $dp[i]=dp[i+1]+2$.

This is not over, since $dp[i+dp[i+1]+2]$ may not be 0, which we also need to take into addition. (We don't need to finish such process recursively since the process is done when we were calculating $dp[i+dp[i+1]+2]$).

Be sure to check every boundaries conditions!!!

Listing 67: Solution Brute force

```

1  int findlength(string s) {
2      int paired = 0;
3      int len = 0;
4      stack<char> stk;
5      for(int i=0; i< s.length(); ++i)
6      {
7          if(s[i]=='(')stk.push('(');
8          else if(!stk.empty()){
9              stk.pop();
10             len+=2;
11             if(stk.empty()&&paired<len)
12                 paired=len;
13         }else{
14             len = 0;
15         }
16     }
17     if(stk.empty() && paired == 0)
18         paired += len;
19     return paired;
20 }
21 int longestValidParentheses(string s) {
22     // Note: The Solution object is instantiated only once.
23     int max = 0;
24     int tmp = 0;
25     for(int i=0; i< s.length(); i++)

```

```
    {
        tmp = findlength(s.substr(i));
        max = max>tmp?max:tmp;
    }
30  return max;
}
```

Listing 68: Solution DP

```
public class Solution {
    public int longestValidParentheses(String s) {
        int[] dp = new int[s.length()];
        int max = 0;
5      for (int idx = s.length() - 2; idx >= 0 ; idx --){
            if (s.charAt(idx) == '('){
                int nextIdx = idx + dp[idx + 1] + 1;
                if (nextIdx >= s.length()){
                    continue;
10                }
                if (s.charAt(nextIdx) == ')'){
                    dp[idx] = dp[idx + 1] + 2;
                    if (nextIdx + 1 <= s.length() - 1
15                      && s.charAt(nextIdx + 1) == '('){
                        dp[idx] += dp[nextIdx + 1];
                    }
                    if (dp[idx] > max){
                        max = dp[idx];
                }
            }
        }
        return max;
25 }
}
```

Problem 55

155 Min Stack

Design a stack that supports push, pop, top, and retrieving the minimum element in constant time.

push(x) -- Push element x onto stack.
pop() -- Removes the element on top of the stack.
top() -- Get the top element.
getMin() -- Retrieve the minimum element in the stack.

<https://leetcode.com/problems/min-stack/>

Solution

We can use a stack with the same size of original stack to maintain the minimum value, however, we don't have to need such a large stack. What we need to do is to record elements less than current minimum value, and take care to check equation when popping take place.

Listing 69: Solution

```
class MinStack {  
    Stack<Integer> origin = new Stack<Integer>();  
    Stack<Integer> min = new Stack<Integer>();  
  
    5    public void push(int x) {  
        origin.push(x);  
        if (min.isEmpty() || x <= min.peek()){  
            min.push(x);  
        }  
    10    }  
  
    public void pop() {  
        int x = origin.pop();  
        15    if (x == min.peek()){  
            min.pop();  
        }  
    }  
  
    20    public int top() {  
        return origin.peek();  
    }  
  
    25    public int getMin() {  
        return min.peek();  
    }  
}
```

Problem 56

239 Sliding Window Maximum

Given an array `nums`, there is a sliding window of size `k` which is moving from the very left of the array to the very right. You can only see the `k` numbers in the window. Each time the sliding window moves right by one position.

For example,

Given `nums = [1,3,-1,-3,5,3,6,7]`, and `k = 3`.

Window position	Max
[1 3 -1] -3 5 3 6 7	3
1 [3 -1 -3] 5 3 6 7	3
1 3 [-1 -3 5] 3 6 7	5
1 3 -1 [-3 5 3] 6 7	5
1 3 -1 -3 [5 3 6] 7	6
1 3 -1 -3 5 [3 6 7]	7

Therefore, return the max sliding window as `[3,3,5,5,6,7]`.

Note:

You may assume `k` is always valid,

ie: `1 <= k <= input array's size` for non-empty array.

Follow up:

Could you solve it in linear time?

Hint:

How about using a data structure such as deque (double-ended queue)?

The queue size need not be the same as the windows size.

Remove redundant elements and the queue

should store only elements that need to be considered.

<https://leetcode.com/problems/sliding-window-maximum/>

Solution

There are many solutions to solve such problem, the simplest but slowest way is using Brute force, which is of $O(nk)$ runtime.

Other ways of solving it are listed as below,

Method 1: 最朴素的解法就是将窗口在数组上进行滑动，每滑动一次求一下窗口的最值。时间复杂度 $O(nk)$ 。

Method 2: 还有就是使用平衡二叉树。

I. 取出数组前`k`个元素，构建平衡二叉树。

II. for `i = 0 → (n-1)`

- 取出平衡二叉树最大值，加入结果集`result`。
- find `nums[i]`，删除这个节点。
- 将`nums[i+k]`出入平衡二叉树。

时间复杂度为 $O(n \log k)$ 。

Method 3: 其实还有更高效的解法就是使用 deque

使用一个 deque，只存储当前窗口中会对之后的选择有影响的元素。插入时每次从队尾进行，每次滑动窗口时只判断队首元素是否超出窗口范围，超出的话则删除。还有就是比较队尾元素与当前元素大小，如果小的话，直接删除队尾元素，否则插入。

判断元素是否超出窗口的范围，需要根据元素的下标。然而根据元素的下标可以确定元素的值，这一可以直接用 queue 保存数组下标的值。

Let's see an example.

实际上一个滑动窗口可以看成是一个队列。当窗口滑动时，处于窗口的第一个数字被删除，同时在窗口的末尾添加一个新的数字。这符合队列的先进先出特性。如果能从队列中找出它的最大数，这个问题也就解决了。

我们实现了一个可以用 $O(1)$ 时间得到最小值的栈。同样，也可以用 $O(1)$ 时间得到栈的最大值。同时我们讨论了如何用两个栈实现一个队列。综合这两个问题的解决方法，我们发现如果把队列用两个栈实现，由于可以用 $O(1)$ 时间得到栈中的最大值，那么也就可以用 $O(1)$ 时间得到队列的最大值，因此总的复杂度也就降到了 $O(n)$ 。我们可以用这个方法来解决这个问题。不过这样就相当于一轮面试的时间内要做两个面试题，时间未必够用。再来看看有没有其它的方法。

下面换一种思路。我们并不把滑动窗口的每个数值都存入队列中，而只把有可能成为滑动窗口最大值的数值存入到一个两端开口的队列。接着以输入数字 $[2, 3, 4, 2, 6, 2, 5, 1]$ 为例进一步分析。

数组的第一个数字是2，把它存入队列中。第二个数字是3，由于它比前一个数字2大，因此2不可能成为滑动窗口中的最大值。2先从队列里删除，再把3存入到队列中。此时队列中只有一个数字3。针对第三个数字4的步骤类似，最终在队列中只剩下一个数字4。此时滑动窗口中已经有3个数字，而它的最大值4位于队列的头部。

接下来处理第四个数字2。2比队列中的数字4小。当4滑出窗口之后2还是有可能成为滑动窗口的最大值，因此把2存入队列的尾部。现在队列中有两个数字4和2，其中最大值4仍然位于队列的头部。

下一个数字是6，由于它比队列中已有的数字4和2都大，因此这时4和2已经不可能成为滑动窗口中的最大值。先把4和2从队列中删除，再把数字6存入队列。这个时候最大值6仍然位于队列的头部。

第六个数字是2，由于它比队列中已有的数字6小，所以2也存入队列的尾部。此时队列中有两个数字，其中最大值6位于队列的头部。

接下来的数字是5，在队列中已有的两个数字6和2里，2小于5，因此2不可能是一个滑动窗口的最大值，可以把它从队列的尾部删除。删除数字2之后，再把数字5存入队列。此时队列里剩下两个数字6和5，其中位于队列头部的是最大值6。

数组最后一个数字是1，把1存入队列的尾部。注意到位于队列头部的数字6是数组的第5个数字，此时的滑动窗口已经不包括这个数字了，因此应该把数字6从队列删除。那么怎么知道滑动窗口是否包括一个数字？应该在队列里存入数字在数组里的下标，而不是数值。当一个数字的下标与当前处理的数字的下标之差大于或者等于滑动窗口的大小时，这个数字已经从滑动窗口中滑出，可以从队列中删除了。

LinkedList in Java has similar interface as deque in C++.

(BTW, there is another way to solve it...) Just think about “PROBLEM 155 Min Stack”, We have already implemented a min stack with $O(1)$ runtime pop(), push() and getMin(), we can accomplish max stack in a similar way, so that we can use the stack to finish this task.

Listing 70: Solution

```
public class Solution {
    public int[] maxSlidingWindow(int[] nums, int k) {
        if (k == 0 || nums.length == 0) {
            return nums;
        }
        LinkedList<Integer> windows = new LinkedList();
        int[] result = new int[nums.length - k + 1];
        for (int idx = 0 ; idx < nums.length ; idx++){
            int cur = nums[idx];
            if (!windows.isEmpty() && windows.getFirst() <= idx - k){
                windows.removeFirst();
            }
            while (!windows.isEmpty() && cur > nums[windows.getLast()]){
                windows.removeLast();
            }
            windows.addLast(idx);
            if (idx >= k - 1){
                result[idx - k + 1] = nums[windows.getFirst()];
            }
        }
        return result;
    }
}
```

Problem 57

189 Rotate Array

Rotate an array of n elements to the right by k steps.

For example, with $n = 7$ and $k = 3$,
the array $[1, 2, 3, 4, 5, 6, 7]$ is rotated to $[5, 6, 7, 1, 2, 3, 4]$.

Note:

Try to come up as many solutions as you can,
there are at least 3 different ways to solve this problem.

Hint:

Could you do it in-place with $O(1)$ extra space?
Related problem: Reverse Words in a String II

<https://leetcode.com/problems/rotate-array/>

Solution 1

Problem ask us to solve the problem in 3 different ways, I'll list my own solutions, and list others in below sections(only description, no code).

Solution 1 is simple, just use another array to simulate the process in description.

Listing 71: Solution 1

```
public class Solution {  
    public void rotate(int[] nums, int k) {  
        int[] result = new int[nums.length];  
        k = k % nums.length;  
5        for (int i = nums.length - k ; i < nums.length ; i++){  
            result[i - nums.length + k] = nums[i];  
        }  
        for (int j = 0 ; j < nums.length - k ; j++){  
            result[k + j] = nums[j];  
10        }  
        for (int i = 0 ; i < nums.length ; i++){  
            nums[i] = result[i];  
        }  
15    }  
}
```

Solution 2

Solution is to move every bit backward for k digits, but which digit should I swap with? This solution forms a cycle, for example, if $n = 9$, $k = 3$, we shall swap $a[0]$ to $a[3]$, then $a[3]$ to $a[6]$, next one should be $a[0]$, which returns to start point. Now we should finish this cycle and move cursor to next digit, namely $a[1]$ to $a[4]$...

Listing 72: Solution 2

```
public class Solution {  
    public void rotate(int[] nums, int k) {
```

```

    int count = 0;
    int start = 0;
5   int last = nums[0];
    k %= nums.length;
    if (k == 0){
        return;
    }
10   while (count < nums.length && start < nums.length - 1){
        int cur = start;
        last = nums[cur];
        do{
            int swap = nums[cur];
            nums[cur] = last;
15             last = swap;
            cur = (cur + k) % nums.length;
            count++;
        } while (cur != start);
20     nums[start] = last;

        start++;
    }
    }
25 }
```

Solution 3

An interesting algorithm, similar to algorithm followed in “Reverse Words in a String II”. Reverse 1st half, reverse 2nd half, reverse the whole sequence.

Listing 73: Solution 3

```

public class Solution {
    public void rotate(int[] nums, int k) {
        k = k % nums.length;
        reverse(nums, 0, nums.length - k - 1);
5       reverse(nums, nums.length - k, nums.length - 1);
        reverse(nums, 0, nums.length - 1);
    }

    private void reverse(int[] nums, int start, int end){
10        int i = start, j = end;
        while (i < j){
            int tmp = nums[i];
            nums[i] = nums[j];
            nums[j] = tmp;
15
            i++;
            j--;
        }
    }
20 }
```


Problem 58

263 Ugly Number

Write a program to check whether a given number is an ugly number.

Ugly numbers are positive numbers whose prime factors only include 2, 3, 5.

For example, 6, 8 are ugly while 14 is not ugly
since it includes another prime factor 7.

Note that 1 is typically treated as an ugly number.

<https://leetcode.com/problems/ugly-number/>

Solution

This is quite an easy problem!

Listing 74: Solution

```
public class Solution {  
    public boolean isUgly(int num) {  
        if (num == 0){  
            return false;  
5        }  
  
        while (num % 2 == 0){  
            num /= 2;  
        }  
10        while (num % 3 == 0){  
            num /= 3;  
        }  
        while (num % 5 == 0){  
            num /= 5;  
15        }  
  
        if (num == 1){  
            return true;  
        }  
20        return false;  
    }  
}
```

Problem 59

264 Ugly Number II

Write a program to find the n-th ugly number.

Ugly numbers are positive numbers whose prime factors only include 2, 3, 5.
For example, 1, 2, 3, 4, 5, 6, 8, 9, 10, 12 is the sequence of the first 10 ugly numbers.

Note that 1 is typically treated as an ugly number.

Hint:

The naive approach is to call isUgly for every number until you reach the nth one. Most numbers are not ugly.
Try to focus your effort on generating only the ugly ones

An ugly number must be multiplied by either 2, 3, or 5 from a smaller ugly number.

The key is how to maintain the order of the ugly numbers.
Try a similar approach of merging from three sorted lists: L1, L2, and L3.

Assume you have U_k , the kth ugly number. Then U_{k+1} must be $\text{Min}(L1 * 2, L2 * 3, L3 * 5)$.

<https://leetcode.com/problems/ugly-number-ii/>

Solution

Solution is as Hint said.

Listing 75: Solution

```
public class Solution {
    public int nthUglyNumber(int n) {
        int nums[] = new int[n];
        nums[0] = 1;
5       int idx2 = 0 , idx3 = 0, idx5 = 0;
        int count = 0;
        while (count < n - 1){
            int val2 = nums[idx2];
            int val3 = nums[idx3];
10          int val5 = nums[idx5];
            int min = Math.min(Math.min(val2 * 2 , val3 * 3) , val5 * 5);
            if (min == val2 * 2){
                idx2 ++;
            }
15          if (min == val3 * 3){
                idx3 ++;
            }
            if (min == val5 * 5){
                idx5++;
20          }
            count ++;
        }
    }
}
```

```
        nums[count] = min;
    }
    return nums[n - 1];
}
25 }
```

Problem 60

313 Super Ugly Number

Write a program to find the nth super ugly number.

Super ugly numbers are positive numbers whose all prime factors are in the given prime list primes of size k.

For example, [1, 2, 4, 7, 8, 13, 14, 16, 19, 26, 28, 32]
is the sequence of the first 12 super ugly numbers
given primes = [2, 7, 13, 19] of size 4.

Note:

- (1) 1 is a super ugly number for any given primes.
- (2) The given numbers in primes are in ascending order.
- (3) $0 < k \leq 100$, $0 < n \leq 106$, $0 < \text{primes}[i] < 1000$.

<https://leetcode.com/problems/ugly-number-ii/>

Solution

This is just a general circumstance of Problem 264 “Ugly Number II”, so just utilize the same algorithm is fine.

Listing 76: Solution

```
public class Solution {
    public int nthSuperUglyNumber(int n, int[] primes) {
        int[] nums = new int[n];
        nums[0] = 1;
5       int count = 0;
        int plen = primes.length;
        int[] indexes = new int[plen];
        while (count < n - 1){
            int min = Integer.MAX_VALUE;
10         for (int i = 0 ; i < plen ; i++){
                if (min > primes[i] * nums[indexes[i]]){
                    min = primes[i] * nums[indexes[i]];
                }
            }
15         for (int i = 0 ; i < plen ; i++){
                if (min == primes[i] * nums[indexes[i]]){
                    indexes[i] ++;
                }
            }
20         count ++;
            nums[count] = min;
        }
        return nums[n - 1];
25    }
}
```

Problem 61

228 Summary Ranges

Given a sorted integer array without duplicates, return the summary of its ranges.

For example, given [0,1,2,4,5,7], return ["0->2","4->5","7"].

<https://leetcode.com/problems/summary-ranges/>

Solution

An easy problem. Simulation solution works.

Listing 77: Solution

```
public class Solution {
    public List<String> summaryRanges(int[] nums) {
        List<String> result = new ArrayList<String>();
        if (nums.length == 0)
            return result;
        int start = nums[0];
        int end = nums[0];
        int last = nums[0];
        for (int idx = 1 ; idx < nums.length ; idx++){
            int cur = nums[idx];
            if (cur == last + 1){
                end = cur;
                last = cur;
                continue;
            }
            else{
                if (start != end)
                    result.add(String.valueOf(start) + "<->"
                        + String.valueOf(end));
                else
                    result.add(String.valueOf(start));
                start = cur;
                end = cur;
            }
            last = cur;
        }
        if (start != end)
            result.add(String.valueOf(start) + "<->" + String.valueOf(end));
        else
            result.add(String.valueOf(start));
        return result;
    }
}
```

Problem 62

88 Merge Sorted Array

Given two sorted integer arrays `nums1` and `nums2`, merge `nums2` into `nums1` as one sorted array.

Note:

You may assume that `nums1` has enough space

(size that is greater or equal to $m + n$) to hold additional elements from `nums2`.

The number of elements initialized in `nums1` and `nums2` are m and n respectively.

<https://leetcode.com/problems/merge-sorted-array/>

Solution

We can solve the problem in-place(using a new array is too simple, so that I won't put it here.), in order to resist conflict, we can start from the end of array, the worst situation is all elements in `s2` is larger than `s1`, even so there won't be conflicts.

Listing 78: Solution

```
public class Solution {
    public void merge(int[] nums1, int m, int[] nums2, int n) {
        if (m == 0){
            for (int i = 0 ; i < nums1.length ; i++){
                nums1[i] = nums2[i];
            }
        }
        if (n == 0){
            return;
        }
        int idx1 = m - 1;
        int idx2 = n - 1;
        int idx = m + n - 1;
        while (idx >= 0){
            int max;
            if (idx2 < 0){
                max = nums1[idx1--];
            }
            else if (idx1 < 0){
                max = nums2[idx2--];
            }
            else if (nums1[idx1] > nums2[idx2]){
                max = nums1[idx1--];
            }
            else{
                max = nums2[idx2--];
            }
            nums1[idx--] = max;
        }
    }
}
```

Problem 63

279 Perfect Squares

Given a positive integer n , find the least number of perfect square numbers (for example, 1, 4, 9, 16, ...) which sum to n .

For example, given $n = 12$, return 3 because $12 = 4 + 4 + 4$; given $n = 13$, return 2 because $13 = 4 + 9$.

<https://leetcode.com/problems/perfect-squares/>

Solution 1

Tags helps a lot. I found a tag named breadth-first search, then I realize that such combination problem which finds minimum count is perfect for BFS. Add to each element *polled* from queue from biggest available square number, if sum is larger then continue, smaller then *offer* to queue, finally equal to get result. Next step is to try the DP tag...

Listing 79: Solution 1

```
public class Solution {
    public int numSquares(int n) {
        Queue<Integer> counts = new LinkedList<Integer>();
        Queue<Integer> values = new LinkedList<Integer>();
        ArrayList<Integer> squares = new ArrayList<Integer>();
        for (int i = 1 ; i <= n ; i++){
            int square = i * i;
            if (square > n){
                break;
            }
            else{
                if (square == n){
                    return 1;
                }
                squares.add(square);
            }
        }
        for (int i = squares.size() - 1 ; i >= 0 ; i--){
            values.offer(squares.get(i));
            counts.offer(1);
        }
        int count = -1;
        while (!values.isEmpty()){
            int value = values.poll();
            count = counts.poll();

            for(int i = squares.size() - 1 ; i >= 0 ; i--){
                if (value < squares.get(i)){
                    continue;
                }
                if (value + squares.get(i) > n){
                    continue;
                }
            }
        }
    }
}
```

```

35         else if (value + squares.get(i) == n){
            return count + 1;
        }
        else{
            values.offer(value + squares.get(i));
            counts.offer(count + 1);
        }
    }
}
return count;
}
45 }
```

Solution 2

Train of thought is similar to 0-1 knapsack problems. $dp[i] = \min dp[i - square_k] + 1$

Listing 80: Solution 2

```

public class Solution {
    public int numSquares(int n) {
        int[] dp = new int[n + 1];
        dp[0] = 1;
        dp[1] = 1;
        ArrayList<Integer> squares = new ArrayList<Integer>();
        for (int i = 1 ; i <= n ; i++){
            if (i * i > n){
                break;
            }
            else{
                squares.add(i * i);
            }
        }
        for (int i = 2 ; i <= n ; i++){
            int min = Integer.MAX_VALUE;
            for (int j = 0 ; j < squares.size() ; j++){
                if (squares.get(j) > i){
                    break;
                }
                else if (squares.get(j) == i){
                    min = 1;
                    break;
                }
                else{
                    if (dp[i - squares.get(j)] + 1 < min){
                        min = dp[i - squares.get(j)] + 1;
                    }
                }
            }
            dp[i] = min;
        }
        return dp[n];
    }
}
35 }
```


Solution 3

Solution 3 is pretty mathematically. Lagrange's four-square theorem states that every natural number can be represented by the sum of no more than 4 square numbers.

还是利用四平方和定理，但是用到了很多技巧，这个解法是参考网上的答案：根据四平方和定理，任意一个正整数均可表示为4个整数的平方和，其实是可以表示为4个以内的平方数之和，那么就是说返回结果只有1,2,3或4其中的一个，首先我们将数简化一下，由于一个数如果含有因子4，那么我们可以把4都去掉，并不影响结果，比如2和8,3和12等等，返回的结果都相同，读者可自行举更多的栗子。还有一个可以化简的地方就是，如果一个数除以8余7的话，那么肯定是由4个完全平方数组成，这里就不证明了，因为我也不会证明，读者可自行举例验证。那么做完两步后，一个很大的数有可能就会变得很小了，大大减少了运算时间，下面我们就来尝试的将其拆为两个平方数之和，如果拆成功了那么就会返回1或2，因为其中一个平方数可能为0。(注：由于输入的n是正整数，所以不存在两个平方数均为0的情况)。注意下面的!!a + !!b这个表达式，可能很多人不太理解这个意思，其实很简单，感叹号!表示逻辑取反，那么一个正整数逻辑取反为0，再取反为1，所以用两个感叹号!!的作用就是看a和b是否为正整数，都为正整数的话返回2，只有一个为正整数的话返回1，

参见代码如下：

Listing 81: Solution 3

```
class Solution {
public:
    int numSquares(int n) {
        while (n % 4 == 0) n /= 4;
        if (n % 8 == 7) return 4;
        for (int a = 0; a * a <= n; ++a) {
            int b = sqrt(n - a * a);
            if (a * a + b * b == n) {
                return !!a + !!b;
            }
        }
        return 3;
    }
};
```

Problem 64

322 Coin Change

You are given coins of different denominations and a total amount of money amount. Write a function to compute the fewest number of coins that you need to make up that amount.

If that amount of money cannot be made up by any combination of the coins, return -1.

Example 1:

```
coins = [1, 2, 5], amount = 11
return 3 (11 = 5 + 5 + 1)
```

Example 2:

```
coins = [2], amount = 3
return -1.
```

Note:

You may assume that you have an infinite number of each kind of coin.

<https://leetcode.com/problems/coin-change/>

Solution

Knapsack problem. Simple but I'm confused!

Listing 82: Solution

```
public class Solution {
    public int coinChange(int[] coins, int amount) {
        int[] dp = new int[amount + 1];
        for (int i = 0 ; i < dp.length ; i++){
5           dp[i] = Integer.MAX_VALUE;
        }
        dp[0] = 0;
        for (int coinIdx = 0 ; coinIdx < coins.length ; coinIdx++){
            for (int money = coins[coinIdx] ; money <= amount ; money++){
10                if (dp[money - coins[coinIdx]] == Integer.MAX_VALUE){
                    continue;
                }
                dp[money] =
15                    Math.min(dp[money], dp[money - coins[coinIdx]] + 1);
            }
        }

        return dp[amount] == Integer.MAX_VALUE ? -1 : dp[amount];
    }
20 }
```

We can also switch the outer and inner for loop, then compared it with Problem 279 "Perfect Squares". Notice something?

Listing 83: Solution1.2

```
public int coinChange(int[] coins, int amount) {  
    int[] dp = new int[amount + 1];  
    dp[0] = 0;  
    for(int i = 1; i <= amount; i++){  
        dp[i] = Integer.MAX_VALUE;  
        for(int j = 0; j < coins.length; j++){  
            // remember to check i - coins[j] >= 0, or you will get array out of bounds  
            if(i - coins[j] >= 0 && dp[i - coins[j]] != Integer.MAX_VALUE) {  
                dp[i] = Math.min(dp[i], dp[i - coins[j]] + 1);  
            }  
        }  
    }  
    return (dp[amount] == Integer.MAX_VALUE) ? -1 : dp[amount];  
}
```

Problem 65

111 Minimum Depth of Binary Tree

Given a binary tree, find its minimum depth.

The minimum depth is the number of nodes along the shortest path from the root node down to the nearest leaf node.

<https://leetcode.com/problems/minimum-depth-of-binary-tree/>

Solution 1

An easy problem. Solution 1 uses DFS. Don't forget to check the condition when one of the current node's child is null while another isn't.

Listing 84: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public int minDepth(TreeNode root) {
        if (root == null){
            return 0;
        }
        15 if (root.left == null && root.right != null){
            return minDepth(root.right) + 1;
        }
        if (root.left != null && root.right == null){
            return minDepth(root.left) + 1;
        }
        20 return Math.min(minDepth(root.left), minDepth(root.right)) + 1;
    }
}
```

Solution 2

BFS is simpler, faster and more natural. We can stop whenever we first meet a node without children. Edge case is root == null.

Listing 85: Solution 2

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     5   TreeNode left;
 *     TreeNode right;

```

```

    *      TreeNode(int x) { val = x; }
    * }
    */
10 public class Solution {
    public int minDepth(TreeNode root) {
        Queue<TreeNode> queue = new LinkedList<TreeNode>();
        if (root == null){
            return 0;
15     }
        queue.add(root);
        int minDepth = 1;
        while (!queue.isEmpty()){
            ArrayList<TreeNode> list = new ArrayList<TreeNode>();
20             int count = 0;
            while (!queue.isEmpty()){
                list.add(queue.poll());
                count++;
            }
25             for (int i = 0 ; i < count ; i++){
                TreeNode cur = list.get(i);
                if (cur == null || (cur.left == null && cur.right == null)){
                    return minDepth;
                }else if (cur.left == null){
30                 queue.offer(cur.right);
                }else if (cur.right == null){
                    queue.offer(cur.left);
                }else{
35                 queue.offer(cur.left);
                    queue.offer(cur.right);
                }
            }
            minDepth++;
        }
40     return minDepth;
    }
}
```

Problem 66

104 Maximum Depth of Binary Tree

Given a binary tree, find its maximum depth.

The maximum depth is the number of nodes along the longest path from the root node down to the farthest leaf node.

<https://leetcode.com/problems/maximum-depth-of-binary-tree/>

Solution 1

Similar to Problem 104 “Minimum Depth of Binary Tree”. Solution 1 uses DFS. It’s interesting that we don’t need to check if one child is null while another not, since we are fetching the maximum value.

Listing 86: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public int maxDepth(TreeNode root) {
        if (root == null){
            return 0;
        }
15         return Math.max(maxDepth(root.left) , maxDepth(root.right)) + 1;
    }
}
```

Solution 2

When using BFS, se stop when all levels are traversed, namely when queue is empty. Edge case is root == null. Don’t forget to check if input parameter is null (when using BFS)!

Listing 87: Solution 2

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public int maxDepth(TreeNode root) {
```

```
15         if (root == null){
            return 0;
        }
        Queue<TreeNode> queue = new LinkedList<TreeNode>();

        int depth = 0;
        queue.offer(root);

20        while (!queue.isEmpty()){
            depth++;
            List<TreeNode> levelNodeList = new ArrayList<TreeNode>();

            while (!queue.isEmpty()){
25                levelNodeList.add(queue.poll());
            }

            for (TreeNode levelNode : levelNodeList){
                if (levelNode.left != null){
30                    queue.offer(levelNode.left);
                }
                if (levelNode.right != null){
                    queue.offer(levelNode.right);
                }
35            }
        }

        return depth;
40    }
```

Problem 67

215 Kth Largest Element in an Array

Find the kth largest element in an unsorted array.

Note that it is the kth largest element in the sorted order, not the kth distinct element.

For example,

Given [3,2,1,5,6,4] and k = 2, return 5.

Note:

You may assume k is always valid, $1 \leq k \leq \text{array's length}$.

<https://leetcode.com/problems/minimum-depth-of-binary-tree/>

Solution 1

This is a good problem! Solution 1 is great, it utilize half of QuickSort – the partition part. During quicksort, when we decide to put a pivot, the pivot is at exactly where it should be in a sorted array(then numbers on both sides are sorted respectively). So we can check if the index of pivot is what we want - k, or more accurate, `nums.length - k`, if so, the question is answered.

Listing 88: Solution 1.1

```
public class Solution {
    private int partition(int[] nums, int low, int high){
        if (low == high){
            return low;
        }
        int i = low;
        int j = high;
        int pivot = nums[i];
        while (i < j ){
            while (i < j && nums[j] >= pivot){
                j--;
            }
            // actually, this is not needed
            if (i == j){
                break;
            }
            nums[i] = nums[j];
            while(i < j && nums[i] <= pivot){
                i++;
            }
            // actually, this is not needed
            if (i == j){
                break;
            }
            nums[j] = nums[i];
        }
        nums[i] = pivot;
        return i;
    }
}
```



```
30 public int findKthLargest(int[] nums, int k) {  
    if (nums.length == 0 || nums.length < k || k == 0){  
        return -1;  
    }  
35 k = nums.length - k;  
    int low = 0 , high = nums.length - 1;  
    int pivotIdx = partition(nums, low, high);  
    while(low < high && pivotIdx != k){  
        if (pivotIdx > k){  
40             high = pivotIdx - 1;  
        }  
        else{  
            low = pivotIdx + 1;  
        }  
45  
        pivotIdx = partition(nums, low, high);  
    }  
    return nums[pivotIdx];  
50 }
```

I found another solution on Internet which also has a fantastic train of thought.(Divide and Conquer, I almost missed this point.)

In this problem, k means the index, namely how many numbers are there that is equal(itself) and bigger than the element. Then the author partitioned two array, storing elements larger and elements smaller,

when length of array of larger numbers is not adequate($k > \text{bigger.length}$), find in larger set for the element that still have k elements larger than itself;

if $k - 1 < \text{bigger.length}$, then find in the smaller number set, k should be transformed into $k - \text{length} - 1$ (no longer need k bigger numbers).

If $(k - 1 == \text{bigger.length})$, we get the result.

Listing 89: Solution 1.2

```
class Solution {  
public:  
    int findKthLargest(vector<int>& nums, int k) {  
        int length = nums.size();  
5        if (length == 1) {  
            return nums[0];  
        }  
        vector<int> left;  
        vector<int> right;  
10  
        for (int index=1; index<length; index++) {  
            if (nums[index] > nums[0]) {  
                right.push_back(nums[index]);  
            }else{  
15                left.push_back(nums[index]);  
            }  
        }  
        length = right.size();  
        if (length >=k) {
```

```
20         return findKthLargest(right , k);
    } else if(length == k-1){
        return nums[0];
    } else{
        return findKthLargest(left , k-length-1);
25    }
    }
};
```

Solution 2

We can also use Heap sort to solve this problem. Since PriorityQueue is made of heap, we use PriorityQueue to solve the problem. An optimization can be take place to reduce $O(n\log n)$ runtime to $O(n\log k)$, to generate a heap with k elements at first, and insert remaining n - k elements into heap if the element is , finally get the element from the top.

Explain it. If we are finding the kth biggest element, we shall build a heap with small root using the first kth elements, so the root of the heap is exactly the kth biggest element(1st smallest) among all elements. For the rest of them, if the element is larger than top element, it will affect the ranking of elements, so poll the top element and add current element into heap; if element is smaller, just ignore it.

Listing 90: Solution 2

```
public class Solution {
    public int findKthLargest(int[] nums, int k) {
        Queue<Integer> heap = new PriorityQueue<Integer>();
        for (int idx = 0 ; idx < k ; idx++){
5            heap.add(nums[idx]);
        }
        for (int idx = k ; idx < nums.length ; idx++){
            if (nums[idx] <= heap.peek()){
                continue;
10            }
            else{
                heap.poll();
                heap.add(nums[idx]);
            }
15        }
        return heap.peek();
    }
}
```

Problem 68

102 Binary Tree Level Order Traversal

Given a binary tree, return the level order traversal of its nodes' values.
(ie, from left to right, level by level).

For example:

Given binary tree {3,9,20,#,#,15,7},

```

    3
   / \
  9  20
   / \
  15  7

```

return its level order traversal as:

```

[
  [3],
  [9,20],
  [15,7]
]

```

confused what "{1,#,2,3}" means?

OJ's Binary Tree Serialization:

The serialization of a binary tree follows a level order traversal,
where '#' signifies a path terminator where no node exists below.

Here's an example:

```

    1
   / \
  2   3
   /
  4
   \
   5

```

The above binary tree is serialized as "{1,2,3,#,#,4,#,#,5}".

<https://leetcode.com/problems/binary-tree-level-order-traversal/>

Solution 1

A standard BFS Problem.

Listing 91: Solution 1

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {

```

```

    public List<List<Integer>> levelOrder(TreeNode root) {
        if (root == null){
            return new ArrayList<List<Integer>>();
        }
15    List<List<Integer>> result = new ArrayList();
    Queue<TreeNode> queue = new LinkedList<TreeNode>();
    queue.add(root);
    while (!queue.isEmpty()){
        List<Integer> list = new ArrayList<Integer>();
20    List<TreeNode> nodeList = new ArrayList<TreeNode>();
        while (!queue.isEmpty()){
            TreeNode node = queue.poll();
            nodeList.add(node);
        }
25    for (TreeNode node : nodeList){
        list.add(node.val);
        if (node.left != null){
            queue.offer(node.left);
        }
30    if (node.right != null){
        queue.offer(node.right);
        }
    }
    result.add(list);
35    }
    return result;
}
}

```

Solution 2

Use recursion(somewhat DFS) to maintain nodes related to its level.

Listing 92: Solution 2

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    List<List<Integer>> result = new ArrayList<List<Integer>>();

    public void traverse(TreeNode root, int level){
        if (root == null){
15        return;
        }
        if (level == result.size()){
            result.add(new ArrayList<Integer>());
        }
20    result.get(level).add(root.val);
    }
}

```

```
        traverse(root.left , level + 1);
        traverse(root.right , level + 1);
    }
25  public List<List<Integer>> levelOrderBottom(TreeNode root) {
        traverse(root , 0);

        return result;
    }
30 }
```

Problem 69

107 Binary Tree Level Order Traversal II

Given a binary tree, return the level order traversal of its nodes' values.
(ie, from left to right, level by level).

For example:

Given binary tree {3,9,20,#,#,15,7},

```

    3
   / \
  9  20
   / \
  15  7

```

return its bottom-up level order traversal as:

```

[
  [15,7],
  [9,20],
  [3]
]

```

confused what "{1,#,2,3}" means?

OJ's Binary Tree Serialization:

The serialization of a binary tree follows a level order traversal,
where '#' signifies a path terminator where no node exists below.

Here's an example:

```

    1
   / \
  2   3
   /
  4
   \
   5

```

The above binary tree is serialized as "{1,2,3,#,#,4,#,#,5}".

<https://leetcode.com/problems/binary-tree-level-order-traversal-ii/>

Solution

Just add a reversion on list among Problem 102 "Binary Tree Level Order Traversal"

Listing 93: Solution

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {

```

```
List<List<Integer>> result = new ArrayList<List<Integer>>();

public void traverse(TreeNode root, int level){
    if (root == null){
15         return;
    }
    if (level == result.size()){
        result.add(new ArrayList<Integer>());
    }
20     result.get(level).add(root.val);

    traverse(root.left, level + 1);
    traverse(root.right, level + 1);
}
25 public List<List<Integer>> levelOrderBottom(TreeNode root) {
    traverse(root, 0);

    reverse(result);
    return result;
30 }

public void reverse(List<List<Integer>> list){
    int i = 0, j = list.size() - 1;
    while (i < j){
35         List<Integer> obj = list.get(i);
        list.set(i, list.get(j));
        list.set(j, obj);

        i++;
40         j--;
    }
}
}
```

Problem 70

103 Binary Tree Zigzag Level Order Traversal

Given a binary tree, return the zigzag level order traversal of its nodes' values.
(ie, from left to right, then right to left for the next level and alternate between).

For example:

Given binary tree {3,9,20,#,#,15,7},

```

      3
     /\
    9 20
   /\  \
  15 7

```

return its zigzag level order traversal as:

```

[
  [3],
  [20,9],
  [15,7]
]

```

<https://leetcode.com/problems/binary-tree-zigzag-level-order-traversal/>

Solution

Not as complicated as what is said on Internet, similar to Problem 102 “Binary Tree Level Order Traversal”, just to change the sequence of printing...

Listing 94: Solution

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public List<List<Integer>> zigzagLevelOrder(TreeNode root) {
        if (root == null){
            return new ArrayList<List<Integer>>();
        }
15    Queue<TreeNode> queue = new LinkedList<TreeNode>();
    List<List<Integer>> result = new ArrayList<List<Integer>>();
    queue.offer(root);
    int direction = 0;
    while (!queue.isEmpty()){
20        List<TreeNode> list = new ArrayList<TreeNode>();
        List<Integer> values = new ArrayList<Integer>();
        while (!queue.isEmpty()){
            list.add(queue.poll());

```



```
    }  
25  
    if (direction == 0){  
        for (TreeNode node : list){  
            values.add(node.val);  
        }  
30    }else{  
        for (int idx = list.size() - 1; idx >= 0 ; idx --){  
            TreeNode node = list.get(idx);  
            values.add(node.val);  
        }  
35    }  
    for (TreeNode node : list){  
        if (node.left != null){  
            queue.offer(node.left);  
        }  
40        if (node.right != null){  
            queue.offer(node.right);  
        }  
    }  
    direction = 1 - direction;  
45    result.add(values);  
}  
    return result;  
}
```

Problem 71

100 Same Tree

Given two binary trees, write a function to check if they are equal or not.

Two binary trees are considered equal if they are structurally identical and the nodes have the same value.

<https://leetcode.com/problems/same-tree/>

Solution

Similar to Problem 101 “Symmetric Tree”, except for the nodes for comparison.

Listing 95: Solution

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public boolean isSameTree(TreeNode p, TreeNode q) {
        if (p == null && q == null){
            return true;
        }else if (p != null && q != null){
15         if (p.val != q.val){
            return false;
        }
        else{
20         return isSameTree(p.left , q.left )
            && isSameTree(p.right , q.right);
        }
        }else {
            return false;
        }
25     }
}
```

Problem 72

101 Symmetric Tree

Given a binary tree, check whether it is a mirror of itself (ie, symmetric around its center).

For example, this binary tree is symmetric:

```

      1
     /\
    2  2
   /\ /\
  3 4 4 3

```

But the following is not:

```

      1
     /\
    2  2
     \  \
      3   3

```

Note:

Bonus points if you could solve it both recursively and iteratively.

<https://leetcode.com/problems/symmetric-tree/>

Solution 1

Recursion solution is kind of simple (but I still can't figure it out at first sight...).

Listing 96: Solution 1

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public boolean checkSymmetric(TreeNode node1 , TreeNode node2){
        if (node1 == null && node2 == null){
            return true;
        }
15     else if (node1 != null && node2 != null && node1.val == node2.val){
        return checkSymmetric(node1.left , node2.right)
            && checkSymmetric(node1.right , node2.left );
        }
        else{
20         return false;
        }
    }

    public boolean isSymmetric(TreeNode root) {

```

```

25         if (root == null){
                return true;
            }
            return checkSymmetric(root.left , root.right);
        }
30    }

```

Solution 2

BFS solution by myself, given the fact that if the tree is symmetric, the indexes of nodes in pair must sum up to $2^{\text{level}} - 1$, here level is the level in tree, root is 0, children of root is 1... index is its serial number from left to right, index of current node is $2 \times \text{parent}$ or $2 \times \text{parent} + 1$ (depending on whether the node is on left side or right side.)

Listing 97: Solution 2.1

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public boolean isSymmetric(TreeNode root) {
        if (root == null){
            return true;
        }
15        if (root.left == null && root.right == null){
            return true;
        }
        Queue queue = new LinkedList<TreeNode>();
        Queue<Integer> indexes = new LinkedList<Integer>();
20        queue.offer(root);
        indexes.offer(0);
        int level = 0;
        while (!queue.isEmpty()){
            ArrayList<TreeNode> list = new ArrayList<TreeNode>();
            ArrayList<Integer> indexList = new ArrayList<Integer>();
25
            while (!queue.isEmpty()){
                list.add(queue.poll());
                indexList.add(indexes.poll());
            }
30            if (list.size() % 2 == 1 && level != 0){
                return false;
            }
            int i = 0 , j = list.size() - 1;
35            while (i < j){
                if (list.get(i).val != list.get(j).val
                    || indexList.get(i) + indexList.get(j) != (int)Math.pow(2, level) - 1){
                    return false;
                }
            }
            level++;
            while (!list.isEmpty()){
                queue.offer(list.poll());
                indexes.offer(indexList.poll());
            }
        }
        return true;
    }
}

```

```

    }
40
    i++;
    j--;
    }
    for (int idx = 0 ; idx < list.size() ; idx++){
45
        TreeNode node = list.get(idx);
        int index = indexList.get(idx);
        if (node.left != null){
            queue.offer(node.left);
            indexes.offer(index * 2);
50
        }
        if (node.right != null){
            queue.offer(node.right);
            indexes.offer(index * 2 + 1);
        }
55
    }
    level++;
    }
    return true;
    }
60
}

```

Solution on Internet, much simpler than my solution...

Listing 98: Solution 2.2

```

/**
 * Definition for binary tree
 * struct TreeNode {
 *     int val;
5
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
10
class Solution {
public:
    bool isSymmetric(TreeNode *root) {
        // Start typing your C/C++ solution below
        // DO NOT write int main() function
15
        if(root==NULL) return true;
        queue<TreeNode*> lt, rt;
        if(root->left) lt.push(root->left);
        if(root->right) rt.push(root->right);
        TreeNode* l;
20
        TreeNode* r;
        while(!lt.empty() && !rt.empty())
        {
            l = lt.front(); lt.pop();
            r = rt.front(); rt.pop();
25
            if(l == NULL && r == NULL) continue;
            if(l == NULL || r == NULL) return false;
            if(l->val != r->val) return false;

```

```
        lt.push(l->left);
        lt.push(l->right);
30      rt.push(r->right);
        rt.push(r->left);
    }
    if(lt.empty() && rt.empty())
        return true;
35    else
        return false;
    }
};
```

Problem 73

62 Unique Paths

A robot is located at the top-left corner of a $m \times n$ grid (marked 'Start' in the diagram below).

The robot can only move either down or right at any point in time. The robot is trying to reach the bottom-right corner of the grid (marked 'Finish' in the diagram below).

Figure is shown below the description.

How many possible unique paths are there?

Above is a 3×7 grid. How many possible unique paths are there?

Note: m and n will be at most 100.

<https://leetcode.com/problems/unique-paths/>



Solution

This is maybe the 1st dp problem I can solve by myself...

Listing 99: Solution

```
public class Solution {  
    public int uniquePaths(int m, int n) {  
        int dp[][] = new int[m][n];  
  
        for (int i = 0 ; i < m ; i++){  
            dp[i][0] = 1;  
        }  
        for (int i = 0 ; i < n ; i++){  
            dp[0][i] = 1;  
        }  
  
        for (int i = 1 ; i < m ; i++){
```

```
15         for (int j = 1 ; j < n ; j ++){
            dp[i][j] = dp[i - 1][j] + dp[i][j - 1];
        }
        return dp[m - 1][n - 1];
20    }
```


Problem 74

63 Unique Paths II

Follow up for "Unique Paths":

Now consider if some obstacles are added to the grids.
How many unique paths would there be?

An obstacle and empty space is marked as 1 and 0 respectively in the grid.

For example,

There is one obstacle in the middle of a 3x3 grid as illustrated below.

```
[
  [0,0,0],
  [0,1,0],
  [0,0,0]
]
```

The total number of unique paths is 2.

Note: m and n will be at most 100.

<https://leetcode.com/problems/unique-paths-ii/>



Solution

Just add checking process to see if there are obstacles, if so, $dp[i][j] = 0$; Mention that in initialization phase(to initialize 1st row and 1st column), when we first met obstacle, the latter grids are all impassable, then dp value should be 0 too.

Listing 100: Solution

```
public class Solution {
    public int uniquePathsWithObstacles(int [][] obstacleGrid) {
        int m = obstacleGrid.length;
        if (m == 0){
            return 0;
        }
        int n = obstacleGrid[0].length;
```

```
10     if (n == 0){
        return 0;
    }
    int dp[][] = new int[m][n];
    for (int i = 0 ; i < m ; i ++){
        if (obstacleGrid[i][0] == 0){
            dp[i][0] = 1;
15         }
        else{
            break;
        }
    }
    for (int i = 0 ; i < n ; i ++){
        if (obstacleGrid[0][i] == 0){
            dp[0][i] = 1;
20         }
        else{
            break;
25         }
    }
    for (int i = 1 ; i < m ; i ++){
        for (int j = 1 ; j < n ; j ++){
30             dp[i][j] = (obstacleGrid[i][j] == 1) ?
                        0 : dp[i - 1][j] + dp[i][j - 1];
        }
    }
    return dp[m - 1][n - 1];
35 }
```

Problem 75

64 Minimum Path Sum

Given a $m \times n$ grid filled with non-negative numbers,
find a path from top left to bottom right which minimizes
the sum of all numbers along its path.

Note: You can only move either down or right at any point in time.

<https://leetcode.com/problems/minimum-path-sum/>

Solution

A bit more complicated than Problem 62 “Unique Paths”, see code in listing 101

Listing 101: Solution

```
public class Solution {  
    public int minPathSum(int[][] grid) {  
        int m = grid.length;  
        if (m == 0){  
            return 0;  
        }  
        int n = grid[0].length;  
        int dp[][] = new int[m][n];  
        dp[0][0] = grid[0][0];  
        for (int i = 1 ; i < m ; i++){  
            dp[i][0] = grid[i][0] + dp[i - 1][0];  
        }  
        for (int i = 1 ; i < n ; i++){  
            dp[0][i] = grid[0][i] + dp[0][i - 1];  
        }  
        for (int i = 1 ; i < grid.length; i++){  
            for (int j = 1 ; j < grid[0].length ; j++){  
                dp[i][j] = Math.min(dp[i - 1][j] , dp[i][j - 1]) + grid[i][j];  
            }  
        }  
        return dp[m - 1][n - 1];  
    }  
}
```

Problem 76

174 Dungeon Game

The demons had captured the princess (P) and imprisoned her in the bottom-right corner of a dungeon.

The dungeon consists of $M \times N$ rooms laid out in a 2D grid.

Our valiant knight (K) was initially positioned in the top-left room and must fight his way through the dungeon to rescue the princess.

The knight has an initial health point represented by a positive integer.

If at any point his health point drops to 0 or below, he dies immediately.

Some of the rooms are guarded by demons, so the knight loses health (negative integers) upon entering these rooms;

other rooms are either empty (0's) or contain magic orbs that increase the knight's health (positive integers).

In order to reach the princess as quickly as possible, the knight decides to move only rightward or downward in each step.

Write a function to determine the knight's minimum initial health so that he is able to rescue the princess.

For example, given the dungeon below, the initial health of the knight must be at least 7 if he follows the optimal path RIGHT\$->\$ RIGHT \$->\$ DOWN \$->\$ DOWN.

```

-----
|-2(K)|  -3 |  3  |
|-----|
|  -5 |    -10 |  1  |
|-----|
| 10  |    30  |-5(P)|
-----

```

Notes:

The knight's health has no upper bound.

Any room can contain threats or power-ups,

even the first room the knight enters

and the bottom-right room where the princess is imprisoned.

<https://leetcode.com/problems/dungeon-game/>

Solution

Much more complicated than Problems 'Unique Path's. You'll find dynamic programming from top left corner is extremely difficult and confusing. Try start from the bottom right corner to get guaranteed health value. Another thing needs to be carefully taken into consideration is when doing DP calculation, since we are starting from bottom right corner(get guaranteed health), extra positive health are useless, since the Knight cannot respawn with extra health, we still need to guarantee his health of his previous steps, so 0 is

enough.

Listing 102: Solution

```
public class Solution {  
    public int calculateMinimumHP(int [][] dungeon) {  
        int m = dungeon.length;  
        if (m == 0){  
5           return 1;  
        }  
        int n = dungeon[0].length;  
  
        int guaranteedHealth [][] = new int[m][n];  
10        guaranteedHealth[m - 1][n - 1] = (dungeon[m - 1][n - 1] > 0)  
            ? 0 : Math.abs(dungeon[m - 1][n - 1]);  
  
        for (int i = m - 2 ; i >= 0 ; i --){  
            int health = guaranteedHealth[i + 1][n - 1] - dungeon[i][n - 1];  
15            guaranteedHealth[i][n - 1] = Math.max(0, health);  
        }  
        for (int i = n - 2 ; i >= 0 ; i --){  
            int health = guaranteedHealth[m - 1][i + 1] - dungeon[m - 1][i];  
            guaranteedHealth[m - 1][i] = Math.max(0, health);  
20        }  
        for (int i = m - 2 ; i >= 0 ; i --){  
            for (int j = n - 2 ; j >= 0 ; j --){  
                guaranteedHealth[i][j] = Math.max(  
25                    Math.min(guaranteedHealth[i + 1][j]  
                        , guaranteedHealth[i][j + 1]) - dungeon[i][j], 0);  
            }  
        }  
  
        return guaranteedHealth[0][0] + 1;  
30    }  
}
```

Problem 77

232 Implement Queue using Stacks

Implement the following operations of a queue using stacks.

```
push(x) -- Push element x to the back of queue.
pop() -- Removes the element from in front of queue.
peek() -- Get the front element.
empty() -- Return whether the queue is empty.
```

Notes:

- You must use only standard operations of a stack
 - which means only push to top, peek/pop from top, size, and is empty operations are valid.

Depending on your language, stack may not be supported natively. You may simulate a stack by using a list or deque (double-ended queue), as long as you use only standard operations of a stack.

You may assume that all operations are valid (for example, no pop or peek operations will be called on an empty queue).

<https://leetcode.com/problems/implement-queue-using-stacks/>

Solution

Use 2 containers to implement another kind of container, just see listings.

Listing 103: Solution

```
class MyQueue {
    Stack<Integer> stk1 = new Stack<Integer>();
    Stack<Integer> stk2 = new Stack<Integer>();
    // Push element x to the back of queue.
5   public void push(int x) {
        stk1.push(x);
    }

    // Removes the element from in front of queue.
10  public void pop() {
        if (stk2.isEmpty()){
            while(!stk1.isEmpty()){
                stk2.push(stk1.pop());
            }
15        }
        stk2.pop();
    }

    // Get the front element.
20  public int peek() {
        if (stk2.isEmpty()){
            while(!stk1.isEmpty()){
                stk2.push(stk1.pop());
            }
        }
        return stk2.peek();
    }
}
```

```
25         }  
        }  
        return stk2.peek();  
    }  
  
30    // Return whether the queue is empty.  
    public boolean empty() {  
        return (stk1.isEmpty() && stk2.isEmpty());  
    }  
}
```

Problem 78

225 Implement Stack using Queues

Implement the following operations of a stack using queues.

`push(x)` -- Push element `x` onto stack.
`pop()` -- Removes the element on top of the stack.
`top()` -- Get the top element.
`empty()` -- Return whether the stack is empty.

Notes:

You must use only standard operations of a queue
-- which means only push to back, peek/pop from front,
size, and is empty operations are valid.

Depending on your language, queue may not be supported natively.
You may simulate a queue by using a list or deque (double-ended queue),
as long as you use only standard operations of a queue.

You may assume that all operations are valid
(for example, no pop or top operations will be called on an empty stack).

<https://leetcode.com/problems/implement-queue-using-stacks/>

Solution

See listings.

Listing 104: Solution

```
class MyStack {
    Queue<Integer> queue1 = new LinkedList<Integer>();
    Queue<Integer> queue2 = new LinkedList<Integer>();

    // Push element x onto stack.
    public void push(int x) {
        queue2.offer(x);

        while (!queue1.isEmpty()){
            queue2.offer(queue1.poll());
        }

        while (!queue2.isEmpty()){
            queue1.offer(queue2.poll());
        }
    }

    // Removes the element on top of the stack.
    public void pop() {
        queue1.poll();
    }
}
```



```
25      // Get the top element.  
      public int top() {  
          return queue1.peek();  
      }  
  
      // Return whether the stack is empty.  
30      public boolean empty() {  
          return (queue1.isEmpty() && queue2.isEmpty());  
      }  
    }
```

Problem 79

148 Sort List

Sort a linked list in $O(n \log n)$ time using constant space complexity.

<https://leetcode.com/problems/sort-list/>

Solution

First, review all kinds of sorting methods!

类别	排序方法	时间复杂度			空间复杂度	稳定性
		平均情况	最好情况	最坏情况	辅助存储	
插入排序	直接插入	$O(n^2)$	$O(n)$	$O(n^2)$	$O(1)$	稳定
	Shell排序	$O(n^{1.3})$	$O(n)$	$O(n^2)$	$O(1)$	不稳定
选择排序	直接选择	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(1)$	不稳定
	堆排序	$O(n \log_2 n)$	$O(n \log_2 n)$	$O(n \log_2 n)$	$O(1)$	不稳定
交换排序	冒泡排序	$O(n^2)$	$O(n)$	$O(n^2)$	$O(1)$	稳定
	快速排序	$O(n \log_2 n)$	$O(n \log_2 n)$	$O(n^2)$	$O(n \log_2 n)$	不稳定
归并排序		$O(n \log_2 n)$	$O(n \log_2 n)$	$O(n \log_2 n)$	$O(n)$	稳定
基数排序		$O(d(r+n))$	$O(d(n+rd))$	$O(d(r+n))$	$O(rd+n)$	稳定

注：基数排序的复杂度中， r 代表关键字的基数， d 代表长度， n 代表关键字的个数。

We can see that Heap sort, Quick sort and Merge sort satisfy the requirements of problem. They all need extra space at first sight, which may not meet the space requirement. However, when we look at the attribute of LinkedList, we can find out that we do not need to allocate extra space to store the sorted list after merging, what we need to do is just to reconnect the nodes.

Listing 105: Solution

```

/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode sortList(ListNode head) {
        if (head == null || head.next == null){
            return head;
        }
        ListNode slow = head;
        ListNode fast = head;
        ListNode prev = slow;
        while (fast != null && fast.next != null){
            prev = slow;
            slow = slow.next;
            fast = fast.next.next;
        }
    }
}

```

```
        prev.next = null;
        ListNode l1 = sortList(head);
        ListNode l2 = sortList(slow);
25      return mergeList(l1, l2);
    }

    public ListNode mergeList(ListNode l1, ListNode l2){
        ListNode dummy = new ListNode(-1);
30      ListNode cur = dummy;
        ListNode cur1 = l1, cur2 = l2;
        while (cur1 != null && cur2 != null){
            if (cur1.val < cur2.val){
35                cur.next = cur1;
                cur1 = cur1.next;
                cur = cur.next;
            }
            else{
40                cur.next = cur2;
                cur2 = cur2.next;
                cur = cur.next;
            }
        }
        while (cur1 != null){
45            cur.next = cur1;
            cur = cur.next;
            cur1 = cur1.next;
        }
        while (cur2 != null){
50            cur.next = cur2;
            cur = cur.next;
            cur2 = cur2.next;
        }
        cur.next = null;
55      return dummy.next;
    }
}
```

Problem 80

147 Insertion Sort List

Sort a linked list using insertion sort.

<https://leetcode.com/problems/insertion-sort-list/>

Solution

It's an easy but a complicated problem which is not easy to be bug-free.

Listing 106: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode insertionSortList(ListNode head) {
        if(head == null || head.next == null){
            return head;
        }
        ListNode dummy = new ListNode(Integer.MIN_VALUE);
        dummy.next = head;
        ListNode curNode = head.next;
        ListNode prevNode = head;

        while (curNode != null){
            // find insertion position
            boolean found = false;
            ListNode searchNode = dummy.next;
            ListNode searchPrevNode = dummy;
            while (searchNode != curNode){
                if (curNode.val < searchNode.val){
                    // reconnect curNode
                    prevNode.next = curNode.next;

                    // connect searchNode & currentNode
                    searchPrevNode.next = curNode;
                    curNode.next = searchNode;

                    curNode = prevNode.next;
                    found = true;
                    break;
                }
                searchPrevNode = searchNode;
                searchNode = searchNode.next;
            }
            if (curNode == searchNode){
                // no need to change
            }
        }
    }
}
```

```
        prevNode = curNode;
        curNode = curNode.next;
        continue;
45     }

    // insert node
    if (found){
        continue;
50     }
    prevNode = curNode;
    curNode = curNode.next;
}

55     return dummy.next;
}
}
```

Problem 81

86 Partition List

Given a linked list and a value x , partition it such that all nodes less than x come before nodes greater than or equal to x .

You should preserve the original relative order of the nodes in each of the two partitions.

For example,

Given $1 \rightarrow 4 \rightarrow 3 \rightarrow 2 \rightarrow 5 \rightarrow 2$ and $x = 3$,
return $1 \rightarrow 2 \rightarrow 2 \rightarrow 4 \rightarrow 3 \rightarrow 5$.

<https://leetcode.com/problems/partition-list/>

Solution

Maintain 2 List, one for elements less than x , one for elements larger or equal to x , finally reconnect them.

Listing 107: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode partition(ListNode head, int x) {
        if (head == null || head.next == null){
            return head;
        }

        ListNode dummySmall = new ListNode(-1);
        ListNode dummyLarge = new ListNode(-1);

        ListNode tailSmall = dummySmall;
        ListNode tailLarge = dummyLarge;

        ListNode curNode = head;
        while (curNode != null){
            if (curNode.val < x){
                tailSmall.next = curNode;
                tailSmall = tailSmall.next;
            }
            else if (curNode.val >= x){
                tailLarge.next = curNode;
                tailLarge = tailLarge.next;
            }

            curNode = curNode.next;
        }
    }
}
```

```
35         tailLarge.next = null;

        tailSmall.next = dummyLarge.next;

        return dummySmall.next;
40    }
```

Problem 82

136 Single Number

Given an array of integers, every element appears twice except for one.
Find that single one.

<https://leetcode.com/problems/single-number/>

Solution 1

Review the whole set of “Single number(I/II/III)”, an interesting set of problems.

First solution is to use a HashSet, if the element appears once, add it into Set; if it has appeared, remove it from Set. Finally the only element remains in Set is the ‘Single Number’.

Listing 108: Solution 1

```
public class Solution {  
    public int singleNumber(int[] nums) {  
        Set<Integer> set = new HashSet<Integer>();  
        for (int num : nums){  
            if (set.contains(num)){  
                set.remove(num);  
            }  
            else{  
                set.add(num);  
            }  
        }  
        Iterator<Integer> ite = set.iterator();  
        if (ite.hasNext()){  
            return ite.next();  
        }  
        else{  
            return 0;  
        }  
    }  
}
```

Solution 2

Bit manipulation is more proper to this SET of problems. Xor is the key to this problem. Apply xor to 2 same numbers will get the result 0, and since xor is commutative($a \text{ xor } b = b \text{ xor } a$) and associative($a \text{ xor } b \text{ xor } c = a \text{ xor } (b \text{ xor } c)$), xor all numbers will get the single number!

In fact, solution in next Problem “Single Number II” is more proper for this SET of problems.

Listing 109: Solution 2

```
public class Solution {  
    public int singleNumber(int[] nums) {  
        if (nums.length == 1){  
            return nums[0];  
        }  
        int number = nums[0];  
        for (int idx = 1; idx < nums.length ; idx++){  
            number ^= nums[idx];  
        }  
    }  
}
```


10

```
        }  
        return number;  
    }  
}
```

Problem 83

137 Single Number II

Given an array of integers, every element appears three times except for one. Find that single one.

Note:

Your algorithm should have a linear runtime complexity.

Could you implement it without using extra memory?

<https://leetcode.com/problems/single-number-ii/>

Solution

Solution 1 of Problem 136 “Single Number” is no longer suitable for this problem. Bit manipulation still works, but this time a bit complicated.

E.g.: array consists of 1,2,2,2,3,3,3, namely 001 010 010 010 011 011 011, for units digit, there are four 1s, so the 1 is from the single number... Follow such train of thought, we can solve this problem.

Listing 110: Solution

```
public class Solution {
    public int singleNumber(int[] nums) {
        int result = 0;
        for (int i = 0 ; i < 32 ; i ++){
            int digit = 0;
            int mask = (1 << i);
            for(int idx = 0 ; idx < nums.length ; idx ++){
                if ((nums[idx] & mask) != 0){
                    digit ++;
                }
            }
            if (digit % 3 != 0){
                result |= mask;
            }
        }
        return result;
    }
}
```

Problem 84

260 Single Number III

Given an array of numbers `nums`, in which exactly two elements appear only once and all the other elements appear exactly twice. Find the two elements that appear only once.

For example:

Given `nums = [1, 2, 1, 3, 2, 5]`, return `[3, 5]`.

Note:

The order of the result is not important. So in the above example, `[5, 3]` is also correct.

Your algorithm should run in linear runtime complexity.

Could you implement it using only constant space complexity?

<https://leetcode.com/problems/single-number-iii/>

Solution

The key to solve this problem is to separate the two single numbers into 2 different groups so that we can apply algorithm from 'Single Number' to each group.

How to achieve that? First we xor all numbers to get a result from which we can conclude a result that when the digit of the result is 1, it means that 2 single numbers differ among this digit, so that in terms of all numbers, the number of 1 should be odd number (consisting 1 single number), so that we can group numbers with digit valued 1 and digit valued 0 on that digit into 2 groups. Finally apply xor algorithm to get the single number from each group.

Listing 111: Solution

```
public class Solution {
    public int[] singleNumber(int[] nums) {
        int total = 0;
        int count1 = 0, count2 = 0;
5      for (int num : nums){
            total ^= num;
        }
        int mDigit = -1;
        for (int i = 0 ; i < 32 ; i++){
10      int mask = (1 << i);
            if ((total & mask) != 0){
                // find the digit
                mDigit = i;
                for (int j = 0 ; j < nums.length ; j++){
15      if ((nums[j] & mask) != 0){
                    count1 ++;
                }
                else{
                    count2 ++;
20      }
                }
            }
        }
```

```
                break;
            }
25     }
    int [] g1 = new int [count1];
    int [] g2 = new int [count2];
    int idx1 = 0 , idx2 = 0;
    int mask = (1 << mDigit);
30     for (int i = 0 ; i < nums.length ; i++){
        if ((nums[i] & mask) != 0){
            g1[idx1] = nums[i];
            idx1 ++;
        }else{
35             g2[idx2] = nums[i];
            idx2 ++;
        }
    }
    int num1 = 0 , num2 = 0;
40     for (int num : g1){
        num1 ^= num;
    }
    for (int num : g2){
        num2 ^= num;
45     }
    int [] result = new int [2];
    result [0] = num1;
    result [1] = num2;
    return result;
50 }
}
```

Problem 85

282 Expression Add Operators

Given a string that contains only digits 0-9 and a target value, return all possibilities to add binary operators (not unary) +, -, or * between the digits so they evaluate to the target value.

Examples:

"123", 6 -> ["1+2+3", "1*2*3"]
"232", 8 -> ["2*3+2", "2+3*2"]
"105", 5 -> ["1*0+5", "10-5"]
"00", 0 -> ["0+0", "0-0", "0*0"]
"3456237490", 9191 -> []

<https://leetcode.com/problems/expression-add-operators/>

Solution

A good DFS problem.

The difficulties focus on those edge cases(String startsWith 0 - 01 is illegal; length of remaining string) and how to deal with multiply operation. In this solution uses a 'diff' to save previous result, for example, 1 - 2 * 3, previous result should be -1(1 - 2), but 1-2*3=-5, which means we should apply -1 - (-2) + (-2 *3), here diff = -2, namely previous number.

However, when a new problem arise, I still can't get the solution method!

Listing 112: Solution

```
public class Solution {
    List<String> result = new ArrayList<String>();

    public List<String> addOperators(String num, int target) {
        dfs(num, 0, target, 0, "");
        return result;
    }

    public void dfs(String remaining, long cur, int target, long diff,
        String expr){
        if (remaining.length() == 0 && target == cur){
            result.add(expr);

            return;
        }

        for (int i = 1; i <= remaining.length() ; i++){
            String sub = remaining.substring(0, i);
            // important
            if (sub.length() > 1 && sub.charAt(0) == '0'){
                return;
            }
            String curRem = remaining.substring(i);
            long subNum = Long.parseLong(sub);

            if (expr.length() == 0){
```

```
        dfs(curRem, subNum, target, subNum, sub);
    }
    else{
30      dfs(curRem, cur + subNum, target, subNum, expr + "+" + sub);
        dfs(curRem, cur - subNum, target, 0 - subNum, expr + "-" + sub);
        dfs(curRem, cur - diff + diff * subNum, target,
35           diff * subNum, expr + "*" + sub);
    }
}
}
```

Problem 86

200 Number of Islands

Given a 2d grid map of '1's (land) and '0's (water), count the number of islands.

An island is surrounded by water and is formed by connecting adjacent lands horizontally or vertically.

You may assume all four edges of the grid are all surrounded by water.

Example 1:

```
11110
11010
11000
00000
Answer: 1
```

Example 2:

```
11000
11000
00100
00011
Answer: 3
```

<https://leetcode.com/problems/number-of-islands/>

Solution

It's just a simple DFS problem...

Listing 113: Solution

```
public class Solution {
    public int numIslands(char[][] grid) {
        int count = 0;
        for (int i = 0 ; i < grid.length ; i++){
            for (int j = 0 ; j < grid[0].length ; j++){
                if (grid[i][j] == '1'){
                    wander(grid, i, j);
                    count++;
                }
            }
        }
        return count;
    }

    public void wander(char[][] grid, int i, int j){
        if (i < 0 || j < 0 || i >= grid.length || j >= grid[0].length){
            return;
        }
        if (grid[i][j] == '1'){
            grid[i][j] = 'X';
            wander(grid, i, j + 1);
        }
    }
}
```

```
25         wander(grid , i + 1 , j);  
            wander(grid , i - 1 , j);  
            wander(grid , i , j - 1);  
        }  
    }  
}
```


Problem 87

287 Find the Duplicate Number

Given an array `nums` containing $n + 1$ integers where each integer is between 1 and n (inclusive),
prove that at least one duplicate number must exist.
Assume that there is only one duplicate number, find the duplicate one.

Note:

You must not modify the array (assume the array is read only).

You must use only constant, $O(1)$ extra space.

Your runtime complexity should be less than $O(n^2)$.

There is only one duplicate number in the array,
but it could be repeated more than once.

<https://leetcode.com/problems/find-the-duplicate-number/>

Solution 1

A pretty interesting problem.

There is a good blog which shows all kinds of solutions for this problem,

see <https://segmentfault.com/a/1190000003817671>.

Binary Search is available with its $O(n \log n)$ runtime and $O(1)$ space, with the idea to binary search for the value from 1 to n to match numbers in array, instead of searching in the array itself.

Listing 114: Solution 1

```
public class Solution {  
    public int findDuplicate(int[] nums) {  
        int min = 1, max = nums.length - 1;  
        while (min <= max){  
5             int mid = (min + max) / 2;  
  
            int count = 0;  
            int countN = 0;  
            for (int num : nums){  
10                 if (num <= mid){  
                     count ++;  
  
                     if (num == mid){  
15                         countN ++;  
                     }  
                }  
            }  
  
            if (countN >= 2){  
20                 return mid;  
            }  
  
            if (count > mid){  
                max = mid;  
25            }else{  
                min = mid + 1;  
            }  
        }  
    }  
}
```

```
        }
    }
30     return -1;
    }
}
```

Solution 2

An $O(n)$ runtime, $O(1)$ space solution, fantastic train of thought!

See <http://keithschwarz.com/interesting/code/?dir=find-duplicate>

Main idea is like “Find Linked List Cycle II”, use ‘slow = array[slow], fast = array[array[fast]]’ mappings to check if there is a cycle and get the entry. The entry is exactly what we are looking for! (Since there is a $\text{array}[i] = \text{array}[j]$, where $i \neq j$).

Listing 115: Solution 2

```
public class Solution {
    public int findDuplicate(int[] nums) {
        int slow = 0;
        int fast = 0;
5
        do{
            slow = nums[slow];
            fast = nums[nums[fast]];
        } while (slow != fast);
10
        fast = 0;
        int prev = 0;

        while(slow != fast){
15
            prev = slow;
            slow = nums[slow];
            fast = nums[fast];
        }

20
        return slow;
        //return nums[prev];
        // Actually, we can return slow directly.
    }
}
```

Problem 88

268 Missing Number

Given an array containing n distinct numbers taken from $0, 1, 2, \dots, n$, find the one that is missing from the array.

For example,
Given `nums = [0, 1, 3]` return `2`.

Note:
Your algorithm should run in linear runtime complexity.
Could you implement it using only constant extra space complexity?

<https://leetcode.com/problems/missing-number/>

Solution - Using Sum

Since there is only one missing element, the different between sum of the array and 0 to n is exactly what is missing.

Binary search similar to Problem “Find the Duplicate Number” is also proper for this problem. (Remember to sort the array first if we don’t want to calculate the count.)

Listing 116: Solution 1

```
public class Solution {  
    public int missingNumber(int[] nums) {  
        int sumExp = 0;  
        int sumAct = 0;  
5         int n = nums.length;  
        for (int i = 0 ; i <= n ; i++){  
            sumExp += i;  
        }  
        for (int num : nums){  
10         sumAct += num;  
        }  
        return sumExp - sumAct;  
    }  
}
```

Solution - Mark Position

Solution 2 is to set up a new array, with a length more than 1 element of former array. In the new array, the elements are the same of (index + 1). Every time we find a number in old array, we set element in corresponding position in new array -1. Finally we iterate the new array, the element which is not changed to -1 is what we want.

Some similar solutions including Using Hashset or sorted array to store the count of each number to see if there is a non-existing element.

Listing 117: Solution 2

```
public class Solution {  
    public int missingNumber(int[] nums) {  
        int[] nums2 = new int[nums.length + 1];  
        for (int i = 0; i < nums2.length; i++) {
```

```
5         nums2[i] = i;
        }
        for (int i = 0; i < nums.length; i++) {
            nums2[nums[i]] = -1;
        }
10        for (int i = 0; i < nums2.length; i++) {
            if (nums2[i] != -1) {
                return i;
            }
        }
15        return -1;
    }
}
```

Solution - Bit Manipulation

This solution is interesting. See the steps first.

1. XOR all the array elements, let the result of XOR be X1.
2. XOR all numbers from 1 to n, let XOR be X2.
3. XOR of X1 and X2 gives the missing number.

Feel familiar? The insight of these steps is exactly like “Single Number”, or to be more specific, the characteristic of XOR operator.

Listing 118: Solution 3

```
#include<stdio.h>
/* getMissingNo takes array and size of array as arguments*/
int getMissingNo(int a[], int n)
{
5     int i;
    int x1 = a[0]; /* For xor of all the elements in array */
    int x2 = 1; /* For xor of all the elements from 1 to n+1 */
    for (i = 1; i < n; i++)
        x1 = x1 ^ a[i];
10    for (i = 2; i <= n+1; i++)
        x2 = x2 ^ i;

    return (x1 ^ x2);
}
15 /*program to test above function */
int main()
{
    int a[] = {1, 2, 4, 5, 6};
    int miss = getMissingNo(a, 5);
20    printf("%d", miss);
    getchar();
}
```

Problem 89

41 First Missing Positive

Given an unsorted integer array, find the first missing positive integer.

For example,

Given [1,2,0] return 3,

and [3,4,-1,1] return 2.

Your algorithm should run in $O(n)$ time and uses constant space.

<https://leetcode.com/problems/first-missing-positive/>

Solution

This is somewhat like the idea of radix sort, to put each element to where it should be ($\text{arr}[\text{idx}] = \text{idx}$), then the 1st grid which doesn't satisfy this attribute is exactly the number we want.

Listing 119: Solution 1.1

```
public class Solution {
    public int firstMissingPositive(int[] nums) {
        for (int i = 0 ; i < nums.length ; i++){
            nums[i] --;
        }
        int idx = 0;
        while (idx < nums.length){
            if (nums[idx] >= nums.length || nums[idx] < 0){
                idx ++;
            }
            else{
                if (nums[idx] == nums[nums[idx]] || nums[idx] == idx){
                    idx ++;
                }
                else{
                    int swap = nums[nums[idx]];
                    nums[nums[idx]] = nums[idx];
                    nums[idx] = swap;
                }
            }
        }

        for (idx = 0 ; idx < nums.length ; idx++){
            if (nums[idx] != idx){
                return idx + 1;
            }
        }
        return nums.length + 1;
    }
}
```

Below is a solution without reducing 1 to all elements.

Listing 120: Solution 1.2

```
public class Solution {  
    public int firstMissingPositive(int[] nums) {  
        int idx = 0;  
        while (idx < nums.length){  
            int value = nums[idx] - 1;  
            if (value >= nums.length || value < 0){  
                idx ++;  
            }  
            else{  
                if (value + 1 == nums[value] || nums[idx] == idx + 1){  
                    idx ++;  
                }  
                else{  
                    nums[idx] = nums[value];  
                    nums[value] = value + 1;  
                }  
            }  
        }  
  
        for (idx = 0 ; idx < nums.length ; idx ++){  
            if (nums[idx] != idx + 1){  
                return idx + 1;  
            }  
        }  
        return nums.length + 1;  
    }  
}
```

Problem 90

98 Validate Binary Search Tree

Given a binary tree, determine if it is a valid binary search tree (BST).

Assume a BST is defined as follows:

The left subtree of a node contains
only nodes with keys less than the node's key.
The right subtree of a node contains
only nodes with keys greater than the node's key.
Both the left and right subtrees must
also be binary search trees.

<https://leetcode.com/problems/validate-binary-search-tree/>

Solution 1

Plenty of methods to solve the problem. First solution is to use recursion to solve the problem. However, code on “Clean Code Handbook” is extremely clean and charming!

Listing 121: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public boolean validBST(TreeNode root, Integer min, Integer max){

        if (root == null){
            return true;
15        }
        if (min != null){
            if (root.val <= min){
                return false;
            }
20        }
        if (max != null){
            if (root.val >= max){
                return false;
            }
25        }

        return validBST(root.left, min, root.val)
            && validBST(root.right, root.val, max);
    }
30

    public boolean isValidBST(TreeNode root) {
```

```
        return validBST(root , null , null);  
    }  
35 }
```

Solution 2

A valid binary search tree has the attribute to keep ascending order when traversing the tree in an inorder order.

Listing 122: Solution 2

```
/**  
 * Definition for a binary tree node.  
 * public class TreeNode {  
 *     int val;  
5  *     TreeNode left;  
 *     TreeNode right;  
 *     TreeNode(int x) { val = x; }  
 * }  
 */  
10 public class Solution {  
    public void inorderTraverse(TreeNode root , List<Integer> traversal){  
        if (root == null){  
            return;  
        }  
15  
        inorderTraverse(root.left , traversal);  
        traversal.add(root.val);  
        inorderTraverse(root.right , traversal);  
    }  
20  
    public boolean isValidBST(TreeNode root) {  
        List<Integer> traversal = new ArrayList<Integer>();  
        inorderTraverse(root , traversal);  
        for (int idx = 1 ; idx < traversal.size() ; idx ++){  
25            if (traversal.get(idx) <= traversal.get(idx - 1)){  
                return false;  
            }  
        }  
30  
        return true;  
    }  
}
```


Problem 91

150 Evaluate Reverse Polish Notation

Evaluate the value of an arithmetic expression in Reverse Polish Notation.

Valid operators are +, -, *, /.

Each operand may be an integer or another expression.

Some examples:

`["2", "1", "+", "3", "*"]` \rightarrow $((2 + 1) * 3)$ \rightarrow 9

`["4", "13", "5", "/", "+"]` \rightarrow $(4 + (13 / 5))$ \rightarrow 6

<https://leetcode.com/problems/evaluate-reverse-polish-notation/>

Solution 1

Reverse Polish expression evaluation is a standard stack problem, normally use stack to solve this.

Listing 123: Solution 1

```
public class Solution {
    public boolean isOperator(String token){
        if (token.equals("+") || token.equals("-") || token.equals("*")
            || token.equals("/")){
5             return true;
        }
        else{
            return false;
        }
10    }

    public int calculate(int num1, int num2, String token){
        if (token.equals("+")){
            return num1 + num2;
15        }
        if (token.equals("-")){
            return num1 - num2;
        }
        if (token.equals("*")){
20            return num1 * num2;
        }
        if (token.equals("/")){
            return num1 / num2;
        }
25        return -1;
    }

    public int evalRPN(String[] tokens) {
30        Stack<String> stack = new Stack<String>();
        for (String token : tokens){
            if (isOperator(token)){
                int num1 = Integer.parseInt(stack.pop());
                int num2 = Integer.parseInt(stack.pop());
```

```

35         int result = calculate(num2, num1, token);

        stack.push(String.valueOf(result));
    }
    else{
40         stack.push(token);
    }
}

    return Integer.parseInt(stack.pop());
45
}
}

```

Solution 2

However, when we want to add a new operator, the structure now is not so extensive since we should modify two parts of program. Use a Map to extend the program(It's a solution followed in "Clean Code Handbook"):

Listing 124: Solution 2

```

interface Operator{
    public int eval(int a , int b);
}

5 public class Solution {
    private static final HashMap<String , Operator> cal
    = new HashMap<String , Operator>(){
        {
            put("+", new Operator(){
10                public int eval(int a , int b){
                    return a + b;
                }
            });
            put("-", new Operator(){
15                public int eval(int a , int b){
                    return a - b;
                }
            });
            put("*", new Operator(){
20                public int eval(int a , int b){
                    return a * b;
                }
            });
            put("/", new Operator(){
25                public int eval(int a , int b){
                    return a / b;
                }
            });
        }
    };
30
}

```

```
public int evalRPN(String [] tokens) {  
    Stack<Integer> stack = new Stack<Integer>();  
35    for (String token : tokens){  
        if (cal.containsKey(token)){  
            int num2 = stack.pop();  
            int num1 = stack.pop();  
            int result = cal.get(token).eval(num1, num2);  
40  
            stack.push(result);  
        }  
        else{  
            stack.push(Integer.parseInt(token));  
45        }  
    }  
  
    return stack.pop();  
50 }  
}
```

Problem 92

130 Surrounded Regions

Given a 2D board containing 'X' and 'O', capture all regions surrounded by 'X'.

A region is captured by flipping all 'O's into 'X's in that surrounded region.

For example,

```
X X X X
```

```
X O O X
```

```
X X O X
```

```
X O X X
```

After running your function, the board should be:

```
X X X X
```

```
X X X X
```

```
X X X X
```

```
X O X X
```

<https://leetcode.com/problems/surrounded-regions/>

Solution

DFS will get StackOverflowException, so use BFS to solve the problem.

(Honestly, I don't know why my DFS solution StackOverflow?!)

Listing 125: Solution

```
class Pair{
    public Pair(int i , int j){
        this.i = i;
        this.j = j;
    }

    public int i, j;
}

public class Solution {
    public void escape(char[][] board, int i , int j){

        int m = board.length;
        int n = board[0].length;

        if (board[i][j] == 'O'){
            Queue<Pair> queue = new LinkedList<Pair>();

            queue.offer(new Pair(i, j));
            while(!queue.isEmpty()){
                Pair pair = queue.poll();
                if (board[pair.i][pair.j] != 'O'){
                    continue;
                }
                else{
                    board[pair.i][pair.j] = 'E';
                }
            }
        }
    }
}
```

```

        if (pair.i < m - 1 && board[pair.i + 1][pair.j] == 'O'){
            queue.offer(new Pair(pair.i + 1 , pair.j));
        } if (pair.j < n - 1 && board[pair.i][pair.j + 1] == 'O'){
            queue.offer(new Pair(pair.i , pair.j + 1));
        } if (pair.i > 0 && board[pair.i - 1][pair.j] == 'O'){
            queue.offer(new Pair(pair.i - 1 , pair.j));
        } if (pair.j > 0 && board[pair.i][pair.j - 1] == 'O'){
            queue.offer(new Pair(pair.i , pair.j - 1));
        }
    }
}

public void solve(char[][] board) {
    int m = board.length;
    if (m == 0){
        return;
    }
    int n = board[0].length;
    for (int i = 0 ; i < m ; i++){
        if (board[i][0] == 'O'){
            escape(board, i , 0);
        }
        if (board[i][n - 1] == 'O'){
            escape(board, i , n - 1);
        }
    }

    for (int i = 0 ; i < n ; i++){
        if (board[0][i] == 'O'){
            escape(board, 0 , i);
        }
        if (board[m - 1][i] == 'O'){
            escape(board, m - 1 , i);
        }
    }

    for (int i = 0 ; i < m ; i++){
        for (int j = 0 ; j < n ; j++){
            if (board[i][j] == 'O'){
                board[i][j] = 'X';
            }
            else if (board[i][j] == 'E'){
                board[i][j] = 'O';
            }
        }
    }
}
}
```

Problem 93

179 Largest Number

Given a list of non negative integers, arrange them such that they form the largest number.

For example, given [3, 30, 34, 5, 9], the largest formed number is 9534330.

Note: The result may be very large, so you need to return a string instead of an integer.

<https://leetcode.com/problems/largest-number/>

Solution

Re-design the compare to judge two strings(A tactic is when we want to get bigger one of s1 and s2, check relation between s1s2 and s2s1, e.g. '93' > '39' so '9' > '3')

Don't forget the condition when result consists several 0s!

Listing 126: Solution

```
public class Solution {
    public void quickSort(int[] nums, int low, int high){
        if (low < 0 || high > nums.length - 1){
            return ;
        }
        if (low > high){
            return ;
        }
        int pivot = nums[low];
        int i = low;
        int j = high;

        while (i < j){
            while (i < j && noLessThan(nums[j], pivot)){
                j --;
            }
            nums[i] = nums[j];
            while (i < j && noLessThan(pivot, nums[i])){
                i ++;
            }
            nums[j] = nums[i];
        }

        nums[i] = pivot;
        quickSort(nums, low, i - 1);
        quickSort(nums, i + 1, high);
    }

    public String largestNumber(int[] nums) {
        quickSort(nums, 0, nums.length - 1);

        StringBuilder sb = new StringBuilder();
```

```

    int idx = nums.length - 1;
35  while (idx >= 0 && nums[idx] == 0){
        idx--;
    }

    while(idx >= 0){
40      sb.append(String.valueOf(nums[idx]));
        idx--;
    }

45  return sb.toString().length() > 0 ? sb.toString() : "0";
}

public boolean noLessThan(int num1 , int num2){
50  String com1 = String.valueOf(num1) + String.valueOf(num2);
    String com2 = String.valueOf(num2) + String.valueOf(num1);
    if (Long.parseLong(com1) >= Long.parseLong(com2)){
        return true;
    }
55  return false;
}
}
```

Problem 94

224 Basic Calculator

Implement a basic calculator to evaluate a simple expression string.

The expression string may contain open (and closing parentheses), the plus + or minus sign -, non-negative integers and empty spaces .

You may assume that the given expression is always valid.

Some examples:

"1 + 1" = 2

" 2-1 + 2 " = 3

"(1+(4+5+2)-3)+(6+8)" = 23

<https://leetcode.com/problems/basic-calculator/>

Solution

A complicated problem.... String evaluation...

Reference: how to get Reverse Polish Expression :

<http://blog.csdn.net/signsmile/article/details/2877729>

Listing 127: Solution

```
public class Solution {
    Stack<String> operators = new Stack<String>();
    Stack<Integer> operands = new Stack<Integer>();
    public boolean isOperator(String op){
5         return (op.equals("+") || op.equals("-")
            || op.equals("(") || op.equals(")"));
    }
    public int basicCal(String operator){
        int num2 = operands.pop();
10        int num1 = operands.pop();
        if (operator.equals("+")){
            return num1 + num2;
        }
        if (operator.equals("-")){
15            return num1 - num2;
        }
        return -1;
    }
    public int calculate(String s) {
20        int idx = 0 ;
        String entity = "";
        while (idx < s.length()){
            char ch = s.charAt(idx);
            if (ch == ' '){
25                idx ++;
                continue;
            }
            String numStr = "";
```



```
30     while(ch != '+' && ch != '-' && ch != '(' && ch != ')') && ch != ' '
    && idx <= s.length() - 1){
        numStr += ch;
        idx ++;
        if (idx <= s.length() - 1){
            ch = s.charAt(idx);
35     }
    }
    if (numStr.length() > 0){
        operands.push(Integer.parseInt(numStr));
    }
40     if (isOperator(String.valueOf(ch))){
        if (ch == '('){
            operators.push(String.valueOf(ch));
        }else if (ch == ')'){
            while (!operators.isEmpty()
45             && !operators.peek().equals("(")){
                String operator = operators.pop();
                int result = basicCal(operator);
                operands.push(result);
            }
            operators.pop();
50         }else{
            // + or -
            while (!operators.isEmpty()
                && (operators.peek().equals("+")
55             || operators.peek().equals("-"))){
                String operator = operators.pop();
                int result = basicCal(operator);
                operands.push(result);
            }
            operators.push(String.valueOf(ch));
60         }
    }

    idx ++;
65 }
while (!operators.isEmpty()){
    String operator = operators.pop();
    int result = basicCal(operator);
70     operands.push(result);
}
return operands.pop();
}
```

Problem 95

227 Basic Calculator II

Implement a basic calculator to evaluate a simple expression string.

The expression string contains only non-negative integers,
+, -, *, / operators and empty spaces .

The integer division should truncate toward zero.

You may assume that the given expression is always valid.

Some examples:

"3+2*2" = 7

" 3/2 " = 1

" 3+5 / 2 " = 5

Note: Do not use the eval built-in library function.

<https://leetcode.com/problems/climbing-stairs/>

Solution

A more complicated “Basic Calculator”, calculate numbers in stack first whose priority is not larger than current operator to push.(e.g.: when current operator to push is *, calculate until the operator on peek of stack is +/-; if the current operator is +/-, calculate all elements in stack.

Listing 128: Solution

```
public class Solution {
    public int getResult(int num1 , int num2 , char op){
        if (op == '+'){
            return num1 + num2;
        } else if (op == '-'){
            return num1 - num2;
        } else if (op == '*'){
            return num1 * num2;
        } else if (op == '/'){
            return num1 / num2;
        }
        return -1;
    }
    public int calculate(String s) {
        Stack<Integer> operands = new Stack<Integer>();
        Stack<Character> operators = new Stack<Character>();
        int idx = 0;
        while (idx < s.length()){
            char curChar = s.charAt(idx);
            while (curChar == ' ' && ++idx < s.length()){
                curChar = s.charAt(idx);
            }
            String numberStr = "";
            while (idx < s.length() && curChar >= '0' && curChar <= '9'){
                numberStr += curChar;
                idx ++;
            }
        }
    }
}
```

```

    if (idx >= s.length()){
        break;
    }
    curChar = s.charAt(idx);
}
if (numberStr.length() > 0){
    int number = Integer.parseInt(numberStr);
    operands.push(number);
}
if (curChar == ' ')
    continue;
if (curChar == '+' || curChar == '-'
    || curChar == '*' || curChar == '/'){
    if (operators.isEmpty()){
        operators.push(curChar);
    } else {
        // char recentOp = operators.peek();
        if (curChar == '+' || curChar == '-'){
            while(!operators.isEmpty()){
                // calculate numbers in stack first
                int num2 = operands.pop();
                int num1 = operands.pop();
                char op = operators.pop();

                int result = getResult(num1, num2, op);
                operands.push(result);
            }
            operators.push(curChar);
        }
        else if ((curChar == '*' || curChar == '/')){
            while (!operators.isEmpty()
                && operators.peek() != '+'
                && operators.peek() != '-'){
                // calculate numbers in stack first
                int num2 = operands.pop();
                int num1 = operands.pop();
                char op = operators.pop();
                int result = getResult(num1, num2, op);
                operands.push(result);
            }
            operators.push(curChar);
        }
    }
}
idx++;
}

while (!operators.isEmpty()){
    int num2 = operands.pop();
    int num1 = operands.pop();
    char op = operators.pop();
    int result = getResult(num1, num2, op);
    operands.push(result);
}
```

```
80         return operands.pop();  
        }  
    }
```

Problem 96

25 Reverse Nodes in k-Group

Given a linked list, reverse the nodes of a linked list k at a time and return its modified list.

If the number of nodes is not a multiple of k then left-out nodes in the end should remain as it is.

You may not alter the values in the nodes, only nodes itself may be changed.

Only constant memory is allowed.

For example,

Given this linked list: 1->2->3->4->5

For k = 2, you should return: 2->1->4->3->5

For k = 3, you should return: 3->2->1->4->5

<https://leetcode.com/problems/reverse-nodes-in-k-group/>

Solution

A bad news is that I debugged it in my Eclipse IDE.....

Listing 129: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
    public ListNode reverse(ListNode start, ListNode end){
        if (start == null){
            return null;
        }
        // tail = start;
        ListNode cur = start;
        ListNode next = cur.next;
        cur.next = null;
        while (next != null && next != end){
            ListNode backup = next.next;
            next.next = cur;
            cur = next;
            next = backup;
        }
        return cur;
    }
}
```

```
public ListNode reverseKGroup(ListNode head, int k) {  
    ListNode cur = head;  
    ListNode last = cur;  
30    ListNode dummy = new ListNode(-1);  
    dummy.next = head;  
    ListNode tail = dummy;  
    while (cur != null){  
        int count = k;  
35        last = cur;  
        while (count-- > 0){  
            cur = cur.next;  
            if (cur == null){  
                break;  
40            }  
        }  
        if (count > 0 && cur == null){  
            // end up ahead of time  
            tail.next = last;  
45            return dummy.next;  
        }  
        ListNode newPart = reverse(last, cur);  
        tail.next = newPart;  
        tail = last;  
50    }  
    return dummy.next;  
}
```

Problem 97

70 Climbing Stairs

You are climbing a stair case. It takes n steps to reach to the top.

Each time you can either climb 1 or 2 steps.

In how many distinct ways can you climb to the top?

<https://leetcode.com/problems/climbing-stairs/>

Solution

A simple dp problem, so similar to Fibonacci problem.

Listing 130: Solution

```
public class Solution {  
    public int climbStairs(int n) {  
        if (n == 0 || n == 1){  
            return 1;  
        }  
        int cs[] = new int[n + 1];  
        cs[0] = 1;  
        cs[1] = 1;  
        for (int idx = 2; idx <= n ; idx++){  
            cs[idx] = cs[idx - 1] + cs[idx - 2];  
        }  
        return cs[n];  
    }  
}
```

Problem 98

3 Longest Substring Without Repeating Characters

Given a string, find the length of the longest substring without repeating characters.

For example, the longest substring without repeating letters for "abcabcbb" is "abc", which the length is 3.

For "bbbbb" the longest substring is "b", with the length of 1.

<https://leetcode.com/problems/longest-substring-without-repeating-characters/>

Solution

Brute force will have $O(n^2)$ runtime complexity, which will get TLE.

So using a clever way of thought. Maintain 2 pointers, one point to the start of sequence while another point to end of sequence. A array of boolean is used to check if current character has exists in previous sequence. If current character doesn't exists, well good news, we can continue iterating and check if we've got the longest subsequence; if it exists, iterate start pivot to recover the states of those characters, make them as they have not been visited, until meet the character same as current character, rearrange the position of pointer, continue iterating.

Don't forget to check if the sequence at end of whole String is the longest.

Listing 131: Solution

```
public class Solution {
    public int lengthOfLongestSubstring(String s) {
        if (s.length() == 0 || s.length() == 1){
            return s.length();
        }
        int start = 0 , end = 0;
        int maxLength = 0;
        boolean exists[] = new boolean[256];
        Arrays.fill(exists , false);

        while (end < s.length()){
            // iterate to find duplicated element
            maxLength = Math.max(maxLength, end - start);
            if (start != end){
                if (!exists[s.charAt(end)]){
                    exists[s.charAt(end)] = true;
                    end ++;
                }
                else{
                    while (s.charAt(start) != s.charAt(end)){
                        exists[s.charAt(start)] = false;
                        start ++;
                    }
                    start ++;
                    end ++;
                }
            }
            else{
                exists[s.charAt(start)] = true;
            }
        }
    }
}
```



```
        end ++;  
30    }  
    }  
  
    maxLength = Math.max(maxLength, s.length() - start);  
35    return maxLength;  
    }  
}
```

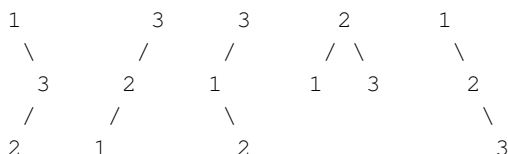
Problem 99

96 Unique Binary Search Trees

Given n , how many structurally unique BST's (binary search trees) that store values $1 \dots n$?

For example,

Given $n = 3$, there are a total of 5 unique BST's.



<https://leetcode.com/problems/unique-binary-search-trees/>

Solution

A picture will be good for our understanding of this problem.

思路

这道题目我想了两天，一直试图找出递推的规律，昨晚才想明白思路有问题，这应该是一道典型的动态规划题目，因为有重叠的子问题，当前决策依赖于子问题的解

我设 $dp[i]$ 表示共有 i 个节点时，能产生的BST树的个数

$n = 0$ 时，空树的个数必然为1，因此 $dp[0] = 1$

$n = 1$ 时，只有1这个根节点，数量也为1，因此 $dp[1] = 1$

$n = 2$ 时，有两种构造方法，如下图所示：

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因此， $dp[2] = dp[0] * dp[1] + dp[1] * dp[0]$

$n = 3$ 时，构造方法如题目给的示例所示， $dp[3] = dp[0] * dp[2] + dp[1] * dp[1] + dp[2] * dp[0]$

同时，当根节点元素为 $1, 2, 3, 4, 5, \dots, i, \dots, n$ 时，基于以下原则的BST树具有唯一性：

以 i 为根节点时，其左子树构成为 $[0, \dots, i-1]$ ，其右子树构成为 $[i+1, \dots, n]$ 构成

因此， $dp[i] = \sum (dp[0 \dots k] * dp[k+1 \dots i]) \quad 0 \leq k < i - 1$

Listing 132: Solution

```

public class Solution {
    public int numTrees(int n) {
        if(n == 0 || n == 1){
            return 1;
        }
        int possibilities[] = new int[n + 1];
        possibilities[0] = 1;
        possibilities[1] = 1;

        for (int rootVal = 2 ; rootVal <= n ; rootVal ++){

```

```

        int sum = 0;
        for (int left = 0 ; left <= rootVal - 1 ; left++){
            int right = rootVal - 1 - left;
            sum += possibilities[left] * possibilities[right];
15    }
        possibilities[rootVal] = sum;
    }

    return possibilities[n];
20 }
}
```

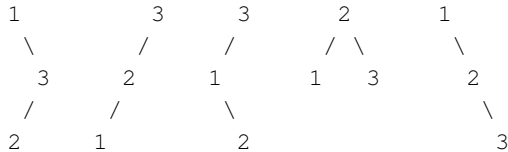
Problem 100

95 Unique Binary Search Trees II

Given n , generate all structurally unique BST's (binary search trees) that store values $1 \dots n$.

For example,

Given $n = 3$, your program should return all 5 unique BST's shown below.



<https://leetcode.com/problems/unique-binary-search-trees-ii/>

Solution

This is an interesting recursion problem since we can return a set of solutions to iteratively add them to left side and right side onto root node.

There are many other problems which set left and right children recursively.

Listing 133: Solution

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public List<TreeNode> generateTree(int start, int end){
        List<TreeNode> uniqueTrees = new ArrayList<TreeNode>();
        if (start > end){
            uniqueTrees.add(null);
15        }
        for (int rootVal = start ; rootVal <= end ; rootVal++){
            List<TreeNode> leftTrees = generateTree(start, rootVal - 1);
            List<TreeNode> rightTrees = generateTree(rootVal + 1, end);
            for (TreeNode leftNode : leftTrees){
20                for (TreeNode rightNode : rightTrees){
                    TreeNode rootNode = new TreeNode(rootVal);
                    rootNode.left = leftNode;
                    rootNode.right = rightNode;
25                uniqueTrees.add(rootNode);
            }
        }
    }
}
  
```

```
        return uniqueTrees;
30    }
    public List<TreeNode> generateTrees(int n) {

        if (n == 0){
            return new ArrayList<TreeNode>();
35        }
        return generateTree(1, n);
    }
}
```

Problem 101

241 Different Ways to Add Parentheses

Given a string of numbers and operators, return all possible results from computing all the different possible ways to group numbers and operators.
The valid operators are +, - and *.

Example 1

Input: "2-1-1".

$((2-1)-1) = 0$
 $(2-(1-1)) = 2$
Output: [0, 2]

Example 2

Input: "2*3-4*5"

$(2*(3-(4*5))) = -34$
 $((2*3)-(4*5)) = -14$
 $((2*(3-4))*5) = -10$
 $(2*((3-4)*5)) = -10$
 $((2*3)-4)*5 = 10$
Output: [-34, -14, -10, -10, 10]

<https://leetcode.com/problems/different-ways-to-add-parentheses/>

Solution

The train of thought is similar to Problem “Unique Binary Search Tree II”, get all possibilities, and combine them all.

Listing 134: Solution

```
public class Solution {
    public List<Integer> diffWaysToCompute(String input) {
        List<Integer> result = new ArrayList<Integer>();
        for (int idx = 0 ; idx < input.length() ; idx++){
            if (input.charAt(idx) == '+' || input.charAt(idx) == '-' ||
                input.charAt(idx) == '*'){
                List<Integer> leftResults =
                    diffWaysToCompute(input.substring(0 , idx));
                List<Integer> rightResults =
                    diffWaysToCompute(input.substring(idx + 1 , input.length()));

                for (int leftResult : leftResults){
                    for (int rightResult : rightResults){
                        switch(input.charAt(idx)){
                            case '+':
                                result.add(leftResult + rightResult);
                                break;
                            case '-':
```

```

        result.add(leftResult - rightResult);
        break;
    default:
        result.add(leftResult * rightResult);
    }
}
}
}
}
}
}
if (result.size() == 0){
    result.add(Integer.parseInt(input));
}
return result;
}
}
```

Problem 102

129 Sum Root to Leaf Numbers

Given a binary tree containing digits from 0-9 only,
each root-to-leaf path could represent a number.

An example is the root-to-leaf path 1->2->3 which represents the number 123.

Find the total sum of all root-to-leaf numbers.

For example,

```
    1
   / \
  2   3
```

The root-to-leaf path 1->2 represents the number 12.

The root-to-leaf path 1->3 represents the number 13.

Return the sum = 12 + 13 = 25.

<https://leetcode.com/problems/sum-root-to-leaf-numbers/>

Solution

A DFS question, mention the edge cases.

Listing 135: Solution

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    int sum = 0;
    public void traverseSum(TreeNode curNode , int curSum){
        if (curNode == null){
            return;
15        }
        curSum = curSum * 10 + curNode.val;
        if (curNode.left == null && curNode.right == null){
            sum += curSum;

20            return;
        }

        traverseSum(curNode.left , curSum);
        traverseSum(curNode.right , curSum);
25    }
}
```



```
30      public int sumNumbers(TreeNode root) {  
          traverseSum(root, 0);  
          return sum;  
      }  
  }
```

Problem 103

173 Binary Search Tree Iterator

Implement an iterator over a binary search tree (BST).
Your iterator will be initialized with the root node of a BST.

Calling `next()` will return the next smallest number in the BST.

Note: `next()` and `hasNext()` should run in average $O(1)$ time and uses $O(h)$ memory, where h is the height of the tree.

<https://leetcode.com/problems/binary-search-tree-iterator/>

Solution

BST is really a tactical data structure. Iterator asks us to return the next smallest element, namely return a ascending list if we call the `next()` function several times. Recall something? The inorder traversal of BST! Then what we only need to do is to repeat what we have done in process of inorder traversal.

Listing 136: Solution

```
/**
 * Definition for binary tree
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class BSTIterator {
    Stack<TreeNode> stack = new Stack<TreeNode>();

    public BSTIterator(TreeNode root) {
        if (root != null){
15         stack.push(root);

        while (root != null && root.left != null){
            stack.push(root.left);
            root = root.left;
20         }
    }

    /** @return whether we have a next smallest number */
    public boolean hasNext() {
25         return (!stack.isEmpty());
    }

    /** @return the next smallest number */
    public int next() {
30         TreeNode nextNode = stack.pop();

        TreeNode rightChild = nextNode.right;
        while (rightChild != null){
            stack.push(rightChild);
        }
    }
}
```

```
        rightChild = rightChild.left;
35    }
    return nextNode.val;
    }
}
/**
40  * Your BSTIterator will be called like this:
    * BSTIterator i = new BSTIterator(root);
    * while (i.hasNext()) v[f()] = i.next();
    */
```

Problem 104

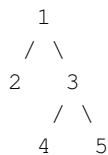
297 Serialize and Deserialize Binary Tree

Serialization is the process of converting a data structure or object into a sequence of bits so that it can be stored in a file or memory buffer, or transmitted across a network connection link to be reconstructed later in the same or another computer environment.

Design an algorithm to serialize and deserialize a binary tree.

There is no restriction on how your serialization/deserialization algorithm should work. You just need to ensure that a binary tree can be serialized to a string and this string can be deserialized to the original tree structure.

For example, you may serialize the following tree



as "[1,2,3,null,null,4,5]", just the same as how LeetCode OJ serializes a binary tree. You do not necessarily need to follow this format, so please be creative and come up with different approaches yourself.

Note: Do not use class member/global/static variables to store states.

Your serialize and deserialize algorithms should be stateless.

<https://leetcode.com/problems/serialize-and-deserialize-binary-tree/>

Solution

Honestly, my solution is not perfect at all(or to say not good at all) since I serialize all elements including null, which is different from LeetCode OJ.

Listing 137: Solution

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Codec {

    // Encodes a tree to a single string.
    public String serialize(TreeNode root) {
        if (root == null){
15         return "[]";
        }
        Queue<TreeNode> queue = new LinkedList<TreeNode>();
        queue.offer(root);
        ArrayList<Integer> valueList = new ArrayList<Integer>();
  
```

```
20     while (!queue.isEmpty()){
        TreeNode curNode = queue.poll();
        if (curNode != null){
            valueList.add(curNode.val);
25     } else{
            valueList.add(null);
            continue;
        }

30         // if (queue.isEmpty()){
        //     // push non-null node
        //     if (curNode.left != null && curNode.right != null){
        //         queue.offer(curNode.left);
        //         queue.offer(curNode.right);
35         //     }else if (curNode.left != null){
        //         queue.offer(curNode.left);
        //     }else if (curNode.right != null){
        //         queue.offer(curNode.right);
        //     }
40         // }else{
        //     // push all children, including null node
        //     queue.offer(curNode.left);
        //     queue.offer(curNode.right);
        // }
45     }

    StringBuilder resultBuilder = new StringBuilder();
    resultBuilder.append("[");
    for (int valueIdx = 0 ; valueIdx < valueList.size() - 1 ; valueIdx ++){
50         if (valueList.get(valueIdx) == null){
            resultBuilder.append("#");
        } else{
            resultBuilder.append(String.valueOf(valueList.get(valueIdx)));
        }
55         resultBuilder.append(",");
    }

    if (valueList.get(valueList.size() - 1) == null){
        resultBuilder.append("#");
60    }else{
        resultBuilder.append(
            String.valueOf(valueList.get(valueList.size() - 1)));
    }
    resultBuilder.append("]");
65

    return resultBuilder.toString();
}

// Decodes your encoded data to tree.
70 public TreeNode deserialize(String data) {
    if (data.equals("[]")){
        return null;
    }
}
```

```
    }

    data = data.substring(1, data.length() - 1);

    Queue<TreeNode> resolvingQueue = new LinkedList<TreeNode>();
    Queue<TreeNode> waitingQueue = new LinkedList<TreeNode>();

    String[] values = data.split(",");
    for (int idxValue = 0 ; idxValue < values.length ; idxValue ++){
        if (values[idxValue].equals("#")){
            waitingQueue.offer(null);
        }else{
            waitingQueue.offer(
                new TreeNode(Integer.parseInt(values[idxValue])));
        }
    }
    TreeNode rootNode = waitingQueue.poll();
    resolvingQueue.offer(rootNode);
    while (!resolvingQueue.isEmpty() && !waitingQueue.isEmpty()){
        TreeNode resolvingNode = resolvingQueue.poll();
        if (resolvingNode != null) {
            TreeNode leftChild = waitingQueue.poll();
            TreeNode rightChild = waitingQueue.poll();

            resolvingNode.left = leftChild;
            resolvingNode.right = rightChild;

            resolvingQueue.offer(leftChild);
            resolvingQueue.offer(rightChild);
        }
    }

    return rootNode;
}

// Your Codec object will be instantiated and called as such:
// Codec codec = new Codec();
// codec.deserialize(codec.serialize(root));
```

Problem 105

284 Peeking Iterator

Given an Iterator class interface with methods: `next()` and `hasNext()`, design and implement a `PeekingIterator` that support the `peek()` operation -- it essentially `peek()` at the element that will be returned by the next call to `next()`.

Here is an example. Assume that the iterator is initialized to the beginning of the list: `[1, 2, 3]`.

Call `next()` gets you 1, the first element in the list.

Now you call `peek()` and it returns 2, the next element. Calling `next()` after that still return 2.

You call `next()` the final time and it returns 3, the last element. Calling `hasNext()` after that should return false.

Hint:

Think of "looking ahead". You want to cache the next element. Is one variable sufficient? Why or why not? Test your design with call order of `peek()` before `next()` vs `next()` before `peek()`. For a clean implementation, check out Google's guava library source code. Follow up: How would you extend your design to be generic and work with all types, not just integer?

<https://leetcode.com/problems/serialize-and-deserialize-binary-tree/>

Solution

My solution is copied from guava source code in some extent, so I will list the standard solution here in Solution 1.2.

Listing 138: Solution 1.1

```
// Java Iterator interface reference:
// https://docs.oracle.com/javase/8/docs/api/java/util/Iterator.html
class PeekingIterator implements Iterator<Integer> {
    Iterator<Integer> iterator;
5    private boolean hasPeeked;
    private Integer peekedElement;

    public PeekingIterator(Iterator<Integer> iterator) {
        // initialize any member here.
10        this.iterator = iterator;
    }

    // Returns the next element in the iteration without advancing the iterator.
    public Integer peek() {
15        if (!hasPeeked){
            if (iterator.hasNext()){
                hasPeeked = true;
            }
        }
    }
}
```

```

        peekedElement = iterator.next();
    }
    20     else{
        // throw new Exception("Nothing to peek");
        peekedElement = null;
    }
}

    25     return peekedElement;
}

// hasNext() and next() should behave the same as in the Iterator interface.
// Override them if needed.
    30     @Override
    public Integer next() {
        if (hasPeeked){
            hasPeeked = false;
            35     return peekedElement;
        }
        else{
            return iterator.next();
        }
    40 }

    @Override
    public boolean hasNext() {
        return hasPeeked || iterator.hasNext();
    45 }
}

```

Listing 139: Solution 1.2

```

/**
 * Implementation of PeekingIterator that avoids peeking unless necessary.
 */
    private static class PeekingImpl<E> implements PeekingIterator<E> {
    5
        private final Iterator<? extends E> iterator;
        private boolean hasPeeked;
        private E peekedElement;

        public PeekingImpl(Iterator<? extends E> iterator) {
            10     this.iterator = checkNotNull(iterator);
        }

        @Override
        15     public boolean hasNext() {
            return hasPeeked || iterator.hasNext();
        }

        @Override
        20     public E next() {
            if (!hasPeeked) {
                return iterator.next();
            }
        }
    }

```



```
    }
    E result = peekedElement;
25    hasPeeked = false;
    peekedElement = null;
    return result;
}

30    @Override
    public void remove() {
        checkState(!hasPeeked, "Can't remove after you've peeked at next");
        iterator.remove();
    }

35    @Override
    public E peek() {
        if (!hasPeeked) {
40            peekedElement = iterator.next();
            hasPeeked = true;
        }
        return peekedElement;
    }
}
```

Problem 106

226 Invert Binary Tree

Invert a binary tree.

```

      4
     / \
    2   7
   / \ / \
  1  3 6  9

```

to

```

      4
     / \
    7   2
   / \ / \
  9  6 3  1

```

Trivia:

This problem was inspired by this original tweet by Max Howell:

Google: 90% of our engineers use the software you wrote (Homebrew),
but you can't invert a binary tree on a whiteboard so fuck off.

<https://leetcode.com/problems/invert-binary-tree/>

Solution 1

A simple problem. Solution 1 is recursive solution.

Listing 140: Solution 1

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public TreeNode invertTree(TreeNode root) {
        if (root == null){
            return null;
        }
15     TreeNode swap = root.left;
        root.left = root.right;
        root.right = swap;

        invertTree(root.left);
        invertTree(root.right);
20     return root;
    }
}

```

Solution 2

Solution 2 is an iterative solution.

Listing 141: Solution 2

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public TreeNode invertTree(TreeNode root) {
        if (root == null){
            return null;
        }
15    Stack<TreeNode> stack = new Stack<TreeNode>();

    stack.push(root);

    while (!stack.isEmpty()){
20        TreeNode curNode = stack.pop();
        TreeNode swap = curNode.left;
        curNode.left = curNode.right;
        curNode.right = swap;

25        if (curNode.left != null){
            stack.push(curNode.left);
        }
        if (curNode.right != null){
30            stack.push(curNode.right);
        }
    }

    return root;
    }
35 }
```

Problem 107

257 Binary Tree Paths

Given a binary tree, return all root-to-leaf paths.

For example, given the following binary tree:

```
    1
   / \
  2   3
   \
    5
```

All root-to-leaf paths are:

```
["1->2->5", "1->3"]
```

<https://leetcode.com/problems/binary-tree-paths/>

Solution

A recursive problem.

Listing 142: Solution

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    List<String> paths = new ArrayList<String>();
    public void traversePath(TreeNode curNode, String prefix){
        if (curNode == null)
            return;
        15 if (curNode.left == null && curNode.right == null)
            paths.add(prefix + curNode.val);
        prefix += curNode.val + "->";
        if (curNode.left != null)
            traversePath(curNode.left , prefix);
        20 if (curNode.right != null)
            traversePath(curNode.right , prefix);
    }
    public List<String> binaryTreePaths(TreeNode root) {
        traversePath(root , "");
        25 return paths;
    }
}
```

Problem 108

144 Binary Tree Preorder Traversal

Given a binary tree, return the preorder traversal of its nodes' values.

For example:

Given binary tree {1,#,2,3},

```
    1
   \
    2
   /
  3
return [1,2,3].
```

Note: Recursive solution is trivial, could you do it iteratively?

<https://leetcode.com/problems/binary-tree-preorder-traversal/>

Solution 1

First come up with the recursive solution.

Listing 143: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public void preorder(TreeNode curNode, List<Integer> result){
        if (curNode != null){
            result.add(curNode.val);
            preorder(curNode.left, result);
            preorder(curNode.right, result);
        }
    }

    public List<Integer> preorderTraversal(TreeNode root) {
        List<Integer> result = new ArrayList<Integer>();
        preorder(root, result);

        return result;
    }
25 }
```

Solution 2

In description, we are asked to finish this task in an iterative way, so we should consider it carefully. Preorder traversal is the easiest traversal in terms of coding. Using a stack to store nodes visited, but remember to push right child first so that in stack they will be popped after left child.

Listing 144: Solution 2

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public List<Integer> preorderTraversal(TreeNode root) {
        if (root == null){
            return new ArrayList<Integer>();
        }
15    List<Integer> result = new ArrayList<Integer>();
    Stack<TreeNode> traversalStack = new Stack<TreeNode>();
    traversalStack.push(root);
    while (!traversalStack.isEmpty()){
        20        TreeNode curNode = traversalStack.pop();

        result.add(curNode.val);

        if (curNode.right != null){
            traversalStack.push(curNode.right);
25        }
        if (curNode.left != null){
            traversalStack.push(curNode.left);
        }
    }
30    return result;
}
}
```

Problem 109

94 Binary Tree Inorder Traversal

Given a binary tree, return the inorder traversal of its nodes' values.

For example:

Given binary tree {1,#,2,3},

```
    1
   \
    2
   /
  3
return [1,3,2].
```

Note: Recursive solution is trivial, could you do it iteratively?

<https://leetcode.com/problems/binary-tree-inorder-traversal/>

Solution 1

First come up with the recursive solution.

Listing 145: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public void inorder(TreeNode curNode, List<Integer> result){
        if(curNode != null){
            inorder(curNode.left, result);
            result.add(curNode.val);
15         inorder(curNode.right, result);
        }
    }

    public List<Integer> inorderTraversal(TreeNode root) {
20         List<Integer> result = new ArrayList<Integer>();
        inorder(root, result);

        return result;
    }
25 }
```

Solution 2

For iterative solution, we use a stack to maintain the order of nodes. This time a bit different- we first push all left nodes of root node iteratively, and when we pop out a `TreeNode`, after adding the value into result list, we should push right child of current node and all its left children into stack.

Listing 146: Solution 2

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public List<Integer> inorderTraversal(TreeNode root) {
        if (root == null){
            return new ArrayList<Integer>();
        }
15    List<Integer> result = new ArrayList<Integer>();
    Stack<TreeNode> traversalStack = new Stack<TreeNode>();

    TreeNode curNode = root;
    while (curNode != null){
20        traversalStack.push(curNode);
        curNode = curNode.left;
    }
    while (!traversalStack.isEmpty()){
        curNode = traversalStack.pop();
25        result.add(curNode.val);

        TreeNode addNode = curNode.right;

        while (addNode != null){
30            traversalStack.push(addNode);
            addNode = addNode.left;
        }
    }

35    return result;
}
```


Problem 110

145 Binary Tree Postorder Traversal

Given a binary tree, return the postorder traversal of its nodes' values.

For example:

Given binary tree {1,#,2,3},

```
    1
   \
    2
   /
  3
return [3,2,1].
```

Note: Recursive solution is trivial, could you do it iteratively?

<https://leetcode.com/problems/binary-tree-postorder-traversal/>

Solution 1

First come up with the recursive solution.

Listing 147: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public void postorderRecursion(TreeNode curNode, List<Integer> result){
        if (curNode != null){
            postorderRecursion(curNode.left, result);
            postorderRecursion(curNode.right, result);
15         result.add(curNode.val);
        }
    }
    public List<Integer> postorderTraversal(TreeNode root) {
        List<Integer> result = new ArrayList<Integer>();
20
        postorderRecursion(root, result);

        return result;
    }
25 }
```

Solution 2

Postorder traversal is more complicated than preorder and inorder. Its complexity lays in the fact that when we first pop the node, we should not print it immediately but should wait for all its children to be

resolved, so we shall set a tag to check if the node is visited, if not, push its children into stack and mark the parent node as visited. So when the next time we meet, we can directly put it into our result list.

Listing 148: Solution 2.1

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 class TagTreeNode{
    TreeNode node;
    boolean visited;
    public TagTreeNode(TreeNode node, boolean visited){
        this.node = node;
        this.visited = visited;
15    }
}

public class Solution {
    public List<Integer> postorderTraversal(TreeNode root) {
20        if (root == null){
            return new ArrayList<Integer>();
        }
        List<Integer> result = new ArrayList<Integer>();
        Stack<TagTreeNode> traversalStack = new Stack<TagTreeNode>();
25        traversalStack.push(new TagTreeNode(root, false));
        while (!traversalStack.isEmpty()){
            TagTreeNode curNode = traversalStack.pop();
            if (curNode.visited){
                result.add(curNode.node.val);
30            }
            else{
                curNode.visited = true;
                traversalStack.push(curNode);
                if (curNode.node.right != null){
35                    traversalStack.push(
                        new TagTreeNode(curNode.node.right, false));
                }
                if (curNode.node.left != null){
40                    traversalStack.push(
                        new TagTreeNode(curNode.node.left, false));
                }
            }
        }
        return result;
45    }
}

```

Another interesting solution is to reverse the result since postorder leads to (left, right, root), when we reverse it, it appears to be (root, right, left), which is so similar to preorder traversal.

Listing 149: Solution 2.2

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public List<Integer> postorderTraversal(TreeNode root) {
        if (root == null){
            return new ArrayList<Integer>();
        }
15    List<Integer> result = new ArrayList<Integer>();
    Stack<TreeNode> traversalStack = new Stack<TreeNode>();
    traversalStack.push(root);
    while (!traversalStack.isEmpty()){
        TreeNode curNode = traversalStack.pop();
20        result.add(curNode.val);
        if (curNode.left != null){
            traversalStack.push(curNode.left);
        }
        if (curNode.right != null){
25            traversalStack.push(curNode.right);
        }
    }
    return reverseList(result);
}
30 private List<Integer> reverseList(List<Integer> list){
    int idxStart = 0 , idxEnd = list.size() - 1;
    while (idxStart < idxEnd){
        Integer swapInt = list.get(idxStart);
        list.set(idxStart, list.get(idxEnd));
35        list.set(idxEnd, swapInt);
        idxStart ++;
        idxEnd --;
    }
    return list;
40 }
}
```

Problem 111

199 Binary Tree Right Side View

Given a binary tree, imagine yourself standing on the right side of it, return the values of the nodes you can see ordered from top to bottom.

For example:

Given the following binary tree,

```

      1             <---
     /  \
    2    3         <---
     \   \
     5    4         <---
  
```

You should return [1, 3, 4].

<https://leetcode.com/problems/binary-tree-postorder-traversal/>

Solution

BFS problem, add value of rightmost node for each level to get the right side view.

Listing 150: Solution

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public List<Integer> rightSideView(TreeNode root) {
        if (root == null){
            return new ArrayList<Integer>();
        }
15    Queue<TreeNode> resolvingNodes = new LinkedList<TreeNode>();
    resolvingNodes.offer(root);
    List<Integer> rightMostNodes = new ArrayList<Integer>();
    while(!resolvingNodes.isEmpty()){
        List<TreeNode> nodeListInCurLevel = new ArrayList<TreeNode>();
20        while(!resolvingNodes.isEmpty()){
            nodeListInCurLevel.add(resolvingNodes.poll());
        }
        rightMostNodes.add(
            nodeListInCurLevel.get(nodeListInCurLevel.size() - 1).val);
25        for (TreeNode curNode : nodeListInCurLevel){
            if (curNode.left != null){
                resolvingNodes.offer(curNode.left);
            }
            if (curNode.right != null){
30                resolvingNodes.offer(curNode.right);
            }
        }
    }
}
  
```

```
35         }  
        }  
        return rightMostNodes;  
    }  
}
```

Problem 112

116 Populating Next Right Pointers in Each Node

Given a binary tree

```
struct TreeLinkNode {
    TreeLinkNode *left;
    TreeLinkNode *right;
    TreeLinkNode *next;
}
```

Populate each next pointer to point to its next right node.

If there is no next right node, the next pointer should be set to NULL.

Initially, all next pointers are set to NULL.

Note:

You may only use constant extra space.

You may assume that it is a perfect binary tree

(ie, all leaves are at the same level, and every parent has two children).

For example,

Given the following perfect binary tree,

```

      1
     / \
    2   3
   / \ / \
  4 5 6 7
```

After calling your function, the tree should look like:

```

      1 -> NULL
     / \
    2 -> 3 -> NULL
   / \ / \
  4->5->6->7 -> NULL
```

<https://leetcode.com/problems/populating-next-right-pointers-in-each-node/>

Solution

Similar to Problem “Binary Tree Right Side View”, use BFS.

Listing 151: Solution

```

/**
 * Definition for binary tree with next pointer.
 * public class TreeLinkNode {
 *     int val;
 *     TreeLinkNode left, right, next;
 *     TreeLinkNode(int x) { val = x; }
 * }
 */
public class Solution {
    public void connect(TreeLinkNode root) {
        if (root == null){
            return;
        }
    }
}
```

```
    }
    Queue<TreeLinkNode> resolvingNodes = new LinkedList<TreeLinkNode>();
15    resolvingNodes.offer(root);
    while (!resolvingNodes.isEmpty()){
        List<TreeLinkNode> nodesCurLevel = new ArrayList<TreeLinkNode>();
        while (!resolvingNodes.isEmpty()){
            nodesCurLevel.add(resolvingNodes.poll());
20        }
        for (int nodeId = 0 ; nodeId < nodesCurLevel.size() ; nodeId++){
            if (nodeId != nodesCurLevel.size() - 1){
                nodesCurLevel.get(nodeId).next =
                    nodesCurLevel.get(nodeId + 1);
25            }
            else{
                nodesCurLevel.get(nodeId).next = null;
            }
            if (nodesCurLevel.get(nodeId).left != null){
30                resolvingNodes.offer(nodesCurLevel.get(nodeId).left);
            }
            if (nodesCurLevel.get(nodeId).right != null){
                resolvingNodes.offer(nodesCurLevel.get(nodeId).right);
            }
35        }
    }
}
```

Problem 113

117 Populating Next Right Pointers in Each Node II

Follow up for problem "Populating Next Right Pointers in Each Node".

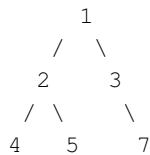
What if the given tree could be any binary tree?
Would your previous solution still work?

Note:

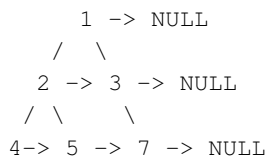
You may only use constant extra space.

For example,

Given the following binary tree,



After calling your function, the tree should look like:



<https://leetcode.com/problems/populating-next-right-pointers-in-each-node-ii/>

Solution

$O(1)$ space is a big challenge. If the tree is a perfect binary tree, an interesting way is to connect left child and right child of each node, and when current node has a next object not pointing to null, connect right child of current node to left child of next node of current node.

However, it is said that the tree is just a common binary tree, things seems a bit more complicated, we shall consider all kinds of circumstances and edge cases, following former train of thought.

Listing 152: Solution

```

/**
 * Definition for binary tree with next pointer.
 * public class TreeLinkNode {
 *     int val;
 *     TreeLinkNode left, right, next;
 *     TreeLinkNode(int x) { val = x; }
 * }
 */
public class Solution {
    public void connect(TreeLinkNode root) {
        TreeLinkNode currentLevelHead = root, nextLevelHead = null;

        while (currentLevelHead != null){
            TreeLinkNode currentNode = currentLevelHead;
            TreeLinkNode prevTail = null;

```



```
20     while (currentNode != null){
        if (currentNode.left != null && currentNode.right != null){
            currentNode.left.next = currentNode.right;
        }
        if (currentNode.left == null && currentNode.right == null){
            currentNode = currentNode.next;
            continue;
        }
25     else{
        TreeLinkNode firstNodeInRow = (currentNode.left == null)
            ? currentNode.right : currentNode.left;
        if (nextLevelHead == null){
            nextLevelHead = firstNodeInRow;
            prevTail = firstNodeInRow;
30        }
        else{
            prevTail.next = firstNodeInRow;
        }
35        while (prevTail.next != null){
            prevTail = prevTail.next;
        }
40        currentNode = currentNode.next;
    }
    currentLevelHead = nextLevelHead;
    nextLevelHead = null;
45 }
}
```

Problem 114

105 Construct Binary Tree from Preorder and Inorder Traversal

Given preorder and inorder traversal of a tree, construct the binary tree.

<https://leetcode.com/problems/-construct-binary-tree-from-preorder-and-inorder-traversal/>

Solution

How to split the two arrays is a important stuff to consider, recursion is a natural way to achieve this.

Listing 153: Solution

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public TreeNode buildTree(int[] preorder, int startPre, int endPre,
        int[] inorder, int startIn, int endIn){
        if (startPre > endPre || startIn > endIn){
            return null;
        }
        15 int split = preorder[startPre];

        TreeNode rootNode = new TreeNode(split);
        if (startPre == endPre){
            20 return rootNode;
        }
        int idx = startIn;
        for (idx = startIn ; idx <= endIn ; idx++){
            if (inorder[idx] == split){
                25 break;
            }
        }
        rootNode.left = buildTree(preorder, startPre + 1,
            startPre + idx - startIn, inorder, startIn, idx - 1);
        30 rootNode.right = buildTree(preorder,
            startPre + idx - startIn + 1, endPre, inorder, idx + 1, endIn);

        return rootNode;
    }
    35 public TreeNode buildTree(int[] preorder, int[] inorder) {
        return buildTree(preorder, 0, preorder.length - 1,
            inorder, 0, inorder.length - 1);
    }
    40 }
```

Problem 115

106 Construct Binary Tree from Inorder and Postorder Traversal

Given inorder and postorder traversal of a tree, construct the binary tree.

Note:

You may assume that duplicates do not exist in the tree.

<https://leetcode.com/problems/-construct-binary-tree-from-inorder-and-postorder-traversal/>

Solution

Similar to Problem 105 “Construct Binary Tree from Preorder and Inorder Traversal”.

Listing 154: Solution

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public TreeNode buildTree(int[] inorder, int startInorder, int endInorder,
        int[] postorder, int startPost, int endPost){
        if (startInorder > endInorder){
            return null;
15        }

        int split = postorder[endPost];

        TreeNode rootNode = new TreeNode(split);
20        if (startInorder == endInorder){
            return rootNode;
        }

        int splitIdx = -1;
25        for (int idx = startInorder; idx <= endInorder; idx++){
            if (inorder[idx] == split){
                splitIdx = idx;
                break;
            }
30        }

        rootNode.left = buildTree(inorder, startInorder, splitIdx - 1,
            postorder, startPost, startPost + splitIdx - 1 - startInorder);
        rootNode.right = buildTree(inorder, splitIdx + 1, endInorder,
35        postorder, startPost + splitIdx - startInorder, endPost - 1);

        return rootNode;
    }
}
```

```
    }  
40    public TreeNode buildTree(int[] inorder, int[] postorder) {  
        return buildTree(inorder, 0, inorder.length - 1,  
            postorder, 0, postorder.length - 1);  
    }  
}
```

Problem 116

242 Valid Anagram

Given two strings *s* and *t*, write a function to determine if *t* is an anagram of *s*.

For example,

s = "anagram", *t* = "nagaram", return true.

s = "rat", *t* = "car", return false.

Note:

You may assume the string contains only lowercase alphabets.

Follow up:

What if the inputs contain unicode characters?

How would you adapt your solution to such case?

<https://leetcode.com/problems/valid-anagram/>

Solution

For ascii characters, maintain 2 arrays containing appearing count for each letter. As for the follow up, can extending the range of array to 65535 work?

Well, there is another solution which uses sorting method. *s.toCharArray()* then *Arrays.sort* to see if sorted *s* and sorted *t* are equal.....

Listing 155: Solution

```
public class Solution {
    public boolean isAnagram(String s, String t) {
        int[] countS = new int[256];
        int[] countT = new int[256];
5       Arrays.fill(countS, 0);
        Arrays.fill(countT, 0);
        if (s.length() != t.length()){
            return false;
        }
10      for (int idxS = 0 ; idxS < s.length() ; idxS++){
            char chS = s.charAt(idxS);
            countS[chS]++;
        }
        for (int idxT = 0 ; idxT < t.length() ; idxT++){
15      char chT = t.charAt(idxT);
            countT[chT]++;
        }
        for (int idx = 0 ; idx < 256 ; idx++){
            if (countS[idx] != countT[idx]){
20      return false;
            }
        }
        return true;
    }
25 }
```

Problem 117

49 Group Anagrams

Given an array of strings, group anagrams together.

For example, given: ["eat", "tea", "tan", "ate", "nat", "bat"],
Return:

```
[
  ["ate", "eat","tea"],
  ["nat","tan"],
  ["bat"]
]
```

Note:

For the return value, each inner list's elements
must follow the lexicographic order.

All inputs will be in lower-case.

<https://leetcode.com/problems/anagrams/>

Solution

I think the solution is somewhat opportunistic... Use Arrays.sort to serve as the fundamental part of the whole program...

However, when I was looking for a more proper solution in the Discuss section, I found that most solutions are of the same trains of thought. In fact, the key insight of solution is not the sorting of charArray but finding a proper hashCode for the Hashmap which stores each group. We may use expressions like "a2b3" to represent each word(don't forget to arrange the order of each letter).

There is another optimal point to check length of each word, words with same length will never be anagrams.

Listing 156: Solution

```
public class Solution {
    public List<List<String>> groupAnagrams(String[] strs) {
        List<List<String>> result = new ArrayList<List<String>>();
        HashMap<String, List<String>> anagramGroups
5         = new HashMap<String, List<String>>();
        Arrays.sort(strs);
        for (String str : strs){
            char[] charArr = str.toCharArray();
            Arrays.sort(charArr);
10         String sortedString = new String(charArr);
            List<String> anaList = new ArrayList<String>();
            if (anagramGroups.containsKey(sortedString)){
                anaList = anagramGroups.get(sortedString);
            }
15         anaList.add(str);
            anagramGroups.put(sortedString, anaList);
        }
        Set<String> anagramKeySet = anagramGroups.keySet();
        for (String anagramKey : anagramKeySet){
20         result.add(anagramGroups.get(anagramKey));
        }
```

```
        }  
        return result;  
    }  
}
```

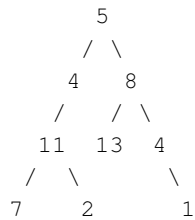
Problem 118

112 Path Sum

Given a binary tree and a sum, determine if the tree has a root-to-leaf path such that adding up all the values along the path equals the given sum.

For example:

Given the below binary tree and sum = 22,



return true, as there exist a root-to-leaf path 5->4->11->2 which sum is 22.

<https://leetcode.com/problems/path-sum/>

Solution

DFS, bring remaining sum into parameter.

Listing 157: Solution

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public boolean findPath(TreeNode curNode, int remainingSum){
        if (curNode == null)
            return false;
        remainingSum -= curNode.val;
15        if (curNode.left == null && curNode.right == null){
            if (remainingSum == 0)
                return true;
            else
                return false;
20        }
        return findPath(curNode.left , remainingSum)
            || findPath(curNode.right , remainingSum);
    }
    public boolean hasPathSum(TreeNode root , int sum) {
25        return findPath(root , sum);
    }
}

```

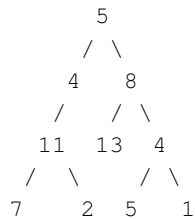

Problem 119

113 Path Sum II

Given a binary tree and a sum, find all root-to-leaf paths where each path's sum equals the given sum.

For example:

Given the below binary tree and sum = 22,



```

return
[
  [5,4,11,2],
  [5,8,4,5]
]

```

<https://leetcode.com/problems/path-sum-ii/>

Solution

I'm using a new ArrayList to do the dfs since List may be changed during recursion. I don't think this is necessary.

Listing 158: Solution

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    List<List<Integer>> result = new ArrayList<List<Integer>>();

    public void findPath(TreeNode curNode, int remainingSum,
15         List<Integer> curPath){
        if (curNode == null){
            return;
        }
        remainingSum -= curNode.val;
        curPath.add(curNode.val);

20         if (curNode.left == null && curNode.right == null){
            if (remainingSum == 0){
                result.add(new ArrayList<Integer>(curPath));
            }
            return;
        }
    }
}

```

```
25         }  
        else{  
            return;  
        }  
    }  
  
30    findPath(curNode.left , remainingSum , new ArrayList(curPath));  
    findPath(curNode.right , remainingSum , new ArrayList(curPath));  
}  
  
35    public List<List<Integer>> pathSum(TreeNode root , int sum) {  
        List<Integer> curPath = new ArrayList<Integer>();  
        findPath(root , sum, curPath);  
  
        return result;  
40    }  
}
```

Problem 120

110 Balanced Binary Tree

Given a binary tree, determine if it is height-balanced.

For this problem, a height-balanced binary tree is defined as a binary tree in which the depth of the two subtrees of every node never differ by more than 1.

<https://leetcode.com/problems/balanced-binary-tree/>

Solution 1

Use maxDepth to see, use a top down recursion, though with $O(n^2)$ runtime complexity.

Be aware of the importance of short-circuit logic

Listing 159: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public int maxDepth(TreeNode curNode){
        if (curNode == null){
            return 0;
        }
15         return Math.max(maxDepth(curNode.left), maxDepth(curNode.right)) + 1;
    }

    public boolean isBalanced(TreeNode root) {
20         if (root == null){
            return true;
        }

        return (Math.abs(maxDepth(root.left) - maxDepth(root.right)) <= 1)
25         && isBalanced(root.left) && isBalanced(root.right);
    }
}
```

Solution 2

Another solution is from “Clean Code Handbook”. A bottom-up recursion, use -1 to indicate that the subtree is already proved non-balanced.

Listing 160: Solution 2

```
/**
```

```

5  * Definition for a binary tree node.
   * public class TreeNode {
   *     int val;
   *     TreeNode left;
   *     TreeNode right;
   *     TreeNode(int x) { val = x; }
   * }
   */
10 public class Solution {
   public int maxDepth(TreeNode curNode){
       if (curNode == null){
           return 0;
       }
15     int lMax = maxDepth(curNode.left);
       if (lMax == -1){
           return -1;
       }
20     int rMax = maxDepth(curNode.right);
       if (rMax == -1){
           return -1;
       }
25     return (Math.abs(lMax - rMax) > 1) ? -1 : Math.max(lMax, rMax) + 1;
   }

   public boolean isBalanced(TreeNode root) {
       if (root == null){
30         return true;
       }

       return maxDepth(root) != -1;
   }
35 }
```

Problem 121

230 Kth Smallest Element in a BST

Given a binary search tree, write a function `kthSmallest` to find the `k`th smallest element in it.

Note:

You may assume `k` is always valid, $1 \leq k \leq$ BST's total elements.

Follow up:

What if the BST is modified (insert/delete operations) often and you need to find the `k`th smallest frequently? How would you optimize the `kthSmallest` routine?

Hint:

Try to utilize the property of a BST.

What if you could modify the BST node's structure?

The optimal runtime complexity is $O(\text{height of BST})$.

<https://leetcode.com/problems/kth-smallest-element-in-a-bst/>

Solution 1

My first solution is to use inorder traversal since the result of such traversal is a list in order, I can easily pick out the `k`th one, which is what I want. However, it cannot solve the follow up problem (or adding values in the sorted list), neither can it be called a 'clean' code and has the optimized complexity of $O(\text{height of BST})$ time.

Listing 161: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    List<Integer> result = new ArrayList<Integer>();
    public void findKth(TreeNode curNode){
        if (curNode == null){
            return;
15         }
        else{
            findKth(curNode.left);
            result.add(curNode.val);
            findKth(curNode.right);
20         }
    }

    public int kthSmallest(TreeNode root, int k) {
        Queue<TreeNode> queue = new LinkedList<TreeNode>();
```

```

25         findKth(root);

        return result.get(k - 1);
    }
30 }

```

Solution 2

However, the problem is definitely a standard binary-search like problem. By judging node number of left subtree and right subtree, we can determine which subtree can be used to continue searching.

More information about the train of thought, see

Add property `size` to BST node. Looks like:

```

class TreeNode(object):
    def __init__(self, x):
        self.val = x
        self.left = None
        self.right = None
        self.size = 1 # the size of subtree whose root is current node

```

When `insert/delete`, update `size` too. Then we can get kth smallest element according to `size`.

Visit from `root` of the `BST`:

- if `root.size == k`, return `root.val`.
- if `root.size > k`, search the `k` smallest element of `root.left`
- if `root.size < k`, search the `k - root.size` smallest element of `root.right`

At most we search `height of BST` times, so it's `O(height of BST)`.

Listing 162: Solution 2

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 class TagTreeNode{
    TreeNode node;
    int size;
    TagTreeNode left;
    TagTreeNode right;
15    public TagTreeNode(TreeNode node, int size){
        this.node = node;
        this.size = size;
    }
}
20 public class Solution {
    public TagTreeNode buildNode(TreeNode root){
        if (root == null){
            return null;

```

```
    }
    TagTreeNode leftTagNode = buildNode(root.left);
    TagTreeNode rightTagNode = buildNode(root.right);
    int lSize = getSize(leftTagNode);
    int rSize = getSize(rightTagNode);
    int newSize = 1 + lSize + rSize;

    TagTreeNode tagNode = new TagTreeNode(root, newSize);
    tagNode.left = leftTagNode;
    tagNode.right = rightTagNode;
    return tagNode;
}

public int getSize(TagTreeNode curNode){
    if (curNode == null){
        return 0;
    }
    return curNode.size;
}

public int kthSmallest(TreeNode root, int k) {
    TagTreeNode tagRoot = buildNode(root);
    int result = binarySearch(tagRoot, k);
    return result;
}

public int binarySearch(TagTreeNode tagNode, int k){
    int size = getSize(tagNode.left);
    if (size + 1 == k){
        return tagNode.node.val;
    }
    else if (size + 1 > k){
        return binarySearch(tagNode.left, k);
    }else{
        return binarySearch(tagNode.right, k - size - 1);
    }
}
}
```

Problem 122

278 First Bad Version

You are a product manager and currently leading a team to develop a new product. Unfortunately, the latest version of your product fails the quality check. Since each version is developed based on the previous version, all the versions after a bad version are also bad.

Suppose you have n versions $[1, 2, \dots, n]$ and you want to find out the first bad one, which causes all the following ones to be bad.

You are given an API `bool isBadVersion(version)` which will return whether version is bad. Implement a function to find the first bad version. You should minimize the number of calls to the API.

<https://leetcode.com/problems/first-bad-version/>

Solution

A very simple binary search problem, BUT! use “ $\text{mid} = \text{low} + (\text{high} - \text{low}) / 2;$ ” instead of “ $\text{mid} = (\text{low} + \text{high}) / 2;$ ” since it may overflow to be negative so that the loop will never end.

Listing 163: Solution

```
/* The isBadVersion API is defined in the parent class VersionControl.
   boolean isBadVersion(int version); */

5 public class Solution extends VersionControl {
    public int firstBadVersion(int n) {
        if (n <= 1){
            return n;
        }
        int low = 1, high = n;
10 while (low < high){
            int mid = low + (high - low) / 2;
            if (isBadVersion(mid)){
                high = mid;
            }
            else{
15 low = mid + 1;
            }
        }
        return low;
20 }
}
```


Problem 123

74 Search a 2D Matrix

Write an efficient algorithm that searches for a value in an $m \times n$ matrix.

This matrix has the following properties:

Integers in each row are sorted from left to right.

The first integer of each row is greater than the last integer of the previous row.

For example,

Consider the following matrix:

```
[
  [1,   3,  5,  7],
  [10, 11, 16, 20],
  [23, 30, 34, 50]
]
```

Given target = 3, return true.

<https://leetcode.com/problems/search-a-2d-matrix/>

Solution

Starting from [0,0] is very difficult to solve the problem since when we've got a larger element, we can go both downward and to the right. Starting from the top right corner. If current element is less than target, go to the left side; if larger, go downward.

I've missed the condition that 'The first integer of each row is greater than the last integer of the previous row.' So I'm using the solution which also goes for the next Problem "Search a 2D Matrix II".

Listing 164: Solution

```
public class Solution {
    public boolean searchMatrix(int [][] matrix, int target) {
        int startRow = 0;
        if (matrix.length == 0)
            return false;
        int startCol = matrix[0].length - 1;
        while (startCol >= 0 && startRow < matrix.length){
            int compare = matrix[startRow][startCol];
            if (compare == target)
                return true;
            else{
                if (target < compare){
                    startCol --;
                }
                else{
                    startRow ++;
                }
            }
        }
        return false;
    }
}
```

Problem 124

240 Search a 2D Matrix II

Write an efficient algorithm that searches for a value in an $m \times n$ matrix.

This matrix has the following properties:

Integers in each row are sorted in ascending from left to right.

Integers in each column are sorted in ascending from top to bottom.

For example,

Consider the following matrix:

```
[
  [1,   4,   7, 11, 15],
  [2,   5,   8, 12, 19],
  [3,   6,   9, 16, 22],
  [10, 13, 14, 17, 24],
  [18, 21, 23, 26, 30]
]
```

Given target = 5, return true.

Given target = 20, return false.

<https://leetcode.com/problems/search-a-2d-matrix-ii/>

Solution

Since I've used the method to solve the previous problem, I won't spend too many words here explaining it.

Listing 165: Solution

```
public class Solution {
    public boolean searchMatrix(int [][] matrix, int target) {
        int curRow = 0;
        if (matrix.length == 0){
5           return false;
        }
        int curCol = matrix[0].length - 1;
        while (curRow < matrix.length && curCol >= 0){
            if (target == matrix[curRow][curCol]){
10              return true;
            }
            else if (target < matrix[curRow][curCol]){
                curCol --;
            }
            else{
15              curRow ++;
            }
        }
        return false;
20    }
}
```

Problem 125

99 Recover Binary Search Tree

Two elements of a binary search tree (BST) are swapped by mistake.

Recover the tree without changing its structure.

Note:

A solution using $O(n)$ space is pretty straight forward.

Could you devise a constant space solution?

<https://leetcode.com/problems/recover-binary-search-tree/>

Solution 1

Regardless of the problem requirements of $O(1)$ space, the problem can be solved using inorder traversal since for BST, the inorder traversal list should be ascending, so if there is one or two mistaken parts, we should swap the values of nodes.

Listing 166: Solution 1

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    List<TreeNode> inorderList = new ArrayList<TreeNode>();

    public void inorderTraverse(TreeNode curNode){
        if (curNode == null){
15             return;
        }

        inorderTraverse(curNode.left);
        inorderList.add(curNode);
20         inorderTraverse(curNode.right);
    }

    public void recoverTree(TreeNode root) {
        inorderTraverse(root);

25         boolean firstToFind = true;
        TreeNode mistakenNode1 = null, mistakenNode2 = null;

        for (int nodeId = 0 ; nodeId < inorderList.size() - 1 ; nodeId++){
30             if (inorderList.get(nodeId).val >= inorderList.get(nodeId + 1).val){
                if (firstToFind){
                    mistakenNode1 = inorderList.get(nodeId);
                    mistakenNode2 = inorderList.get(nodeId + 1);
                }
            }
        }
        // Swap the values of the two nodes
        int temp = mistakenNode1.val;
        mistakenNode1.val = mistakenNode2.val;
        mistakenNode2.val = temp;
    }
}
```

```

        firstToFind = false;
35         }else{
            mistakenNode2 = inorderList.get(nodeIdx + 1);
        }
    }
}

40 // swap
int swapVal = mistakenNode1.val;
mistakenNode1.val = mistakenNode2.val;
mistakenNode2.val = swapVal;
45 }
}

```

Solution 2

Since the requirement is $O(1)$ space, we need to think something strange to us – the Morris traversal. See figures:

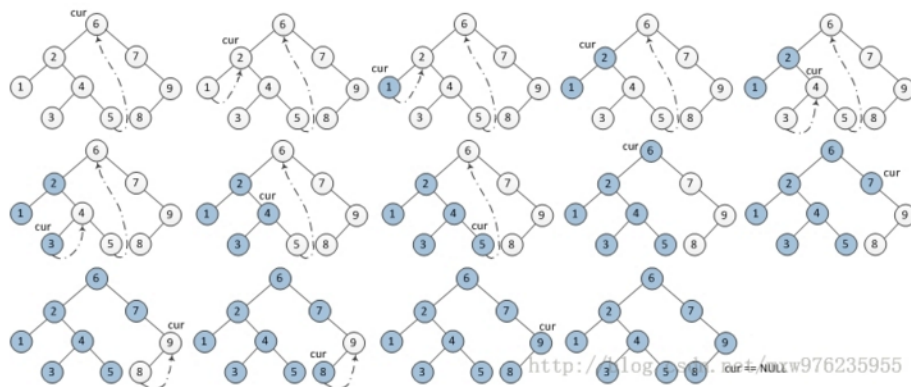
一、中序遍历

步骤:

1. 如果当前节点的左孩子为空，则输出当前节点并将其右孩子作为当前节点。
2. 如果当前节点的左孩子不为空，在当前节点的左子树中找到当前节点在中序遍历下的前驱节点。
 - a) 如果前驱节点的右孩子为空，将它的右孩子设置为当前节点。当前节点更新为当前节点的左孩子。
 - b) 如果前驱节点的右孩子为当前节点，将它的右孩子重新设为空（恢复树的形状）。输出当前节点。当前节点更新为当前节点的右孩子。
3. 重复以上1、2直到当前节点为空。

图示:

下图为每一步迭代的结果（从左至右，从上到下），cur代表当前节点，深色节点表示该节点已输出



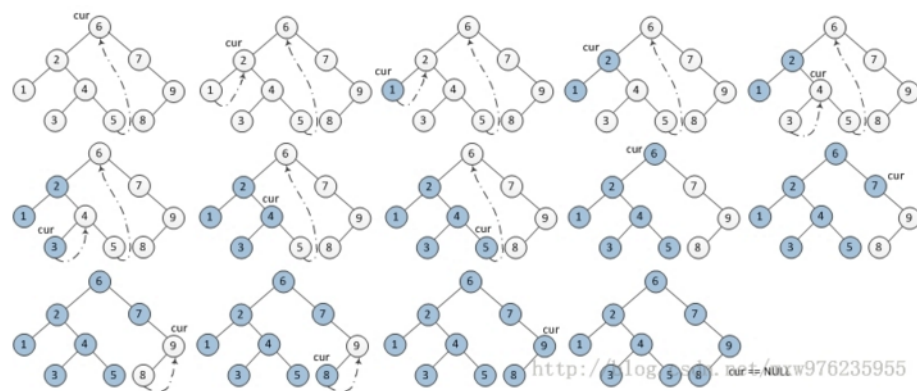
二、前序遍历

前序遍历与中序遍历相似，代码上只有一行不同，不同就在于输出的顺序。

步骤：

1. 如果当前节点的左孩子为空，则输出当前节点并将其右孩子作为当前节点。
2. 如果当前节点的左孩子不为空，在当前节点的左子树中找到当前节点在中序遍历下的前驱节点。
 - a) 如果前驱节点的右孩子为空，将它的右孩子设置为当前节点。输出当前节点（**在这里输出，这是与中序遍历唯一点不同**）。当前节点更新为当前节点的左孩子。
 - b) 如果前驱节点的右孩子为当前节点，将它的右孩子重新设为空。当前节点更新为当前节点的右孩子。
3. 重复以上1、2直到当前节点为空。

图示：



三、后序遍历

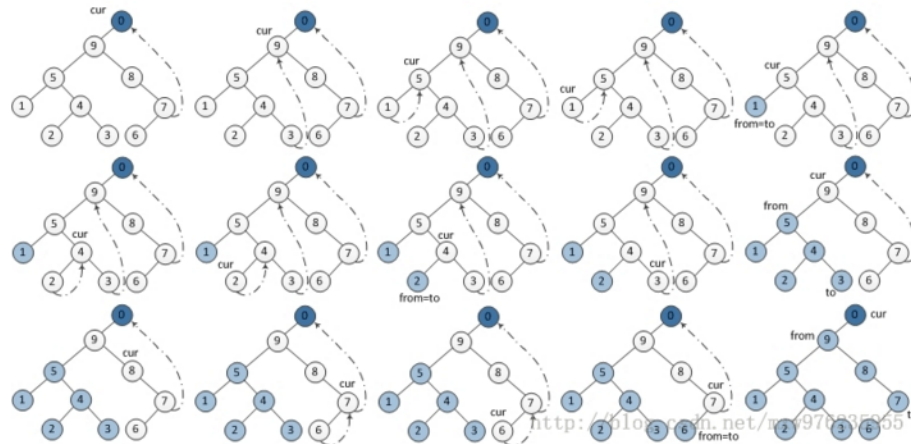
后续遍历稍显复杂，需要建立一个临时节点dump，令其左孩子是root。并且还需要一个子过程，就是倒序输出某两个节点之间路径上的各个节点。

步骤：

当前节点设置为临时节点dump。

1. 如果当前节点的左孩子为空，则将其右孩子作为当前节点。
2. 如果当前节点的左孩子不为空，在当前节点的左子树中找到当前节点在中序遍历下的前驱节点。
 - a) 如果前驱节点的右孩子为空，将它的右孩子设置为当前节点。当前节点更新为当前节点的左孩子。
 - b) 如果前驱节点的右孩子为当前节点，将它的右孩子重新设为空。倒序输出从当前节点的左孩子到该前驱节点这条路径上的所有节点。当前节点更新为当前节点的右孩子。
3. 重复以上1、2直到当前节点为空。

图示：



So simulating the process is the key to achieve this. In the algorithm, when we want to 'print', it means it's the time for us to refresh the value of previous node and make the check if the array is still sorted.

Listing 167: Solution 2

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
public class Solution {
    TreeNode mistakenNode1 = null, mistakenNode2 = null;
    TreeNode prev = null;
    boolean firstFound = true;
    public void recoverTree(TreeNode root) {
        TreeNode curNode = root;
        while (curNode != null){
            if (curNode.left == null){
                // print
            }
        }
    }
}

```

```
20         if (prev != null){
            if (curNode.val < prev.val){
                if (firstFound){
                    mistakenNode1 = prev;
                    mistakenNode2 = curNode;
                    firstFound = false;
25                 }
                else{
                    mistakenNode2 = curNode;
                }
            }
        }
        prev = curNode;
        curNode = curNode.right;
    }
    else{
35        //curNode.left != null, so find rightmost(predecessor)
        // node in left subtree
        TreeNode movingNode = curNode.left;
        while (movingNode.right != null && movingNode.right != curNode){
            movingNode = movingNode.right;
40        }
        if (movingNode.right == null){
            movingNode.right = curNode;
            curNode = curNode.left;
            // This is important, we should
            // not add this statement since this is not the right time
            // To be more specific, now curNode = curNode.left,
            // so the prev should points to curNode.left.predecessor,
            // which is exactly what prev used to be
            // prev = movingNode;
45        }
        else{
            movingNode.right = null;
            prev = movingNode;
            if (prev != null){
55                if (curNode.val < prev.val){
                    if (firstFound){
                        mistakenNode1 = prev;
                        mistakenNode2 = curNode;
                        firstFound = false;
60                    }
                    else{
                        mistakenNode2 = curNode;
                    }
                }
            }
65        }
        prev = curNode;
        curNode = curNode.right;
    }
}
70 }
```

```
75      // swap
      int swapVal = mistakenNode1.val;
      mistakenNode1.val = mistakenNode2.val;
      mistakenNode2.val = swapVal;
    }
  }
```

Solution 3

And a recursion solution ... I can hardly understand it.

Listing 168: Solution 3

```
/**
 * Definition for binary tree
 * struct TreeNode {
 *     int val;
 *     TreeNode *left;
 *     TreeNode *right;
 *     TreeNode(int x) : val(x), left(NULL), right(NULL) {}
 * };
 */
10 class Solution {
public:
    TreeNode *p,*q;
    TreeNode *prev;
    15 void recoverTree(TreeNode *root)
    {
        p=q=prev=NULL;
        inorder(root);
        swap(p->val,q->val);
    }
    20 void inorder(TreeNode *root)
    {
        if(root->left) inorder(root->left);
        if(prev!=NULL&&(prev->val>root->val))
        {
            25 if(p==NULL)p=prev;
            q=root;
        }
        prev=root;
        if(root->right) inorder(root->right);
    30 }
};
```

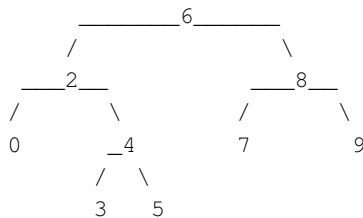

Problem 126

235 Lowest Common Ancestor of a Binary Search Tree

Given a binary search tree (BST), find the lowest common ancestor (LCA) of two given nodes in the BST.

According to the definition of LCA on Wikipedia:

The lowest common ancestor is defined between two nodes v and w as the lowest node in T that has both v and w as descendants (where we allow a node to be a descendant of itself).



For example, the lowest common ancestor (LCA) of nodes 2 and 8 is 6. Another example is LCA of nodes 2 and 4 is 2, since a node can be a descendant of itself according to the LCA definition.

<https://leetcode.com/problems/lowest-common-ancestor-of-a-binary-search-tree/>

Solution 1

LCA is a typical kind of problem. If possible, review Tarjan & find-union.

First solution makes use of the characteristics of binary SEARCH tree, when finding the lowest ancestor, the value of the ancestor node v_{an} should be $v_{min} \leq v_{an} \leq v_{max}$.

Listing 169: Solution 1

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public TreeNode lowestCommonAncestor(TreeNode root, TreeNode p, TreeNode q){
        TreeNode curNode = root;
        int minVal = Math.min(p.val, q.val);
        int maxVal = Math.max(p.val, q.val);
15
        while (!(curNode.val >= minVal && curNode.val <= maxVal)){
            if (curNode.val > maxVal){
                curNode = curNode.left;
            }
20            if (curNode.val < minVal){

```

```

        curNode = curNode.right;
    }
}
25     return curNode;
}
}

```

Solution 2

Solution use the path to solve the problem. Still, we are using the feature of BST. In the next problem we will show a more general algorithm which can be applied to Binary tree.

Listing 170: Solution 2

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public List<TreeNode> getPath(TreeNode startNode, TreeNode target){
        TreeNode curNode = startNode;
        List<TreeNode> result = new ArrayList<TreeNode>();

15        while (curNode != target){
            result.add(curNode);
            if (curNode.val > target.val){
                curNode = curNode.left;
            }else if (curNode.val < target.val){
20                curNode = curNode.right;
            }else{
                break;
            }
        }
25        result.add(target);

        return result;
    }

30    public TreeNode lowestCommonAncestor(TreeNode root, TreeNode p, TreeNode q) {
        List<TreeNode> pathP = new ArrayList<TreeNode>();
        List<TreeNode> pathQ = new ArrayList<TreeNode>();

35        pathP = getPath(root, p);
        pathQ = getPath(root, q);

        for (int idx = 0 ; idx < Math.min(pathP.size(), pathQ.size()) ; idx ++){
            if (pathP.get(idx) != pathQ.get(idx)){

```

```
40         return pathP.get(idx - 1);  
        }  
    }  
    return (pathP.size() > pathQ.size()) ? pathQ.get(pathQ.size() - 1)  
        : pathP.get(pathP.size() - 1) ;  
45 }  
}
```

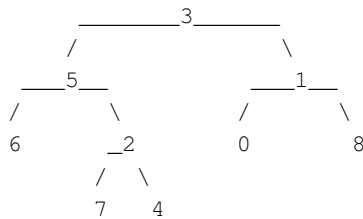
Problem 127

236 Lowest Common Ancestor of a Binary Tree

Given a binary tree, find the lowest common ancestor (LCA) of two given nodes in the tree.

According to the definition of LCA on Wikipedia:

The lowest common ancestor is defined between two nodes v and w as the lowest node in T that has both v and w as descendants (where we allow a node to be a descendant of itself).



For example, the lowest common ancestor (LCA) of nodes 5 and 1 is 3.

Another example is LCA of nodes 5 and 4 is 5, since a node can be a descendant of itself according to the LCA definition.

<https://leetcode.com/problems/lowest-common-ancestor-of-a-binary-tree/>

Solution 1

Since this is just a binary tree, we can only use the characteristic that the two target nodes are respectively at the left and right subtree of their LCA. It's a bottom-up solution.

Listing 171: Solution 1

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public TreeNode findChild(TreeNode curNode, TreeNode targetP, TreeNode targetQ){
        if (curNode == null){
            return null;
        }
15     if (curNode == targetP || curNode == targetQ){
        return curNode;
    }
    TreeNode leftChild = findChild(curNode.left, targetP, targetQ);
    TreeNode rightChild = findChild(curNode.right, targetP, targetQ);
20     if (leftChild != null && rightChild != null){
        return curNode;
    }
    return (leftChild == null) ? rightChild : leftChild;
}
  
```

```

    }
25    public TreeNode lowestCommonAncestor(TreeNode root , TreeNode p, TreeNode q) {
        TreeNode lca = findChild(root , p, q);
        return lca;
    }
}

```

Solution 2

Still, use path to check LCA, but this time a little complicated.

Listing 172: Solution 2

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5    *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public List<TreeNode> findPath(TreeNode curNode, TreeNode targetNode,
        List<TreeNode> curPath){
        if (curNode == null){
            return null;
15        }
        curPath.add(curNode);
        if (curNode == targetNode){
            return new ArrayList<>(curPath);
        }
        List<TreeNode> leftPath = findPath(curNode.left , targetNode, curPath);
        List<TreeNode> rightPath = findPath(curNode.right , targetNode, curPath);
        curPath.remove(curPath.size() - 1);
        return (leftPath == null) ? rightPath : leftPath;
20    }

25    public TreeNode lowestCommonAncestor(TreeNode root , TreeNode p, TreeNode q) {
        List<TreeNode> pathP = new ArrayList<TreeNode>();
        List<TreeNode> pathQ = new ArrayList<TreeNode>();
        pathP = findPath(root , p, new ArrayList<TreeNode>());
        pathQ = findPath(root , q, new ArrayList<TreeNode>());
30
        for (int idx = 0 ; idx < Math.min(pathP.size() , pathQ.size()) ; idx ++){
            if (pathP.get(idx) != pathQ.get(idx)){
                return pathP.get(idx - 1);
35            }
        }
        // lca is the tail of shorter path
        return (pathP.size() > pathQ.size()) ? pathQ.get(pathQ.size() - 1)
            : pathP.get(pathP.size() - 1);
40    }
}

```

Problem 128

222 Count Complete Tree Nodes

Given a complete binary tree, count the number of nodes.

Definition of a complete binary tree from Wikipedia:

In a complete binary tree every level, except possibly the last, is completely filled, and all nodes in the last level are as far left as possible.

It can have between 1 and 2^h nodes inclusive at the last level h .

<https://leetcode.com/problems/count-complete-tree-nodes/>

Solution 1

Recursive solution. Firstly, simply applying recursion to check all nodes in $O(n)$ time will get TLE, so we need optimization. Given the tree is a complete binary tree, then when left height of left subtree equals to right height of right subtree, the tree is a full binary tree, so the number of nodes can be worked out easily through $2^{\text{height}} - 1$.

Note: this is not enough. Using inefficient function like 'Math.pow' will still get TLE, use $1 \ll \text{height}$ instead.

Besides, there is another thing we can optimize is to save the depth of subtree calculated since when we move curNode to its left child, the height of the new left subtree is $\text{height}(\text{former}) - 1$, unnecessary to calculate again.

Listing 173: Solution 1

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public int getLeftHeight(TreeNode curNode){
        if (curNode == null){
            return 0;
        }
15    int height = 0;
    while (curNode != null){
        height ++;
        curNode = curNode.left ;
    }
20    return height;
}
    public int getRightHeight(TreeNode curNode){
        if (curNode == null){
            return 0;
25        }
        int height = 0;

```

```
        while (curNode != null){
            height ++;
            curNode = curNode.right;
30    }
        return height;
    }
    public int countNodes(TreeNode root) {
        int count = 0;
35    if (root == null){
        return 0;
    }
    if (getLeftHeight(root.left) == getRightHeight(root.right)){
        return (1 << getLeftHeight(root)) - 1;
40    }
    else{
        return 1 + countNodes(root.left) + countNodes(root.right);
    }
    }
45 }
```

Solution 2

Iterative solution from <https://segmentfault.com/a/1190000003818177>. Look at the left maximum height of left subtree and right subtree, if left equals right, then left tree is full, the end node is at the right subtree and if left is larger, it means that the end node is in left subtree.

Listing 174: Solution 2

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public int getLeftHeight(TreeNode curNode){
        if (curNode == null){
            return 0;
        }
15    int height = 0;

    while (curNode != null){
        height ++;
        curNode = curNode.left;
20    }

    return height;
}

25 public int countNodes(TreeNode root) {
    int count = 0;
```

```
    if (root == null){
        return 0;
    }

    30    TreeNode curNode = root;
    while (curNode != null){
        int lHeight = getLeftHeight(curNode.left);
        35    int rHeight = getLeftHeight(curNode.right);

        if (lHeight > rHeight){
            // +1 means adding the root node
            count += ((1 << rHeight) - 1) + 1;
            curNode = curNode.left;
        }
        40    else {
            // lHeight == rHeight
            count += ((1 << lHeight) - 1) + 1;
            curNode = curNode.right;
        }
        45    }
    }

    return count;
}
50 }
```


Problem 129

108 Convert Sorted Array to Binary Search Tree

Given a complete binary tree, count the number of nodes.

Given an array where elements are sorted in ascending order, convert it to a height balanced BST.

<https://leetcode.com/problems/convert-sorted-array-to-binary-search-tree/>

Solution

A binary search problem. Root node is the middle element where left subtree and right subtree are formed from elements from both sides of middle elements.

Listing 175: Solution

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public TreeNode buildNode(int[] nums, int start, int end){
        if (start > end){
            return null;
        }
15     int mid = start + (end - start) / 2;
    TreeNode rootNode = new TreeNode(nums[mid]);
    rootNode.left = buildNode(nums, start, mid - 1);
    rootNode.right = buildNode(nums, mid + 1, end);

20     return rootNode;
}

    public TreeNode sortedArrayToBST(int[] nums) {
        return buildNode(nums, 0, nums.length - 1);
25    }
}
```

Problem 130

109 Convert Sorted List to Binary Search Tree

Given a singly linked list where elements are sorted in ascending order, convert it to a height balanced BST.

<https://leetcode.com/problems/convert-sorted-list-to-binary-search-tree/>

Solution

Certainly we can use the solution in previous “Convert Sorted Array to Binary Search Tree”, however, it will cost us extra runtime to get the middle element of list(in array, finding a element is of $O(1)$ runtime).

So we go for another direction. We can find out that the sorted list is exactly what an inorder traversal of a tree looks like, so we transform the print(or other) operations with iterating the list and form a rootNode.

There is another problem which is related but (much?) more difficult than this one, of “Convert Binary Search Tree (BST) to Sorted Doubly-Linked List”, worth pondering.

See <http://articles.leetcode.com/convert-binary-search-tree-bst-to/>

Listing 176: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
/**
10  * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
15  *     TreeNode(int x) { val = x; }
 * }
 */
public class Solution {
    private ListNode currentListNode;
20    public TreeNode buildTree(int low, int high){
        if (low > high){
            return null;
        }
        int mid = low + (high - low) / 2;
25    TreeNode leftChild = buildTree(low, mid - 1);

        TreeNode rootNode = new TreeNode(currentListNode.val);
        currentListNode = currentListNode.next;

        rootNode.left = leftChild;
30    TreeNode rightChild = buildTree(mid + 1, high);
        rootNode.right = rightChild;
        return rootNode;
    }
}
```

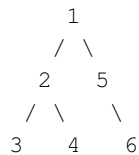
```
    }  
35  
    public TreeNode sortedListToBST(ListNode head) {  
        if (head == null){  
            return null;  
        }  
40        currentListNode = head;  
        ListNode curNode = head;  
        int length = 0;  
        while (curNode != null){  
            length ++;  
45            curNode = curNode.next;  
        }  
  
        return buildTree(0 , length - 1);  
    }  
50 }
```

Problem 131

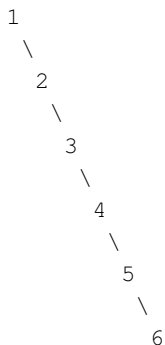
114 Flatten Binary Tree to Linked List

Given a binary tree, flatten it to a linked list in-place.

For example,
Given



The flattened tree should look like:



Hints:

If you notice carefully in the flattened tree,
each node's right child points to the next node of a pre-order traversal.

<https://leetcode.com/problems/flatten-binary-tree-to-linked-list/>

Solution

The solution of this one is in fact direct. We can easily notice that the new tree structure follows the preorder traversal order. Then we just:

1. check if `curNode.left == null`, if not, recurse on its right node(since in preorder traversal, right node comes latter from root node)
2. if `curNode.left != null`, then left child should be the next element in preorder traversal, so `root.left = null`, `root.right = formerLeft`. BUT what about the former right node? Remember that former right node always follows after all nodes in left subtree of rootNode when doing preorder traversal, so we just append them to the right child of the rightMost node of formerLeft(namely current `root.right`).
3. repeat 1 & 2

Listing 177: Solution

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
```

```

5      *      int val;
      *      TreeNode left;
      *      TreeNode right;
      *      TreeNode(int x) { val = x; }
      *  }
      */
10 public class Solution {
      public void flatten(TreeNode root) {
          if (root == null){
              return;
          }

15          if (root.left != null){
              TreeNode formerLeft = root.left;
              TreeNode formerRight = root.right;
              root.left = null;
20              root.right = formerLeft;

              TreeNode rightMost = formerLeft;
              while(rightMost.right != null){
                  rightMost = rightMost.right;
25              }
              rightMost.right = formerRight;
          }
          flatten(root.right);
30      }
}
```

Problem 132

341 Flatten Nested List Iterator

Given a nested list of integers, implement an iterator to flatten it.

Each element is either an integer, or a list --
whose elements may also be integers or other lists.

Example 1:

Given the list `[[1,1],2,[1,1]]`,

By calling `next` repeatedly until `hasNext` returns false,
the order of elements returned by `next` should be: `[1,1,2,1,1]`.

Example 2:

Given the list `[1,[4,[6]]]`,

By calling `next` repeatedly until `hasNext` returns false,
the order of elements returned by `next` should be: `[1,4,6]`.

<https://leetcode.com/problems/flatten-nested-list-iterator/>

Solution

Seemingly I list 2 solutions here, they are in nature similar, the difference just lies in when we flatten the list (to flatten the whole list or flatten it every time we pop an element out of the stack.)

Listing 178: Solution - Flatten it using Stack

```
/**
 * // This is the interface that allows for creating nested lists.
 * // You should not implement it, or speculate about its implementation
 * public interface NestedInteger {
5  *
 *      // @return true if this NestedInteger holds a single integer, rather than a nested list
 *      public boolean isInteger();
 *
 *      // @return the single integer that this NestedInteger holds, if it holds a single integer
10  *      // Return null if this NestedInteger holds a nested list
 *      public Integer getInteger();
 *
 *      // @return the nested list that this NestedInteger holds, if it holds a nested list
 *      // Return null if this NestedInteger holds a single integer
15  *      public List<NestedInteger> getList();
 *  }
 */
public class NestedIterator implements Iterator<Integer> {
    private Stack<NestedInteger> stack = new Stack<NestedInteger>();
20  private NestedInteger curInt;

    public NestedIterator(List<NestedInteger> nestedList) {
        push(nestedList);
    }
25 }
```

```

    @Override
    public Integer next() {
        return curInt.getInteger();
    }

    @Override
    public boolean hasNext() {
        while (!stack.isEmpty()) {
            curInt = stack.pop();
            if (curInt.isInteger()) {
                return true;
            }
            push(curInt.getList());
        }
        return false;
    }

    private void push(List<NestedInteger> list) {
        for (int nIdx = list.size() - 1; nIdx >= 0; nIdx --) {
            stack.push(list.get(nIdx));
        }
    }
}

/**
 * Your NestedIterator object will be instantiated and called as such:
 * NestedIterator i = new NestedIterator(nestedList);
 * while (i.hasNext()) v[f()] = i.next();
 */

```

Listing 179: Solution - Flatten the whole list at first sight

```

/**
 * // This is the interface that allows for creating nested lists.
 * // You should not implement it, or speculate about its implementation
 * public interface NestedInteger {
 *
 *     // @return true if this NestedInteger holds a single integer, rather than a nested list
 *     public boolean isInteger();
 *
 *     // @return the single integer that this NestedInteger holds, if it holds a single integer
 *     // Return null if this NestedInteger holds a nested list
 *     public Integer getInteger();
 *
 *     // @return the nested list that this NestedInteger holds, if it holds a nested list
 *     // Return null if this NestedInteger holds a single integer
 *     public List<NestedInteger> getList();
 * }
 */
public class NestedIterator implements Iterator<Integer> {
    List<Integer> list = new ArrayList<Integer>();
    Iterator<Integer> ite;

    public NestedIterator(List<NestedInteger> nestedList) {

```

```

    for (NestedInteger nestedInt : nestedList){
        flatten(nestedInt);
25    }
    ite = list.iterator();
}

private void flatten(NestedInteger nestedInt){
30    if (nestedInt.isInteger()){
        list.add(nestedInt.getInteger());
    }else{
        List<NestedInteger> nestedList = nestedInt.getList();
        for (NestedInteger nested : nestedList){
35            flatten(nested);
        }
    }
}

@Override
40 public Integer next() {
    return ite.next();
}

@Override
45 public boolean hasNext() {
    return ite.hasNext();
}

50 }

/**
 * Your NestedIterator object will be instantiated and called as such:
 * NestedIterator i = new NestedIterator(nestedList);
55 * while (i.hasNext()) v[f()] = i.next();
 */
```


Problem 133

50 Pow(x, n)

Implement `pow(x, n)`.

<https://leetcode.com/problems/powx-n/>

Solution

Implementing using for loop will get TLE(I can hardly imagine how slow loops can be). Moreover, the edge cases are pretty interesting.

So use the formula in figure:

$$a^n = \begin{cases} a^{n/2} \cdot a^{n/2} & n \text{ 为偶数} \\ a^{(n-1)/2} \cdot a^{(n-1)/2} \cdot a & n \text{ 为奇数} \end{cases}$$

Lots of things to consider!! BTW, use `n >>> 1`, not `n >> 1` to substitute `n / 2`.

Listing 180: Solution

```

public class Solution {
    public double myPow(double x, int n) {
        if (x - 0.0000001 == 0){
            return 0.0;
        }
        if (n < 0){
            return (1 / pow(x, (long)(-n)));
        }
        return pow(x, (long)n);
    }
    public double pow(double x, long n){
        if (n == 0){
            return 1.0;
        }
        if (n == 1){
            return x;
        }
        double sqrtVal = pow(x, (n >>> 1));
        if ((n & 1) == 1){
            // odd
            return x * sqrtVal * sqrtVal;
        }
        else{
            return sqrtVal * sqrtVal;
        }
    }
}

```

Problem 134

69 Sqrt(x)

Implement `int sqrt(int x)`.

Compute and return the square root of `x`.

<https://leetcode.com/problems/sqrtx/>

Solution

Brute force will definitely TLE. Use binary search to solve the problem. Be aware of the final result returned. For example, $\text{sqrt}(3) = 1$, but $1 * 1 = 1 < 3$ ($3/1 > 1$)

Listing 181: Solution

```
public class Solution {
    public int mySqrt(int x) {
        if (x == 0 || x == 1){
            return x;
        }
        int start = 1;
        int end = x - 1;
        int mid = start + (end - start) / 2;
        while (start <= end){
            mid = start + (end - start) / 2;
            if (mid == x / mid){
                return mid;
            }
            else{
                if (mid > x / mid){
                    end = mid - 1;
                }
                else{
                    start = mid + 1;
                }
            }
        }
        if (mid > x / mid){
            return mid - 1;
        }
        else{
            return mid;
        }
    }
}
```

Problem 135

65 Valid Number

Validate if a given string is numeric.

Some examples:

"0" => true

" 0.1 " => true

"abc" => false

"1 a" => false

"2e10" => true

Note: It is intended for the problem statement to be ambiguous.

You should gather all requirements up front before implementing one.

<https://leetcode.com/problems/valid-number/>

Solution

RIIts difficulty lies in so many edge cases of the problem.

Listing 182: Solution

```
public class Solution {
    public boolean isNumber(String s) {
        int idx = 0;
        // skip spaces in the front
5       while(idx < s.length() && s.charAt(idx) == ' '){
            idx ++;
        }

        if (idx < s.length() && (s.charAt(idx) == '+' || s.charAt(idx) == '-')){
10          idx ++;
        }
        boolean isNumeric = false;
        while (idx < s.length() && Character.isDigit(s.charAt(idx))){
15          isNumeric = true;
            idx ++;
        }

        if (idx < s.length() && s.charAt(idx) == '.'){
20          idx ++;
            while (idx < s.length() && Character.isDigit(s.charAt(idx))){
                isNumeric = true;
                idx ++;
            }
        }

25
        if (idx < s.length() && isNumeric && (s.charAt(idx) == 'e'
        || s.charAt(idx) == 'E')){
            idx ++;
            isNumeric = false;
30          if (idx < s.length() && (s.charAt(idx) == '+'
            || s.charAt(idx) == '-')){
                idx ++;
            }
        }
    }
}
```

```
        }  
35         while (idx < s.length() && Character.isDigit(s.charAt(idx))) {  
            isNumeric = true;  
            idx ++;  
        }  
40     }  
  
    // skip spaces in the end  
    while (idx < s.length() && s.charAt(idx) == ' '){  
        idx ++;  
    }  
45    return isNumeric && idx == s.length();  
    }  
}
```

Problem 136

124 Binary Tree Maximum Path Sum

Given a binary tree, find the maximum path sum.

For this problem, a path is defined as any sequence of nodes from some starting node to any node in the tree along the parent-child connections. The path does not need to go through the root.

For example:

Given the below binary tree,

```
      1
     / \
    2   3
Return 6.
```

<https://leetcode.com/problems/binary-tree-maximum-path-sum/>

Solution

It's hard to distinguish whether the problem is hard or not. The idea of solution is the most important.

Listing 183: Solution

```
/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
 *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    int maxSum = Integer.MIN_VALUE;
    public int findMax(TreeNode root){
        if (root == null){
            return 0;
        }
        15     int leftMax = findMax(root.left);
        int rightMax = findMax(root.right);
        maxSum = Math.max(maxSum , root.val + leftMax + rightMax);
        int maxPath = Math.max(leftMax , rightMax) + root.val;
        20     return maxPath > 0 ? maxPath : 0;
    }
    public int maxPathSum(TreeNode root) {
        findMax(root);
        25     return maxSum;
    }
}
```

Problem 137

26 Remove Duplicates from Sorted Array

Given a sorted array, remove the duplicates in place such that each element appear only once and return the new length.

Do not allocate extra space for another array, you must do this in place with constant memory.

For example,

Given input array nums = [1,1,2],

Your function should return length = 2, with the first two elements of nums being 1 and 2 respectively. It doesn't matter what you leave beyond the new length.

<https://leetcode.com/problems/remove-duplicates-from-sorted-array/>

Solution

Keep a clean thought will help you solve the problem quickly.

Listing 184: Solution

```
public class Solution {  
    public int removeDuplicates(int[] nums) {  
        if (nums.length == 0){  
            return 0;  
        }  
        int formerElement = nums[0];  
        int length = 1;  
  
        for (int idx = 0 ; idx < nums.length ; idx ++){  
            if (nums[idx] != formerElement){  
                nums[length++] = nums[idx];  
                formerElement = nums[idx];  
            }  
        }  
        return length;  
    }  
}
```

Problem 138

80 Remove Duplicates from Sorted Array II

Follow up for "Remove Duplicates":
What if duplicates are allowed at most twice?

For example,

Given sorted array `nums = [1,1,1,2,2,3]`,

Your function should return `length = 5`,
with the first five elements of `nums` being 1, 1, 2, 2 and 3.
It doesn't matter what you leave beyond the new length.

<https://leetcode.com/problems/remove-duplicates-from-sorted-array-ii/>

Solution

Add a flag to check if we still allow a same element exists.

Listing 185: Solution

```
public class Solution {
    public int removeDuplicates(int[] nums) {
        if (nums.length == 0){
            return nums.length;
        }
        int length = 1;
        int formerElement = nums[0];
        boolean allowAnother = true;

        for (int idx = 1 ; idx < nums.length ; idx ++){
            if (formerElement != nums[idx]){
                formerElement = nums[idx];
                nums[length++] = nums[idx];
                allowAnother = true;
            }
            else{
                // former == current
                if (allowAnother){
                    nums[length++] = nums[idx];
                    allowAnother = false;
                }
            }
        }

        return length;
    }
}
```

Problem 139

27 Remove Element

Given an array and a value, remove all instances of that value in place and return the new length.

The order of elements can be changed. It doesn't matter what you leave beyond the new length.

<https://leetcode.com/problems/remove-element/>

Solution

So stupid to get 2 WA and a RE on such problem.....

Listing 186: Solution

```
public class Solution {  
    public int removeElement(int[] nums, int val) {  
        int idx = 0;  
        int length = nums.length;  
  
        while (idx < length){  
            if (nums[idx] != val){  
                idx ++;  
                continue;  
            }  
            while (length > 0 && idx < length && nums[idx] == val){  
                int swap = nums[idx];  
                nums[idx] = nums[length - 1];  
                nums[length - 1] = swap;  
                length --;  
            }  
            idx ++;  
        }  
  
        return length;  
    }  
}
```


Problem 140

283 Move Zeroes

Given an array `nums`, write a function to move all 0's to the end of it while maintaining the relative order of the non-zero elements.

For example, given `nums = [0, 1, 0, 3, 12]`,
after calling your function, `nums` should be `[1, 3, 12, 0, 0]`.

Note:

1. You must do this in-place without making a copy of the array.
2. Minimize the total number of operations.

<https://leetcode.com/problems/move-zeroes/>

Solution

An easy problem.

Listing 187: Solution

```
public class Solution {  
    public void moveZeroes(int[] nums) {  
        int nonZeroLength = 0;  
        int idx = 0;  
        while (idx < nums.length){  
            if (nums[idx] != 0){  
                if (idx != nonZeroLength){  
                    nums[nonZeroLength] = nums[idx];  
                    nums[idx] = 0;  
                }  
                nonZeroLength ++;  
            }  
            idx ++;  
        }  
    }  
}
```

Problem 141

138 Copy List with Random Pointer

A linked list is given such that each node contains an additional random pointer which could point to any node in the list or null.

Return a deep copy of the list.

<https://leetcode.com/problems/copy-list-with-random-pointer/>

Solution 1

Solution 1 makes use of HashMap, mapping original node and copied node so that we can get the random node by querying from HashMap. Use a 2-pass solution to respectively deal with next pointer and random pointer.

Listing 188: Solution 1

```
/**
 * Definition for singly-linked list with a random pointer.
 * class RandomListNode {
 *     int label;
 *     RandomListNode next, random;
 *     RandomListNode(int x) { this.label = x; }
 * };
 */
public class Solution {
    public RandomListNode copyRandomList(RandomListNode head) {
        if (head == null){
            return null;
        }
        HashMap<RandomListNode, RandomListNode> copyMap
        = new HashMap<RandomListNode, RandomListNode>();

        RandomListNode curNode = head;
        RandomListNode dummy = new RandomListNode(-1);
        RandomListNode curCopyNode = dummy;

        while (curNode != null){
            RandomListNode newCopyNode = new RandomListNode(curNode.label);
            curCopyNode.next = newCopyNode;
            copyMap.put(curNode, newCopyNode);
            curNode = curNode.next;
            curCopyNode = curCopyNode.next;
        }

        curNode = head;
        curCopyNode = dummy;

        while (curNode != null){
            curCopyNode.next.random = copyMap.get(curNode.random);
            curNode = curNode.next;
            curCopyNode = curCopyNode.next;
        }
    }
}
```

```

    }

    return dummy.next;
}
40 }
```

Solution 2

Solution 2 changes the structure of the list. See figure



So that we can just use

$node.next.random = node.random.next$

Listing 189: Solution 2

```

/**
 * Definition for singly-linked list with a random pointer.
 * class RandomListNode {
 *     int label;
5  *     RandomListNode next, random;
 *     RandomListNode(int x) { this.label = x; }
 * };
 */
public class Solution {
10     public RandomListNode copyRandomList(RandomListNode head) {
        if (head == null){
            return null;
        }
        RandomListNode curNode = head;
15     while(curNode != null){
        RandomListNode copyNode = new RandomListNode(curNode.label);
        RandomListNode nextNode = curNode.next;
        curNode.next = copyNode;
        copyNode.next = nextNode;

        curNode = curNode.next.next;
    }

    curNode = head;
25     while (curNode != null){
        curNode.next.random = (curNode.random != null)
            ? curNode.random.next : null;
        curNode = curNode.next.next;
    }

30     //recover the original structure, seperate the two list
    curNode = head;
    RandomListNode dummy = new RandomListNode(-1);
    RandomListNode tail = dummy;
```

```
35         while (curNode != null){
            RandomListNode copyNode = curNode.next;
            RandomListNode nextNode = copyNode.next;
            tail.next = copyNode;
40             tail = tail.next;

            curNode.next = nextNode;
            curNode = nextNode;
        }
45     return dummy.next;
    }
```

Problem 142

207 Course Schedule

There are a total of n courses you have to take, labeled from 0 to $n - 1$.

Some courses may have prerequisites, for example to take course 0 you have to first take course 1, which is expressed as a pair: [0,1]

Given the total number of courses and a list of prerequisite pairs, is it possible for you to finish all courses?

For example:

2, [[1,0]]

There are a total of 2 courses to take.

To take course 1 you should have finished course 0. So it is possible.

2, [[1,0],[0,1]]

There are a total of 2 courses to take.

To take course 1 you should have finished course 0, and to take course 0 you should also have finished course 1. So it is impossible.

Note:

The input prerequisites is a graph represented by a list of edges, not adjacency matrices.

Hints:

1. This problem is equivalent to finding if a cycle exists in a directed graph. If a cycle exists, no topological ordering exists and therefore it will be impossible to take all courses.
2. Topological Sort via DFS - A great video tutorial (21 minutes) on Coursera explaining the basic concepts of Topological Sort.
3. Topological sort could also be done via BFS.

<https://leetcode.com/problems/course-schedule/>

Solution 1

A graph problem, use topological sort to solve. Equivalent to problems which identify whether there is a cycle in the graph. Use adjacency list to store edges for a easier calculation on indegrees. BFS sort.

Listing 190: Solution 1

```
public class Solution {
    public boolean canFinish(int numCourses, int[][] prerequisites) {
        if (prerequisites.length == 0){
            return true;
        }
        int[] indegree = new int[numCourses];
        List<List<Integer>> adjacency = new ArrayList<List<Integer>>(numCourses);
        for (int i = 0 ; i < numCourses ; i++){
            adjacency.add(new ArrayList<Integer>());
        }
    }
}
```

```

10     }

    for (int i = 0 ; i < prerequisites.length ; i++){
        adjacency.get(prerequisites[i][1]).add(prerequisites[i][0]);
        indegree[prerequisites[i][0]] ++;
15     }

    Queue<Integer> outQueue = new LinkedList<Integer>();

    for (int courseId = 0 ; courseId < numCourses; courseId++){
20         if (indegree[courseId] == 0){
            outQueue.offer(courseId);
        }
    }

25     int courseCount = 0;
    while (!outQueue.isEmpty()){
        int courseId = outQueue.poll();

        List<Integer> latterCourses = adjacency.get(courseId);
30         for (int latterCourseId : latterCourses){
            indegree[latterCourseId] --;
            if (indegree[latterCourseId] == 0){
                outQueue.offer(latterCourseId);
            }
35         }

        courseCount ++;
    }

40     return courseCount == numCourses;

    }
}

```

Solution 2

Solution 2 uses DFS. We should explain the usage of array 'visited'. For value 0, it means the node has not been visited at all; value 1 means that the adjacent nodes has already been visited; while value -1 means the node has been visited while the adjacencies are not totally visited.

Listing 191: Solution 2

```

public class Solution {

    private boolean traverse(int courseId, ArrayList<Integer>[] adjacency
5        ,int[] visited ){
        visited[courseId] = -1;
        for (int adjacencyEle : adjacency[courseId]){
            if (visited[adjacencyEle] == -1){
                return false;
            }
10         else if (visited[adjacencyEle] == 0){

```

```
        if (!traverse(adjacencyEle, adjacency, visited)){
            return false;
        }
    }
15     }

    visited[courseIdx] = 1;
    return true;
}
20

public boolean canFinish(int numCourses, int [][] prerequisites) {
    ArrayList<Integer>[] adjacency = new ArrayList[numCourses];
    for (int i = 0 ; i < numCourses ; i++){
        adjacency[i] = new ArrayList<Integer>();
25     }
    for (int i = 0 ; i < prerequisites.length ; i++){
        adjacency[prerequisites[i][1]].add(prerequisites[i][0]);
    }

30     int [] visited = new int[numCourses];
    for (int courseIdx = 0 ; courseIdx < numCourses ; courseIdx++){
        if (visited[courseIdx] == 0){
            if (!traverse(courseIdx, adjacency, visited)){
35                 return false;
            }
        }
    }

40     return true;
}
}
```

Problem 143

210 Course Schedule II

There are a total of n courses you have to take, labeled from 0 to $n - 1$.

Some courses may have prerequisites, for example to take course 0 you have to first take course 1, which is expressed as a pair: [0,1]

Given the total number of courses and a list of prerequisite pairs, return the ordering of courses you should take to finish all courses.

There may be multiple correct orders, you just need to return one of them. If it is impossible to finish all courses, return an empty array.

For example:

2, [[1,0]]

There are a total of 2 courses to take.

To take course 1 you should have finished course 0.

So the correct course order is [0,1]

4, [[1,0],[2,0],[3,1],[3,2]]

There are a total of 4 courses to take.

To take course 3 you should have finished both courses 1 and 2.

Both courses 1 and 2 should be taken after you finished course 0.

So one correct course order is [0,1,2,3]. Another correct ordering is [0,2,1,3].

Note:

The input prerequisites is a graph represented by a list of edges, not adjacency matrices.

Hints:

1. This problem is equivalent to finding if a cycle exists in a directed graph. If a cycle exists, no topological ordering exists and therefore it will be impossible to take all courses.
2. Topological Sort via DFS - A great video tutorial (21 minutes) on Coursera explaining the basic concepts of Topological Sort.
3. Topological sort could also be done via BFS.

<https://leetcode.com/problems/course-schedule-ii/>

Solution

Similar to “Course Schedule”, add a step to record each step, namely record the sequence polled from queue.

Listing 192: Solution

```
public class Solution {  
    public int[] findOrder(int numCourses, int[][] prerequisites) {  
        // ADD this statement will get a WA  
        // if (prerequisites.length == 0){  
5         //     return new int[numCourses];  
    }  
}
```



```
// }
int[] indegree = new int[numCourses];

List<List<Integer>> adjacency = new ArrayList<List<Integer>>();
10 for (int i = 0 ; i < numCourses ; i++){
    adjacency.add(new ArrayList<Integer>());
}

15 for (int preIdx = 0 ; preIdx < prerequisites.length ; preIdx++){
    adjacency.get(prerequisites[preIdx][1]).
    add(prerequisites[preIdx][0]);
    indegree[prerequisites[preIdx][0]]++;
}

20 Queue<Integer> startingCourses = new LinkedList<Integer>();
for (int courseIdx = 0 ; courseIdx < numCourses ; courseIdx++){
    if (indegree[courseIdx] == 0){
        startingCourses.offer(courseIdx);
    }
25 }

int[] result = new int[numCourses];
int courseNum = 0;
while (!startingCourses.isEmpty()){
30     int courseIdx = startingCourses.poll();
    for (int latterCourseIdx : adjacency.get(courseIdx)){
        if (--indegree[latterCourseIdx] == 0){
            startingCourses.offer(latterCourseIdx);
        }
35     }

    result[courseNum] = courseIdx;
    courseNum++;
}
40 return (courseNum == numCourses) ? result : new int[0];
}
```

Problem 144

1 Two Sum

Given an array of integers, return indices of the two numbers such that they add up to a specific target.

You may assume that each input would have exactly one solution.

Example:

Given `nums = [2, 7, 11, 15]`, `target = 9`,

Because `nums[0] + nums[1] = 2 + 7 = 9`,
return `[0, 1]`.

<https://leetcode.com/problems/two-sum/>

Solution

It's a very interesting series of problems.

For all kSum conclusion, there is a good blog here

<http://www.sigmainfy.com/blog/summary-of-ksum-problems.html>

As for this problem, there are mainly two ways, except for the Brute force method. First way is to use HashMap.

Solution utilizes two pointers to get the **unique** result also works. Sorting the list first is needed as a prerequisite. However, we need to keep the original index of number array. So here we just list the HashMap solution.

Listing 193: Solution

```
public class Solution {
    public int[] twoSum(int[] nums, int target) {
        int[] result = new int[2];
        HashMap<Integer, Integer> valueIdxMap = new HashMap<Integer, Integer>();

        for (int idx = 0 ; idx < nums.length ; idx++){
            if (valueIdxMap.containsKey(target - nums[idx])){
                result[0] = valueIdxMap.get(target - nums[idx]);
                result[1] = idx;
                break;
            }
            valueIdxMap.put(nums[idx], idx);
        }

        return result;
    }
}
```

Problem 145

15 3Sum

Given an array S of n integers, are there elements a , b , c in S such that $a + b + c = 0$? Find all unique triplets in the array which gives the sum of zero.

Note:

Elements in a triplet (a,b,c) must be in non-descending order. (ie, $a \leq b \leq c$)
The solution set must not contain duplicate triplets.

For example, given array $S = \{-1\ 0\ 1\ 2\ -1\ -4\}$,

A solution set is:

$(-1, 0, 1)$

$(-1, -1, 2)$

<https://leetcode.com/problems/3sum/>

Solution

Finally we can use the 2 pointers solution.

Note that in the problem, the description request us to prevent from producing duplicate triples, so we need a Set or just skip those duplicate elements to solve this problem.

Listing 194: Solution

```
public class Solution {
    public List<List<Integer>> threeSum(int[] nums) {
        List<List<Integer>> results = new ArrayList<List<Integer>>();
        Arrays.sort(nums);

        for (int idx1 = 0 ; idx1 < nums.length - 2 ; idx1 ++){
            while (idx1 > 0 && idx1 < nums.length - 2
                && nums[idx1] == nums[idx1 - 1]){
                idx1 ++;
            }
            int val1 , val2 , val3;
            val1 = nums[idx1];

            int idx2 = idx1 + 1, idx3 = nums.length - 1;
            while (idx2 < idx3){
                val2 = nums[idx2];
                val3 = nums[idx3];

                int sum = val1 + val2 + val3;

                if (sum == 0){
                    List<Integer> triple = new ArrayList<Integer>();
                    triple.add(val1);
                    triple.add(val2);
                    triple.add(val3);

                    results.add(triple);
                }
            }
        }
    }
}
```

```

    while(idx2 < idx3 && nums[idx3] == nums[idx3 - 1]){
        idx3 --;
    }
    idx3 --;
    idx2++;
    while(idx2 < idx3 && nums[idx2 - 1] == nums[idx2]){
        idx2 ++;
    }
    }else if (sum > 0){
        while(idx2 < idx3 && nums[idx3] == nums[idx3 - 1]){
            idx3 --;
        }
        idx3 --;
    }else{
        idx2 ++;
        while(idx2 < idx3 && nums[idx2 - 1] == nums[idx2]){
            idx2 ++;
        }
    }
    }
    }

    return results;
}
}
```

Problem 146

16 3Sum Closest

Given an array S of n integers, find three integers in S such that the sum is closest to a given number, $target$.
Return the sum of the three integers.
You may assume that each input would have exactly one solution.

For example, given array $S = \{-1\ 2\ 1\ -4\}$, and $target = 1$.

The sum that is closest to the target is 2. $(-1 + 2 + 1 = 2)$.

<https://leetcode.com/problems/3sum-closest/>

Solution

Personally, I think this is a bit easier than 3sum since I don't need to care about the duplicates.

Listing 195: Solution

```
public class Solution {
    public int threeSumClosest(int[] nums, int target) {
        Arrays.sort(nums);

        int diff = Integer.MAX_VALUE;
        int closest = -1;

        for (int idx1 = 0 ; idx1 < nums.length - 2 ; idx1++){
            int idx2 = idx1 + 1;
            int idx3 = nums.length - 1;
            while (idx2 < idx3){
                int sum = nums[idx1] + nums[idx2] + nums[idx3];
                int curDiff = Math.abs(sum - target);
                if (curDiff == 0){
                    return target;
                }
                if (curDiff < diff){
                    diff = curDiff;
                    closest = sum;
                }
                if (sum < target){
                    idx2++;
                }
                if (sum > target){
                    idx3--;
                }
            }
        }

        return closest;
    }
}
```

Problem 147

18 4Sum

Given an array S of n integers, are there elements a , b , c , and d in S such that $a + b + c + d = \text{target}$?

Find all unique quadruplets in the array which gives the sum of target.

Note:

Elements in a quadruplet (a, b, c, d) must be in non-descending order.

(ie, $a \leq b \leq c \leq d$)

The solution set must not contain duplicate quadruplets.

For example, given array $S = \{1\ 0\ -1\ 0\ -2\ 2\}$, and $\text{target} = 0$.

A solution set is:

$(-1, 0, 0, 1)$

$(-2, -1, 1, 2)$

$(-2, 0, 0, 2)$

<https://leetcode.com/problems/4sum/>

Solution

Solving the kSum problem! Just one try!!!(Though not bug free...)

Listing 196: Solution

```
public class Solution {
    public List<List<Integer>> kSum(int k, int[] nums, int target,
        int start, int end){
        List<List<Integer>> result = new ArrayList<List<Integer>>();
5       if (k == 2){
            return twoSum(nums, target, start, end);
        }
        else{
            for (int idx = start; idx <= end; idx++){
10              while (idx > start && idx <= end && nums[idx] == nums[idx - 1]){
                  idx++;
              }
              if (idx > end){
                  break;
15              }
              int curEle = nums[idx];
              List<List<Integer>> oldList = kSum(k - 1, nums,
                  target - curEle, idx + 1, end);
              if (oldList != null){
20                  result.addAll(combine(curEle, oldList));
              }
            }
        }
        return (result.size() != 0) ? result : null;
25    }
    public List<List<Integer>> twoSum(int[] nums, int target,
        int start, int end){
        List<List<Integer>> result = new ArrayList<List<Integer>>();
```

```

    int idx1 = start;
    int idx2 = end;
    while (idx1 < idx2){
        int val1 = nums[idx1];
        int val2 = nums[idx2];
        int sum = val1 + val2;
        if (sum == target){
            List<Integer> newPair = new ArrayList<Integer>();
            newPair.add(val1);
            newPair.add(val2);
            result.add(newPair);
            idx2 --;
            while(idx2 >= start && idx2 < end
                && nums[idx2] == nums[idx2 + 1]){
                idx2 --;
            }
            idx1 ++;
            while (idx1 > start && idx1 <= end
                && nums[idx1] == nums[idx1 - 1]){
                idx1 ++;
            }
        }
        else if (sum > target){
            idx2 --;
        }
        else{
            idx1 ++;
        }
    }
    return (result.size() == 0) ? null : result;
}

public List<List<Integer>> combine(int ele, List<List<Integer>> oldList){
    List<List<Integer>> newList = new ArrayList<List<Integer>>();
    for (List<Integer> oldItem : oldList){
        List<Integer> newItem = new ArrayList<Integer>();
        newItem.add(ele);
        newItem.addAll(oldItem);
        newList.add(newItem);
    }
    return newList;
}

public List<List<Integer>> fourSum(int[] nums, int target) {
    Arrays.sort(nums);
    List<List<Integer>> result = kSum(4, nums, target, 0, nums.length - 1);
    return result == null ? new ArrayList<List<Integer>>() : result ;
}
}
```

Problem 148

55 Jump Game

Given an array of non-negative integers, you are initially positioned at the first index of the array.

Each element in the array represents your maximum jump length at that position.

Determine if you are able to reach the last index.

For example:

A = [2,3,1,1,4], return true.

A = [3,2,1,0,4], return false.

<https://leetcode.com/problems/jump-game/>

Solution

A greedy or DP problem. The maximum steps we can go on is `Math.max(curMax, nums[idx])`, don't forget to reduce `curMax` by 1 simulating every step we proceed.

Listing 197: Solution

```
public class Solution {
    public boolean canJump(int[] nums) {
        if (nums.length == 0 || nums.length == 1){
            return true;
        }
        int[] dp = new int[nums.length];
        int curMax = nums[0] + 1;

        for (int idx = 0 ; idx < nums.length ; idx++){
            curMax--;
            curMax = Math.max(nums[idx] , curMax);

            if (curMax == 0){
                return false;
            }
            if (curMax + idx >= nums.length - 1){
                return true;
            }
        }

        return true;
    }
}
```

A more concise version. Do not need to deduct 1, only need to check if we can reach the first elements finally.

Listing 198: Solution1.1

```
public class Solution {  
    public boolean canJump(int [] nums) {  
        int des = nums.length - 1;  
        for(int i = des - 1; i >= 0; i--){  
            if(nums[i] >= des - i){  
                des = i;  
            }  
        }  
        return des <= 0;  
    }  
}
```

Problem 149

45 Jump Game II

Given an array of non-negative integers,
you are initially positioned at the first index of the array.

Each element in the array represents
your maximum jump length at that position.

Your goal is to reach the last index in the minimum number of jumps.

For example:
Given array A = [2,3,1,1,4]

The minimum number of jumps to reach the last index is 2.
(Jump 1 step from index 0 to 1, then 3 steps to the last index.)

Note:
You can assume that you can always reach the last index.

<https://leetcode.com/problems/jump-game-ii/>

Solution

Still greedy, but needs some time to figure out the correctness...(Honestly, I don't think this is a complete greedy problem, since the solution is not that greedy. It's more like a DP problem.)

Let's say the range of the current jump is [curBegin, lastMax], curMax is the farthest point that all points in [curBegin, lastMax] can reach. Once the current point reaches lastMax, then trigger another jump, and set the new lastMax with curMax.

Listing 199: Solution

```
public class Solution {  
    public int jump(int[] nums) {  
        int curMax = 0;  
        int lastMax = 0;  
        int stepCnt = 0;  
  
        for (int currentIdx = 0 ; currentIdx < nums.length ; currentIdx ++){  
            // Of no use since you can always reach the last index  
            // if (i < curMax){  
            //     return -1;  
            // }  
  
            if (currentIdx > lastMax){  
                lastMax = curMax;  
                stepCnt ++;  
            }  
  
            curMax = Math.max(curMax, currentIdx + nums[currentIdx]);  
        }  
  
        return stepCnt;  
    }  
}
```

```
}  
}
```

Problem 150

120 Triangle

Given a triangle, find the minimum path sum from top to bottom.
Each step you may move to adjacent numbers on the row below.

For example, given the following triangle

```
[
  [2],
  [3,4],
  [6,5,7],
  [4,1,8,3]
]
```

The minimum path sum from top to bottom is 11 (i.e., $2 + 3 + 5 + 1 = 11$).

Note:

Bonus point if you are able to do this using only $O(n)$ extra space,
where n is the total number of rows in the triangle.

<https://leetcode.com/problems/triangle/>

Solution

A dp problem, I think it is maybe in an easier level of dp problems.

As for the bonus, manipulating on triangle variable itself is a fine idea. Just remember to mention the boundaries.

Listing 200: Solution

```
public class Solution {
    public int jump(int[] nums) {
        int curMax = 0;
        int lastMax = 0;
        int stepCnt = 0;

        for (int currentIdx = 0 ; currentIdx < nums.length ; currentIdx ++){
            // Of no use since you can always reach the last index
            // if (i < curMax){
            //     return -1;
            // }

            if (currentIdx > lastMax){
                lastMax = curMax;
                stepCnt ++;
            }

            curMax = Math.max(curMax, currentIdx + nums[currentIdx]);
        }

        return stepCnt;
    }
}
```

Problem 151

187 Repeated DNA Sequences

All DNA is composed of a series of nucleotides abbreviated as A, C, G, and T, for example: "ACGAATTCG". When studying DNA, it is sometimes useful to identify repeated sequences within the DNA.

Write a function to find all the 10-letter-long sequences (substrings) that occur more than once in a DNA molecule.

For example,

Given s = "AAAAACCCCCAAAAACCCCCCAAAAAGGGTTT",

Return:
["AAAAACCCCC", "CCCCCAAAAA"].

<https://leetcode.com/problems/repeated-dna-sequences/>

Solution

An interesting problem but not that easy. Using String Hashmap will get MLE. So we need to reduce the size of our data. That's when "Bit Manipulation" comes into sight. Since there are only 4 kinds of chars, we can mark them as '00','01','10','11', so that we can use a 20-bit integer as the hash code of each 10-digit string, as it won't cost as much memory as String. See code in listing.

Listing 201: Solution

```
public class Solution {
    public List<String> findRepeatedDnaSequences(String s) {
        List<String> result = new ArrayList<String>();
        if (s.length() <= 9){
            return new ArrayList<String>();
        }

        HashMap<Character, Integer> digitMapping
            = new HashMap<Character, Integer>();
        digitMapping.put('A', 0);
        digitMapping.put('C', 1);
        digitMapping.put('G', 2);
        digitMapping.put('T', 3);

        HashMap<Integer, Integer> countMapping
            = new HashMap<Integer, Integer>();

        int hash = 0;
        int curLen = 0;
        for (int strIdx = 0 ; strIdx < s.length() ; strIdx++){
            curLen++;

            hash = (hash << 2) + digitMapping.get(s.charAt(strIdx));
            if (curLen > 9) {
                hash = (hash & ((1 << 20) - 1));
            }
        }

        for (int i = 0; i < countMapping.size(); i++){
            int key = countMapping.keySet().toArray()[i];
            int value = countMapping.get(key);
            if (value > 1){
                result.add(s.substring(strIdx - 9, strIdx));
            }
        }

        return result;
    }
}
```

```
30         if (countMapping.containsKey(hash)
31             && countMapping.get(hash) == 1){
32             countMapping.put(hash, countMapping.get(hash) + 1);
33             String repeatedStr = s.substring(strIdx - 9, strIdx + 1);
34             result.add(repeatedStr);
35         }
36         else if (!countMapping.containsKey(hash)){
37             countMapping.put(hash, 1);
38         }
39     }
40     return result;
}
```

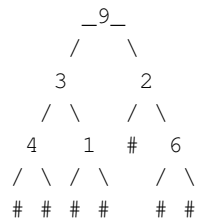
Problem 152

331 Verify Preorder Serialization of a Binary Tree

One way to serialize a binary tree is to use pre-order traversal.

When we encounter a non-null node, we record the node's value.

If it is a null node, we record using a sentinel value such as #.



For example, the above binary tree can be serialized

to the string "9,3,4,#,#,1,#,#,2,#,6,#,#", where # represents a null node.

Given a string of comma separated values,

verify whether it is

a correct preorder traversal serialization of a binary tree.

Find an algorithm without reconstructing the tree.

Each comma separated value in the string must be

either an integer or a character '#' representing null pointer.

You may assume that the input format is always valid,

for example it could never contain two consecutive commas such as "1,,3".

Example 1:

```
"9,3,4,#,#,1,#,#,2,#,6,#,#"
```

Return true

Example 2:

```
"1,#"
```

Return false

Example 3:

```
"9,#,#,1"
```

Return false

<https://leetcode.com/problems/verify-preorder-serialization-of-a-binary-tree/>

Solution

A stack problem may be more suitable for the orientation of this problem. The train of thought is somewhat like the Evaluation of Reverse Polish Expression.

When scrape nodes from a complete subtree together, we may combine them into a '#' to simplify the tree structure, finally we should only get a single '#' in the stack.

To be more specific, we first push every element we get into stack; if depth of current stack is larger or equal to 3 and the top 2 elements are both '#' when the 3rd element is not(which will never happen, even if the original string contains 3 successive '#', the first 2 will be combined into one single '#'). We pop the 2 '#'s out along with the 3rd non- '#' element and push an extra '#' instead.

Repeat the whole process.

Finally check if the remaining element in Stack is '#'.

Listing 202: Solution

```
public class Solution {
    public boolean isValidSerialization(String preorder) {
        Stack<String> serialization = new Stack<String>();
        String[] parts = preorder.split(",");
5      for (String part : parts){
            if (!part.equals("#")){
                serialization.push(part);
            }
            else{
10             if (serialization.isEmpty()
                || !(serialization.peek().equals("#"))){
                    serialization.push(part);
                }
            else{
15             while (!serialization.isEmpty()
                && serialization.peek().equals("#")){
                    String sharp = serialization.pop();

                    if (serialization.isEmpty()){
20                     return false;
                    }
                    String formerElement = serialization.pop();
                    if (formerElement.equals("#")){
                        return false;
25                     }
                    }
                serialization.push("#");
            }
        }
30    }

    if (serialization.isEmpty()){
        return false;
    }
35    if (!serialization.pop().equals("#")
        || !serialization.isEmpty()){
        return false;
    }
    return true;
40 }
}
```


Problem 153

72 Edit Distance

Given two words word1 and word2,
find the minimum number of steps required to convert word1 to word2.
(each operation is counted as 1 step.)

You have the following 3 operations permitted on a word:

- a) Insert a character
- b) Delete a character
- c) Replace a character

<https://leetcode.com/problems/edit-distance/>

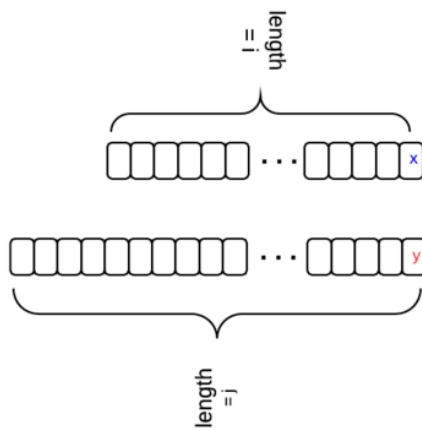
Solution

A representative DP problem.

The train of thought is listed as below:

Let $dp[i][j]$ stands for the edit distance between two strings with length i and j , i.e., $word1[0, \dots, i-1]$ and $word2[0, \dots, j-1]$.

There is a relation between $dp[i][j]$ and $dp[i-1][j-1]$. Let's say we transform from one string to another. The first string has length i and its last character is "x"; the second string has length j and its last character is "y". The following diagram shows the relation.



1. if $x == y$, then $dp[i][j] == dp[i-1][j-1]$
2. if $x \neq y$, and we insert y for $word1$, then $dp[i][j] = dp[i][j-1] + 1$
3. if $x \neq y$, and we delete x for $word1$, then $dp[i][j] = dp[i-1][j] + 1$
4. if $x \neq y$, and we replace x with y for $word1$, then $dp[i][j] = dp[i-1][j-1] + 1$
5. When $x \neq y$, $dp[i][j]$ is the min of the three situations.

Initial condition:

$dp[i][0] = i$, $dp[0][j] = j$

Listing 203: Solution

```
public class Solution {
    public int minDistance(String word1, String word2) {
```

```
int [][] minDis = new int[word1.length() + 1][word2.length() + 1];

5   for (int idx1 = 0 ; idx1 <= word1.length() ; idx1++){
        minDis[idx1][0] = idx1;
    }
    for (int idx2 = 0 ; idx2 <= word2.length() ; idx2++){
10      minDis[0][idx2] = idx2;
    }
    for (int idx1 = 1 ; idx1 <= word1.length() ; idx1++){
        for (int idx2 = 1 ; idx2 <= word2.length() ; idx2++){
            int diff = 0;
            if (word1.charAt(idx1 - 1) != word2.charAt(idx2 - 1)){
15              diff = 1;
            }
            minDis[idx1][idx2] = Math.min(minDis[idx1 - 1][idx2 - 1]
                + diff , Math.min(minDis[idx1 - 1][idx2] + 1
                , minDis[idx1][idx2 - 1] + 1));
20        }
    }

    return minDis[word1.length()][word2.length()];
25 }
```

Problem 154

319 Bulb Switcher

There are n bulbs that are initially off. You first turn on all the bulbs. Then, you turn off every second bulb. On the third round, you toggle every third bulb (turning on if it's off or turning off if it's on). For the i th round, you toggle every i bulb. For the n th round, you only toggle the last bulb. Find how many bulbs are on after n rounds.

Example:

Given $n = 3$.

At first, the three bulbs are [off, off, off].
After first round, the three bulbs are [on, on, on].
After second round, the three bulbs are [on, off, on].
After third round, the three bulbs are [on, off, off].

So you should return 1, because there is only one bulb is on.

<https://leetcode.com/problems/bulb-switcher/>

Solution

This is a very very very very interesting problem! Simulating the process is too time-consuming, so we need to find another way.

It can be easily discovered that when we switch the light for odd times, the state of light will change, while switching for even times won't. Now consider all numbers, for example, $12 = 1*12; 2*6; 3*4$, will be switched 6 times, its state won't change; however, as for $25 = 1*25; 5*5$, it will change its state!!

So the problem is exactly asking us to find how many square numbers are in the range of 1 to n .

Listing 204: Solution

```
public class Solution {  
    public int bulbSwitch(int n) {  
        return new Double(Math.sqrt(n)).intValue();  
    }  
}
```

5

Problem 155

42 Trapping Rain Water

Given n non-negative integers representing an elevation map where the width of each bar is 1, compute how much water it is able to trap after raining.

For example,

Given $[0,1,0,2,1,0,1,3,2,1,2,1]$, return 6.

The above elevation map is represented by array

$[0,1,0,2,1,0,1,3,2,1,2,1]$. In this case,

6 units of rain water (blue section) are being trapped.

<https://leetcode.com/problems/trapping-rain-water/>



Solution

Use 2 pointers is the key to this problem. One from left, one from right, find the lower value for the candidate and find 'traps' whose depth is lower than current minimum(between left and right) and add the depth to amount of total water, until finding a new peak.

Listing 205: Solution

```
public class Solution {
    public int trap(int[] height) {
        int startIdx = 0 , endIdx = height.length - 1;

        if (height.length == 0){
            return 0;
        }
        int lHeight = height[startIdx];
        int rHeight = height[endIdx];
        int lowerHeight = 0;
        int totalWater = 0;

        while(startIdx < endIdx){
            lowerHeight = Math.min(lHeight , rHeight);
            if (lHeight == lowerHeight){
                lHeight = height[++startIdx];
            }
            if (rHeight == lowerHeight){
                rHeight = height[--endIdx];
            }
            totalWater += lowerHeight * (endIdx - startIdx);
        }
        return totalWater;
    }
}
```

```
20         while (startIdx < endIdx && lowerHeight >= lHeight){
           totalWater += lowerHeight - lHeight;
           startIdx ++;
           lHeight = height[startIdx];
        }
      }
      else if (rHeight == lowerHeight){
25         rHeight = height[--endIdx];
         while (startIdx < endIdx && lowerHeight >= rHeight){
           totalWater += lowerHeight - rHeight;
           endIdx --;
           rHeight = height[endIdx];
        }
30     }
  }

  return totalWater;
35 }
```

Problem 156

11 Container With Most Water

Given n non-negative integers a_1, a_2, \dots, a_n , where each represents a point at coordinate (i, a_i) . n vertical lines are drawn such that the two endpoints of line i is at (i, a_i) and $(i, 0)$. Find two lines, which together with x-axis forms a container, such that the container contains the most water.

<https://leetcode.com/problems/container-with-most-water/>

Solution

Somewhat like 2Sum, use 2 pointers to achieve this. The greedy algorithm needs proving.

Listing 206: Solution

```
public class Solution {  
    public int maxArea(int[] height) {  
        int max = -1;  
        int i = 0 , j = height.length - 1;  
5        while (i < j){  
            max = Math.max(max, (j - i) * Math.min(height[i] , height[j]));  
            if (height[i] < height[j])  
                i ++;  
            else  
10                j --;  
        }  
        return max;  
    }  
}
```

Problem 157

51 N-Queens

The n-queens puzzle is the problem of placing n queens on an n*n chessboard such that no two queens attack each other.

Given an integer n, return all distinct solutions to the n-queens puzzle.

Each solution contains a distinct board configuration of the n-queens' placement, where 'Q' and '.' both indicate a queen and an empty space respectively.

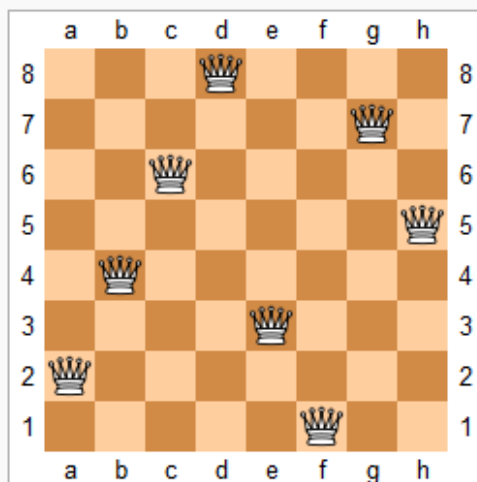
For example,

There exist two distinct solutions to the 4-queens puzzle:

```
[
  [".Q..", // Solution 1
   "...Q",
   "Q...",
   "..Q."],

  [ "..Q.", // Solution 2
    "Q...",
    "...Q",
    ".Q.."]
]
```

<https://leetcode.com/problems/n-queens/>



One solution to the eight queens puzzle

Solution

"N-queens" is a representative series of problem in backtracking problems.

However the train of thought is not that difficult, instead it's somewhat straightforward.

In the next problem 'N-queen-ii', another solution with $O(n)$ space will be displayed.

Listing 207: Solution

```
public class Solution {
    List<List<String>> puzzles = new ArrayList<List<String>>();

    public void solveNQueens(int n, int curRow, int [][] puzzle){
5         if (curRow == n){
            puzzles.add(generatePuzzle(puzzle));
            return ;
        }
        for (int col = 0 ; col < n ; col ++){
10             if (isValid(puzzle, curRow, col)){
                puzzle[curRow][col] = 1;
                solveNQueens(n, curRow + 1, puzzle);
                puzzle[curRow][col] = 0;
            }
15         }
    }

    public List<List<String>> solveNQueens(int n) {
        int [][] puzzle = new int[n][n];
20
        solveNQueens(n, 0, puzzle);

        return puzzles;
    }

25    public boolean isValid(int [][] puzzle, int row, int col){
        // check row
        for (int i = 0 ; i < col ; i++){
            if (puzzle[row][i] == 1){
30                 return false;
            }
        }

        // check column
35        for (int i = 0 ; i < row ; i++){
            if (puzzle[i][col] == 1){
                return false;
            }
        }

        // check top left
40        for (int i = 1 ; i <= Math.min(row, col) ; i++){
            if (puzzle[row - i][col - i] == 1){
                return false;
            }
        }

        // check top right
45        for (int i = 1 ; i <= Math.min(row, puzzle.length - col - 1) ; i++){
            if (puzzle[row - i][col + i] == 1){
50                 return false;
            }
        }
    }
}
```



```
        return true;
    }

55     public List<String> generatePuzzle(int [][] puzzle){
        List<String> result = new ArrayList<String>();
        for (int i = 0 ; i < puzzle.length ; i++){
            String newLine = "";
60             for (int j = 0 ; j < puzzle[0].length ; j++){
                if (puzzle[i][j] == 0){
                    newLine += ".";
                }
                else{
65                     newLine += "Q";
                }
            }
            result.add(newLine);
        }
70     return result;
    }
}
```

Problem 158

52 N-Queens II

Follow up for N-Queens problem.

Now, instead outputting board configurations,
return the total number of distinct solutions.

<https://leetcode.com/problems/n-queens-ii/>

Solution

This is not much harder than N-Queen, the ideas are the same. But this time a $O(n)$ space solution will be displayed. It doesn't use a two-dimensional array to keep the condition of the puzzle, instead use an $a[i]$ array where $a[i]$ representing the index of i^{th} row.

So how to judge whether the position is available? The array is of 1-dimension, so there won't be conflict in rows, checking whether columns conflict are also not difficult, what about diagonalization? Observation gives us the answer that when two point fit the attribute that $x_1 + y_1 = x_2 + y_2$ OR $x_1 - y_1 = x_2 - y_2$, eg (1,2) & (2,1) OR (1,2) & (2,3), they are on the same diagonal line.

Listing 208: Solution

```
public class Solution {  
    int count = 0;  
  
    public void solveNQueens(int total, int row, int[] cols){  
5        if (total == row){  
            count++;  
            return;  
        }  
        for (int col = 0 ; col < total; col ++){  
10            if (isValid(cols, row, col)){  
                cols[row] = col;  
                solveNQueens(total, row + 1, cols);  
                cols[row] = -1;  
            }  
15        }  
    }  
  
    public boolean isValid(int[] cols, int row, int col){  
20        for (int i = 0 ; i < row ; i++){  
            // check same col  
            if (cols[i] == col){  
                return false;  
            }  
25            // check top left & top right  
            if (row - col == i - cols[i] || row + col == i + cols[i]){  
                return false;  
            }  
        }  
30        return true;  
    }  
}
```

```
35     public int totalNQueens(int n) {  
        int [] cols = new int[n];  
        for (int i = 0 ; i < n ; i++){  
            cols[i] = -1;  
        }  
        solveNQueens(n, 0, cols);  
40     return count;  
    }  
}
```

Problem 159

217 Contains Duplicate

Given an array of integers, find if the array contains any duplicates.
Your function should return true if any value appears
at least twice in the array,
and it should return false if every element is distinct.

<https://leetcode.com/problems/contains-duplicate/>

Solution

A HashSet is enough for the problem series I and II.

Listing 209: Solution

```
public class Solution {  
    public boolean containsDuplicate(int[] nums) {  
        Set<Integer> uniqueElements = new HashSet<Integer>();  
        for (int num : nums){  
5           if (uniqueElements.contains(num)){  
                return true;  
            }  
            uniqueElements.add(num);  
        }  
10        return false;  
    }  
}
```

Problem 160

219 Contains Duplicate II

Given an array of integers and an integer k ,
find out whether there are two distinct indices i and j
in the array such that $\text{nums}[i] = \text{nums}[j]$
and the difference between i and j is at most k .

<https://leetcode.com/problems/contains-duplicate-ii/>

Solution

Either use a HashMap to record the latest index of element or use a k -length slide window is fine. Here I'm using a k -length slide window(made of HashSet) to solve the problem.

Listing 210: Solution

```
public class Solution {  
    public boolean containsNearbyDuplicate(int[] nums, int k) {  
        Set<Integer> kSlidingWindow = new HashSet<Integer>();  
        for (int idx = 0 ; idx < nums.length ; idx++){  
            if (kSlidingWindow.contains(nums[idx])){  
                return true;  
            }  
            kSlidingWindow.add(nums[idx]);  
            if (kSlidingWindow.size() > k){  
                kSlidingWindow.remove(nums[idx - k]);  
            }  
        }  
        return false;  
    }  
}
```

Problem 161

220 Contains Duplicate III

Given an array of integers,
find out whether there are two distinct indices
i and j in the array such that the difference
between nums[i] and nums[j] is at most t and
the difference between i and j is at most k.

<https://leetcode.com/problems/contains-duplicate-iii/>

Solution

Use a existence data structure TreeSet, which has both the attribute of Set and Binary Search Tree, a perfect match of this problem.

Listing 211: Solution

```
public class Solution {  
    public boolean containsNearbyAlmostDuplicate(int[] nums, int k, int t) {  
        TreeSet<Integer> almostUniqueElements = new TreeSet<Integer>();  
  
        for (int idx = 0 ; idx < nums.length ; idx ++){  
            if ((almostUniqueElements.floor(nums[idx]) != null &&  
                (nums[idx] <= t + almostUniqueElements.floor(nums[idx]))) ||  
                (almostUniqueElements.ceiling(nums[idx]) != null &&  
                (almostUniqueElements.ceiling(nums[idx]) <= t + nums[idx] ))  
            ){  
                return true;  
            }  
  
            almostUniqueElements.add(nums[idx]);  
            if (almostUniqueElements.size() > k){  
                almostUniqueElements.remove(nums[idx - k]);  
            }  
        }  
  
        return false;  
    }  
}
```

Problem 162

75 Sort Colors

Given an array with n objects colored red, white or blue, sort them so that objects of the same color are adjacent, with the colors in the order red, white and blue.

Here, we will use the integers 0, 1, and 2 to represent the color red, white, and blue respectively.

Note:

You are not suppose to use the library's sort function for this problem.

Follow up:

A rather straight forward solution is a two-pass algorithm using counting sort.

First, iterate the array counting number of 0's, 1's, and 2's, then overwrite array with total number of 0's, then 1's and followed by 2's.

Could you come up with an one-pass algorithm using only constant space?

<https://leetcode.com/problems/sort-colors/>

Solution

Use two pointers to mark where the index of 0 and 2 are(In fact, there are 3 pointers, the main idx pointer marks where 1s are.)

In LintCode, there is a Sort Colors II, for k colors. The key insight is like bucket sorting.

Listing 212: Solution

```
public class Solution {
    public void sortColors(int[] nums) {
        int idxZero = 0 , idxTwo = nums.length - 1;
        int idx = 0;
        while (idx <= idxTwo){
            if (nums[idx] == 0){
                swap(nums, idx, idxZero);
                idx ++;
                idxZero ++;
            }
            else if (nums[idx] == 2){
                swap(nums, idx, idxTwo);
                // idx ++;
                idxTwo --;
            }
            else{
                idx ++;
            }
        }
    }
}
```

```
25      public void swap(int [] nums, int idx1, int idx2){  
          int swap = nums[idx1];  
          nums[idx1] = nums[idx2];  
          nums[idx2] = swap;  
      }  
  }
```


Problem 163

39 Combination Sum

Given a set of candidate numbers (C) and a target number (T),
find all unique combinations in C where the candidate numbers sums to T.

The same repeated number may be chosen from C unlimited number of times.

Note:

All numbers (including target) will be positive integers.

Elements in a combination (a1, a2, ... , ak) must be in non-descending order.

(ie, a1 <= a2 <= ... <= ak).

The solution set must not contain duplicate combinations.

For example, given candidate set 2,3,6,7 and target 7,

A solution set is:

[7]

[2, 2, 3]

<https://leetcode.com/problems/combination-sum/>

Solution

A dfs problem. Remember to sort the array first since *Elements in a combination (a1, a2, ..., ak) must be in non-descending order.*

Listing 213: Solution

```
public class Solution {  
    List<List<Integer>> result = new ArrayList<List<Integer>>();  
    public List<List<Integer>> combinationSum(int[] candidates, int target) {  
        Arrays.sort(candidates);  
5         List<Integer> combination = new ArrayList<Integer>();  
        findSum(candidates, target, 0, combination);  
        return result;  
    }  
    public void findSum(int[] candidates, int target, int curIdx,  
10    List<Integer> curCombination){  
        int curSum = 0;  
        for (int com : curCombination)  
            curSum += com;  
        if (curSum > target)  
15            return;  
        if (curSum == target){  
            result.add(new ArrayList<Integer>(curCombination));  
            return;  
        }  
20        for (int idx = curIdx ; idx < candidates.length; idx ++){  
            curCombination.add(candidates[idx]);  
            findSum(candidates, target, idx, curCombination);  
            curCombination.remove(curCombination.size() - 1);  
        }  
25    }  
}
```

Problem 164

40 Combination Sum II

Given a collection of candidate numbers (C) and a target number (T), find all unique combinations in C where the candidate numbers sums to T.

Each number in C may only be used once in the combination.

Note:

All numbers (including target) will be positive integers.

Elements in a combination (a1, a2, ... , ak)

must be in non-descending order. (ie, a1 <= a2 <= ... <= ak).

The solution set must not contain duplicate combinations.

For example, given candidate set 10,1,2,7,6,1,5 and target 8,

A solution set is:

```
[1, 7]
[1, 2, 5]
[2, 6]
[1, 1, 6]
```

<https://leetcode.com/problems/combination-sum-ii/>

Solution

More difficult than “Combination Sum”, mainly on how to deal with duplicates. Except for set the startIdx to lastIdx + 1 to avoid using the same elements, we should also use a while loop to avoid using elements with the same value after we have already combined them.

Listing 214: Solution

```
public class Solution {
    List<List<Integer>> combinations = new ArrayList<List<Integer>>();

    public List<List<Integer>> combinationSum2(int[] candidates, int target) {
5        Arrays.sort(candidates);
        List<Integer> curCombination = new ArrayList<Integer>();
        // curIdx is initially -1 since the array should start iterating from 0
        // and -1 + 1 = 0
        findSum(candidates, target, -1, curCombination);
10        return combinations;
    }

    private void findSum(int[] candidates, int target, int curIdx,
15        List<Integer> curCombination){
        int curSum = 0;
        for (int element : curCombination){
            curSum += element;
        }
        if (target < curSum){
20            return;
        }
        if (target == curSum){
            combinations.add(new ArrayList<Integer>(curCombination));
            return;
        }
    }
}
```

```
25     }

    for (int idx = curIdx + 1 ; idx < candidates.length ; idx++){
        curCombination.add(candidates[idx]);
        findSum(candidates, target, idx, curCombination);
30     curCombination.remove(curCombination.size() - 1);
        // To avoid duplicates, think of it.
        // For example, 11167,
        // after 1->11->111->1116->11167->1117-> next what?
        // (116..., if we don't add such while loop,
35     //         it will continue to add 111...)
        while (idx < candidates.length - 1
                && candidates[idx] == candidates[idx + 1]){
            idx++;
        }
40     }
}
}
```

Problem 165

216 Combination Sum III

Find all possible combinations of k numbers that add up to a number n, given that only numbers from 1 to 9 can be used and each combination should be a unique set of numbers.

Ensure that numbers within the set are sorted in ascending order.

Example 1:

Input: k = 3, n = 7

Output:

[[1,2,4]]

Example 2:

Input: k = 3, n = 9

Output:

[[1,2,6], [1,3,5], [2,3,4]]

<https://leetcode.com/problems/combination-sum-iii/>

Solution

Don't know if it is my illusion, I think this Problem is a bit easier than "Combination Sum II" since it doesn't need to be processed to avoid duplicates.

Listing 215: Solution

```
public class Solution {
    List<List<Integer>> result = new ArrayList<List<Integer>>();
    public List<List<Integer>> combinationSum3(int k, int n) {
        List<Integer> curCombination = new ArrayList<Integer>(k);
5         for (int i = 0 ; i < k ; i++){
            curCombination.add(0);
        }
        findSum(k , 0 , 0 , n, curCombination);

10        return result;
    }
    private void findSum(int maxDepth, int curDepth,
        int curNum, int target, List<Integer> curCombination){
        int curSum = 0;
15        for (int element : curCombination){
            curSum += element;

            if (curSum > target){
```

```
20         return;
    }
    if (curDepth == maxDepth){
        if (curSum == target){
            result.add(new ArrayList(curCombination));
25         }
        return;
    }

    for (int num = curNum + 1 ; num <= 9 ; num++){
30         curCombination.set(curDepth, num);
        findSum(maxDepth, curDepth + 1, num, target, curCombination);
        curCombination.set(curDepth, 0);
    }
35     }
}
```

Problem 166

48 Rotate Image

You are given an $n * n$ 2D matrix representing an image.

Rotate the image by 90 degrees (clockwise).

Follow up:

Could you do this in-place?

<https://leetcode.com/problems/rotate-image/>

Solution

Finding the rule of swapping is the key to solving the problem. (Honestly, my solution is not exactly in-place since I used a boolean array to check if the number in grid has been changed. I'll list the proper code downward my solution.

However, the world has taught us not to under-estimate every problem, there is also another interesting way of solving the problem: when we regard the matrix as a paper, we fold the paper up and down, then fold it along the diagonal, we will realize this is exactly the result we want!

Listing 216: Solution

```
public class Solution {
    public void rotate(int[][] matrix) {
        int n = matrix.length;
        if (n == 0){
            return;
        }
        boolean[][] swapped = new boolean[n][n];
        for (int i = 0 ; i < n ; i++){
            for (int j = 0 ; j < n ; j++){
                if (swapped[i][j]){
                    continue;
                }
                int startX = i , startY = j;
                int swapX = startY , swapY = n - 1 - startX;
                swapped[i][j] = true;
                int swap = matrix[startX][startY];

                while (startX != swapX || startY != swapY){
                    int tmp = matrix[swapX][swapY];
                    matrix[swapX][swapY] = swap;
                    swap = tmp;
                    swapped[swapX][swapY] = true;
                    // swap swapX & swapY
                    swapX = swapX + swapY;
                    swapY = swapX - swapY;
                    swapX = swapX - swapY;
                    swapY = n - 1 - swapY;
                }
                matrix[startX][startY] = swap;
            }
        }
    }
}
```

```
    }  
    }  
}  
  
35 // Solution WITHOUT boolean array  
public class Solution {  
    public void rotate(int [][] matrix) {  
        int n = matrix.length;  
        int limit = (n-1)/2;  
40        for(int i=0;i<= limit; i++){  
            for(int j=i;j<n-1-i;j++){  
                int temp = matrix[i][j];  
                matrix[i][j] = matrix[n-1-j][i];  
45                matrix[n-1-j][i] = matrix[n-1-i][n-1-j];  
                matrix[n-1-i][n-1-j] = matrix[j][n-1-i];  
                matrix[j][n-1-i] = temp;  
            }  
        }  
50    }  
}
```

Problem 167

53 Maximum Subarray

Find the contiguous subarray within an array
(containing at least one number) which has the largest sum.

For example, given the array `[-2,1,-3,4,-1,2,1,-5,4]`,
the contiguous subarray `[4,-1,2,1]` has the largest sum = 6.

More practice:

If you have figured out the $O(n)$ solution,
try coding another solution using the divide and conquer approach,
which is more subtle.

<https://leetcode.com/problems/maximum-subarray/>

Solution

This is a conventional DP problem, it's worth us carefully thinking about it.

What's the first solution comes into mind at first sight? Enumeration, of course, the Brute force method. However, when we talk about its complexity, we may need $O(n)$ to iterate both i and j , along with an extra $O(n)$ effort to sum up $a[i...j]$, $O(N^3)$ in total.

Listing 217: Guide 1

```
for(int i = 1; i <= n; i++){
    for(int j = i; j <= n; j++){
        int sum = 0;
        for(int k = i; k <= j; k++)
5         sum += a[k];
        max = Max(max, sum);
    }
}
```

How to optimize? Can we calculate the sum during the process of enumeration? The sum is actually `formerSum + a[idx-1]`.

Listing 218: Guide 2

```
for(int i = 1; i <= n; i++){
    int sum = 0;
    for(int j = i; j <= n; j++){
        sum += a[j];
5         max = Max(max, sum);
    }
}
```

Divide and Conquer can even promote our solution further more. We may find a middle point, and the maximum subarray sum lies either on left, right, or go through the middle element(since we are finding consecutive sequence). Then the steps of algorithm is to separate subarray according to a middle element and calculate optimal sum `sum1` and `sum2` start from middle element(With the time complexity $2 * T(n/2) + O(n) = O(n \log n)$);

Finally we can use DP to solve it. Let $dp[i]$ be the maximum we can get until idx i, so $dp[i] = \max(dp[i - 1] + \text{nums}[i], \text{nums}[i])$.

Listing 219: Solution

```
public class Solution {
    public int maxSubArray(int[] nums) {
        int[] maxDp = new int[nums.length];
        if (nums.length == 0){
5           return 0;
        }
        int totalMax = nums[0];
        maxDp[0] = nums[0];

10        for (int idx = 1 ; idx < nums.length ; idx++){
            maxDp[idx] = Math.max(maxDp[idx - 1] + nums[idx], nums[idx]);
            totalMax = Math.max(maxDp[idx], totalMax);
        }

15        return totalMax;
    }
}
```

Solution Follow-up

The follow up need us to do it in a divide and conquer way which is mentioned above, as listing, the code is copied from Internet, respectively calculating sums $[\text{startIdx}, \text{currentIdx}]$, $[\text{currentIdx} + 1, \text{endIdx}]$ and $[\text{startIdx}, \text{endIdx}]$, use a binary search method to decide the value of currentIdx (so the runtime complexity is $O(n \log n)$).

Listing 220: Solution Follow-up

```
public class Solution {
    public int maxSum(int[] A, int left, int right)
    {
        if( left == right ){
5           return A[left];
        }
        int center = (left + right) / 2;
        int maxLeftSum = maxSum( A, left, center);
        int maxRightSum = maxSum( A, center+1, right);

10        int maxLeft = Integer.MIN_VALUE, tempLeft = 0;
        int maxRight = Integer.MIN_VALUE, tempRight = 0;
        for (int i=center; i>=left; --i){
            tempLeft += A[i];
            if (tempLeft > maxLeft){
15                maxLeft = tempLeft;
            }
        }
        for (int i=center+1; i<=right; ++i){
20            tempRight += A[i];
            if (tempRight > maxRight){
                maxRight = tempRight;
            }
        }
        return Math.max(maxLeftSum, maxRightSum);
    }
}
```

```
    }  
25    }  
  
    int maxCenterSum = maxLeft + maxRight;  
  
    return maxCenterSum > maxLeftSum ?  
30    (maxCenterSum > maxRightSum ?  
        maxCenterSum : maxRightSum) : maxLeftSum > maxRightSum ?  
        maxLeftSum : maxRightSum;  
    }  
  
35    public int maxSubArray(int [] A){  
        int len = A.length;  
        return maxSum(A,0 , len -1);  
    }  
}
```

Problem 168

152 Maximum Product Subarray

Find the contiguous subarray within an array
(containing at least one number) which has the largest product.

For example, given the array `[2,3,-2,4]`,
the contiguous subarray `[2,3]` has the largest product = 6.

<https://leetcode.com/problems/maximum-product-subarray/>

Solution

The solution of this problem is a bit like “Maximum Subarray”, except for the fact that we use an extra array to store the minimum product until `idx` since minimum value may also produce maximum value since negatives multiply negatives equal positives.

Listing 221: Solution

```
public class Solution {
    public int maxProduct(int[] nums) {
        int[] curMin = new int[nums.length];
        int[] curMax = new int[nums.length];
5         if (nums.length == 0){
            return 0;
        }
        int totalMax = nums[0];
        curMax[0] = curMin[0] = nums[0];
10
        for (int idx = 1 ; idx < nums.length ; idx++){
            curMax[idx] = Math.max(
                Math.max(curMax[idx - 1] * nums[idx],
                    curMin[idx - 1] * nums[idx]),
15                 nums[idx]);
            curMin[idx] = Math.min(
                Math.min(curMax[idx - 1] * nums[idx],
                    curMin[idx - 1] * nums[idx]),
                nums[idx]);
20             totalMax = Math.max(curMax[idx], totalMax);
        }

        return (int)totalMax;
    }
25 }
```

Problem 169

238 Product of Array Except Self

Given an array of n integers where $n > 1$, `nums`,
return an array `output` such that `output[i]`
is equal to the product of all the elements of `nums` except `nums[i]`.

Solve it without division and in $O(n)$.

For example, given `[1,2,3,4]`, return `[24,12,8,6]`.

Follow up:

Could you solve it with constant space complexity?

(Note: The output array does not count as extra space
for the purpose of space complexity analysis.)

<https://leetcode.com/problems/product-of-array-except-self/>

Solution

Solution using $O(n)$ space is trivial, so I just put the $O(1)$ space (except for the result array) solution here. It's like multiplying a matrix, from start of array to the end, `result[i] = result[i - 1] * a[i - 1]` (avoiding multiplying `a[i]`), then backward, `result[i] = result[i] * a[i + 1]`.

Listing 222: Solution

```
public class Solution {
    public int[] productExceptSelf(int[] nums) {
        if (nums.length == 0){
            return new int[0];
        }
        int[] products = new int[nums.length];
        products[0] = 1;
        for (int idx = 1 ; idx < nums.length ; idx++){
            products[idx] = products[idx - 1] * nums[idx - 1];
        }

        int product = 1;
        for (int idx = nums.length - 2 ; idx >= 0 ; idx --){
            product *= nums[idx + 1];
            products[idx] *= product;
        }

        return products;
    }
}
```

Problem 170

133 Clone Graph

Clone an undirected graph.

Each node in the graph contains a label and a list of its neighbors.

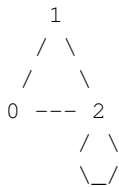
OJ's undirected graph serialization:

Nodes are labeled uniquely.

We use # as a separator for each node, and ,
as a separator for node label and each neighbor of the node.
As an example, consider the serialized graph {0,1,2#1,2#2,2}.

The graph has a total of three nodes,
and therefore contains three parts as separated by #.

First node is labeled as 0. Connect node 0 to both nodes 1 and 2.
Second node is labeled as 1. Connect node 1 to node 2.
Third node is labeled as 2. Connect node 2 to node 2 (itself),
thus forming a self-cycle.
Visually, the graph looks like the following:



<https://leetcode.com/problems/clone-graph/>

Solution

Using a HashMap to map original node and cloned node(just like what we did in 'Copy List With Random Pointer'), then use BFS or DFS to clone the whole graph. For practise, I list DFS solution here.

Listing 223: Solution

```

/**
 * Definition for undirected graph.
 * class UndirectedGraphNode {
 *     int label;
 *     List<UndirectedGraphNode> neighbors;
 *     UndirectedGraphNode(int x) { label = x; neighbors
 *         = new ArrayList<UndirectedGraphNode>(); }
 * };
 */
10 public class Solution {
    HashMap<UndirectedGraphNode, UndirectedGraphNode> map
        = new HashMap<UndirectedGraphNode, UndirectedGraphNode>();
    public UndirectedGraphNode cloneGraph(UndirectedGraphNode node) {
        if (node == null){
15         return null;
        }
    }
}

```

```
    }

    UndirectedGraphNode newNode = new UndirectedGraphNode(node.label);
    if (!map.containsKey(node)){
20         map.put(node, newNode);
    }
    else{
        return map.get(node);
    }
25     for (UndirectedGraphNode neighborNode : node.neighbors){
        if (map.containsKey(neighborNode)){
            map.get(node).neighbors.add(map.get(neighborNode));
        }
        else{
30             map.get(node).neighbors.add(cloneGraph(neighborNode));
        }
    }

    return newNode;
35 }
}
```

Problem 171

153 Find Minimum in Rotated Sorted Array

Suppose a sorted array is rotated at some pivot unknown to you beforehand.

(i.e., 0 1 2 4 5 6 7 might become 4 5 6 7 0 1 2).

Find the minimum element.

You may assume no duplicate exists in the array.

<https://leetcode.com/problems/find-minimum-in-rotated-sorted-array/>

Solution

A binary search problem. Of course we can do it in $O(n)$ runtime by iterating the list linearly, but binary search is of course better for its $O(\log n)$ runtime.

Listing 224: Solution

```
public class Solution {  
    public int findMin(int[] nums) {  
        if (nums.length == 0){  
            return -1;  
        }  
        int start = 0, end = nums.length - 1;  
  
        while (start < end){  
            int mid = start + (end - start) / 2;  
            if (nums[start] == nums[mid] && nums[start] > nums[start + 1]){  
                return nums[start + 1];  
            }  
            else{  
                if (nums[start] < nums[mid]){  
                    start = mid;  
                }  
                else{  
                    end = mid;  
                }  
            }  
        }  
  
        return nums[0];  
    }  
}
```

Problem 172

154 Find Minimum in Rotated Sorted Array II

Follow up for "Find Minimum in Rotated Sorted Array":
What if duplicates are allowed?

Would this affect the run-time complexity? How and why?
Suppose a sorted array is rotated at some pivot unknown to you beforehand.

(i.e., 0 1 2 4 5 6 7 might become 4 5 6 7 0 1 2).

Find the minimum element.

The array may contain duplicates.

<https://leetcode.com/problems/find-minimum-in-rotated-sorted-array-ii/>

Solution

This one is following previous problem "Find Minimum in Rotated Sorted Array", and though there are only several lines of code that are different, the problem is much harder than previous one. Take care of the condition when there are plenty of duplicates, for example, when the list is 10,1,10,10,10, this is when $a[\text{low}] < a[\text{high}]$, which is not we want when using binary search(Think of why?).

Listing 225: Solution

```
public class Solution {
    public int findMin(int[] nums) {
        if (nums.length == 0)
            return 0;
5       int start = 0 , end = nums.length - 1;
        while (start < end){
            // this is weird since we should always consider the condition
            // that a[l] > a[r]
            if (nums[start] < nums[end])
10              return nums[start];
            int mid = start + (end - start) / 2;
            if (mid == start){
                if (nums[start] > nums[start + 1])
                    return nums[start + 1];
15            }
            if (nums[start] > nums[mid])
                end = mid;
            else if (nums[end] < nums[mid])
                start = mid;
20            else
                start ++;
        }
        return nums[0];
    }
25 }
```


Problem 173

33 Search in Rotated Sorted Array

Suppose a sorted array is rotated at some pivot unknown to you beforehand.

(i.e., 0 1 2 4 5 6 7 might become 4 5 6 7 0 1 2).

You are given a target value to search.

If found in the array return its index, otherwise return -1.

You may assume no duplicate exists in the array.

<https://leetcode.com/problems/search-in-rotated-sorted-array/>

Solution

A complicated binary search problem, need to check if the target is in monotone interval(always increasing) or on the other side. Don't forget to check if the element matches start or end pointer of list.

Listing 226: Solution

```
public class Solution {
    public int search(int[] nums, int target) {
        int start = 0, end = nums.length - 1;
        while (start <= end){
5           int mid = start + (end - start) / 2;
            if (nums[mid] == target){
                return mid;
            }
            if (nums[start] < nums[end]){
10              if (target > nums[mid]){
                  start = mid + 1;
              }
              else{
                  end = mid - 1;
15              }
            }
            else{
                if (nums[start] < nums[mid]){
                    if (target == nums[start]){
20                       return start;
                    }
                    if (target > nums[start] && target < nums[mid]){
                        end = mid - 1;
                    }
                    else{
25                       start = mid;
                    }
                }
            }
            else{
30              if (target == nums[end]){
                  return end;
              }
              if (target > nums[mid] && target < nums[end]){
```

```
35         start = mid + 1;
36     }
37     else{
38         end = mid - 1;
39     }
40 }
41 }
42 return -1;
43 }
44 }
```

Problem 174

81 Search in Rotated Sorted Array II

Follow up for "Search in Rotated Sorted Array":

What if duplicates are allowed?

Would this affect the run-time complexity? How and why?

Write a function to determine if a given target is in the array.

<https://leetcode.com/problems/search-in-rotated-sorted-array-ii/>

Solution

With the experience of last problem "Search in Rotated Sorted Array", solving this problem is not that difficult. Still, when we find `nums[start] == nums[mid] == nums[end]`, we just increase start by 1 since we are not sure where our target is ([10,1,10,10,10] & [10,10,10,1,10]).

What's more, in this solution I slightly simplify my code in previous problem.

Listing 227: Solution

```
public class Solution {
    public boolean search(int[] nums, int target) {
        int start = 0, end = nums.length - 1;
        while (start <= end) {
            int mid = start + (end - start) / 2;
            if (target == nums[mid]) {
                return true;
            }
            if (nums[start] < nums[end]) {
                if (target > nums[mid]) {
                    start = mid + 1;
                }
                else {
                    end = mid - 1;
                }
            }
            else {
                if (nums[start] > nums[mid]) {
                    if (target > nums[mid] && target <= nums[end]) {
                        start = mid + 1;
                    }
                    else {
                        end = mid - 1;
                    }
                }
                else if (nums[start] < nums[mid]) {
                    if (target >= nums[start] && target < nums[mid]) {
                        end = mid - 1;
                    }
                    else {
                        start = mid + 1;
                    }
                }
            }
        }
    }
}
```

```
    }  
    else{  
        start ++;  
    }  
    }  
}  
  
return false;  
}  
}
```

Problem 175

76 Minimum Window Substring

Given a string S and a string T,
find the minimum window in S which
will contain all the characters in T in complexity $O(n)$.

For example,
S = "ADOBECODEBANC"
T = "ABC"
Minimum window is "BANC".

Note:
If there is no such window in S
that covers all characters in T, return the empty string "".

If there are multiple such windows,
you are guaranteed that there will always be
only one unique minimum window in S.

<https://leetcode.com/problems/minimum-window-substring/>

Solution

A troublesome problem. Use two pointers to maintain current window, containing all characters we need in t string. And under the guaranteeing circumstances that our window contains all digits we want, we may shrink the window by moving the left pointer to right; when the window is no longer adequate, we move the right pointer forward to meet the requirement.

Listing 228: Solution

```
public class Solution {  
    public String minWindow(String s, String t) {  
        if (s.length() == 0 || t.length() == 0){  
            return "";  
        }  
        int[] require = new int[256];  
        boolean[] needLetter = new boolean[256];  
  
        for (int i = 0 ; i < t.length() ; i++){  
            needLetter[t.charAt(i)] = true;  
            require[t.charAt(i)] ++;  
        }  
        int count = t.length();  
        int minLength = Integer.MAX_VALUE;  
        String minStr = "";  
  
        int i = 0;  
        int j = -1;  
  
        while (i < s.length() && j < s.length()){  
            if (count > 0){  
                j ++;
```

```

    if (j >= s.length()){
        break;
25    }
    require[s.charAt(j)] --;
    if (needLetter[s.charAt(j)] && require[s.charAt(j)] >= 0){
        count --;
    }
30    }else{
        int curLength = j - i + 1;
        if (curLength < minLength){
            minLength = curLength;
            minStr = s.substring(i, j + 1);
35        }

        require[s.charAt(i)] ++;
        if (needLetter[s.charAt(i)] && require[s.charAt(i)] > 0){
            count ++;
40        }
        i ++;
    }
    }
45    return minStr;
}
}
```

Problem 176

78 Subsets

Given a set of distinct integers, `nums`,
return all possible subsets.

Note:

Elements in a subset must be in non-descending order.

The solution set must not contain duplicate subsets.

For example,

If `nums = [1,2,3]`, a solution is:

```
[
  [3],
  [1],
  [2],
  [1,2,3],
  [1,3],
  [2,3],
  [1,2],
  []
]
```

<https://leetcode.com/problems/subsets/>

Solution 1

Take `[1,2,3]` for example.

Initially, the result is empty, namely `[]`.

Then, we add 1 to result, combine the result with former ones, that is `[],[1]`

Next, we add 2 to every list in result, also combine the new and old list together, `[],[1],[2],[1,2]`

Finally, it's 3's turn, we get `[],[1],[2],[1,2],[3],[1,3],[2,3],[1,2,3]` as our final result.

Listing 229: Solution 1

```
public class Solution {
    public List<List<Integer>> subsets(int[] nums) {
        Arrays.sort(nums);
        List<List<Integer>> result = new ArrayList<List<Integer>>();
5         result.add(new ArrayList<Integer>());

        for (int num : nums){
            int resultLength = result.size();
            for (int idx = 0 ; idx < resultLength ; idx++){
10                 List<Integer> newList = new ArrayList(result.get(idx));
                    newList.add(num);
                    // Collections.sort(newList);
                    result.add(newList);
            }
15         }

        return result;
    }
}
```

```
}

```

Solution 2

The second solution is a DFS method. Interesting since it utilized 2 same function call to simulate picking or not picking element at curIdx.

Listing 230: Solution 2

```
public class Solution {
    List<List<Integer>> result = new ArrayList<List<Integer>>();

    public void findSubsets(List<Integer> curList, int[] nums, int curIdx){
5         if (curIdx == nums.length){
            result.add(new ArrayList<Integer>(curList));
            return;
        }
        curList.add(nums[curIdx]);
10        findSubsets(curList, nums, curIdx + 1);
        curList.remove(curList.size() - 1);
        findSubsets(curList, nums, curIdx + 1);
    }
    public List<List<Integer>> subsets(int[] nums) {
15        Arrays.sort(nums);
        findSubsets(new ArrayList<Integer>(), nums, 0);

        return result;
    }
20 }
```


Problem 177

90 Subsets II

Given a collection of integers that might contain duplicates, `nums`, return all possible subsets.

Note:

Elements in a subset must be in non-descending order.

The solution set must not contain duplicate subsets.

For example,

If `nums = [1,2,2]`, a solution is:

```
[
  [2],
  [1],
  [1,2,2],
  [2,2],
  [1,2],
  []
]
```

<https://leetcode.com/problems/subsets-ii/>

Solution 1

The train of thought is similar to Solution 1 in Problem “Subsets I” (Subsets II can also be solved using Solution 2). The key is on how to deal with the duplicates. Mention that when current element equals to last element, eg, 2 & 2 in [1,2,2], the new list generated are only from what has just been generated in last round.

Listing 231: Solution 1

```
public class Solution {
    List<List<Integer>> result = new ArrayList<List<Integer>>();
    public List<List<Integer>> subsetsWithDup(int[] nums) {
        if (nums.length == 0){
            return result;
        }
        Arrays.sort(nums);
        List<List<Integer>> generatedList = new ArrayList<List<Integer>>();
        result.add(new ArrayList<Integer>());
        for (int eleIdx = 0 ; eleIdx < nums.length ; eleIdx ++){
            int num = nums[eleIdx];
            if (eleIdx > 0 && nums[eleIdx] == nums[eleIdx - 1]){
                List<List<Integer>> newGeneratedList = new ArrayList<List<Integer>>();
                for (int formerIdx = 0 ; formerIdx < generatedList.size() ; formerIdx ++){
                    List<Integer> formerElement = new ArrayList(generatedList.get(formerIdx));
                    formerElement.add(num);
                    newGeneratedList.add(formerElement);
                }
                result.addAll(newGeneratedList);
                generatedList = newGeneratedList;
            }
        }
    }
}
```

```

        else{
            generatedList.clear();
            for (int formerIdx = 0 ; formerIdx < result.size() ; formerIdx ++){
25         List<Integer> formerElement = new ArrayList(result.get(formerIdx));
                formerElement.add(num);
                generatedList.add(formerElement);
            }
            result.addAll(generatedList);
30     }
    }

    return result;
35 }
}

```

Solution 2

The backtracking method. Dealing with duplicates are also the key insight. The train of thought is somewhat like solution in “Combination Sum II”.

Listing 232: Solution 2

```

public class Solution {
    List<List<Integer>> result = new ArrayList<List<Integer>>();
    public void findSubsets(List<Integer> curList, int[] nums, int curIdx){
        if (curIdx == nums.length){
5         result.add(new ArrayList(curList));
            return;
        }
        curList.add(nums[curIdx]);
        findSubsets(curList, nums, curIdx + 1);
10     curList.remove(curList.size() - 1);
        int diff = 1;
        while (curIdx + diff < nums.length
            && nums[curIdx + diff - 1] == nums[curIdx + diff]){
            diff++;
15     }
        findSubsets(curList, nums, curIdx + diff);
    }
    public List<List<Integer>> subsetsWithDup(int[] nums) {
        Arrays.sort(nums);
20     findSubsets(new ArrayList<Integer>(), nums, 0);
        return result;
    }
}

```

Problem 178

303 Range Sum Query - Immutable

Given an integer array `nums`,
find the sum of the elements between indices `i` and `j` ($i \leq j$), inclusive.

Example:

Given `nums = [-2, 0, 3, -5, 2, -1]`

`sumRange(0, 2) -> 1`

`sumRange(2, 5) -> -1`

`sumRange(0, 5) -> -3`

Note:

You may assume that the array does not change.

There are many calls to `sumRange` function.

<https://leetcode.com/problems/range-sum-query-immutable/>

Solution

Not a bad problem. Since the `sumRange` function may be called often, adding them every time we call the function sounds a bad idea. Use a sum array to store the sums for convenience. Mention the boundaries.

Listing 233: Solution

```
public class NumArray {
    private int[] sums;
    public NumArray(int[] nums) {
        sums = new int[nums.length];
5       int sum = 0;
        for (int idx = 0 ; idx < nums.length ; idx++){
            sum += nums[idx];
            sums[idx] = sum;
        }
10    }

    public int sumRange(int i, int j) {
        if (i == 0){
            return sums[j];
15        }
        else{
            return sums[j] - sums[i - 1];
        }
20    }
}

// Your NumArray object will be instantiated and called as such:
25 // NumArray numArray = new NumArray(nums);
// numArray.sumRange(0, 1);
// numArray.sumRange(1, 2);
```

Problem 179

307 Range Sum Query - Mutable

Given an integer array `nums`, find the sum of the elements between indices `i` and `j` ($i \leq j$), inclusive.

The `update(i, val)` function modifies `nums` by updating the element at index `i` to `val`.

Example:

Given `nums = [1, 3, 5]`

`sumRange(0, 2) -> 9`

`update(1, 2)`

`sumRange(0, 2) -> 8`

Note:

The array is only modifiable by the `update` function.

You may assume the number of calls to `update` and `sumRange` function is distributed evenly.

<https://leetcode.com/problems/range-sum-query-mutable/>

Solution

Using solution in former problem “Range Sum Query - Immutable” will get TLE, so try some other solutions.

Make use of a freaking awesome data structure of niubility! Binary Index Tree(or Fenwick Tree)!

For more reference on index trees, see

<http://www.doc88.com/p-9475113829694.html>

<http://blog.csdn.net/int64ago/article/details/7429868>

<https://www.topcoder.com/community/data-science/data-science-tutorials/binary-indexed-tree/>

Something need to be considered or mistakes I have made:

`nums[0]` is not in BIT since `lowbit(0) == 0` will cause a infinite loop, so `nums[0]` is resolved seperately.

I forgot to update `nums` array but just update BIT, which brings me confusion on what is wrong with my solution.

Be aware of the boundaries.

Listing 234: Solution

```
public class NumArray {
    private int[] biTree;
    private int[] nums;

    5     private int lowbit(int x){
        return x & (-x);
    }
    public NumArray(int[] nums) {
        10         if (nums.length == 0){
            return;
        }
        this.nums = nums;
        biTree = new int[nums.length];
        biTree[0] = nums[0];
        15         for (int idx = 1 ; idx < nums.length ; idx ++){
```

```
        for (int j = idx - lowbit(idx) + 1; j <= idx ; j++){
            biTree[idx] += nums[j];
        }
    }
20 }
    void update(int i, int val) {
        int diff = val - nums[i];
        nums[i] = val;
        while (i < nums.length){
25             biTree[i] += diff;
            i += lowbit(i);
            if (i == 0){
                break;
            }
30         }
    }

    private int getSum(int idx){
        int sum = 0;
35         while (idx > 0){
            sum += biTree[idx];
            idx -= lowbit(idx);
        }
        sum += biTree[0];
40
        return sum;
    }

    public int sumRange(int i, int j) {
        if (i == 0){
45             return getSum(j);
        }
        else{
            return getSum(j) - getSum(i - 1);
        }
50     }
}

// Your NumArray object will be instantiated and called as such:
55 // NumArray numArray = new NumArray(nums);
// numArray.sumRange(0, 1);
// numArray.update(1, 10);
// numArray.sumRange(1, 2);
```

Problem 180

304 Range Sum Query 2D - Immutable

Given a 2D matrix 'matrix', find the sum of the elements inside the rectangle defined by its upper left corner (row1, col1) and lower right corner (row2, col2).

Range Sum Query 2D

The above rectangle (with the red border) is defined by

(row1, col1) = (2, 1) and (row2, col2) = (4, 3), which contains sum = 8.

Example:

```
Given matrix = [
  [3, 0, 1, 4, 2],
  [5, 6, 3, 2, 1],
  [1, 2, 0, 1, 5],
  [4, 1, 0, 1, 7],
  [1, 0, 3, 0, 5]
]
```

```
sumRegion(2, 1, 4, 3) -> 8
```

```
sumRegion(1, 1, 2, 2) -> 11
```

```
sumRegion(1, 2, 2, 4) -> 12
```

Note:

You may assume that the matrix does not change.

There are many calls to sumRegion function.

You may assume that row1 row2 and col1 col2.

<https://leetcode.com/problems/range-sum-query-2d-immutable/>

Solution

An easy dp problem but costs me half an hour and 10 WA or RE since I made a mistake on dealing with the array index. Reserve an buffer is a good choice.

Listing 235: Solution

```
public class NumMatrix {
    int [][] sums;

    public NumMatrix(int [][] matrix) {
5        if (matrix == null || matrix.length == 0) {
            return;
        }
        sums = new int[matrix.length + 1][matrix[0].length + 1];
        sums[1][1] = matrix[0][0];
10        for (int row = 1; row < matrix.length; row++) {
            sums[row + 1][1] += sums[row][1] + matrix[row][0];
        }
        for (int col = 1; col < matrix[0].length; col++) {
            sums[1][col + 1] += sums[1][col] + matrix[0][col];
15        }

        for (int row = 1; row < matrix.length; row++) {
```

```
20         for (int col = 1; col < matrix[0].length; col++) {
            sums[row + 1][col + 1] = sums[row][col + 1] + sums[row + 1][col]
            - sums[row][col] + matrix[row][col];
        }
    }

25     public int sumRegion(int row1, int col1, int row2, int col2) {
        return sums[row2 + 1][col2 + 1] - sums[row2 + 1][col1]
            - sums[row1][col2 + 1] + sums[row1][col1];
    }
30 }

// Your NumMatrix object will be instantiated and called as such:
// NumMatrix numMatrix = new NumMatrix(matrix);
// numMatrix.sumRegion(0, 1, 2, 3);
35 // numMatrix.sumRegion(1, 2, 3, 4);
```

Problem 181

198 House Robber

You are a professional robber planning to rob houses along a street.
Each house has a certain amount of money stashed,
the only constraint stopping you from robbing each of them
is that adjacent houses have security system connected and
it will automatically contact the police if two adjacent houses
were broken into on the same night.

Given a list of non-negative integers representing
the amount of money of each house,
determine the maximum amount of money you can rob tonight
without alerting the police.

<https://leetcode.com/problems/house-robber/>

Solution

A bit like finding the maximum increasing sequence, the dp function is exactly $\max(a[i-1], a[i-2] + \text{nums}[i])$, but take care of the initialization part (I made a mistake here), $a[0] = \text{nums}[0]$, $a[1] = \text{Math.max}(\text{nums}[0], \text{nums}[1])$;

Listing 236: Solution

```
public class Solution {  
    public int rob(int[] nums) {  
        if (nums.length == 0){  
            return 0;  
        }  
        if (nums.length == 1){  
            return nums[0];  
        }  
        if (nums.length == 2){  
            return Math.max(nums[0], nums[1]);  
        }  
        int[] sums = new int[nums.length];  
  
        sums[0] = nums[0];  
        sums[1] = Math.max(nums[0], nums[1]);  
        for (int idx = 2; idx < nums.length; idx++){  
            sums[idx] = Math.max(sums[idx - 2] + nums[idx], sums[idx - 1]);  
        }  
        return sums[nums.length - 1];  
    }  
}
```


Problem 182

213 House Robber II

Note: This is an extension of House Robber.

After robbing those houses on that street,
the thief has found himself a new place for his thievery
so that he will not get too much attention.
This time, all houses at this place are arranged in a circle.
That means the first house is the neighbor of the last one.
Meanwhile, the security system for these houses remain
the same as for those in the previous street.

Given a list of non-negative integers representing
the amount of money of each house,
determine the maximum amount of money you can rob tonight
without alerting the police.

<https://leetcode.com/problems/house-robber-ii/>

Solution

A cyclic DP problem, we can separate the problem into 2 linear DP problems by assuming whether rob the first house(so that the last can't be robbed) or don't rob it.

Listing 237: Solution

```
public class Solution {  
    public int rob(int[] nums) {  
        if (nums.length == 0)  
            return 0;  
        if (nums.length == 1)  
            return nums[0];  
        int[] sums = new int[nums.length];  
        // not robbing the 1st house  
        sums[0] = 0;  
        sums[1] = nums[1];  
        for (int idx = 2; idx < nums.length; idx++){  
            sums[idx] = Math.max(sums[idx - 1], sums[idx - 2] + nums[idx]);  
        }  
        int max = sums[nums.length - 1];  
        // robbing the 1st house  
        sums[0] = nums[0];  
        sums[1] = Math.max(nums[0], nums[1]);  
        for (int idx = 2; idx < nums.length - 1; idx++){  
            sums[idx] = Math.max(sums[idx - 1], sums[idx - 2] + nums[idx]);  
        }  
        return Math.max(max, sums[nums.length - 2]);  
    }  
}
```

Problem 183

337 House Robber III

The thief has found himself a new place for his thievery again. There is only one entrance to this area, called the "root." Besides the root, each house has one and only one parent house. After a tour, the smart thief realized that "all houses in this place forms a binary tree". It will automatically contact the police if two directly-linked houses were broken into on the same night.

Determine the maximum amount of money the thief can rob tonight without alerting the police.

Example 1:

```

      3
     /\
    2  3
     \  \
      3  1
  
```

Maximum amount of money the thief can rob = 3 + 3 + 1 = 7.

Example 2:

```

      3
     /\
    4  5
   /\  \
  1  3  1
  
```

Maximum amount of money the thief can rob = 4 + 5 = 9.

<https://leetcode.com/problems/house-robber-iii/>

Solution

When the thief see a house, he has 2 choices, rob it or don't rob it. When he decides to rob it, things are simple, he cannot rob the 2 children houses; but not robbing is a bit complicated, he need to re-decide when he meets the 2 children houses whether he should rob it.

Listing 238: Solution

```

/**
 * Definition for a binary tree node.
 * public class TreeNode {
 *     int val;
5  *     TreeNode left;
 *     TreeNode right;
 *     TreeNode(int x) { val = x; }
 * }
 */
10 public class Solution {
    public int rob(TreeNode root) {
        if (root == null){
            return 0;
        }
15    int exc = robExcludeRoot(root);
  
```

```
        int inc = robIncludeRoot(root);
        return Math.max(exc, inc);
    }

20    private int robExcludeRoot(TreeNode root) {
        if (root == null){
            return 0;
        }
        return rob(root.left) + rob(root.right);
25    }

    private int robIncludeRoot(TreeNode root){
        if (root == null){
            return 0;
30        }
        return root.val + robExcludeRoot(root.left) + robExcludeRoot(root.right);
    }
}
```

Problem 184

10 Regular Expression Matching

Implement regular expression matching with support for '.' and '*'.

'.' Matches any single character.

'*' Matches zero or more of the preceding element.

The matching should cover the entire input string (not partial).

The function prototype should be:

```
bool isMatch(const char *s, const char *p)
```

Some examples:

```
isMatch("aa","a") -> false
```

```
isMatch("aa","aa") -> true
```

```
isMatch("aaa","aa") -> false
```

```
isMatch("aa","a*") -> true
```

```
isMatch("aa",".*") -> true
```

```
isMatch("ab",".*") -> true
```

```
isMatch("aab","c*a*b") -> true
```

<https://leetcode.com/problems/regular-expression-matching/>

Solution

A very complicated backtracking problem, takes me 1 and a half hour to finish it (and even in an IDE!!!).....

Both this problem and next problem needs to be supplemented with missing DP or recursion solutions.

See <http://www.xuebuyuan.com/1936978.html>, good illustration.

Listing 239: Solution 1.1

```
public class Solution {
    private boolean isMatch(String s, int idxS, String p, int idxP) {
        if (idxS < s.length() && idxP >= p.length()) {
            return false;
        }
        if (idxP == p.length() && idxS == s.length()) {
            return true;
        }
        if ((idxP < p.length() - 1) && p.charAt(idxP + 1) == '*') {
            if (idxS < s.length()
                && (p.charAt(idxP) == s.charAt(idxS) || (p.charAt(idxP) == '.'))) {
                return isMatch(s, idxS + 1, p, idxP + 2)
                    || isMatch(s, idxS + 1, p, idxP)
                    || isMatch(s, idxS, p, idxP + 2);
            } else {
                return isMatch(s, idxS, p, idxP + 2);
            }
        }
        if (idxS < s.length()
            && (p.charAt(idxP) == '.' || s.charAt(idxS) == p.charAt(idxP))) {

```

```
        return isMatch(s, idxS + 1, p, idxP + 1);
    }
    return false;
25 }

    public boolean isMatch(String s, String p) {
        // if (s.length() == 0)
        return isMatch(s, 0, p, 0);
30 }
}
```

Append a rather simple code...

Listing 240: Solution 1.2

```
public class Solution {
    public boolean isMatch(String s, String p) {
        for(int i = 0; i < p.length(); s = s.substring(1)) {
            char c = p.charAt(i);
5           if(i + 1 >= p.length() || p.charAt(i + 1) != '*' )
                i++;
            else if(isMatch(s, p.substring(i + 2)))
                return true;

10           if(s.isEmpty() || (c != '.' && c != s.charAt(0)))
                return false;
        }

        return s.isEmpty();
15    }
}
```

Problem 185

44 Wildcard Matching

Implement wildcard pattern matching with support for '?' and '*'.

'?' Matches any single character.

'*' Matches any sequence of characters (including the empty sequence).

The matching should cover the entire input string (not partial).

The function prototype should be:

```
bool isMatch(const char *s, const char *p)
```

Some examples:

```
isMatch("aa","a") -> false
```

```
isMatch("aa","aa") -> true
```

```
isMatch("aaa","aa") -> false
```

```
isMatch("aa", "*") -> true
```

```
isMatch("aa", "a*") -> true
```

```
isMatch("ab", "?*") -> true
```

```
isMatch("aab", "c*a*b")    false
```

<https://leetcode.com/problems/wildcard-matching/>

Solution

Displaying the DP solution.

Listing 241: Solution

```
public class Solution {
    public boolean isMatch(String s, String p) {
        if (s.length() == 0){
            int idx = 0;
            while (idx < p.length() && p.charAt(idx) == '*'){
                idx++;
            }

            if (idx == p.length()){
                return true;
            }

            return false;
        }

        boolean dp[][] = new boolean[s.length() + 1][p.length() + 1];

        dp[0][0] = true;
        for (int j = 1 ; j <= p.length(); j ++ ){
            for (int i = 1 ; i <= s.length() ; i++){
                if (dp[i - 1][j - 1] && (s.charAt(i - 1) == p.charAt(j - 1)
                    || p.charAt(j - 1) == '?')){
                    dp[i][j] = true;
                }
            }
        }
    }
}
```

```
25         if (p.charAt(j - 1) == '*' ){
30             if (dp[i - 1][j - 1]){
35                 for (int k = i - 1 ; k <= s.length() ; k ++){
40                     dp[k][j] = true;
45                     break;
35             }
40             else if (dp[i][j - 1]){
45                 for (int k = i ; k <= s.length(); k ++){
40                     dp[k][j] = true;
45                     break;
35             }
40         }
45     }

    return dp[s.length()][p.length()];
}
```

Problem 186

73 Set Matrix Zeroes

Given a $m \times n$ matrix, if an element is 0,
set its entire row and column to 0. Do it in place.

Follow up:

Did you use extra space?

A straight forward solution using $O(mn)$ space is probably a bad idea.

A simple improvement uses $O(m + n)$ space, but still not the best solution.

Could you devise a constant space solution?

<https://leetcode.com/problems/set-matrix-zeroes/>

Solution

Solution using $O(mn)$ space is too trivial. First show the result using $O(mn)$ space, then the $O(1)$ space solution.

Listing 242: Solution 1.1

```
public class Solution {
    public void setZeroes(int [][] matrix) {
        int rowLength = matrix.length;
        if (rowLength == 0){
5           return ;
        }
        int colLength = matrix[0].length;
        boolean rowZero [] = new boolean[rowLength];
        boolean colZero [] = new boolean[colLength];

10
        for (int rowIdx = 0 ; rowIdx < rowLength ; rowIdx ++){
            for (int colIdx = 0 ; colIdx < colLength; colIdx ++){
                if(matrix[rowIdx][colIdx] == 0){
                    rowZero[rowIdx] = true;
                    colZero[colIdx] = true;
15                }
            }
        }

        for (int rowIdx = 0 ; rowIdx < rowLength ; rowIdx ++){
20            if (rowZero[rowIdx]){
                for (int colIdx = 0 ; colIdx < colLength ; colIdx ++){
                    matrix[rowIdx][colIdx] = 0;
                }
            }
25        }

        for (int colIdx = 0 ; colIdx < colLength ; colIdx ++){
            if (colZero[colIdx]){
30                for (int rowIdx = 0 ; rowIdx < rowLength ; rowIdx ++){
                    matrix[rowIdx][colIdx] = 0;
                }
            }
        }
    }
}
```



```

35     }
    }
}

```

The $O(1)$ space solution. Use the 1st(index 0) row and 1st column to mark 0 when there is a grid valued 0 with the same rowIdx or colIdx. Attention! We should deal with the 1st row and column separately, or it will be influenced and make all grids 0!

Listing 243: Solution 1.2

```

public class Solution {
    public void setZeroes(int [][] matrix) {
        int m = matrix.length;
        if (m == 0){
5           return;
        }
        int n = matrix[0].length;
        boolean zeroInFirstRow = false;
        boolean zeroInFirstCol = false;
10        for (int rowIdx = 0 ; rowIdx < m ; rowIdx ++){
            if (matrix[rowIdx][0] == 0){
                zeroInFirstCol = true;
                break;
            }
        }
15        for (int colIdx = 0 ; colIdx < n ; colIdx ++){
            if (matrix[0][colIdx] == 0){
                zeroInFirstRow = true;
                break;
            }
20        }

        for (int rowIdx = 0 ; rowIdx < m ; rowIdx ++){
            for (int colIdx = 0 ; colIdx < n ; colIdx ++){
25                if (matrix[rowIdx][colIdx] == 0){
                    matrix[rowIdx][0] = 0;
                    matrix[0][colIdx] = 0;
                }
            }
30        }

        for (int rowIdx = 1 ; rowIdx < m ; rowIdx ++){
            for (int colIdx = 1 ; colIdx < n ; colIdx ++){
35                if (matrix[rowIdx][0] == 0 || matrix[0][colIdx] == 0){
                    matrix[rowIdx][colIdx] = 0;
                }
            }
        }

40        if (zeroInFirstRow){
            for (int colIdx = 0 ; colIdx < n ; colIdx ++){
                matrix[0][colIdx] = 0;
            }
        }
    }
}

```

```
    }  
45    if (zeroInFirstCol){  
        for (int rowIdx = 0 ; rowIdx < m ; rowIdx ++){  
            matrix[rowIdx][0] = 0;  
        }  
50    }  
    }  
}
```

Problem 187

329 Longest Increasing Path in a Matrix

Given an integer matrix,
find the length of the longest increasing path.

From each cell, you can either move to four directions:
left, right, up or down.

You may NOT move diagonally or move outside of the boundary
(i.e. wrap-around is not allowed).

Example 1:

```
nums = [
  [9,9,4],
  [6,6,8],
  [2,1,1]
]
Return 4
The longest increasing path is [1, 2, 6, 9].
```

Example 2:

```
nums = [
  [3,4,5],
  [3,2,6],
  [2,2,1]
]
Return 4
The longest increasing path is [3, 4, 5, 6].
Moving diagonally is not allowed.
```

<https://leetcode.com/problems/longest-increasing-path-in-a-matrix/>

Solution

A DP(?), DFS and memoization problem. Not that difficult.

Listing 244: Solution

```
public class Solution {
    int [][] directions = {{0,1},{0,-1},{1,0},{-1,0}};

    public int findLongestPath(int [][] matrix, int [][] pathLen, int x, int y){
5      if (pathLen[x][y] != 0){
          return pathLen[x][y];
        }
        pathLen[x][y] = 1;
        for (int dirIdx = 0 ; dirIdx < directions.length ; dirIdx ++){
10         int i = x + directions[dirIdx][0];
            int j = y + directions[dirIdx][1];

            if (i < 0 || j < 0 || i >= matrix.length
                || j >= matrix[0].length || matrix[i][j] <= matrix[x][y]){
```

```
15         continue;
        }
        else{
            pathLen[x][y] = Math.max(pathLen[x][y]
20                , findLongestPath(matrix, pathLen, i, j) + 1);
        }
    }

    return pathLen[x][y];
}

25 public int longestIncreasingPath(int [][] matrix) {
    int m = matrix.length;
    if (m == 0){
        return 0;
30    }
    int n = matrix[0].length;
    int [][] pathLen = new int[m][n];

    int max = Integer.MIN_VALUE;

35    for (int i = 0 ; i < m ; i ++){
        for (int j = 0 ; j < n ; j ++){
            max = Math.max(max, findLongestPath(matrix, pathLen, i, j));
40        }
    }

    return max;
}
}
```

Problem 188

93 Restore IP Addresses

Given a string containing only digits,
restore it by returning all possible valid IP address combinations.

For example:

Given "25525511135",

return ["255.255.11.135", "255.255.111.35"]. (Order does not matter)

<https://leetcode.com/problems/restore-ip-addresses/>

Solution

Honestly, this is not such a difficult question, but there are so many edge cases in this problem, which make us busy debugging when we meet a new edge case.

Take care of those:

- (1) numbers starts with '0' except for 0 itself;
- (2) length out of bounds
- (3) more than 4 numbers generated
- (4) number more than 255
- (5) number with the length of more than 3

Listing 245: Solution

```
public class Solution {  
    List<String> result = new ArrayList<String>();  
  
    private void findIpAddresses(String s, int startIdx, List<Integer> curList){  
5         if (startIdx > s.length()){  
            return;  
        }  
        if (curList.size() >= 4 && startIdx < s.length()){  
            return;  
10        }  
        if (curList.size() == 4 && startIdx == s.length()){  
            result.add(formIpString(curList));  
            return;  
15        }  
  
        for (int endIdx = startIdx + 1 ;  
            endIdx <= Math.min(s.length(), startIdx + 3) ; endIdx++){  
            String newNum = s.substring(startIdx, endIdx);  
            if (endIdx > startIdx + 1 && s.charAt(startIdx) == '0'){  
20                return;  
            }  
            if (Integer.parseInt(newNum) > 255){  
                return;  
            }  
25            curList.add(Integer.parseInt(newNum));  
            findIpAddresses(s, endIdx, curList);  
            curList.remove(curList.size() - 1);  
        }  
    }  
}
```

```
    }  
    }  
30  
    private String formIpString(List<Integer> list){  
        return list.get(0) + "." + list.get(1) + "." +  
            list.get(2) + "." + list.get(3);  
    }  
35  
    public List<String> restoreIpAddresses(String s) {  
        findIpAddresses(s, 0, new ArrayList<Integer>());  
        return result;  
    }  
40 }
```

Problem 189

131 Palindrome Partitioning

Given a string *s*, partition *s* such that every substring of the partition is a palindrome.

Return all possible palindrome partitioning of *s*.

For example, given *s* = "aab",
Return

```
[
  ["aa","b"],
  ["a","a","b"]
]
```

<https://leetcode.com/problems/palindrome-partitioning/>

Solution

A not so difficult backtracking problem, add string to list if the string is palindrome.

Listing 246: Solution

```
public class Solution {
    List<List<String>> result = new ArrayList<List<String>>();

    private boolean isPalindrome(String str){
5         int start = 0 , end = str.length() - 1;
        while (start < end){
            if (str.charAt(start) != str.charAt(end)){
                return false;
            }
10         start ++;
            end --;
        }
        return true;
    }

15     private void findPalindrome(String str , int startIdx , List<String> curList){
        if (startIdx == str.length()){
            result.add(new ArrayList<String>(curList));
            return;
20        }
        for (int endIdx = startIdx + 1 ; endIdx <= str.length() ; endIdx ++){
            String newStr = str.substring(startIdx , endIdx);
            if (isPalindrome(newStr)){
                curList.add(newStr);
25                findPalindrome(str , endIdx , curList);
                curList.remove(curList.size() - 1);
            }
        }
30    }
}
```

```
    public List<List<String>> partition(String s) {  
        findPalindrome(s, 0, new ArrayList<String>());  
        return result;  
    }  
35 }
```


Problem 190

132 Palindrome Partitioning II

Given a string *s*, partition *s* such that every substring of the partition is a palindrome.

Return the minimum cuts needed for a palindrome partitioning of *s*.

For example, given *s* = "aab",

Return 1 since the palindrome partitioning ["aa","b"] could be produced using 1 cut.

<https://leetcode.com/problems/palindrome-partitioning-ii/>

Solution

Solution from previous Problem “Palindrome Partitioning” will get TLE with long sequence with so many palindromes within it.(Even when I changed the train of thought to find the longest palindrome first and do the pruning when result list size is larger than min)

So DP will be the alternative.

Listing 247: Solution

```
public class Solution {
    public int minCut(String s) {
        boolean[][] isPal = new boolean[s.length()][s.length()];
        int cut[] = new int[s.length()];
        isPal[0][0] = true;
        cut[0] = 0;
        for (int i = 0 ; i < s.length() ; i++){
            cut[i] = i;
            for (int j = 0 ; j <= i ; j++){
                if (s.charAt(i) == s.charAt(j)
                    && (i - j <= 1 || isPal[j + 1][i - 1])){
                    isPal[j][i] = true;
                    if (j == 0)
                        // no need to cut
                        cut[i] = 0;
                    else
                        cut[i] = Math.min(cut[i] , cut[j - 1] + 1);
                }
            }
        }
        return cut[s.length() - 1];
    }
}
```

Problem 191

38 Count and Say

The count-and-say sequence is the sequence of integers beginning as follows:

1, 11, 21, 1211, 111221, ...

1 is read off as "one 1" or 11.

11 is read off as "two 1s" or 21.

21 is read off as "one 2, then one 1" or 1211.

Given an integer n, generate the nth sequence.

Note: The sequence of integers will be represented as a string.

<https://leetcode.com/problems/count-and-say/>

Solution

Attention for the easy problems, not as easy as I thought.

Listing 248: Solution

```
public class Solution {
    private String generateNumber(String formerStr){
        String newStr = "";
        char compare = formerStr.charAt(0);
        int dupCount = 1;
        for (int idx = 1 ; idx < formerStr.length() ; idx++){
            while (idx < formerStr.length() && formerStr.charAt(idx) == compare){
                dupCount++;
                idx++;
            }
            newStr = newStr + dupCount + compare;
            if (idx < formerStr.length()){
                compare = formerStr.charAt(idx);
                dupCount = 1;
            }
        }
        if(dupCount == 1)
            newStr += "1" + compare;
        return newStr;
    }
    public String countAndSay(int n) {
        String initialStr = "1";
        while (n > 1){
            initialStr = generateNumber(initialStr);
            n--;
        }
        return initialStr;
    }
}
```

Problem 192

324 Wiggle Sort II

Given an unsorted array `nums`,
reorder it such that `nums[0] < nums[1] > nums[2] < nums[3] ...`.

Example:

- (1) Given `nums = [1, 5, 1, 1, 6, 4]`,
one possible answer is `[1, 4, 1, 5, 1, 6]`.
(2) Given `nums = [1, 3, 2, 2, 3, 1]`,
one possible answer is `[2, 3, 1, 3, 1, 2]`.

Note:

You may assume all input has valid answer.

Follow Up:

Can you do it in $O(n)$ time and/or in-place with $O(1)$ extra space?

<https://leetcode.com/problems/wiggle-sort-ii/>

Solution

Using `wiggleIndex` to re-index the array is the key to solving this. the formula is $(2 * i + 1) \% (n|1)$.
`A[0,1,2,3,4]` will be mapped to `a[1,3,5,0,2,4]`, respectively.

Listing 249: Solution

```
public class Solution {
    private int wiggleIndex(int idx, int length){
        return (2 * idx + 1) % (length | 1);
    }

    private int partitionMedian(int[] nums){
        int low = 0, high = nums.length - 1, targetIdx = nums.length / 2;

        while(true){
            int pivotIdx = -1;
            int pivot = nums[low];
            int start = low, end = high;

            while (low < high){
                while (low < high && nums[high] <= pivot){
                    high--;
                }
                nums[low] = nums[high];
                while (low < high && nums[low] >= pivot){
                    low++;
                }
                nums[high] = nums[low];
            }
            if (low == high){
                pivotIdx = low;
            }
            nums[pivotIdx] = pivot;
        }
    }
}
```

```
        if (pivotIdx == targetIdx){
            return pivot;
        }

        if (pivotIdx > targetIdx){
            high = pivotIdx - 1;
            low = start;
        }
        else{
            low = pivotIdx + 1;
            high = end;
        }
    }
}

public void wiggleSort(int[] nums) {
    int median = partitionMedian(nums);
    int current = 0, low = nums.length - 1, high = 0;
    while (current <= low){
        if (nums[wiggleIndex(current, nums.length)] == median){
            current++;
        }
        else if (nums[wiggleIndex(current, nums.length)] > median){
            int swap = nums[wiggleIndex(current, nums.length)];
            nums[wiggleIndex(current, nums.length)]
                = nums[wiggleIndex(high, nums.length)];
            nums[wiggleIndex(high, nums.length)] = swap;
            high++;
            current++;
        }
        else{
            int swap = nums[wiggleIndex(current, nums.length)];
            nums[wiggleIndex(current, nums.length)]
                = nums[wiggleIndex(low, nums.length)];
            nums[wiggleIndex(low, nums.length)] = swap;
            low--;
        }
    }
}
```

Problem 193

289 Game of Life

According to the Wikipedia's article:

"The Game of Life, also known simply as Life, is a cellular automaton devised by the British mathematician John Horton Conway in 1970."

Given a board with m by n cells, each cell has an initial state live (1) or dead (0). Each cell interacts with its eight neighbors (horizontal, vertical, diagonal) using the following four rules (taken from the above Wikipedia article):

Any live cell with fewer than two live neighbors dies, as if caused by under-population.
Any live cell with two or three live neighbors lives on to the next generation.
Any live cell with more than three live neighbors dies, as if by over-population..
Any dead cell with exactly three live neighbors becomes a live cell, as if by reproduction.
Write a function to compute the next state (after one update) of the board given its current state.

Follow up:

Could you solve it in-place?

Remember that the board needs to be updated at the same time:
You cannot update some cells first and then use their updated values to update other cells.

In this question, we represent the board using a 2D array.

In principle, the board is infinite, which would cause problems when the active area encroaches the border of the array.
How would you address these problems?

<https://leetcode.com/problems/game-of-life/>

Solution

Bit manipulation is the first thought which comes into my mind, use 1st digit as next state and 2nd digit for current state.

Listing 250: Solution

```
public class Solution {
    int [][] directions = {{-1,0},{0,-1},{-1,-1},{1,0},
                           {0,1},{1,1},{-1,1},{1,-1}};
    public void gameOfLife(int [][] board) {
5        if (board.length == 0){
            return;
        }
        for(int i = 0 ; i < board.length ; i++){
            for (int j = 0 ; j < board[0].length ; j++){
10                int liveNeighbors = 0;
```

```
15         for (int dirIdx = 0 ; dirIdx < 8 ; dirIdx ++){
16             int x = i + directions[dirIdx][0];
17             int y = j + directions[dirIdx][1];
18             if (x < 0 || y < 0 || x >= board.length
19                 || y >= board[0].length){
20                 continue;
21             }
22             liveNeighbors += board[x][y] & 1;
23         }
24
25         if (liveNeighbors < 2 || liveNeighbors > 3){
26             board[i][j] = (0 << 1) + board[i][j];
27         } else if (liveNeighbors == 2){
28             board[i][j] = (board[i][j] << 1) + board[i][j];
29         } else if (liveNeighbors == 3){
30             board[i][j] = (1 << 1) + board[i][j];
31         }
32     }
33 }
34
35 for (int i = 0 ; i < board.length ; i ++){
36     for (int j = 0 ; j < board[0].length ; j++){
37         board[i][j] = (board[i][j] >>> 1);
38     }
39 }
40 }
```

Problem 194

165 Compare Version Numbers

Compare two version numbers version1 and version2.

If version1 > version2 return 1, if version1 < version2 return -1, otherwise return 0.

You may assume that the version strings are

non-empty and contain only digits and the . character.

The . character does not represent

a decimal point and is used to separate number sequences.

For instance, 2.5 is not "two and a half" or

"half way to version three",

it is the fifth second-level revision of the second first-level revision.

Here is an example of version numbers ordering:

0.1 < 1.1 < 1.2 < 13.37

<https://leetcode.com/problems/compare-version-numbers/>

Solution

Though simple, do not miss the edge cases when '1.0.0' equals '1'

Listing 251: Solution

```
public class Solution {
    public int compareVersion(String version1, String version2) {
        int idx1 = 0;
        int idx2 = 0;

        while (idx1 < version1.length() && idx2 < version2.length()){
            int num1 = 0, num2 = 0;
            while (idx1 < version1.length() && version1.charAt(idx1) != '.'){
                num1 = num1 * 10 + (version1.charAt(idx1) - '0');
                idx1++;
            }
            while (idx2 < version2.length() && version2.charAt(idx2) != '.'){
                num2 = num2 * 10 + (version2.charAt(idx2) - '0');
                idx2++;
            }

            if (num1 > num2){
                return 1;
            }
            else if (num2 > num1){
                return -1;
            }

            idx1++;
            idx2++;
        }
        if (idx1 < version1.length()){
```

```
30         while (idx1 < version1.length()
           && (version1.charAt(idx1) == '.'
           || version1.charAt(idx1) == '0')){
           idx1 ++;
           }
           if (idx1 < version1.length()){
           return 1;
35       }
       }
       else if (idx2 < version2.length()){
       while (idx2 < version2.length()
           && (version2.charAt(idx2) == '.'
           || version2.charAt(idx2) == '0')){
40           idx2 ++;
           }
           if (idx2 < version2.length()){
           return -1;
45       }
       }
       return 0;
50   }
```


Problem 195

58 Length of Last Word

Given a string *s* consists of upper/lower-case alphabets and empty space characters ' ', return the length of last word in the string.

If the last word does not exist, return 0.

Note: A word is defined as a character sequence consists of non-space characters only.

For example,
Given *s* = "Hello World",
return 5.

<https://leetcode.com/problems/length-of-last-word/>

Solution

Counting from backward will benefit a lot.

Listing 252: Solution

```
public class Solution {  
    public int lengthOfLastWord(String s) {  
        int len = s.length();  
        int idx = len - 1;  
        int count = 0;  
        while (idx >= 0){  
            // skip spaces in the end  
            while (idx >= 0 && s.charAt(idx) == ' '){  
                idx --;  
            }  
            if (idx < 0){  
                return 0;  
            }  
            while (idx >= 0 && s.charAt(idx) != ' '){  
                count ++;  
                idx --;  
            }  
            break;  
        }  
        return count;  
    }  
}
```

Problem 196

118 Pascal's Triangle

Given numRows, generate the first numRows of Pascal's triangle.

For example, given numRows = 5,
Return

```
[
  [1],
  [1,1],
  [1,2,1],
  [1,3,3,1],
  [1,4,6,4,1]
]
```

<https://leetcode.com/problems/pascals-triangle/>

Solution

Such an easy problem.

Listing 253: Solution

```
public class Solution {
    public List<List<Integer>> generate(int numRows) {
        if (numRows == 0){
            return new ArrayList<List<Integer>>();
        }
        List<List<Integer>> result = new ArrayList<List<Integer>>();
        List<Integer> oldList = new ArrayList<Integer>();
        if (numRows != 0){
            oldList.add(1);
        }
        result.add(new ArrayList<Integer>(oldList));
        for (int i = 0; i < numRows - 1 ; i++){
            List<Integer> newList = new ArrayList<Integer>();
            for (int newIdx = 0 ; newIdx < oldList.size() + 1 ; newIdx++){
                if (newIdx == 0 || newIdx == oldList.size()){
                    newList.add(1);
                }
                else{
                    newList.add(oldList.get(newIdx - 1) + oldList.get(newIdx));
                }
            }
            result.add(new ArrayList<Integer>(newList));
            oldList = newList;
        }
        return result;
    }
}
```

Problem 197

119 Pascal's Triangle II

Given an index k , return the k th row of the Pascal's triangle.

For example, given $k = 3$,
Return $[1, 3, 3, 1]$.

Note:

Could you optimize your algorithm to use only $O(k)$ extra space?

<https://leetcode.com/problems/pascals-triangle-ii/>

Solution

Such an easy problem, too.

Listing 254: Solution

```
public class Solution {  
    public List<Integer> getRow(int rowIndex) {  
        List<Integer> list = new ArrayList<Integer>(rowIndex + 1);  
        for (int i = 0 ; i <= rowIndex ; i++){  
5           list.add(0);  
        }  
        for (int i = 0 ; i <= rowIndex ; i++){  
            list.set(i, 1);  
            for (int j = i - 1 ; j > 0 ; j --){  
10                list.set(j, list.get(j) + list.get(j - 1));  
            }  
        }  
        return list;  
    }  
15 }
```

Problem 198

6 ZigZag Conversion

The string "PAYPALISHIRING" is written in a zigzag pattern on a given number of rows like this:
(you may want to display this pattern in a fixed font for better legibility)

```
P   A   H   N
A P L S I I G
Y   I   R
```

And then read line by line: "PAHNAPLSIIGYIR"

Write the code that will take a string and make this conversion given a number of rows:

```
string convert(string text, int numRows);
convert("PAYPALISHIRING", 3) should return "PAHNAPLSIIGYIR".
```

<https://leetcode.com/problems/zigzag-conversion/>

Solution

Such an easy problem, three, find the rule and the problem can be solved easily.

Listing 255: Solution

```
public class Solution {
    public String convert(String s, int numRows) {
        String result = "";
        if (s.length() == 0 || numRows <= 1){
            return s;
        }
        for (int row = 0 ; row < numRows ; row++){
            for (int col = row ; col < s.length() ; col += (numRows - 1) * 2){
                result += s.charAt(col);

                if (row != 0 && row != numRows - 1){
                    // except for 1st and last row
                    int right = 2 * row;
                    int left = 2 * (numRows - 1) - right;
                    if (col + left < s.length()){
                        result += s.charAt(col + left);
                    }
                }
            }
        }

        return result;
    }
}
```

Problem 199

273 Integer to English Words

Convert a non-negative integer to its english words representation.
Given input is guaranteed to be less than $2^{31} - 1$.

For example,

123 -> "One Hundred Twenty Three"

12345 -> "Twelve Thousand Three Hundred Forty Five"

1234567 ->

"One Million Two Hundred Thirty Four Thousand Five Hundred Sixty Seven"

Hint:

Did you see a pattern in dividing the number into chunk of words? For example, 123 and 123000.

Group the number by thousands (3 digits).

You can write a helper function that takes a number less than 1000 and convert just that chunk to words.

There are many edge cases. What are some good test cases?

Does your code work with input such as 0? Or 1000010?

(middle chunk is zero and should not be printed out)

<https://leetcode.com/problems/integer-to-english-words/>

Solution

Hint is enough for solving the problem, honestly. But, never enough for all the edge cases(especially when involved with spaces).

Listing 256: Solution

```
public class Solution {
    String[] underTwenty = { "", "One", "Two", "Three", "Four", "Five", "Six",
                             "Seven", "Eight", "Nine", "Ten", "Eleven", "Twelve", "Thirteen",
                             "Fourteen", "Fifteen", "Sixteen", "Seventeen", "Eighteen",
5      "Nineteen" };
    String[] productTen = { "", "", "Twenty", "Thirty", "Forty", "Fifty", "Sixty",
                             "Seventy", "Eighty", "Ninety" };
    String[] thousands = { "", "Thousand", "Million", "Billion" };

10    public String numberToWords(int num) {
        String result = "";
        int cur = num;
        int thousandCnt = 0;
        if (num == 0){
15      return "Zero";
        }
        List<String> numbers = new ArrayList<String>();
        while (cur > 0){
            int qou = cur / 1000;
20      int rem = cur % 1000;

            numbers.add(transform(rem));
            if (qou != 0){
```

```

        thousandCnt ++;
25    }
    cur = qou;
}
for (int i = thousandCnt ; i >= 0 ; i --){
    String part = numbers.get(i);
30    if (part.trim().length() != 0){
        result = result + part + thousands[i] + " ";
    }
}

35    int endIdx = result.length() - 1;
    while (result.charAt(endIdx) == ' '){
        endIdx --;
    }
    result = result.substring(0, endIdx + 1);
40    return result;
}

private String transform(int num) {
    String result = "";
    int hun = num / 100;
45    int tens = num % 100;
    int ten = tens / 10;
    int one = num % 10;
    if (hun > 0) {
        result += underTwenty[hun] + " Hundred ";
50    }
    if (tens < 20 && tens != 0) {
        result += underTwenty[tens] + " ";
    } else if (tens != 0){
        result += productTen[ten] + " " +
55        ((one > 0) ? underTwenty[one] + " " : "");
    }

    return result;
}
60 }
```

Problem 200

310 Minimum Height Trees

For a undirected graph with tree characteristics,
 we can choose any node as the root.
 The result graph is then a rooted tree.
 Among all possible rooted trees,
 those with minimum height are called minimum height trees (MHTs).
 Given such a graph,
 write a function to find all the MHTs and return a list of their root labels.

Format

The graph contains n nodes which are labeled from 0 to $n - 1$.
 You will be given the number n and a list of undirected edges
 (each edge is a pair of labels).

You can assume that no duplicate edges will appear in edges.
 Since all edges are undirected, $[0, 1]$ is the same as $[1, 0]$
 and thus will not appear together in edges.

Example 1:

Given $n = 4$, edges = $[[1, 0], [1, 2], [1, 3]]$

```

      0
      |
      1
     / \
    2   3
return [1]
```

Example 2:

Given $n = 6$, edges = $[[0, 3], [1, 3], [2, 3], [4, 3], [5, 4]]$

```

    0  1  2
     \ | /
      3
      |
      4
      |
      5
return [3, 4]
```

Hint:

How many MHTs can a graph have at most?

Note:

(1) According to the definition of tree on Wikipedia:
 ``a tree is an undirected graph in which any two vertices
 are connected by exactly one path. In other words,

```
any connected graph without simple cycles is a tree.''
```

(2) The height of a rooted tree is the number of edges on the longest downward path between the root and a leaf.

<https://leetcode.com/problems/minimum-height-trees/>

Solution

The idea is like topological sort, since this is an undirected graph, find the node with degree of 1, remove them from the graph, until 1 or 2 nodes left.

Back news is: I used 'while (remaining > 2)' to check if I get the result, which means I'm not totally understanding it.

Take care when $n == 1$.

Listing 257: Solution

```
public class Solution {
    public List<Integer> findMinHeightTrees(int n, int[][] edges) {
        List<Integer> result = new ArrayList<Integer>();
        if (n == 1){
5           result.add(0);
           return result;
        }

        int degree[] = new int[n];
10       List<Integer> adjacency = new ArrayList<Integer>[n];
        for (int i = 0 ; i < n ; i++){
            adjacency[i] = new ArrayList<Integer>();
        }

15       for (int[] edge : edges){
            int node1 = edge[0];
            int node2 = edge[1];
            degree[node1]++;
            degree[node2]++;
20       adjacency[node1].add(node2);
            adjacency[node2].add(node1);
        }

        Queue<Integer> leaves = new LinkedList<Integer>();
25       for (int i = 0 ; i < n ; i++){
            if (degree[i] == 1){
                leaves.offer(i);
            }
        }

30       int remaining = n;
        while (remaining > 2){
            remaining -= leaves.size();
            List<Integer> leavesList = new ArrayList<Integer>();
35       while (!leaves.isEmpty()){
                leavesList.add(leaves.poll());
            }
            for (int leafIdx : leavesList){
```



```

        List<Integer> neighbors = adjacency[leafIdx];
40      for (int neighborIdx : neighbors){
        degree[neighborIdx] --;
        if (degree[neighborIdx] == 1){
            leaves.offer(neighborIdx);
        }
45      }
    }
    while (!leaves.isEmpty()){
50      result.add(leaves.poll());
    }
    return result;
}
}
```

Problem 201

300 Longest Increasing Subsequence

Given an unsorted array of integers,
find the length of longest increasing subsequence.

For example,

Given [10, 9, 2, 5, 3, 7, 101, 18],

The longest increasing subsequence is [2, 3, 7, 101],
therefore the length is 4.

Note that there may be more than one LIS combination,
it is only necessary for you to return the length.

Your algorithm should run in $O(n^2)$ complexity.

Follow up: Could you improve it to $O(n \log n)$ time complexity?

<https://leetcode.com/problems/longest-increasing-subsequence/>

Solution

As for the $O(n^2)$ solution, it is an easy DP problem, start of LCS(Longest Common Substring/Subsequence).

However, optimizing it to $O(n \log n)$ is tricky.

First, the $O(n^2)$ solution.

Listing 258: Solution 1.1

```
public class Solution {  
    public int lengthOfLIS(int[] nums) {  
        if (nums.length == 0 || nums.length == 1){  
            return nums.length;  
        }  
        int dp[] = new int[nums.length];  
        int max = 0;  
        for (int idx = 0 ; idx < nums.length ; idx++){  
            dp[idx] = 1;  
            for (int formerIdx = 0 ; formerIdx < idx ; formerIdx++){  
                if (nums[idx] > nums[formerIdx]){  
                    dp[idx] = Math.max(dp[idx], dp[formerIdx] + 1);  
                }  
            }  
            max = Math.max(max, dp[idx]);  
        }  
        return max;  
    }  
}
```

Then the $O(n \log n)$ solution. The key insight is to find the $O(\log n)$ factor. Below are train of thought from Discuss module:(or refer to this blog:)

<https://segmentfault.com/a/1190000003819886>

Runtime: To get an $O(n \log n)$ runtime, I'm going to create a second list S . (Stick with me for now – I'll get rid of it in a minute to get $O(1)$ space.) I'll do a single pass through $nums$, and as I look at each element:

The length of S will be equal to the length of the longer subsequence I've found to that point.

The last element of S will be the last element of that subsequence. (However, the earlier elements may no longer be part of that sequence – S is not actually the subsequence itself.)

At the end, the length of S will be our solution.

S will be sorted at all times. Each new element is inserted into S , replacing the smallest element in S that is not smaller than it (which we can find with a binary search). If that element is larger than the last element of S , then we extend S by one – maintaining both properties.

For example, if

$nums = [5,6,7,1,2,8,3,4,0,5,9]$

then after we process the 7:

$S = [5,6,7]$

after we process the 2:

$S = [1,2,7]$

after we process the 8:

$S = [1,2,7,8]$

Then we process the 3:

$S = [1,2,3,8]$

We process the 4:

$S = [1,2,3,4]$

and now the next three elements:

$S = [0,2,3,4,5,9]$

S is not the actual subsequence, but it is the right length (end ends in the right number).

We are making n passes, and doing a binary search each time. So $O(n \log n)$ time.

Space: Assuming we are allowed to destroy the list, we don't need S . Since S will never be larger than the number of elements we have looked at, and we only need to look at each element once, we can just use the beginning of $nums$ for S (keeping track of the size of " S " in a separate variable).

Listing 259: Solution 1.2

```

public class Solution {
    public int lengthOfLIS(int[] nums) {
        if (nums.length == 0 || nums.length == 1){
            return nums.length;
5        }
        int[] tails = new int[nums.length];
        int len = 0;
        tails[len++] = nums[0];

10        for (int idx = 1 ; idx < nums.length ; idx++){
            if (nums[idx] < tails[0]){
                tails[0] = nums[idx];
            } else if (nums[idx] > tails[len - 1]){
                tails[len++] = nums[idx];
15            } else{
                int fitPos = binarySearch(tails, nums[idx], len);
                tails[fitPos] = nums[idx];
            }
        }

20        return len;
    }
}

```

```

    }

    private int binarySearch(int[] tails, int target, int curLen){
25      int start = 0, end = curLen - 1;
        while (start < end){
            int mid = (end - start) / 2 + start;
            if (tails[mid] == target){
                return mid;
30            }
            else if (tails[mid] > target){
                end = mid;
            }
            else{
35                start = mid + 1;
            }
        }
        if (start == end){
            return start;
40        }
        else{
            return (tails[start] > target) ? start : end;
        }
    }
45 }

```

Moreover, if we want to recover the actual subsequence, we may need to do some extra work. My solution is to maintain a LinkedList for each position, when there is a element bigger than former ones coming, add -1 to its head, the number of -1 depends on the maximum depth of former linked list. See code in listing 260

Listing 260: Solution 1.2 - Recover actual subsequence

```

import java.io.*;
import java.util.*;
public class LIS {
    public static void main(String[] args){
5        try {
            BufferedReader br = new BufferedReader(
                new InputStreamReader(System.in));
            int count = Integer.parseInt(br.readLine());
            int nums[] = new int[count];
10            for (int i = 0 ; i < count ; i++){
                nums[i] = Integer.parseInt(br.readLine());
            }

            List<List<Integer>> seq = new ArrayList<List<Integer>>();

15            List<Integer> initialList = new ArrayList<Integer>();
            initialList.add(nums[0]);
            seq.add(initialList);

20            for (int i = 1; i < count ; i++){
                List<Integer> lastEle = seq.get(seq.size() - 1);
                if (nums[i] > lastEle.get(lastEle.size() - 1)){

```

```

        List<Integer> endList = new ArrayList<Integer>();
        int maxLevel = getMaxLevel(seq, seq.size());
25     for (int idx = 0 ; idx < maxLevel ; idx++){
            endList.add(-1);
        }
        endList.add(nums[i]);
        seq.add(endList);
30     }
    else if (nums[i] == lastEle.get(lastEle.size() - 1)){
        continue;
    }
    else{
35         insertIntoSeq(seq, nums[i]);
    }
}

    int maxLevel = getMaxLevel(seq, seq.size());
40     System.out.println(seq.size());
    for (int seqIdx = seq.size() - 1 ; seqIdx >= 0 ; seqIdx --){
        List<Integer> curList = seq.get(seqIdx);
        if (curList.size() - 1 <= maxLevel){
            maxLevel = curList.size() - 1;
45         }
        System.out.print(curList.get(maxLevel) + " ");
    }
    br.close();
} catch (IOException e) {
50     e.printStackTrace();
}

private static int getMaxLevel(List<List<Integer>> seq, int idx) {
55     int maxLevel = -1;
    for (int seqIdx = 0 ; seqIdx < idx ; seqIdx++){
        maxLevel = Math.max(maxLevel, seq.get(seqIdx).size() - 1);
    }
    return maxLevel;
60 }

private static void insertIntoSeq(List<List<Integer>> seq, int num) {
    int start = 0 , end = seq.size() - 1;
    while (start < end){
65         int mid = start + (end - start) / 2;
        int cur = seq.get(mid).get(seq.get(mid).size() - 1);
        if (num < cur){
            end = mid;
        }
70         else{
            start = mid + 1;
        }
    }
    int maxLevel = getMaxLevel(seq, start);
75     int sub = maxLevel - seq.get(start).size();

```

```

        List<Integer> newList = new ArrayList<Integer>();
        for (int i = 0 ; i < sub ; i ++){
            newList.add(-1);
        }
80     for (int ele : seq.get(start)){
            newList.add(ele);
        }
        newList.add(num);
        seq.set(start , newList);
85     }
}
```

Problem 202

334 Increasing Triplet Subsequence

Given an unsorted array return whether an increasing subsequence of length 3 exists or not in the array.

Formally the function should:

Return true if there exists i, j, k

such that $arr[i] < arr[j] < arr[k]$ given $0 \leq i < j < k \leq n-1$ else return false.

Your algorithm should run in $O(n)$ time complexity and $O(1)$ space complexity.

Examples:

Given [1, 2, 3, 4, 5],
return true.

Given [5, 4, 3, 2, 1],
return false.

<https://leetcode.com/problems/increasing-triplet-subsequence/>

Solution

Use idea of Solution 1.2 of Problem “Longest Increasing Subsequence”, by having smaller, larger integers to store current minimum element and second smallest element after the minimum element, if there is another element larger than larger(2nd smallest after smallest/ or maybe before smallest, but has another smaller element ahead), there will be an increasing sequence with the length of 3.

Listing 261: Solution

```
public class Solution {
    public boolean increasingTriplet(int[] nums) {
        if (nums.length < 3){
            return false;
        }
        int len = 1;
        int smaller = nums[0];
        int larger = -1;

        for (int idx = 1 ; idx < nums.length ; idx++){
            if (nums[idx] < smaller){
                smaller = nums[idx];
            }
            else if (nums[idx] > smaller){
                if (len == 1){
                    larger = nums[idx];
                    len++;
                }
                if (len == 2){
                    if (nums[idx] > larger){
                        return true;
                    }
                    else{
                        larger = nums[idx];
                    }
                }
            }
        }
    }
}
```

```

25     }
    }
}

return false;

30 }
```


Problem 203

208 Implement Trie (Prefix Tree)

Implement a trie with insert, search, and startsWith methods.

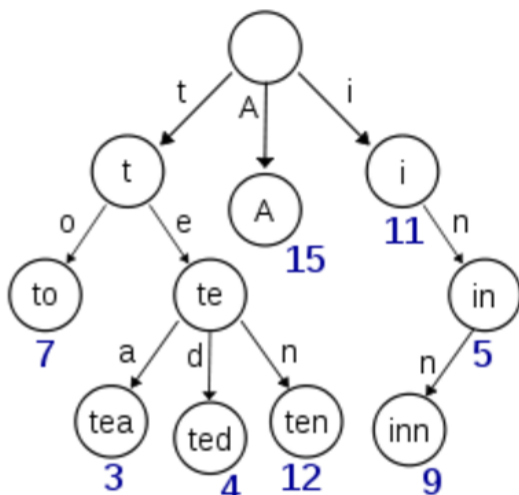
<https://leetcode.com/problems/implement-trie-prefix-tree/>

Solution

The problem itself is not difficult at all, but Trie(Prefix Tree) is what we need to know.

Trie树又被称为字典树、前缀树，是一种用于快速检索的多叉树。Trie树可以利用字符串的公共前缀来节省存储空间。

但如果系统存在大量没有公共前缀的字符串，相应的Trie树将非常消耗内存。（下图为Wiki上的Trie树示意图，<https://en.wikipedia.org/wiki/Trie>）



Trie树的基本性质如下：

- 根节点不包括字符，除根节点外每个节点包括一个支付。
- 从根节点到某一节点，路径上经过的字符连接起来，即为对应的字符串。
- 每个节点的所有子节点包含的字符串各不相同。

本题中的Trie树可以简单实现如下：

Trie树节点数据结构定义如下：

1. val表示该节点对应的字符
2. 子节点简单用一个数组表示，这样实现比较简单，但比较耗费内存
3. isWord标记从根节点到该节点是否表示一个单词，还是另一单词的前缀。

Listing 262: Solution

```
class TrieNode {  
    // Initialize your data structure here.  
    TrieNode[] children = new TrieNode[26];  
    boolean isWord;  
  
    public TrieNode() {  
        this.isWord = false;  
        for (int i = 0 ; i < 26 ; i ++){  
            children[i] = null;  
        }  
    }  
}
```

```
    }  
}  
  
15 public class Trie {  
    private TrieNode root;  
  
    public Trie() {  
        root = new TrieNode();  
    }  
20  
    // Inserts a word into the trie.  
    public void insert(String word) {  
        if (word.length() == 0){  
            return;  
25        }  
        TrieNode cur = this.root;  
  
        for (int wordIdx = 0 ; wordIdx < word.length() ; wordIdx ++){  
            int childrenIdx = word.charAt(wordIdx) - 'a';  
30            if (cur.children[childrenIdx] == null){  
                cur.children[childrenIdx] = new TrieNode();  
            }  
  
            cur = cur.children[childrenIdx];  
35        }  
        cur.isWord = true;  
    }  
  
    // Returns if the word is in the trie.  
40    public boolean search(String word) {  
        if (word.length() == 0){  
            return false;  
        }  
        TrieNode cur = this.root;  
45  
        for (int wordIdx = 0 ; wordIdx < word.length() ; wordIdx ++){  
            int childrenIdx = word.charAt(wordIdx) - 'a';  
            if (cur.children[childrenIdx] == null){  
                return false;  
50            }  
            cur = cur.children[childrenIdx];  
        }  
        return cur.isWord;  
    }  
55  
    // Returns if there is any word in the trie  
    // that starts with the given prefix.  
    public boolean startsWith(String prefix) {  
        if (prefix.length() == 0){  
60            return false;  
        }  
        TrieNode cur = this.root;  
        for (int prefixIdx = 0 ; prefixIdx < prefix.length() ; prefixIdx ++){
```

```
        int childrenIdx = prefix.charAt(prefixIdx) - 'a';
65      if (cur.children[childrenIdx] == null){
          return false;
        }
        cur = cur.children[childrenIdx];
70      }

      return true;
    }
}

75 // Your Trie object will be instantiated and called as such:
// Trie trie = new Trie();
// trie.insert("somestring");
// trie.search("key");
```

Problem 204

211 Add and Search Word - Data structure design

Design a data structure that supports the following two operations:

```
void addWord(word)
bool search(word)
search(word) can search a literal word or a regular expression string
containing only letters a-z or ..
A . means it can represent any one letter.
```

For example:

```
addWord("bad")
addWord("dad")
addWord("mad")
search("pad") -> false
search("bad") -> true
search(".ad") -> true
search("b..") -> true
```

Note:

You may assume that all words are consist of lowercase letters a-z.

You should be familiar with how a Trie works.

If not, please work on this problem: Implement Trie (Prefix Tree) first.

<https://leetcode.com/problems/add-and-search-word-data-structure-design/>

Solution

Not the same as Problem “Implement Trie”, a bit more difficult, use a recursive solution to solve the problem.

Listing 263: Solution

```
class TrieNode{
    TrieNode[] children = new TrieNode[26];
    boolean isWord;
    char val;

    public TrieNode(){
        isWord = false;
        for (int childrenIdx = 0 ;
            childrenIdx < children.length ; childrenIdx++){
            children[childrenIdx] = null;
        }
    }

    public TrieNode(char val){
        this();
        this.val = val;
    }
}
```

```
public class WordDictionary {
20     TrieNode root = new TrieNode();
    // Adds a word into the data structure.
    public void addWord(String word) {
        TrieNode cur = root;
        for (int idx = 0 ; idx < word.length() ; idx++){
25             int cIdx = word.charAt(idx) - 'a';
            if (cur.children[cIdx] == null){
                cur.children[cIdx] = new TrieNode(word.charAt(idx));
            }
            cur = cur.children[cIdx];
30        }
        cur.isWord = true;
    }

    // Returns if the word is in the data structure. A word could
35    // contain the dot character '.' to represent any one letter.
    public boolean search(String word) {
        TrieNode cur = root;
        return match(cur, word, 0);
    }

40    private boolean match(TrieNode cur, String word, int startIdx){
        boolean result = false;

        if (cur == null){
45            return false;
        }

        if (startIdx >= word.length()){
            return cur.isWord;
50        }
        char curCh = word.charAt(startIdx);
        if (curCh == '.'){
            for (TrieNode select : cur.children){
                if (select != null){
55                    result = result | match(select, word, startIdx + 1);
                    if (result){
                        return true;
                    }
                }
            }
60        }
        else{
            cur = cur.children[curCh - 'a'];
            if (cur == null){
65                return false;
            }
            else{
                char compareCh = cur.val;
                // Is this condition check necessary?
70                if (curCh == compareCh){
                    return match(cur, word, startIdx + 1);
                }
            }
        }
    }
}
```

```
75         }
           else{
               return false;
           }
       }
   }

   return result;
}

// Your WordDictionary object will be instantiated and called as such:
// WordDictionary wordDictionary = new WordDictionary();
85 // wordDictionary.addWord("word");
// wordDictionary.search("pattern");
```

Problem 205

79 Word Search

Given a 2D board and a word, find if the word exists in the grid.

The word can be constructed from letters of sequentially adjacent cell, where "adjacent" cells are those horizontally or vertically neighboring. The same letter cell may not be used more than once.

For example,
Given board =

```
[
  ['A','B','C','E'],
  ['S','F','C','S'],
  ['A','D','E','E']
]
word = "ABCCED", -> returns true,
word = "SEE", -> returns true,
word = "ABCB", -> returns false.
```

<https://leetcode.com/problems/word-search/>

Solution

Problem 79.

Solution is using straightforward backtracking method, be sure to return as soon as we get result, otherwise will get a TLE error.(I stuck in misusing || as |).

Listing 264: Solution

```
public class Solution {
    // int[][] directions = {{1,0},{-1,0},{0,1},{0,-1}};
    private boolean wordSearch(char[][] board, boolean[][] visited,
5      int curX, int curY, String word, int wordIdx){
        int m = board.length;
        int n = board[0].length;

        if (wordIdx >= word.length()){
            return true;
10        }
        if (curX < 0 || curY < 0 || curX >= board.length || curY >= board[0].length
            || word.charAt(wordIdx) != board[curX][curY]
            || visited[curX][curY] == true){
            return false;
15        }

        visited[curX][curY] = true;

        boolean result=wordSearch(board,visited, curX-1, curY, word, wordIdx+1)
20          || wordSearch(board,visited, curX+1, curY, word, wordIdx+1)
          || wordSearch(board,visited, curX, curY-1, word, wordIdx+1)
          || wordSearch(board,visited, curX, curY+1, word, wordIdx+1);
```

```
        visited[curX][curY] = false;

25     return result;
    }

    public boolean exist(char[][] board, String word) {
        if (board.length == 0 || board[0].length == 0){
30             return false;
        }
        boolean[][] visited = new boolean[board.length][board[0].length];

        for (int i = 0 ; i < board.length ; i++){
35             for (int j = 0 ; j < board[0].length ; j++){
                if (wordSearch(board, visited, i, j, word, 0)){
                    return true;
                }
            }
40         }

        return false;
    }
}
```


Problem 206

212 Word Search II

Given a 2D board and a list of words from the dictionary, find all words in the board.

Each word must be constructed from letters of sequentially adjacent cell, where "adjacent" cells are those horizontally or vertically neighboring. The same letter cell may not be used more than once in a word.

For example,

Given words = ["oath", "pea", "eat", "rain"] and board =

```
[
  ['o','a','a','n'],
  ['e','t','a','e'],
  ['i','h','k','r'],
  ['i','f','l','v']]
```

Return ["eat", "oath"].

Note:

You may assume that all inputs are consist of lowercase letters a-z.

[click to show hint.](#)

You would need to optimize your backtracking to pass the larger test. Could you stop backtracking earlier?

If the current candidate does not exist in all words' prefix, you could stop backtracking immediately.

What kind of data structure could answer such query efficiently?

Does a hash table work? Why or why not? How about a Trie?

If you would like to learn how to implement a basic trie, please work on this problem: [Implement Trie \(Prefix Tree\)](#) first.

<https://leetcode.com/problems/word-search-ii/>

Solution

This is the first problem I cannot finish for the first time!! Because I even cannot finish the testing of big data in Eclipse! More than 65536 bytes! And it stuck my computer!

For optimizations, see

<https://leetcode.com/discuss/77851/java-15ms-easiest-solution-100-00%25>

Some notes:

Combing them, Trie is the natural choice. Notice that:

TrieNode is all we need. search and startsWith are useless.

No need to store character at TrieNode. c.next[i] != null is enough.

Never use c1 + c2 + c3. Use StringBuilder.

No need to use $O(n^2)$ extra space visited[m][n].

No need to use StringBuilder. Storing word itself at leaf node is enough.

No need to use HashSet to de-duplicate. Use "one time search" trie.

Some optimizations:

1. 59ms: Use search and startsWith in Trie class like this popular solution.
2. 33ms: Remove Trie class which unnecessarily starts from root in every dfs call.
3. 30ms: Use w.toCharArray() instead of w.charAt(i).
4. 22ms: Use StringBuilder instead of c1 + c2 + c3.
5. 20ms: Remove StringBuilder completely by storing word instead of boolean in TrieNode.
6. 20ms: Remove visited[m][n] completely by modifying board[i][j] = '#' directly.
7. 18ms: check validity, e.g., if(i > 0) dfs(...), before going to the next dfs.
8. 17ms: De-duplicate c - a with one variable i.
9. 15ms: Remove HashSet completely. dietpepsi's idea is awesome.

I'll list my solution and optimized solution on my note.

Listing 265: My Solution

```
class TrieNode {
    TrieNode[] next = new TrieNode[26];
    String word;
5 }

public class Solution {

    public List<String> findWords(char[][] board, String[] words) {
10         if (board.length == 0 || board[0].length == 0){
            return new ArrayList<String>();
        }

        List<String> res = new ArrayList<>();
        TrieNode root = buildTrie(words);
15         for(int i = 0; i < board.length; i++) {
            for(int j = 0; j < board[0].length; j++) {
                wordSearch(board, i, j, root, res);
            }
20         }
        return res;
    }

    private void wordSearch(char[][] board, int curX, int curY,
25     TrieNode curTrie, List<String> res){
        if (curX < 0 || curY < 0
            || curX >= board.length || curY >= board[0].length
            || board[curX][curY] == '#'){
            return;
30         }

        char curC = board[curX][curY];
```

```

    curTrie = curTrie.next[curC - 'a'];
    if (curTrie == null){
35         return;
    }
    if (curTrie.word != null){
        res.add(curTrie.word);
        curTrie.word = null;
40         // cannot return here! may have a word whose prefix is also a word
        // return;
    }
    board[curX][curY] = '#';
    wordSearch(board, curX + 1, curY, curTrie, res);
45    wordSearch(board, curX - 1, curY, curTrie, res);
    wordSearch(board, curX, curY + 1, curTrie, res);
    wordSearch(board, curX, curY - 1, curTrie, res);

    board[curX][curY] = curC;
50 }

public TrieNode buildTrie(String[] words) {
    TrieNode root = new TrieNode();
    for(String w : words) {
55         TrieNode p = root;
        for(char c : w.toCharArray()) {
            int i = c - 'a';
            if(p.next[i] == null) p.next[i] = new TrieNode();
            p = p.next[i];
60         }
        p.word = w;
    }
    return root;
}
65 }

```

And a much better solution.

Listing 266: Better Solution

```

public List<String> findWords(char[][] board, String[] words) {
    List<String> res = new ArrayList<>();
    TrieNode root = buildTrie(words);
    for(int i = 0; i < board.length; i++) {
5         for(int j = 0; j < board[0].length; j++) {
            dfs(board, i, j, root, res);
        }
    }
    return res;
10 }

public void dfs(char[][] board, int i, int j, TrieNode p, List<String> res) {
    char c = board[i][j];
    if(c == '#' || p.next[c - 'a'] == null) return;
15    p = p.next[c - 'a'];
    if(p.word != null) { // found one

```

```
        res.add(p.word);
        p.word = null;    // de-duplicate
    }

    board[i][j] = '#';
    if(i > 0) dfs(board, i - 1, j, p, res);
    if(j > 0) dfs(board, i, j - 1, p, res);
    if(i < board.length - 1) dfs(board, i + 1, j, p, res);
    if(j < board[0].length - 1) dfs(board, i, j + 1, p, res);
    board[i][j] = c;
}

public TrieNode buildTrie(String[] words) {
    TrieNode root = new TrieNode();
    for(String w : words) {
        TrieNode p = root;
        for(char c : w.toCharArray()) {
            int i = c - 'a';
            if(p.next[i] == null) p.next[i] = new TrieNode();
            p = p.next[i];
        }
        p.word = w;
    }
    return root;
}

class TrieNode {
    TrieNode[] next = new TrieNode[26];
    String word;
}
```

(BTW. I solved this problem using his TrieNode...)

Problem 207

139 Word Break

Given a string `s` and a dictionary of words `dict`,
determine if `s` can be segmented into a space-separated sequence
of one or more dictionary words.

For example, given
`s = "leetcode"`,
`dict = ["leet", "code"]`.

Return `true` because `"leetcode"` can be segmented as `"leet code"`.

<https://leetcode.com/problems/word-break/>

Solution

A DP problem, not so difficult if you know the key idea of solution.

The state transition equation is like $dp[i+word.length()] = dp[i]$, where `s.substring(i, i+word.length()).equals(word)`.

Listing 267: Solution

```
public class Solution {
    public boolean wordBreak(String s, Set<String> wordDict) {
        boolean dp[] = new boolean[s.length() + 1];
        dp[0] = true;
        for (int i = 0 ; i < s.length() ; i++){
            if (dp[i]){
                for (String word : wordDict){
                    if (i + word.length() <= s.length()){
                        String compare = s.substring(i, i + word.length());
                        if (compare.equals(word)){
                            dp[i + word.length()] = true;
                        }
                    }
                }
            }
        }

        return dp[s.length()];
    }
}
```

Problem 208

140 Word Break II

Given a string *s* and a dictionary of words *dict*,
add spaces in *s* to construct a sentence where
each word is a valid dictionary word.

Return all such possible sentences.

For example, given
s = "catsanddog",
dict = ["cat", "cats", "and", "sand", "dog"].

A solution is ["cats and dog", "cat sand dog"].

<https://leetcode.com/problems/word-break-ii/>

Solution

First have a glance at Problem “Word Break”, on how to finish the split process.

For this problem, we should also maintain the dp array, but this time not boolean, but String, a list of String to be more specific. *dp[i]* is the list of words in dictionary ends at index *i* with similar train of thought as “Word Break”;

Finally, use dfs to combine all possibilities.

Listing 268: Solution

```
public class Solution {
    List<String> result = new ArrayList<String>();
    private void findWordCombinations(List<String> curResult, int endIdx,
        String s, List<String>[] wordSnippets){
5         if (endIdx > s.length()){
            return;
        }
        if (endIdx == 0){
            String combine = "";
10         for (int wordIdx = curResult.size() - 1 ;
                wordIdx >= 0 ; wordIdx --){
            combine += curResult.get(wordIdx);
            if (wordIdx != 0){
                combine += " ";
15            }
        }
        result.add(combine);
        return;
    }
20    if (wordSnippets[endIdx] == null){
        return;
    }
    for (String word : wordSnippets[endIdx]){
        curResult.add(word);
25        findWordCombinations(curResult, endIdx - word.length()
            , s, wordSnippets);
        curResult.remove(curResult.size() - 1);
    }
}
```

```
    }  
}  
30 public List<String> wordBreak(String s, Set<String> wordDict) {  
    List<String>[] wordSnippets = new ArrayList[s.length() + 1];  
    wordSnippets[0] = new ArrayList<String>();  
    for (int i = 0 ; i < s.length() ; i++){  
        if (wordSnippets[i] != null){  
35         // Should have a precessor  
            for (String word : wordDict){  
                int wordLen = word.length();  
                if (wordLen + i <= s.length()){  
                    String part = s.substring(i, i + wordLen);  
40                    if (part.equals(word)){  
                        if (i + wordLen > s.length()){  
                            continue;  
                        }  
                        if (wordSnippets[i + wordLen] == null){  
45                            wordSnippets[i + wordLen]  
                                = new ArrayList<String>();  
                        }  
                        wordSnippets[i + wordLen].add(part);  
50                    }  
                }  
            }  
        }  
    }  
    findWordCombinations(new ArrayList<String>(),  
55        s.length() , s, wordSnippets);  
    return result;  
}
```

Problem 209

127 Word Ladder

Given two words (beginWord and endWord),
and a dictionary's word list,
find the length of shortest transformation sequence
from beginWord to endWord, such that:

Only one letter can be changed at a time
Each intermediate word must exist in the word list
For example,

Given:
beginWord = "hit"
endWord = "cog"
wordList = ["hot","dot","dog","lot","log"]
As one shortest transformation is
"hit" -> "hot" -> "dot" -> "dog" -> "cog",
return its length 5.

Note:
Return 0 if there is no such transformation sequence.
All words have the same length.
All words contain only lowercase alphabetic characters.

<https://leetcode.com/problems/word-ladder/>

Solution

First we may realize that this is a graph(undirected) problem using BFS. Two nodes are neighbors when they have an distance of 1(Only one letter differs). So we may write code as below:

Listing 269: TLE Solution

```
public class Solution {  
    public int ladderLength(String start, String end, Set<String> dict) {  
  
        //create graph  
        5      HashMap<String, ArrayList<String>> graph =  
            new HashMap<String, ArrayList<String>>();  
  
        graph.put(start, new ArrayList<String>());  
        graph.put(end, new ArrayList<String>());  
        10     for(String d : dict) {  
            graph.put(d, new ArrayList<String>());  
        }  
        for(String s : graph.keySet()) {  
            ArrayList<String> list = graph.get(s);  
            15     for(String t : graph.keySet()) {  
                if(getDiff(s,t) == 1) {  
                    list.add(t);  
                }  
            }  
        }  
        20     }  
        // use BFS to traverse the node in the graph, we begin with "start"
```



```

    int step = 0;
    HashSet<String> visited = new HashSet<String>();
    ArrayList<String> firstLevel = new ArrayList<String>(graph.get(start));
25
    while (firstLevel.size() != 0) {
        step++;
        ArrayList<String> nextLevel = new ArrayList<String>();
        for (String s : firstLevel) {
30            if (s.equals(end)) return step + 1;
            visited.add(s);
            nextLevel.addAll(graph.get(s));
        }
        firstLevel.clear();
35        for (String t : nextLevel) {
            if (!visited.contains(t)) {
                firstLevel.add(t);
            }
        }
40        nextLevel.clear();
    }
    return 0;
}

45 public int getDiff(String w1, String w2) {
    int count = 0;
    for (int i = 0; i < w1.length(); i++) {
        if (w1.charAt(i) != w2.charAt(i)) {
50            count++;
        }
    }
    return count;
}
}

```

A simplified method is to avoid process that generates a graph. Still a BFS framework, but this time a different method. Substitute each letter of a word from 'a' to 'z' to see if generated word exists in dictionary, which is in nature same as former solution, but has a better performance when the count of word is large (since few of them are of 1 distance, we waste our time comparing them) and word is short.

However, a one-pass solution will also get TLE error. So we need to optimize it further. Do it with 2 pass. Go from both sourceSet and targetSet, set the shorter one to 'sourceSet', guaranteeing that we are always searching from a shorter set, which will accelerate our execution.

Listing 270: Solution

```

public class Solution {
    public int ladderLength(String start, String end, Set<String> dict) {
        Set<String> sourceSet = new HashSet<String>();
        Set<String> targetSet = new HashSet<String>();
5        sourceSet.add(start);
        targetSet.add(end);
        return findLadder(sourceSet, targetSet, dict, 1);
    }
}

```

```
10     private int findLadder(Set<String> sourceSet ,
        Set<String> targetSet , Set<String> dict , int depth){
        if (sourceSet.size() == 0){
            return 0;
        }
15     if (targetSet.size() < sourceSet.size()){
        return findLadder(targetSet , sourceSet , dict , depth);
    }
    for (String s : sourceSet){
        dict.remove(s);
20    }
    for (String s : targetSet){
        dict.remove(s);
    }
    Set<String> newSet = new HashSet<String>();
25    for (String source : sourceSet){
        for (int strIdx = 0 ; strIdx < source.length() ; strIdx ++){
            // If I put
            // char[] sourceCharArr = source.toCharArray();
            // here, I will get TLE?!
30        for (char subChar = 'a' ; subChar <= 'z' ; subChar ++){
            if (source.charAt(strIdx) == subChar){
                continue;
            }
            // This statement...
35        char[] sourceCharArr = source.toCharArray();
            sourceCharArr[strIdx] = subChar;
            String newSource = new String(sourceCharArr);
            for (String target : targetSet){
                if (target.equals(newSource)){
40                    return depth + 1;
                }
            }
            if (dict.contains(newSource)){
                newSet.add(newSource);
45            }
        }
    }
50    return findLadder(newSet , targetSet , dict , depth + 1);
}
```

Problem 210

126 Word Ladder II

Given two words (beginWord and endWord),
and a dictionary's word list,
find all shortest transformation sequence(s)
from beginWord to endWord, such that:

Only one letter can be changed at a time
Each intermediate word must exist in the word list
For example,

```
Given:
beginWord = "hit"
endWord = "cog"
wordList = ["hot", "dot", "dog", "lot", "log"]
Return
[
  ["hit", "hot", "dot", "dog", "cog"],
  ["hit", "hot", "lot", "log", "cog"]
]
```

Note:

All words have the same length.
All words contain only lowercase alphabetic characters.

<https://leetcode.com/problems/word-ladder-ii/>

Solution

Similar to Problem “Word Ladder”, still a two-pass method, set beginWords the shorter set which will reduce the time burden.

Building the result path is another tricky part(which is maybe more tricky). Generating a HashMap with key of source String and value of strings it can generated by just changing one letter of the word is able to solve the problem. Just DFS the Map, we will get the final result.

Listing 271: Solution

```
public class Solution {
    List<List<String>> result = new ArrayList<List<String>>();
    public List<List<String>> findLadders(String beginWord, String endWord,
                                         Set<String> wordList) {
5      Map<String, List<String>> succs = new HashMap<String, List<String>>();
      Set<String> beginWords = new HashSet<String>();
      Set<String> endWords = new HashSet<String>();
      beginWords.add(beginWord);
      endWords.add(endWord);
10     bfsLadders(beginWords, endWords, wordList, succs, false);

      List<String> pathStarts = new ArrayList<String>();
      pathStarts.add(beginWord);
      buildPath(beginWord, endWord, succs, new ArrayList<String>(pathStarts));
15     return result;
    }
    private void bfsLadders(Set<String> beginWords, Set<String> endWords,
```

```

        Set<String> dict , Map<String , List<String>> succs , boolean reversed){

20     if (beginWords.isEmpty()){
        return;
    }
    if (beginWords.size() > endWords.size()){
        bfsLadders(endWords, beginWords, dict , succs , !reversed);
25     return;
    }
    for (String beginWord : beginWords){
        dict.remove(beginWord);
    }
30     for (String endWord : endWords){
        dict.remove(endWord);
    }
    Set<String> generated = new HashSet<String>();
    boolean found = false;
35     for (String beginWord : beginWords){
        for (int strIdx = 0 ; strIdx < beginWord.length() ; strIdx++){
            char[] wordCharArr = beginWord.toCharArray();
            for (char subChar = 'a' ; subChar <= 'z' ; subChar++){
                // Don't add this!
                // if (wordCharArr[strIdx] == subChar){
40                 //     continue;
                // }
                wordCharArr[strIdx] = subChar;
                String newStr = new String(wordCharArr);
                String key = reversed ? newStr : beginWord;
                String value = reversed ? beginWord : newStr;
                List<String> list = succs.containsKey(key) ?
                    succs.get(key) : new ArrayList<String>();

50                 if (endWords.contains(newStr)){
                    found = true;

                    list.add(value);
                    succs.put(key, list);

55                 }
                else{
                    if (!dict.contains(newStr)){
                        continue;
                    }
60                 else{
                    if (!found){
                        generated.add(newStr);

                        list.add(value);
                        succs.put(key, list);
65                 }
                }
            }
        }
    }
70 }

```

```
    }

    if (found){
        return;
75    }
    bfsLadders(generated, endWords, dict, succs, reversed);
}

private void buildPath(String beginWord, String endWord,
80    Map<String, List<String>> succs, List<String> curList){
    if (beginWord.equals(endWord)){
        result.add(new ArrayList(curList));
        return;
    }

85    if (!succs.containsKey(beginWord)){
        return;
    }

90    for (String next : succs.get(beginWord)){
        curList.add(next);
        buildPath(next, endWord, succs, curList);
        curList.remove(curList.size() - 1);
    }

95    }
}
```

Problem 211

5 Longest Palindromic Substring

Given a string S,
 find the longest palindromic substring in S.
 You may assume that the maximum length of S is 1000,
 and there exists one unique longest palindromic substring.

<https://leetcode.com/problems/longest-palindromic-substring/>

Solution - DP

First solution is DP solution with $O(n^2)$, it can output a result through 'Run'(within an acceptable time limit, I guess), but will get TLE when submit.

Mention the variables in loop - len and i rather than i and j since we should get dp value of shorter string first to get longer ones.

Listing 272: Solution 1

```

public class Solution {
    public String longestPalindrome(String s) {
        int len = s.length();
        if (len == 0 || len == 1){
5           return s;
        }
        int [][] dp = new int[len][len];
        int max = 1;
        int startIdx = -1;
10        for (int i = 0 ; i < len ; i++){
            dp[i][i] = 1;
            if (max == 1){
                startIdx = i;
            }
15            if (i > 0 && s.charAt(i) == s.charAt(i - 1)){
                dp[i - 1][i] = 2;
                max = 2;
                startIdx = i - 1;
            }
20        }
        for (int subLen = 3; subLen <= len ; subLen++){
            for (int i = 0 ; i < len - subLen + 1 ; i++){
                int j = i + subLen - 1;
                if (s.charAt(i) == s.charAt(j) && dp[i + 1][j - 1] != 0){
25                    dp[i][j] = Math.max(dp[i][j] , dp[i + 1][j - 1] + 2);
                    if (max < dp[i][j]){
                        max = dp[i][j];
                        startIdx = i;
                    }
30                }
            }
        }
        return s.substring(startIdx , startIdx + max);
    }
35 }

```

Solution - Expanding

Solution is an $O(n^2)$ runtime solution. The train of thought is interesting and reversing. Instead of finding the longest, the solution start from the middle of palindrome and form a palindrome from the middle. Note that there are two ways of expanding the middle element(s), one from single element, while another from two adjacent elements.

Listing 273: Solution 2

```

public class Solution {
    public String longestPalindrome(String s) {
        if (s.length() == 0 || s.length() == 1){
            return s;
5        }
        int max = 0;
        int startIdx = -1;
        for (int i = 0 ; i < s.length() ; i++){
            int singleMax = expandToSides(s, i, i);
10            if (max < singleMax){
                max = singleMax;
                startIdx = i - max / 2;
            }
            if (i > 0){
15                int doubleMax = expandToSides(s, i - 1, i);
                if (max < doubleMax){
                    max = doubleMax;
                    startIdx = i - max / 2;
                }
20            }
        }
        return s.substring(startIdx, startIdx + max);
    }
    public int expandToSides(String s, int left, int right){
25        int start = left, end = right;
        int max = 0;
        while (start >= 0 && end < s.length()){
            if (s.charAt(start) == s.charAt(end)){
                max = Math.max(end - start + 1, max);
30            }
            else{
                break;
            }
            start --;
            end ++;
35        }
        return max;
    }
}

```

Solution - Manacher

Such a neat, elegant, excellent algorithm that I still cannot totally understand it...

Though there is a blog for Manacher in Chinese, I think the English version in Leetcode articles is much better.(Maybe with the example?) I'll post the screenshots and code here for future use.

An O(N) Solution (Manacher's Algorithm):

First, we transform the input string, S, to another string T by inserting a special character '#' in between letters. The reason for doing so will be immediately clear to you soon.

For example: S = "abaaba", T = "#a#b#a#a#b#a#".

To find the longest palindromic substring, we need to expand around each T_i such that $T_{i-d} \dots T_{i+d}$ forms a palindrome. You should immediately see that d is the length of the palindrome itself centered at T_i .

We store intermediate result in an array P, where $P[i]$ equals to the length of the palindrome centers at T_i . The longest palindromic substring would then be the maximum element in P.

Using the above example, we populate P as below (from left to right):

```
T = # a # b # a # a # b # a #
P = 0 1 0 3 0 1 6 1 0 3 0 1 0
```

Looking at P, we immediately see that the longest palindrome is "abaaba", as indicated by $P_6 = 6$.

Did you notice by inserting special characters (#) in between letters, both palindromes of odd and even lengths are handled gracefully? (Please note: This is to demonstrate the idea more easily and is not necessarily needed to code the algorithm.)

Now, imagine that you draw an imaginary vertical line at the center of the palindrome "abaaba". Did you notice the numbers in P are symmetric around this center? That's not only it, try another palindrome "aba", the numbers also reflect similar symmetric property. Is this a coincidence? The answer is yes and no. This is only true subjected to a condition, but anyway, we have great progress, since we can eliminate recomputing part of $P[i]$'s.

Let us move on to a slightly more sophisticated example with more some overlapping palindromes, where S = "babcbabcbaccba".

```
index      2      9      11     13      20
var        L      i'     C      i      R

T = # b # a # b # c # b # a # b # c # b # a # c # c # b # a #
P = 0 1 0 3 0 1 0 7 0 1 0 9 0 ?
```

Above image shows T transformed from S = "babcbabcbaccba". Assumed that you reached a state where table P is partially completed. The solid vertical line indicates the center (C) of the palindrome "abcbabcb". The two dotted vertical line indicate its left (L) and right (R) edges respectively. You are at index i and its mirrored index around C is i'. How would you calculate $P[i]$ efficiently?

Assume that we have arrived at index $i = 13$, and we need to calculate $P[13]$ (indicated by the question mark ?). We first look at its mirrored index i' around the palindrome's center C, which is index $i' = 9$.

```
index      2      9      11     13      20
var        L      i'     C      i      R

T = # b # a # b # c # b # a # b # c # b # a # c # c # b # a #
P = 0 1 0 3 0 1 0 7 0 1 0 9 0 ?
```

The two green solid lines above indicate the covered region by the two palindromes centered at i and i'. We look at the mirrored index of i around C, which is index i'. $P[i'] = P[9] = 1$. It is clear that $P[i]$ must also be 1, due to the symmetric property of a palindrome around its center.

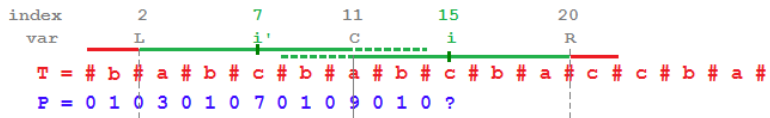
As you can see above, it is very obvious that $P[i] = P[i'] = 1$, which must be true due to the symmetric property around a palindrome's center. In fact, all three elements after C follow the symmetric property (that is, $P[12] = P[10] = 0$, $P[13] = P[9] = 1$, $P[14] = P[8] = 0$).

```
index      2      7      11     15      20
var        L      i'     C      i      R

T = # b # a # b # c # b # a # b # c # b # a # c # c # b # a #
P = 0 1 0 3 0 1 0 7 0 1 0 9 0 1 0 ?
```

Now we are at index $i = 15$, and its mirrored index around C is $i' = 7$. Is $P[15] = P[7] = 7$?

Now we are at index $i = 15$. What's the value of $P[i]$? If we follow the symmetric property, the value of $P[i]$ should be the same as $P[i'] = 7$. But this is wrong. If we expand around the center at T_{15} , it forms the palindrome "a#b#c#b#a", which is actually shorter than what is indicated by its symmetric counterpart. Why?



Colored lines are overlaid around the center at index i and i' . Solid green lines show the region that must match for both sides due to symmetric property around C . Solid red lines show the region that might not match for both sides. Dotted green lines show the region that crosses over the center.

It is clear that the two substrings in the region indicated by the two solid green lines must match exactly. Areas across the center (indicated by dotted green lines) must also be symmetric. Notice carefully that $P[i']$ is 7 and it expands all the way across the left edge (L) of the palindrome (indicated by the solid red lines), which does not fall under the symmetric property of the palindrome anymore. All we know is $P[i] \geq 5$, and to find the real value of $P[i]$ we have to do character matching by expanding past the right edge (R). In this case, since $P[21] \neq P[1]$, we conclude that $P[i] = 5$.

Let's summarize the key part of this algorithm as below:

```
if  $P[i'] \leq R - i$ ,
then  $P[i] \leftarrow P[i']$ 
else  $P[i] \geq P[i']$ . (Which we have to expand past the right edge ( $R$ ) to find  $P[i]$ ).
```

See how elegant it is? If you are able to grasp the above summary fully, you already obtained the essence of this algorithm, which is also the hardest part.

The final part is to determine when should we move the position of C together with R to the right, which is easy:

```
If the palindrome centered at  $i$  does expand past  $R$ , we update  $C$  to  $i$ , (the center of this new palindrome), and extend  $R$  to the new palindrome's right edge.
```

In each step, there are two possibilities. If $P[i] \leq R - i$, we set $P[i]$ to $P[i']$ which takes exactly one step. Otherwise we attempt to change the palindrome's center to i by expanding it starting at the right edge, R . Extending R (the inner while loop) takes at most a total of N steps, and positioning and testing each centers take a total of N steps too. Therefore, this algorithm guarantees to finish in at most $2 \cdot N$ steps, giving a linear time solution.

Listing 274: Solution 3

```
public class Solution {
    public String longestPalindrome(String s) {
        if (s.length() <= 2){
            if (s.length() != 2 || s.charAt(1) == s.charAt(0)){
                return s;
            }
            else{
                return s.substring(0, 1);
            }
        }
        int maxlen = 1;
        String curStr = preprocess(s);
        int[] lens = new int[2 * curStr.length() + 1];
        int rightBound = 0;
        int mid = 1;
        for (int idx = 1 ; idx < lens.length ; idx++){
            lens[idx] = (rightBound > idx)
                ? Math.min(lens[2 * mid - idx], rightBound - idx) : 0;
```

```
20         while (idx + lens[idx] < curStr.length() && idx - lens[idx] >= 0
                && curStr.charAt(idx + lens[idx])
                == curStr.charAt(idx - lens[idx])){
                lens[idx] ++;
            }
            if (idx + lens[idx] > rightBound){
25                 rightBound = idx + lens[idx];
                mid = idx;
            }
        }
        int maxPos = 0;
30        for (int lenIdx = 0 ; lenIdx < lens.length ; lenIdx ++){
            int len = lens[lenIdx];
            if (maxLen < len - 1){
                maxLen = len - 1;
                maxPos = lenIdx;
35            }
        }
        String result = s.substring((maxPos - maxLen)/2,
                maxLen + (maxPos - maxLen) / 2);
        return result;
40    }

    private String preprocess(String curStr){
        StringBuilder sb = new StringBuilder();
        sb.append("#");
45        for (int idxCh = 0 ; idxCh < curStr.length() ; idxCh ++){
            sb.append(curStr.charAt(idxCh));
            sb.append("#");
        }
        return sb.toString();
50    }
}
```

Problem 212

221 Maximal Square

Given a 2D binary matrix filled with 0's and 1's, find the largest square containing all 1's and return its area.

For example, given the following matrix:

```
1 0 1 0 0
1 0 1 1 1
1 1 1 1 1
1 0 0 1 0
Return 4.
```

<https://leetcode.com/problems/maximal-square/>

Solution

A simple DP problem. Take care of the input: not `int[][]` matrix, but `char[][]` matrix.

Listing 275: Solution

```
public class Solution {
    public int maximalSquare(char [][] matrix) {
        int m = matrix.length;
        if (m == 0){
5           return 0;
        }
        int n = matrix[0].length;
        int dp [][] = new int[m][n];
        int maxSide = 0;
10        for (int i = 0 ; i < m ; i++){
            dp[i][0] = matrix[i][0] - '0';
        }
        for (int i = 0 ; i < n ; i++){
            dp[0][i] = matrix[0][i] - '0';
15        }
        for (int i = 1; i < m ; i++){
            for (int j = 1 ; j < n ; j++){
                if (matrix[i][j] == '1'){
20                     dp[i][j] = Math.min(dp[i - 1][j - 1],
                        Math.min(dp[i - 1][j], dp[i][j - 1])) + 1;
                }
            }
        }
        for (int i = 0 ; i < m ; i++){
25            for (int j = 0 ; j < n ; j++){
                maxSide = Math.max(maxSide, dp[i][j]);
            }
        }
        return maxSide * maxSide;
30    }
}
```

Problem 213

84 Largest Rectangle in Histogram

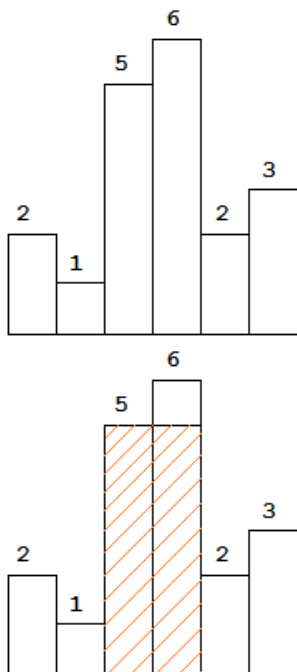
Given n non-negative integers representing the histogram's bar height where the width of each bar is 1, find the area of largest rectangle in the histogram.

Above is a histogram where width of each bar is 1, given height = [2,1,5,6,2,3].

The largest rectangle is shown in the shaded area, which has area = 10 unit.

For example,
Given heights = [2,1,5,6,2,3],
return 10.

<https://leetcode.com/problems/largest-rectangle-in-histogram/>



Solution

A simple train of thought is like Problem “Longest Palindromic Substring”, by starting from middle and expand to both sides. However, when there is a set of big data with similar values, we have to compute nearly all values for n^2 , which will of course lead to TLE.

Listing 276: TLE solution

```
public class Solution {  
    public int largestRectangleArea(int[] heights) {  
        int max = 0;  
        for (int i = 0 ; i < heights.length ; i ++){
```

```

5      int height = heights[i];
      int left = i , right = i;
      while (left > 0 && heights[left - 1] >= height){
          left --;
      }
10     while (right < heights.length - 1
          && heights[right + 1] >= height){
          right ++;
      }
      int width = right - left + 1;
15
      max = Math.max(width * height , max);
  }

  return max;
20 }
}

```

There is another brilliant train of thought of dealing with this problem.

Use a ‘ascending stack’ to solve the problem. Iterate along array, when heights[idx] is larger than top element of stack, push it; otherwise pop all elements larger than current element, and since popped elements are in descending order, we can just use (height of popped element * width++) to get a potential maximum rectangle. The width of rectangle is another difficulty, see figure.

这道题brute force的方法很直接，就是对于每一个窗口，找到其中最低的高度，然后求面积，去其中最大的矩形面积。总共有 n^2 个窗口，找最低高度是 $O(n)$ 的操作，所以复杂度是 $O(n^3)$ 。

接下来我们讨论一种比较容易理解的思路，就是从每一个bar往两边走，以自己的高度为标准，直到两边低于自己的高度为止，然后用自己的高度乘以两边走的宽度得到矩阵面积。因为我们对于任意一个bar都计算了以自己为目标高度的最大矩阵，所以最好的结果一定会被取到。每次往两边走的复杂度是 $O(n)$ ，总共有 n 个bar，所以时间复杂度是 $O(n^2)$ 。因为代码比较简单，这里就不列出来了哈。

最后我们谈谈最优的解法，思路跟Longest Valid Parentheses比较类似，我们要维护一个栈，这个栈从低向上的高度是依次递增的，如果遇到当前bar高度比栈顶元素低，那么就出栈直到满足条件，过程中检测前面满足条件的矩阵。关键问题就在于在出栈时如何确定当前出栈元素的所对应的高度的最大范围是多少。答案跟Longest Valid Parentheses中括号的范围相似，就是如果栈已经为空，说明到目前为止所有元素（当前下标元素除外）都比出栈元素高度要大（否则栈中肯定还有元素），所以矩阵面积就是高度乘以当前下标。如果栈不为空，那么就是从当前栈顶元素的下一个到当前下标的元素之前都比出栈元素高度大（因为栈顶元素第一个比当前出栈元素小的）。具体的实现代码如下：

Listing 277: Solution

```

public class Solution {
    public int largestRectangleArea(int[] heights) {
        Stack<Integer> ascendingStack = new Stack<Integer>();
5      if (heights.length == 0){
          return 0;
        }
        if (heights.length == 1){
          return heights[0];
10     }

        int max = 0;
        for (int i = 0 ; i < heights.length ; i++){
            if (ascendingStack.isEmpty()
15             || heights[i] > heights[ascendingStack.peek()]){
                ascendingStack.push(i);
                continue;
            }
        }
    }
}

```

```
    }
    else{
20      while (!ascendingStack.isEmpty()
        && heights[ascendingStack.peek()] >= heights[i]){
        int heightIdx = ascendingStack.pop();
        int height = heights[heightIdx];
        int width = (ascendingStack.isEmpty())
25        ? i : i - ascendingStack.peek() - 1;
        max = Math.max(width * height , max);
      }
      ascendingStack.push(i);
    }
30  }

  while (!ascendingStack.isEmpty()){
    int heightIdx = ascendingStack.pop();
    int height = heights[heightIdx];
35    int width = (ascendingStack.isEmpty())
      ? heights.length: heights.length - ascendingStack.peek() - 1;
    max = Math.max(width * height , max);
  }

40  return max;
}
}
```

Problem 214

85 Maximal Rectangle

Given a 2D binary matrix filled with 0's and 1's,
find the largest rectangle containing all ones and return its area.

<https://leetcode.com/problems/maximal-rectangle/>

Solution

Inspired by Problem “Largest Rectangle in Histogram”, the problem can be transformed to a problem doing m (m is the total number of rows) times with algorithm in previous problem. (Think of why)

Listing 278: Solution

```
public class Solution {
    public int maximalRectangle(char[][] matrix) {
        Stack<Integer> stack = new Stack<Integer>();
        int m = matrix.length;
        if (m == 0){
            return 0;
        }
        int n = matrix[0].length;

        int[][] nums = new int[m][n];

        for (int i = m - 1 ; i >= 0 ; i --){
            for (int j = 0 ; j < n ; j ++){
                if (i == m - 1){
                    nums[i][j] = matrix[i][j] - '0';
                }
                else{
                    if (matrix[i][j] == '1'){
                        nums[i][j] = nums[i + 1][j] + (matrix[i][j] - '0');
                    }
                    else{
                        nums[i][j] = 0;
                    }
                }
            }
        }

        int max = 0;

        for (int i = m - 1 ; i >= 0 ; i --){
            stack = new Stack<Integer>();
            for (int j = 0 ; j <= n ; j ++){
                if (j != n && (stack.isEmpty()
                    || nums[i][stack.peek()] < nums[i][j])){
                    stack.push(j);
                    continue;
                }
                else{
                    while (!stack.isEmpty()
                        && (j == n || nums[i][stack.peek()] >= nums[i][j])){
```

```

    int idx = stack.pop();
    int height = nums[i][idx];
    int width = (stack.isEmpty())? j: j - stack.peek() - 1;
    max = Math.max(width * height, max);
45     }

    if (j < n){
        stack.push(j);
50     }
    }
    }
    return max;
55 }
}
```


Problem 215

43 Multiply Strings

Given two numbers represented as strings,
return multiplication of the numbers as a string.

Note: The numbers can be arbitrarily large and are non-negative.

<https://leetcode.com/problems/multiply-strings/>

Solution

I unexpectedly solve the problem in just one attempt!

However, I missed the given condition that numbers are non-negative, which means I forgot to consider the negative numbers at first sight (but I do consider the 0 result).

Listing 279: Solution

```
public class Solution {
    public String multiply(String num1, String num2) {
        if (num1.length() < num2.length()) {
            return multiply(num2, num1);
        }

        if (num1.equals("0") || num2.equals("0")) {
            return "0";
        }

        int[] numArr1 = new int[num1.length()];
        int[] numArr2 = new int[num2.length()];
        for (int idx = 0 ; idx < num1.length() ; idx++) {
            numArr1[idx] = num1.charAt(idx) - '0';
        }
        for (int idx = 0 ; idx < num2.length() ; idx++) {
            numArr2[idx] = num2.charAt(idx) - '0';
        }

        int[] result = new int[num1.length() + num2.length()];

        for (int num2Idx = num2.length() - 1 ; num2Idx >= 0 ; num2Idx--) {
            int carry = 0;
            int resultCarry = 0;
            int endIdx = num2.length() - 1 - num2Idx;
            for (int num1Idx = num1.length() - 1 ; num1Idx >= 0 ; num1Idx--) {
                int multiplyResult = numArr1[num1Idx] * numArr2[num2Idx];
                int digit = (multiplyResult + carry) % 10;
                carry = (multiplyResult + carry) / 10;
                int addResult = digit + result[endIdx];
                result[endIdx] = (addResult + resultCarry) % 10;
                resultCarry = (addResult + resultCarry) / 10;
                endIdx++;
            }
            if (carry != 0 || resultCarry != 0) {
                result[endIdx] += resultCarry + carry;
            }
        }
    }
}
```

```

        if (result[endIdx] >= 10){
            result[endIdx + 1] = result[endIdx];
            result[endIdx] %= 10;
40        }
    }
    String finalResult = "";
    boolean firstDigit = false;
45    for (int idx = result.length - 1; idx >= 0 ; idx --){
        if (!firstDigit && result[idx] != 0){
            firstDigit = true;
        }
        if (!firstDigit && result[idx] == 0){
50            continue;
        }
        finalResult += String.valueOf(result[idx]);
    }
55    return finalResult;
}
}
```

Problem 216

29 Divide Two Integers

Divide two integers without using multiplication, division and mod operator.

If it is overflow, return MAX_INT.

<https://leetcode.com/problems/divide-two-integers/>

Solution

At first sight I thought this is a problem similar to “Multiply Strings”, however, this is indeed a Math and Binary Search problem.

Without division and multiplication, remaining operators are +, -, and bit operators. The most straightforward way is to reduce divisor from dividend until remaining number is less than divisor, which will get TLE.

Have you noticed that, $\ll 1$ is the same as $* 2$? That’s the key insight to this problem. See code in listing.(There are still many edge cases!)

Something to take care. See in the listing, use ‘divisor \leq (dividend $\gg 1$)’ instead of ‘dividend < divisor’, the latter will take 903ms to calculate Integer.MAX_VALUE / 2 while the former takes 0ms...

Listing 280: Solution

```
public class Solution {
    public int divide(int dividend, int divisor) {
        if (divisor == 0){
            return Integer.MAX_VALUE;
        }

        int result = 0;

        if(dividend==Integer.MIN_VALUE){
            result = 1;
            dividend += Math.abs(divisor);
        }
        if(divisor==Integer.MIN_VALUE)
            return result;

        // boolean negative = ((dividend^divisor)>>>31==1)?true:false;
        boolean negative = (dividend > 0 && divisor < 0 ||
            dividend < 0 && divisor > 0);
        dividend = Math.abs(dividend);
        divisor = Math.abs(divisor);

        if (dividend == 0 || dividend < divisor){
            return negative ? -result : result;
        }

        int count = 0;
        // while (dividend < divisor){
        while (divisor <= (dividend >> 1)){
```

```
30         divisor <<= 1;
           count ++;
       }
       while (count >= 0){
           while (dividend >= divisor){
35               dividend -= divisor;
               result += (1 << count);
           }
           divisor >>= 1;
           count --;
40       }
       if (result == Integer.MIN_VALUE && !negative){
           return Integer.MAX_VALUE;
       }
45       return negative ? -result : result;
   }
}
```

Problem 217

36 Valid Sudoku

Determine if a Sudoku is valid, according to:
Sudoku Puzzles - The Rules.

The Sudoku board could be partially filled,
where empty cells are filled with the character '.'.

A partially filled sudoku which is valid.

Note:

A valid Sudoku board (partially filled)
is not necessarily solvable.
Only the filled cells need to be validated.

<https://leetcode.com/problems/valid-sudoku/>

5	3			7				
6			1	9	5			
	9	8					6	
8				6				3
4			8		3			1
7				2				6
	6					2	8	
			4	1	9			5
				8			7	9

Solution

It's an easy problem, use 3 arrays to validate value on rows, columns and grids(with 3 * 3 size) are used.
Description in Note makes the problem even easier.

Listing 281: Solution

```

public class Solution {
    int max_value = 9;
    public boolean isValidSudoku(char[][] board) {
        boolean[][] rowFilled = new boolean[board.length][max_value];
        boolean[][] colFilled = new boolean[board[0].length][max_value];
        boolean[][] gridFilled = new boolean[9][max_value];

        for (int i = 0; i < board.length; i++){
            for (int j = 0 ; j < board[0].length ; j++){
                if (board[i][j] == '.'){
                    continue;
                }
                int curVal = board[i][j] - '0' - 1;
                if (rowFilled[i][curVal] || colFilled[j][curVal]
                    || gridFilled[i / 3 * 3 + j / 3][curVal]){
                    return false;
                }
            }
        }
    }
}

```

```
                }  
                rowFilled[i][curVal] = true;  
                colFilled[j][curVal] = true;  
20         gridFilled[i / 3 * 3 + j / 3][curVal] = true;  
            }  
        }  
        return true;  
25    }  
}
```

Problem 218

37 Sudoku Solver

Write a program to solve a Sudoku puzzle by filling the empty cells.

Empty cells are indicated by the character '.'.

You may assume that there will be only one unique solution.

A sudoku puzzle...

...and its solution numbers marked in red.

<https://leetcode.com/problems/sudoku-solver/>

5	3		7					
6			1	9	5			
	9	8				6		
8			6					3
4			8	3				1
7			2					6
	6					2	8	
			4	1	9			5
			8			7	9	

5	3	4	6	7	8	9	1	2
6	7	2	1	9	5	3	4	8
1	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9	2	4	8	5	6
9	6	1	5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9

Solution

Just like "N-Queens", pre-processing will much accelerate the speed.

Listing 282: Solution

```

public class Solution {
    boolean[][] rowFilled = new boolean[9][9];
    boolean[][] colFilled = new boolean[9][9];
    boolean[][] gridFilled = new boolean[9][9];

    private boolean fill(char[][] board, int pos){
        int curX = pos / 9;
        int curY = pos % 9;
        if (board[curX][curY] == '.'){
            for (int ava = 0 ; ava < 9 ; ava++){
                if (rowFilled[curX][ava] || colFilled[curY][ava]
                    || gridFilled[curX / 3 * 3 + curY / 3][ava]){
                    continue;
                }
                else{
                    // have a try
                    rowFilled[curX][ava] = true;
                    colFilled[curY][ava] = true;
                    gridFilled[curX / 3 * 3 + curY / 3][ava] = true;
                    board[curX][curY] = (char)('0' + 1 + ava);
                    if (pos == 80){
                        return true;
                    }
                }
            }
        }
    }
}

```

```
25         if (!fill(board, pos + 1)){
           //recover
           board[curX][curY] = '.';
           rowFilled[curX][ava] = false;
           colFilled[curY][ava] = false;
30           gridFilled[curX / 3 * 3 + curY / 3][ava] = false;
           }
           else{
               return true;
           }
35       }
       }
       return false;
   }
   else{
40       if (pos == 80){
           return true;
       }

       return fill(board, pos + 1);
45   }
}

public void solveSudoku(char[][] board) {
   for (int i = 0 ; i < 9 ; i++){
       for (int j = 0 ; j < 9 ; j++){
50           if (board[i][j] == '.'){
               continue;
           }
           int curVal = board[i][j] - '0' - 1;
55           rowFilled[i][curVal] = true;
           colFilled[j][curVal] = true;
           gridFilled[i / 3 * 3 + j / 3][curVal] = true;
       }
   }
60   fill(board, 0);
}
}
```


Problem 219

34 Search for a Range

Given a sorted array of integers,
find the starting and ending position of a given target value.

Your algorithm's runtime complexity must be in the order of $O(\log n)$.

If the target is not found in the array, return `[-1, -1]`.

For example,
Given `[5, 7, 7, 8, 8, 10]` and target value `8`,
return `[3, 4]`.

<https://leetcode.com/problems/search-for-a-range/>

Solution

Binary search problem. The key to problem is to do binary search twice.

Listing 283: Solution

```
public class Solution {
    public int[] searchRange(int[] nums, int target) {
        int[] result = new int[2];
        int left = find(nums, target, false);
        int right = find(nums, target, true);
        result[0] = left;
        result[1] = right;
        return result;
    }

    private int find(int[] nums, int target, boolean forward){
        int start = 0, end = nums.length - 1;

        while (start <= end){
            int mid = start + (end - start) / 2;
            if (nums[mid] == target){
                if (forward){
                    if (mid == nums.length - 1 || nums[mid] < nums[mid + 1]){
                        return mid;
                    }
                    else{
                        start = mid + 1;
                    }
                }
                else{
                    if (mid == 0 || nums[mid - 1] < nums[mid]){
                        return mid;
                    }
                    else{
                        end = mid - 1;
                    }
                }
            }
            else if (nums[mid] < target){
```

```
35         start = mid + 1;
        }
        else {
            end = mid - 1;
        }
    }
    return -1;
}
```

Problem 220

209 Minimum Size Subarray Sum

Given an array of n positive integers and a positive integer s , find the minimal length of a subarray of which the sum $\geq s$. If there isn't one, return 0 instead.

For example, given the array $[2,3,1,2,4,3]$ and $s = 7$, the subarray $[4,3]$ has the minimal length under the problem constraint.

[click to show more practice.](#)

More practice:

If you have figured out the $O(n)$ solution, try coding another solution of which the time complexity is $O(n \log n)$.

<https://leetcode.com/problems/search-for-a-range/>

Solution 1

Use 2 pointers to maintain the sum and check if less than current minimum condition.
 $O(n)$ solution.

Listing 284: Solution 1

```
public class Solution {
    public int minSubArrayLen(int s, int[] nums) {
        if (nums.length == 0){
            return 0;
        }
        int low = 0 , high = 0;
        int sum = 0;
        int minLen = Integer.MAX_VALUE;
        while (high < nums.length){
            if (sum >= s){
                minLen = Math.min(high - low , minLen);
                sum -= nums[low];
                low ++;
            }
            else{
                sum += nums[high];
                high ++;
            }
        }
        boolean changed = false;
        while (low < nums.length && sum >= s){
            changed = true;
            sum -= nums[low];
            low ++;
            if (sum < s){
                break;
            }
        }
        if (changed){
            minLen = Math.min(minLen, high - low + 1);
        }
    }
}
```

```
    }  
    return minLen == Integer.MAX_VALUE ? 0 : minLen;  
}  
}
```

Solution 2

Second solution is to use binary search to find proper length to calculate the sum. See listing.

The $O(n \log n)$ solution.

Listing 285: Solution 2

```
public class Solution {  
    public int minSubArrayLen(int s, int[] nums) {  
        int min = Integer.MAX_VALUE;  
        int i = 1, j = nums.length;  
        while (i <= j){  
            int mid = i + (j - i) / 2;  
            if (hasWindow(nums, mid, s)){  
                min = Math.min(min, mid);  
                j = mid - 1;  
            }  
            else{  
                i = mid + 1;  
            }  
        }  
        return min == Integer.MAX_VALUE ? 0 : min;  
    }  
    public boolean hasWindow(int[] nums, int curLen, int s){  
        int sum = 0;  
        for (int i = 0 ; i < nums.length ; i++){  
            if (i >= curLen){  
                sum -= nums[i - curLen];  
            }  
            sum += nums[i];  
            if (sum >= s){  
                return true;  
            }  
        }  
        return false;  
    }  
}
```

Problem 221

292 Nim Game

You are playing the following Nim Game with your friend:
There is a heap of stones on the table,
each time one of you take turns to remove 1 to 3 stones.
The one who removes the last stone will be the winner.
You will take the first turn to remove the stones.

Both of you are very clever and have optimal strategies for the game.
Write a function to determine whether
you can win the game given the number of stones in the heap.

For example, if there are 4 stones in the heap,
then you will never win the game:
no matter 1, 2, or 3 stones you remove,
the last stone will always be removed by your friend.

Hint:

If there are 5 stones in the heap, could you figure out
a way to remove the stones such that you will always be the winner?

<https://leetcode.com/problems/nim-game/>

Solution

An interesting problem with a solution with only one line!! Solved by myself!

Listing 286: Solution

```
public class Solution {  
    public boolean canWinNim(int n) {  
        return (n % 4 != 0);  
    }  
}
```

5

Problem 222

290 Word Pattern

Given a pattern and a string str, find if str follows the same pattern.

Here follow means a full match, such that there is a bijection between a letter in pattern and a non-empty word in str.

Examples:

pattern = "abba", str = "dog cat cat dog" should return true.

pattern = "abba", str = "dog cat cat fish" should return false.

pattern = "aaaa", str = "dog cat cat dog" should return false.

pattern = "abba", str = "dog dog dog dog" should return false.

Notes:

You may assume pattern contains only lowercase letters,
and str contains lowercase letters separated by a single space.

<https://leetcode.com/problems/word-pattern/>

Solution

Not difficult, solved by Set and a array(similar to a Hashmap)

Someone solved it by a Set and 2 Hashmaps...

Listing 287: Solution

```
public class Solution {
    public boolean wordPattern(String pattern, String str) {
        String[] strs = new String[26];
        String[] parts = str.split(" ");
        5      if (parts.length != pattern.length()) {
            return false;
        }
        Set<String> strOcc = new HashSet<String>();
        10     for (int i = 0 ; i < 26 ; i++){
            strs[i] = "";
        }
        for (int idx = 0 ; idx < pattern.length() ; idx++){
            int curOffset = pattern.charAt(idx) - 'a';
            15     if (strs[curOffset].length() == 0){
                if (strOcc.contains(parts[idx]))
                    return false;
                else
                    strOcc.add(parts[idx]);
                strs[curOffset] = parts[idx];
                20     continue;
            }
            else{
                if (!strs[curOffset].equals(parts[idx]))
                    25     return false;
            }
        }
        return true;
    }
}
```

Problem 223

316 Remove Duplicate Letters

Given a string which contains only lowercase letters,
remove duplicate letters so that every letter appear once and only once.
You must make sure your result is the smallest
in lexicographical order among all possible results.

Example:

Given "bcabc"

Return "abc"

Given "cbacdcbc"

Return "acdb"

<https://leetcode.com/problems/remove-duplicate-letters/>

Solution

List the train of thought as below:

1. Calculate a array of count where that how many times each letter appears.
2. Find a non-repeating sequence, break the iterating when any character has no successors using the count array maintained.
3. Get the smallest letter in sequence, remove letters in front of it and all same letters behind the element.
4. Iterate to get next element.

Solution description in Chinese can be seen

<https://segmentfault.com/a/1190000004188227>.

Listing 288: Solution

```
public class Solution {
    public String removeDuplicateLetters(String s) {
        if (s.length() == 0)
            return "";
        int[] count = new int[26];
        for (int i = 0 ; i < s.length() ; i++)
            count[s.charAt(i) - 'a']++;
        char minChar = 'z';
        int minPos = 0;
        for (int i = 0 ; i < s.length() ; i++){
            if (minChar > s.charAt(i)){
                minPos = i;
                minChar = s.charAt(i);
            }
            count[s.charAt(i) - 'a']--;
            if (count[s.charAt(i) - 'a'] == 0)
                break;
        }
        return minChar + removeDuplicateLetters(
            s.substring(minPos + 1).replaceAll("" + minChar, ""));
    }
}
```

```
}  
}
```


Problem 224

89 Gray Code

The gray code is a binary numeral system where two successive values differ in only one bit.

Given a non-negative integer n representing the total number of bits in the code, print the sequence of gray code.

A gray code sequence must begin with 0.

For example, given $n = 2$, return $[0,1,3,2]$.

Its gray code sequence is:

00 - 0

01 - 1

11 - 3

10 - 2

Note:

For a given n , a gray code sequence is not uniquely defined.

For example, $[0,2,3,1]$ is also a valid gray code sequence according to the above definition.

For now, the judge is able to judge based on one instance of gray code sequence. Sorry about that.

<https://leetcode.com/problems/gray-code/>

Solution

Consider when $n=2$ and $n=3$

When $n=2$, the sequence is 0,1,3,2, namely 00,01,11,10

When $n=3$, the sequence is 000,001,011,010,110,111,101,100, just the sequence from 2 and add 1 to the front reversely.

Listing 289: Solution 1.1

```
public class Solution {
    public List<Integer> grayCode(int n) {
        List<Integer> result = new ArrayList<Integer>();
        result.add(0);
5       for (int i = 0 ; i < n ; i ++){
            int high = 1 << i;
            for (int j = result.size() - 1; j >= 0; j --){
                result.add(result.get(j) + high);
            }
10      }

        return result;
    }
}
```

There is even a solution looks like this.

Listing 290: Solution 1.2

```
public class Solution {  
    public List<Integer> grayCode(int n) {  
        List<Integer> gray = new ArrayList<Integer>();  
        for(int i = 0 ; i < 1<<n ; i++) {  
5           int temp = i>>1;  
            gray.add(i^temp);  
        }  
        return gray;  
10    }  
}
```

Another alternative solution.

Gray Code, 每次看每次都不记得。写下来让自己好温习。

Gray Code 0 = 0, 下一项是toggle最右边的bit(LSB), 再下一项是toggle最右边值为 "1" bit的左边一个bit。然后重复

如: 3bit

Gray Code: 000, 001, 011, 010, 110, 111, 101, 100, 最右边值为 "1" 的bit在最左边了, 结束。

Binary : 000, 001, 010, 011, 100, 101, 110, 111

再者就是Binary Code 转换为Gray Code了。

如:

Binary Code : 1011 要转换成Gray Code

1011 = 1 (照写第一位), 1(第一位与第二位异或 $1 \oplus 0 = 1$), 1(第二位与第三位, $0 \oplus 1 = 1$), 0 ($1 \oplus 1 = 0$) = 1110

其实就等于 $(1011 \gg 1) \oplus 1011 = 1110$

有了上面的等式写code就简单了

Problem 225

134 Gas Station

There are N gas stations along a circular route,
where the amount of gas at station i is gas[i].

You have a car with an unlimited gas tank and
it costs cost[i] of gas to travel from station i
to its next station (i+1).

You begin the journey with an empty tank at one of the gas stations.

Return the starting gas station's index
if you can travel around the circuit once, otherwise return -1.

Note:

The solution is guaranteed to be unique.

<https://leetcode.com/problems/gas-station/>

Solution 1

First, get all diffs by gas[i] & cost[i], which tells us how much oil will be left after one station.

Then, do you realize that the problem has been transformed into a problem finding the sub-array with maximum sum, but in a **circular** array.

(BTW, how to deal with subarray problems, refer to

<http://blog.csdn.net/hcbbt/article/details/10454947>)

Back to the problem. We have known how to solve sub-array problems, how about circular array?!

Don't be afraid(since I'm stuck for such a long while). The difficulty lays on how to find max sum of subarray starts from latter half of array and ends at former half of it. How about get the minimum of current array? The total sum is given, minimum sum means maximum on the other side. What we need to do differently from sequential problems is to check which one is larger, max or totalSum-min.

Listing 291: Solution 1

```
public class Solution {  
    public int canCompleteCircuit(int[] gas, int[] cost) {  
        int[] diff = new int[gas.length];  
        int[] sum = new int[gas.length];  
5         int totalGas = 0;  
        // Of course I can combine the calculation of totalSum  
        // with calculating max and min  
        for (int i = 0 ; i < gas.length ; i++){  
            diff[i] = gas[i] - cost[i];  
10         totalGas += diff[i];  
        }  
        if (totalGas < 0){  
            return -1;  
        }  
15         int max = Integer.MIN_VALUE, min = Integer.MAX_VALUE;  
        int maxIdx = 0, minIdx = 0;  
        int curMax = 0 , curMin = 0;  
        for (int i = 0 ; i < gas.length ; i++){  
            if (curMax < 0){  
20                 curMax = diff[i];  
            }
```

```
        maxIdx = i;
    }
    else{
        curMax += diff[i];
25    }
    if (curMin > 0){
        curMin = diff[i];
    }
    else{
30        curMin += diff[i];
    }
    if (curMin < min){
        min = curMin;
        minIdx = (i + 1) % gas.length;
35    }
    if (curMax > max){
        max = curMax;
    }
}
40    return (max > totalGas - min) ? maxIdx : minIdx;
}
}
```

Solution 2

Another solution is the jump solution.

Listing 292: Solution 2

```
public class Solution {
    public int canCompleteCircuit(int[] gas, int[] cost) {
        int[] diff = new int[gas.length];
        int totalGas = 0;
5        for (int i = 0 ; i < gas.length ; i++){
            diff[i] = gas[i] - cost[i];
            totalGas += diff[i];
        }
        if (totalGas < 0){
10            return -1;
        }
        int startIdx = 0;
        int sum = 0;
        for (int i = 0 ; i < gas.length ; i++){
15            sum += diff[i];
            if (sum < 0){
                startIdx = (i + 1) % gas.length;
                sum = 0;
            }
        }
20        return startIdx;
    }
}
```

Problem 226

135 Candy

There are N children standing in a line.
Each child is assigned a rating value.

You are giving candies to these children
subjected to the following requirements:

Each child must have at least one candy.
Children with a higher rating get more candies than their neighbors.
What is the minimum candies you must give?

<https://leetcode.com/problems/candy/>

Solution 1

Solution 1 is greedy. Iterating from front to end, when there is a pair which $\text{ratings}[i] < \text{ratings}[i + 1]$, give $\text{candies}[i] = 1$ OR stay unchanged; give $\text{candies}[i + 1] = \text{candies}[i] + 1$; Then iterate backward, doing the similar operation, except for the condition when the i^{th} people has already been given the candies, that's when we give the man $\text{max}(\text{candies}[i], \text{candies}[i + 1] + 1)$.

Finally, don't forget to give all children who aren't given any candies 1 candy for comfort.(Consider the input of [2,2,1])

Listing 293: Solution 1

```
public class Solution {
    public int candy(int[] ratings) {
        if (ratings.length == 0){
            return 0;
        }
        if (ratings.length == 1){
            return 1;
        }
        int candies[] = new int[ratings.length];
        Arrays.fill(candies, 1);
        for (int i = 0 ; i < ratings.length - 1 ; i++){
            if (ratings[i] < ratings[i + 1]){
                candies[i + 1] = candies[i] + 1;
            }
        }
        for (int i = ratings.length - 1 ; i > 0 ; i--){
            if (ratings[i] < ratings[i - 1]){
                candies[i - 1] = Math.max(candies[i] + 1, candies[i - 1]);
            }
        }

        int sum = 0;
        for (int i = 0 ; i < candies.length ; i++){
            sum += candies[i];
        }

        return sum;
    }
}
```

```
}

```

Solution 2

Solution 2 is tactical and can avoid using extra $O(n)$ space. Let's consider the non-descending condition - [1,2,3,4,5] OR [1,2,3,3,4], when the sequence is ascending, we don't need to worry about the candy-giving thing, since the result is obvious after one pass - [1,2,3,4,5] & [1,2,3,1,2]. Now consider something more complicated - [1,5,4,3,2], still we give kid_0 (with rating of 1) 1 candy, kid_1 2 candies temporarily, then we meet kid_2 with 4 ratings, give him 1 (we'll know why afterwards), and now we are in a descending sequence with length of 2 by now.

Then the important part. We meet kid_3 with 3 ratings, which is still in the descending sequence, and that's when we know we have underestimated kid_1 , so we give him an extra candy, namely $total++$; don't forget kid_3 , we give him 2 candies (here we are balancing kid_2 and kid_3 , in fact kid_2 is the one who should get 2 candies by now, not kid_3 , since we are calculating the total amount of candies, so it doesn't matter who get the candy.) Similarly, when kid_4 arrives, we give him (or kid_2 , to be more precious) 3 candies and compensate kid_1 again.

Listing 294: Solution 2

```

public class Solution {
    public int candy(int[] ratings) {
        if (ratings.length == 0){
            return 0;
        }
        if (ratings.length == 1){
            return 1;
        }
        int sum = 1;
        int descLen = 0;
        int ascCandyCnt = 1;
        int descCandyCnt = 0;
        for (int i = 1 ; i < ratings.length ; i++){
            if (ratings[i] > ratings[i - 1]){
                if (descLen != 0){
                    ascCandyCnt = 1;
                }
                ascCandyCnt ++;
                sum += ascCandyCnt;
                descLen = 0;
                descCandyCnt = 0;
            }
            else if (ratings[i] == ratings[i - 1]){
                sum ++;
                ascCandyCnt = 1;
                descLen = 0;
                descCandyCnt = 0;
            }
            else{
                // ratings[i] < ratings[i - 1]
                descLen ++;
                if (ascCandyCnt <= descLen){
                    sum ++;
                }
                descCandyCnt ++;
                sum += descCandyCnt;
            }
        }
    }
}

```

```
        }  
    }  
    return sum;  
}  
}
```

Problem 227

205 Isomorphic Strings

Given two strings s and t,
determine if they are isomorphic.

Two strings are isomorphic if
the characters in s can be replaced to get t.

All occurrences of a character must be replaced with
another character while preserving the order of characters.
No two characters may map to the same character but
a character may map to itself.

For example,
Given "egg", "add", return true.

Given "foo", "bar", return false.

Given "paper", "title", return true.

Note:
You may assume both s and t have the same length.

<https://leetcode.com/problems/isomorphic-strings/>

Solution

An easy problem using HashMap..... Wait a minute! HashMaps(or a map with a set)! Consider when the input is "ab" & "aa".

Listing 295: Solution

```
public class Solution {  
    public boolean isIsomorphic(String s, String t) {  
        if (s.length() == 0){  
            return true;  
5         }  
        Map<Character, Character> map = new HashMap<Character, Character>();  
        Map<Character, Character> reverseMap  
            = new HashMap<Character, Character>();  
  
10         for (int idx = 0 ; idx < s.length() ; idx++){  
            char charS = s.charAt(idx);  
            char charT = t.charAt(idx);  
  
            if (map.containsKey(charS)){  
15                 char compare = map.get(charS);  
                 if (compare != charT){  
                     return false;  
                 }  
            }  
20         else{  
             if (reverseMap.containsKey(charT)){  
                 return false;  
             }  
         }  
     }  
}
```



```
25         }
        map.put(charS, charT);
        reverseMap.put(charT, charS);
    }
}

30     return true;
}
}
```

Problem 228

4 Median of Two Sorted Arrays

There are two sorted arrays `nums1` and `nums2` of size `m` and `n` respectively.

Find the median of the two sorted arrays.

The overall run time complexity should be $O(\log(m+n))$.

<https://leetcode.com/problems/median-of-two-sorted-arrays/>

Solution 1

Honestly, this is not a correct solution though it got AC on OJ. The complexity is $O(m+n)$ rather than $O(\log(m+n))$.

Listing 296: Solution 1

```
public class Solution {
    public double findMedianSortedArrays(int[] nums1, int[] nums2) {
        int len1 = nums1.length;
        int len2 = nums2.length;
5       int mid = (len1 + len2) / 2, midVal = 0;
        int mid2 = mid + 1, mid2Val = 0;
        if (mid == 0){
            return (nums1.length == 0) ? nums2[0] : nums1[0];
        }
10      int idx1 = 0, idx2 = 0;
        int idx = 0;
        while (idx1 < nums1.length && idx2 < nums2.length){
            int compare1 = nums1[idx1], compare2 = nums2[idx2];
            idx++;
15          int compare = 0;
            if (compare1 <= compare2){
                compare = compare1;
                idx1++;
            }
20          else{
                compare = compare2;
                idx2++;
            }
            if (idx == mid){
25              midVal = compare;
            }
            if (idx == mid2){
                mid2Val = compare;
                break;
30          }
        }
        if (idx < mid2){
            while (idx1 < nums1.length){
35              idx++;
                if (idx == mid){
                    midVal = nums1[idx1];
                }
                if (idx == mid2){
```

```

        mid2Val = nums1[idx1];
40    }
        idx1++;
    }
    while (idx2 < nums2.length){
        idx++;
45    if (idx == mid){
        midVal = nums2[idx2];
        }
        if (idx == mid2){
            mid2Val = nums2[idx2];
50        }
        idx2++;
    }
}

55
if ((len1 + len2) % 2 == 0){
    return (midVal + mid2Val) * 1.0 / 2;
}
else{
60    return mid2Val;
}
}
}

```

Solution 2

Solution 2 makes use of Binary Search, with a pretty complicated process on calculating the related index. The train of thought is quite interesting.

For detailed description on solution , visit <http://blog.csdn.net/zxzxy1988/article/details/8587244>

Listing 297: Solution 2

```

public class Solution {
    private double findMedian(int[] nums1, int[] nums2,
                               int startIdx1, int startIdx2, int remaining){
        if (nums1.length - startIdx1 < nums2.length - startIdx2){
5            return findMedian(nums2, nums1, startIdx2, startIdx1, remaining);
        }
        // Since length1 > length2, no need to add this snippet
        // if (startIdx1 == nums1.length){
        //     return nums2[startIdx2 + remaining - 1];
10        // }
        if (startIdx2 == nums2.length){
            return nums1[startIdx1 + remaining - 1];
        }
        if (remaining == 1){
15            return Math.min(nums1[startIdx1], nums2[startIdx2]);
        }
        int offset2 = Math.min(nums2.length - startIdx2, remaining / 2);
        int offset1 = remaining - offset2;
        if (nums1[startIdx1 + offset1 - 1] == nums2[startIdx2 + offset2 - 1]){

```

```
20         return nums1[startIdx1 + offset1 - 1];
    }
    else if (nums1[startIdx1 + offset1 - 1] > nums2[startIdx2 + offset2 - 1]){
        return findMedian(nums1, nums2,
            startIdx1, startIdx2 + offset2, remaining - offset2);
25    }
    else{
        return findMedian(nums1, nums2,
            startIdx1 + offset1, startIdx2, remaining - offset1);
    }
30 }

public double findMedianSortedArrays(int[] nums1, int[] nums2) {
    int m = nums1.length, n = nums2.length;
    if ((m + n) % 2 != 0){
35         return findMedian(nums1, nums2, 0, 0, (m + n) / 2 + 1);
    }
    else{
        return (findMedian(nums1, nums2, 0, 0, (m + n) / 2)
            + findMedian(nums1, nums2, 0, 0, (m + n) / 2 + 1)) / 2;
40    }
}
}
```

Problem 229

295 Find Median from Data Stream

Median is the middle value in an ordered integer list.

If the size of the list is even, there is no middle value.

So the median is the mean of the two middle value.

Examples:

[2,3,4] , the median is 3

[2,3], the median is $(2 + 3) / 2 = 2.5$

Design a data structure that supports the following two operations:

`void addNum(int num)` - Add a integer number from the data stream to the data structure.

`double findMedian()` - Return the median of all elements so far.

For example:

```
add(1)
add(2)
findMedian() -> 1.5
add(3)
findMedian() -> 2
```

<https://leetcode.com/problems/find-median-from-data-stream/>

Though almost all solutions on Internet like CSDN has only one way to solve this - using two heaps, one max heap and one min heap, this is still a problem worth 2+ hours to think carefully about it, to be honest.

Solution 1 - Array & List

Honestly, this is not a correct solution though it got AC on OJ. The complexity is $O(m + n)$ rather than $O(\log(m + n))$.

Unsorted array is enough for the problem. Partition is what we usually do in QuickSort, so as to finish the task of finding the median. Partition is of $O(n)$ runtime, where inserting to unsorted array only costs $O(1)$. This method will get TLE in OJ.

Listing 298: Solution 1.1 - Unsorted Array

```
public class MedianFinder {
    static final int MAXLEN = 100000;
    int[] nums = new int[MAXLEN];
    int length = 0;
5    // Adds a number into the data structure.
    public void addNum(int num) {
        nums[length++] = num;
    }

10    // Returns the median of current data stream
    public double findMedian() {
        if (length % 2 == 0){
            return (partition(nums, 0, length - 1, length / 2 - 1)
15                + partition(nums, 0, length - 1, length / 2)) / 2;
        }
    }
}
```

```

        else{
            return partition(nums, 0, length - 1, length / 2);
        }
    }

    private double partition(int[] nums, int start, int end, int target){
        // if (start > end){
        //     return -1;
        // }
        int pivot = nums[start];
        int low = start, high = end;
        while (start < end){
            while (start < end && nums[end] >= pivot){
                end--;
            }
            nums[start] = nums[end];

            while (start < end && nums[start] <= pivot){
                start++;
            }
            nums[end] = nums[start];
        }
        nums[start] = pivot;
        if (start == target){
            return pivot;
        }
        else if (start > target){
            return partition(nums, low, start - 1, target);
        }
        else{
            return partition(nums, start + 1, high, target);
        }
    }
};

// Your MedianFinder object will be instantiated and called as such:
// MedianFinder mf = new MedianFinder();
// mf.addNum(1);
// mf.findMedian();

```

Sorted array may make getting median faster(but inserting elements to container slower). Inserting an element to the array and keep it sorted needs $O(1)$ runtime, but getting median only need $O(1)$ runtime. This solution gets AC in OJ.

Listing 299: Solution 1.2 - Sorted Array

```

public class MedianFinder {
    static final int MAXLEN = 100000;
    int[] nums = new int[MAXLEN];
    int length = 0;
    // Adds a number into the data structure.

```

```

10     public void addNum(int num) {
        int pos = 0;
        for (int idx = length - 1 ; idx >= 0 ; idx --){
            if (nums[idx] > num){
                nums[idx + 1] = nums[idx];
            }
            else{
                pos = idx + 1;
                break;
            }
        }
        nums[pos] = num;
        length ++;
    }

20     // Returns the median of current data stream
    public double findMedian() {
        if (length % 2 == 1){
            return nums[length / 2] * 1.0;
        }
        else{
            return (nums[length / 2 - 1] + nums[length / 2]) * 1.0 / 2;
        }
    }
30 };

    // Your MedianFinder object will be instantiated and called as such:
    // MedianFinder mf = new MedianFinder();
    // mf.addNum(1);
35 // mf.findMedian();

```

What's more, we can use Sorted List to solve the problem, adding 2 pointers pointing at 2 medians can reduce the complexity of getting median to $O(1)$ (when the length is odd, the pointers point at the same element.)

Solution 2 - Heaps(Max & Min) in form of PriorityQueue

Solution is more proper to solve the problem.

Maintain 2 heaps, one max and another min. Keep them have the same size or max heap is 1 element larger than min heap in size. We store the biggest k elements in min heap while the others(k OR k + 1) in max heap. Then either the median is root of max heap OR average of roots from both trees.

How to maintain their sizes? We always push eleme

Listing 300: Solution 2 - Heaps

```

    public class MedianFinder {
        PriorityQueue<Integer> minHeap = new PriorityQueue<Integer>();
        PriorityQueue<Integer> maxHeap =
        new PriorityQueue<Integer>(10000, new Comparator<Integer>(){
5            @Override
            public int compare(Integer x, Integer y){
                return y - x;
            }
        });
10    // Adds a number into the data structure.

```

```

    public void addNum(int num) {
        if (minHeap.size() == maxHeap.size()) {
            maxHeap.offer(num);
        }
15      else {
            minHeap.offer(num);
        }
        if (!minHeap.isEmpty() && (maxHeap.peek() > minHeap.peek())) {
            minHeap.offer(maxHeap.poll());
20      maxHeap.offer(minHeap.poll());
        }
    }

    // Returns the median of current data stream
25  public double findMedian() {
        if ((minHeap.size() + maxHeap.size()) % 2 == 0) {
            return (maxHeap.peek() + minHeap.peek()) * 1.0 / 2;
        }
        else {
30      return maxHeap.peek();
        }
    }
};

35 // Your MedianFinder object will be instantiated and called as such:
// MedianFinder mf = new MedianFinder();
// mf.addNum(1);
// mf.findMedian();

```

Solution3 - Trees(BST, Binary Search Tree OR AVL Tree)

Since in Java there are no built-in structures similar to BST or AVL, so I won't list the code here, just think of how to finish out task.

We may add a field in TreeNode to make searching for median faster(to $O(\log n)$).

To avoid the worst condition of BST, we may use AVL- the highly balanced tree structure to finish what BST need to do.

In summary, all sorts of methods can be concluded as figure below,

Type of Data Container	Time to Insert	Time to Get Median
Unsorted Array	$O(1)$	$O(n)$
Sorted Array	$O(n)$	$O(1)$
Sorted List	$O(n)$	$O(1)$
Binary Search Tree	$O(\log n)$ on average, $O(n)$ for the worst case	$O(\log n)$ on average, $O(n)$ for the worst case
AVL	$O(\log n)$	$O(1)$
Max and Min Heap	$O(\log n)$	$O(1)$

Problem 230

17 Letter Combinations of a Phone Number

Given a digit string, return all possible letter combinations that the number could represent.

A mapping of digit to letters (just like on the telephone buttons) is given below.

Input: Digit string "23"

Output: ["ad", "ae", "af", "bd", "be", "bf", "cd", "ce", "cf"].

Note:

Although the above answer is in lexicographical order, your answer could be in any order you want.

<https://leetcode.com/problems/letter-combinations-of-a-phone-number/>



Solution

A backtracking problem, similar to Problem "Combinations".

Listing 301: Solution

```
public class Solution {
    String[] mapping = {"", "", "abc", "def", "ghi", "jkl", "mno", "pqrs", "tuv", "wxyz"};
    List<String> result = new ArrayList<String>();

    5     public List<String> letterCombinations(String digits) {
        traversePhone(digits, 0, new ArrayList<Character>());
        return result;
    }

    10    private void traversePhone(String digits, int depth, List<Character> list){
        if (digits.length() == 0){
            return;
        }
        15    if (depth == digits.length()){
            String newStr = "";
            for (char ele : list){
                newStr += ele;
            }
            result.add(newStr);
            20    return;
        }
    }
```

```
    }
    String possibilities = mapping[digits.charAt(depth) - '0'];
    if (possibilities.length() == 0){
        return;
25    }
    else{
        for (int idx = 0 ; idx < possibilities.length() ; idx ++){
            list.add(possibilities.charAt(idx));
            traversePhone(digits , depth + 1, list);
30            list.remove(list.size() - 1);
        }
    }
}
35 }
```

Problem 231

71 Simplify Path

Given an absolute path for a file (Unix-style), simplify it.

For example,

path = "/home/", => "/home"

path = "/a/./b/../../c/", => "/c"

click to show corner cases.

Corner Cases:

Did you consider the case where path = "/../"?

In this case, you should return "/".

Another corner case is the path might

contain multiple slashes '/' together, such as "/home//foo/".

In this case, you should ignore redundant slashes

and return "/home/foo".

<https://leetcode.com/problems/simplify-path/>

Solution

It's obvious that the problem is proper to be solved using Stack, there are 2 things to be taken into consideration:

1. Mind those corner cases.(See detailed cases in problem description)
2. Don't forget to reverse those remaining paths in Stack, /home/test will be in the order of "test -> home" in Stack.

Listing 302: Solution

```
public class Solution {
    public String simplifyPath(String path) {
        if (path.length() == 0){
            return "/";
        }
        String result = "";
        Stack<String> stack = new Stack<String>();
        int startIdx = 0;

        while (startIdx < path.length() && path.charAt(startIdx) == '/'){
            startIdx++;
        }

        while (startIdx < path.length()){
            int curIdx = startIdx;
            while (curIdx < path.length() && path.charAt(curIdx) != '/'){
                curIdx++;
            }
            int endIdx = curIdx;
            String curNode = path.substring(startIdx, endIdx);
            if (!curNode.equals(".")){
                if (curNode.equals("..")){
                    if (!stack.isEmpty()){

```

```

        stack.pop();
25         }
        }
        else{
            stack.push(curNode);
        }
30     }
    while (curIdx < path.length() && path.charAt(curIdx) == '/'){
        curIdx++;
    }
    startIdx = curIdx;
35 }
Stack<String> resultStack = new Stack<String>();
while (!stack.isEmpty()){
    resultStack.push(stack.pop());
}
40 while (!resultStack.isEmpty()){
    result += "/" ;
    result += resultStack.pop();
}
if (result.length() == 0){
45     result += "/" ;
}

    return result;
}
50 }
```

Problem 232

22 Generate Parentheses

Given n pairs of parentheses,
write a function to generate all combinations of well-formed parentheses.

For example, given $n = 3$, a solution set is:

"((()))", "(()())", "(())()", "()(())", "()()()"

<https://leetcode.com/problems/generate-parentheses/>

Solution

An easy backtracking problem, when remaining left parentheses number equals right number, we must pick left one so that there won't be illegal pairs; when left less than right, we can pick either left or right parentheses.

Listing 303: Solution

```
public class Solution {  
    List<String> result = new ArrayList<String>();  
    private void dfs(String curStr, int leftRem, int rightRem){  
        if (leftRem == 0 && rightRem == 0){  
            result.add(curStr);  
            return;  
        }  
        if (leftRem > rightRem){  
            return;  
        }  
        else if (leftRem == rightRem){  
            dfs(curStr + "(", leftRem - 1, rightRem);  
        }  
        else{  
            // leftRem < rightRem  
            if (leftRem > 0){  
                dfs(curStr + "(", leftRem - 1, rightRem);  
            }  
            dfs(curStr + ")", leftRem, rightRem - 1);  
        }  
    }  
    public List<String> generateParenthesis(int n) {  
        dfs("", n, n);  
        return result;  
    }  
}
```

Problem 233

338 Counting Bits

Given a non negative integer number num.
For every numbers i in the range $0 \leq i \leq \text{num}$
calculate the number of 1's in their binary representation and
return them as an array.

Example:

For num = 5 you should return [0,1,1,2,1,2].

Follow up:

It is very easy to come up with a solution
with run time $O(n \cdot \text{sizeof}(\text{integer}))$.

But can you do it in linear time $O(n)$ /possibly in a single pass?

Space complexity should be $O(n)$.

Can you do it like a boss? Do it without using

any builtin function like `__builtin_popcount` in c++ or in any other language.

Hint:

You should make use of what you have produced already.

Divide the numbers in ranges like [2-3], [4-7], [8-15] and so on.

And try to generate new range from previous.

Or does the odd/even status of the number

help you in calculating the number of 1s?

<https://leetcode.com/problems/counting-bits/>

Solution

An easy backtracking problem, when remaining left parentheses number equals right number, we must pick left one so that there won't be illegal pairs; when left less than right, we can pick either left or right parentheses.

Listing 304: Solution 1.1

```
public class Solution {
    public int[] countBits(int num) {
        int[] result = new int[num + 1];
        if (num == 0){
5           return result;
        }
        int tail = 1;
        result[0] = 0;
        int count = 0;
10        while (tail <= num){
            int prevTail = tail;
            for (int i = 0 ; i < prevTail ; i++){
                result[tail++] = result[i] + 1;
                if (tail > num){
15                     return result;
                }
            }
        }
    }
}
```

```
    }  
    return result;  
20 }  
}
```

Another solution, the train of thought is similar in nature with mine, simpler.

Listing 305: Solution 1.2

```
public int [] countBits(int num) {  
    int [] result = new int[num + 1];  
    result[0] = 0;  
    if(num == 0){ return result;}  
5    int leftMostOne = 1; // will increase as 2, 4, 8, 16, 32 ....  
    for(int i = 1; i <= num; i++){  
        if(i == 2 * leftMostOne){ leftMostOne = i;}  
        result[i] = 1 + result[i - leftMostOne];  
    }  
10    return result;  
}
```

Problem 234

299 Bulls and Cows

You are playing the following Bulls and Cows game with your friend:
You write down a number and ask your friend to guess what the number is.
Each time your friend makes a guess, you provide a hint that indicates
how many digits in said guess match your secret number exactly
in both digit and position (called "bulls") and
how many digits match the secret number but
locate in the wrong position (called "cows").
Your friend will use successive guesses and hints to eventually
derive the secret number.

For example:

Secret number: "1807"

Friend's guess: "7810"

Hint: 1 bull and 3 cows. (The bull is 8, the cows are 0, 1 and 7.)

Write a function to return a hint according to the secret number
and friend's guess, use A to indicate the bulls and B to indicate the cows.
In the above example, your function should return "1A3B".

Please note that both secret number and friend's guess may
contain duplicate digits, for example:

Secret number: "1123"

Friend's guess: "0111"

In this case, the 1st 1 in friend's guess is a bull,
the 2nd or 3rd 1 is a cow, and your function should return "1A1B".
You may assume that the secret number and your friend's guess only
contain digits, and their lengths are always equal.

<https://leetcode.com/problems/bulls-and-cows/>

Solution

I'm following an easy simulation, deal with 'A' & 'B' separately, but I think there is space for optimization.
Last FREE Easy Problem solved till 3rd, Apr., 2016.

Listing 306: Solution

```
public class Solution {
    public String getHint(String secret, String guess) {
        int[] count = new int[10];
        boolean[] visited = new boolean[secret.length()];
5       int cntA = 0, cntB = 0;
        for (int i = 0 ; i < secret.length() ; i++){
            if (secret.charAt(i) == guess.charAt(i)){
                cntA++;
                visited[i] = true;
10          }
            else{
                count[secret.charAt(i) - '0']++;
            }
        }
    }
}
```



```
    }  
15    for (int i = 0 ; i < secret.length() ; i ++){  
        if (!visited[i]){  
            char guessChar = guess.charAt(i);  
            if (count[guessChar - '0'] > 0){  
                count[guessChar - '0']--;  
20                cntB ++;  
            }  
        }  
    }  
25    return cntA + "A" + cntB + "B";  
}
```

Problem 235

162 Find Peak Element

A peak element is an element that is greater than its neighbors.

Given an input array where $\text{num}[i] \neq \text{num}[i+1]$,
find a peak element and return its index.

The array may contain multiple peaks,
in that case return the index to any one of the peaks is fine.

You may imagine that $\text{num}[-1] = \text{num}[n] = -\infty$ (since I cannot type the symbol here)

For example, in array $[1, 2, 3, 1]$, 3 is a peak element
and your function should return the index number 2.

Note:
Your solution should be in logarithmic complexity.

<https://leetcode.com/problems/find-peak-element/>

Solution 1

$O(n)$ method, simple simulation, sequential method.

Listing 307: Solution 1

```
public class Solution {
    public int findPeakElement(int[] nums) {
        if (nums.length == 0){
            return -1;
5        }
        for (int i = 0 ; i < nums.length ; i++){
            if (i == 0 && i == nums.length - 1){
                return 0;
            }
10         else if (i == 0){
            if (nums[0] > nums[1]){
                return 0;
            }
        }
15         else if (i == nums.length - 1){
            if (nums[i] > nums[i - 1]){
                return i;
            }
        }
20         else{
            if (nums[i] > nums[i - 1] && nums[i] > nums[i + 1]){
                return i;
            }
        }
25     }
    return -1;
}
```

```
}

```

Solution 2

Binary search will reduce the complexity to $O(\log n)$.

Listing 308: Solution 2

```
public class Solution {
    public int findPeakElement(int[] nums) {
        int low = 0;
        int high = nums.length - 1;
5       while (low < high) {
            int mid = low + (high - low) / 2;
            if (nums[mid] > nums[mid + 1])
                high = mid;
            else
10             low = mid + 1;
        }
        return low;
    }
}
```

Problem 236

318 Maximum Product of Word Lengths

Given a string array `words`, find the maximum value of `length(word[i]) * length(word[j])` where the two words do not share common letters. You may assume that each word will contain only lower case letters. If no such two words exist, return 0.

Example 1:

Given ["abcw", "baz", "foo", "bar", "xtfn", "abcdef"]

Return 16

The two words can be "abcw", "xtfn".

Example 2:

Given ["a", "ab", "abc", "d", "cd", "bcd", "abcd"]

Return 4

The two words can be "ab", "cd".

Example 3:

Given ["a", "aa", "aaa", "aaaa"]

Return 0

No such pair of words.

<https://leetcode.com/problems/maximum-product-of-word-lengths/>

Solution

Use 26-bit integers to represent each word(since a & aa means the same to us in this problem). For example, 'abd' is 0000..001011, then we can know if two integers carry out bit and & calculation and get a non-zero result, the two have some letters in common and vice versa.

Listing 309: Solution

```
public class Solution {
    public int maxProduct(String[] words) {
        int[] bitMap = new int[words.length];
        for (int idx = 0 ; idx < words.length; idx++){
            String word = words[idx];
            for(int strIdx = 0 ; strIdx < word.length() ; strIdx++){
                char curChar = word.charAt(strIdx);
                bitMap[idx] |= 1 << curChar - 'a';
            }
        }
        int max = 0;
        for (int i = 0 ; i < words.length ; i++){
            for (int j = i + 1 ; j < words.length ; j++){
                if ((bitMap[i] & bitMap[j]) == 0){
                    max = Math.max(words[i].length() * words[j].length(), max);
                }
            }
        }
        return max;
    }
}
```

Problem 237

164 Maximum Gap

Given an unsorted array,
find the maximum difference between the successive elements
in its sorted form.

Try to solve it in linear time/space.

Return 0 if the array contains less than 2 elements.

You may assume all elements in the array are
non-negative integers and fit in the 32-bit signed integer range.

<https://leetcode.com/problems/maximum-gap/>

Solution

Hard problem is worthy of its name... Though this problem is not so difficult to use Bucket Sort, there are so many edge cases and annoying index calculation.

Back to the problem, the maximum gap can only appear between maximum and minimum value of 2 non-empty adjacent buckets, we can use this idea to solve the problem.

Listing 310: Solution

```
class Bucket{
    int max = Integer.MIN_VALUE;
    int min = Integer.MAX_VALUE;
    int size = 0;
5    List<Integer> list = new ArrayList<Integer>();
}
public class Solution {
    public int maximumGap(int[] nums) {
        if (nums.length < 2){
10            return 0;
        }
        if (nums.length == 2){
            return Math.abs(nums[0] - nums[1]);
        }
15        int totalMax = Integer.MIN_VALUE;
        int totalMin = Integer.MAX_VALUE;
        for (int i = 0 ; i < nums.length ; i++){
            totalMax = Math.max(totalMax, nums[i]);
            totalMin = Math.min(totalMin, nums[i]);
20        }

        int interval = (int) Math.ceil(((double)(totalMax + 1 - totalMin)
            / (nums.length - 1)));
        int bucketCnt = (int) Math.ceil(((double)
25            (totalMax + 1 - totalMin) / interval));
        Bucket[] buckets = new Bucket[bucketCnt];
        for (int i = 0 ; i < buckets.length ; i++){
            buckets[i] = new Bucket();
        }
    }
}
```

```
    }  
30  
    if (interval == 0){  
        return 0;  
    }  
    for (int i = 0 ; i < nums.length ; i++){  
35        int idx = (nums[i] - totalMin) / interval;  
        buckets[idx].size ++;  
        buckets[idx].list.add(nums[i]);  
        buckets[idx].max = Math.max(buckets[idx].max, nums[i]);  
        buckets[idx].min = Math.min(buckets[idx].min, nums[i]);  
40    }  
  
    int maxGap = 0;  
    int lastMax = Integer.MIN_VALUE;  
    for (int i = 0 ; i < buckets.length ; i++){  
45        if (buckets[i].size == 0){  
            continue;  
        } else{  
            if (lastMax == Integer.MIN_VALUE){  
                lastMax = buckets[i].max;  
50            } else{  
                maxGap = Math.max(maxGap, buckets[i].min - lastMax);  
                lastMax = buckets[i].max;  
            }  
        }  
55    }  
  
    return maxGap;  
}
```

Problem 238

332 Reconstruct Itinerary

Given a list of airline tickets represented by pairs of departure and arrival airports [from, to], reconstruct the itinerary in order.
All of the tickets belong to a man who departs from JFK.
Thus, the itinerary must begin with JFK.

Note:

If there are multiple valid itineraries, you should return the itinerary that has the smallest lexical order when read as a single string.

For example,

the itinerary ["JFK", "LGA"] has a smaller lexical order than ["JFK", "LGB"].

All airports are represented by three capital letters (IATA code).
You may assume all tickets form at least one valid itinerary.

Example 1:

```
tickets
= [["MUC", "LHR"], ["JFK", "MUC"], ["SFO", "SJC"], ["LHR", "SFO"]]
Return ["JFK", "MUC", "LHR", "SFO", "SJC"].
```

Example 2:

```
tickets
= [["JFK", "SFO"], ["JFK", "ATL"], ["SFO", "ATL"], ["ATL", "JFK"], ["ATL", "SFO"]]
Return ["JFK", "ATL", "JFK", "SFO", "ATL", "SFO"].
Another possible reconstruction is ["JFK", "SFO", "ATL", "JFK", "ATL", "SFO"].
But it is larger in lexical order.
```

<https://leetcode.com/problems/reconstruct-itinerary/>

Solution

A clear DFS problem in adjacency graph. However, I truly struggled for a long while on the 'smallest lexical order'. So sort them!

Listing 311: Solution

```
public class Solution {
    public List<String> findItinerary(String[][] tickets) {
        if (tickets.length == 0) {
            return new ArrayList<String>();
        }
        Map<String, List<String>> map = new HashMap<String, List<String>>();
        for (String[] ticket : tickets) {
            String from = ticket[0], to = ticket[1];
            List<String> oldList = map.containsKey(from) ? map.get(from)
            : new ArrayList<String>();
            oldList.add(to);
            map.put(from, oldList);
        }
        for (String key : map.keySet()) {
            Collections.sort(map.get(key));
        }
    }
}
```

```
    }
    List<String> result = new ArrayList<String>();
    exploreItinerary(map, "JFK", result, tickets.length + 1);
    return result;
20 }

private boolean exploreItinerary(Map<String, List<String>> map,
    String curPort, List<String> result, int targetSize) {
    result.add(curPort);
25    if (result.size() == targetSize) {
        return true;
    }
    if (map.get(curPort) != null){
        // Note: you can remove elements when iterating using index
30    for (int idx = 0 ; idx < map.get(curPort).size() ; idx++) {
        String to = map.get(curPort).get(idx);
        map.get(curPort).remove(idx);
        if (exploreItinerary(map, to, result, targetSize)) {
            return true;
35        }
        map.get(curPort).add(to);
        Collections.sort(map.get(curPort));
    }
    }
40    result.remove(result.size() - 1);
    return false;
}
}
```


Problem 239

91 Decode Ways

A message containing letters from A-Z is
being encoded to numbers using the following mapping:

'A' -> 1
'B' -> 2
...
'Z' -> 26

Given an encoded message containing digits,
determine the total number of ways to decode it.

For example,
Given encoded message "12", it could be decoded as
"AB" (1 2) or "L" (12).

The number of ways decoding "12" is 2.

<https://leetcode.com/problems/decode-ways/>

Solution

```
dp[i] = ((last_digit==0) ? dp[i - 1] : 0) + ((10<=last_2_digits<=26) ? dp[i - 2] : 0);
```

Listing 312: Solution

```
public class Solution {  
    public int numDecodings(String s) {  
        if (s.length() == 0 || s.charAt(0) == '0'){  
            return 0;  
        }  
        if(s.length() == 1){  
            return 1;  
        }  
        int[] dp = new int[s.length()];  
        dp[0] = 1;  
        if (Integer.parseInt(s.substring(0, 2)) <= 26  
            && Integer.parseInt(s.substring(0, 2)) >= 10){  
            if (Integer.parseInt(s.substring(0, 2)) % 10 == 0)  
                dp[1] = 1;  
            else  
                dp[1] = 2;  
        }  
        else{  
            if (Integer.parseInt(s.substring(0, 2)) % 10 == 0)  
                dp[1] = 0;  
            else  
                dp[1] = 1;  
        }  
        for (int idx = 2; idx < s.length() ; idx++){  
            int singleDigit = Integer.parseInt(s.substring(idx , idx + 1));  
            int doubleDigits = Integer.parseInt(s.substring(idx - 1 , idx + 1));
```

```
        if (singleDigit != 0){
            dp[idx] += dp[idx - 1];
30    }
        if ((doubleDigits >= 10) && (doubleDigits <= 26)){
            dp[idx] += dp[idx - 2];
        }
35    }

    return dp[s.length() - 1];
}
}
```



```
        long intL = numL / denL;
        long remL = numL % denL;

25        result += intL;
        if (remL == 0){
            return result;
        }
        else{
30            result += ".";
        }

        Map<Long, Integer> remPos = new HashMap<Long, Integer>();
        // remPos.put(remL, 0);
35        String decimal = "";
        int count = 0;
        while (remL != 0){
            if (remPos.containsKey(remL)){
                int latterIdx = remPos.get(remL);
40                String part1 = decimal.substring(0, latterIdx);
                String part2 = decimal.substring(latterIdx);
                result += part1 + "(" + part2 + ")";
                return result;
            }
45            remPos.put(remL, count ++);
            remL *= 10;
            intL = remL / denL;
            decimal += intL;
            remL = remL % denL;
50        }
        result += decimal;
        return result;
    }
}
```

Problem 241

301 Remove Invalid Parentheses

Remove the minimum number of invalid parentheses in order to make the input string valid. Return all possible results.

Note: The input string may contain letters other than the parentheses (and).

Examples:

```
"()()()" -> ["()()()", "(()())"]
"(a)()()" -> ["(a)()()", "(a())"]
")(" -> [""]
```

<https://leetcode.com/problems/remove-invalid-parentheses/>

Solution - BFS

The problem can be solved in both DFS & BFS way.

Basic train of thought is to remove every character in string iteratively to check if generated string is valid.

For BFS, it is a natural idea when we want to find the minimum of invalid parentheses.

Listing 314: Solution 1

```
public class Solution {
    public List<String> removeInvalidParentheses(String s) {
        Queue<String> queue = new LinkedList<String>();
        queue.offer(s);
        Set<String> visited = new HashSet<String>();
        Set<String> queueVisited = new HashSet<String>();
        List<String> result = new ArrayList<String>();
        if (isValid(s)){
            result.add(s);
            return result;
        }
        boolean found = false;
        while (!found && !queue.isEmpty()){
            List<String> curLevel = new ArrayList<String>();
            while (!found && !queue.isEmpty()){
                curLevel.add(queue.poll());
            }
            for (String cur : curLevel){
                for (int splitIdx = 0 ; splitIdx < cur.length() ; splitIdx ++){
                    if (cur.charAt(splitIdx) != '('
                        && cur.charAt(splitIdx) != ')'){
                        continue;
                    }
                    String newStr = cur.substring(0, splitIdx)
                        + cur.substring(splitIdx + 1);
                    if (queueVisited.contains(newStr)){
                        continue;
                    }
                }
            }
        }
    }
}
```

```

30         if (isValid(newStr)){
            if (!visited.contains(newStr)){
                visited.add(newStr);
                result.add(newStr);
                found = true;
35         }
        }
        queueVisited.add(newStr);
        queue.offer(newStr);
    }
40 }

    if (result.size() == 0){
        result.add("");
45    }
    return result;
}

private boolean isValid(String str){
    int count = 0;
50    for (int idx = 0 ; idx < str.length() ; idx++){
        if (str.charAt(idx) == '('){
            count++;
        }
        if (str.charAt(idx) == ')'){
55            if (count == 0){
                return false;
            }
            count--;
        }
60    }

    return count == 0;
}
}

```

Solution - DFS

The idea is searched from Internet, copied here,

Key Points:

Generate unique answer once and only once, do not rely on Set.

Do not need preprocess.

Runtime 3 ms.

Explanation: We all know how to check a string of parentheses is valid using a stack. Or even simpler use a counter. The counter will increase when it is '(' and decrease when it is ')'. Whenever the counter is negative, we have more ')' than '(' in the prefix.

To make the prefix valid, we need to remove a ')'. The problem is: which one? The answer is any one in the prefix. However, if we remove any one, we will generate duplicate results, for example: s = ()), we can remove s[1] or s[2] but the result is the same (). Thus, we restrict ourself to remove the first) in a series of consecutive)s.

After the removal, the prefix is then valid. We then call the function recursively to solve the rest of the string. However, we need to keep another information: the last removal position. If we do not have this

position, we will generate duplicate by removing two ')' in two steps only with a different order. For this, we keep tracking the last removal position and only remove ')' after that.

Now one may ask. What about '('? What if s = '(()()' in which we need remove '('? The answer is: do the same from right to left. However a cleverer idea is: reverse the string and reuse the code! Here is the final implement in Java.

<https://leetcode.com/discuss/81478/easy-short-concise-and-fast-java-dfs-3-ms-solution>

Listing 315: Solution 2

```

public class Solution {
    public List<String> removeInvalidParentheses(String s) {
        List<String> result = new ArrayList<String>();
        dfs(s, 0, 0, new char[] { '(', ')' }, result);
5         return result;
    }

    private void dfs(String s, int prevI, int prevJ, char[] flag,
10         List<String> result){
        int halfCnt = 0;
        for (int i = prevI ; i < s.length() ; i++){
            if (s.charAt(i) == flag[0]){
                halfCnt ++;
            }
15         if (s.charAt(i) == flag[1]){
            halfCnt --;
        }
        if (halfCnt >= 0){
            continue;
20        }
        for (int j = prevJ ; j <= i ; j++){
            if (s.charAt(j) == flag[1]
                && (j == prevJ || flag[1] != s.charAt(j - 1))){
                dfs(s.substring(0, j) + s.substring(j + 1, s.length()),
25                 i, j, flag, result);
            }
        }
        // I don't quite understand why here is a 'return' statement
        return;
30    }

    String reversed = new StringBuilder(s).reverse().toString();
    if (flag[0] == '('){
        dfs(reversed, 0, 0, new char[] { ')', '(' }, result);
35    } else{
        result.add(reversed);
    }
}

```

Problem 242

201 Bitwise AND of Numbers Range

Given a range [m, n] where $0 \leq m \leq n \leq 2147483647$,
return the bitwise AND of all numbers in this range, inclusive.

For example, given the range [5, 7], you should return 4.

<https://leetcode.com/problems/bitwise-and-of-numbers-range/>

Solution

The problem is really interesting, it can finally be transformed to a problem to find common prefix of m and n in binary representation and fill 0 to the end to cover the position.

Listing 316: Solution 1.1

```
public class Solution {
    public int rangeBitwiseAnd(int m, int n) {
        int bit = 0;
        while(m != n) {
            m >>= 1;
            n >>= 1;
            bit ++;
        }
        return m << bit;
    }
}

// OR

public class Solution {
    public int rangeBitwiseAnd(int m, int n) {
        int mask = Integer.MAX_VALUE;
        while((mask & m) != (mask & n)) {
            mask = mask << 1;
        }
        return mask & m;
    }
}
```

Another solution is as below, which is in nature the same as previous method.

Listing 317: Solution 1.2

```
public class Solution {
    public int rangeBitwiseAnd(int m, int n) {
        while(n > m){
            n = n & (n - 1);
        }
        return n;
    }
}
```


Problem 243

151 Reverse Words in a String

Given an input string, reverse the string word by word.

For example,

Given s = "the sky is blue",
return "blue is sky the".

Update (2015-02-12):

For C programmers: Try to solve it in-place in O(1) space.

Clarification:

Q: What constitutes a word?

A: A sequence of non-space characters constitutes a word.

Q: Could the input string contain leading or trailing spaces?

A: Yes. However, your reversed string should not contain leading or trailing spaces.

Q: How about multiple spaces between two words?

A: Reduce them to a single space in the reversed string.

<https://leetcode.com/problems/reverse-words-in-a-string/>

Solution

The $O(n)$ solution (can only in C since C regards string as char array). Reverse each word, and then reverse the whole sentence.

Listing 318: Solution 1.1

```
public class Solution {  
    public String reverseWords(String s) {  
        s = reverse(s);  
        int curIdx = 0;  
        String revStr = "";  
        while (curIdx < s.length()) {  
            while (curIdx < s.length() && s.charAt(curIdx) == ' ') {  
                curIdx++;  
            }  
            if (curIdx == s.length()) {  
                break;  
            }  
            int startIdx = curIdx;  
            while (curIdx < s.length() && s.charAt(curIdx) != ' ') {  
                curIdx++;  
            }  
            int endIdx = curIdx;  
            String newPart = reverse(s.substring(startIdx, endIdx));  
            revStr += newPart;  
  
            while (curIdx < s.length() && s.charAt(curIdx) == ' ') {  
                curIdx++;  
            }  
        }  
    }  
}
```

```
25         if (curIdx < s.length()) {
            revStr += " ";
        }
    }
    return revStr;
}

30 public String reverse(String s) {
    char[] array = s.toCharArray();
    int start = 0, end = s.length() - 1;
    while (start < end) {
35         char tmp = array[start];
        array[start] = array[end];
        array[end] = tmp;
        start++;
        end--;
40     }
    return new String(array);
}
}
```

Well, in ‘Discuss’ section, solutions make use of built-in functions, which I think is not that proper....(But really concise and short).

Listing 319: Solution 1.2

```
public String reverseWords(String s) {
    String[] temp = s.split(" ");
    StringBuilder sb = new StringBuilder();
    for(int i = temp.length - 1; i >= 0; i --){
5        sb.append(temp[i].trim()+" ");
    }
    return sb.toString().trim();
}
```

Problem 244

128 Longest Consecutive Sequence

Given an unsorted array of integers,
find the length of the longest consecutive elements sequence.

For example,

Given [100, 4, 200, 1, 3, 2],

The longest consecutive elements sequence is [1, 2, 3, 4].

Return its length: 4.

Your algorithm should run in $O(n)$ complexity.

<https://leetcode.com/problems/longest-consecutive-sequence/>

Solution

Brute force method is to use Sorting methods to sort the array first, but with at least $O(n \log n)$,

For $O(n)$ solution, not so difficult, just use a Set to store all elements in the array, iterate each of them and find if there are ascending and descending elements to update the maximum consecutive length, and remove them from Set to avoid duplicate calculation.

Listing 320: Solution

```
public class Solution {
    public int longestConsecutive(int[] nums) {
        if (nums.length == 0)
            return 1;
5       Set<Integer> set = new HashSet<Integer>();
        for (int num : nums)
            set.add(num);
        int maxLen = 0;
        for (int num : nums){
10         if (!set.contains(num))
            continue;
            set.remove(num);
            int len = 1;
            int cur = num;
15         while (set.contains(--cur)){
            len ++;
            set.remove(cur);
        }
            cur = num;
20         while (set.contains(++cur)){
            len ++;
            set.remove(cur);
        }
            maxLen = Math.max(maxLen , len );
25     }
        return maxLen;
    }
}
```

Problem 245

312 Burst Balloons

Given n balloons, indexed from 0 to $n-1$.
 Each balloon is painted with a number on it represented by array `nums`.
 You are asked to burst all the balloons.
 If the you burst balloon i you will get
 $\text{nums}[\text{left}] * \text{nums}[i] * \text{nums}[\text{right}]$ coins.
 Here `left` and `right` are adjacent indices of i .
 After the burst, the `left` and `right` then becomes adjacent.

Find the maximum coins you can collect by bursting the balloons wisely.

Note:

- (1) You may imagine $\text{nums}[-1] = \text{nums}[n] = 1$.
 They are not real therefore you can not burst them.
- (2) $0 \leq n \leq 500$, $0 \leq \text{nums}[i] \leq 100$

Example:

Given `[3, 1, 5, 8]`

Return 167

```

nums = [3,1,5,8] --> [3,5,8] --> [3,8] --> [8] --> []
coins = 3*1*5      + 3*5*8      + 1*3*8      + 1*8*1      = 167

```

<https://leetcode.com/problems/burst-balloons/>

Solution

A excellent DP problem.

The state transition equation is $\text{dp}[l][r] = \text{nums}[l] * \text{nums}[m] * \text{nums}[r] + \text{dp}[l][m] + \text{dp}[m][r]$, $l < m < r$;

For detailed solution description, see

<http://bookshadow.com/weblog/2015/11/30/leetcode-burst-balloons/>

Listing 321: Solution

```

public class Solution {
    public int maxCoins(int[] nums) {
        int[] totalNums = new int[nums.length + 2];
        totalNums[0] = totalNums[totalNums.length - 1] = 1;
5      for (int i = 0 ; i < nums.length ; i++){
            totalNums[i + 1] = nums[i];
        }
        if (nums.length == 0){
            return 0;
10      }
        int dp[][] = new int[totalNums.length][totalNums.length];
        for (int len = 2 ; len < totalNums.length ; len++){
            for (int start = 0 ; start < totalNums.length - len ; start++){
                int end = start + len;
15      for (int mid = start + 1 ; mid < end ; mid++){

```

```
                dp[start][end] = Math.max(dp[start][end],
                totalNums[start] * totalNums[mid] * totalNums[end]
                + dp[start][mid] + dp[mid][end]);
            }
        }
    }

    return dp[0][totalNums.length - 1];
}
```

Problem 246

330 Patching Array

Given a sorted positive integer array `nums` and an integer `n`, add/patch elements to the array such that any number in range `[1, n]` inclusive can be formed by the sum of some elements in the array.
Return the minimum number of patches required.

Example 1:

`nums = [1, 3], n = 6`
Return 1.

Combinations of `nums` are `[1], [3], [1,3]`,
which form possible sums of: 1, 3, 4.

Now if we add/patch 2 to `nums`, the combinations are:

`[1], [2], [3], [1,3], [2,3], [1,2,3]`.

Possible sums are 1, 2, 3, 4, 5, 6,

which now covers the range `[1, 6]`.

So we only need 1 patch.

Example 2:

`nums = [1, 5, 10], n = 20`
Return 2.

The two patches can be `[2, 4]`.

Example 3:

`nums = [1, 2, 2], n = 5`
Return 0.

<https://leetcode.com/problems/patching-array/>

Solution

The idea is not so easy to figure out.

First, we sort the list from smallest to largest.

Then iterate the list(or iterate through 1 to required max-sum).

We maintain a maximum sum we can reach currently.

When we meet `nums[i] <= sum`, eg: `[1,2,10]`, `curSum = 0` at first. `1 <= 0+1`, then `curSum += 1`; then `2 <= 1+1`, `curSum += 2`;

When `nums[i] > sum`, eg, `curSum=4` & `nums[i]=10` in previous example, now we cannot reach sum more than `curSum`, we then add `curSum+1` to our list to get more conditions for sum. What is the maximum sum we can reach now? The answer is `curSum*2+1(2*(curSum+1) - 1)`.

Note that a common pitfall is that sum may exceed max int, so we use long instead of int. (Consider the test case `[1,2,31,33]`, 2147483647)

Listing 322: Solution

```
public class Solution {
    public int minPatches(int[] nums, int n) {
        int count = 0;
        // Arrays.sort(nums);
        long maxSum = 0;
```

```

    for (int i = 0 ; i < nums.length ; i++){
        if (nums[i] <= maxSum + 1){
            maxSum += nums[i];
        }
10     else{
        while (nums[i] > maxSum + 1){
            count ++;
            maxSum = 2 * maxSum + 1;

15         if (maxSum >= n){
            return count;
        }

        if (nums[i] <= maxSum + 1){
            maxSum += nums[i];
20         }
        }

        if (maxSum >= n){
            return count;
25         }
    }
    if (maxSum == 0){
        maxSum = 1;
        count ++;
30    }
    while (maxSum < n){
        maxSum = 2 * maxSum + 1;
        count++;
    }
35    return count;
}
}

// A more concise version
40 class Solution {
public:
    int minPatches(vector<int>& nums, int n) {
        int cnt = 0, i = 0;
        for (long known_sum = 1; known_sum <= n; ) {
45         if (i < nums.size() && nums[i] <= known_sum) {
            known_sum += nums[i++];
        }
        else {
            known_sum <<= 1;
50         cnt++;
        }
    }
    return cnt;
}
55 };

```

Problem 247

306 Additive Number

Additive number is a string whose digits can form additive sequence.

A valid additive sequence should contain at least three numbers.

Except for the first two numbers,
each subsequent number in the sequence must be the sum of the preceding two.

For example:

"112358" is an additive number because the digits can
form an additive sequence: 1, 1, 2, 3, 5, 8.

$1 + 1 = 2$, $1 + 2 = 3$, $2 + 3 = 5$, $3 + 5 = 8$

"199100199" is also an additive number, the additive sequence is:

1, 99, 100, 199.

$1 + 99 = 100$, $99 + 100 = 199$

Note: Numbers in the additive sequence cannot have leading zeros,
so sequence 1, 2, 03 or 1, 02, 3 is invalid.

Given a string containing only digits '0'-'9',
write a function to determine if it's an additive number.

Follow up:

How would you handle overflow for very large input integers?

<https://leetcode.com/problems/additive-number/>

Solution

Brute force, iterate to get 2 starting adjacent strings, add them until we find(or can't find) correct answer.

But, except for using 'long' type, this solution is not able to handle too large integer.(Or maybe BigInteger or applying addition on long number sequence?)

Listing 323: Solution

```
public class Solution {
    public boolean isAdditiveNumber(String num) {
        for(int i = 1 ; i < num.length() ; i++){
            for (int j = i + 1 ; j <= num.length() ; j++){
                String str1 = num.substring(0, i);
                String str2 = num.substring(i, j);
                if (str1.length() > 1 && str1.charAt(0) == '0'){
                    continue;
                }
                if (str2.length() > 1 && str2.charAt(0) == '0'){
                    continue;
                }
                if (str2.equals("")){
                    continue;
                }
                long num1 = Long.parseLong(str1);
```



```
20      long num2 = Long.parseLong(str2);
      long sum = num1 + num2;
      String prefix = str1 + str2;
      if (prefix.equals(num)){
          continue;
      }

25      while (prefix.length() < num.length()){
          prefix += sum;
          if (!num.startsWith(prefix)
              || prefix.length() > num.length()){
              break;
          }
30          num1 = num2;
          num2 = sum;
          sum = num1 + num2;
      }
      if (num.equals(prefix)){
          return true;
35      }
    }
  }
  return false;
40 }
}
```

Problem 248

56 Merge Intervals

Given a collection of intervals,
merge all overlapping intervals.

For example,

Given `[1,3],[2,6],[8,10],[15,18]`,
return `[1,6],[8,10],[15,18]`.

<https://leetcode.com/problems/merge-intervals/>

Solution

Sort all intervals according to their start number(write a new Comparator). Then there are 3(or 2, to be more specific) conditions, eg, `[1,100]` VS `[90,101]` or `[90,99]` with `[1,10]` VS `[90,100]`.

For those `cur.end >= next.start`, combine the 2 intervals, set `cur.end = Math.max(next.end, cur.end)`; for opposite conditions, set `cur = new Node(next.start, next.end)`;

Listing 324: Solution

```

/**
 * Definition for an interval.
 * public class Interval {
 *     int start;
 *     int end;
 *     Interval() { start = 0; end = 0; }
 *     Interval(int s, int e) { start = s; end = e; }
 * }
 */
10 public class Solution {
    public List<Interval> merge(List<Interval> intervals) {
        if (intervals.size() <= 1){
            return intervals;
        }
15    List<Interval> result = new ArrayList<Interval>();
    Collections.sort(intervals, new Comparator<Interval>(){
        @Override
        public int compare(Interval i1, Interval i2){
            return i1.start - i2.start;
20        }
    });
    Interval cur = new Interval(intervals.get(0).start, intervals.get(0).end);
    for (int idx = 1 ; idx < intervals.size() ; idx++){
        Interval interval = intervals.get(idx);
25        if (cur.end >= interval.start){
            cur.end = Math.max(cur.end, interval.end);
        }
        else{
            result.add(cur);
30            cur = new Interval(interval.start, interval.end);
        }
    }
}

```

```
35         result.add(cur);  
           return result;  
    }
```

Problem 249

57 Insert Interval

Given a set of non-overlapping intervals,
insert a new interval into the intervals (merge if necessary).

You may assume that the intervals were initially sorted
according to their start times.

Example 1:

Given intervals [1,3],[6,9], insert and merge [2,5] in as [1,5],[6,9].

Example 2:

Given [1,2],[3,5],[6,7],[8,10],[12,16],
insert and merge [4,9] in as [1,2],[3,10],[12,16].

This is because the new interval [4,9]
overlaps with [3,5],[6,7],[8,10].

<https://leetcode.com/problems/insert-interval/>

Solution

Iterate every element in former List. Three conditions are waiting:

1. $\text{cur.end} < \text{new.start}$, add to list, continue
2. $\text{cur.start} > \text{new.end}$, break, this is where we should insert our new Interval(maybe not the initial one)
3. cur interleaves with new, set $\text{new.start} = \min(\text{cur.start}, \text{new.start})$, $\text{new.end} = \max(\text{cur.end}, \text{new.end})$, continue;

Finally don't forget to add **new** and those Intervals after **new** ($\text{cur.start} > \text{new.end}$) to end of result list.
(BTW, there is another solution, first use binary search to decide where to insert the new Interval, then use the same idea in Problem "Merge Intervals", to re-merge the list.)

Listing 325: Solution

```
/**
 * Definition for an interval.
 * public class Interval {
 *     int start;
5  *     int end;
 *     Interval() { start = 0; end = 0; }
 *     Interval(int s, int e) { start = s; end = e; }
 * }
 */
10 public class Solution {
    public List<Interval> insert(List<Interval> intervals, Interval newInterval){
        List<Interval> result = new ArrayList<Interval>();
        if (intervals.size() == 0){
            result.add(newInterval);
15    return result;
    }
```

```
    }
    Interval compare = new Interval(newInterval.start , newInterval.end);
    int endIdx = -1;
    for (int i = 0 ; i < intervals.size() ; i ++){
20         Interval cur = intervals.get(i);
        if (cur.end < compare.start){
            result.add(cur);
        }
        else if (cur.start > compare.end){
25             endIdx = i;
            break;
        }
        else{
            compare.start = Math.min(compare.start , cur.start);
            compare.end = Math.max(compare.end , cur.end);
30         }
    }
    result.add(compare);
    if (endIdx != -1){
35         for (int i = endIdx ; i < intervals.size() ; i ++){
            result.add(intervals.get(i));
        }
    }
    return result;
40 }
}
```

Problem 250

87 Scramble String

Given a string `s1`, we may represent it as a binary tree by partitioning it to two non-empty substrings recursively.

Below is one possible representation of `s1 = "great"`:

```

      great
     /   \
    gr    eat
   / \   / \
  g  r e at
           / \
          a  t

```

To scramble the string, we may choose any non-leaf node and swap its two children.

For example, if we choose the node "gr" and swap its two children, it produces a scrambled string "rgeat".

```

      rgeat
     /   \
    rg    eat
   / \   / \
  r  g e at
           / \
          a  t

```

We say that "rgeat" is a scrambled string of "great".

Similarly, if we continue to swap the children of nodes "eat" and "at", it produces a scrambled string "rgtae".

```

      rgtae
     /   \
    rg    tae
   / \   / \
  r  g ta e
           / \
          t  a

```

We say that "rgtae" is a scrambled string of "great".

Given two strings `s1` and `s2` of the same length, determine if `s2` is a scrambled string of `s1`.

<https://leetcode.com/problems/scramble-string/>

Solution - Recursion

First the recursive solution. Still use an example to help understanding.

For `s1 = 'great'` and `s2 = 'rgtae'`. We split each string to 2 parts, the split point can be every letter in word. When we split `s1` to 'gr' and 'eat', while `s2` to 'rg' and 'tae' (along with 'ae' and 'rgt'), the scramble can only be possible to take place between 'gr' and 'rg' (or maybe 'ae', recurse along this idea until the 2

words equal ,length not equals, or letters they contain are not the same(aab and bab/c, etc). If we don't use the last condition to prune, we will get TLE(using bit operation neither, see comments in code).

Listing 326: Solution 1

```

public class Solution {
    public boolean isScramble(String s1, String s2) {
        if (s1.length() != s2.length()){
            return false;
        }
        if (s1.equals(s2)){
            return true;
        }
        // former solution, much more slower than saved array
        // int num1 = 0, num2 = 0;
        // for (int i = 0 ; i < s1.length() ; i++){
        //     num1 |= (1 << (s1.charAt(i) - 'a'));
        //     num2 |= (1 << (s2.charAt(i) - 'a'));
        // }
        // if (num1 != num2){
        //     return false;
        // }
        int[] saved = new int[26];
        for(int idx = 0 ; idx < s1.length() ; idx++){
            saved[s1.charAt(idx) - 'a'] += 1;
            saved[s2.charAt(idx) - 'a'] -= 1;
        }
        for(int i = 0 ; i < 26 ; i++){
            if(saved[i] != 0) return false;
        }

        boolean result = false;
        for (int i = 1 ; i < s1.length() ; i++){
            String left1 = s1.substring(0, i), right1 = s1.substring(i);
            String left2 = s2.substring(0, i), right2 = s2.substring(i);
            String left2Rev = s2.substring(0, s1.length() - i),
                right2Rev = s2.substring(s1.length() - i);
            result |= (isScramble(left1, left2) && isScramble(right1, right2))
                || (isScramble(left1, right2Rev) && isScramble(right1, left2Rev));
            if (result == true){
                return result;
            }
        }
        return result;
    }
}

```

Solution - DP

Second solution makes use of Dynamic Programming, with the state transition equation:

The idea is search from the Internet. For $dp[k][i][j]$, it means whether $s1[i \text{ to } i+k-1]$ equals to $s2[j \text{ to } j+k-1]$

See <http://www.acmerblog.com/leetcode-solution-scramble-string-6224.html>

Listing 327: Solution 2

```
public class Solution {
    public boolean isScramble(String s1, String s2) {
        boolean dp[][][] = new boolean[s1.length() + 1][s1.length()][s1.length()];
        for (int i = 0 ; i < s1.length() ; i++){
            for (int j = 0 ; j < s1.length() ; j++){
                dp[1][i][j] = (s1.charAt(i) == s2.charAt(j));
            }
        }

        for (int k = 1 ; k <= s1.length() ; k++){
            for (int i = 0 ; i <= s1.length() - k ; i++){
                for (int j = 0 ; j <= s1.length() - k ; j++){
                    for (int inner = 1 ; inner < k ; inner++){
                        dp[k][i][j] |= dp[inner][i][j]
                                && dp[k - inner][i + inner][j + inner];
                        dp[k][i][j] |= dp[inner][i][j + k - inner]
                                && dp[k - inner][i + inner][j];
                        if (dp[k][i][j]){
                            break;
                        }
                    }
                }
            }
        }

        return dp[s1.length()][0][0];
    }
}
```


Problem 251

315 Count of Smaller Numbers After Self

You are given an integer array `nums` and you have to return a new counts array.

The counts array has the property where `counts[i]` is the number of smaller elements to the right of `nums[i]`.

Example:

Given `nums = [5, 2, 6, 1]`

To the right of 5 there are 2 smaller elements (2 and 1).

To the right of 2 there is only 1 smaller element (1).

To the right of 6 there is 1 smaller element (1).

To the right of 1 there is 0 smaller element.

Return the array `[2, 1, 1, 0]`.

<https://leetcode.com/problems/count-of-smaller-numbers-after-self/>

Solution - Merge Sort

Merge sort is a typical solution on finding inversions. On Book “Offer”, there is a similar problem of finding all inversions in an array which uses merge sort too.

An Example, For `[2,4],[1,3]`, we first compare 4 & 3, when 4 is bigger, smaller count of 4 is 2(1 and 3, `index(3) - index(1) + 1`), similarly, smaller count of 2 is 1...

Listing 328: Solution 1

```
class Number{
    int num;
    int index;

    public Number(int num, int index){
        this.num = num;
        this.index = index;
    }
}

public class Solution {
    public List<Integer> countSmaller(int[] nums) {
        int[] count = new int[nums.length];
        Number[] numbers= new Number[nums.length];
        for (int i = 0 ; i < nums.length ; i++){
            Number number = new Number(nums[i] , i);
            numbers[i] = number;
        }
        mergeSort(numbers, 0, numbers.length - 1, count);

        List<Integer> result = new ArrayList<Integer>();
        for (int c : count){
            result.add(c);
        }
    }
}
```

```

25     return result;
    }

    private void mergeSort(Number[] numbers, int start, int end, int[] count){
        if (start >= end){
30             return;
        }
        int mid = start + (end - start) / 2;
        mergeSort(numbers, start, mid, count);
        mergeSort(numbers, mid + 1, end, count);
35     mergeCombine(numbers, start, mid, mid + 1, end, count);
    }

    private void mergeCombine(Number[] numbers, int start1,
        int end1, int start2, int end2, int[] count){
        Number[] result = new Number[end2 - start1 + 1];
40     int idx = result.length - 1;
        int idxLeft = end1, idxRight = end2;
        while(idxLeft >= start1 && idxRight >= start2){
            if (numbers[idxRight].num >= numbers[idxLeft].num){
45                 result[idx --] = numbers[idxRight --];
            }
            else{
                count[numbers[idxLeft].index] += idxRight - start2 + 1;
                result[idx --] = numbers[idxLeft --];
            }
50     }
        while (idxLeft >= start1){
            count[numbers[idxLeft].index] += idxRight - start2 + 1;
            result[idx --] = numbers[idxLeft --];
        }
55     while (idxRight >= start2){
            result[idx --] = numbers[idxRight --];
        }

        for (int i = 0 ; i < result.length ; i++){
60             numbers[start1 + i] = result[i];
        }
    }
}

```

Solution - BIT

For BIT, we firstly sort and discrete the original array to make array in BIT much shorter. Then we assign each element a index according to its position after sorted, eg: for [5,2,6,1], there will be $1 - > 1, 2 - > 2, 3 - > 5, 4 - > 6$.

Then it's time for BIT to take effect. We iterate from tail of original array, now sum of discrete[nums[i]] is count of its smaller elements, and we may then update the BIT since we have a new element counting from right.

Since we are counting 'smaller' rather than 'not bigger', we have to minus 1 to discrete[num[i]] when doing the sum operation, which is a bit forgettable.

Listing 329: Solution 2

```
class BITree{
    int capacity;
    int[] elements;

5     private int lowbit(int num){
        return num & (-num);
    }

    public BITree(int capacity){
10        this.capacity = capacity;
        elements = new int[capacity + 1];
    }

    public void add(int index, int value){
15        while (index <= capacity){
            elements[index] += value;
            index += lowbit(index);
            if (index == 0){
                break;
20            }
        }
    }

    public int sum(int index){
        int sum = 0;
25        while (index > 0){
            sum += elements[index];
            index -= lowbit(index);
        }
        return sum;
30    }
}

public class Solution {
    public List<Integer> countSmaller(int[] nums){
        // Sort and Discretize
35        int[] clone = new int[nums.length];
        for (int numIdx = 0 ; numIdx < nums.length ; numIdx++){
            clone[numIdx] = nums[numIdx];
        }
        // Sorting
40        Arrays.sort(clone);
        // Discretization
        Map<Integer, Integer> discreteMap = new HashMap<Integer, Integer>();
        int count = 0;
        for (int dIdx = 0 ; dIdx < clone.length ; dIdx++){
45            if (dIdx > 0 && clone[dIdx] == clone[dIdx - 1]){
                continue;
            }
            // discrete index is from 1~n not 0~n
            // since tree.elements[0] is
50            // to store the count less than minimum value in array
            discreteMap.put(clone[dIdx], ++count);
        }
        Integer[] result = new Integer[nums.length];
```

```
55     BITree tree = new BITree(discreteMap.keySet().size());
    for (int numIdx = nums.length - 1 ; numIdx >= 0 ; numIdx --){
        // here -1 means finding smaller elements
        result[numIdx] = tree.sum(discreteMap.get(nums[numIdx]) - 1);
        tree.add(discreteMap.get(nums[numIdx]) , 1);
    }
60     return new ArrayList<Integer>(Arrays.asList(result));
}
}
```

There is also a BST solution in

<http://whatsme.net/2015/12/16/Count-of-Smaller-Numbers-After-Self/>

Segment tree:

<https://leetcode.com/discuss/73917/accepted-c-solution-using-segment-tree>

Divide and Conquer:

<https://leetcode.com/discuss/73326/divide-and-conquer-like-merge-sort-order-statistics-tree>

Problem 252

327 Count of Range Sum

Given an integer array `nums`,
return the number of range sums that lie in `[lower, upper]` inclusive.
Range sum `S(i, j)` is defined as
the sum of the elements in `nums` between indices `i` and `j` (`i ≤ j`), inclusive.

Note:

A naive algorithm of $O(n^2)$ is trivial.
You MUST do better than that.

Example:

Given `nums = [-2, 5, -1]`, `lower = -2`, `upper = 2`,
Return 3.
The three ranges are : `[0, 0]`, `[2, 2]`, `[0, 2]`
and their respective sums are: `-2`, `-1`, `2`.

<https://leetcode.com/problems/count-of-range-sum/>

Solution - Merge Sort

The Merge sort train of thought is a bit like Problem “Count of Smaller Numbers After Self”. We first sum up `[0,x]` elements to get a `sums` array with the length of `n` (which means, to be more specific, `sums[0] = nums[0]`, `sums[1] = nums[0] + [nums[1]`, ..., `sums[n] = nums[0] + nums[1] + nums[2] + ...`).

Then apply Merge Sort among `sums` array. During merging, count the sums satisfy the `[lower, upper]` constraint. How?

Let's see what we are looking for. What is the value of sum from index `i` to `j`? `sums[i to j] = sums[j] - sums[i - 1]`. So what we do is to check every possible index pair `(i,j)` to check if `sums[i to j]` satisfy the constraint.

So why Merge Sort?(Honestly, because of its $O(\log n)$ runtime complexity.) And how? Let example talk.

Suppose we are doing merge sort at a circumstance where the `sums` array looks like `[(2,5,9),(6,8,11)]` with `upper=7` and `lower=5`, then we may combine the array to `[2,5,6,8,9,11]`... Wait! We have something else to do. When checking 2, we should also check 6,8 and maybe 11. See `6-2=4 < 5`, continue when `8-2=6 > 5`, stop and mark `lowerIdx=idx(8)`, then `11-2 > 7`, mark `upperIdx=idx(11)`, then we add count 1 to `totalCnt`.

Namely, the merge sort based solution counts the answer while doing the merge. During the merge stage, we have already sorted the left half `[start, mid)` and right half `[mid, end)`. We then iterate through the left half with index `i`. For each `i`, we need to find two indices `k` and `j` in the right half where

`j` is the first index satisfy `sums[j] - sums[i] > upper` and

`k` is the first index satisfy `sums[k] - sums[i] >= lower`.

Then the number of sums in `[lower, upper]` is `j-k`. We also use another index `t` to copy the elements satisfy `sums[t] < sums[i]` to a cache in order to complete the merge sort.

Finally, don't forget to use long type since the last test data will be out of bound.

Listing 330: Solution 1

```
public class Solution {
    public int countRangeSum(int[] nums, int lower, int upper) {
        if (nums.length == 0){
            return 0;
        }
        long[] sums = new long[nums.length];
```

```

    sums[0] = nums[0];
    int count = 0;
    for (int i = 1 ; i < nums.length ; i++){
10         sums[i] = sums[i - 1] + nums[i];
    }
    //      System.out.println("Initial state");
    for (int i = 0 ; i < nums.length ; i++){
15 //      System.out.print(nums[i] + " ");
        if (sums[i] >= lower && sums[i] <= upper){
            count ++;
        }
    }
20 //      System.out.println();
    //      System.out.println("After Transformation.");
    return count + mergeSort(sums, 0, sums.length - 1, lower, upper);
}

25 private int mergeSort(long[] sums, int startIdx,
    int endIdx, long lower, long upper){
    if (startIdx >= endIdx){
        return 0;
    }
30 int midIdx = startIdx + (endIdx - startIdx) / 2;

    int count = mergeSort(sums, startIdx, midIdx, lower, upper)
        + mergeSort(sums, midIdx + 1, endIdx, lower, upper);
    int lowerIdx = midIdx + 1, upperIdx = midIdx + 1;
35 int mergeIdx = midIdx + 1;
    int cacheIdx = 0;
    long cache[] = new long[endIdx - startIdx + 1];
    for (int leftIdx = startIdx ; leftIdx <= midIdx ; leftIdx++){
        while (lowerIdx <= endIdx
40             && sums[lowerIdx] - sums[leftIdx] < lower){
            lowerIdx ++;
        }
        upperIdx = lowerIdx;
        while (upperIdx <= endIdx
45             && sums[upperIdx] - sums[leftIdx] <= upper){
            upperIdx ++;
        }
        while (mergeIdx <= endIdx
            && sums[mergeIdx] < sums[leftIdx]){
50             cache[cacheIdx ++] = sums[mergeIdx ++];
        }
        count += upperIdx - lowerIdx;
        cache[cacheIdx ++] = sums[leftIdx];
    }
55 while (mergeIdx <= endIdx){
    cache[cacheIdx ++] = sums[mergeIdx ++];
}
    for (int recoverIdx = 0 ;
        recoverIdx < endIdx - startIdx + 1 ; recoverIdx++){
```

```

60         sums[recoverIdx + startIdx] = cache[recoverIdx];
        }
        //         for (int sum : sums){
        //             System.out.print(sum + " ");
        //         }
65     //         System.out.println();
        return count;
    }
}

```

Solution - BST

I didn't understand the solution at first sight, but honestly this is more intuitive. We are looking for a range sum from 'lower' to 'upper', so when we consider the 'sums' array in previous solution, we are looking for a $lower \leq sums[j] - sums[i - 1] \leq upper$. So we build a BST to store each sums, before inserting a new sums[i], we check in the tree that is there any previous sums X which satisfy $X \geq sums[j] - upper$ and $X \leq sums[j] - lower$. We used an exclusive method to get final result.

Listing 331: Solution 2

```

class BSTNode{
    BSTNode left;
    BSTNode right;
    int count;
5   int leftCnt;
    int rightCnt;
    long val;
    public BSTNode(long val){
        this.count = 1;
10        this.val = val;
    }
}

public class Solution {
    private int countLower(BSTNode node, long val){
15        if (node == null){
            return 0;
        }
        if (node.val == val){
            return node.leftCnt;
20        }
        else if (node.val > val){
            return countLower(node.left, val);
        }
        else{
25            return node.leftCnt + node.count + countLower(node.right, val);
        }
    }

    private int countUpper(BSTNode node, long val){
30        if (node == null){
            return 0;
        }
        if (node.val == val){
            return node.rightCnt;

```

```

35     }
    else if (node.val < val){
        return countUpper(node.right , val);
    }
    else {
40         return node.rightCnt + node.count + countUpper(node.left , val);
    }
}
private BSTNode insert(BSTNode node, long val){
    if (node == null){
45         return new BSTNode(val);
    }
    if (node.val == val){
        node.count ++;
        return node;
50    }
    else if (node.val > val) {
        node.leftCnt ++;
        node.left = insert(node.left , val);
    }
55    else{
        node.rightCnt ++;
        node.right = insert(node.right , val);
    }
    return node;
60 }

private int countInRange(BSTNode root, long lower, long upper){
    int total = root.count + root.leftCnt + root.rightCnt;
    int lowerCnt = countLower(root, lower);
65    int upperCnt = countUpper(root, upper);
    return total - lowerCnt - upperCnt;
}

public int countRangeSum(int[] nums, int lower, int upper) {
70    long sums[] = new long[nums.length + 1];
    for (int i = 0 ; i < nums.length ; i++){
        sums[i + 1] = sums[i] + nums[i];
    }

75    BSTNode root = new BSTNode(sums[0]);
    int total = 0;
    for (int i = 1 ; i < sums.length ; i++){
        total += countInRange(root, sums[i] - upper, sums[i] - lower);
        insert(root, sums[i]);
80    }

    return total > 0 ? total : 0;
}
}

```

Solution - Others(BIT)

Honestly, I don't quite understand the insight of this solution, just paste here the code searched from

Discuss part for further usage.

Listing 332: Solution 3

```
public class Solution {
    public int countRangeSum(int[] nums, int lower, int upper) {
        List<Long> cand = new ArrayList<>();
        cand.add(Long.MIN_VALUE); // make sure no number gets a 0-index.
5        cand.add(0L);
        long[] sum = new long[nums.length + 1];
        for (int i = 1; i < sum.length; i++) {
            sum[i] = sum[i - 1] + nums[i - 1];
            cand.add(sum[i]);
10        cand.add(lower + sum[i - 1] - 1);
            cand.add(upper + sum[i - 1]);
        }
        Collections.sort(cand); // finish discretization

15        int[] bit = new int[cand.size()];
        for (int i = 0; i < sum.length; i++) plus(bit,
            Collections.binarySearch(cand, sum[i]), 1);
        int ans = 0;
        for (int i = 1; i < sum.length; i++) {
20            plus(bit, Collections.binarySearch(cand, sum[i - 1]), -1);
            ans += query(bit, Collections.binarySearch(cand,
                upper + sum[i - 1]));
            ans -= query(bit, Collections.binarySearch(cand,
                lower + sum[i - 1] - 1));
25        }
        return ans;
    }

    private void plus(int[] bit, int i, int delta) {
30        for (; i < bit.length; i += i & -i) bit[i] += delta;
    }

    private int query(int[] bit, int i) {
        int sum = 0;
35        for (; i > 0; i -= i & -i) sum += bit[i];
        return sum;
    }
}
```

Problem 253

30 Substring with Concatenation of All Words

You are given a string, 's', and a list of words, 'words', that are all of the same length.
Find all starting indices of substring(s) in s that is a concatenation of each word in words exactly once and without any intervening characters.

For example, given:

s: "barfoothefoobarman"
words: ["foo", "bar"]

You should return the indices: [0,9].
(order does not matter).

<https://leetcode.com/problems/substring-with-concatenation-of-all-words/>

Solution

Just a Brute force method, the condition that all words are of the same length makes the problem much easier. Keep a given length of sequence and check every word in it if they match our given **words** array.

Listing 333: Solution

```
public class Solution {  
    public List<Integer> findSubstring(String s, String[] words) {  
        List<Integer> result = new ArrayList<Integer>();  
        if (words.length == 0){  
5           return new ArrayList<Integer>();  
        }  
        Map<String, Integer> mapping = new HashMap<String, Integer>();  
        Map<String, Integer> current = new HashMap<String, Integer>();  
        // words in 'words' array may be duplicate  
10       for (String word : words){  
            if (mapping.containsKey(word)){  
                mapping.put(word, mapping.get(word) + 1);  
            }  
            else{  
15                mapping.put(word, 1);  
            }  
        }  
        int wordCnt = words.length;  
        int wordLen = words[0].length();  
20       // Since words are of the same length, we split each substrings.  
        for (int i = 0 ; i <= s.length() - wordCnt * wordLen; i++){  
            current.clear();  
            int j = 0;  
            for (j = 0 ; j < wordCnt ; j++){  
25                String substring = s.substring(i + j * wordLen,  
                    i + (j + 1) * wordLen);  
                if (!mapping.containsKey(substring)){  
                    break;  
                }  
            }  
        }  
    }  
}
```

```
30         if (current.containsKey(substring)){
           current.put(substring, current.get(substring) + 1);
        }
        else{
           current.put(substring, 1);
35        }
        if (current.get(substring) > mapping.get(substring)){
           break;
        }
    }
40    if (j == wordCnt){
        result.add(i);
    }
}
45    return result;
}
}
```

Problem 254

149 Max Points on a Line

Given n points on a 2D plane,
find the maximum number of points that lie on the same straight line.

<https://leetcode.com/problems/max-points-on-a-line/>

Solution

Main idea is to use the slope to form a map to check the number of points on same line. Iterate both i and j from 0 to n(which I think has lots of duplicate calculation).

Other difficulties are caused by corner cases: points with same x; same points; when all points are in the same position...

Listing 334: Solution

```
/**
 * Definition for a point.
 * class Point {
 *     int x;
 *     int y;
 *     Point() { x = 0; y = 0; }
 *     Point(int a, int b) { x = a; y = b; }
 * }
 */
10 public class Solution {
    public int maxPoints(Point[] points) {
        if (points.length <= 1){
            return points.length;
        }
15    Map<Double, Integer> slopeCnt = new HashMap<Double, Integer>();
    int max = 0;
    for (int i = 0 ; i < points.length ; i ++){
        slopeCnt.clear();
        Point pointX = points[i];
20        int duplicate = 0;
        for (int j = 0 ; j < points.length ; j++){
            if (i == j){
                continue;
            }
25            Point pointY = points[j];
            if (pointX.x == pointY.x && pointX.y == pointY.y){
                duplicate ++;
                continue;
            }
30            else if (pointX.x == pointY.x){
                slopeCnt.put(Double.MAX_VALUE,
                    slopeCnt.containsKey(Double.MAX_VALUE)
                        ? slopeCnt.get(Double.MAX_VALUE) + 1 : 1);
            }
35            else{
                double slope =
```

```

        (pointY.y - pointX.y) * 1.0 / (pointY.x - pointX.x);
        slopeCnt.put(slope, slopeCnt.containsKey(slope)
            ? slopeCnt.get(slope) + 1 : 1);
40    }
    }
    // System.out.println("Current pos: " + pointX.x + " , " + pointX.y);
    for (Double key : slopeCnt.keySet()) {
        // System.out.println("\t" + key + " : " + slopeCnt.get(key));
45    max = Math.max(max, slopeCnt.get(key) + duplicate + 1);
    }
    if (slopeCnt.keySet().size() == 0) {
        max = Math.max(max, duplicate + 1);
    }
50 }

    return max;
}
}
```

Problem 255

115 Distinct Subsequences

Given a string S and a string T,
count the number of distinct subsequences of T in S.

A subsequence of a string is a new string
which is formed from the original string by deleting some
(can be none) of the characters without disturbing
the relative positions of the remaining characters.
(ie, "ACE" is a subsequence of "ABCDE" while "AEC" is not).

Here is an example:

S = "rabbbit", T = "rabbit"

Return 3.

<https://leetcode.com/problems/distinct-subsequences/>

Solution

Honestly, this is not a difficult DP problem. The state transition equation is simple:

Let $dp[i][j]$ represents match count of first i th letters in S with first j th letters in T.

When $s.charAt(i-1) == t.charAt(j-1)$, then $dp[i][j] = dp[i-1][j-1] + dp[i-1][j]$; Eg, for 2nd a in rrraa(t is ra), dp equals 3 + 3, where first 3 represents how many 'ra' the first 'rrra' can produce; where last 3 represents how many can 'rrrXa(current a)' get.

While when $s.charAt(i-1) != t.charAt(j-1)$, $dp[i][j] = dp[i-1][j]$;

Listing 335: Solution

```

public class Solution {
    public int numDistinct(String s, String t) {
        int lenS = s.length();
        int lenT = t.length();
5       int [][] dp = new int[s.length() + 1][t.length() + 1];
        for (int i = 0 ; i <= lenS ; i++){
            dp[i][0] = 1;
        }
        for (int j = 1 ; j <= lenT ; j++){
10         for (int i = 1 ; i <= lenS ; i++){
            if (s.charAt(i - 1) == t.charAt(j - 1)){
                dp[i][j] = dp[i - 1][j] + dp[i - 1][j - 1];
            }
            else{
15                 dp[i][j] = dp[i - 1][j];
            }
        }
    }
    return dp[lenS][lenT];
20 }
}

```

Problem 256

97 Interleaving String

Given s_1 , s_2 , s_3 , find whether s_3 is formed by the interleaving of s_1 and s_2 .

For example,

Given:

$s_1 = \text{"aabcc"}$,

$s_2 = \text{"dbbca"}$,

When $s_3 = \text{"aadbcbcbac"}$, return true.

When $s_3 = \text{"aadbcbaccc"}$, return false.

<https://leetcode.com/problems/interleaving-string/>

Solution

Another DP problem involving String issues.

How to find the state transition equation is a bit difficult but the equation is not difficult to understand. Though there are 3 strings, s_1 , s_2 and s_3 , $\text{len}(s_1) + \text{len}(s_2) = \text{len}(s_3)$ (otherwise the function must return false), consequently there are only 2 independent variables, so we let boolean $\text{dp}[i][j]$ represents if first i digits of s_1 interleaved with first j digits of s_2 comes up with first $i+j$ digits of s_3 .

Thus, $\text{dp}[i][j] = (s_1.\text{charAt}(i-1) == s_3.\text{charAt}(i+j-1) \ \&\& \ \text{dp}[i-1][j]) \ || \ (s_2.\text{charAt}(j-1) == s_3.\text{charAt}(i+j-1) \ \&\& \ \text{dp}[i][j-1])$;

Listing 336: Solution

```
public class Solution {
    public boolean isInterleave(String s1, String s2, String s3) {
        boolean[][] dp = new boolean[s1.length() + 1][s2.length() + 1];
        if (s1.length() + s2.length() != s3.length()) {
5             return false;
        }
        dp[0][0] = true;
        for (int i = 0 ; i <= s1.length(); i++){
            for (int j = 0; j <= s2.length() ; j++){
10                 if (i > 0 && s1.charAt(i - 1)
                    == s3.charAt(i + j - 1) && dp[i - 1][j]){
                    dp[i][j] = true;
                }
                if (j > 0 && s2.charAt(j - 1)
15                 == s3.charAt(i + j - 1) && dp[i][j - 1]){
                    dp[i][j] = true;
                }
            }
        }
20
        return dp[s1.length()][s2.length()];
    }
}
```

Problem 257

335 Self Crossing

You are given an array x of n positive numbers.
 You start at point $(0,0)$ and moves $x[0]$ metres to the north,
 then $x[1]$ metres to the west, $x[2]$ metres to the south,
 $x[3]$ metres to the east and so on.
 In other words,
 after each move your direction changes counter-clockwise.

Write a one-pass algorithm with $O(1)$ extra space to determine,
 if your path crosses itself, or not.

Example 1:

Given $x = [2, 1, 1, 2]$,

```

---
|   |
---+-->
|

```

Return true (self crossing)

Example 2:

Given $x = [1, 2, 3, 4]$,

```

-----
|       |
|       |
|       |
|       |
----->$

```

Return false (not self crossing)

Example 3:

Given $x = [1, 1, 1, 1]$,

```

---
|   |
---+-->
|

```

Return true (self crossing)

<https://leetcode.com/problems/self-crossing/>

Solution

Well, it is more like a math problem or pattern-finding problem.

The key is to find out what are the serial numbers of lines that can form a cross with the 1st line (1st line is just a representation of all lines since the process of validation can be cyclic, which means we can regard 2nd line as the '1st' line).

Just list the conclusion,

there are 3 situation where 2 lines form a cross with each other,

1.

```

      x(1)
      ---
x(2) |   | x(0)
      ---+-->
      x(3) |

```

2.

```

      x(1)
      -----
      |           | x(0)
x(2) |           ^
      |           | x(4)
      |           |
      -----|
      x(3)

```

3.

```

      x(1)
      -----
      |           | x(0)
x(2) |           <|---|
      |           x(5) | x(4)
      -----
      x(3)

```

Where each conditions can be derived by given circumstances. See code in listing.

Listing 337: Solution

```

class Move{
    int fromX, toX;
    int fromY, toY;
    public Move(int fromX, int fromY, int toX, int toY){
5         this.fromX = fromX;
        this.fromY = fromY;
        this.toX = toX;
        this.toY = toY;
    }
10 }

public class Solution {
    public boolean isSelfCrossing(int[] x) {
        if (x.length <= 3){
            return false;
15        }
        for (int moveIdx = 0 ; moveIdx < x.length ; moveIdx ++){
            // 1st line cross 4th line
            if (moveIdx >= 3){
                if (x[moveIdx - 3] >= x[moveIdx - 1]
20                 && x[moveIdx] >= x[moveIdx - 2]){
                    return true;
                }
            }
            // 1st cross 5th
25            if (moveIdx >= 4){
                if (x[moveIdx - 3] == x[moveIdx - 1]

```

```

        && x[moveIdx - 2] <= x[moveIdx - 4] + x[moveIdx]){
            return true;
        }
    }
    // 1st cross 6th
    if (moveIdx >= 5){
        if (x[moveIdx - 4] <= x[moveIdx - 2]
            && x[moveIdx - 1] <= x[moveIdx - 3]
            // No need to check
            // &&x[moveIdx - 5] <= x[moveIdx - 3]
            // && x[moveIdx] <= x[moveIdx - 2]
            && x[moveIdx - 2] <= x[moveIdx - 4] + x[moveIdx]
            && x[moveIdx - 3] <= x[moveIdx - 5] + x[moveIdx - 1]){
                return true;
            }
        }
    }
    return false;
}
}
```

Problem 258

68 Text Justification

Given an array of words and a length L,
format the text such that each line has
exactly L characters and is fully (left and right) justified.

You should pack your words in a greedy approach;
that is, pack as many words as you can in each line.
Pad extra spaces ' ' when necessary so that each line
has exactly L characters.

Extra spaces between words should be distributed
as evenly as possible.
If the number of spaces on a line do not divide evenly between words,
the empty slots on the left will be assigned more spaces
than the slots on the right.

For the last line of text, it should be left justified and
no extra space is inserted between words.

For example,
words: ["This", "is", "an", "example", "of", "text", "justification."]
L: 16.

Return the formatted lines as:

```
[
  "This    is    an",
  "example  of text",
  "justification. "
]
```

Note: Each word is guaranteed not to exceed L in length.

Corner Cases:

A line other than the last line might contain only one word.

What should you do in this case?

In this case, that line should be left-justified.

<https://leetcode.com/problems/text-justification/>

Solution

Just a simulation problem, but a bit complicated. Be aware of the last line and number of spaces in each line(I used Math.ceil to get spaceCnt each time, then remove spaceCnt from total space count).

Listing 338: Solution

```
public class Solution {
    public List<String> fullJustify(String[] words, int maxWidth) {
        if (words.length == 0){
            return new ArrayList<String>();
        }
        List<String> result = new ArrayList<String>();
```

```
boolean lastLine = false;
int curIdx = 0;
while (!lastLine && curIdx < words.length){
10     int curWidth = 0;
    int cntInLine = 0;
    int startIdx = curIdx;
    String newLine = "";
    while (curIdx < words.length &&
15         curWidth + words[curIdx].length() + curIdx - startIdx
            <= maxWidth){
        curWidth += words[curIdx++].length();
    }
    System.out.println(curIdx);
    if (curIdx == words.length){
20         lastLine = true;
        for (int i = startIdx ; i < curIdx ; i++){
            newLine += words[i];
            if (i != curIdx - 1){
25                 newLine += " ";
            }
        }

        int spaceCnt = maxWidth - newLine.length();
        for (int spaceIdx = 0 ; spaceIdx < spaceCnt ; spaceIdx++){
30             newLine += " ";
        }
    }
    else{
35         int spaceCnt = maxWidth - curWidth;
        for (int i = startIdx; i < curIdx ; i++){
            newLine += words[i];
            if (i != curIdx - 1){
                int curSpaceCnt =
40                 (int) Math.ceil(spaceCnt * 1.0 / (curIdx - i - 1));
                for (int spaceIdx = 0 ;
                    spaceIdx < curSpaceCnt ; spaceIdx++){
                    newLine += " ";
                }
45                 spaceCnt -= curSpaceCnt;
            }
        }
        for (int spaceIdx = 0 ; spaceIdx < spaceCnt; spaceIdx++){
50             newLine += " ";
        }
    }
    result.add(newLine);
}

55 return result;
}
```

Problem 259

218 The Skyline Problem

A city's skyline is the outer contour of the silhouette formed by all the buildings in that city when viewed from a distance. Now suppose you are given the locations and height of all the buildings as shown on a cityscape photo (Figure A), write a program to output the skyline formed by these buildings collectively (Figure B).

Buildings Skyline Contour

The geometric information of each building is represented by a triplet of integers $[Li, Ri, Hi]$, where Li and Ri are the x coordinates of the left and right edge of the i th building, respectively, and Hi is its height. It is guaranteed that $0 \leq Li, Ri \leq INT_MAX$, $0 < Hi \leq INT_MAX$, and $Ri - Li > 0$. You may assume all buildings are perfect rectangles grounded on an absolutely flat surface at height 0.

For instance, the dimensions of all buildings in Figure A are recorded as: $[[2, 9, 10], [3, 7, 15], [5, 12, 12], [15, 20, 10], [19, 24, 8]]$.

The output is a list of "key points" (red dots in Figure B) in the format of $[[x_1, y_1], [x_2, y_2], [x_3, y_3], \dots]$ that uniquely defines a skyline.

A key point is the left endpoint of a horizontal line segment. Note that the last key point, where the rightmost building ends, is merely used to mark the termination of the skyline, and always has zero height.

Also, the ground in between any two adjacent buildings should be considered part of the skyline contour.

For instance, the skyline in Figure B should be represented as: $[[2, 10], [3, 15], [7, 12], [12, 0], [15, 10], [20, 8], [24, 0]]$.

Notes:

The number of buildings in any input list is guaranteed to be in the range $[0, 10000]$.

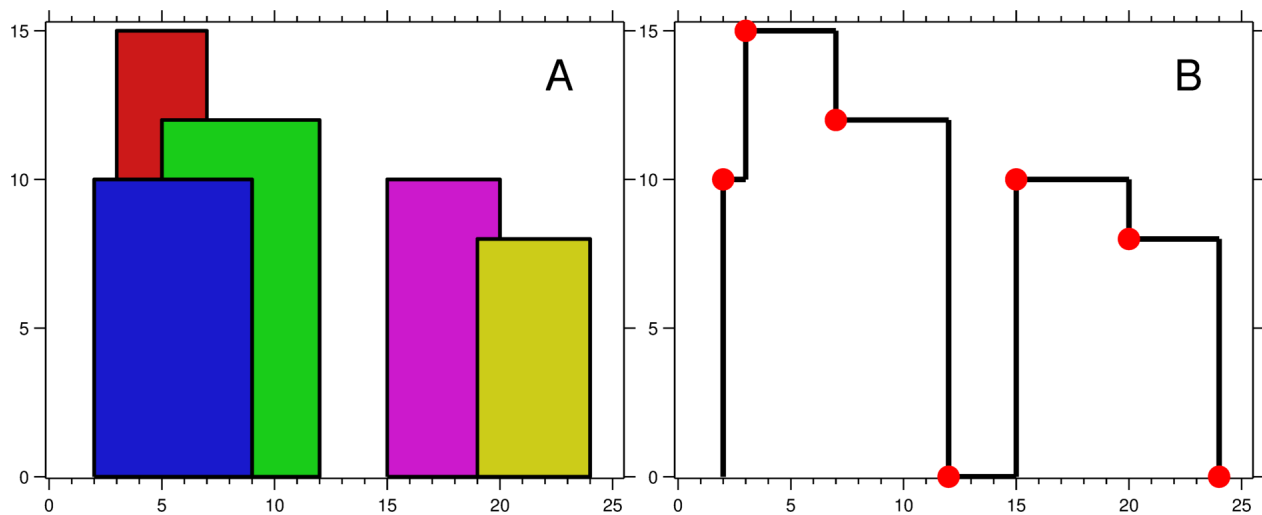
The input list is already sorted in ascending order by the left x position Li .

The output list must be sorted by the x position.

There must be no consecutive horizontal lines of equal height in the output skyline.

For instance, $[[...[2, 3], [4, 5], [7, 5], [11, 5], [12, 7], ...]]$ is not acceptable; the three lines of height 5 should be merged into one in the final output as such: $[[...[2, 3], [4, 5], [12, 7], ...]]$

<https://leetcode.com/problems/the-skyline-problem/>



Solution

Not so difficult if we know the routine to solve this problem.

1. We use left top and right top corner of rectangle to represent itself, for simplicity.
2. Sort the Points(corners) according to x-axis value, from small to large ones, identifying the left corner and right corner.
3. Iterate each Point, when we meet a left corner, add it into a max heap by its y-axis value. When the value we add is larger than former peek from heap, we may find a new 'shadow', so we add the (x,y) pair into result; when there is a right corner, pop it from the heap. When new peek is smaller than the element we pop, this is a new identifying point, add it into result.
4. Don't forget to deal with some corner cases, imagine when input is $[[1,2,1],[2,3,1]]$ OR $[[1,2,1],[1,2,2],[1,2,3]]$ OR $[[1,2,2],[1,2,3],[1,2,1]]$, what will happen?

Listing 339: Solution

```

class Point{
    int x, y;
    boolean left;
    int index;
    public Point(int x, int y , boolean left , int index){
        this.x = x;
        this.y = y;
        this.left = left;
        this.index = index;
    }
}

public class Solution {
    public List<int>> getSkyline(int [][] buildings) {
        if (buildings.length == 0){
            return new ArrayList<int []>();
        }
        List<int []> result = new ArrayList<int []>();
        List<Point> points = new ArrayList<Point>();
        for (int i = 0 ; i < buildings.length ; i++){
            int [] building = buildings[i];
            Point left = new Point(building[0], building[2], true, i);

```

```
        Point right = new Point(building[1], building[2], false, i);
        points.add(left);
        points.add(right);
25    }
    Queue<Integer> heap = new PriorityQueue<Integer>(
        buildings.length, new Comparator<Integer>(){
            public int compare(Integer i1, Integer i2){
30                return i2 - i1;
            }
        });

    heap.offer(0);
    Collections.sort(points, new Comparator<Point>(){
35        public int compare(Point p1, Point p2){
            if (p1.x == p2.x){
                if (p1.left && !p2.left){
                    return -1;
                }
40                else if (!p1.left && p2.left){
                    return 1;
                }
                else{
                    return 0;
45                }
            }
            return p1.x - p2.x;
        }
    });
50    for (Point point : points){
        if (point.left){
            if (point.y > heap.peek()){
                if (result.size() > 0
55                    && result.get(result.size() - 1)[0] == point.x){
                    result.remove(result.size() - 1);
                }
                result.add(new int[] {point.x, point.y});
            }
            heap.offer(point.y);
60        }else{
            heap.remove(point.y);
            if (point.y > heap.peek()){
                if (result.size() > 0
65                    && result.get(result.size() - 1)[0] == point.x){
                    result.remove(result.size() - 1);
                }
                result.add(new int[] {point.x, heap.peek()});
            }
        }
70    }
    return result;
}
```

Problem 260

336 Palindrome Pairs

Given a list of unique words.
Find all pairs of distinct indices (i, j) in the given list,
so that the concatenation of the two words,
i.e. words[i] + words[j] is a palindrome.

Example 1:

```
Given words = ["bat", "tab", "cat"]
Return [[0, 1], [1, 0]]
The palindromes are ["battab", "tabbat"]
```

Example 2:

```
Given words = ["abcd", "dcba", "lls", "s", "sssll"]
Return [[0, 1], [1, 0], [3, 2], [2, 4]]
The palindromes are ["dcbabcd", "abcdcba", "slls", "llssssll"]
```

<https://leetcode.com/problems/palindrome-pairs/>

Solution

First of all I must state that I don't think the problem is correct. When my input is ["aaa", "aaa", "aaa", "aa", "aa"] I should have gotten all combinations as my answer, but the expected answer is not what I want.

In order to pass the judgement, I have to do it in the way OJ advises.

When si+sj is palindrome, for example "as" and "llsa", there must be 'll' be a palindrome and 'as' be reversed by 'sa'. That's our key insight of solution. Iterate every word, split it into 2 parts, left and right, check if left OR right is palindrome and if it does, reverse the remaining part to see if it matches any other strings in original list.

Listing 340: Solution

```
class Solution{
    public List<List<Integer>> palindromePairs(String[] words) {
        List<List<Integer>> ret = new ArrayList<>();
        Map<String, Integer> wordIdxs = new HashMap<String, Integer>();
5      for (int wordIdx = 0 ; wordIdx < words.length ; wordIdx ++){
            wordIdxs.put(words[wordIdx], wordIdx);
        }
        for (int curIdx = 0 ; curIdx < words.length ; curIdx ++){
            String curStr = words[curIdx];
10         // String reverse = new StringBuilder(curStr).reverse().toString();
            for (int splitIdx = 0 ; splitIdx <= curStr.length() ; splitIdx ++){
                String left = curStr.substring(0, splitIdx);
                String right = curStr.substring(splitIdx);
                if (isPalindrome(left)){
15                 String rightRev =
                    new StringBuilder(right).reverse().toString();
                    if (wordIdxs.containsKey(rightRev)
                        && wordIdxs.get(rightRev) != curIdx){
                        List<Integer> pair = new ArrayList<Integer>();
20                 pair.add(wordIdxs.get(rightRev));
                        pair.add(curIdx);
                    }
                }
            }
        }
        return ret;
    }
}
```



```
                ret.add(pair);
            }
        }
    }
    if (isPalindrome(right)){
        String leftRev
            = new StringBuilder(left).reverse().toString();
        if (wordIdxs.containsKey(leftRev) && wordIdxs.get(leftRev)
            != curIdx && right.length() != 0){
            List<Integer> pair = new ArrayList<Integer>();
            pair.add(curIdx);
            pair.add(wordIdxs.get(leftRev));
            ret.add(pair);
        }
    }
}

return ret;
}

private boolean isPalindrome(String str) {
    int left = 0;
    int right = str.length() - 1;
    while (left <= right) {
        if (str.charAt(left++) != str.charAt(right--)) return false;
    }
    return true;
}
}
```

Problem 261

146 LRU Cache

Design and implement a data structure for Least Recently Used (LRU) cache. It should support the following operations: get and set.

```
get(key) - Get the value (will always be positive)
           of the key if the key exists in the cache, otherwise return -1.
set(key, value) - Set or insert the value
                  if the key is not already present.
                  When the cache reached its capacity,
                  it should invalidate the least recently used item
                  before inserting a new item.
```

<https://leetcode.com/problems/lru-cache/>

Solution

A interesting and not so difficult problem. With plenty kinds of solutions. The most convenient method I've ever seen is to extend LinkedHashMap directly and overwrite some functions, which I think is a bit tricky.

Use a HashMap to trace each element in cache and out of cache. For each CacheEntry, assign its previous entry and next entry. A trick is to assign initially a dummy head and dummy tail node to quickly acquire the least recently used entry(by tail.prev) and most recently used(by head.next).

Listing 341: Solution

```
class CacheEntry{
    int value, key;
    CacheEntry pre, next;
    public CacheEntry(int key, int value){
5         this.key = key;
          this.value = value;
    }
}
public class LRUCache {
10     CacheEntry dummyHead;
    CacheEntry dummyTail;
    int capacity;
    int curSize;
    HashMap<Integer, CacheEntry> map = new HashMap<Integer, CacheEntry>();
15
    public LRUCache(int capacity) {
        this.capacity = capacity;
        curSize = 0;
        dummyHead = new CacheEntry(-1, -1);
20         dummyTail = new CacheEntry(-1, -1);
        dummyHead.next = dummyTail;
        dummyTail.pre = dummyHead;
    }

25     public int get(int key) {
        if (map.containsKey(key)){
```

```
        CacheEntry entry = map.get(key);
        moveToHead(entry);

30         return entry.value;
    }
    return -1;
}

35 public void set(int key, int value) {
    if (map.containsKey(key)){
        CacheEntry entry = map.get(key);
        entry.value = value;
        moveToHead(entry);
40        map.put(key, entry);
    }else{
        CacheEntry entry = new CacheEntry(key, value);
        if (curSize < capacity){
            curSize++;
            addToHead(entry);
45        }
        else{
            map.remove(dummyTail.pre.key);
            removeEntry(dummyTail.pre);
50            addToHead(entry);
        }
        map.put(key, entry);
    }
}

55 private void moveToHead(CacheEntry node){
    removeEntry(node);
    addToHead(node);
}

60 private void removeEntry(CacheEntry node){
    // remove node
    node.next.pre = node.pre;
    node.pre.next = node.next;
}

65 private void addToHead(CacheEntry node){
    node.next = dummyHead.next;
    dummyHead.next.pre = node;
    node.pre = dummyHead;
    dummyHead.next = node;
70 }
}
```

Problem 262

214 Shortest Palindrome

Given a string *S*, you are allowed to convert it to a palindrome by adding characters in front of it.

Find and return the shortest palindrome you can find by performing this transformation.

For example:

Given "aacecaaa", return "aaacecaaa".

Given "abcd", return "dcbabcd".

<https://leetcode.com/problems/shortest-palindrome/>

Solution

A Brute force method is to substring from (0, s.length()) to (0,1), if the substring result is palindrome, then stop checking immediately(to get the shortest result) and reverse the remaining string when palindrome is cut from original string. Then attach the reversed string to the head of original string, which is exactly our result. Worst complexity $O(n^2)$, which will get TLE.

Another train of thought is to use the KMP algorithm.

For a quicker review of KMP, visit <http://www.cnblogs.com/c-cloud/p/3224788.html>.

For Chinese description of this problem, visit <https://segmentfault.com/a/1190000003059361>

Actually, the next[] array is enough for this problem. First we append reversed source string to tail of original string to get quite a long string for a result. Apply KMP on new string, get the largest element in next array, which is exactly the reversed part of original string.

BUT! We are extending the next array to be twice length as what former string has, which may lead to a error state that our longest prefix & suffix go through both reversed string and original string, when we only want to find longest prefix/suffix in original string. So we add a special character between s and reverse_s(here we use '#').

In conclusion, our solution can be divided into 3 steps.

1. Reverse string s as rev_s
2. Append rev_s to tail of s, namely form s#rev_s
3. Apply KMP to new string, get the maximum common length of prefix and suffix, such segment of string is actually palindrome(think of why.)
4. Split rev_s to get s.length() - common_length substring, attach it to head of string s

Listing 342: Solution

```
public class Solution {
    public String shortestPalindrome(String s) {
        if (s.length() <= 1){
            return s;
        }
        String revS = new StringBuilder(s).reverse().toString();
```

```
String longStr = s + "#" + revS;
int patternIdx = 0;
int[] next = new int[longStr.length()];
10 next[0] = 0;
for (int strIdx = 1 ; strIdx < longStr.length() ; strIdx++){
    while (patternIdx > 0 && longStr.charAt(strIdx)
        != longStr.charAt(patternIdx)){
15         patternIdx = next[patternIdx - 1];
    }
    if (longStr.charAt(strIdx) == longStr.charAt(patternIdx)){
        patternIdx++;
    }
    next[strIdx] = patternIdx;
20    // System.out.println(strIdx + " : " + patternIdx);
}
return revS.substring(0, revS.length() - next[longStr.length() - 1]) + s;
}
```

Problem 263

321 Create Maximum Number

Given two arrays of length m and n with digits 0-9 representing two numbers.
Create the maximum number of length k $k \leq m + n$
from digits of the two. The relative order of the digits
from the same array must be preserved.
Return an array of the k digits.
You should try to optimize your time and space complexity.

Example 1:

```
nums1 = [3, 4, 6, 5]
nums2 = [9, 1, 2, 5, 8, 3]
k = 5
return [9, 8, 6, 5, 3]
```

Example 2:

```
nums1 = [6, 7]
nums2 = [6, 0, 4]
k = 5
return [6, 7, 6, 0, 4]
```

Example 3:

```
nums1 = [3, 9]
nums2 = [8, 9]
k = 3
return [9, 8, 9]
```

<https://leetcode.com/problems/create-maximum-number/>

Solution

A Greedy problem, paste the solution in Chinese down here since I may not explain it correctly and accurately.

贪心，开一个栈，每次都尽量往里塞最大的数，最后这个栈就是结果。

假设需要取 n 位数字，每一轮入栈前要先维护栈的状态，如果同时满足以下条件就把栈顶的元素弹出：

1. 栈不为空；
2. 栈顶元素小于当前的元素；
3. 栈顶弹出后，后面剩下没处理的子串加上栈的length要够 n 位，

看这个例子[3,2,1,4,5]， $n = 3$ 。

前三轮没有疑问，栈里放入3, 2, 1， $stack = [3, 2, 1]$ 。

第四轮循环的时候4比3, 2, 1都大，但不能把它们都弹出，因为不够3位了，只弹出2个， $stack = [3, 4]$ 。

最后5放入栈， $stack = [3, 4, 5]$ 。

第三点对应代码中 " $stack.length + len - i > n$ " 这个条件。

分别取得了两个最大的子数，合并的操作就类似于merge sort。

双指针，一开始都指向0，放入一个大的，指针++，最后把剩下的都放入结果。

要注意可能有相同的数字，此时要看后面的数来决定哪个比较大，举例来说，[6, 1] 和 [6, 7]应该拿[6, 7]中的6。

Listing 343: Solution

```
public class Solution {
    public int [] maxNumber(int [] nums1, int [] nums2, int k) {
        int [] result = new int [k];
        if (nums1.length + nums2.length < k){
5             return result;
        }
        for (int length1 = 0 ;
            length1 <= Math.min(nums1.length , k) ; length1++){
            int length2 = k - length1;
10             if (length2 > nums2.length){
                continue;
            }

            int [] list1 = getKLenStack(nums1, length1);
15             int [] list2 = getKLenStack(nums2, length2);
            int [] newResult = new int [k];
            mergeStacks(list1 , list2 , newResult);

            for (int resIdx = 0 ; resIdx < newResult.length ; resIdx++){
20                 if (newResult[resIdx] == result[resIdx]){
                    continue;
                }
                else if (newResult[resIdx] > result[resIdx]){
                    result = newResult;
25                     break;
                }
                else{
                    break;
                }
            }
30        }
```

```
    }
    return result;
}

35 private void mergeStacks(int[] list1, int[] list2, int[] result){

    int idx = 0;
    int idx1 = 0, idx2 = 0;
    while (idx1 < list1.length && idx2 < list2.length){
40         if (list1[idx1] > list2[idx2]){
            result[idx++] = list1[idx1++];
        }
        else if (list1[idx1] < list2[idx2]){
            result[idx++] = list2[idx2++];
45         }
        else{
            int next1 = list1[idx1];
            int next2 = list2[idx2];
            int count = 1;
50             while (next1 == next2){
                if (idx1 + count < list1.length){
                    next1 = list1[idx1 + count];
                }
                if (idx2 + count < list2.length){
55                 next2 = list2[idx2 + count];
                }
                count++;
                if (idx1 + count > list1.length
                    && idx2 + count > list2.length){
60                     break;
                }
            }
            if (next1 >= next2){
                result[idx++] = list1[idx1++];
65             }
            else{
                result[idx++] = list2[idx2++];
            }
        }
70     }

    while (idx1 < list1.length){
        result[idx++] = list1[idx1++];
    }
    while (idx2 < list2.length){
75         result[idx++] = list2[idx2++];
    }
}

private int[] getKLenStack(int[] nums, int length){
80     int[] result = new int[length];
    Stack<Integer> stack = new Stack<Integer>();
    int curLen = 0;
    for (int numIdx = 0 ; numIdx < nums.length ; numIdx++){
```



```

    int num = nums[numIdx];
85    if (nums.length + stack.size() - numIdx <= length){
        stack.push(num);
        result[curLen] = num;
        curLen++;
        continue;
90    }
    while (!stack.isEmpty()
           && num > stack.peek()
           && stack.size() + nums.length - numIdx > length){
95        stack.pop();
        curLen--;
    }
    if (curLen < length){
100        stack.push(num);
        result[curLen] = num;
        curLen++;
    }
    }
    return result;
105 }
```

Problem 264

343 Integer Break

Given a positive integer n ,
break it into the sum of at least two positive integers
and maximize the product of those integers.
Return the maximum product you can get.

For example, given $n = 2$,
return 1 ($2 = 1 + 1$); given $n = 10$, return 36 ($10 = 3 + 3 + 4$).

Note: you may assume that n is not less than 2.

Hint:

There is a simple $O(n)$ solution to this problem.
You may check the breaking results of n ranging from
7 to 10 to discover the regularities.

<https://leetcode.com/problems/integer-break/>

Solution

Take as many 3 as we can, the proof process can be found below.

However, I'm not using such train of thought, but iterate each possible count of sum to get maximum product.

这道题的难点其实在于证明为什么拆出足够多的 3 就能使得乘积最大。下面我就试着证明一下。

首先证明拆出的因子大于 4 是不行的。设 x 是一个因子, $x > 4$, 那么可以将这个因子再拆成两个因子 $x - 2$ 和 2, 易证 $(x - 2) \times 2 > x$ 。所以不能有大于 4 的因子。

4 这个因子也是可有可无的, $4 = 2 + 2$, $4 = 2 \times 2$ 。因此 4 这个因子可以用两个 2 代替。

除非没有别的因子可用, 1 也不能选作因子。一个数 n 当它大于 3 时, 有 $(x - 2) \times 2 > (x - 1) \times 1$ 。

这样呢, 就只剩下 2 和 3 这两个因子可以选了。下面再证明 3 比 2 好。

一个数 $x = 3m + 2n$, 那么 $f = 3^m \times 2^n = 3^m \times 2^{\frac{x-3m}{2}}$ 可以对它取个对数。

$$\left. \begin{aligned} \ln f &= m \ln 3 + n \ln 2 \\ &= m \ln 3 + \frac{x-3m}{2} \ln 2 \\ &= \frac{x}{2} \ln 2 + \left(\ln 3 - \frac{3}{2} \ln 2 \right) m \end{aligned} \right\}$$

其中 $\ln 3 - \frac{3}{2} \ln 2 > 0$ 所以 f 是 m 的增函数, 也就是说 m 越大越好。所以 3 越多越好。

再多说一句, 如果拆出的因子不限于整数的话, 可以证明 $e = 2.718\dots$ 是最佳的选择。感兴趣的可以试着证明一下。

Listing 344: Solution

```
public class Solution {
    public int integerBreak(int n) {
        int max = -1;
        for (int digit = 2 ; digit <= n ; digit++){
            int product = 1;
```

```
10         int base = n / digit;
        int rem = n % digit;
        for (int prodIdx = 0 ; prodIdx < digit ; prodIdx ++){
            product *= (rem <= 0) ? base : (base + 1);
            rem --;
        }
        max = Math.max(product , max);
    }
    return max;
15 }
}
```

Problem 265

344 Reverse String

Write a function that takes a string as input and returns the string reversed.

Example:

Given s = "hello", return "olleh".

<https://leetcode.com/problems/reverse-string/>

Solution

An easy problem with plenty of ways to solve it.

Listing 345: Solution

```
public class Solution {  
    public String reverseString(String s) {  
        char[] charArr = s.toCharArray();  
        int i = 0 , j = s.length() - 1;  
5         while(i < j){  
            char tmp = charArr[i];  
            charArr[i] = charArr[j];  
            charArr[j] = tmp;  
            i++;  
10            j--;  
        }  
        return new String(charArr);  
    }  
15 }
```

Problem 266

345 Reverse Vowels of a String

Write a function that takes a string as input and reverse only the vowels of a string.

Example 1:

Given s = "hello", return "holle".

Example 2:

Given s = "leetcode", return "leotcede".

<https://leetcode.com/problems/reverse-vowels-of-a-string/>

Solution

An easy problem either, a bit bit more difficult than previous problem 'Reverse String'.

Listing 346: Solution

```
public class Solution {  
    public String reverseVowels(String s) {  
        Set<Character> set = new HashSet<Character>();  
        set.add('A'); set.add('a');  
        set.add('E'); set.add('e');  
        set.add('I'); set.add('i');  
        set.add('O'); set.add('o');  
        set.add('U'); set.add('u');  
        char[] charArr = s.toCharArray();  
        int i = 0, j = s.length() - 1;  
        while(i < j){  
            while (i < j && !(set.contains(s.charAt(i)))){  
                i++;  
            }  
            while (i < j && !(set.contains(s.charAt(j)))){  
                j--;  
            }  
            if (i >= j){  
                break;  
            }  
            char tmp = charArr[i];  
            charArr[i] = charArr[j];  
            charArr[j] = tmp;  
            i++;  
            j--;  
        }  
        return new String(charArr);  
    }  
}
```

Problem 267

349 Intersection of Two Arrays

Given two arrays, write a function to compute their intersection.

Example:

Given nums1 = [1, 2, 2, 1], nums2 = [2, 2], return [2].

Note:

Each element in the result must be unique.

The result can be in any order.

<https://leetcode.com/problems/intersection-of-two-arrays/>

Solution

Using HashSet to find intersection and avoid duplicate.

Listing 347: Solution

```
public class Solution {  
    public int[] intersection(int[] nums1, int[] nums2) {  
        Set<Integer> set1 = new HashSet<Integer>();  
        Set<Integer> result = new HashSet<Integer>();  
5         for (int num1 : nums1){  
            set1.add(num1);  
        }  
        for (int num2 : nums2){  
            if (set1.contains(num2)){  
10             result.add(num2);  
            }  
        }  
        int[] resultArr = new int[result.size()];  
        int idx = 0;  
15         for (int res : result){  
            resultArr[idx] = res;  
            idx ++;  
        }  
        return resultArr;  
20     }  
}
```

Problem 268

350 Intersection of Two Arrays II

Given two arrays, write a function to compute their intersection.

Example:

Given nums1 = [1, 2, 2, 1], nums2 = [2, 2], return [2, 2].

Note:

Each element in the result should appear as many times as it shows in both arrays.
The result can be in any order.

Follow up:

What if the given array is already sorted? How would you optimize your algorithm?
What if nums1's size is small compared to nums2's size? Which algorithm is better?
What if elements of nums2 are stored on disk, and the memory is limited
such that you cannot load all elements into the memory at once?

<https://leetcode.com/problems/intersection-of-two-arrays-ii/>

Solution 1

The most direct solution is to use a HashMap for 1 array recording each number for occurrence and match another, maintaining the count of number in 2nd array.

Listing 348: Solution 1

```
public class Solution {  
    public int[] intersect(int[] nums1, int[] nums2) {  
        if (nums1.length == 0 || nums2.length == 0){  
            return new int[0];  
        }  
        ArrayList<Integer> result = new ArrayList<Integer>();  
        HashMap<Integer, Integer> map1 = new HashMap<Integer, Integer>();  
        for (int num1 : nums1){  
            if (map1.containsKey(num1)){  
                map1.put(num1, map1.get(num1) + 1);  
            }  
            else{  
                map1.put(num1, 1);  
            }  
        }  
        for (int num2 : nums2){  
            if (map1.containsKey(num2) && map1.get(num2) > 0){  
                map1.put(num2, map1.get(num2) - 1);  
                result.add(num2);  
            }  
        }  
        int[] res = new int[result.size()];  
        for (int idx = 0 ; idx < result.size() ; idx ++){  
            res[idx] = result.get(idx);  
        }  
        return res;  
    }  
}
```

Solution 2

Another train of thought is by sorting and merging, can be achieved using either `Arrays.sort` or manually `quickSort`.

Listing 349: Solution 2

```
public class Solution {
    public int[] intersect(int[] nums1, int[] nums2) {
        if (nums1.length == 0 || nums2.length == 0){
            return new int[0];
        }
        Arrays.sort(nums1);
        Arrays.sort(nums2);
        List<Integer> result = new ArrayList<Integer>();
        int idx1 = 0; int idx2 = 0;
        while (idx1 < nums1.length && idx2 < nums2.length){
            if (nums1[idx1] == nums2[idx2]){
                result.add(nums1[idx1]);
                idx1++;
                idx2++;
            }
            else if (nums1[idx1] > nums2[idx2]){
                idx2++;
            }
            else{
                idx1++;
            }
        }
        int[] res = new int[result.size()];
        for (int idx = 0 ; idx < result.size(); idx++){
            res[idx] = result.get(idx);
        }
        return res;
    }
}
```

However, I still feel that we need further investigation on that, on one hand we are not quite sure about the questions in Follow up, while on the other there is no difference between I and II on algorithm, which may indicates that there must be somewhere we can optimize.

Problem 269

371 Sum of Two Integers

Calculate the sum of two integers a and b, but you are not allowed to use the operator + and -.

Example:

Given a = 1 and b = 2, return 3.

<https://leetcode.com/problems/sum-of-two-integers/>

Solution

It's more a mathematical(bit manipulation) problem. For $a + b$, if there is no carry, the result is exactly $a \oplus b$ by bit; similarly, if no carry, the new carry is $(a \& b) \ll 1$, adding them still get sum, which begins a new iteration. Loop until new carry == 0, where $a \oplus b$ gets final sum.

Listing 350: Solution

```
public class Solution {  
    public int getSum(int a, int b) {  
        int sum = 0, carry = -1;  
        while (carry != 0){  
5           sum = (a ^ b);  
           carry = (a & b) << 1;  
           a = sum;  
           b = carry;  
        }  
10        return sum;  
    }  
}
```

Problem 270

383 Ransom Note

Given an arbitrary ransom note string and another string containing letters from all the magazines, write a function that will return true if the ransom note can be constructed from the magazines ; otherwise, it will return false.

Each letter in the magazine string can only be used once in your ransom note.

Note:

You may assume that both strings contain only lowercase letters.

```
canConstruct("a", "b") - false
canConstruct("aa", "ab") - false
canConstruct("aa", "aab") - true
```

<https://leetcode.com/problems/ransom-note/>

Solution

Array is enough.

Listing 351: Solution

```
public class Solution {
    public boolean canConstruct(String ransomNote, String magazine) {
        int[] letters = new int[257];
        char[] magChars = magazine.toCharArray();
        5      for (char magCh : magChars){
            letters[magCh] ++;
        }
        char[] ransomChars = ransomNote.toCharArray();
        10     for (char ransomCh : ransomChars){
            if (letters[ransomCh] == 0){
                return false;
            }
            else{
                letters[ransomCh] --;
            }
        }
        15
        return true;
    }
    20 }
```

Problem 271

373 Find K Pairs with Smallest Sums

You are given two integer arrays `nums1` and `nums2` sorted in ascending order and an integer `k`.

Define a pair (u,v) which consists of one element from the first array and one element from the second array.

Find the `k` pairs $(u_1,v_1), (u_2,v_2) \dots (u_k,v_k)$ with the smallest sums.

Example 1:

Given `nums1 = [1,7,11]`, `nums2 = [2,4,6]`, `k = 3`

Return: `[1,2], [1,4], [1,6]`

The first 3 pairs are returned from the sequence:

`[1,2], [1,4], [1,6], [7,2], [7,4], [11,2], [7,6], [11,4], [11,6]`

Example 2:

Given `nums1 = [1,1,2]`, `nums2 = [1,2,3]`, `k = 2`

Return: `[1,1], [1,1]`

The first 2 pairs are returned from the sequence:

`[1,1], [1,1], [1,2], [2,1], [1,2], [2,2], [1,3], [1,3], [2,3]`

Example 3:

Given `nums1 = [1,2]`, `nums2 = [3]`, `k = 3`

Return: `[1,3], [2,3]`

All possible pairs are returned from the sequence:

`[1,3], [2,3]`

<https://leetcode.com/problems/find-k-pairs-with-smallest-sums/>

Solution

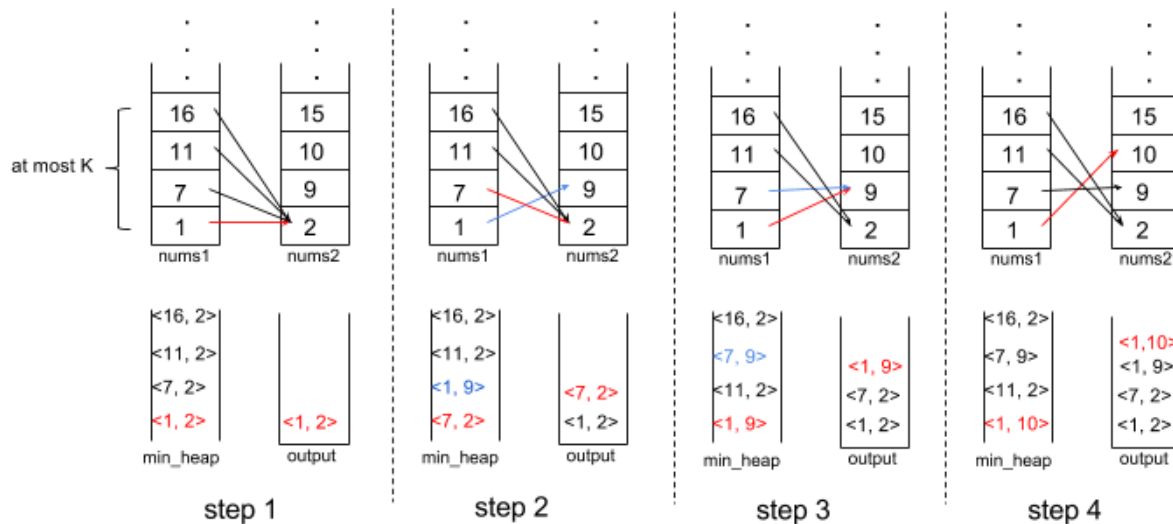
Heap is the key to solution. However, this is not adequate, the train of thought is still a bit complicated.

See explanation and figure from Discussion section.

Basic idea: Use `min_heap` to keep track on next minimum pair sum, and we only need to maintain `K` possible candidates in the data structure.

Some observations: For every numbers in `nums1`, its best partner(yields min sum) always strats from `nums2[0]` since arrays are all sorted; And for a specific number in `nums1`, its next candidate could be `[this specific number] + nums2[current_associated_index + 1]`, unless out of boundary;)

Here is a simple example demonstrate how this algorithm works.



Figure_1. A simple example when $K = 4$ and `nums1={1,7,11,16,...}`, `nums2={2,9,10,15,...}`

The run time complexity is $O(k \log k)$ since `queue.size` $\leq k$ and we do at most k loop.

Listing 352: Solution

```

public class Solution {
    public List<int[]> kSmallestPairs(int[] nums1, int[] nums2, int k) {
        List<int[]> result = new ArrayList<int[]>();
        PriorityQueue<int[]> queue = new PriorityQueue<>(
5         (int[] a, int[] b) -> (a[0] + a[1] - b[0] - b[1]));
        if (nums1.length == 0 || nums2.length == 0 || k == 0) {
            return result;
        }
        for (int idx = 0; idx < nums1.length; idx++) {
10         queue.offer(new int[] { nums1[idx], nums2[0], 0 });
        }
        while (result.size() < k && !queue.isEmpty()) {
            int[] cur = queue.poll();
            result.add(new int[] { cur[0], cur[1] });
15         if (cur[2] == nums2.length - 1) {
                continue;
            }
            queue.offer(new int[] { cur[0], nums2[cur[2] + 1], cur[2] + 1 });
        }
20         return result;
    }
}

```

Problem 272

374 Guess Number Higher or Lower

We are playing the Guess Game. The game is as follows:

I pick a number from 1 to n. You have to guess which number I picked.

Every time you guess wrong, I'll tell you whether the number is higher or lower.

You call a pre-defined API `guess(int num)` which returns 3 possible results (-1, 1, or 0):

-1 : My number is lower

1 : My number is higher

0 : Congrats! You got it!

Example:

n = 10, I pick 6.

Return 6.

<https://leetcode.com/problems/guess-number-higher-or-lower/>

Solution

A standard binary search problem. Easy to implement and solved.

Listing 353: Solution

```
/* The guess API is defined in the parent class GuessGame.
   @param num, your guess
   @return -1 if my number is lower, 1 if my number is higher, otherwise return 0
   int guess(int num); */
5
public class Solution extends GuessGame {
    public int guessNumber(int n) {
        int low = 1, high = n;
        int mid = low + (high - low) / 2;
10    int flag = guess(mid);
        while (flag != 0 && low < high){
            if (flag == 1){
                low = mid + 1;
            }
15        else if (flag == -1){
            high = mid - 1;
        }
        mid = low + (high - low) / 2;
20        flag = guess(mid);
    }
    return mid;
}
```

Problem 273

357 Count Numbers with Unique Digits

Given a non-negative integer n , count all numbers with unique digits, x , where $0 \leq x < 10^n$.

Example:

Given $n = 2$, return 91. (The answer should be the total numbers in the range of $0 \leq x < 100$, excluding [11,22,33,44,55,66,77,88,99])

Hint:

A direct way is to use the backtracking approach.

Backtracking should contains three states which are

(the current number, number of steps to get that number and a bitmask which represent which number is marked as visited so far in the current number).

Start with state (0,0,0) and count all valid number

till we reach number of steps equals to 10^n .

This problem can also be solved using a dynamic programming approach and some knowledge of combinatorics.

Let $f(k)$ = count of numbers with unique digits with length equals k .

$f(1) = 10, \dots, f(k) = 9 * 9 * 8 * \dots (9 - k + 2)$

[The first factor is 9 because a number cannot start with 0].

<https://leetcode.com/problems/count-numbers-with-unique-digits/>

Solution 1

Backtracking solution, recurse on each digit from 0 to 9, use a int act as bitmap to see which number has been used. Be careful when checking if 0 used. I'm using the method that when used == 0, other numbers have not been used, which also means that 0 chosen is the 1st digit, only in this case shouldn't we mark number 0 as used.

Listing 354: Solution 1 - Backtracking

```
public class Solution {
    int count = 0;
    public int countNumbersWithUniqueDigits(int n) {
        int used = 0;
5       if (n == 0){
            return 1;
        }
        if (n > 10){
10          return countNumbersWithUniqueDigits(10);
        }

        countNumber(1, n, used);
        return count;
    }
15    private void countNumber(int digitIdx, int n, int used){
        if (digitIdx > n){
            count ++;
        }
    }
}
```

```
        return;
    }
20    for (int digit = 0 ; digit <= 9 ; digit++){
        int mask = (1 << digit);
        if ((mask & used) != 0){
            continue;
        }
25    if (digit != 0 || used != 0){
        used |= mask;
    }

    countNumber(digitIdx + 1, n, used);
30    used &= (~mask);
    }
    }
}
```

Solution 2

Use math method as hint says. For $n = 1$, $f(1) = 10$ [0,1,2,...,9] For $n = 2$, $f(2) = 9 * 9$ [1st-digit: 1-9 & 2nd-digit: 0-9 except for 1st-digit] ... Finally, sum them up.

Listing 355: Solution 2 - Math

```
public class Solution {
    public int countNumbersWithUniqueDigits(int n) {
        if (n == 0){
            return 1;
5        }
        int sum = 10, product = 9;
        for (int digit = 2 ; digit <= Math.min(n, 10) ; digit++){
            product *= (11 - digit);
            sum += product;
10        }

        return sum;
    }
}
```

Problem 274

357 Count Numbers with Unique Digits

Given a non-negative integer n , count all numbers with unique digits, x , where $0 \leq x < 10^n$.

Example:

Given $n = 2$, return 91. (The answer should be the total numbers in the range of $0 \leq x < 100$, excluding [11,22,33,44,55,66,77,88,99])

Hint:

A direct way is to use the backtracking approach.

Backtracking should contains three states which are

(the current number, number of steps to get that number and a bitmask which represent which number is marked as visited so far in the current number).

Start with state (0,0,0) and count all valid number

till we reach number of steps equals to 10^n .

This problem can also be solved using a dynamic programming approach and some knowledge of combinatorics.

Let $f(k)$ = count of numbers with unique digits with length equals k .

$f(1) = 10, \dots, f(k) = 9 * 9 * 8 * \dots (9 - k + 2)$

[The first factor is 9 because a number cannot start with 0].

<https://leetcode.com/problems/count-numbers-with-unique-digits/>

Solution 1

Backtracking solution, recurse on each digit from 0 to 9, use a int act as bitmap to see which number has been used. Be careful when checking if 0 used. I'm using the method that when used == 0, other numbers have not been used, which also means that 0 chosen is the 1st digit, only in this case shouldn't we mark number 0 as used.

Listing 356: Solution 1 - Backtracking

```
public class Solution {
    int count = 0;
    public int countNumbersWithUniqueDigits(int n) {
        int used = 0;
5         if (n == 0){
            return 1;
        }
        if (n > 10){
10            return countNumbersWithUniqueDigits(10);
        }

        countNumber(1, n, used);
        return count;
    }
15    private void countNumber(int digitIdx, int n, int used){
        if (digitIdx > n){
            count ++;
        }
    }
}
```



```

        return;
    }
20    for (int digit = 0 ; digit <= 9 ; digit++){
        int mask = (1 << digit);
        if ((mask & used) != 0){
            continue;
        }
25    if (digit != 0 || used != 0){
        used |= mask;
    }

    countNumber(digitIdx + 1, n, used);
30    used &= (~mask);
    }
}
}

```

Solution 2

Use math method as hint says. For $n = 1$, $f(1) = 10$ [0,1,2,...,9] For $n = 2$, $f(2) = 9 * 9$ [1st-digit: 1-9 & 2nd-digit: 0-9 except for 1st-digit] ... Finally, sum them up.

Listing 357: Solution 2 - Math

```

public class Solution {
    public int countNumbersWithUniqueDigits(int n) {
        if (n == 0){
            return 1;
5        }
        int sum = 10, product = 9;
        for (int digit = 2 ; digit <= Math.min(n, 10) ; digit++){
            product *= (11 - digit);
            sum += product;
10        }

        return sum;
    }
}

```

Problem 275

367 Valid Perfect Square

Given a positive integer num, write a function which returns True if num is a perfect square else False.

Note: Do not use any built-in library function such as sqrt.

Example 1:

Input: 16
Returns: True
Example 2:

Input: 14
Returns: False

<https://leetcode.com/problems/valid-perfect-square/>

Solution 1

As for a binary search problem, I can't tell this should be classified as Medium difficulty. Using division instead of multiplication is a trick to prevent overflow.

Listing 358: Solution 1 - Binary Search

```
public class Solution {  
    public boolean isPerfectSquare(int num) {  
        if (num == 1){  
            return true;  
        }  
        int low = 1, high = num;  
        int mid = low + (high - low) / 2;  
        while (low <= high){  
            int div = num / mid;  
            if (mid == div){  
                if (num % div == 0){  
                    return true;  
                }  
            }  
            else{  
                return false;  
            }  
            else if (mid < div){  
                low = mid + 1;  
            }  
            else{  
                high = mid - 1;  
            }  
            mid = low + (high - low) / 2;  
        }  
        return false;  
    }  
}
```

Solution 2

A interesting thing is that:

$$1 = 1 \quad 4 = 1 + 3 \quad 9 = 1 + 3 + 5 \quad \dots$$

Listing 359: Solution 2 - Math.sum

```
public class Solution {  
    public boolean isPerfectSquare(int num) {  
        int i = 1;  
        while (num > 0) {  
            num -= i;  
            i += 2;  
        }  
        return num == 0;  
    }  
}
```

Solution 3

The 3rd solution is a math method. I don't quite get the point so just leave the solution here.

Listing 360: Solution 3 - Math.Newton

```
class Solution {  
public:  
    bool isPerfectSquare(int n) {  
        float diff;  
        float x, x_prev = n;  
        do{  
            x = x_prev - (x_prev*x_prev - n)/(2*x_prev);  
            diff = x_prev - x;  
            x_prev = x;  
        }  
        while(diff >= 1);  
        // if n is a perfect square, its square root must be the floor of x  
        x = (int) x;  
        if (x*x == n)  
            return true;  
        else  
            return false;  
    }  
};
```

Problem 276

376 Wiggle Subsequence

A sequence of numbers is called a wiggle sequence if the differences between successive numbers strictly alternate between positive and negative.

The first difference (if one exists) may be either positive or negative.

A sequence with fewer than two elements is trivially a wiggle sequence.

For example, `[1,7,4,9,2,5]` is a wiggle sequence because

the differences `(6,-3,5,-7,3)` are alternately positive and negative.

In contrast, `[1,4,7,2,5]` and `[1,7,4,5,5]` are not wiggle sequences,

the first because its first two differences are positive and

the second because its last difference is zero.

Given a sequence of integers, return the length of the longest subsequence that is a wiggle sequence. A subsequence is obtained by deleting some number of elements (eventually, also zero) from the original sequence, leaving the remaining elements in their original order.

Examples:

Input: `[1,7,4,9,2,5]`

Output: 6

The entire sequence is a wiggle sequence.

Input: `[1,17,5,10,13,15,10,5,16,8]`

Output: 7

There are several subsequences that achieve this length. One is `[1,17,10,13,10,16,8]`.

Input: `[1,2,3,4,5,6,7,8,9]`

Output: 2

Follow up:

Can you do it in $O(n)$ time?

<https://leetcode.com/problems/valid-perfect-square/>

Solution 1

A dp solution with $O(n^2)$, unfortunately this is the worst solution in time complexity.

Listing 361: Solution 1 - DP

```
public class Solution {
    public int wiggleMaxLength(int[] nums) {
        if (nums.length == 0){
            return 0;
        }
        int maxLength = -1;
        // last diff, 0 for positive, 1 for negative
        int[][] dp = new int[nums.length][2];

        for (int idx = 0; idx < nums.length ; idx++){
            dp[idx][0] = 1;
            dp[idx][1] = 1;
            for (int formerIdx = idx - 1 ; formerIdx >= 0 ; formerIdx --){
```

```
15         int maxPositive = -1, maxNegative = -1;
16         int diff = nums[idx] - nums[formerIdx];
17         if (diff > 0){
18             maxPositive = Math.max(dp[formerIdx][1], maxPositive);
19         }
20         else if (diff < 0){
21             maxNegative = Math.max(dp[formerIdx][0], maxPositive);
22         }
23         else{
24             continue;
25         }
26         dp[idx][0] = Math.max(dp[idx][0], maxPositive + 1);
27         dp[idx][1] = Math.max(dp[idx][1], maxNegative + 1);
28         maxLength = Math.max(maxLength, Math.max(dp[idx][0], dp[idx][1]));
29     }
30 }
31
32 return Math.max(maxLength, 1);
33 }
```

Solution 2

A $O(n)$ solution can be found in Greedy or Math. As for greedy, just find ; while as for math, just find the switching points, that is exactly the number of maxLength.

Listing 362: Solution 2 - Greedy

```
// not yet completed
```

Listing 363: Solution 2 - Math

```
// not yet completed
```

Problem 277

387 First Unique Character in a String

Given a string, find the first non-repeating character in it and return it's index. If it doesn't exist, return -1.

Examples:

```
s = "leetcode"
return 0.
```

```
s = "loveleetcode",
return 2.
```

Note: You may assume the string contain only lowercase letters.

<https://leetcode.com/problems/first-unique-character-in-a-string/>

Solution

A simple problem, array counting its count and index is enough.

Listing 364: Solution

```
public class Solution {
    public int firstUniqChar(String s) {
        int[] appearances = new int[26];
        for (int aIdx = 0 ; aIdx < 26; aIdx++){
            appearances[aIdx] = -1;
        }
        for (int idx = 0 ; idx < s.length() ; idx++){
            int letterIdx = s.charAt(idx) - 'a';
            if (appearances[letterIdx] == -1){
                appearances[letterIdx] = idx;
            }
            else if (appearances[letterIdx] >= 0){
                appearances[letterIdx] = -2;
            }
        }
        int minIdx = Integer.MAX_VALUE;
        for (int aIdx = 0 ; aIdx < 26 ; aIdx++){
            if (appearances[aIdx] >= 0){
                minIdx = Math.min(appearances[aIdx], minIdx);
            }
        }
        return (minIdx == Integer.MAX_VALUE) ? -1 : minIdx;
    }
}
```

Problem 278

365 Water and Jug Problem

You are given two jugs with capacities x and y litres.
There is an infinite amount of water supply available.
You need to determine whether it is possible to measure exactly z litres using these two jugs.

If z liters of water is measurable,
you must have z liters of water contained within one or both buckets
by the end.

Operations allowed:

Fill any of the jugs completely with water.
Empty any of the jugs.
Pour water from one jug into another till the other jug is completely full or the first jug itself is empty.

Example 1: (From the famous "Die Hard" example)

Input: $x = 3, y = 5, z = 4$

Output: True

Example 2:

Input: $x = 2, y = 6, z = 5$

Output: False

<https://leetcode.com/problems/water-and-jug-problem/>

Solution

If you know the solution, the problem is not so difficult. Just check if z is divisible by $\text{gcd}(x,y)$.

BUT! Mention so may edge cases, like $x == 0, y == 0, z == 0$, and $x + y < z$ (For example, for 1,2,5, though $1+2+2=5$, it cannot be achieved)

Listing 365: Solution

```
public class Solution {
    public boolean canMeasureWater(int x, int y, int z) {
        if (z == 0){
            return true;
        }
        if (x + y < z){
            return false;
        }
        if (x == 0 && y == 0){
            return false;
        }
        int cd = gcd(x, y);
        return (z % cd == 0);
    }

    private int gcd(int x, int y){
```

```
    if (x < y){  
        return gcd(y, x);  
    }  
20    if (y == 0){  
        return x;  
    }  
    int quo = 0, rem = 0;  
25    do{  
        rem = x % y;  
        x = y;  
        y = rem;  
    }  
    while (rem != 0);  
30    return x;  
    }  
}
```


Problem 279

372 Super Pow

Your task is to calculate $ab \bmod 1337$ where a is a positive integer and b is an extremely large positive integer given in the form of an array.

Example1:

$a = 2$
 $b = [3]$

Result: 8

Example2:

$a = 2$
 $b = [1,0]$

Result: 1024

<https://leetcode.com/problems/super-pow/>

Solution 1

Fast Exponentiation is one solution, just aware of using high-precision solution to deal with the div 2 calculation. And, remember to mod 1337 everywhere.

Listing 366: Solution 1 - Fast Exponentiation

```
public class Solution {
    public int superPow(int a, int[] b) {
        if (isZero(b)){
            return 1;
        }
        boolean isEven = (b[b.length - 1] % 2 == 0);
        int product = superPow(a, div2(b)) % 1337;
        if (isEven){
            return product * product % 1337;
        }
        else{
            return (product * product % 1337 * (a % 1337)) % 1337;
        }
    }
    private boolean isZero(int[] b){
        for (int digit : b){
            if (digit != 0){
                return false;
            }
        }
        return true;
    }
    private int[] div2(int[] b){
        int carry = 0;
        for (int idxB = 0 ; idxB < b.length ; idxB ++){
            b[idxB] = carry * 10 + b[idxB];
        }
    }
}
```

```

        carry = b[idxB] % 2;
        b[idxB] /= 2;
    }
30     return b;
    }
}

```

Solution 2

I didn't get the solution... Copy the solution here, for future inspiration.

The main idea is cashed on the repeated pattern of the remainder of a^b .

As long as we know the length of the pattern m, we just have to find an index point of this pattern based on b mod m.

In addition, if $a > 1337$, we can let $a = a \bmod 1337$.

Because if we let $a = (1337x + c)$ where $c = a \bmod 1337$,

$(1337x + c)(1337x + c)(1337x + c) \dots (1337x + c) \bmod 1337 == c \dots c \bmod 1337$.

It iterates at most 1337 times to find the circle. And find the remainder to the loop size. And then look up in the history to find the final answer.

Listing 367: Solution 2 - the Remainder Repeat Pattern

```

public class Solution {
    int DIV = 1337;
    List<Integer> findLoop(int a){
        List<Integer> index = new ArrayList<>();
5         boolean[] set = new boolean[DIV];
        int rem = a % DIV;
        while ( ! set[rem] ) {
            set[rem]=true;
            index.add(rem);
10            rem = (rem*a) % DIV;
        }
        return index;
    }
    int modBy(int[] b, int m){
15         int rem = 0;
        for (int i=0; i < b.length; i++) {
            rem = (rem*10+b[i]) % m;
        }
        return rem;
20    }
    public int superPow(int a, int[] b) {
        if (a==0 || a==DIV || b==null || b.length == 0) return 0;
        if (a==1) return 1;
        if (a > DIV) return superPow( a % DIV, b);
25        List<Integer> index = findLoop(a);
        int loopsize = index.size();
        int rem = modBy(b, loopsize);
        rem = rem==0? loopsize: rem;
        return index.get(rem-1);
30    }
}

```

Problem 280

386 Lexicographical Numbers

Given an integer n , return $1 - n$ in lexicographical order.

For example, given 13, return: [1,10,11,12,13,2,3,4,5,6,7,8,9].

Please optimize your algorithm to use less time and space.

The input size may be as large as 5,000,000.

<https://leetcode.com/problems/lexicographical-numbers/>

Solution 1

Just a dfs problem, not difficult at all. However, optimization on the way.

Listing 368: Solution 1 - DFS

```
public class Solution {
    List<Integer> result = new ArrayList<Integer>();
    public List<Integer> lexicalOrder(int n) {
        traverse(n, 0);
        return result;
    }
    private boolean traverse(int n, int curNum){
        if (curNum > n){
            return false;
        }
        if (curNum != 0){
            result.add(curNum);
        }
        for (int i = 0 ; i <= 9 ; i++){
            if (curNum == 0 && i == 0){
                continue;
            }
            if (!traverse(n, curNum * 10 + i)){
                break;
            }
        }
        return true;
    }
}
```

Solution 2

To be continued... Needs optimization.

Problem 281

347 Top K Frequent Elements

Given a non-empty array of integers, return the k most frequent elements.

For example,

Given [1,1,1,2,2,3] and k = 2, return [1,2].

Note:

You may assume k is always valid, $1 \leq k \leq$ number of unique elements.

Your algorithm's time complexity must be better than $O(n \log n)$, where n is the array's size.

<https://leetcode.com/problems/top-k-frequent-elements/>

Solution

For $O(n \log n)$, heap sort, merge sort and quick sort come into mind. When PriorityQueue provide us with built-in heap sort, we use it to solve the problem.

Listing 369: Solution

```
public class Solution {  
    public List<Integer> topKFrequent(int[] nums, int k) {  
        List<Integer> result = new ArrayList<Integer>();  
        PriorityQueue<Entry> queue = new PriorityQueue<Entry>(  
5          (a,b) -> b.count - a.count);  
        HashMap<Integer, Integer> map = new HashMap<Integer, Integer>();  
        for (int num : nums){  
            if (map.containsKey(num)){  
                map.put(num, map.get(num) + 1);  
10            }  
            else{  
                map.put(num, 1);  
            }  
        }  
15        for (int key : map.keySet()){  
            Entry entry = new Entry();  
            entry.value = key;  
            entry.count = map.get(key);  
            queue.offer(entry);  
20        }  
  
        for (int idx = 0 ; idx < k ; idx ++){  
            Entry entry = queue.poll();  
            result.add(entry.value);  
25        }  
        return result;  
    }  
}  
class Entry{  
30    public int value;  
    public int count;  
}
```

Problem 282

388 Longest Absolute File Path

Suppose we abstract our file system by a string in the following manner:

The string "dir\n\tsubdir1\n\tsubdir2\n\t**t**\tfile.ext" represents:

```
dir
  subdir1
  subdir2
  file.ext
```

The directory `dir` contains an empty sub-directory `subdir1` and a sub-directory `subdir2` containing a file `file.ext`.

[illegible]

```
dir
  subdir1
    file1.ext
    subsubdir1
  subdir2
    subsubdir2
    file2.ext
```

The directory `dir` contains two sub-directories `subdir1` and `subdir2`.
`subdir1` contains a file `file1.ext` and
an empty second-level sub-directory `subsubdir1`.
`subdir2` contains a second-level sub-directory `subsubdir2` containing
a file `file2.ext`.

We are interested in finding the longest (number of characters) absolute path to a file within our file system.

For example, in the second example above,
the longest absolute path is "dir/subdir2/subsubdir2/file2.ext",
and its length is 32 (not including the double quotes).

Given a string representing the file system in the above format, return the length of the longest absolute path to file in the abstracted file system. If there is no file in the system, return 0.

Note:

The name of a file contains at least a . and an extension.

The name of a directory or sub-directory will not contain a ..

Time complexity required: $O(n)$ where n is the size of the input string.

Notice that `a/aa/aaa/file1.txt` is not the longest file path, if there is another path `aaaaaaaaaaaaaaaaaaaaaaa/sth.png`.

<https://leetcode.com/problems/longest-absolute-file-path/>

Solution

Used number of '\t' as array index, a 1D array is enough to store the value since the structure is like a tree.

Listing 370: Solution

```
/**
 * Definition for singly-linked list.
 * public class ListNode {
 *     int val;
 *     ListNode next;
 *     ListNode(int x) { val = x; }
 * }
 */
public class Solution {
10     public ListNode reverse(ListNode start, ListNode end){
        if (start == null)
            return null;
        ListNode cur = start, next = cur.next;
        cur.next = null;
15     while (next != null && next != end){
        ListNode backup = next.next;
        next.next = cur;
        cur = next;
        next = backup;
20     }
    return cur;
}

    public ListNode reverseKGroup(ListNode head, int k) {
        ListNode cur = head;
25     ListNode last = cur;
        ListNode dummy = new ListNode(-1);
        dummy.next = head;
        ListNode tail = dummy;
        while (cur != null){
30             int count = k;
            last = cur;
            while (count-- > 0){
                cur = cur.next;
                if (cur == null)
35                     break;
            }
            if (count > 0 && cur == null){
                tail.next = last;
                return dummy.next;
40            }
            ListNode newPart = reverse(last, cur);
            tail.next = newPart;
            tail = last;
        }
45     return dummy.next;
    }
}
```

Problem 283

354 Russian Doll Envelopes

You have a number of envelopes with widths and heights given as a pair of integers (w, h).

One envelope can fit into another if and only if both the width and height of one envelope is greater than the width and height of the other envelope.

What is the maximum number of envelopes can you Russian doll? (put one inside other)

Example:

Given envelopes = [[5,4],[6,4],[6,7],[2,3]],
the maximum number of envelopes you can Russian doll is
3 ([2,3] => [5,4] => [6,7]).

<https://leetcode.com/problems/russian-doll-envelopes/>

Solution

A very interesting problem. First, sort the envelopes according to no matter width or height, let's say width, when 2 envelopes are of the same width, put the one with larger height in the front. Then the problem is converted to 'Longest Increasing Subsequence' (and now we know why we put larger ones in front, since envelopes with same width cannot enclose each other, we use the sequence of larger -> smaller to avoid enclosing envelopes with same width). For 'Longest Increasing Subsequence', there is a $O(n^2)$ DP solution and $O(n \log n)$ Binary Search solution.

Listing 371: Solution

```
public class Solution {
    public int maxEnvelopes(int[][] envelopes) {
        quickSort(envelopes, 0, envelopes.length - 1);
        List<Integer> increasing = new ArrayList<Integer>();
5        for (int idx = 0 ; idx < envelopes.length ; idx++){
            if (increasing.size() == 0
                || envelopes[idx][1] > increasing.get(increasing.size() - 1)){
                increasing.add(envelopes[idx][1]);
            }
10           else{
                // binary search
                int properIdx = binarySearch(increasing, 0,
                    increasing.size() - 1, envelopes[idx][1]);
                if (properIdx < increasing.size()){
15                     increasing.set(properIdx, envelopes[idx][1]);
                }
            }
        }
        return increasing.size();
20    }

    private void quickSort(int[][] envelopes, int start, int end){
        if (start >= end){
```

```

    return;
25 }
    int pivot = envelopes[start][0];
    int pivot2 = envelopes[start][1];
    int low = start, high = end;
    while (low < high){
30     while (low < high && (envelopes[high][0] > pivot
        || (envelopes[high][0] == pivot && envelopes[high][1] <= pivot2))){
        high--;
    }
    envelopes[low][0] = envelopes[high][0];
35 envelopes[low][1] = envelopes[high][1];

    while (low < high && (envelopes[low][0] < pivot
        || (envelopes[low][0] == pivot && envelopes[low][1] >= pivot2))){
        low++;
40    }
    envelopes[high][0] = envelopes[low][0];
    envelopes[high][1] = envelopes[low][1];
    }
    envelopes[low][0] = pivot;
45 envelopes[low][1] = pivot2;

    quickSort(envelopes, start, low - 1);
    quickSort(envelopes, low + 1, end);
}

50 private int binarySearch(List<Integer> list, int start, int end, int target){
    int low = start, high = end;
    int mid = start + (end - start) / 2;
    while (low < high){
55     mid = low + (high - low) / 2;
        if (list.get(mid) < target){
            low = mid + 1;
        }
        else if (list.get(mid) > target){
60             high = mid - 1;
        }
        else{
            return mid;
        }
65    }
    if (list.get(low) < target){
        return low + 1;
    }
    else{
70     return low;
    }
}
}
```