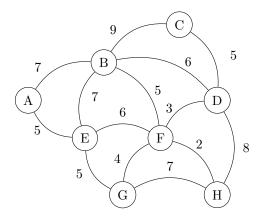
Data Structures and Algorithms Spring 2024 — Problem Sets

by Kirill Efimovich

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Week 12. Problem set

1. Run Prim-Jarník algorithm [Cormen, §21.2] on the following graph, starting at vertex C. Assuming that the algorithm is using Binary heap implementation [Cormen, §6] of a priority queue, show the state of the Binary heap after each iteration of the algorithm (i.e. after adding each new vertex to the MST). The graph contains 8 vertices, which means that your solution must provide 8 states of the Binary heap. Each heap state must be represented as an array. No justification is required.



Answer (on next page):

| Current vertex | Possible vertices | Weights |
|----------------|-------------------|------------|
| С | B, D | 9, 5 |
| D | B, F, H | 6, 3, 8 |
| F | B, E, G, H | 5, 6, 4, 2 |
| Н | B, E, G | 5, 6, 4 |
| G | В, Е | 5, 5 |
| В | A, E | 7, 5 |
| Е | A | 5 |
| A | | |

2. Suppose that all edge weights in a graph are integers in the range from 1 to |V|. How fast can you make Kruskal's algorithm run by modifying it somehow? What if the edge weights are integers in the range from 1 to W for some constant W? Justify your answer in at most two paragraphs.

Answer:

Kruskal's algorithm sorts edges in nondecreasing order by weight. If the edge weights are integers in the range 1 to V, we can use Counting-Sort to sort the edges in (V+E) time. Then Kruskal's algorithm will run in $O(V+E+V\log V)=O(E+V\log V)$ time. If the edge weights are integers in the range from 1 to W for some constant W, we can use Counting-Sort to sort the edges in (W+E) time and Kruskal's algorithm will run in $O(W+E+V\log V)$ time.

References

[Cormen] T. H. Cormen, C. E. Leiserson, R. L. Rivest and C. Stein. Introduction to Algorithms, Fourth Edition. The MIT Press 2022