Compilers (Racket)
CSCI P423/523, Fall 2022

Midterm

Name: _____

This exam has 9 questions, for a total of 100 points.

1. [10 points] Given the grammar below for expressions, indicate which of the following programs are in the language specified by the grammar. That is, which can be parsed as the *exp* non-terminal.

```
exp ::= int | (read) | (- exp) | (if exp exp exp) | var | (eq? exp exp)
      | (let ([var exp]) exp)
```

   1. `(- (if (eq? (read) 0) 10 20))`

   2. ```
      (let ([x (if (eq? (read) 0) (- 10))])
        (+ x 10))
      ```

   3. `(eq? (- (- (- (read)))) (- (- (- 10))))`

   4. `(let ([x (if (eq? (read) 0) 10 (- 10))]))`

   5. `(- (read))`

Solution: 2 points each

1. Yes.
2. No, one-armed `if` is not an *exp*.
3. Yes.
4. No, the `let` is incomplete, missing its body.
5. Yes.
      ```

2. [12 points] Convert the following program to its Abstract Syntax Tree representation (see the grammar for $\mathcal{L}_{\mathsf{If}}$ in the Appendix of this exam) and draw the tree.

```
(let ([x (if (eq? (read) 0) 5 (- (read)))])
   (+ x 42))
```

Solution:

```
(Program '()
 (Let 'x
  (If (Prim 'eq? (list (Prim 'read '()) (Int 0)))
      (Int 5)
      (Prim '- (list (Prim 'read '()))))
  (Prim '+ (list (Var 'x) (Int 42)))))
```

3. [12 points] The following is a partial impelmentation of the type checker for expressions of the $\mathcal{L}_{\mathsf{if}}^{mon}$ language, which includes integers, Booleans, conditionals, and several primitive operations. The `env` parameter is a dictionary that maps every in-scope variable to a type. Fill in the blanks of this type checker.

```
(define (type-check-exp env)
  (lambda (e)
    (match e
      [(Bool b) (values (Bool b) 'Boolean)]
      [(Let x e body)
       (define-values (e^ Te) ((type-check-exp env) e))
       (define-values (b Tb) ((type-check-exp ___(a)___) body))
       (values (Let x e^ b) ___(b)___)]
      [(If cnd thn els)
       (define-values (cnd^ Tc) ((type-check-exp env) cnd))
       (define-values (thn^ Tt) ___(c)___)
       (define-values (els^ Te) ((type-check-exp env) els))
       (check-type-equal? Tc ___(d)___)
       (check-type-equal? Tt ___(e)___)
       (values (If cnd^ thn^ els^) ___(f)___)]
      ...)))
```

Solution: (2 points each)

```
(a) (dict-set env x Te)
(b) Tb
(c) ((type-check-exp env) thn)
(d) 'Boolean
(e) Te
(f) Te  (or Tt)
```

4. [12 points] Fill in the blanks to complete the cases for `If` in the following implementation of `rco-exp` and `rco-atom` (Remove Complex Operands), which translate from the $\mathcal{L}_{\mathsf{If}}$ language into $\mathcal{L}_{\mathsf{if}}^{mon}$. The grammars for these languages can be found in the Appendix of this exam. Recall that `rco-atom` must produce an atomic expression and an association list of variables and expressions. `rco-exp` returns an expression (which does not have to be atomic).

```
(define (rco-atom e)
  (match e
    [(Let x rhs body)
     (define new-rhs (rco-exp rhs))
     (define-values (new-body body-ss) (rco-atom body))
     (values new-body (append ___(a)___ body-ss))]
    [(Bool b) (values (Bool b) '())]
    [(If cnd thn els)
     (define if-exp (If ___(b)___ (rco-exp thn) (rco-exp els)))
     (define tmp (gensym 'tmp))
     (values ___(c)___ `((,tmp . ,___(d)___)))]
    ...))

(define (rco-exp e)
  (match e
    [(Let x rhs body)
     (Let x (rco-exp rhs) (rco-exp body))]
    [(Bool b) (Bool b)]
    [(If cnd thn els)
     (define cnd^ ___(e)___)
     (define thn^ (rco-exp thn))
     (define els^ (rco-exp els))
     ___(f)___]
    ...))
```

Solution: (2 points each)

```
(a) `((,x . ,new-rhs))
(b) (rco-exp cnd)
(c) (Var tmp)
(d) if-exp
(e) (rco-exp cnd)
(f) (If cnd^ thn^ els^)
```

5. [10 points] Translate the following $\mathcal{L}_{\mathsf{if}}^{mon}$ program into $\mathcal{C}_{\mathsf{lf}}$. The grammar for $\mathcal{C}_{\mathsf{lf}}$ is in the Appendix of this exam.

```
(if (let ([tmp7 (read)])
       (eq? tmp7 0))
    (let ([tmp8 (read)])
       (- tmp8))
    (read))
```

Solution: (Approx. 1 point per statement.)

```
start:
    tmp7 = (read);
    if (eq? tmp7 0)
       goto block9;
    else
       goto block0;
block0:
    return (read);
block9:
    tmp8 = (read);
    return (- tmp8);
```

6. [14 points] Given the following psuedo-x86 program, compile it to an equivalent and complete x86 program, using stack locations (not registers) for the variables. Your answer should be given in the AT&T syntax that the GNU assembler expects for .s files.

```
start:
    callq read_int
    movq %rax, x
    movq $-4, t0
    movq t0, t1
    addq x, t1
    movq t1, %rdi
    callq print_int
    movq $0, %rax
    jmp conclusion
```

Solution:

```
        .globl main
main:
    pushq %rbp
    movq %rsp, %rbp
    subq $16, %rsp
    jmp start
start:
    callq read_int
    movq %rax, -16(%rbp)
    movq $-4, -8(%rbp)
    movq -16(%rbp), %rax
    addq %rax, -8(%rbp)
    movq -8(%rbp), %rdi
    callq print_int
    movq $0, %rax
    jmp conclusion
conclusion:
    addq $16, %rsp
    popq %rbp
    retq
```

Rubric:

- prelude (3 points)
- correct use of the stack for the variables (2 points)
- call to `read_int` and move from `rax` (2 point)
- at most one memory argument per instruction (2 points)
- move to `rdi` and call to `print_int` (2 points)
- conclusion (3 points)

7. [10 points] Apply liveness analysis to the following pseudo-x86 program to determine the set of live locations before and after every instruction. (The callee and caller saved registers are listed in the Appendix of this exam.)

```
start:
    movq $0, sum

    movq $5, i

    jmp block.0


block.0:
    cmpq $0, i

    jg block.2

    jmp block.3


block.3:
    jmp block.1
```

```
block.2:
    addq i, sum

    subq $1, i

    jmp block.0


block.1:
    movq sum, %rdi

    callq print_int

    movq $0, %rax

    jmp conclusion
```

Solution:

```
block.1:
                {sum}
  movq sum, %rdi
                {%rdi}
  callq print_int
                {}
  movq $0, %rax
                {%rax}
  jmp conclusion
                {%rax}
block.2:
                {sum, i}
  addq i, sum
                {sum, i}
  subq $1, i
                {sum, i}
  jmp block.0
                {sum, i}
```

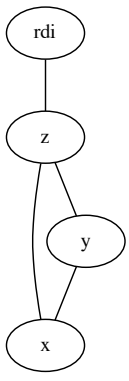
```
block.3:
                {sum}
  jmp block.1
                {sum}
block.0:
                {sum, i}
  cmpq $0, i
                {sum, i}
  jg block.2
                {sum, i}
  jmp block.3
                {sum, i}
start:
                {}
  movq $0, sum
                {sum}
  movq $5, i
                {sum, i}
  jmp block.0
                {sum, i}
```

8. [10 points] Given the following results from liveness analysis, draw the interference graph. (The callee and caller saved registers are listed in the Appendix of this exam.)

```
start:
                {}
  callq _read_int
                {%rax}
  movq %rax, x
                {x}
  movq x, y
                {y, x}
  addq $1, y
                {y, x}
  movq y, z
                {y, x, z}
  addq $1, z
                {y, z, x}
  cmpq $0, x
                {y, x, z}
  je block.1
                {y, x, z}
  jmp block.2
                {y, z, x}
```

```
block.1:
                {x, z}
  movq x, %rdi
                {%rdi, z}
  callq print_int
                {z}
  jmp block.0
                {z}
block.2:
                {y, z}
  movq y, %rdi
                {%rdi, z}
  callq print_int
                {z}
  jmp block.0
                {z}
block.0:
                {z}
  movq z, %rdi
                {%rdi}
  callq print_int
                {}
  movq $0, %rax
                {%rax}
  jmp conclusion
                {%rax}
```

Solution:



9. [10 points] Fill in the blanks to complete the following graph coloring algorithm.

```
(define (make-pqueue <=? [init '()]) ...)
(define (pqueue-push! q key) ...)
(define (pqueue-pop! q) ...)
(define (pqueue-decrease-key! q node) ...)
(define (pqueue-count q) ...)

(define (color-graph interfere-graph move-graph info)
  (define locals (dict-keys (dict-ref info 'locals-types)))
  (define unavailable-colors (make-hash))
  (define (compare u v)
    (>= (set-count (hash-ref unavailable-colors u))
        (set-count (hash-ref unavailable-colors v))))
  (define Q (make-pqueue ___(a)___))
  (define pq-node (make-hash))
  (define color (make-hash))
  (for ([r registers-for-alloc])
    (hash-set! color r (register->color r)))
  (for ([x locals])
       (define adj-reg
          (filter (lambda (u) (set-member? registers u))
                  (get-neighbors interfere-graph x)))
       (define adj-colors (list->set (map register->color adj-reg)))
       (hash-set! unavailable-colors x adj-colors)
       (hash-set! pq-node x ___(b)___))
  (while ___(c)___
         (define v (pqueue-pop! Q))
         (define move-related
           (sort (filter (lambda (x) (>= x 0))
                         (map (lambda (k) (hash-ref color k -1))
                              (get-neighbors move-graph v)))
                 <))
         (define c (choose-color v (hash-ref unavailable-colors v)
                                 move-related info))
         (hash-set! color v c)
         (for ([u ___(d)___])
              (when (not (set-member? registers u))
                    (hash-set! unavailable-colors u ___(e)___)
                    (pqueue-decrease-key! Q (hash-ref pq-node u)))))
  color)
```

Solution: (2 points each)

```
(a) compare
(b) (pqueue-push! Q x)
(c) (> (pqueue-count Q) 0)
(d) (in-neighbors interfere-graph v)
(e) (set-add (hash-ref unavailable-colors u) c)
```

Appendix

The caller-saved registers are:

```
rax rcx rdx rsi rdi r8 r9 r10 r11
```

and the callee-saved registers are:

```
rsp rbp rbx r12 r13 r14 r15
```

Grammar for $\mathcal{L}_{\mathsf{If}}$

type	::=	`Integer`
op	::=	`read` | `+` | `-`
exp	::=	`(Int` *int*`)` | `(Prim` *op* `(`*exp* `...))`
exp	::=	`(Var` *var*`)` | `(Let` *var exp exp*`)`
type	::=	`Boolean`
bool	::=	`#t` | `#f`
cmp	::=	`eq?` | `<` | `<=` | `>` | `>=`
op	::=	*cmp* | `and` | `or` | `not`
exp	::=	`(Bool` *bool*`)` | `(If` *exp exp exp*`)`
$\mathcal{L}_{\mathsf{If}}$	::=	`(Program '()` *exp*`)`

Grammar for $\mathcal{L}_{\mathsf{if}}^{mon}$

atm	::=	`(Int` *int*`)` | `(Var` *var*`)`
exp	::=	*atm* | `(Prim 'read ())`
	|	`(Prim '-` (*atm*)) | `(Prim '+` (*atm atm*)) | `(Prim '-` (*atm atm*))
	|	`(Let` *var exp exp*`)`
atm	::=	`(Bool` *bool*`)`
exp	::=	`(Prim not` (*atm*)) | `(Prim` *cmp* (*atm atm*)) | `(If` *exp exp exp*`)`
$\mathcal{L}_{\mathsf{if}}^{mon}$	::=	`(Program () `*exp*`)`

Grammar for $\mathcal{C}_{\mathsf{If}}$

atm	::=	`(Int` *int*`)` | `(Var` *var*`)`
exp	::=	*atm* | `(Prim 'read ())` | `(Prim '-` (*atm*))
	|	`(Prim '+` (*atm atm*)) | `(Prim '-` (*atm atm*))
stmt	::=	`(Assign (Var` *var*`)` *exp*`)`
tail	::=	`(Return` *exp*`)` | `(Seq` *stmt tail*`)`
atm	::=	`(Bool` *bool*`)`
cmp	::=	`eq?` | `<` | `<=` | `>` | `>=`
exp	::=	`(Prim 'not` (*atm*)) | `(Prim '`*cmp* (*atm atm*))
tail	::=	`(Goto` *label*`)`
	|	`(IfStmt (Prim` *cmp* (*atm atm*)) `(Goto` *label*`) (Goto` *label*`))`
$\mathcal{C}_{\mathsf{If}}$	::=	`(CProgram` *info* `((`*label* `.` *tail*`) ...))`