### Theoretical computer science

Ttutorial - week 8

March 11, 2021

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# Agenda

Operations on DPDA

**▶** Union

- Union
- ► Intersection

- ► Union
- Intersection
- Difference

- ► Union
- ► Intersection
- Difference
- ► Complement

Suppose  $L_1$  and  $L_2$  are both languages over the alphabet A. If  $x \in A^*$ , then knowing whether  $x \in L_1$  and whether  $x \in L_2$  is enough to determine whether  $x \in L_1 \cup L_2$ .

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▶  $L_1 \cup L_2$ ?

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- $ightharpoonup L_1 \cup L_2$ ?
- $ightharpoonup L_1 \cap L_2$ ?

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- $ightharpoonup L_1 \cup L_2$ ?
- $ightharpoonup L_1 \cap L_2$ ?
- $ightharpoonup L_1 \backslash L_2$ ?

Suppose A and B are both languages over the alphabet  $\Sigma = \{a, b, c\}$ :

$$AnBn = \{a^n b^n | n \ge 1\}$$

$$AnB2n = \{a^n b^{2n} | n \ge 1\}$$

What is the language recognized by  $AnBn \cup AnB2n$ ?

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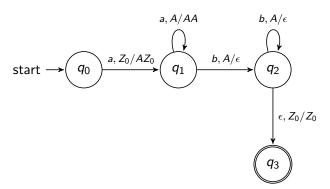
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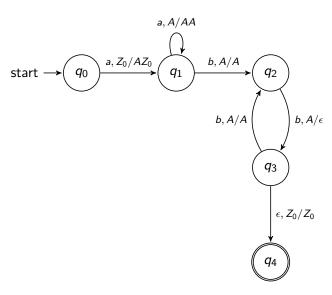
$$AnBn \cup AnB2n = \{a^n b^n | n \ge 1\} \cup \{a^n b^{2n} | n \ge 1\}$$

This language cannot be recognized by DPDA, thus the class of languages recognized by PDA is not closed under intersection

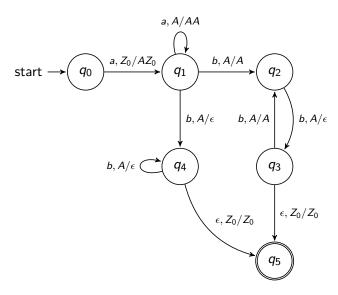
#### AnBn



#### AnB2n



# AnBn ∪ AnB2n (non-deterministic)



Suppose A and B are both languages over the alphabet

$$\Sigma = \{a, b, c\}$$
:

$$A = \{a^n b^n c^m | n, m \ge 0\}$$

$$B = \{a^m b^n c^n | n, m \ge 0\}$$

What is the language recognized by  $A \cap B$ ?

Suppose A and B are both languages over the alphabet  $\Sigma = \{a, b, c\}$ :

$$A = \{a^n b^n c^m | n, m \ge 0\}$$

$$B = \{a^m b^n c^n | n, m \ge 0\}$$

What is the language recognized by  $A \cap B$ ?

$$A \cap B = \{a^n b^n c^n | n \ge 0\}$$

This language cannot be recognized by DPDA, thus the class of languages recognized by PDA is not closed under intersection

Suppose A and B are both languages over the alphabet  $\Sigma = \{a, b, c\}$ :

$$A = \{a^n b^n c^m | n, m \ge 0\}$$
$$B = \{a^m b^n c^n | n, m > 0\}$$

For complement, note that  $A \cap B = \overline{\overline{A} \cup \overline{B}}$ 

- ► If context-free languages were closed under complement, they would also be closed under intersection
- Therefore context-free languages are not closed under complementation because they are not closed under intersection

Languages accepted by Deterministic PDA have closure under ...?

► Union -

- Union not closed
- Intersection -

- Union not closed
- Intersection not closed
- Difference -

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- Complement -

- Union not closed
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# Wrap up

▶ What have you learnt today?

## Wrap up

- ► What have you learnt today?
- ▶ What for this could be useful?