# Theoretical Computer Science Lab Session 1

January 28, 2021

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## Agenda

- ► Introduction (rules of the game!)
- Preliminaries Sets
- Exercises on Formal Languages

Laboratory Exercises: There are weekly laboratory exercises.

Assessment: Mid-term Exam (25%), Final Exam (25%),

Technical Assignments (45%), and Labs (5%).

Extra points: Tutorials (5%).

#### Group switching

Switching lab groups is allowed under the following conditions:

- ► The group size limit is 30 students.
- Students who switch a lab group are not eligible for TA points.

#### Policy and Procedures on Cheating and Plagiarism

Exam policy: If two or more students are caught communicating for any reason during exams (including mid-terms) they will be asked to leave the room and their exam will be failed. Same will happen for unauthorized use of devices.

#### Policy and Procedures on Cheating and Plagiarism

Report policy: If a submitted report contains work other than student's one it is necessary to explicitly acknowledge the source. It is encouraged to refer and quote other works, but it has to be made clear which words and ideas are property and creation of the student, and which ones have come from others (which must not correspond to more than 30% of the work). If two or more reports show evidence of being produced by unauthorized cooperative work, i.e. copied from fellow students, they will be all failed without further investigation on who produced the results and who actually copied.

Preliminaries - Sets

#### Sets

A finite set can be described, at least in principle, by listing its elements:  $A = \{1, 2, 4, 8\}$  says that A is the set whose elements are 1, 2, 4, and 8.

For infinite (even for finite sets if they have more than just a few elements) sets ellipses (...) are sometimes used to describe how the elements might be listed:

$$B = \{0, 3, 6, 9, \ldots\}$$
  
 
$$C = \{13, 14, 15, \ldots, 71\}$$

#### Sets

A more reliable way is to give the property that characterises their elements (also called set comprehension). Sets  $B = \{0, 3, 6, 9, \ldots\}$  and  $C = \{13, 14, 15, \ldots, 71\}$  can be described as

 $B = \{x \mid x \text{ is a non-negative integer multiple of 3}\}$   $C = \{x \mid x \text{ is an integer and } 13 \leq x \leq 71\}$ It reads: "B is the set of all x such that x is a non-negative integer multiple of 3"

# Sets (exercise)

What are the sets D and E?:

```
D = \{\{x\} \mid x \text{ is a non-negative integer such that } x \leq 4\}

E = \{3i + 5j \mid i \text{ and } j \text{ are non-negative integers}\}
```

# Sets (exercise)

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$$D = \{\{0\}, \{1\}, \{2\}, \{3\}, \{4\}\}\}$$

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$$D = \{\{0\}, \{1\}, \{2\}, \{3\}, \{4\}\} \}$$
  
$$E = \{0, 3, 5, 6, 8, 9, 10, \ldots\}$$

# Sets (operations)

- For any set A, the statement that x is an element of A is written  $x \in A$ .
- ▶  $A \subseteq B$  means that A is a subset of B: every element of A is an element of B.
- ▶ ∅ denotes the empty set: the set with no elements.

To show that two sets A and B are the same, we must show that A and B have exactly the same elements, i.e.  $A \subseteq B$  and  $B \subseteq A$ .

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To show that two sets A and B are the same, we must show that A and B have exactly the same elements, i.e.  $A \subseteq B$  and  $B \subseteq A$ . Are the following statements true?

```
 \{0,1\} = \{1,0\} \\ \{0,1,2,1,0\} = \{1,1,1,1,0,2,2\}
```

# Sets (operations)

For two sets A and B, we can define their union  $A \cup B$ , their intersection  $A \cap B$ , and their difference  $A \setminus B$  (sometimes denoted as A - B), as follows<sup>1</sup>:

$$A \cup B = \{x \mid x \in A \lor x \in B\}$$
  

$$A \cap B = \{x \mid x \in A \land x \in B\}$$
  

$$A \backslash B = \{x \mid x \in A \land x \notin B\}$$

 $<sup>^{1}\</sup>lor$  and  $\land$  denote the logical 'or' and logical 'and' respectively.

# Sets (Union of any number of sets) - Notation

If  $A_0$ ,  $A_1$ ,  $A_2$ , ... are sets, the union of these sets can be denoted as

$$\bigcup \{A_i \mid i \ge 0\} = \{x \mid x \in A_i \text{ for at least one } i \text{ with } i \ge 0\}$$

or

$$\bigcup_{i=0}^{\infty} A_i$$

# Sets (Power Sets)

For a set A, the set of all subsets of A is called the power set. Can be denoted as  $\mathcal{P}(A)$  or as  $2^A$ .

Power set of set  $\{a, b, c\}$  is

$$\{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$$

For a set A, the set  $\mathcal{P}(A)$  has exactly  $2^n$  elements, where n is the cardinality of A.

Exercise Session - Languages

## Notation and Terminology

```
a finite set of symbols, e.g. \{a, b\}, or \{0, 1\}.
  Alphabet:
                Normally denoted by \Sigma
              a string over an alphabet (\Sigma) is a finite sequence of
     String:
               symbols in \Sigma.
     length: for a string x, |x| is the number of symbols of x.
empty string: is the null string over \Sigma. It is denoted as \epsilon. By
                definition, |\epsilon| = 0
Set of all strings: the set of all strings over \Sigma is denoted by \Sigma^*,
               e.g. for the alphabet A = \{a, b\}
               A^* = \{\epsilon, a, b, aa, ab, bb, aaa, aab, \ldots\}
```

## Concatenation of strings

If x and y are two strings over an alphabet, the concatenation xy (sometimes denoted as  $x \cdot y$ ) consists of the symbols of x followed by those of y:

$$x = ab$$
  
 $y = bab$   
 $xy = abbab$ 

Concatenation is an associative operation: (xy)z = x(yz) for all possible strings x, y, and z.

Languages are sets.

- Poperations on languages are ways of constructing new languages: for two languages  $L_1$  and  $L_2$  over the alphabet  $\Sigma$ ,  $L_1 \cup L_2$ ,  $L_1 \cap L_2$ , and  $L_1 \setminus L_2$  are also languages over  $\Sigma$ .
- ▶ String operation of concatenation is also used to construct new languages: if  $L_1$  and  $L_2$  are both languages over  $\Sigma$ , the concatenation of  $L_1$  and  $L_2$  is the language

$$L_1L_2 = \{xy \mid x \in L_1, y \in L_2\}$$

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Example:

$$\{a,aa\}\{\epsilon,b,ab\} =$$

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Example:

$$\{a,aa\}\{\epsilon,b,ab\}=\{a,ab,aab,aa,aab,aaab\}$$

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Example:

$$\{a,aa\}\{\epsilon,b,ab\}=\{a,ab,aab,aa,aab,aaab\}$$

Is this statement true?

$$L_1L_2=L_2L_1$$

#### Exponential notation

The concatenation of k copies of a single symbol a, a single string s, or a single language L is defined as:

If k = 0, then

$$a^k = \{\epsilon\}$$

If k > 0, then

$$a^k = aa \dots a$$

where there are k occurrences of a, similarly for  $s^k$  and  $L^k$ . In the case where L is simply the alphabet  $\Sigma$ ,

$$\Sigma^k = \{ x \in \Sigma^* \mid |x| = k \}$$

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Example:

$$\begin{array}{l} \Sigma = \{0,1\} \\ \Sigma^2 = \end{array}$$

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$$\Sigma^k = \{ x \in \Sigma^* \mid |x| = k \}$$

Example:

$$\begin{split} \Sigma &= \{0,1\} \\ \Sigma^2 &= \{00,01,10,11\} \end{split}$$

# Exercises (1)

Construct the power set for the following sets:

```
i \{a, b\}
   ii \{0,1\} \cup \{1,2\}
  iii \{z\}
  iv \{0,1,2,3,4\} \cap \{1,3,5,a\}
   \vee \{0,1,2,3\} \setminus \{1,3,5,a\}
  vi Ø
Determine the following languages over the alphabet \Sigma = \{0,1\}
 vii \Sigma^0
viii \Sigma^4
  ix \mathcal{P}(\Sigma)
   \times \mathcal{P}(\Sigma^*)
```

# Solution (1)

Construct the power set for the following sets:

```
\begin{split} & \mathrm{i} \ \mathcal{P}(\{a,b\}) = \{\emptyset,\{a\},\{b\},\{a,b\}\} \\ & \mathrm{ii} \ \mathcal{P}(\{0,1\} \cup \{1,2\}) = \mathcal{P}(\{0,1,2\}) = \\ & \{\emptyset,\{0\},\{1\},\{2\},\{0,1\},\{0,2\},\{1,2\},\{0,1,2\}\} \\ & \mathrm{iii} \ \mathcal{P}(\{z\}) = \{\emptyset,\{z\}\} \\ & \mathrm{iv} \ \mathcal{P}(\{0,1,2,3,4\} \cap \{1,3,5,a\}) = \\ & \mathcal{P}(\{1,3\}) = \{\emptyset,\{1\},\{3\},\{1,3\}\} \\ & \mathrm{v} \ \mathcal{P}(\{0,1,2,3\} \backslash \{1,3,5,a\}) = \mathcal{P}(\{0,2\}) = \{\emptyset,\{0\},\{2\},\{0,2\}\} \\ & \mathrm{vi} \ \mathcal{P}(\emptyset) = \{\emptyset\} \end{split}
```

# Solution (1)

```
Determine the following languages over the alphabet \Sigma = \{0,1\}
 vii \Sigma^0 = \{\epsilon\}
viii
\Sigma^4 = \{0000, 0001, 0010, 0011,
0100, 0101, 0110, 0111,
 1000, 1001, 1010, 1011,
1100, 1101, 1110, 1111}
```

ix 
$$\mathcal{P}(\Sigma) = \{\emptyset, \{0\}, \{1\}, \{0, 1\}\}\$$
  
  $\times \mathcal{P}(\Sigma^*) = (\text{Infinite set which includes empty set and all combinations of 0s and 1s})$ 

# Exercises (2)

Find a possible alphabet for the following languages<sup>2</sup>

- i The language  $L = \{oh, ouch, ugh\}$
- ii The language  $L = \{apple, pear, 4711\}$
- iii The language of all binary strings

Determine what the Kleene star operation produces over the following alphabets:

- iv  $\Sigma = \{0, 1\}$
- $v \Sigma = \{a\}$
- vi  $\Sigma = \emptyset$  (the empty alphabet)

 $<sup>^{2}</sup>$ A word foo should be interpreted as a string of characters f, o, and o.

# Solution (2)

Find a possible alphabet for the following languages<sup>3</sup>

- i The language  $L = \{oh, ouch, ugh\}$ :  $\Sigma = \{o, h, u, c, g\}$
- ii The language  $L = \{apple, pear, 4711\}$ :  $\Sigma = \{a, p, l, e, r, 4, 7, 1\}$
- iii The language of all binary strings:  $\Sigma = \{0, 1\}$

Determine what the Kleene star operation produces over the following alphabets:

- iv  $\Sigma = \{0,1\}$ : All binary strings
- v  $\Sigma = \{a\}$ : All strings which contain nothing but a's
- vi  $\, \Sigma = \emptyset$  (the empty alphabet): the language that contains only the empty string

 $<sup>^{3}</sup>$ A word foo should be interpreted as a string of characters f, o, and o.

# Exercises (3)

State the alphabet  $\Sigma$  for the following languages:

i 
$$L = \Sigma^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, \ldots\}$$

ii 
$$L = \Sigma^* = \{\epsilon, a, aa, aaa, aaaa, \ldots\}$$

Assuming that  $\Sigma = \{0,1\}$  , construct complement languages for the following:

iii 
$$\overline{\{010,101,11\}}$$

iv 
$$\overline{\Sigma^* \setminus \{110\}}$$

State the following languages explicitly

v 
$$\mathcal{P}(\{a,b\})\backslash\mathcal{P}(\{a,c\})$$

vi  $\{x \mid x, y \in \mathbb{N} \land \exists y : y < 10 \land (y+2=x)\}$  ( $\mathbb{N}$  is the set of all non-negative integers)

# Solution (3)

State the alphabet  $\Sigma$  for the following languages:

```
\begin{array}{l} \text{i} \ \ L = \Sigma^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, \ldots\} \\ \Sigma = \{0, 1\} \\ \text{ii} \ \ L = \Sigma^* = \{\epsilon, a, aa, aaa, aaaa, \ldots\} \\ \Sigma = \{a\} \end{array}
```

Assuming that  $\Sigma=\{0,1\},$  construct complement languages for the following:

# Solution (3)

State the following languages explicitly

$$V \mathcal{P}(\{a,b\}) \backslash \mathcal{P}(\{a,c\})$$
$$L = \{\{b\}, \{a,b\}\}$$

vi  $\{x \mid x, y \in \mathbb{N} \land \exists y : y < 10 \land (y+2=x)\}$  ( $\mathbb{N}$  is the set of all non-negative integers)

$$L = \{2, 3, 4, 5, 6, 7, 8, 9, 10, 11\}$$