

# Chapter 1

## Library SemanticEquivalence

```
Require Import Coq.Init.Peano.
Require Import Coq.Init.Nat.
Require Import Coq.Arith.EqNat.
Require Import Coq.Arith.PeanoNat.
Include Coq.Init.Nat.
Check _/_.

Require Import List.
Import ListNotations.
Check [1].

Definition stack := list nat.

Inductive bit : Set :=
  | I : bit
  | O : bit.

Definition bitstr := list bit.

Inductive expr (ops : Set) (lits : Set) : Set :=
  | Lit : lits → expr ops lits
  | Bin : ops → expr ops lits → expr ops lits → expr ops lits.

Arguments Lit {ops} {lits}.
Arguments Bin {ops} {lits}.

Check Lit 1.

Inductive semantics (L : Set) (D : Set) : Type :=
  sem_f : (L → D) → semantics L D.

Arguments sem_f {L} {D}.

Inductive semantic_evaluation {L D : Set} (sem : L → D) | (e : L) (d : D) : Prop
  := sem_eval : sem e = d → semantic_evaluation e d.

Definition sem {L D} (s : semantics L D) : L → D :=
```

let 'sem\_f sem\_fun := s in sem\_fun.

Notation "[[ x ]][ s ]" := (sem s x).

Fixpoint *expr\_semantics'*

{ops lits : Set} {D : Set} (sem\_lit : lits → D)  
 (sem\_op : ops → D → D → D) (e : expr ops lits) : D :=  
 let expr\_sem := *expr\_semantics'* sem\_lit sem\_op in  
 match e with  
 | Lit x ⇒ sem\_lit x  
 | Bin op e1 e2 ⇒ sem\_op op (expr\_sem e1) (expr\_sem e2)  
 end.

Definition *expr\_semantics*

{ops lits D : Set} (sem\_lit : lits → D) (sem\_op : ops → D → D → D) : semantics (expr ops lits) D :=  
 sem\_f (*expr\_semantics'* sem\_lit sem\_op).

Inductive *expr\_semantics\_definition* (lits ops D : Set) : Set :=

| *expr\_sem\_def\_parts* : ∀ (sem\_lit : lits → D) (sem\_op : ops → D → D → D),  
*expr\_semantics\_definition* lits ops D.

Inductive *bin\_ops* : Set := *bin\_add*.

Definition *expr\_bits* := *expr bin\_ops bitstr*.

Fixpoint *expr\_bits\_sem\_lits* (bs : bitstr) : nat :=

match bs with  
 | I :: bs' ⇒ 1 + (2 × *expr\_bits\_sem\_lits* bs')  
 | O :: bs' ⇒ 2 × *expr\_bits\_sem\_lits* bs'  
 | [] ⇒ 0  
 end.

Definition *expr\_bits\_sem\_op* (o : bin\_ops) : nat → nat → nat :=

match o with  
 | *bin\_add* ⇒ fun a b ⇒ a + b  
 end.

Definition *bin\_expr\_sems* := *expr\_semantics expr\_bits\_sem\_lits expr\_bits\_sem\_op*.

Check *bin\_expr\_sems* : semantics *expr\_bits* nat.

Definition *eval\_expr* := *expr\_semantics' expr\_bits\_sem\_lits expr\_bits\_sem\_op*.

Fixpoint *bits\_to\_nat* (b : bitstr) : nat :=

match b with  
 | I :: b' ⇒ S (2 × *bits\_to\_nat* b')  
 | O :: b' ⇒ (2 × *bits\_to\_nat* b')  
 | [] ⇒ 0  
 end.

Definition *nat\_to\_bit* ( $n : \text{nat}$ ) : *bit* :=

```

match n with
| 0 => O
| 1 => I
| _ => O
end.

```

Fixpoint *bit\_inc* ( $bs : \text{bitstr}$ ) : *bitstr* :=

```

match bs with
| [] => [I]
| I :: bs' => O :: bit_inc bs'
| O :: bs' => I :: bs'
end.

```

Fixpoint *nat\_to\_bits* ( $n : \text{nat}$ ) : *bitstr* :=

```

match n with
| 0 => [O]
| S n' => bit_inc (nat_to_bits n')
end.

```

Lemma *bits\_to\_nat\_function* :  $\forall (b1\ b2 : \text{bitstr}),\ b1 = b2 \rightarrow \text{bits\_to\_nat } b1 = \text{bits\_to\_nat } b2$ .

Fixpoint *bit\_add* ( $b1\ b2 : \text{bitstr}$ ) : *bitstr* :=

```

match b1, b2 with
| I :: b1', b :: b2' => bit_inc (b :: bit_add b1' b2')
| O :: b1', b :: b2' => b :: bit_add b1' b2'
| [], _ => b2
| _, [] => b1
end.

```

Notation " $x +. + y$ " := (*bit\_add*  $x\ y$ ) (at level 40, left associativity).

Definition *binary* := *bitstr*.

Lemma *binary\_add\_unit\_left* :  $\forall (b : \text{binary}),\ b +. + [] = b$ .

Lemma *binary\_add\_unit\_right* :  $\forall (b : \text{binary}),\ [] +. + b = b$ .

Lemma *binary\_add\_id\_left* :  $\exists (e : \text{binary}),\ \forall (b : \text{binary}),\ e +. + b = b$ .

Lemma *binary\_add\_id\_right* :  $\exists (e : \text{binary}),\ \forall (b : \text{binary}),\ b +. + e = b$ .

Lemma *add\_id* :  $\forall (n : \text{nat}),\ n + 0 = n$ .

Lemma *add\_succ\_assoc* :  $\forall (n\ m : \text{nat}),\ S\ (n + m) = (S\ n) + m$ .

Lemma *bits\_nat\_succ\_eq* :  $\forall (n : \text{nat}),\ \text{bit\_inc } (\text{nat\_to\_bits } n) = \text{nat\_to\_bits } (S\ n)$ .

Lemma *nat\_bits\_succ\_eq* :  $\forall (b : \text{bitstr}),\ S\ (\text{bits\_to\_nat } b) = \text{bits\_to\_nat } (\text{bit\_inc } b)$ .

Notation " $[[\ x\ ]]$ " := (*bits\_to\_nat*  $x$ ).

Lemma *binary\_add\_linearity\_1* :

$(\forall (b : \text{bitstr}), \text{bits\_to\_nat } ([O] ++ b) = \text{bits\_to\_nat } b \wedge \text{bits\_to\_nat } (b ++ [O]) = \text{bits\_to\_nat } b) \rightarrow$

$\forall (n : \text{nat}), \text{bits\_to\_nat } (\text{nat\_to\_bits } n ++ [O]) = n + \text{bits\_to\_nat } [O].$

Definition *str1* := *nat\_to\_bits* 2.

Definition *str2* := *nat\_to\_bits* 1.

Theorem *bit\_nat\_correspondance* :  $\forall (b \ c : \text{bitstr}), [[\ b \ ++ \ c \ ]] = [[\ b \ ]] + [[\ c \ ]].$

Theorem *nat\_to\_bit\_add\_hom* :  $\forall (n \ m : \text{nat}), \text{nat\_to\_bits } (n + m) = \text{bit\_add } (\text{nat\_to\_bits } n) (\text{nat\_to\_bits } m).$

Theorem *bit\_to\_nat\_add\_hom* :  $\forall (a \ b : \text{bitstr}), \text{bits\_to\_nat } (\text{bit\_add } a \ b) = (\text{bits\_to\_nat } a) + (\text{bits\_to\_nat } b).$

Theorem *idempotent\_nat* :  $\forall (n : \text{nat}), \text{bits\_to\_nat } (\text{nat\_to\_bits } n) = n.$

Theorem *thm\_idk* :  $\forall (n : \text{nat}) (b : \text{bitstr}), n = \text{bits\_to\_nat } b \rightarrow \text{nat\_to\_bits } n = b.$

Theorem *idempotent\_bits* :  $\forall (b : \text{bitstr}) (n : \text{nat}), n = \text{bits\_to\_nat } b \rightarrow \text{nat\_to\_bits } (\text{bits\_to\_nat } b) = b.$

Theorem *bits\_num\_bijection* :  $\forall (n : \text{nat}), \exists (b : \text{bitstr}),$

Theorem *bits\_additive\_identity* :

Theorem *list\_sem\_correctness* :  $\forall (b : \text{bitstr}) (n : \text{nat}), \text{bits\_to\_nat } b = n \rightarrow \text{eval } (\text{Lit } b) = n.$

Theorem *bit\_expr\_add\_hom* :  $\forall (b1 \ b2 : \text{bitstr}) (n \ m : \text{nat}), \text{eval } (\text{Bin } \text{bin\_add } b1 \ b2) = \text{eval } ()$