⇒ If we able to convert that parallel schedule in any of the 6 forms of serial schedule, then we can say that it is a Serializable.

⇒ Why we get 6 ~~forms~~ ! →

we have 3 values (T_1, T_2 & T_3).

So,

$3! = 6$ ways.

---

(1) Conflict Equivalent Schedules ! →

R(A)	R(A)	} Non - Conflict pair.
R(A)	W(A)	} Conflict pairs
W(A)	R(A)	
W(A)	W(A)	
R(B)	R(A)	} Non - Conflict pair.
W(B)	R(A)	
R(B)	W(A)	
W(A)	W(B)	

(2) How to Convert : →

If we have adjacent non - conflict pair, then swap their pair.

Ex :

T_1	T_2
	R(A)
R(B)	

⇒

T_1	T_2
R(B)	
	R(A)

Q:- To check Conflict Equivalent :-→

$S \equiv S'$ → check

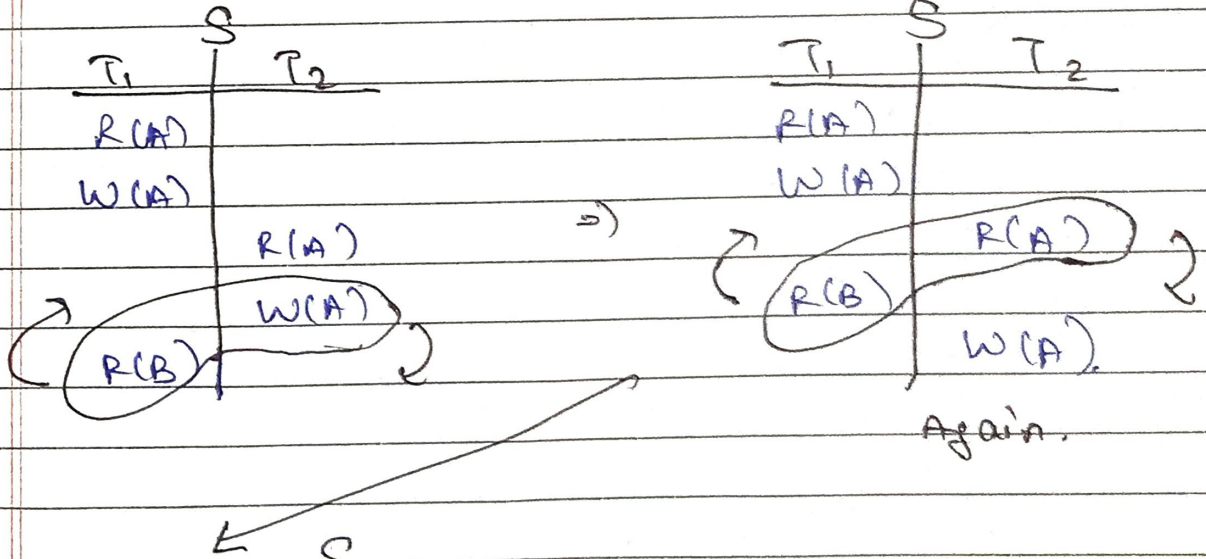
S

T_1	T_2
R(A)	
W(A)	
	R(A)
	W(A)
R(B)	

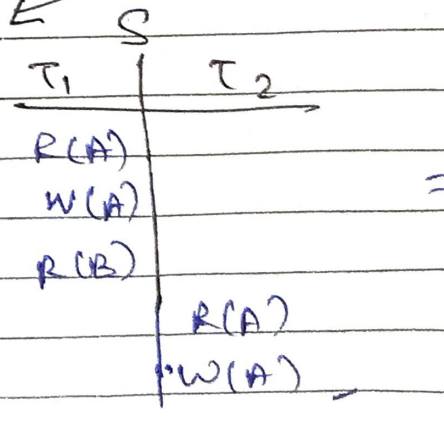
S'

T_1	T_2
R(A)	
W(A)	
R(B)	
	R(A)
	W(A)

Sol":- So, In S, we have adjacent non-conflict pair, so swap them.



Again.



Now Serial ← Schedule.

$= S'$

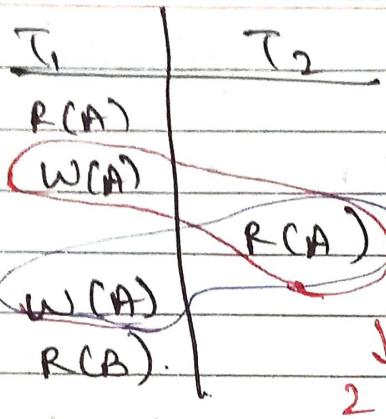
Hence,

$S \equiv S'$

Hence.

S & S' are conflict Equivalent Schedule.

Ex:-



2 adjacent pairs, But they are conflict pairs. i.e, no change in positions.

Note:-> $S \xrightarrow{CE} S' \longrightarrow$ Serializable i.e, (Serial Schedule)

(CE: conflict equivalent)

(i - 81) videos end here.