

# CSC324: Principles of Programming Languages

## Lecture 8

Wednesday, 3 June, 2020

# Comments on lab 4

- People seemed to like the recitation format! We'll stick with it going forward.
- If you didn't find the earlier labs a good use of your time, you might like the new format better.

# Comments on exercise 3

Many of you only implemented functionality that deconstructs only one kind of Expr

```
26 data Expr
27   = Number Integer      -- ^ A numeric literal.
28   | Identifier String   -- ^ An identifier.
29   deriving (Show, Eq)
30
31 data Binding
32   = Binding String Expr -- ^ A binding of an identifier name to an expression.
33   deriving (Show, Eq)
34
35 someFunc Binding s (Number n) = ...
36 -- What about if the expression isn't a number?
37
38
```

# Comments on exercise 3

Many of you only implemented functionality that deconstructs only one kind of Expr

```
26 data Expr
27   = Number Integer      -- ^ A numeric literal.
28   | Identifier String   -- ^ An identifier.
29   deriving (Show, Eq)
30
31 data Binding
32   = Binding String Expr -- ^ A binding of an identifier name to an expression.
33   deriving (Show, Eq)
34
35 someFunc Binding s (Number n) = ...
36 someFunc Binding s (Identifier id) = ...
37
38
```

# In previous classes...

... we discussed how **quotation** allows a strictly-evaluated language to change an expression's evaluation semantics, and how we call an expression with non-standard evaluation rules a **special form**.

Today, we'll talk about how to implement special forms of our own

Zzzzzz.....

How long will it take for  
this call to snooze to  
return?

Who knows??? We don't  
have units, argh

```
#lang racket
```

```
; sleep pauses the execution of the program  
; for the amount of time given.
```

```
;   
; int -> void  
(define (sleep n)  
  ...)
```

```
; it's 12:00:00 right now. When do we resume?  
(sleep 60)
```

# Zzzzzzz.....

One way to solve this is to make it explicit in the documentation how long we'll sleep for

```
#lang racket
```

```
; sleep pauses the execution of the program  
; for the amount of time given, in seconds.
```

```
;
```

```
; int -> void
```

```
(define (sleep n)  
  ...)
```

```
; it's 12:00:00 right now.
```

```
(sleep 60)
```

```
; it's now 12:01:00.
```

# Zzzzzz.....

Maybe we could define a time-period datatype to let callers assign a unit to the time period?

Let's use the template for time-period to figure out how to write snooze.

```
#lang racket
```

```
; a time-period is a unit of time, and is one of:  
; 'h : a symbol representing an hour (60 * 60 seconds)  
; 'm : a symbol representing a minute (60 seconds)  
; 's : a symbol representing a second  
; 'ms : a symbol representing a millisecond (1/1000 seconds)
```

```
; snooze pauses the execution of the program for the  
; amount of time given.
```

```
;
```

```
; int time-period -> void
```

```
(define (snooze n tp)
```

```
  (let [(in-seconds (match tp  
                        ('h (* n 60 60))  
                        ('m (* n 60))  
                        ('s n)  
                        ('ms (/ n 1000)))))]
```

```
    (sleep in-seconds)))
```

```
(snooze 60 'ms) ; now we know how long we'll snooze for
```



# Zzzzzz.....

Here, we used a **symbol**, which is a particular kind of **value**, to indicate what unit of time we are operating on.

```
#lang racket
```

```
; a time-period is a unit of time, and is one of:  
; 'h : a symbol representing an hour (60 * 60 seconds)  
; 'm : a symbol representing a minute (60 seconds)  
; 's : a symbol representing a second  
; 'ms : a symbol representing a millisecond (1/1000 seconds)
```

```
; snooze pauses the execution of the program for the  
; amount of time given.
```

```
;   
;   
; int time-period -> void  
(define (snooze n tp)  
  (let [(in-seconds (match tp  
                        ('h (* n 60 60))  
                        ('m (* n 60))  
                        ('s n)  
                        ('ms (/ n 1000)))]  
    (sleep in-seconds)))
```

```
(snooze 60 'ms) ; now we know how long we'll snooze for
```

# Symbols vs identifiers

A **symbol**, which is a particular kind of **value**. It has meaning within the context of a given program in the language.

An identifier is a **syntactic category** within the context of the language itself. We have seen examples of identifiers:

- language reserved words (keywords): `lambda`, `define`, `let`, ...
- bound variables: `(let ((x 31337)) (+ x 42))`

# Symbols vs identifiers

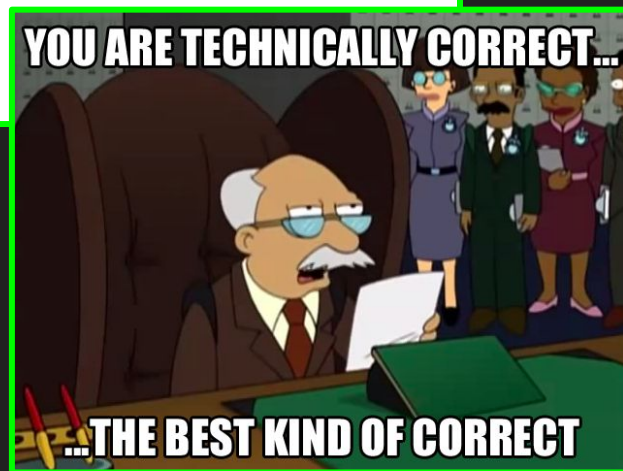
A **symbol**, which is a particular kind of **value**. It has meaning within the context of a given program in the language.

An identifier is a **syntactic category** within the context of the language itself

```
(snooze 60 'ms) ; second argument is a symbol, versus...  
(snooze 60 ms) ; ...a nicer syntax like this??
```

# Time unit identifiers, round one

```
; a time-period is a unit of time, and is one of:  
; 'h : an identifier representing an hour (60 * 60 seconds)  
; 'm : an identifier representing a minute (60 seconds)  
; 's : an identifier representing a second  
; 'ms : an identifier representing a millisecond (1/1000 seconds)  
(define h 'h)  
(define m 'm)  
(define s 's)  
(define ms 'ms)
```



# Time unit identifiers, round one

```
; a time-period is a unit of time, and is one of:  
the time-identifiers are: (60 * 60 seconds)  
the time-identifiers are: (60 seconds)  
the time-identifiers are: (1/1000 seconds)  
(let [(d 31)  
      (m "december")  
      (y 1974)]  
  (... (snooze 60 m))) ; oh no!
```



# Syntax transformers

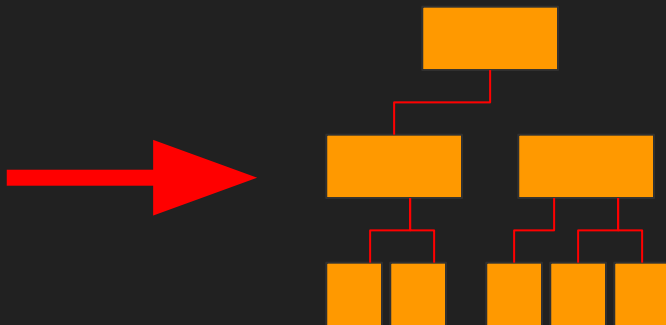
**Definition:** A **syntax transformer** is a special kind of function *that runs at compile time*, consuming certain code as input and rewriting it as new code.

- Syntax transformers are also known informally as **macros**.
- Syntax transformers are also known as **syntactic forms** or just **forms**

When we talk about "special forms" with special evaluation rules, then, we know that this happens because the **syntax is being transformed to accommodate the special evaluation rules**.

```
; my great program  
  
(define x ...)  
(define y ...)  
  
(define-macro foo ...)  
(define (bar z) ...)
```

Source code



AST

**>> (main)**  
**42**

Output

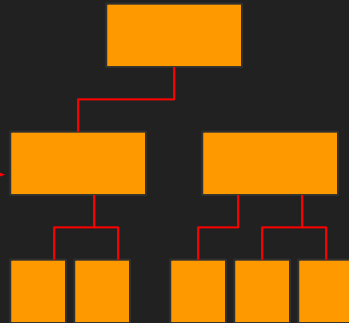
```
; my great program  
  
(define x ...)  
(define y ...)  
  
(define-macro foo ...)  
(define (bar z) ...)
```

Source code



```
; my great program  
  
(define x ...)  
(define y ...)  
  
(define (foo a b) ...)  
(define (bar z) ...)
```

Macro expanded  
source



AST



>> (main)  
42

Output



# Metaprogramming!!!

Creating a **custom syntax** within the context of a snooze call is our first step into **metaprogramming a domain-specific language**.

Our solution to this problem will be to write a **syntax transformer** that transforms the syntax we would like snooze to have, to syntax that Racket already understands

```
(snooze 60 m)    ; to Racket, this will appear as (snooze 60 'm)
(snooze 200 ms) ; to Racket, this will appear as (snooze 200 'ms)
(snooze 60 ns)  ; should emit a syntax error before Racket tries to eval it
```

# Syntax transformers

A macro is like a **function that consumes syntax and produces syntax**.

Unlike ordinary functions, which run when they're called during program execution, macros "expand" **at compilation time**, before the final program is compiled and executed.

```
(snooze 60 m)    ; to Racket, this will appear as (snooze 60 'm)
(snooze 200 ms) ; to Racket, this will appear as (snooze 200 'ms)
(snooze 60 ns)  ; should emit a syntax error before Racket tries to eval it
```

# define-syntax-rule

```
; (define-syntax-rule  
  (rule-name <pattern>)  
  <template>)
```

```
(define-syntax-rule  
  (here-is-a-new-function (name arg) body)  
  (define (name arg) body))  
  
(here-is-a-new-function (add1 n) (+ 1 n))
```

# define-syntax-rule

```
;(define-syntax-rule  
  (rule-name <pattern>  
    <template>))
```

```
(define-syntax-rule  
  (here-is-a-new-function name arg body)  
  (define (name arg) body))  
  
(here-is-a-new-function add1 n (+ 1 n))
```

# define-syntax-rule vs define-syntax

The define-syntax-rule form binds a macro that matches a single pattern (in this case (snooze n tp))

There is a more general form.

Similar distinction to (define (f x y) ...) vs (define f (lambda (x y) ...))

```
(define-syntax-rule (snooze n tp)
  (let* [(in-seconds (match 'tp
                        ('h (* n 60 60))
                        ('m (* n 60))
                        ('s n)
                        ('ms (/ n 1000))))])
    in-seconds))
```

```
(define-syntax snooze
  (syntax-rules ()
    ((snooze n tp)
     (let* [(in-seconds (match 'tp
                              ('h (* n 60 60))
                              ('m (* n 60))
                              ('s n)
                              ('ms (/ n 1000))))])
       in-seconds))))
```

# define-syntax-rule vs define-syntax

```
(define-syntax the-function-named
  (syntax-rules (of is)
    [(the-function-named name of arg is body) (define (name arg) body)]))

(the-function-named add1 of n is (+ 1 n))
```

# define-syntax-rule

```
(define-syntax-rule (snooze n tp)
  (let* [(in-seconds (match 'tp
                        ('h (* n 60 60))
                        ('m (* n 60))
                        ('s n)
                        ('ms (/ n 1000))))])
    in-seconds))

(let [(m "hello")]
  (snooze 60 m))
```

# A macro for list comprehension

```
→ ~ python3
Python 3.7.6 (default, Dec 30 2019, 19:38:26)
[Clang 11.0.0 (clang-1100.0.33.16)] on darwin
Type "help", "copyright", "credits" or "license" for more information.
>>> [2 * i + 3 for i in range(10)]
[3, 5, 7, 9, 11, 13, 15, 17, 19, 21]
```

```
→ ~ ghci
GHCi, version 8.8.3: https://www.haskell.org/ghc/ :? for help
Loaded package environment from /Users/ntaylor/.ghc/x86_64-darwin-8.8.3/environments/default
Prelude> [2 * i + 3 | i <- [0..9]]
[3,5,7,9,11,13,15,17,19,21]
Prelude> █
```



# A macro for list comprehension

```
→ ~ python3
```

```
Python 3.7.6 (default, Dec 30 2019, 19:38:26)
```

```
[Clang 11.0.0 (clang-1100.0.33.16)] on darwin
```

```
Ty
```

```
>>
```

```
[3
```

```
; Our goal:
```

```
(list-comp (+ (* 2 i) 3) : i <- (range 0 10))
```

```
Loaded package environment from /Users/ncaytor/.ghc/x86_64-darwin-8.8.5/environmentm  
ents/default
```

```
Prelude> [2 * i + 3 | i <- [0..9]]
```

```
[3,5,7,9,11,13,15,17,19,21]
```

```
Prelude> █
```

## A macro that functions like cond

```
(cond [(= n 1) "one"]  
      [(= n 2) "two"]  
      [(= n 3) "many"]  
      [else "lots"])
```

```
(if (= n 1)  
    "one"  
    (if (= n 2)  
        "two"  
        (if (= n 3)  
            "many"  
            "lots"))))
```

## A macro that functions like cond


```
(cond [(= n 1) "one"]  
      rest ...)
```

```
(if (= n 1)  
    "one"  
    rest ...)
```

# A macro that functions like cond

```
(define-syntax my-cond
  (syntax-rules (else)
    [(my-cond [else <val>])
     <val>]
    [(my-cond [<test> <val>] <next> ...)
     (if <test> <val> (my-cond <next> ...))]))

(my-cond [(= n 1) "one"]
         [(= n 2) "two"]
         [(= n 3) "many"]
         [else "lots"])
```



<rest> ... says that we match any number of occurrences of a pattern that looks like <rest> (that is, are just a single identifier)

# A deeper dive into "..."

- On the left hand side of a syntax-rules arm

`<lhs-pattern> ...` says that we match any number of occurrences of a pattern that looks like `<lhs-pattern>`

- On the right hand side of a syntax-rules arm

`<rhs-pattern> ...` says: for every pattern matched with `<lhs-pattern> ...` , substitute `<lhs-pattern>` with `<rhs-pattern>` .

... "..." behaves a bit like `map` !!!

# A deeper dive into "..."

```
#lang racket
```

```
; double-all should be a macro that takes a list, and  
; returns the list of all elements multiplied by two, quoted.  
;
```

```
; you would probably not write this as a macro, but we're  
; doing so to see how ... works.
```

```
(double-all (list 1 (- 3 1) 3))
```

```
; should produce
```

```
'((* 2 1) (* 2 (- 3 1)) (* 2 3))
```

# A deeper dive into "..."

```
#lang racket

(define-syntax double-all
  (syntax-rules ()
    [(double-all (list <val> ...)) (list '(* 2 <val>) ...)]))

; double-all should be a macro that takes a list, and
; returns the list of all elements multiplied by two, quoted.
;
; you would probably not write this as a macro, but we're
; doing so to see how ... works.
(double-all (list 1 (- 3 1) 3))

; should produce
'((* 2 1) (* 2 (- 3 1)) (* 2 3))
```

# A deeper dive into "..."

```
#lang racket

(define-syntax double-all
  (syntax-rules ()
    [(double-all (list <val> ...)) (list '(* 2 <val>) ...)]))

; double-all should be a macro that takes a list, and
; returns the list of all elements multiplied by two, quoted.
;
; you would probably not write this as a macro, but we're
; doing so to see how ... works.
(double-all (list 1 (- 3 1) 3))

; should produce
'((* 2 1) (* 2 (- 3 1)) (* 2 3))

; Does the macro produce the same thing as the following map?
; Why, or why not? [note: (quasiquote (* 2 (unquote x))) = `(* 2 ,x)]
(map (λ (x) (quasiquote (* 2 (unquote x)))) (list 1 (- 3 1) 3))
```



# Macros versus functions

We use functions to manipulate common features from **computation**.

- Functions **evaluate** as the program is evaluated (at run-time)
- Pattern matching deconstructs pieces of data
- Functions operate on arbitrarily-long data by recursing or HOFs like `map`

We use macros to manipulate common features from **syntax**.

- Macros **expand** as the program is parsed (at compile-time)
- Pattern matching deconstructs pieces of code
- Macros operate on arbitrarily-long data by recursing or "..."