

CS 331: Algorithms and Complexity (Spring 2024)
Unique Number: 50930, 50935 50940, 50945

Assignment 5

Due on Tuesday, 5 March, by 11.59pm

Problem 1

(8 pts)

(a) (6 pts)

A naive solution would be to split into 3 cases, one for each of the 3 possible operations

(1) Adding a gap to the first string

(2) Adding a gap to the second string

(3) Including characters in both strings

$$OPT(i, j) = \min(\alpha_{x_i y_i} + OPT(i - 1, j - 1), \delta + OPT(i - 1, j), \delta + OPT(i, j - 1))$$

This yields a time complexity of $O(3^{m+n})$

Example: (CH, EN)

First layer: (C, E), (C, -), (-, E)

Second layer: (CH, EN), (CH, E-), (C-, EN), (CH, EN), (CH, --), (C-, -E), (-C, EN), (-C, E-), (--, EN)

(b) (2 pts)

1.

-	A	L	G	O
T	1	2	3	4
E	2	2	3	4
S	3	3	3	4
T	4	4	4	4

*Alignment*₁: ALGO

*Alignment*₂: TEST

2. The cost is the value in the bottom right corner, 4

Cost is 4 since there are 4 sets of characters that are different

Problem 2

(12 pts)

- (a) Indices (i, j) store the minimum number of cuts needed to make the substring from i to j a set of palindromes.

```
def min_palindrome(s):
    sections = [∞] * len(s)
    for j in [1, len(s)]:
        min_sections = j
        for i in [1, j]:
            if isPali(i, j):
                min_sections = min(min_sections, sections[i - 1] + 1)
        sections[j] = min_sections
```

I'll test by running it on the string 'coffee'

Each iteration runs with a fixed endpoint j and a increasing start point i.

First iteration: (1 → 1, 1)

1	∞	∞	∞	∞	∞
---	---	---	---	---	---

Second iteration: (1 → 2, 2)

1	2	∞	∞	∞	∞
---	---	---	---	---	---

Third iteration: (1 → 3, 3)

1	2	3	∞	∞	∞
---	---	---	---	---	---

Fourth iteration: (1 → 4, 4)

1	2	3	3	∞	∞
---	---	---	---	---	---

Fifth iteration: (1 → 5, 5)

1	2	3	3	4	∞
---	---	---	---	---	---

Sixth iteration: (1 → 6, 6)

1	2	3	3	4	4
---	---	---	---	---	---

(b)

Problem 3

(10 pts)

- (a) Since we can't include the direct children of a manager, then we need to include subtrees 3 layers deep.

This isn't greedy since local best choices don't always lead to the global best choice.

- (b) I used memoization to store the maximum enjoyment for each person.

```
Map(person, enjoyment) = {}
def max_enjoyment(person):
    if person is None:
        return 0
    if person in Map:
        return Map(person)
```

```
ce = 0
for child in person.children:
    ce += max_enjoyment(child)

gce = 0
for child in person.children:
    for grandchild in child.children:
        gce += max_enjoyment(grandchild)
Map(tree, max(ce, gce + person.enjoyment))
return max(ce, gce + person.enjoyment)
```

Time complexity is $O(n)$