

CS 331: Algorithms and Complexity (Spring 2024)

Unique numbers: 50930, 50935 50940, 50945

Classwork 1

(a) Prove that $n^3 - n$ is divisible by 3 for all positive integers.

Sol.

Base case: $n = 1$, $1^3 - 1$ is divisible by 3, 0 divisible by 3.

Induction hypothesis: Assume that the formula holds for $n = k$, i.e., $k^3 - k$ is divisible by 3

Induction Step: show that the formula holds for $n = k + 1$, i.e., $(k + 1)^3 - (k + 1)$ is divisible by 3

$$\begin{aligned}(k + 1)^3 - (k + 1) &= (k^3 + 3k^2 + 3k + 1) - (k + 1) \\ &= (k^3 - k) + 3k^2 + 3k \\ &= (k^3 - k) + 3(k^2 + k)\end{aligned}$$

But we know that $(k^3 - k)$ is divisible by 3 by induction hypothesis, and $3(k^2 + k)$ is divisible by 3

Conclusion: $(k + 1)^3 - (k + 1)$ is divisible by 3

Since if a is divisible by b and c is divisible by b , then $a + c$ is divisible by b

(b) The Fibonacci numbers $F_n \geq 0$ are the numbers of a famous sequence defined by

$$F_0 = 0$$

$$F_1 = 1$$

...

$$F_n = F_{n-1} + F_{n-2}.$$

Prove that for all $n \geq 0$, $F_n < 2^n$.

Sol.

Define $P(n) = F_n < 2^n$.

Base case: We need two base cases since the definition of F_n requires two values.

Base Case 1: $n = 0$. $F_0 = 0$, which is less than $2^0 = 1$.

Base Case 2: $n = 1$. $F_1 = 1$, which is less than $2^1 = 2$.

Induction hypothesis: Assume that the formula holds for $n = k$, i.e., $P(0)$, $P(1)$, ..., $P(k)$ are all true.

Induction Step: show that the formula holds for $n = k + 1$, i.e., $P(k + 1)$, or $F_{k+1} < 2^{k+1}$.
 $F_{k+1} = F_k + F_{k-1} < 2^k + 2^{k-1} < 2^k + 2^k = 2(2^k) = 2^{k+1}$.

By strong induction, for all $n \geq 0$, $F_n < 2^n$.

(c) You were asked to provide an algorithm that finds the maximum value in a given array, A , with n elements. The algorithm is supposed to scan through each location in the array and report the maximum value seen so far. Algorithm ?? is purported to solve this problem. Show that this is indeed the case by providing a mathematical proof of correctness for Algorithm ?. I.e., show that it finds the maximum value in A . Note: the array index starts from 0.

Algorithm 1: Find maximum in array

Initially $index = 0$ and $max = A[index]$;

while $index < n$ **do**

if $A[index] > max$ **then**

$max = A[index]$;

 Increase $index$ by 1;

report max ;

Sol: Base case: If $n = 1$, then the body of the while loop doesn't get executed and we report the maximum value as $A[0]$.

Induction hypothesis: Assume that max stores the maximum value of $A[0..n - 2]$ (which covers $n-1$ elements of the array A).

Induction step: Show that max stores the maximum value of $A[0..n - 1]$ (which covers n

elements of the array A).

We know from the induction hypothesis that at the $(n - 1)$ -th iteration, max stores the maximum value of $A[0..n - 2]$. In the n -th and final iteration, if $A[n - 1] > \text{max}$, then max will store $A[n - 1]$. Otherwise, max will keep the maximum of $A[0..n - 2]$ since $A[n - 1]$ is not bigger. Therefore, upon termination of the final iteration, max will store the maximum of $A[0..n - 1]$.