## **Network Level**

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#### Content

- Network Level
  - IPv4 Problem
    - Context
    - Characteristics
    - Subnets
  - Private Networks
  - ICMP
  - Address Resolution
  - IPv6 overview
    - Details -> Future Course

### Context

#### Initial Situation

Before the Internet, only nodes from the same network could communicate with each other

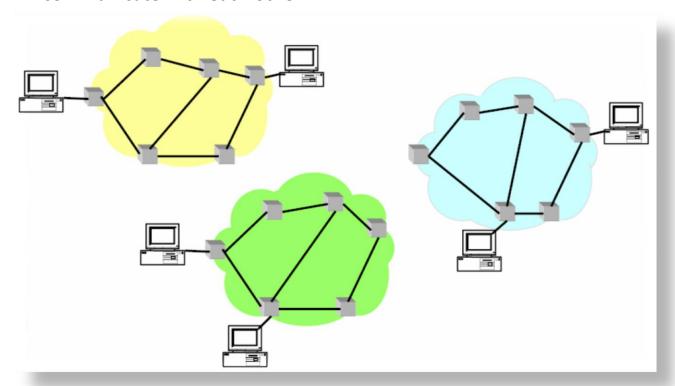


Figure: Individual Network

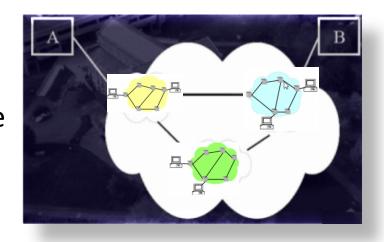
## Considerations

#### Problem

— How to carry packages in a heterogeneous environment?

#### Heterogeneity

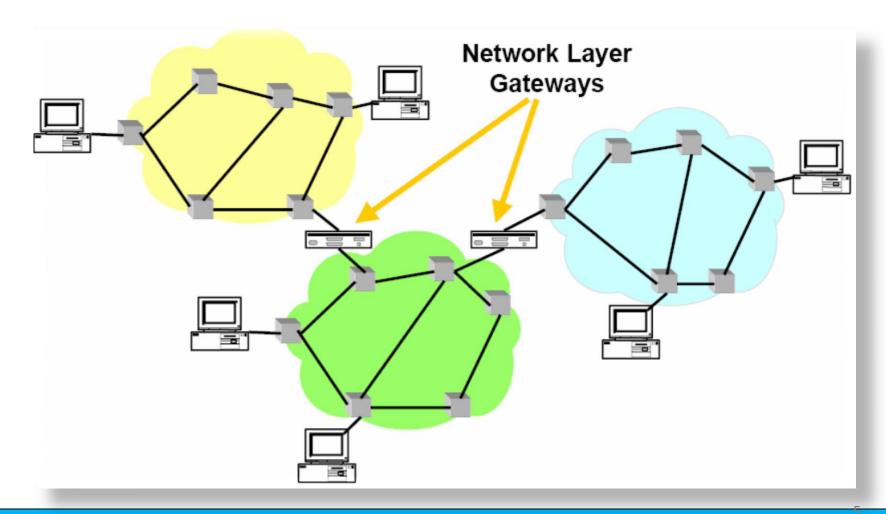
- At lower levels: how to make the interconnection of a large number of independent networks?
- At higher levels: how to provide support for a wide variety of applications?



• **Scaling:** how could we handle a large number of nodes and applications in such a system of interconnected networks?

## Solution

IP – Internet Protocol



### **Network Level**

• IP protocol is used for autonomous systems (AS - Autonomous Systems) in order to interconnect

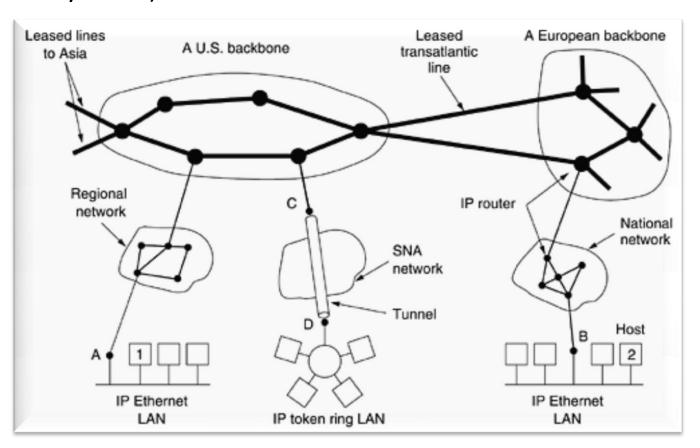


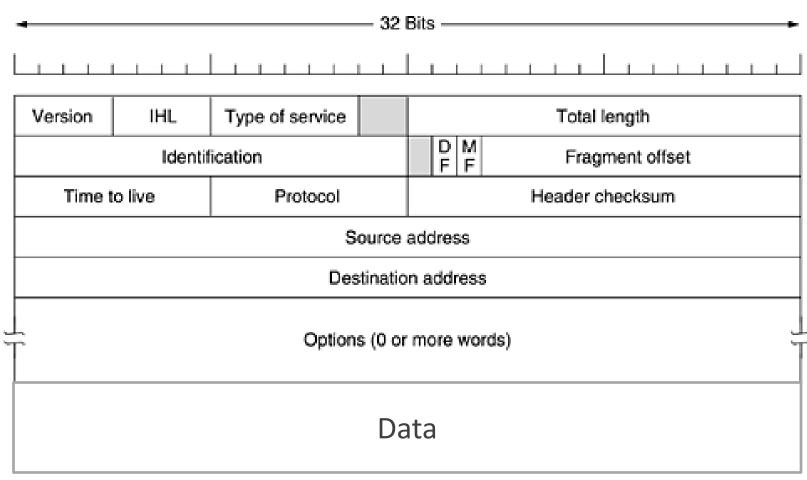
Figure: Internet - collection of interconnected networks

[Computer Networks, 2003 Andrew S. Tanenbaum]

### **Network Level**

- Role: offers connectionless services to transport datagrams from source to destination; source and destination can be in different networks
- Each datagram is independent from the others
- This level does not guarantee the right transmission (loss, multiplier,...)
- +...Future Course

#### IPv4 Datagram



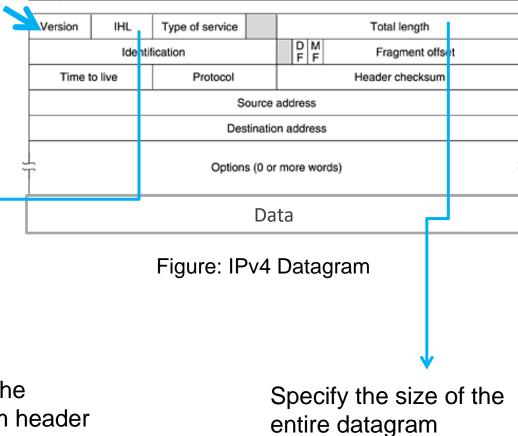
[Computer Networks, 2003 Andrew S. Tanenbaum]

#### **IPv4** Datagram

 Common values for Version field are:

> • 4 – IP Protocol (RFC 791)

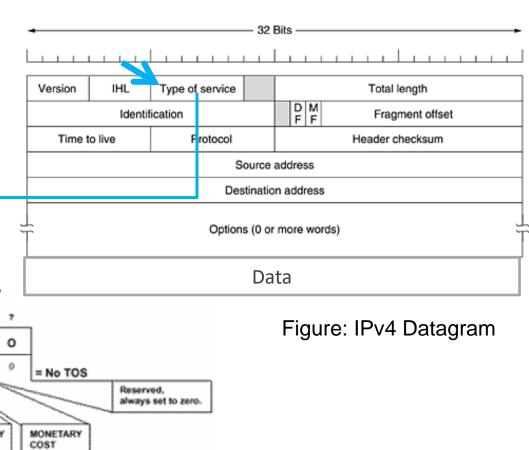
(6 for IPv6 protocol (RFC 1883))

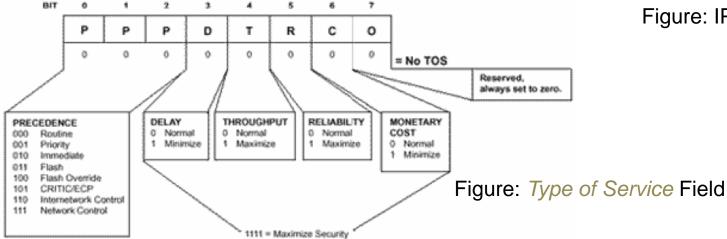


entire datagram

#### IPv4 Datagram

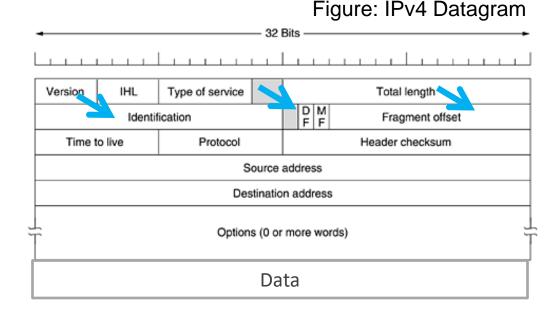
Type of service the
field allows the host
to communicate to
the subnet (e.g.
routers) what type of
service is desired





#### IPv4 Datagram

 Identification field allows the host to identify if the received segment is part of one datagram



#### Flags:

- DF (Don't Fragment) rooters can't fragment the indicated datagram
- MF (More Fragments) signals that the packet is a fragment, followed by others; last fragment has this bit 0
- Fragment offset field represents the fragment placed in a datagram

#### IP Datagram

- Datagram's Fragmentation:
  - Each fragment has the same structure as the IP datagram
  - Reassembly of datagrams is performed by the receiver
  - If a fragment of a datagram is lost, the datagram is destroyed (an ICMP - Internet Control Message Protocol message is sent to the sender)
  - Fragmentation mechanism has been used for some attacks

     firewall piercing (a "special" fragment is considered as
     part of a connection already established, so that it can pass
     through a firewall)

#### IP Datagram

- Datagram Filtering:
  - It is accomplished by a *firewall*: it allows access from the outside in the internal network, according to some policy, certain types of packets ((used by certain protocols / services)
    - Forestall a series of attacks regarding security
    - The firewall can be software or hardware
    - The firewall can function as a proxy or a gateway

#### Proxy- role and architecture:

- Indirect access to other networks (Internet) to hosts on the local network (via proxy)
- The proxy allows an Internet connection sharing
- The proxy can be software or hardware
- May play roles such us: firewall or cache server
- Used to improve the performance (e.g., caching, flow control), filtering messages, ensuring anonymity

#### IPv4 Datagram

 TTL (Time to Live) field specifies the lifetime of the package;

the number is

decremented by every
router through which
the packet passes

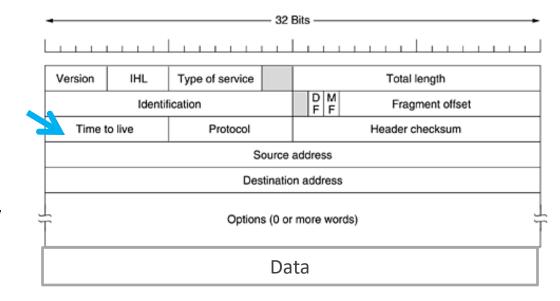
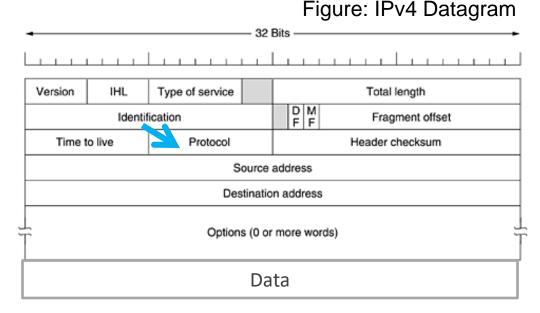


Figure: IPv4 Datagram

#### IPv4 Datagram

 Protocol field specifies the protocol (from the superior level) intended to process the datagram



- 1 ICMP (Internet Control Message Protocol)
- 2 IGMP (Internet Group Message Protocol)
- 6 TCP (Transmission Control Protocol)
- 17 UDP (User Datagram Protocol)
- ... etc.(RFC 1700)

#### IPv4 Datagram

 Header checksum field used for detection; if an error occurs the datagram is destroyed

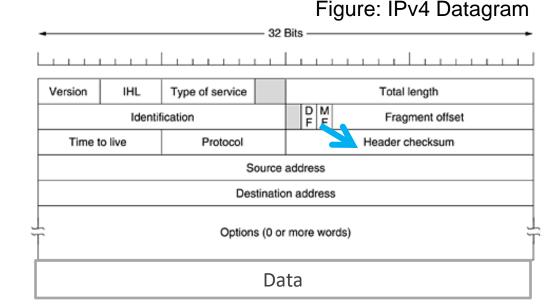
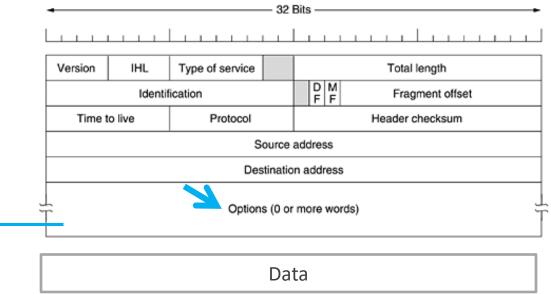


Figure: IPv4 Datagram

- IPv4 Datagram
  - Options Field



Options	Details
Security	Mention if the datagram is a "secret" one
Strict source routing	Show full path to go
Loose source routing	Indicates a list of routers that should not be skipped
Record route	Each rooter adds its own IP
Timestamp	Each rooter adds its own IP and a timestamp

#### IPv4 Datagram

 Source address and Destination address fields indicate the source address and the destination address

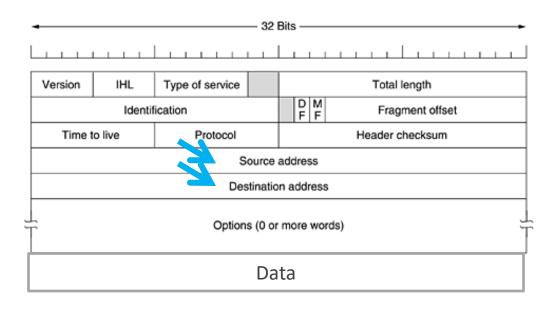
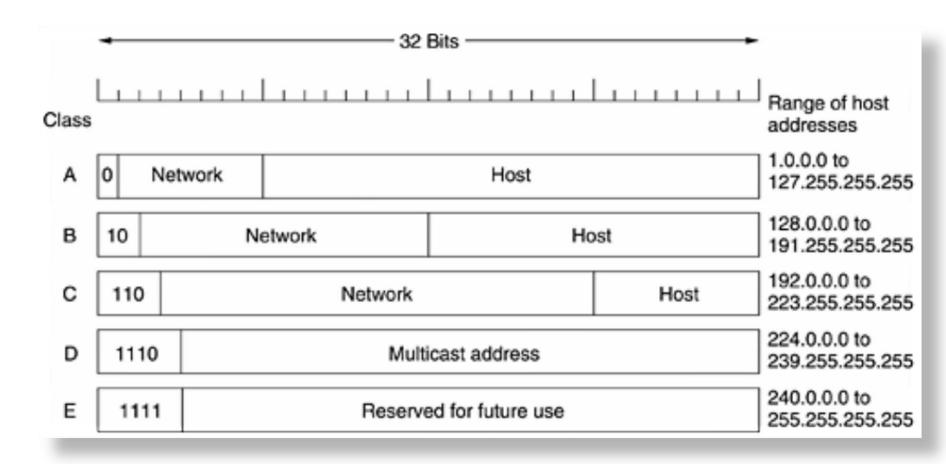


Figure: IPv4 Datagram

- Each IP address includes a network identifier(NetID) and a host identifier (HostID)
- Each network interface has a single IPv4 address
- An IPv4 address has a length of 32 bits
- Initially (RFC 791) there was a division into network classes: A,B,C,D,E

#### IPv4 Addresses



[Computer Networks, 2003 Andrew S. Tanenbaum]

- Class A: 128 possible networks, 2<sup>24</sup> hosts/network
- Class B: 2<sup>14</sup> possible networks, 2<sup>16</sup> hosts/network
- Class C: over 2 million networks, 255 hosts/network
- Network Identifier(NetID) is assigned by a central authority (NIC – Network Information Center)
- Host Identifier(HostID) is assigned locally by a network administrator
- Example: 85.122.23.145 Class A (in decimal notation convention)
  - 0101 0101 0111 1010 0001 0111 1001 0001
- For IPv6, hexadecimal representation is recommended

- An interface network has assigned a unique IP address
- A host can have multiple NICs, therefore it has multiple IP addresses
- The hosts of the same network have the same network identifier (the same NetID)
- Broadcast addresses have HostID's bites equaled to 1
- The IP address in which all HostID's bites are 0 is called a network address – refers to the hole network
  - Example: 85.122.23.0 (network address for a host such us 85.122.23.145 and 85.122.23.1)
- 127.0.0.1 loopback address (localhost)

- From the address space, some addresses are reserved: (RFC 1918):
  - 0.0.0.0 0.255.255.255
  - 10.0.0.0 10.255.255.255 (private addresses)
  - 127.0.0.0 127.255.255.255 (*loopback* addresses)
  - 172.16.0.0 172.31.255.255 (private addresses)
  - 192.168.0.0 192.168.255.255 (private addresses)
- Private addresses: addresses that are not accessible to the outside (the "real"Internet), but only in the organization's intranet

## **Private Networks**

#### Aspects:

- The exponential growth of the hosts number
- Not all hosts offer resources available on the Internet Solution: NAT (Network Address Translation) – RFC 3022, 4008
- The private addresses can be reused (RFC 1918)
- It is based on replacing the private IP address with a public IP address (IP masquerading)

## **Private Networks**

#### Functionality:

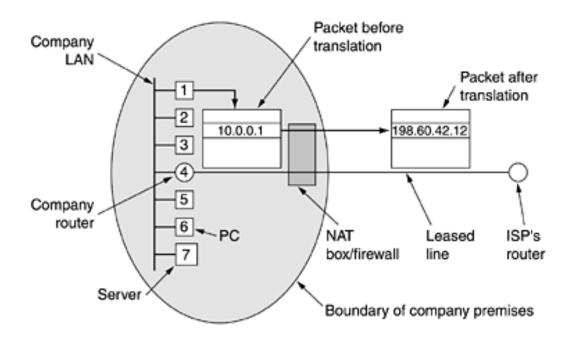


Figure: NAT mechanism

- Routers normally ignore datagrams containing private addresses => private IP addresses can be used in the organization's intranet
- Access to the outside (the "real" Internet ) is achieved via a gate (mediating gateway) that rewrites the source IP addresses / destination

[Computer Networks, 2003 Andrew S. Tanenbaum]

#### Subnets using network masks

- It appeared as a solution to the problem of IP address space's exhaustion
- Simplify Routing
- Subnets cannot be detected externally

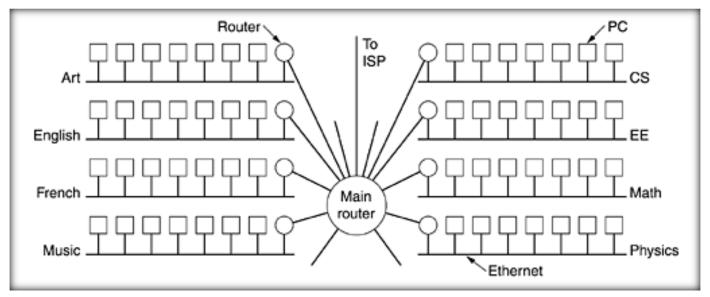
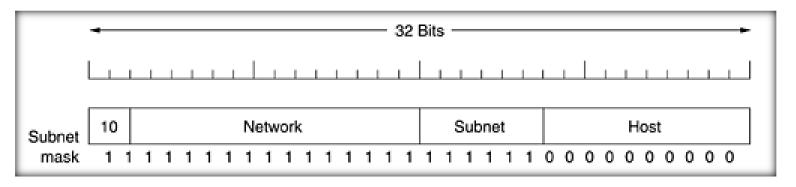


Figure: A campus network

#### Subnets using network masks

- Division into subnetworks will be made via the network mask(netmask): NetID bits are 1, HostID bits are 0
- Subnet identifier (SubnetID) is generally used to group computers based on physical topology



Example. One way to create a subnet in a B network

- Subnets using network masks
  - Example:
    - Let's consider the IP address: 160.0.6.7
    - 10100000 00000000 00000110 00000111
    - Network mask: 255.255.252.0
    - 11111111 11111111 111111100 00000000



Network address: 160.0.4.0

10100000 00000000 00000100 00000000

Network address = network mask AND IP address

- Default subnet masks:
  - 255.0.0.0 Class A
  - 255.255.0.0 Class B
  - 255.255.255.0 Class C

- Convention mark: x.x.x.x/m means that we apply an m bits mask to the IP specified x.x.x.x address
- Example:
  - 10.0.0.0/12 it applies a 12-bits mask to 10.0.0.0 address, we select possible values for the last 20 bits (=32-12)
  - 85.122.16.0/20 it applies a 20-bit mask to 85.122.16.0 address

### **Network Level**

- Protocols
  - ICMP (RFC 792)
  - ARP (RFC 826)
  - RARP (RFC 903)
  - BOOTP (RFC 951,1048,1084)
  - DHCP
- From IPv4 to IPv6

#### ICMP – Internet Control Message Protocol

- Used to exchange control messages
- Use IP
- ICMP messages are processed by the IP software, not by the user processes
- Messages types

Message Type	Description
8 Echo Request	Ask if a host is active
0 Echo Replay	"Yes, I'm active"
3 Destination Unreachable	The package can't be delivered (e.g. DF is set)
5 Redirect	The message is not correctly routed
11 Time Exceeded	Time elapsed (TTL=0) <- (e.g. loop, congestions, low values for time)
etc (RFC 792)	http://www.iana.org/assignments/icmp-parameters

## Protocolul ICMP

#### Used by:

- ping command (Packet Internet Gropher)
- traceroute command
  - A package with TTL=1 (1 hop) is sent
  - The first router ignores the packet and sends back an ICMP message "time-to-live exceeded"
  - A package with TTL=2 is sent (2 hops)
  - The second router ignores the packet and sends back an ICMP message "time-to-live exceeded"
  - Repeat until it has received a response from the destination or has reached the maximum number of hops

## Address resolution

#### IP addresses <-> hardware addresses (physical)

- The process of finding the hardware address of a host, knowing its IP address is called address resolution(address resolution) – ARP protocol (RFC 826)
  - ARP –broadcast protocol (each host receives a request for a physical address, and the answer is given by the one in question)
- The process of finding the IP address based on the hardware address is called reverse address resolution –RARP Protocol (RFC 903)
  - Used to boot workstations without disks
  - BOOTP (RFC 951,1048,1084)
  - DHCP (Dynamic Host Configuration Protocol) RFC 2131,2132

- Context:
  - Issues in IPv4 addresses world:
    - The exponential growth of the hosts` number
    - Very large routing tables
    - Complex configurations, more and more users (and increasing)
    - Failure to ensure QoS
  - Pressure from mobile operators

- Objectives for a new protocol:
  - Support for billions of hosts
  - Reducing routing tables
  - Simplifying Protocol
  - Support for mobile hosts
  - Compatibility with the old IP
  - Support for future developments of the Internet
  - RFC 2460, 2553



• 6 June 2012

#### Aspects:

- IPv6 addresses are 16 bytes in length 2<sup>128</sup> addresses
- Note: 16 hexadecimal numbers, 2 digits each, separated by ":"
  - Example: 2001:0db8:0000:0000:0000:0000:1428:57ab
  - If one or more groups of 4 digits is 0000, the zeros may be omitted and replaced (once) with "::"
  - Example: 2001:0db8::1428:57ab
- To maintain compatibility, public IP addresses are considered a subset of IPv6 address space
- IPv4 addresses in IPv6 can be written as: 10.0.0.1 -> ::10.0.0.1or 0:0:0:0:0:0:0:1

#### ICMPv6

- ICMP provides functions (reporting data transmission, errors, etc.) plus:
  - Neighbor Discovery (Neighbor Discovery Protocol NDP) Replaces the ARP
  - Multicast listener discovery(Multicast Listener Discovery) – replaces IGMP (Internet Group Management Protocol)
- Details in RFC 4443

• ... More -> Future Course

# Summary

- Network Level
  - IPv4 Problem
    - Context
    - Characteristics
    - Subnets
  - Private Networks
  - ICMP
  - Address Resolution
  - IPv6 overview
    - Details -> Future Course



# Questions?