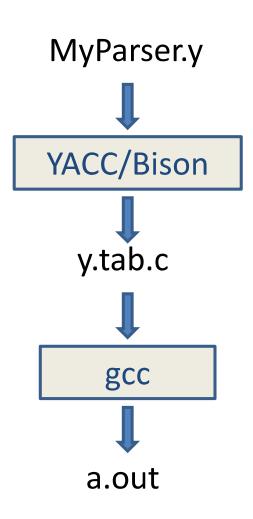
CSE 310

YACC (or Bison)

YACC, Bison

- Yet Another Compilers Compiler
- Unix utility that parses a token stream produced by lex according to specified LALR(1) context free grammar
- We will use Bison which is YACC compatible

How to use YACC



yacc file containing grammar rules

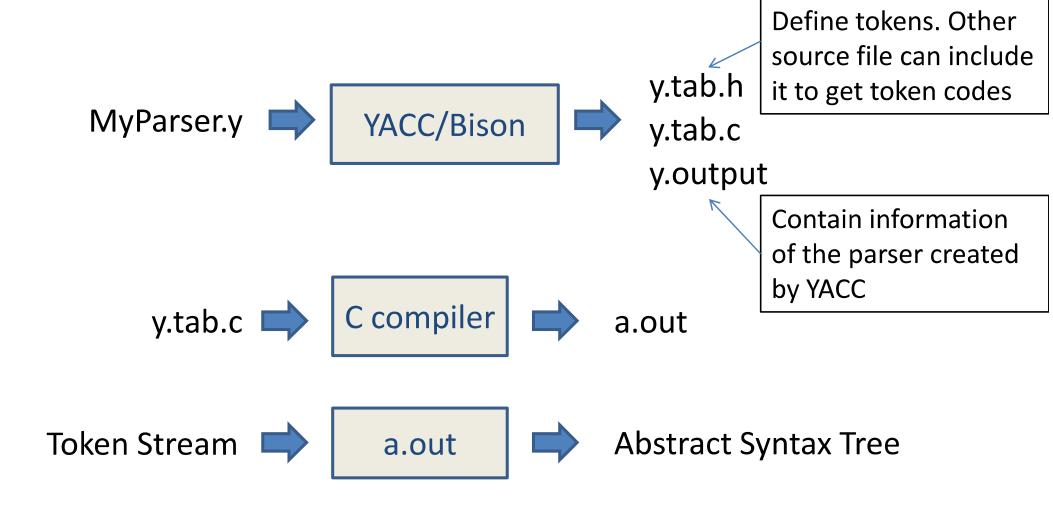
Yacc or Bison program

Parser source code created by yacc Contains a yyparse () function which is invoked to do the parsing task.

C compiler

Executable program to parse according to grammar given in Myparser.y

How to use YACC



Basics: How Bison Works?

- Perform Shift/reduce parsing
 - So that's bottom-up parsing?
- Maintains set of states, reflecting one or more partially parsed rules
- After reading a token it may take two possible actions
 - Shift: If the token cannot complete any rule, shift the token in internal stack
 - Reduce: If a rule can be completed, then pop all R.H.S. symbol from the stack and push L.H.S. symbol

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
stack:
<empty>
```

a = 7; b = 3 + a + 2

input:

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

SHIFT!

stack:

NAME

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
stack:
NAME '='
```

```
input: 7; b = 3 + a + 2
```

```
stack:
NAME '=' 7
```

```
input:
; b = 3 + a + 2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

REDUCE!

```
stack:
```

NAME '=' exp

```
input:
; b = 3 + a + 2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

REDUCE!

stack:

stmt

```
input:
; b = 3 + a + 2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
stack:
stmt ';'
```

```
input: b = 3 + a + 2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
stack:
stmt ';' NAME
```

```
input: = 3 + a + 2
```

```
stack:
stmt ';' NAME '='
```

```
input: 3 + a + 2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
stack:
stmt ';' NAME '=' NUMBER
```

```
input: + a + 2
```

REDUCE!

```
stack:
stmt ';' NAME '=' exp
```

```
input: + a + 2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
stack:
stmt ';' NAME '=' exp '+'
```

```
input:
a + 2
```

```
stack:
stmt ';' NAME '=' exp '+'
NAME
```

```
input:
+ 2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

REDUCE!

```
stack:
stmt ';' NAME '=' exp '+'
exp
```

```
input:
+ 2
```

REDUCE!

```
stack:
stmt ';' NAME '=' exp
```

```
input:
+ 2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
stack:
stmt ';' NAME '=' exp '+'
```

```
input:
2
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
stack:
stmt ';' NAME '=' exp '+'
NUMBER
```

```
input:
<empty>
```

REDUCE!

```
stack:
stmt ';' NAME '=' exp '+'
exp
```

```
input:
<empty>
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

REDUCE!

```
stack:
stmt ';' NAME '=' exp
```

```
input:
<empty>
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
REDUCE!
```

```
stack:
stmt ';' stmt
```

```
input:
<empty>
```

```
stmt: stmt ';' stmt
| NAME '=' exp

exp: exp '+' exp
| exp '-' exp
| NUMBER
```

```
REDUCE!
```

stack:

stmt

```
input:
<empty>
```

```
DONE!
```

stack:

stmt

```
input:
<empty>
```

How parser get Tokens?

 We constructed lexical analyzer that generate tokens!

 So somehow our parser should communicate with scanner!

Scanner Parser Interaction

• Parser assumes the existance of yylex() function

This is called by yyparse()

But how they agree on Token names?

Scanner Parser Interaction

- Parser assumes the existance of yylex () function
- This is called by yyparse()
- But how they agree on Token names?
 - Define token names in YACC program.
 - If you compile using "yacc –d" or "bison –d", it will produce a y.tab.h file
 - This contains an enumeration of token definition
 - Include this y.tab.h file in your lex source file

Scanner Parser Interaction

- With each token scanner can send some value associated with it using global variable yylval.
- Default type is int
- You can redefine its type
 - We will see example
- yytext is also available to your parser

YACC file format

- An yacc file has .y extension
- Three Sections.

```
/**** Definition Section ******/
%%
/**** Rules Section ******/
%%
/**** User SubRoutines ******/
```

Looks Familiar??

Definition Section

- Any code within %{ and %} is copied in the output file of yacc program.
- Name of the tokens (%token)
- Associativity and precedence rules (%left, %right, %nonassoc)
- Name of the start symbol (%start)
- Data type of value associated with each token (%union)
- Several other things

Definition Section

```
왕 {
#include <stdio.h>
#include <stdlib.h>
왕}
                         Terminal
%token ID NUM
                          Non
%type variable
                         Terminal
%start expr
```

Rules Section

Grammar rules and corresponding actions in C code.

```
expr : expr '+' term
                         \{ \$\$ = \$1 + \$3; \}
                          \{ $$ = $1; \}
       term
term : term '*' factor { $$ = $1 * $3; }
                          \{ $$ = $1; \}
       factor
factor : '(' expr ')' { $$ = $2; }
         ID
         NUM
```

User Subroutine Section

C codes copied directly in the y.tab.c file

Usually contains user defined main function

Call yyparse () from main function

Token

• In yacc file token defintions:

%token NUM

• In y.tab.h file:

#define NUM 258

- The lex file can return NUM
- Definitions usually starts from 258 in y.tab.h

yylval

- Global variable that can be used to return some values along with token
- Declare it in lex file
- Data type is YYSTYPE which is defined as int by default
- Lex Program:

yylval

 If different token requires different values, you can define a union in yacc file and associate appropriate values for a Token

• Use union

See an example!

Precedence & Associativity

```
%left \+', \-'
%left \*', '/'
%nonassoc UMINUS
응응
expr : expr '+' expr
                                    \{\$\$=\$1+\$2;\}
      | expr '-' expr
                                    \{\$\$=\$1-\$2;\}
        expr '*' expr
                                    {$$=$1*$2;}
      | expr '/' expr
                                    \{\$\$=\$1/\$2;\}
      | '-' expr %prec UMINUS
                                   \{\$\$=-\$2;\}
        MUM
```

Accessing Value Stack

We can use \$ to access value returned by lex

Recursive Grammar

Left recursion

```
list:
   item
   | list ',' item
   ;
```

Right recursion

```
list:
   item
   | item ',' list
;
```

YACC, Bison prefers left recursion. Why?

Conflicts

 Conflicts arise when there is more than one way to proceed with parsing.

- Two types:
 - shift-reduce [default action: shift]
 - reduce-reduce [default: reduce with the first rule listed]

Conflicts

Reduce/Reduce Conflict:

```
A : B | C
B : 'X'
C : 'X'
```

Shift/Reduce Conflict:

```
Stmt: IF '('exp')' stmt | IF '('exp')' stmt else stmt
```

Handling Conflicts

• See the y.output file for conflict details

Think out why conflict occurred

Rewrite grammar accordingly

int yyparse()

• Called once from main()

- Repeatedly calls yylex() until done:
 - On syntax error, calls yyerror (char *s)
 - Returns 0 if all of the input was processed;
 - Returns 1 if aborting due to syntax error.

Error Reporting

Define the yyerror (char *) function

 Setting '%error-verbose' in definition section can produce more information about error

Error Handling

- The "token" 'error' is reserved for error handling:
 - can be used in rules;
 - suggests places where errors might be detected and recovery can occur.

Example:

```
stmt : IF '(' expr ')' stmt | IF '(' error ')' stmt | FOR ... | ...
```

Error Handling

When an error occurs, the parser:

- pops its stack until it enters a state where the token 'error' is legal;
- then behaves as if it saw the token 'error'
 - performs the action encountered;
 - resets the lookahead token to the token that caused the error.

If no 'error' rules specified, processing halts.

YACC Declaration Summary

`%start'

Specify the grammar's start symbol

`%union'

Declare the collection of data types that semantic values may have

`%token'

Declare a terminal symbol (token type name) with no precedence or associativity specified

`%type'

Declare the type of semantic values for a nonterminal symbol

YACC Declaration Summary

`%right'

Declare a terminal symbol (token type name) that is right-associative

'%left'

Declare a terminal symbol (token type name) that is left-associative

'%nonassoc'

Declare a terminal symbol (token type name) that is nonassociative (using it in a way that would be associative is a syntax error, ex: x op. y op. z is syntax error)

Reference

- Flex & Bison by John Levine
- Lecture on YACC by Tanvir Ahmed Khan
- Powerpoint slide from www.csie.ntu.edu.tw/~b92006/YACC-present-2005.ppt
- Powerpoint slide from www.cs.arizona.edu/~debray/Teaching/CSc453/DOCS/yacc%20tutorial.ppt

Thank You!