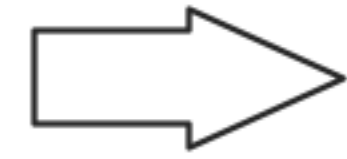


W

Is *QuickSort* always $\mathcal{O}(n \log(n))$? Is *MergeSort* always $\mathcal{O}(n \log(n))$? Justify your answer.
Hint: Consider what this depends on



[0,1,2,3,4,5]
pivot = 0

NO

Mergesort is
always
 $\mathcal{O}(n \log n)$ as it
roughly splits
the list into
half n times

↑ Real

merge sort is always $\mathcal{O}(n \log(n))$ (cuz learned this in lec) and quicksort can be $\mathcal{O}(n^2)$ because the paritaton has the possiblity to always be the smallest in the list

Mergesort is always $n \log n$, quicksort depends on ur pivots (worst case n^2)

Quick sort is on average $\mathcal{O}(n \log(n))$ but not always $\mathcal{O}(n \log(n))$, this depends on how pivot is chosen. Merge sort is always $\mathcal{O}(n \log(n))$ because you always split the list in half.

Quicksort worst case: $\mathcal{O}(n^2)$

quicksort can be unevenly split - worst case of $\mathcal{O}(n^2)$ fr

Dr. Halstead and Dr. Manning are arguing about the new hospital database system sorting upgrades. Dr. Manning argues Quicksort is the faster option, while Dr. Halstead is a firm believer in the efficiency mergesort. With the help of April, Dr. Choi runs some tests with both sorting algorithms on the hospital dataset. Here is what you know about the dataset:

- 1 99% of the time, the hospitals dataset is mostly sorted
- 2 All the sorting algorithms are run in their non-in-place variants
- 3 Every other part of their software is as optimized as it can get - There are no slowdowns other than the choice of algorithm

Note: Since these are doctors and not computer scientists, they do not know asymptotic notation (which would've been the ideal compromise for Will and Natalie)
Given what we know about the hospital database, what did Dr. Choi recommend for the tests?

just use
stalinsort
 $O(1)$ 100%

Merge
sort is
W

let $\epsilon > 0$

bogosort
is the way

Quicksort be the
play because u can
pick a pivot in the
middle and sort the
list in fewer calls
than compared to
mergesort where
you always split list
into chunks sized-2

bubblesort
tho?

quick
sort
 \ggg

quicksort has better
constants than
mergesort, and
since the database
is mostly sorted,
then using the
middle pivot would
almost always be
efficient

quick sort should be
better, because the
mostly sorted
nature of the input
should speed up the
partition operation,
especially if we
select the middle
element as the pivot

Since the list is sorted
roughly, depending on
the choice of pivot, I
think quick sort would
be slower than merge
sort most definitely!

no school on friday!!!!

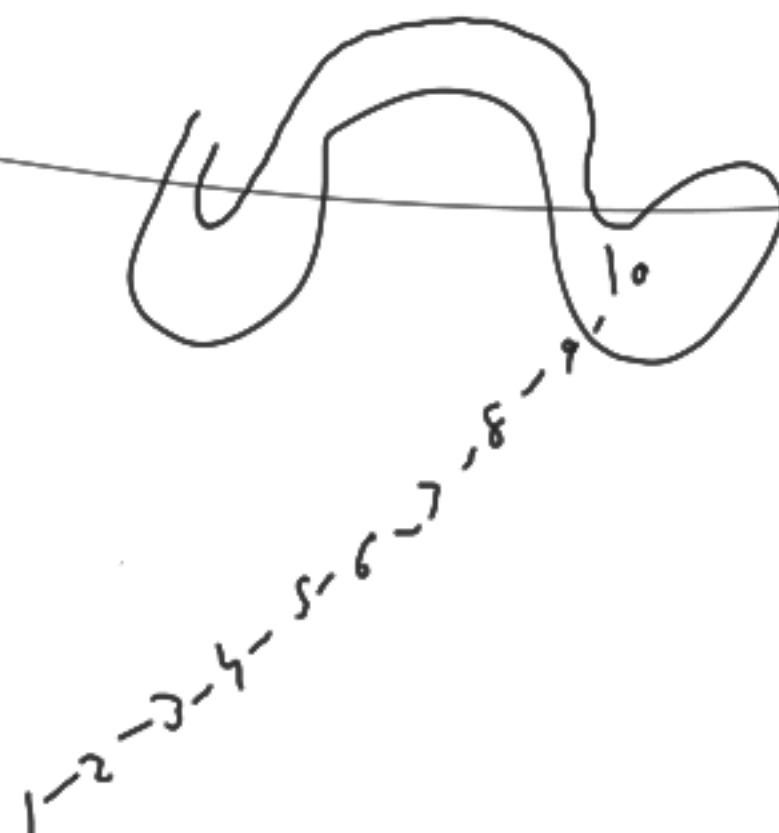
T1:
 $O(n)$

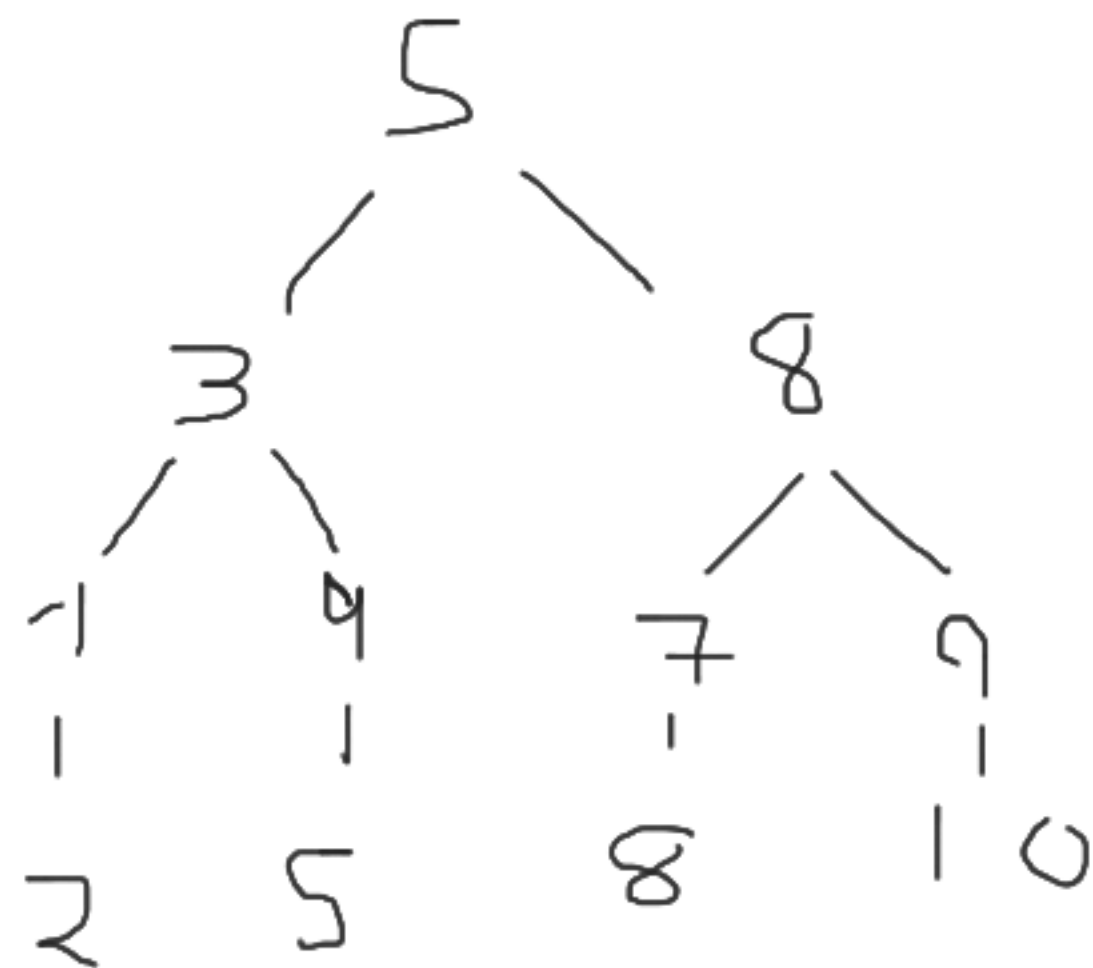
T2:
 $O(n)$

T3:
 $O(\log n)$

1) $O(n)$ 2)
 $O(n)$ 3)
 $O(\log n)$

**Shuffle
is the
faster!**





1, 2, 3, 4, 5, 6, 7, 8, 9, 10

1, 2, 3, 4, 5

6, 7, 8, 9, 10

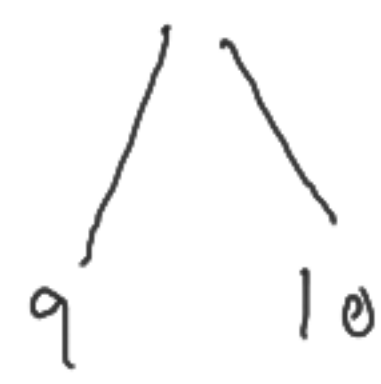
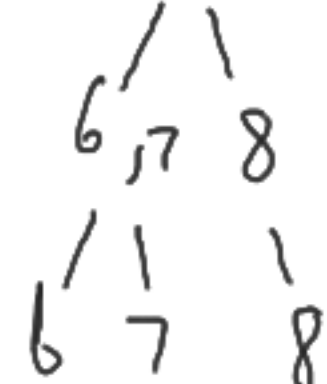
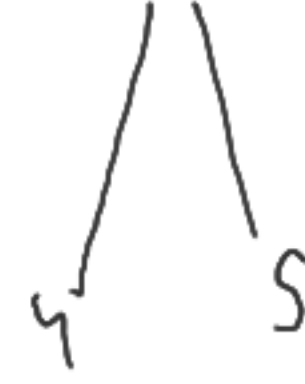
C S

1, 2, 3

4, 5

6, 7, 8

9, 10



1, 2, 3

4, 5

6, 7, 8

9, 10

1, 2, 3, 4, 5

6, 7, 8, 9, 10

1, 2, 3, 4, 5, 6, 7, 8, 9, 10

















