Obfuscation

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Obfuscation Goal



The goal of code obfuscation is to make a program more

- difficult to understand for a human than the original code
- difficult to analyse automatically for a reverse engineering tool in order to hide some asset in the program.

Vague definition of difficult:

 The obfuscated program requires more human time, more money, or more computing power to analyse than the original program.

Example: obfuscated code



```
public class C {
 static Object get0(Object[] I) {
  Integer I7, I6, I4, I3; int t9, t8;
  17=new Integer(9);
  for (;;) {
   if (((Integer)I[0]).intValue()%((Integer)I[1]).intValue()==0)
       {t9=1; t8=0;} else {t9=0; t8=0;}
   14=new Integer(t8);
   16=new Integer(t9);
   if ((I4.intValue()^I6.intValue())!=0)
     return new Integer(((Integer)I[1]).intValue());
   else {
     if ((((I7.intValue() + I7.intValue() *I7.intValue()) %2!=0)?0:1)!=1)
        return new Integer(0);
     I3=new Integer(((Integer)I[0]).intValue()%
           ((Integer) I[1]).intValue());
     I[0]=new Integer(((Integer)I[1]).intValue());
     I[1]=new Integer(I3.intValue());
```

Example: original code



- An obfuscation tool turns the original code into obfuscated code.
- We want that obfuscating transformations make the program as hard to understand as possible.

```
public class C {
   static int gcd(int x, int y) {
     int t;
     while (true) {
       boolean b = x % y == 0;
       if (b) return y;
       t = x % y; x = y; y = t;
     }
}
```

This code computes the Greatest Common Denominator of its arguments...

Code Obfuscations



- Lexical transformations
 - Modify variable names (Identifier-Renaming)
 - Strings scrambling
- Control-flow transformations
 - Opaque Predicates -Redundant Code
 - Increase Indirection Levels
- Data transformations
 - Modify data structures
- Anti-disassembly
- Anti-debugging
- Code/Data Encoding or encryption

Identifier Renaming



Basic Obfuscation based on lexical transformations of all names (classes, methods, variables) into random strings with no semantics

```
public void addUserToList(String strRoomName, String strUser)
    RoomTabItem tab = getRoom(strRoomName);
    if(tab != null)
       tab.addUserToList(strUser);
public void removeUserFromList(String strRoomName, String strUser)
    RoomTabItem tab = getRoom(strRoomName);
    if(tab != null)
       tab.removeUserFromList(strUser);
```

```
public void k(String s, String s1)
    h h1 = h(s);
    if(h1 != null)
        h1.k(s1);
public void 1(String s, String s1)
    h h1 = h(s);
    if(h1 != null)
        h1.1(s1);
```

Lexical Transformations



- The Stunnix obfuscator targets at obfuscating only the layout of the JavaScript code
- As the obfuscator parses the code, it removes spaces, comments and new line feeds
- While doing so, as it encounters user defined names, it replaces them with some random string
- It replaces print strings with their hexadecimal values
- It replaces integer values with complex equations

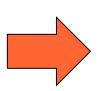
From Stunnix to Javascript



Actual code:

```
function foo( arg1)
{
  var myVar1 = "some string"; //first
  comment
  var intVar = 24 * 3600; //second
  comment
  /* here is a
  multi-line comment */
  document. write( "vars are:" +
    myVar1 + " " + intVar + " " + arg1)
  ;
};
```

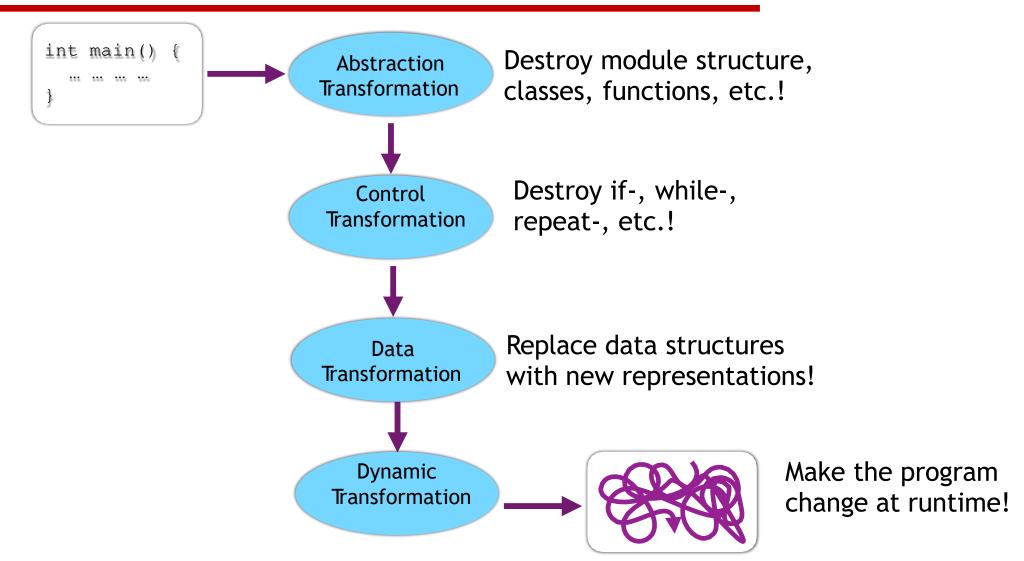
Obfuscated code:



```
function z001c775808(
z3833986e2c) { var
z0d8bd8ba25=
"\x73\x6f\x6d\x65\x20\x73\x74\
x72\x69\x6e\x67"; var
z0ed9bcbcc2= (0x90b+785-
0xc04)* (0x1136+6437-0x1c4b);
document. write(
"\x76\x61\x72\x73\x20\x61\x72\
x65\x3a"+ z0d8bd8ba25+ "\x20"+
z0ed9bcbcc2+ "\x20"+
z3833986e2c);};
```

What is Obfuscation?





Example



```
int main() {
   int y = 6;
   y = foo(y);
   bar (y, 42);
int foo(int x)
   return x*7;
void bar(int x, int z) {
   if (x==z)
      printf("%i\n",x);
```

Abstract Transformation



```
int main() {
   int y = 6;
   y = foo(y);
   bar(y, 42);
int foo(int x)
   return x*7;
void bar(int x, int z) {
   if (x==z)
      printf("%i\n",x);
```

Abstraction Transformation



```
int main() {
   int y = 6;
   y = foobar(y, 99, 1);
   foobar(y, 42, 2);
int foobar(int x, int z, int s)
   if (s==1)
      return x*7; else
   if (s==2)
      if (x==z)
         printf("%i \n",x);
```

Data Transformation



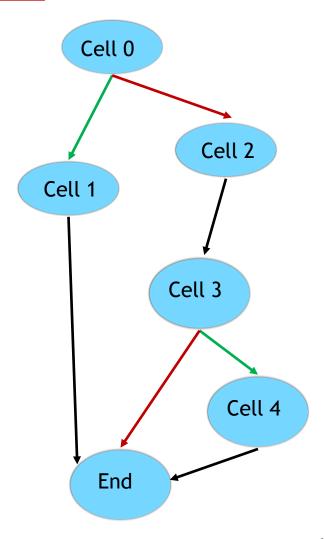
```
Abstraction
                                   Transformation
int main() {
   int y = 6;
   y = foobar(y, 99, 1);
   foobar (y, 42, 2);
                                      Data
                                   Transformation
int foobar(int x, int z, int s) {
   if (s==1)
       return x*7; else if
    (s==2)
       if (x==z)
           printf("%i \n",x);
```

```
int main () {
    int y = 12;
    y = foobar(y, 99, 1);
    foobar (y, 36, 2);
int foobar(int x, int z, int s) {
    if (s==1)
        return (x*7)%51;
     else if (x==z) {
      int x2=x*x%51,x3=x2*x%51;
      int x4=x2*x2%51,x8=x4*x4%51;
      int x11=x8*x3%51;
      printf("%i\n",x11);
```

Control Flow Transformation



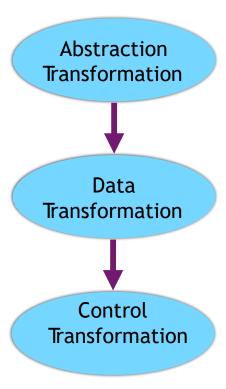
```
int main () {
    int y = 12;
    y = foobar(y, 99, 1);
    foobar (y, 36, 2);
int foobar(int x, int z, int s) {
      if (s==1) /cell 0
             return (x*7) %51; /cell 1
      else /cell 2
      if (x==z) { /cell 3
      int x2=x*x%51,x3=x2*x%51; /cell 4
      int x4=x2*x2%51,x8=x4*x4%51;
      int x11=x8*x3%51;
      printf("%i\n",x11);
  /end
```



Control Flow Transformation



```
int foobar(int x, int z, int s) {
char* next= &&cell0; int retVal = 0;
cell0: {next=(s==1)? &&cell1: &&cell2;
qoto *next; }
cell1: {retVal=(x*7) %51; goto end;}
cell2: \{next=(s==2)? \&\&cell3: \&\&end; \}
qoto *next; }
cell3: {next=(x==z)? &&cell4: &&end;
goto *next; }
cell4: {
      int x2=x*x%51, x3=x2*x%51;
      int x4=x2*x2\$51, x8=x4*x4\$51;
      int x11=x8*x3 % 51;
      printf("%i \n", x11); goto end;
 end: return retVal;
```





Tigress

Tigress Obfuscator



- Open-Source obfuscator for C language from University of Arizona
- It works on linux, Darwin (Mac O.S.), Android
- It works with Intel, ARM, web assembly,
- Can deal with 32 or 64 instruction set,
- Can be compiled with gcc, clang, emcc

References

Tigress - https://tigress.wtf

Clang- https://clang.llvm.org/

EMCC - https://emscripten.org/

Web Assembly - https://webassembly.org/

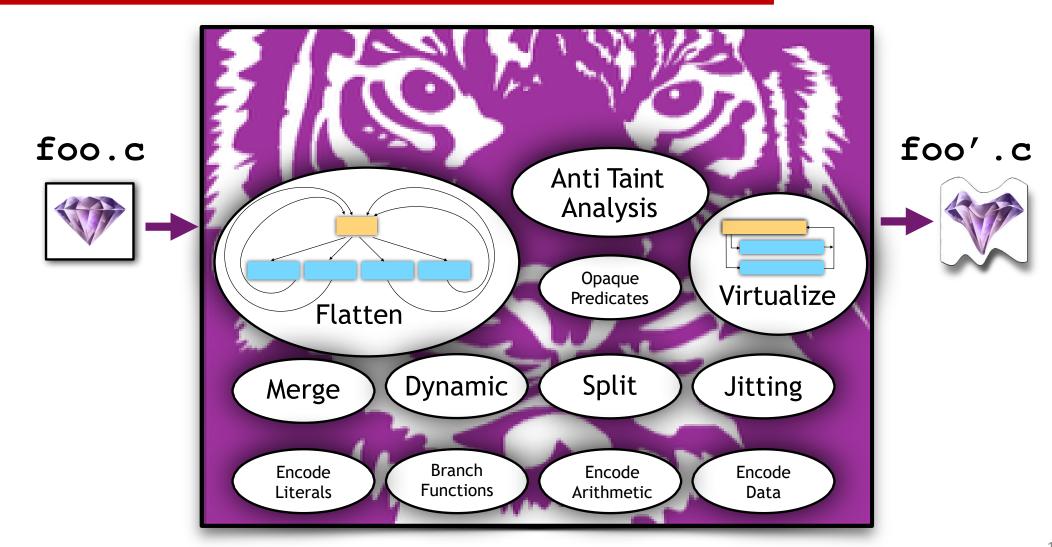






Tigress Obfuscations





Tigress obfuscation script

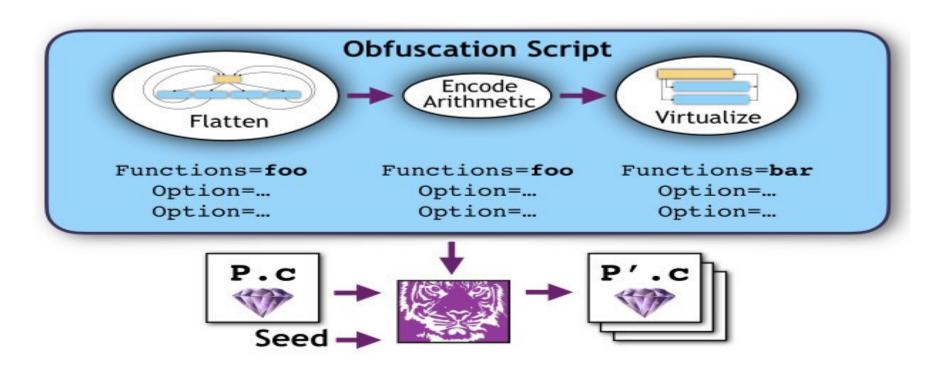


- Tigress is a diversifying virtualizer/obfuscator for the C language that supports many protections against both static and dynamic reverse engineering and de-virtualization attacks.
- Tigress is a source-to-source transformer it takes a C source program as input and returns a new C program as output.

Obfuscation Script



 An obfuscation script (actually, a long sequence of command line options) describes the sequences of transformations that should be applied to the functions of the program:



Example of Tigress transformations



- Transform a function bar in a program foo.c using Tigress.
- The first transformation flattens the code, the second turns the flattened function into an interpreter:

```
tigress --Verbosity=1 --Environment=x86 64:Linux:Gcc:4.6 --Seed=42 \
--Transform=Flatten \
   --Functions=bar \
   --FlattenDispatch=switch \
                                 --FlattenObfuscateNext=false \
   --FlattenRandomizeBlocks=true \
                                       --FlattenConditionalKinds=branch,compute,flag \
   --FlattenImplicitFlowNext=true \
 --Transform=Virtualize \
   --Functions=bar \ --VirtualizeDispatch=direct \
 --Transform=CleanUp \
    --CleanUpKinds=annotations,constants,names \
 --out=foo out.c foo.c
```

Web Assembly



- The typical way to protect code running in the browser from attacks by the user is to write the code in Javascript and transform it with a Javascript obfuscator.
- An alternative approach is to write the code in C, transform it with Tigress, and to compile the resulting code using <u>Emscripten</u> (a C-to-WASM compiler) to a WebAssembly/html/Javascript package.

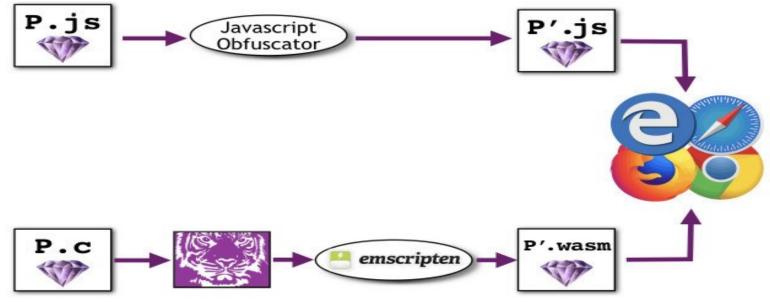
Web Assembly



- The WebAssembly/html/Javascript package can be loaded by most browsers.
- The advantage is that the resulting code can run at near-native speeds

Tigress supports more powerful obfuscating transformations than most

Javascript obfuscators.



Obfuscating Arithmetic

Encoding Integer Arithmetic



•
$$x+y = x - \neg y - 1$$

•
$$x+y = (x \oplus y) + 2 \cdot (x \wedge y)$$

$$\bullet x + y = (x \lor y) + (x \land y)$$

$$\bullet x + y = 2 \cdot (x \lor y) - (x \bigoplus y)$$

Logic Operator	Math symbol	C language operator
OR	V	i
XOR	\oplus	٨
AND	٨	&
NOT	٦	!

Example



One possible encoding of

$$z = x+y+w$$

• is

$$z = (((x ^ y) + ((x & y) << 1)) | w) + (((x ^ y) + ((x & y) << 1)) & w);$$

Many others are possible, which is good for diversity.

Exercise



- The virtualizer's add instruction handler could still be identified by the fact that it uses a + operator!
- Try adding an arithmetic transformer:

```
tigress --Environment=x86_64:Linux:Gcc:4.6\
    --Transform=Virtualize \
         --Functions=fib \
         --VirtualizeDispatch=switch\
         --Transform=EncodeArithmetic \
          --Functions=fib \
          --out=fib5.c fib.c
```

 What differences do you notice between before and after arithmetic encoding?

Virtualized Code



```
int fib(int n ) {
  while (1) {
    switch (*( 1 fib $pc[0])) {
    case PlusA: {
     (1 \text{ fib } \$sp[0] + -1) -> int =
        (1 \text{ fib } \$sp[0] + -1) -> int
        (1 \text{ fib } \$sp[0] + 0) -> int;
   break;
```

Virtualization and Obfuscating Arithmetic University



```
int fib(int n ) {
  while (1) {
    switch (*( 1 fib $pc[0])) {
    case PlusA: {
     (1 \text{ fib } \$sp[0] + -1) -> int =
        ((1 \text{ fib } \$sp[0] + -1) -> int
         (1 \text{ fib } \$sp[0] + 0) -> int)
         (((1 \text{ fib } \$sp[0] + -1) -> int)
         (1 \text{ fib } \$sp[0] + 0) -> int) << 1);
     break;
```

Virtualization and Obfuscating Arithmetic University



```
x+y = (x \oplus y) + 2 \cdot (x \wedge y)
int fib(int
  while (1) {
     switch (*( 1 fib $pc[0])) {
     case PlusA: {
     (1 \text{ fib } \$sp[0] + -1) -> int =
         ((1 \text{ fib } \$sp[0] + -1) -> int
         (1 \text{ fib } \$sp[0] + 0) -> int)
         (((1 \text{ fib } \$sp[0] + -1) -> int)
         (1 \text{ fib } \$sp[0] + 0) -> int) << 1
     break;
```

Opaque Expressions

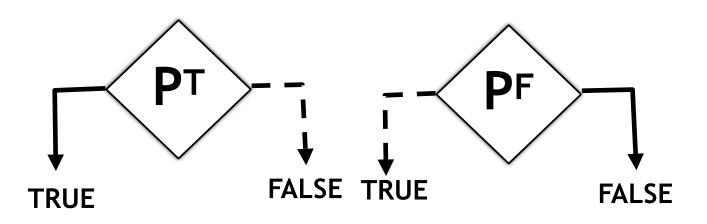
Opaque Expressions



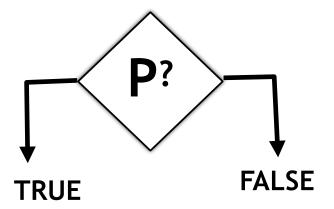
 An expression whose value is known to you as the defender (at obfuscation time) but which is difficult for an attacker to figure out

Opaque Predicates





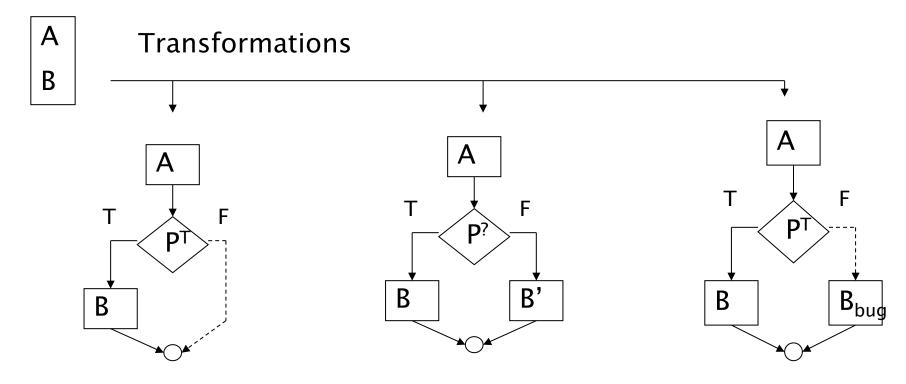
Opaquely true/ false predicate



Opaquely indeterminate predicate

Opaque Predicates Strategies





P^T: Predicate always TRUE

TRUE or FALSE

B and B' have same effect

P?: Predicate can be

P^T: Predicate always TRUE

B_{bug:}A buggy version of B

Inserting Bogus Control Flow





```
if (x[k] == E=1)
  R = (s*y) % n
else
  R = s;
s = R*R % n;
L = R;
```

Inserting Bogus Control Flow



```
if (x[k] == 1)
   R = (s*y) % n
else
   R = s;
s = R*R % n;
L = R;
```



```
if (x[k] == 1)
  R = (s*y) % n
else
else
   s = R*R * n;
L = R;
```

Inserting Bogus Control Flow



```
if (x[k] == 1)
   R = (s*y) % n
else
   R = s;
s = R*R % n;
L = R;
```



```
if (x[k] == 1)
  R = (s*y) % n
else
else
   s = (R%n) * (R%n) %n;
L = R;
```

Opaque Predicates



- Predicate is opaque if its outcome is known at obfuscation time, but difficult to deduce later
- Implementation of opaque predicates uses pointer-based structures
 - Static code analysis is harder
 - But Dynamic Analysis (debugging) can identify the actual running paths and the ones which are never executed!

Exercise

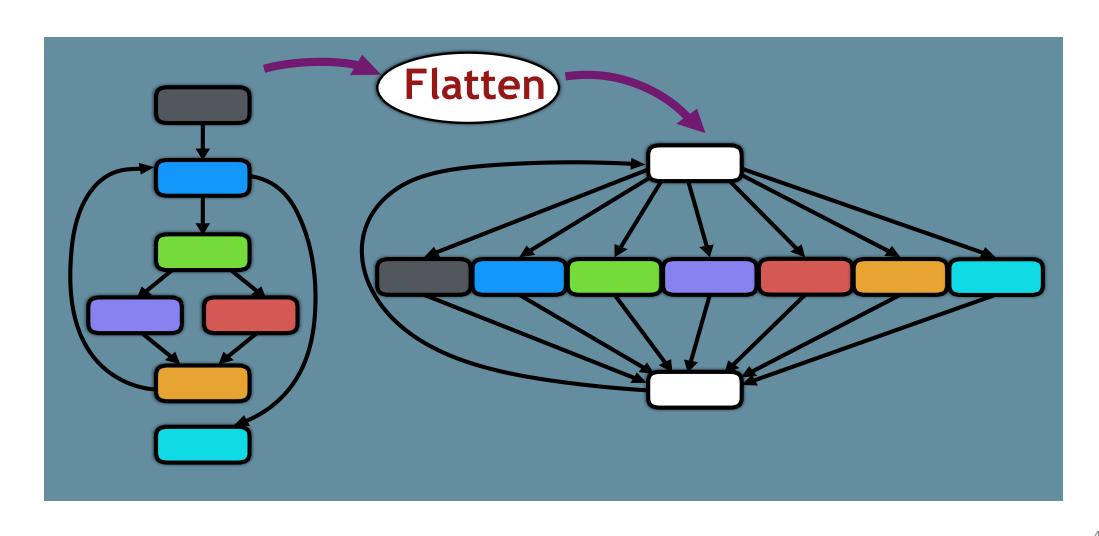


```
*********************
 1) Opaque Predicates
********************
tigress --Environment=x86 64:Linux:Gcc:4.6 --Seed=0 \
  --Transform=InitEntropy \
     --InitEntropyKinds=vars \
  --Transform=InitOpaque \
     --Functions=main\
     --InitOpaqueCount=2\
     --InitOpaqueStructs=list,array \
  --Transform=AddOpaque\
     --Functions=fib\
     --AddOpaqueKinds=question \
     --AddOpaqueCount=10 \
   --out=fib1.c fib.c
```

Control Flow Flattening

Control Flow Flattening

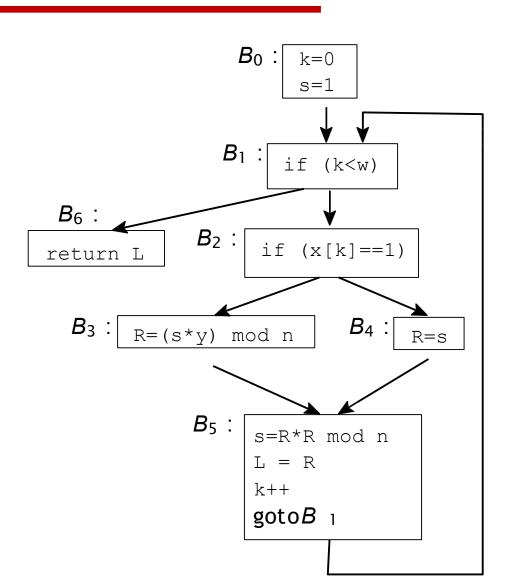




Example



```
int modexp(
   int y, int x[],
   int w, int n) {
   int R, L;
   int k=0; int s=0;
   while (k < w) {
      if (x[k] == 1)
         R = (s*y) % n
      else
         R = s;
      s = R*R % n;
      L = R;
      k++;
   return L;
```



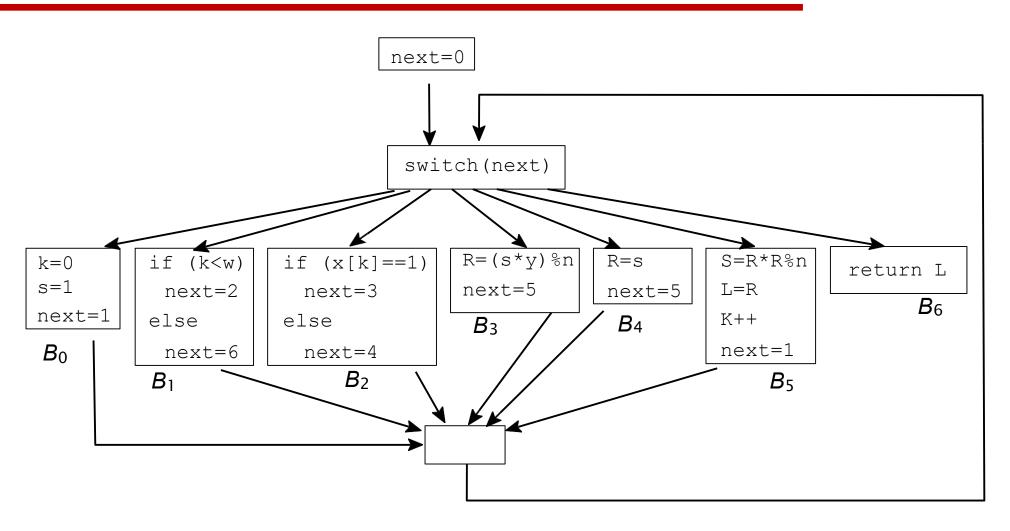
Example flattened



```
int modexp(int y, int x[], int w, int n) {
int R, L, k, s;
   int next=0;
  for(;;)
      switch(next) {
         case 0 : k=0; s=1; next=1; break;
         case 1 : if (k<w) next=2; else next=6; break;
         case 2 : if (x[k]==1) next=3; else next=4; break;
         case 3 : R=(s*y)%n; next=5; break;
         case 4 : R=s; next=5; break;
         case 5 : s=R*R%n; L=R; k++; next=1; break;
         case 6 : return L;
```

Flattened Control Flow Graph

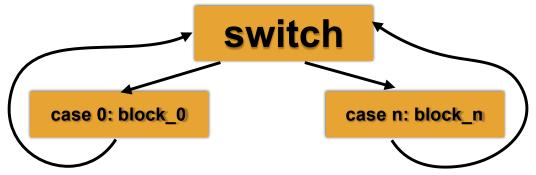




Flattening Algorithm



- 1. Construct the CFG
- 2. Add a new variable int next=0;
- 3. Create a switch inside an infinite loop, where every basic block is a case:

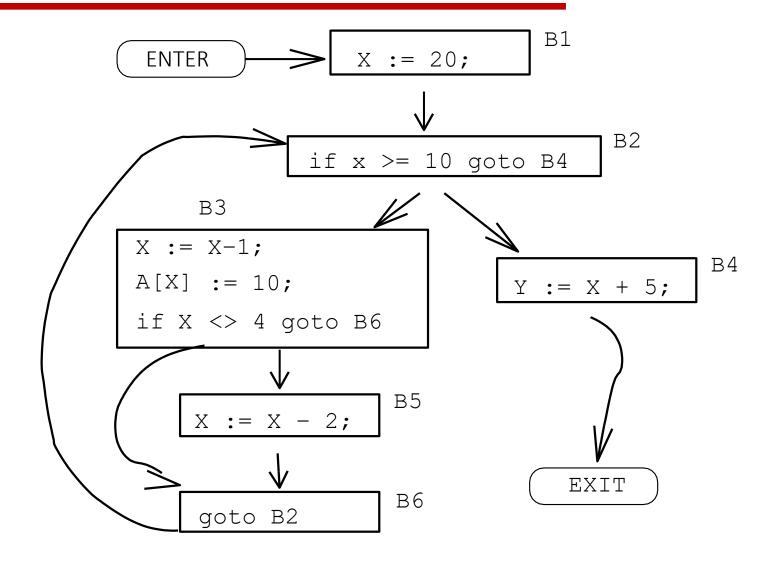


4. Add code to update the **next** variable

```
case n: {
    if (expression)
        next = ...
    else
        next = ...
}
```

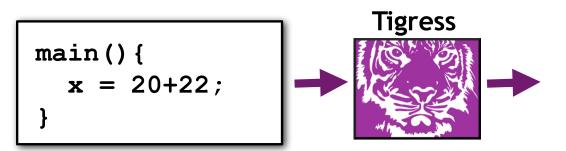
Flatten this CFG! Work with your friends! University



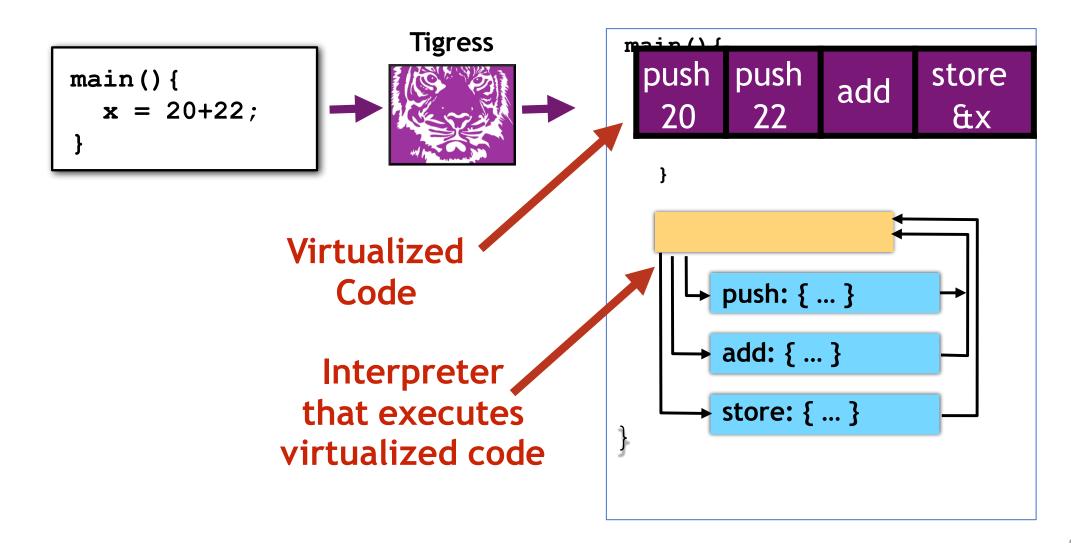


Virtualization Obfuscation





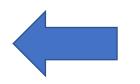






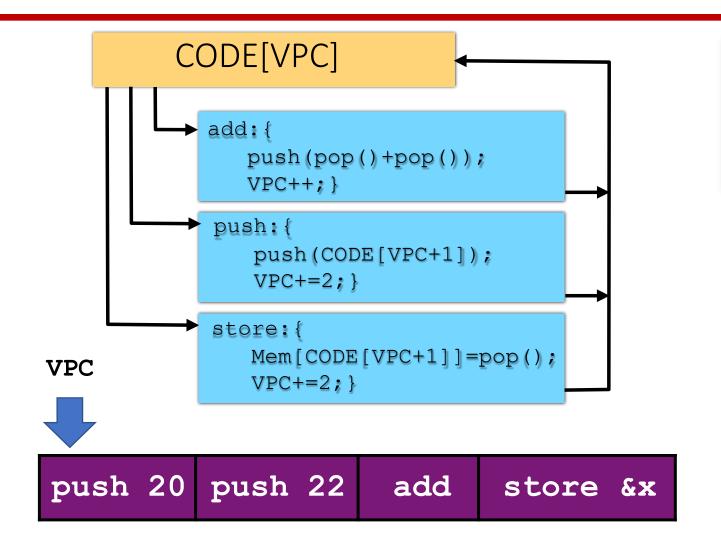
- Example: Translate code into high-level bytecode
- Bytecode has no registries, it only works with stack-based operations
- Create a virtual machine that implements such bytecode into platformspecific assembly
- Virtual Program Counter (VPC) iterates over the bytecode and execute it
- Stack Pointer (SP) points to the top of the Stack

```
push 20 push 22 add store &x
```

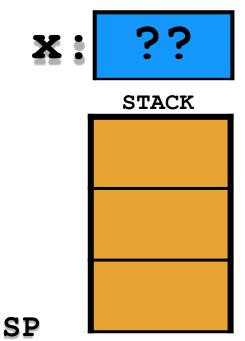


```
main() {
   int x;
   x = 20 + 22;
}
```



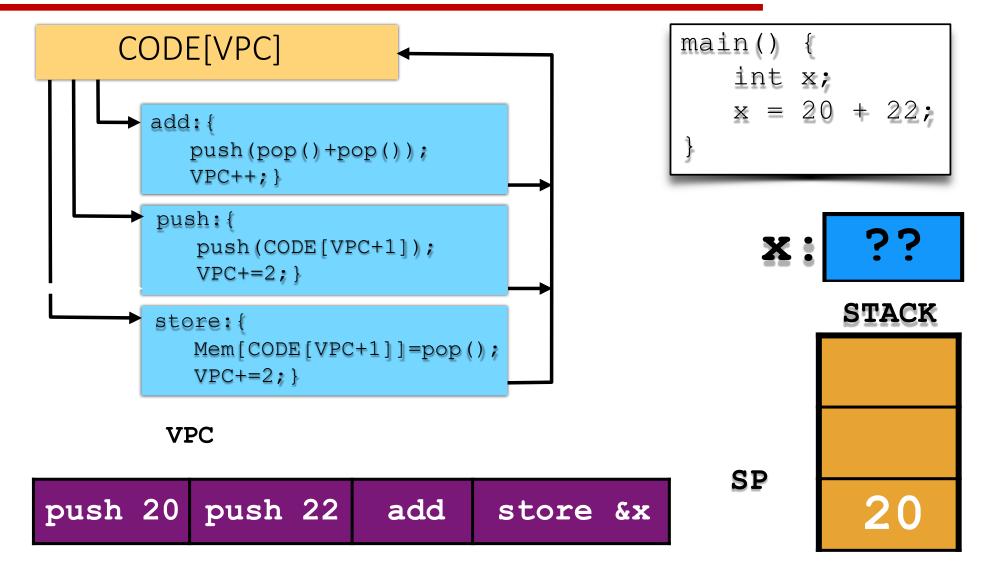


```
main() {
   int x;
   x = 20 + 22;
}
```



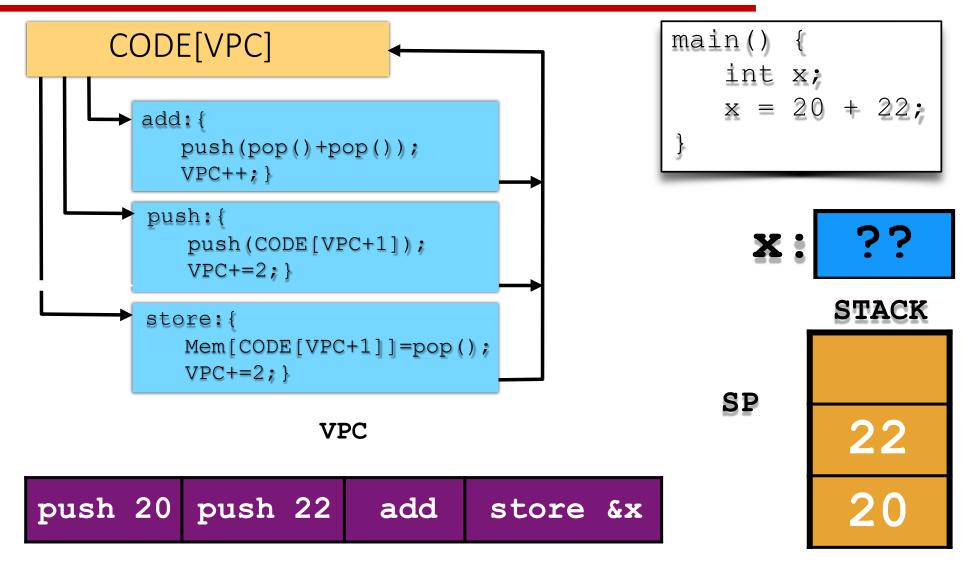
Example: Execute "Push 20"





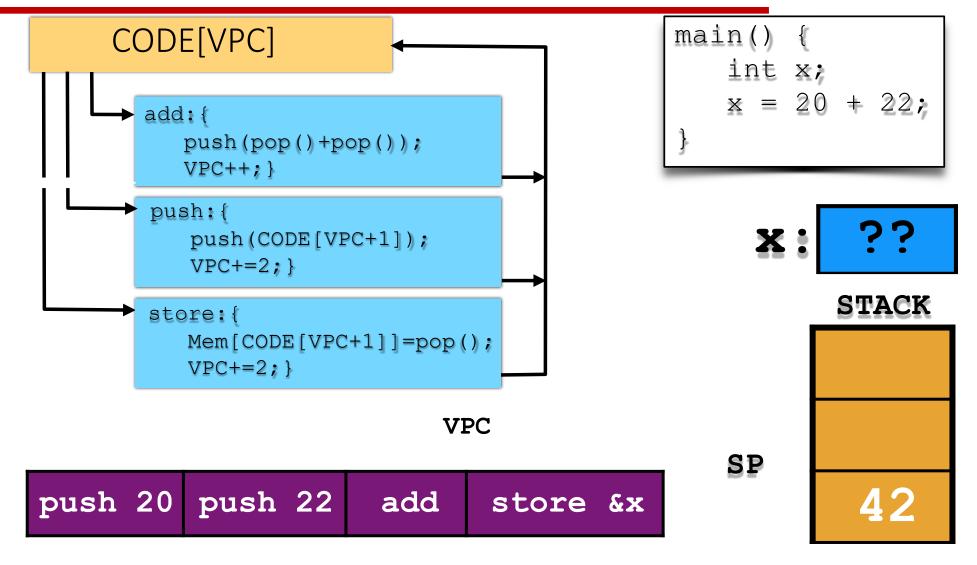
Example: Execute "Push 22"





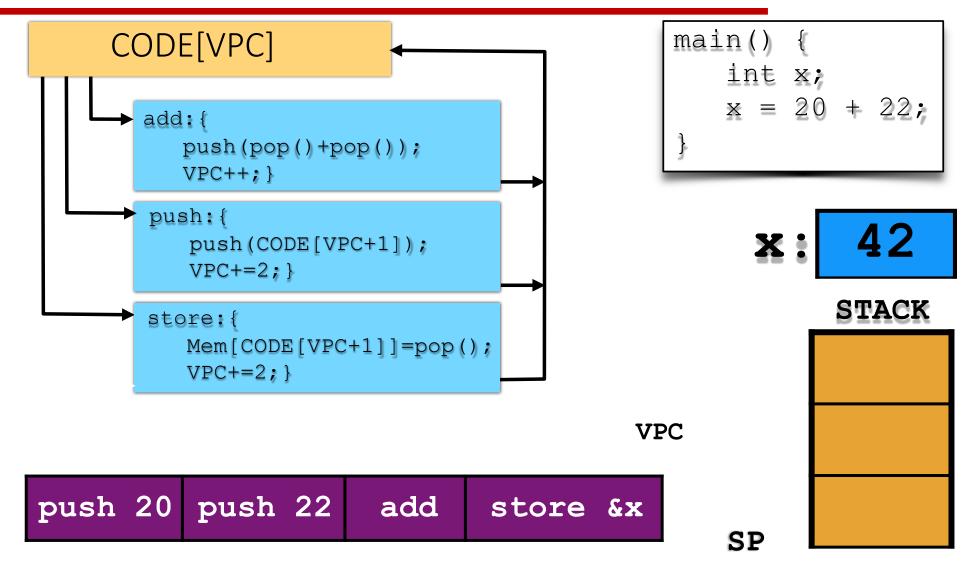
Example: add two numbers





Example: Store the result in memory





Liger VM



Here is a virtual machine that contains all the software tools you need for this course, developed by University of Arizona

https://ligerlabs.org/software.html

- Download it this week, and follow the instructions
- Try Tigress obfuscator examples

Tigress Installation



- In case you want to install Tigress only:
 - 1. Make sure you have the following installed:
 - gcc (or a drop-in replacement such as clang)
 - perl
 - Bash
 - 2. Download Tigress from https://tigress.wtf/
 - 3. Unzip the zip file
 - 4. Set this environment variable:
 - export TIGRESS_HOME=PATH-TO-THIS-DIRECTORY/tigress/3.1
 - 5. Add this directory to your PATH:
 - export PATH= ...: PATH-TO-THIS-DIRECTORY/tigress/3.1:...

Tigress Installation



Here's a simple test program you can use to try out Tigress: /bin/cp PATH-TO-THIS-DIRECTORY/tigress/3.1/test1.c.

```
Here's the Tigress command:
tigress --Environment=ENV \
--Transform=Virtualize \
--Functions=main,fib,fac \
--out=result.c test1.c
```

ENV depends on your platform; it should be one of x86_64:Linux:Gcc:4.6 x86_64:Darwin:Clang:5.1 armv7:Linux:Gcc:4.6 armv8:Linux:Gcc:4.6

This should give you an obfuscated program in result.c.

Exercises with Tigress



- Download these two files from our Moodle page or from here:
 - http://tigress.cs.arizona.edu/exercises.txt
 - http://tigress.cs.arizona.edu/fib.c
- The file exercises.txt contains Tigress commands I want you to run.
- The commands can be long and I don't want you to have to type them yourselves.
- The file fib.c contains the C file I want you to work on.

Example in exercises.txt

