Q1. Inventory control with forecasts

(a): Single period

Primitives *D*: Demand, which is a random variable

$$D = \begin{cases} P_L & \text{with probability q} \\ P_S & \text{with probability 1 - q} \end{cases}$$

p: Sale price per unit

c: Ordering cost per unit

s: Salvage value per unit

b = p - c: Backorder cost per unit

h = c - s: Inventory holding cost per unit

Q: Inventory decision.

The profit for this problem can be written as:

$$\pi(Q) = E[p.min\{D, Q\} - c.Q + s(Q - D)^{+}]$$

$$\pi(Q) = pE[D] - (p - c)E[(D - Q)^{+}] + (c - s)E[(Q - D)^{+}]$$

$$\pi(Q) = pE[D] - bE[(D - Q)^{+}] - hE[(Q - D)^{+}]$$

We then expand the formula to account for the expectation of demand D which takes value P_L with probability q:

$$\pi(Q) = pE[D] - b[q.E_{P_L}[(D-Q)^+] + (1-q).E_{P_S}[(D-Q)^+]] - h[qE_{P_L}[(Q-D)^+] + (1-q)E_{P_S}[(Q-D)^+]]$$

We take the derivative of the profit with respect to quantity in order to minimise, and set it equal to zero (with P_L and P_S the probability distributions in the large and small demand cases respectively):

$$\pi'(Q^*) = -b[qP_L(D > Q^*) + (1-q)P_S(D > Q^*)] + h[qP_L(Q^* > D) + (1-q)P_S(Q^* > D) = 0$$

$$-bq(1 - P_L(D \leq Q^*) - b(1 - q)(1 - P_S(D \leq Q^*)) + hqP_L(D \leq Q^*) + h(1 - q)P_S(D \leq Q^*) = 0$$

We can then obtain that the optimal decision Q^* satisfies:

$$qP_L(D \leqslant Q^*) + (1 - q)P_S(D \leqslant Q^*) = \frac{b}{b + h}$$

(b) Multi-period problem

Primitives

 x_k : Inventory level at period k

 u_k : Decision variable over the number of units to order.

 w_k : Demand at period k.

 $y_{k+1} = \xi$ where ξ takes the value "large demand" with probability q:

$$\xi_k = \begin{cases} L & \text{with probability q} \\ S & \text{with probability 1 - q} \end{cases}$$

c: ordering cost per unit

p: Sale price per unit

s: Salvage value

b = p - c: Backorder cost per unit

h = c - s: Inventory holding cost per unit

Dynamics:

$$\begin{aligned} x_{k+1} &= x_k + u_k - w_k \\ y_{k+1} &= \xi_k \\ g_N(x_N) &= -s(x_N) \\ g_k(x_k, u_k, w_k) &= cu_k - pw_k + bmax\{w_k - x_k - u_k, 0\} + hmax\{x_k + u_k - w_k\} \end{aligned}$$

DP algorithm:
$$J_N(x_N) = -s(x_N)$$

 $J_k(x_k) = \min_{u_k \ge 0} E_{w_k, \xi_k} [g_k(x_k, u_k, w_k) + J_{k+1}(x_k, u_k, w_k)]$

We can introduce the variable $z_k = x_k + u_k$, and rewrite our DP equation as follows:

$$J_k(x_k) = \min_{z_k \geqslant x_k} [G_k(z_k)] - cx_k$$

where
$$G_k(z_k) = cz_k + bE[max\{0, w_k - z_k\}] + hE[max\{0, z_k - w_k\}] + E[J_{k+1}(z_k - w_k)]$$

Once again we can expand this expression to take into account the two different values of ξ for large and small demand:

$$G_k(z_k) = cz_k + bqE[\max\{0, w_k - z_k\} | \xi = L] + b(1 - q)E[\max\{0, w_k - z_k\} | \xi = S] + hqE[\max\{0, z_k - w_k\} | \xi = L] + h(1 - q)E[\max\{0, z_k - w_k\} | \xi = S] + E[J_{k+1}(z_k - w_k)]$$

 $G_k(z_k)$ is a convex function for all k, and there will be a solution S_k which minimises G(z) such that the optimal ordering policy becomes:

 $u_k^* = \mu_k^* = S_k - x_k$ if $S_k > x_k$, and 0 otherwise.

Q2. Inventory Pooling

Primitives

D = demands

Q = quantity ordered

P = price

h = inventory costs = c-s

b = backholding costs = p-c

First we will show that $Q_p^* = \sqrt{n}Q^* + \mu(n - \sqrt{n})$

$$G(Q) = hE[(Q - D)]^{+}bE[(D - Q)]^{+}$$

$$G'(Q^{*}) = hP(D \leq Q^{*}) - b(1 - P[(D \leq Q^{*})]) = 0$$

$$P(D \leq Q^{*}) = \frac{b}{h + b}$$

Now, for the pooling we would obtain the same:

$$P(\sum_{L=1}^{n} D_i \leqslant Q_p^*) = \frac{b}{b+R}$$

Since this $\sum_{i=1}^{n} D_i = \sqrt{n}D_i + \mu(n - \sqrt{n}),$

$$P(\sqrt{n}D_1 + \mu(n - \sqrt{n}) \leqslant Q_p^*) = \frac{b}{b + R} \Longleftrightarrow P(D_1 \leqslant \frac{1}{\sqrt{n}}(Q_p^* - \mu(n - \sqrt{n}))) = \frac{b}{b + h}$$

which implies that $\frac{1}{\sqrt{n}}(Q_p^* - \mu(n - \sqrt{n}))) = Q^*$. Thus, $Q_P^* = \sqrt{n}Q^* + \mu(n - \sqrt{n})$. Next, we will apply the hint to prove the desired result:

$$nG(Q^*) = n[hE[(Q^* - D)^+] + bE[(D - Q^*)^+]]$$
$$G(Q_p^*) = [hE[(Q_p^* - \sum_{i=1}^n D_i)^+] + bE[(\sum_{i=1}^n D_i - Q_p^*)^+]]$$

Since $\sum_{i=1}^{n} D_i = \sqrt{n}D_i + \mu(n-\sqrt{n})$ and $Q_P^* = \sqrt{n}Q^* + \mu(n-\sqrt{n})$, we can show that:

$$G(Q_P^*) = [hE[(\sqrt{n}Q^* + \mu(n - \sqrt{n}) - \sqrt{n}D_i - \mu(n - \sqrt{n}))^+]$$

$$+bE[\sqrt{n}D_i + \mu(n - \sqrt{n}) - \sqrt{n}Q^* + \mu(n - \sqrt{n})^+]]$$

$$= \sqrt{n}hE[(Q^* - D_i)^+] + \sqrt{n} bE[(D_i - Q^*)^+] = \frac{nG(Q^*)}{\sqrt{n}}$$

Q.E.D.

1 Q4 An Investment Problem

Primitives:

 x_k : wealth at the beginning of period k.

 u_k : amount invested.

 w_k : outcome (return).

$$w_k = \begin{cases} 2u_k & \text{w.p. } p \\ 0 & \text{w.p. } 1 - p \end{cases}$$
 where $\frac{1}{2} .$

Constrains:

 $u_k \geqslant 0$,

 $u_k \leqslant x_k$.

Dynamics:

$$x_{k+1} = x_k - u_k + w_k,$$

Cost:

$$g_N(x_N) = x_N$$

$$g_k(x_k, u_k, w_k) = 0.$$

DP algorithm:

To simplify th $J_N(x_N) = \log(x_N)$. $J_k(x_k) = \max_{u_k \in \mathcal{U}_k} \mathbb{E}_{w_k} [J_{k+1}(x_k - u_k + w_k)].$

We will first prove that $J_k(x_k) = A_k + \log(x_k)$ by induction (where log is the logarithm in base e).

$$J_{N-1}(x_{N-1}) = \max_{u_{n-1}} \mathbb{E}_{w_{N-1}} \left[\log(x_{N-1} - u_{N-1} + w_{N-1}) \right] =$$
 (1)

$$= \max_{u_{N-1}} \mathbb{E}_{w_{N-1}} [\log(x_{N-1} - u_{N-1} + w_{N-1})] =$$
 (2)

$$= \max_{u_{N-1}} \left[p \log(x_{N-1} + u_{N-1}) + (1-p) \log(x_{N-1} - u_{N-1}) \right]. \tag{3}$$

Now, we derive and apply FOC to find the maximum,

$$\frac{p}{x_{N-1} + u_{N-1}} - \frac{1-p}{x_{N-1} - u_{N-1}} = 0 \tag{4}$$

$$px_{N-1} - pu_{N-1}^* = x_{N-1} + u_{N-1}^* - px_{N-1} - pu_{N-1}^*, (5)$$

solving we obtain $u_{N-1}^* = (2p-1)x_{N-1}$, which makes sense, because 2p-1 is between 0 and 1 for $\frac{1}{2} . Plugging into the expression of <math>J_{N-1}(x_{N-1})$

$$J_{N-1}(x_{N-1}) = p\log(x_{N-1} + (2p-1)x_{N-1}) + (1-p)\log(x_{N-1} - (2p-1)x_{N-1}) =$$
(6)

$$= p\log(2px_{N-1}) + (1-p)\log(2(1-p)x_{N-1}) =$$
(7)

$$= p \log(2p) + p \log(x_{N-1}) + (1-p) \log(2(1-p)) + (1-p) \log(x_{N-1}) = (8)$$

$$= p\log(2p) + (1-p)\log(2(1-p)) + \log(x_{N-1}). \tag{9}$$

So, for N-1 we have $A_{N-1} = p \log(2p) + (1-p) \log(2(1-p))$. Now, suppose $J_{k+1}(x_{k1}) = A_{k+1} + \log(x_{k+1})$ let's prove it for k:

$$J_k(x_k) = \max_{u_k} \mathbb{E}_{w_k} \left[A_{k+1} + \log(x_{k+1}) \right] = \tag{10}$$

$$= \max_{u_k} \mathbb{E}_{w_k} \left[A_{k+1} + \log(x_k - u_k + w_k) \right] = \tag{11}$$

$$= \max_{u_k} \left[A_{k+1} + p \log(x_k + u_k) + (1-p) \log(x_k - u_k) \right]. \tag{12}$$

Applying first order conditions we would obtain the same equation as before $(A_{k+1}$ does not depend on u_k), so $u_k = (2p-1)x_k$. Plugging in the $J_k(x_k)$ we obtain the same way as before

$$J_k(x_k) = A_{k+1} + p\log(2p) + (1-p)\log(2(1-p)) + \log(x_k).$$
(13)

So, $J_k(x_k) = A_k + \log(x_k)$, where $A_k = A_{k+1} + p\log(2p) + (1-p)\log(2(1-p))$. So, the policy will be to invest $(2p-1)x_k$ of wealth each period k.