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Depth First Search: Suppose G is a connected undirected graph. An edge whose removal disconnects the graph is called a *bridge*. Either prove the following statement or provide a counter-example: every bridge e must be an edge in a depth-first search tree of G .

Solution

This is true. Suppose it were not true. Consider the edge e that connects vertices u and v , where e is a bridge. That is, if e were removed from the graph G , then we would have a partition of the vertices of G such that no edge crosses the partition. Start a depth first search in the partition that includes u . Eventually, since G is connected, this DFS must reach vertex u . When it does, it will examine all of the vertices adjacent to u and incorporate any who have not yet been “touched” into the DFS-tree. This includes v . We are guaranteed that v has not yet been examined since the edge (u, v) is the only edge that crosses our partition of vertices, and we started our DFS in the partition that included u . Therefore, the edge (u, v) must be included in the DFS, giving us our contradiction.