## **Type Theorists HATE Him!**

Learn this ONE WEIRD TRICK to fake dependent types in a language that doesn't support them

LEARN THE TRUTH NOW

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## **Dependent types**

### Dependent types... in Idris

### Dependent types... in Haskell?

### Dependent types... in Haskell?

<interactive>:1:28: error: parse error on input 'if'



### There's hope yet

- A Haskell can support dependent types, if you're willing to squint
- Ne'll need to enable a modest number of GHC extensions

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- A Haskell can support dependent types, if you're willing to squint
- Ne'll need to enable a modest number of GHC extensions

```
{-# LANGUAGE DefaultSignatures
                                         #-}
{-# LANGUAGE EmptyCase
                                         #-}
{-# LANGUAGE ExistentialQuantification
                                         #-}
{-# LANGUAGE FlexibleContexts
                                         #-}
{-# LANGUAGE FlexibleInstances
                                         #-}
{-# LANGUAGE GADTs
                                         #-}
{-# LANGUAGE InstanceSigs
                                         #-}
{-# LANGUAGE KindSignatures
                                         #-}
{-# LANGUAGE RankNTypes
                                         #-}
{-# LANGUAGE ScopedTypeVariables
                                         #-}
{-# LANGUAGE TemplateHaskell
                                         #-}
{-# LANGUAGE TypeFamilies
                                         #-}
{-# LANGUAGE TypeInType
                                         #-}
{-# LANGUAGE TypeOperators
                                         #-}
{-# LANGUAGE UndecidableInstances
                                         #-}
```

```
data family Sing :: k -> Type
```

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```

```
data instance Sing :: Bool -> Type where
   SFalse :: Sing False
   STrue :: Sing True
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data instance Sing :: Bool -> Type where
   SFalse :: Sing False
   STrue :: Sing True

data instance Sing (z :: Bool)
   = (z ~ False) => SFalse
   | (z ~ True) => STrue
```

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data instance Sing :: Bool -> Type where
   SFalse :: Sing False
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data instance Sing (z :: Bool)
   = (z ~ False) => SFalse
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```
data family Sing :: k -> Type

data Nat = Z | S Nat

data instance Sing :: Nat -> Type where
   SZ :: Sing Z
   SS :: Sing (n :: Nat) -> Sing (S n)
```

```
data Nat = Z | S Nat
data instance Sing :: Nat -> Type where
  SZ :: Sing Z
  SS :: Sing (n :: Nat) -> Sing (S n)
data instance Sing (z :: Nat)
  = (z \sim Z) \Rightarrow SZ
  | forall (n :: Nat). (z \sim S n) \Rightarrow SS n
```

data family Sing :: k -> Type

### Dependent types... in Idris (redux)

```
type family
  If (c :: Bool) (t :: k) (f :: k) :: k where
  If True  t f = t
  If False  t f = f
```

```
type family
 If (c :: Bool) (t :: k) (f :: k) :: k where
 If True t f = t
 If False t f = f
mkSingle :: Sing (x :: Bool) ->
            If x Nat [Nat]
mkSingle STrue = 0
mkSingle SFalse = []
```

type  $\Pi = Sing$ 

type  $\Pi = Sing$ 

Dependent pattern matching

Dependent pattern matching... in Idris

```
data Vect : Nat -> Type -> Type where
  Nil : Vect Z a
  Cons : a -> Vect n a -> Vect (S n) a
```

Dependent pattern matching... in Idris

```
data Vect : Nat -> Type -> Type where
  Nil : Vect Z a
  Cons : a -> Vect n a -> Vect (S n) a

len : Vect n a -> Nat
len Nil = 0
len (Cons x xs) = 1 + len xs
```

Dependent pattern matching... in Idris

```
data Vect : Nat -> Type -> Type where
  Nil : Vect Z a
  Cons : a -> Vect n a -> Vect (S n) a

len : {n : Nat} -> Vect n a -> Nat
len {n=Z} Nil = 0
len {n=S k} (Cons x xs) = 1 + len xs
```

```
data Vect :: Nat -> Type -> Type where
  Nil :: Vect Z a
  Cons :: a -> Vect n a -> Vect (S n) a
```

```
class SingI (a :: k) where
  sing :: Sing (a :: k)

instance SingI Z where
  sing = SZ

instance SingI n => SingI (S n) where
  sing = SS sing
```

- No With enough elbow grease, one can simulate a great deal of dependently typed code
- Impress your friends at the bar! Be the envy of your family!
- http://hackage.haskell.org/package/singletons

## Any questions?