Deriving-via

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We present a new Haskell language extension that miraculously solves all problems in generic programming that ever existed.

ACM Reference Format:

"These types we write down they're not just names for data representations in memory, they're tags that queue in mathematical structures that we exploit."

1 INTRODUCTION

It is common folklore that Monoids can be lifted over Applicatives,

```
instance (Appliative f, Monoid a) ⇒ Monoid (f a) where
mempty :: f a
mempty = pure mempty
mappend :: f a → f a → f a
mappend = liftA2 mappend
```

Conor McBride calls this "routine programming" using ${\tt Monoid}$ and ${\tt Applicative}$ as building blocks. 2

But this instance is undesirable for multiple reasons (TODO: more reasons, rewrite)

- It overlaps with every Monoid instance over an applied type.
- "Structure of the f is often considered more significant that that of x."
- It may not be the desired Monoid: Some constructors have an 'inherent monoidal structure', most notably the *free monoid* (lists: [a]) where we prioritize the list structure and not that of the elements.

Lists are in fact an instance of a wholly separate way of defining Monoids based on Alternative

```
instance Alternative f ⇒ Monoid (f a) where
mempty :: f a
mempty = empty
```

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XXXX-XXXX/2017/11-ART

https://doi.org/10.1145/nnnnnn.nnnnnnn

¹Taken from unknown position: https://www.youtube.com/watch?v=3U3lV5VPmOU

²http://strictlypositive.org/Idiom.pdf

³Much of this is stolen from Conor: https://personal.cis.strath.ac.uk/conor.mcbride/so-pigworker.pdf

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```
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        mappend :: f a \rightarrow f a \rightarrow f a
        mappend = (<|>)
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      So what are our options.
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        An unfortunate solution is to duplicate code
      instance Monoid a >> Monoid (IO a) where
55
        mempty = pure mempty
        mappend = liftA2 mappend
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      instance (Monoid a, Monoid b) ⇒ Monoid (a, b) where
        mempty = pure mempty
59
        mappend = liftA2 mappend
61
      instance Monoid b \Rightarrow Monoid (a \Rightarrow b) where
62
        mempty = pure mempty
63
        mappend = liftA2 mappend
```

but this quickly becomes unviable as Num, Floating and Fractional which amount to around 50 methods lifted in the exact same way. Another solution provided by Conal Elliott is to use preprocessor.

But we already have

2 EXAMPLES

3 FORMALISM

4 ADVANCED USES

 \bullet Avoiding orphan instances Before we had a Monoid (IO a) instance, we could not write 5

```
newtype Plugin = Plugin (IO (String → IO ()))
deriving Monoid
```

deriving via enables us to override and insert arbitrary instances adding the following line

```
via App IO (String -> App IO ())
```

• Asymptotic improvement For representable functors the definitions of m *> $_$ = m and $_$ <* m = m become O(1) 6

4.1 Generalized GeneralizedNewtypeDeriving

- 4.2 DeriveAnyClass
- 5 LIMITATIONS, CONCLUSIONS AND FUTURE WORK
- 6 RELATED WORK

⁴https://hackage.haskell.org/package/applicative-numbers

⁵http://www.haskellforall.com/2014/07/equational-reasoning-at-scale.html

⁶Edward Kmett: https://ghc.haskell.org/trac/ghc/ticket/10892?cversion=0&cnum_hist=4#comment:4