IT Security



Chapter 3: Checksums and Digital Signatures Part 1

- Hash functions
- Message Authentication Code MAC
- Signature procedure
- PKI



What do we want to learn?

- What is the meaning of hash functions?
- What is a cryptographic hash function, such as MAC?
- How does a digital signature work?
- What is the legal significance of electronic signatures?
- What do you have to consider in the practical implementation?
- What are the components of a PKI?
- How is a certificate structured, how do you verify it?





How do I ensure the integrity of my data?



- Problem:
 - I have a lot of data and want a short checksum to make sure the data has not been modified.
- Possible solutions
 - Hashing
 - Cryptographic hashing
 - Digital signatures

Not safe against tampering!

Hash Functions

BUT: Hash functions are not encryption!

One-way encryption functions:

Mapping of variable length input to "unique" output with 0 fixed length.



- Serves as checksum, fingerprint, message digest
- Properties
 - small change in message results in large change in hash value
 - message cannot be reconstructed from hash value (function cannot be inverted)
 - hash value as unique as possible collision resistant: it is practically impossible to find two values x1 and x2 with f(x1) = f(x2)
 - hash value can be calculated fast and by everyone

Not recomended anymore

Algorithms: MD5 (128 bit), SHA-1 (160 bit),
SHA-2 (SHA-256 (256 bit), SHA-384, SHA-512, SHA-224)
SHA-3



Security of hash functions

- Insecure "classic" methods: MD5, SHA1, RIPEMD
- Recommendations and key lengths for cryptographic methods by BSI

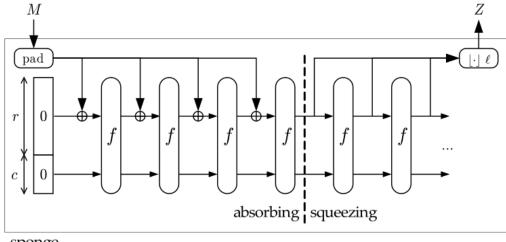
Methods	Recommended algorithms
cryptographic strong hash functions	SHA-256, SHA-512/256, SHA-384 und SHA-512, SHA3-256, SHA3-384, SHA3-512
MAC methods	HMAC>= 128, CMAC>=128, GMAC>=128
Signature methods	RSA 3000, DSA 3000, ECDSA 250, Merkle signatures

https://www.bsi.bund.de/SharedDocs/Downloads/EN/BSI/Publications/TechGuidelines/TG02102/BSI-TR-02102-1.pdf? blob=publicationFile&v=6



New hash standard since 2012: SHA-3 (Keccak)

- Winner from a competition by the NIST (National Institute of Standards and Technology)
 http://csrc.nist.gov/groups/ST/hash/sha-3/
 https://keccak.team/index.html
- Important requirement: hash must be quickly computable
- Output size is variable (224, 256, 384, 512 bit)
- Security can be influenced by capacity using a sponge function (more security -> less performance)

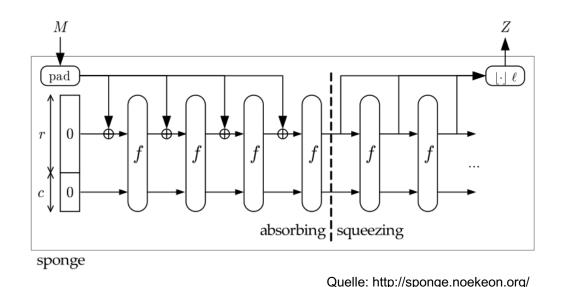


sponge

Quelle: http://sponge.noekeon.org/



Details of SHA-3 (Keccak)



Absorbing:

- state is divided into two parts: Capacity (c bits) and Bit-Rate (r bits).
- input blocks M are padded (filled up)
- r-bit part is XORed with state, c-bit part remains unchanged
- then state is mixed with f (Keccak permutation).
- Squeezing
 - Hash value is extracted from r-bit portion of state
 - If length is not sufficient, state is permuted several times and further outputs are extracted after each permutation.
- Visualization see CrypTool 2 (https://www.cryptool.org/de/ct2/)

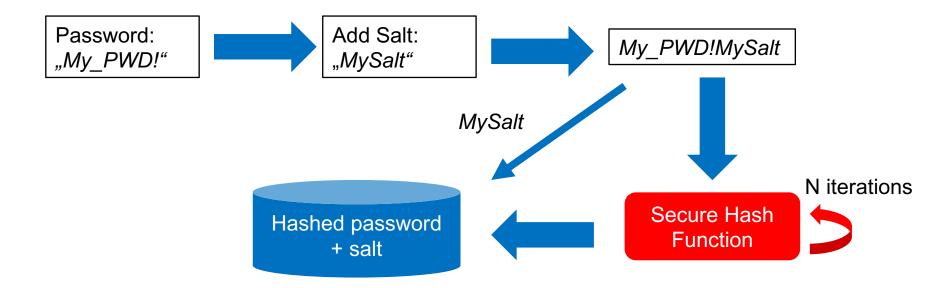


Hash functions for password hashing

- Problems
 - normal Hash functions are very fast and support brute force attacks
 - same password results in same hash value
- Criteria for Secure Hash functions https://www.password-hashing.net/cfh.html
 - add cost parameter to tune time and/or space usage
 - provide cryptographic security (no speed up for hackers possible)
 - combine it with a **Salt** to produce different hash values for same password **Salt**: random data as additional input for the hash function, generate new salt for each password
- Secure Password Hashing Algorithms
 - Bcrypt, Scrypt, Argon2, PBKDF2



Process for secure password hashing



Data authentication

- Data authentication means cryptographic procedures that guarantee that transmitted or stored data has not been modified by unauthorized persons.
 - it is based on cryptographic keys that are used to calculate checksums.
 - there are symmetric and asymmetric procedures
- Two security goals can be achieved with data authentication
 - ensuring the integrity of data
 - securing the non-repudiation of a message
 - → this is only possible with digital signatures

https://www.bsi.bund.de/SharedDocs/Downloads/EN/BSI/Publications/TechGuidelines/TG02102/BSI-TR-02102-1.pdf? blob=publicationFile&v=6



Message Authentication Code (MAC)

- A hash function that additionally uses a secret key
- The key is known only to the two communication partners
- This enables the detection of unauthorized modifications (integrity)
- This enables the verification if data have an authentic origin (authenticity)

First attempt for a MAC

- Simple concatenation hash(Secret || Daten) = MAC
- Problem:
 - if the attacker knows | length(secret || data)
 - then he can compute for some hash functions hash(secret || data || more data) without knowing the secret!
- This is called: **Length extension attack**https://en.wikipedia.org/wiki/Length_extension_attack
- Vulnerable are e.g. SHA1, SHA2, not vulnerable is SHA3
- Problem: you must create a hash, which can only be calculated with a secret.
 - one possible solution: HMAC

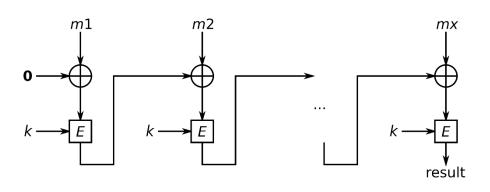
HMAC (Hashed MAC)

- HMAC is a popular MAC
- \vdash $HMAC(K,m) = H((K \oplus opad) \parallel H((K \oplus ipad) \parallel m))$
 - K is the secret key
 - block size is 64 bytes: either pad K with zeros or reduce to 64 bytes with a hash function.
 - popad, ipad are constants in block size
- A video illustrating HMAC
 - https://www.youtube.com/watch?v=BjInMA-b8ZE



CBC-MAC (Cipher Block Chaining MAC)

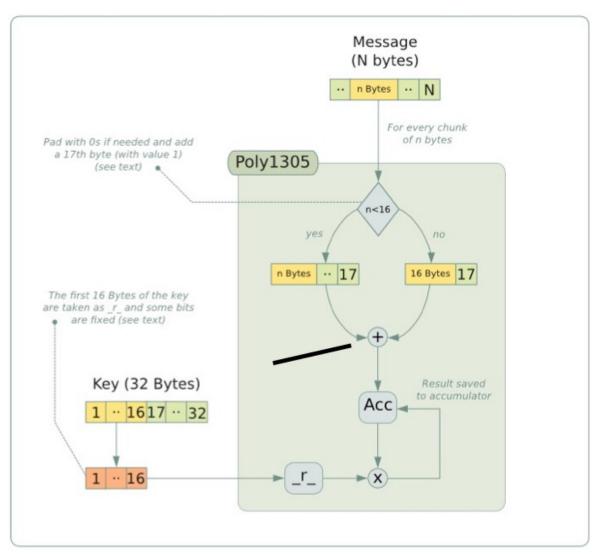
- CBC-MAC is a technique to calculate a MAC using a block cipher method.
- It uses block cipher method in CBC mode
- The IV is set to 0, the MAC is the last encrypted block
- $c_1 = (m_1 \oplus 0) \oplus K \qquad c_x = (m_x \oplus c_{x-1}) \oplus K$
- CBC-MAC is only safe for messages of fixed/known length, otherwise a length extension attack is possible.



Quelle: By Benjamin D. Esham (bdesham) - Own work based on: Cbcmac.png by en:User:PeterPearson.Own work by bdesham using: Inkscape., Public Domain, https://commons.wikimedia.org/w/index.php?curid=2277179

Poly1305 MAC

- Poly1305 is a MAC which uses a 32 Byte secret key and generates a 16 Byte authenticator
- Poly1305 is faster than HMAC



Quelle: https://www.adalabs.com/adalabs-chacha20poly1305.pdf

Applications and limitations of MAC

- Modern ciphering methods combine encryption with MAC calculation
 - Counter Mode CTR with CBC-MAC
 - ▶ **GCM** (see chapter 2: GCM is a cipher with MAC calculation)
 - ChaCha20 with Poly1305
- Application of MAC
 - attack detection for file systems (Filesystem Intrusion Detection)
 - securing software packages (patch, update)
 - securing of communication protocols (e.g. TLS)
- A MAC secures the integrity and authenticity of a message
 - but it lacks the evidential value for non-repudiation
 - it cannot be verified by a third-party authority
 - it is based on symmetric cryptography
 - it lacks certificates



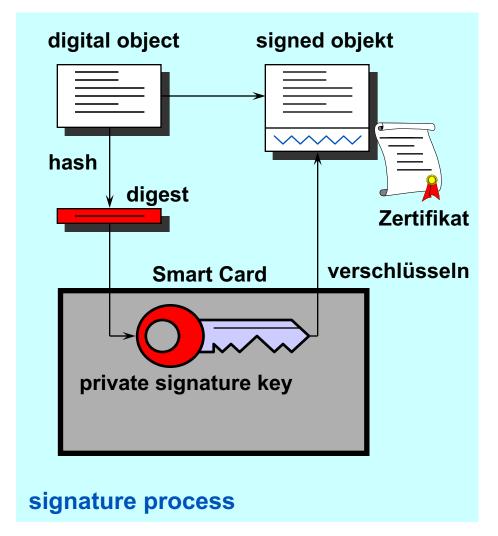
Digital signatures ensure the integrity of data and prove authenticity

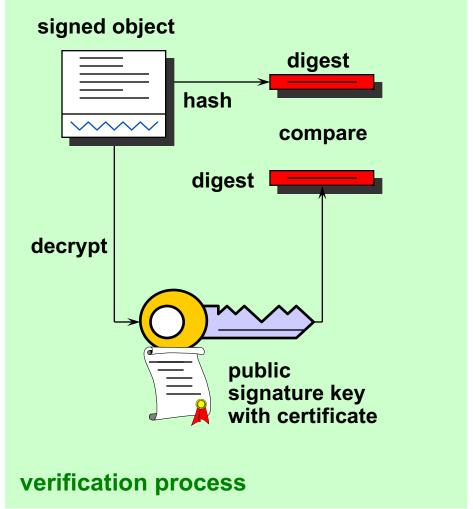
- A Digital Signature contains
 - time from a timestamp service
 - person or place (server name)
 - signed digest of the document
- A Digital Signature proves
 - integrity of the document (what)
 - who signed and when





The procedure of the Digital Signature





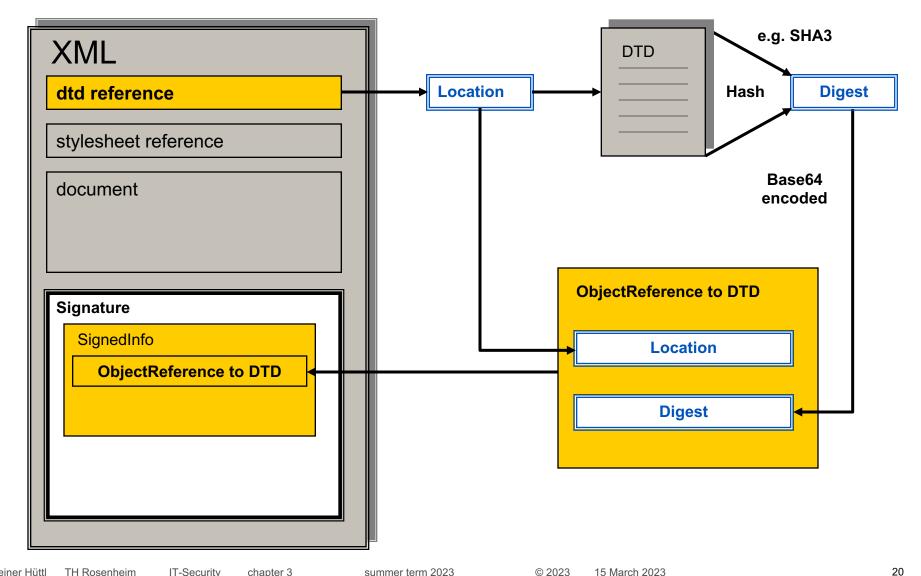


Standard for XML signatures

- Defines rules and syntax for signature
 - of whole XML documents
 - of parts of XML documents
 - any other files
 - literature: http://www.w3.org/Signature/
- Three ways to integrate XML signatures
 - Detached Signature: The signature is detached from the document and not embedded in the document.
 - **Enveloped Signature**: The signature is embedded in the document.
 - **Enveloping Signature**: The signature has the function of an envelope. It encloses the entire XML document.



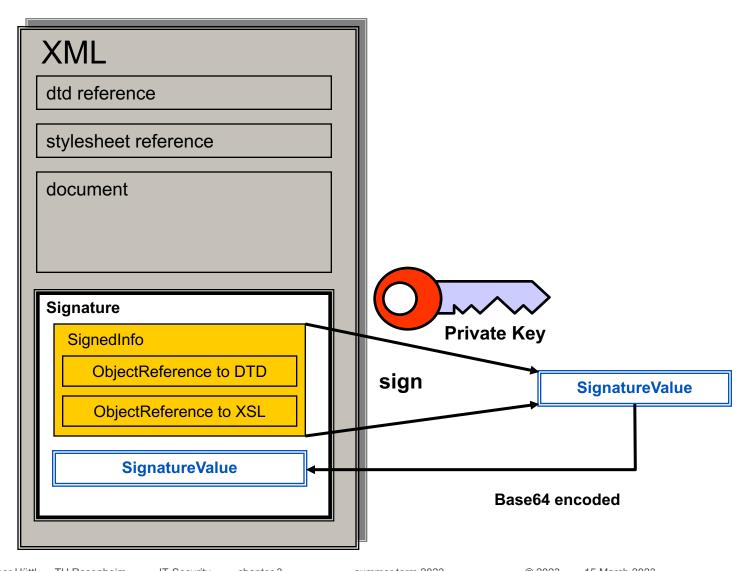
XML-Signature: Digest calculation, reference creation



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XML-Signature: Signature of the objekt references





Components of an XML signature

```
Document to sign
<?xml version="1.0" encoding="UTF-8"?>
<Envelope xmlns="urn:envelope">
                                                                                  Processing information
 <Signature xmlns="http://www.w3.org/2000/09/xmldsig#">
   <SignedInfo>
    <CanonicalizationMethod Algorithm="http://www.w3.org/TR/2001/REC-xml-c14n- 20010315#WithComments"/>
    <SignatureMethod Algorithm="http://www.w3.org/2000/09/ xmldsig#rsa-sha1"/>
                                                                                   Referenced data with
    <Reference URI="">
      <Transforms>
                                                                                   processing information
        <Transform Algorithm="http://www.w3.org/2000/09/ xmldsig#enveloped-signature"/>
                                                                                   and digest
      </Transforms>
      <DigestMethod Algorithm="http://www.w3.org/2000/09/ xmldsig#sha1"/>
      <DigestValue>uoogbWYa5VCgcJCbuvmBKgm17vY=</DigestValue>
    </Reference>
   </SignedInfo>
   <SignatureValue> KedJuTob5gtvYx9gM3k3gm7kbLBwVbEQRl26S2tmXjgNND7MRGtoew== </SignatureValue>
   <KevInfo>
                                                                         Base64 encoded signature value
    <KeyValue>
      <RSAKeyValue>
                                                                         over the SignedInfo element
        <Modulus>
           4llzOY3Y9fXoh3Y5f06wBbtTg94Pt6vcfcd1KQ0FLm0S36aGJtTSb6pYKfyX7PgCUQ8wgL6xUJ5GRPEsu9gyz8
           ZobwfZsGCsvu40CWoT9fcFBZPfXro1Vtlh/xl/yYHm+Gzqh0Bw76xtLHSfLfpVOrmZdwKmSFKMTvNXOFd0V18=
        </Modulus>
        <Exponent>AQAB</Exponent>
                                                                         Information about the key to verify
      </RSAKeyValue>
                                                                         the signature
    </KeyValue>
   </KeyInfo>
 </Signature>
                                XML tags for XML signature
</Envelope>
```

Canonicalization

XML documents with semantically the same content can be represented in different ways:

```
<myelement attr="123"/>
<myelement attr="123"></myelement>
```

- This gives the signature a completely different value
- Therefore, a standardized representation is necessary: Canonization
- Literature: https://www.w3.org/TR/xml-c14n/

