IT Security

Chapter 2: Encryption (Part 1)

- Symmetric encryption
- Base64 encoding
- Asymmetric encryption
- Practical aspects of encryption
- Random numbers
- Key Management

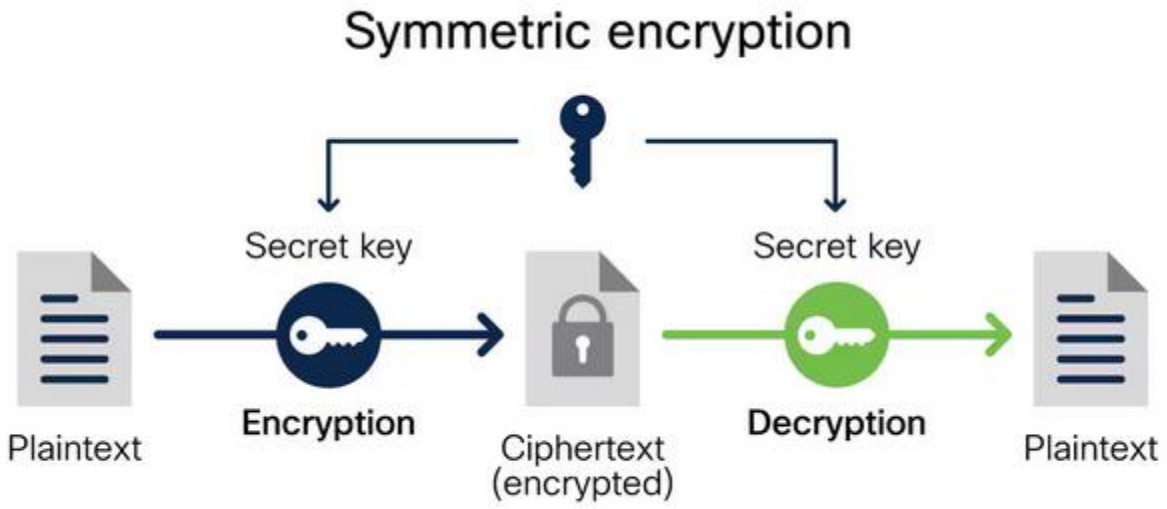
Prof. Dr. Reiner Hüttl, FH Rosenheim, © 2023, 15.03.23

What do we want to learn?

- What are the types of encryption?`
- What are the most important algorithms and procedures?
- How can I use encryption in practice?
- What do I have to pay attention to so that encryption really brings more security?

Symmetric encryption

- One central key for encryption and decryption
- Disadvantages: exchange of secret key required
- Advantages:
  - Very fast,
  - implementable in HW



Quelle: https://www.cisco.com/c/en/us/products/security/encryption-explained.html#~q-a

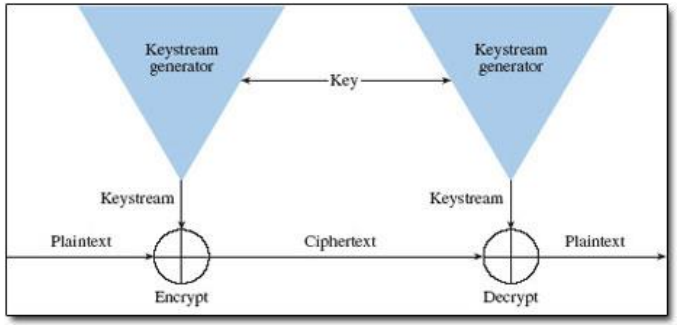
Symmetric Cipher Variants

- **Stream ciphers**
  - Take a key of fixed length
  - Generate a data stream of arbitrary length, consisting of pseudo-random numbers, from the key
  - Convert each bit of the plaintext to ciphertext using XOR
  - Examples: RC4 (UNSECURE since 2013!!!), ChaCha20
- **Block ciphers**
  - Process plaintext and ciphertext in blocks (e.g. 64 Bit, 128 Bit)
  - Have a mode (block cipher mode)
  - Examples: DES, DES3 (UNSECURE!!), AES, Twofish, Serpent
- **Padding**
  - Filling up the last incomplete block with "regular" patterns.
  - Reason: Some algorithms (mostly block ciphers) require full blocks
  - Examples: PKCS5 padding, W3C padding, ISO padding, ESP padding.

irregular but reconstructable

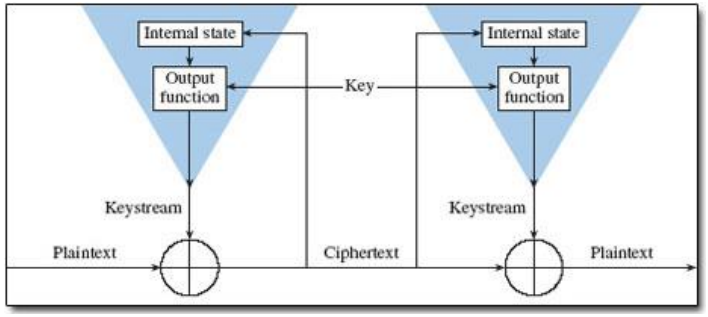
Stream cipher variants

- Synchronous stream ciphering
  - Generates the key stream independently of the clear text or key text



From Schneier, *1996, Figure 9.6*

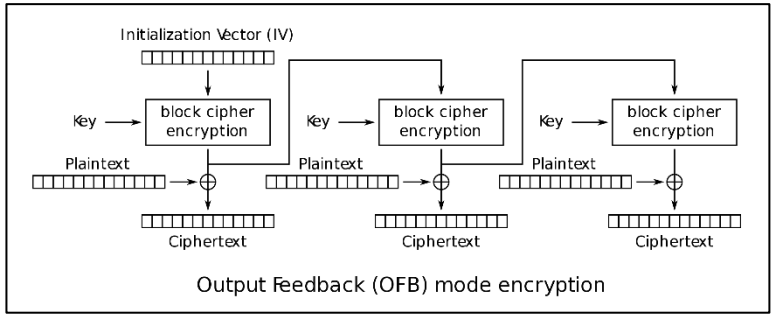
- Self-synchronizing stream ciphering
  - Key stream depends on previous encrypted bit

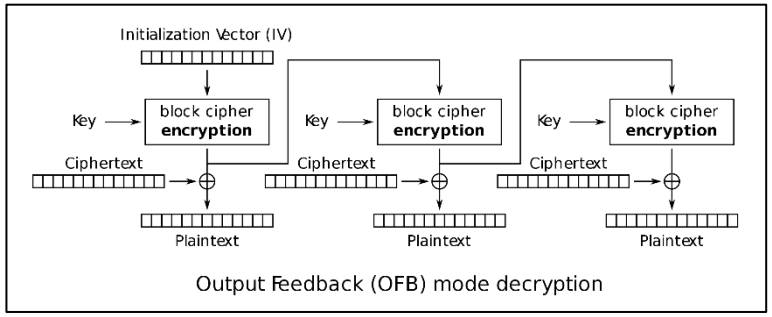


From Schneier, *1996, Figure 9.8*

Block ciphers can be implemented as stream ciphers

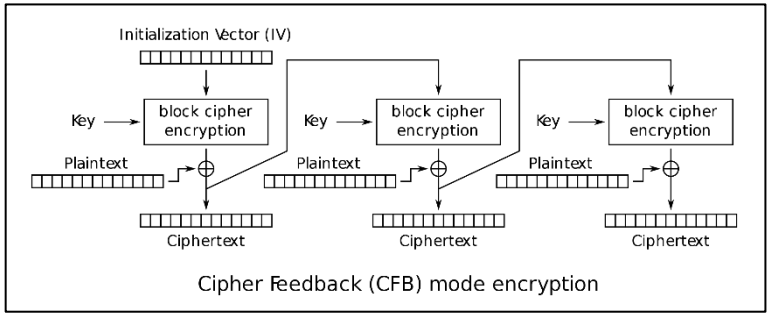
- **OFB, Output Feedback Mode**
- Is a synchronous stream ciphering mode
- The key stream can be precalculated

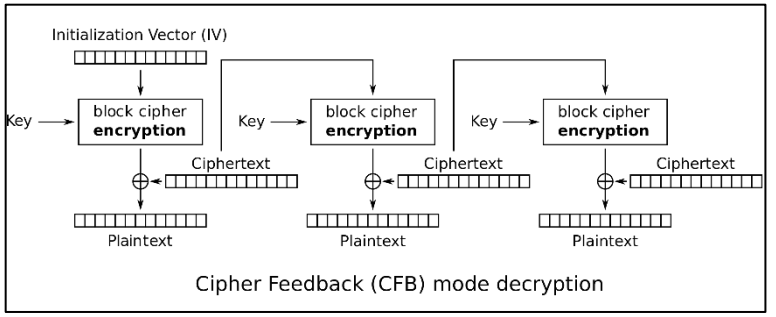




https://de.wikipedia.org/wiki/Output_Feedback_Mode

- **CFB Cipher Feedback Mode**
- Self-synchronizing stream ciphering
- Errors in IV or ciphertext affect only two blocks





https://de.wikipedia.org/wiki/Cipher_Feedback_Mode

The standard encryption method Advanced Encryption Standard (AES)

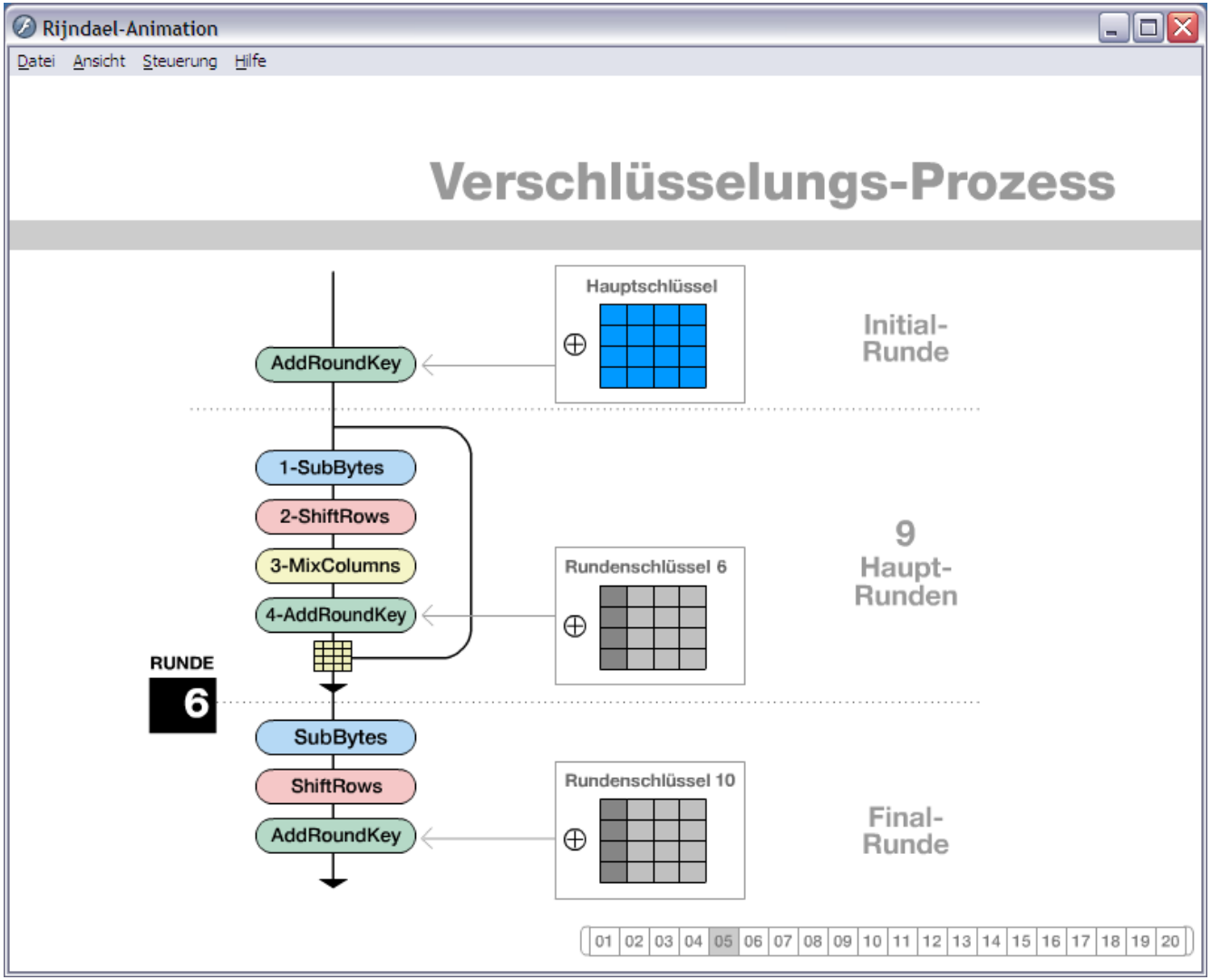
- Requirements for AES in competition from NIST 2001
  - Block cipher with 128, 192 and 256 bits
  - Mathematical justification of security
  - Simplicity of design
  - Flexibility in block sizes and key lengths
  - Efficiency
  - Easy to implement in HW and SW

Old standard, unsafe today, do not use anymore!!!

- Official successor to DES
  - Based on Rijndael algorithm
  - Selected from 15 proposals in an open process
  - All design criteria published
  - Encrypts blocks with fixed length of 128 bits (16 bytes)
  - Key length optionally 128, 192 or 256 bits

recommended key length

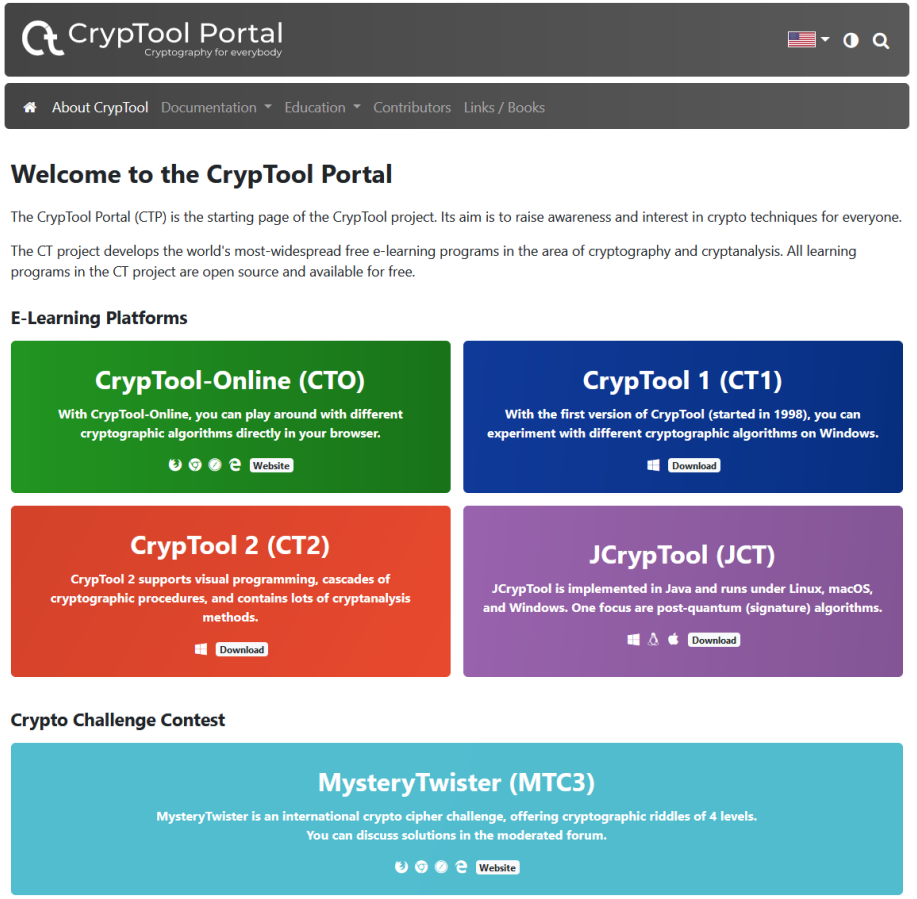
The Algorithm of AES



Quelle: Cryptool Portal,
https://www.cryptool.org/en/cto/aes-animation

Small exercise: View simulation of AES

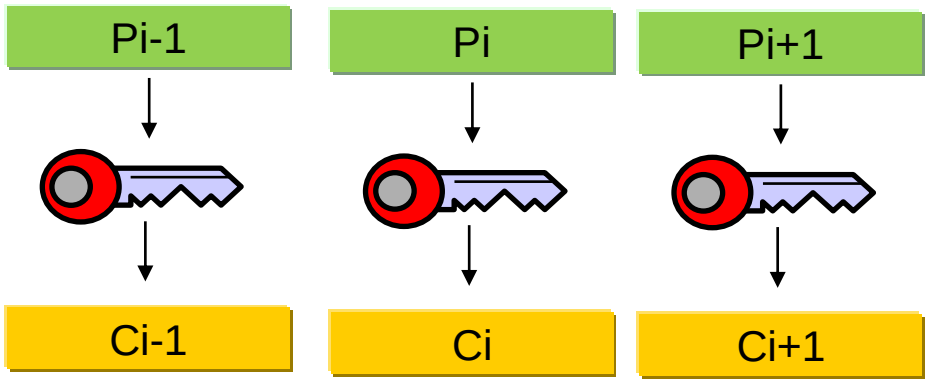
- Go to CryptTool Portal https://www.cryptool.org/en/
- The CT project develops free e-learning programs in the area of cryptography and cryptanalysis
- Prof. Bernhard Esslinger, University of Siegen
- Go to CrypTool-Online (CTO)
  Go to AES Animation
  https://www.cryptool.org/en/cto/aes-animation
- Go through the animation step by step



Operating mode for block cipher: ECB

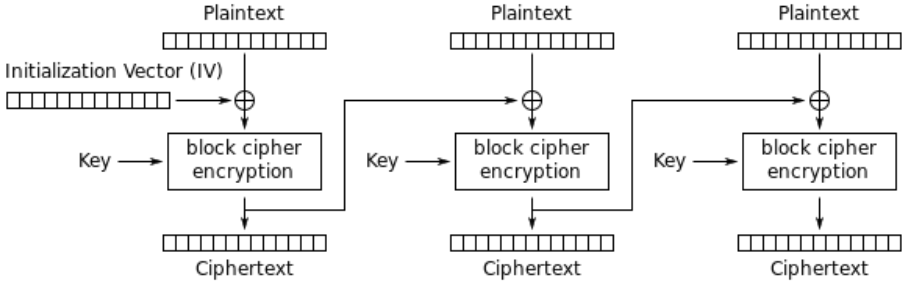
- A symmetric method (e.g., AES) can be operated in different modes

- **Electronic Code Book (ECB)**
  - A plaintext block is encoded into a ciphertext block
  - Same plaintext blocks generate same ciphertext blocks
  - All blocks can be ciphered independently of each other
  - Vulnerable for cryptanalysis (stereotyped beginnings and endings)
  - Simple bit errors have no influence on other blocks
  - If bits are lost, resynchronization of block boundaries is required

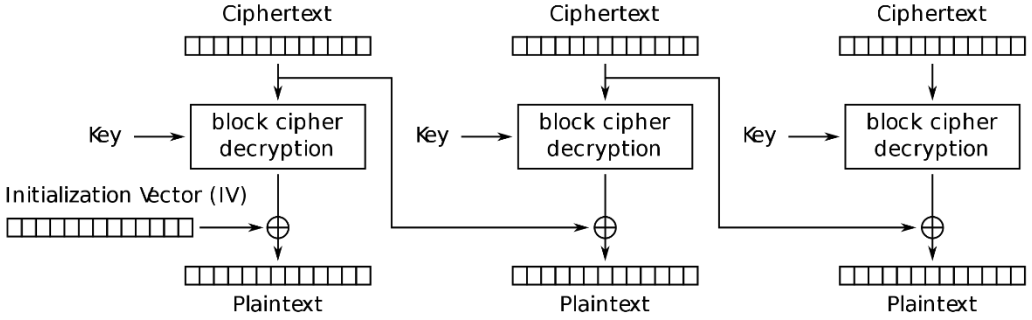


Operating mode for block cipher: CBC

- **Cipher Block Chaining (CBC)**
  - Feedback: encryption of a block depends on predecessor blocks
  - Plaintext is XORed with previous ciphertext block before encryption
  - **Initialization vector** for first block necessary:
    not secret, uniquely reconstruct able for decryption, different for each message
  - Problem: error propagation
    Bit error: 1-bit error in ciphertext leads to error in block and following block
    Synchronization error: no recovery anymore



Cipher Block Chaining (CBC) mode encryption

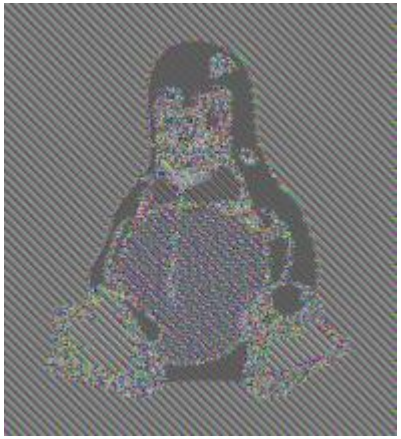


Cipher Block Chaining (CBC) mode decryption

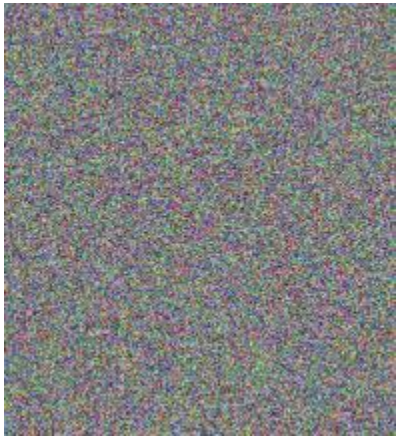
$$C_i = E_K(P_i \oplus C_{i-1})$$
$$P_i = C_{i-1} \oplus D_K(C_i)$$

ECB vs. CBC

- Popular error with encryption: Use wrong mode
- **ECB is bad**, CBC is better
- But much better is GCM (comes later)



ECB



CBC

Attention when choosing the initialization vector

- Initialization vector
  - Does not have to be kept secret (just send it along with the ciphertext)
  - Should be random (e.g. SecureRandom)
  - Should be different for each message
  - Is the same size as the block size, for AES **always** 128 bit
- Popular errors
  - Use Key as IV
    - This may cause the key to be read in the worst case [1]
  - Set IV to 0 or other fixed value
    - IV brings randomness into the system - therefore roll dice randomly!
  - Generate IV from password
    - Better chosse IV randomly

[1] https://crypto.stackexchange.com/questions/16161/problems-with-using-aes-key-as-iv-in-cbc-mode

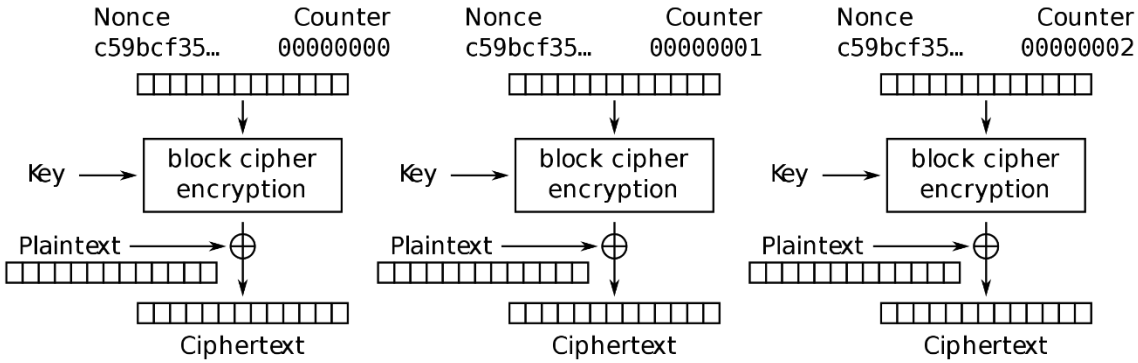
Operating mode CTR Counter Block Mode

- For each block a new unpredictable keystream block is calculated:
  Keystream block = IV(Nonce) + current counter + encryption key

Nonce = Number used only once

- Advantages:
  - Key does not depend on the key of the previous block
  - Encryption and decryption can be performed in parallel
  - Key stream can be precomputed (when operating as a stream cipher)
  - Bit errors affect only one block

$$C_i = P_i \oplus E_K(ctr_i)$$



Counter (CTR) mode encryption

Application of cryptography is difficult

- I use AES/CBC, all good? **NO!**
  - Attacker can modify ciphertext and thus influence plaintext
  - Attacker can perform **Padding Oracle Attack** and read plaintext, see https://www.arxumpathsecurity.de/blog/2019/10/16/cbc-mode-is-malleable-dont-trust-it-for-authentication
- Other popular error
  - Never use CBC, CTR, etc. without authentication!
  - Otherwise attacker can read / modify plaintext
- That‘s why: Secure encryption with authentication
  - MAC Message Authentication Code
  - MAC = Hash function using a secret key

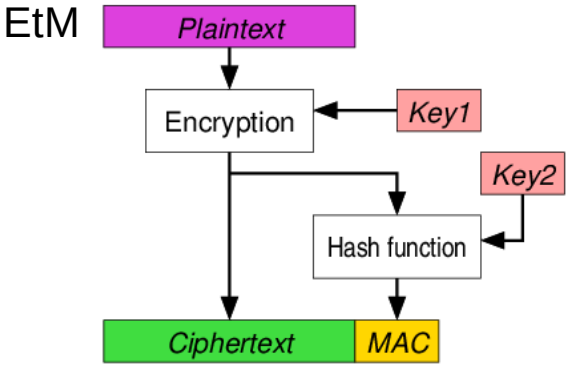
Authentication
Goal of secure identification of a person or a machine

MAC is discussed in more detail in the Digital Signatures chapter

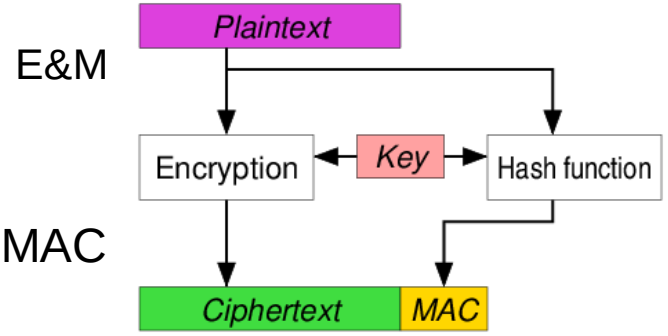
Authenticated Encryption

- Encryption methods are called authenticated if not only the confidentiality but also the integrity of the data to be encrypted is protected.

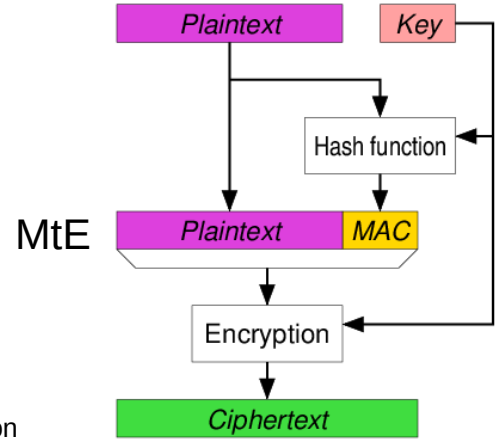
- Encrypt-then-MAC
  - MAC on ciphertext
  - **High Security** if both keys are different



- Encrypt-and-MAC
  - MAC on plaintext
  - Plaintext is encrypted without MAC



- MAC-then-Encrypt
  - MAC on plaintext
  - MAC and plaintext are encrypted



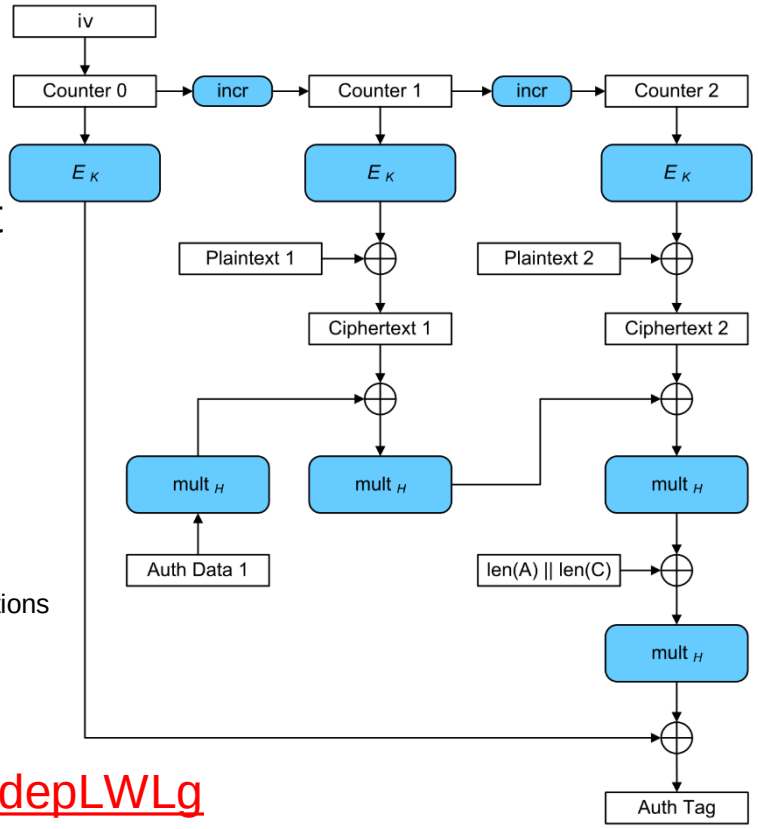
Quelle: https://en.wikipedia.org/wiki/Authenticated_encryption

Authenticated Encryption

- Encrypt-then-MAC
  - AES+CBC+HMAC-SHA256
    https://www.mkammerer.de/blog/encrypt-something-with-aes-how-hard-can-it-be/
  - ChaCha20 (encrypt) + Poly1305 (MAC)
    https://github.com/phxql/chacha20-poly1305-java
- Popular erros
  - Use same key for AES and HMAC
    - Is ok in certain cases, if possible: don't do it
  - MAC-then-encrypt
    - AES(HMAC(Plaintext) || Plaintext)
    - There are padding attacks (e.g. with SSL)
  - MAC-and-encrypt
    - AES(plaintext) || HMAC(plaintext)
    - No integrity on ciphertext

Operating mode GCM Galois Counter Mode

- Authenticated encryption mode with associated data (AEAD)
  - Fast, as parallelization is possible
  - Can also be used as stand-alone MAC
  - Accepts IV (nonce) of any length
  - Auth Tag contains Auth Data, IV and ciphertext



Ek Encryption with Key k

Len (a) = len (Auth Data)

Len (C) = len (ciphertext)

Mult H Galois field multiplications

- Video illustrating GCM

https://www.youtube.com/watch?v=R2SodepLWLg

https://de.qwe.wiki/wiki/Galois/Counter_Mode

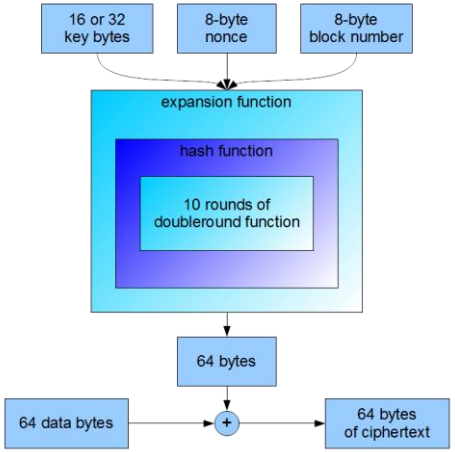
AEAD Authenticated Encryption with Associated Data

- AEAD combines confidentiality, authenticity and integrity
  - is automatically secure against ciphertext modification
  - Combine encryption and MAC into one integrated protocol
  - e.g. AES-GCM
    - Uses a 96 bit nonce
    - Nonce (IV) need not be random, but must NEVER be reused with same key
  - Further variant AES-EAX

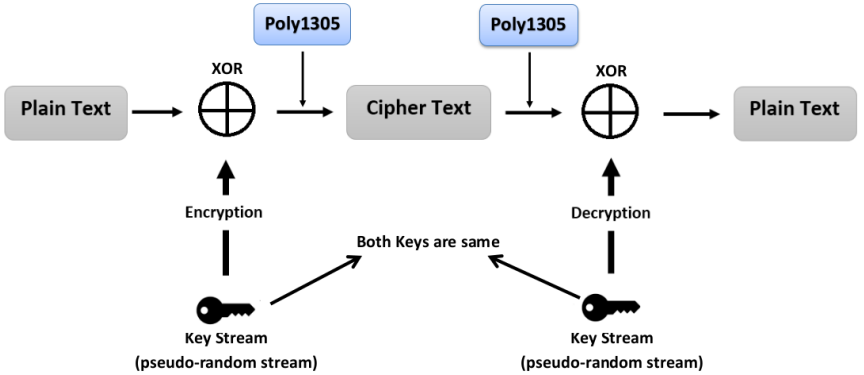
- Video illustrating AEAD
  https://www.youtube.com/watch?v=od44W45sCQ4

AEAD with stream cipher ChaCha20 and checksum Poly1305

- ChaCha20
  - Symmetric stream cipher based on Salsa20 by Daniel Bernstein
  - Implementation easier than AES-GCM
  - Used in IPSec, TLS, OpenSSL, OpenSSH
- Poly1305 is much faster as MAC than HMAC



Block Diagram of Salsa20 Algorithm
http://www.crypto-it.net/eng/symmetric/salsa20.html



ChaCha20 Poly1305 Encryption and Decryption scheme
https://javainterviewpoint.com/chacha20-poly1305-encryption-and-decryption/

Base64 Encoding

Base64 is an encoding and not an encryption!

- For transmission of binary data in sequences of 8-bit bytes across channels that only reliable support text content, e.g. the World Wide Web.
- Representation of binary data with 64 printable ASCII characters.
- Principle:
  - Split 24 bits into 4 parts of 6 bits
  - Extend each 6 bit sequence to 8 bits
- Disadvantage: message becomes 33% longer
- Application:
  - Keys, signatures and certificates are often stored or transmitted BASE64 encoded.
  - E-Mail Attachments

Example for Base64 Encoding

Base64 Alphabet

Wert	Zeichen	Wert	Zeichen	Wert	Zeichen	Wert	Zeichen
0	A	17	R	34	i	51	z
1	B	18	S	35	j	52	0
2	C	19	T	36	k	53	1
3	D	20	U	37	l	54	2
4	E	21	V	38	m	55	3
5	F	22	W	39	n	56	4
6	G	23	X	40	o	57	5
7	H	24	Y	41	p	58	6
8	I	25	Z	42	q	59	7
9	J	26	a	43	r	60	8
10	K	27	b	44	s	61	9
11	L	28	c	45	t	62	+
12	M	29	d	46	u	63	/
13	N	30	e	47	v	64	=
14	O	31	f	48	w		
15	P	32	g	49	x		
16	Q	33	h	50	y		

Beispiel abcde			
Quelle	**Binärdarstellung**	**Base64**	
a	01100001		
b	01100010		
c	01100011		
24-Bit	011000010110001001100011		
6-Bit	011000	24	Y
	010110	22	W
	001001	9	J
	100011	35	j
d	01100100		
e	01100101		
	____		
24-Bit	011001000110010100____		
6-Bit	011001	25	Z
	000110	6	G
	010100	20	U
	____		=
YWJjZGU=			