Exercise 6.1 (Computing). Determine the Cholesky decompositions of the following matrices. That is, a lower triangular matrix G with  $A = GG^T$ .

a) 
$$A = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 8 & 12 \\ 3 & 12 & 27 \end{pmatrix}$$

b) 
$$A = \begin{pmatrix} 1 & -2 & 3 & -4 \\ -2 & 8 & -8 & 14 \\ 3 & -8 & 11 & -14 \\ -4 & 14 & -14 & 35 \end{pmatrix}$$

#### Suggested Solution.

(a) Step k = 1:

(i) 
$$g_{11} = \sqrt{a_{11}^{(1)}} = \sqrt{1} = 1$$

(ii) 
$$g_{21} = \frac{a_{ik}^{(k)}}{g_{kk}} = \frac{a_{21}^{(1)}}{g_{11}} = \frac{2}{1} = 2,$$

$$g_{31} = \frac{a_{31}^{(1)}}{g_{11}} = \frac{3}{1} = 3$$

(i) 
$$g_{kk} := \sqrt{a_{kk}^{(k)}}$$

(ii) 
$$g_{ik} := \frac{a_{ik}^{(k)}}{g_{kk}}$$
 für  $i \ge k+1$ 

$$\begin{array}{l} \text{for } k=1,\ldots,n:\\ (\mathrm{i}) \ \ g_{kk} \coloneqq \sqrt{a_{kk}^{(k)}}\\ \\ (\mathrm{ii}) \ \ g_{ik} \coloneqq \frac{a_{ik}^{(k)}}{g_{kk}} \ \text{für } i \geq k+1\\ \\ (\mathrm{iii}) \ \ a_{ij}^{(k+1)} \coloneqq a_{ij}^{(k)} - g_{ik}g_{jk} \ \text{für } i,j \geq k+1\\ \\ \text{endfor} \end{array}$$

(iii) 
$$(g_{i1}g_{j1})_{2 \le i,j \le 3} = \binom{2}{3}(2;3) = \binom{g_{21}^2 \times g_{31}^2}{g_{31}g_{21}g_{31}^2} = \binom{4}{6}\binom{6}{9}$$
  

$$\Rightarrow (a_{ij}^{(2)})_{2 \le i,j \le 3} = (a_{ij}^{(1)})_{2 \le i,j \le 3} - (g_{i1}g_{j1})_{2 \le i,j \le 3} = \binom{8}{12}\binom{12}{27} - \binom{4}{6}\binom{6}{9} = \binom{4}{6}\binom{6}{18}$$

Step k=2:

(i) 
$$g_{22} = \sqrt{a_{22}^{(2)}} = \sqrt{4} = 2$$

(ii) 
$$g_{32} = \frac{a_{32}^{(2)}}{g_{22}} = \frac{6}{2} = 3$$

(iii)  $(q_{i2}q_{i2})_{3 \le i, i \le 3} = (q_{32}^2) = (9)$  and thus

$$(a_{ij}^{(3)})_{3 \le i,j \le 3} = (a_{ij}^{(2)})_{3 \le i,j \le 3} - (g_{i1}g_{j1})_{3 \le i,j \le 3} = (18) - (9) = (9)$$

Step k = 3:

(i) 
$$g_{33} = \sqrt{a_{33}^{(3)}} = \sqrt{9} = 3$$

Finally we obtain  $A = GG^t$  with  $G = \begin{pmatrix} 1 & 0 & 0 \\ 2 & 2 & 0 \\ 3 & 3 & 3 \end{pmatrix}$ .

(b) From

$$A = \begin{pmatrix} 1 & -2 & 3 & -4 \\ -2 & 8 & -8 & 14 \\ 3 & -8 & 11 & -14 \\ -4 & 14 & -14 & 35 \end{pmatrix}$$

we obtain

# Step k=1:

(i) 
$$g_{11} = \sqrt{a_{11}^{(1)}} = \sqrt{1} = 1$$

(ii) 
$$g_{21} = \frac{a_{ik}^{(k)}}{g_{kk}} = \frac{a_{21}^{(1)}}{g_{11}} = \frac{-2}{1} = -2,$$

$$g_{31} = \frac{a_{31}^{(1)}}{g_{11}} = \frac{3}{1} = 3$$

$$g_{41} = \frac{a_{41}^{(1)}}{g_{11}} = \frac{-4}{1} = -4$$

for 
$$k = 1, \ldots, n$$
:

(i) 
$$g_{kk} := \sqrt{a_{kk}^{(k)}}$$

(ii) 
$$g_{ik} := \frac{a_{ik}^{(k)}}{g_{kk}}$$
 für  $i \ge k+1$ 

$$\begin{cases} \text{for } k=1,\dots,n: \\ \text{(i) } g_{kk} \coloneqq \sqrt{a_{kk}^{(k)}} \\ \\ \text{(ii) } g_{ik} \coloneqq \frac{a_{ik}^{(k)}}{g_{kk}} \text{ für } i \geq k+1 \\ \\ \text{(iii) } a_{ij}^{(k+1)} \coloneqq a_{ij}^{(k)} - g_{ik}g_{jk} \text{ für } i,j \geq k+1 \\ \\ \text{endfor} \end{cases}$$

(iii) 
$$(g_{i1}g_{j1})_{2 \le i,j \le 4} = \begin{pmatrix} -2 \\ 3 \\ -4 \end{pmatrix} \begin{pmatrix} -2; 3; -4 \end{pmatrix} = \begin{pmatrix} g_{21}^2 & \star & \star \\ g_{31}g_{21} & g_{31}^2 & \star \\ g_{41}g_{21} & g_{41}g_{31} & g_{41}^2 \end{pmatrix} = \begin{pmatrix} 4 & -6 & 8 \\ -6 & 9 & -12 \\ 8 & -12 & 16 \end{pmatrix}$$

$$\Rightarrow (a_{ij}^{(2)})_{2 \le i,j \le 4} = (a_{ij}^{(1)})_{2 \le i,j \le 4} - (g_{i1}g_{j1})_{2 \le i,j \le 4} = \begin{pmatrix} 8 & -8 & 14 \\ -8 & 11 & -14 \\ 14 & -14 & 35 \end{pmatrix} - \begin{pmatrix} 4 & -6 & 8 \\ -6 & 9 & -12 \\ 8 & -12 & 16 \end{pmatrix} = \begin{pmatrix} 4 & \star & \star \\ -2 & 2 & -2 \\ 6 & -2 & 19 \end{pmatrix}$$

#### Step k=2:

(i) 
$$g_{22} = \sqrt{a_{22}^{(2)}} = \sqrt{4} = 2$$

(ii) 
$$g_{32} = \frac{a_{32}^{(2)}}{g_{22}} = \frac{-2}{2} = -1$$

$$g_{42} = \frac{a_{42}^{(2)}}{q_{22}} = \frac{6}{2} = 3$$

(iii) 
$$(g_{i2}g_{j2})_{3 \le i,j \le 4} = \begin{pmatrix} g_{32}^2 & \star \\ g_{42}g_{32} & g_{42}^2 \end{pmatrix} = \begin{pmatrix} 1 & -3 \\ -3 & 9 \end{pmatrix}$$
 and thus

$$(a_{ij}^{(3)})_{3 \le i,j \le 4} = (a_{ij}^{(2)})_{3 \le i,j \le 4} - (g_{i1}g_{j1})_{3 \le i,j \le 4} = \begin{pmatrix} 2 & -2 \\ -2 & 19 \end{pmatrix} - \begin{pmatrix} 1 & -3 \\ -3 & 9 \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 10 \end{pmatrix}$$

## Step k = 3:

(i) 
$$g_{33} = \sqrt{a_{33}^{(3)}} = \sqrt{1} = 1$$

(ii) 
$$g_{43} = \frac{a_{43}^{(3)}}{g_{33}} = \frac{1}{1} = 1$$

(iii) 
$$(g_{i3}g_{j3})_{4 \le i,j \le 4} = (g_{43}^2) = (1)$$
 and thus

$$A^{(4)} = (a_{ij}^{(4)})_{4 \le i,j \le 4} = (a_{ij}^{(3)})_{4 \le i,j \le 4} - (g_{i1}g_{j1})_{4 \le i,j \le 4} = (10) - (1) = (9)$$

## Step k = 4:

(i) 
$$g_{44} = \sqrt{a_{44}^{(4)}} = \sqrt{9} = 3$$

Finally we obtain 
$$A = GG^t$$
 with  $G = \begin{pmatrix} 1 & 0 & 0 & 0 \\ -2 & 2 & 0 & 0 \\ 3 & -1 & 1 & 0 \\ -4 & 3 & 1 & 3 \end{pmatrix}$ .

The Octave command chol(A, 'lower')=chol(A)' computes G.

**Exercise 6.2** (Octave). Implement the Cholesky method in Octave, i.e. write a program that requires a matrix A as input and - if A is spd - returns the decomposition matrix G as output, so that  $A = GG^T$  holds. Test the program with the matrices from Exercise 6.1.

## Suggested Solution.

```
function G = ex6_2_cholesky(A)
    % Initialize G as a zero matrix.
    [n,n] = size(A);
    G = zeros(n);
    for k=1:n
        % Step k part (i)
        G(k,k) = sqrt(A(k,k));
        for i=k+1:n
            G(i,k)=A(i,k)/G(k,k);
        end
        % Step k part (ii)
        % Compute auxiliary matrix A(k)
        % Overwrite A thereby
        A(k+1:n,k+1:n) = A(k+1:n,k+1:n) - G(k+1:n,k)*transpose(G(k+1:n,k));
        end
endfunction
```

**Exercise 6.3.** Show that for all regular  $A \in \mathbb{R}^{n \times n}$ 

$$\operatorname{cond}_2(A) \leq \operatorname{cond}_F(A) \leq n \operatorname{cond}_2(A)$$
.

Show that the last inequality is sharp.

**Suggested Solution.** For all regular A we know that

$$||A||_2 \le ||A||_F$$

from Exercise 4.2 (d) and

$$||A||_F \le \sqrt{n}||A||_2$$

from Exercise 5.3 (b). From this we obtain for all regular A

$$\operatorname{cond}_{2}(A) = \|A\|_{2} \|A^{-1}\|_{2} \le \underbrace{\|A\|_{F} \|A^{-1}\|_{F}}_{=\operatorname{cond}_{F}(A)} \le \sqrt{n} \|A\|_{2} \sqrt{n} \|A^{-1}\|_{2} = n \operatorname{cond}_{2}(A).$$

The last inequality is sharp for A := E:

$$\underbrace{\|E\|_F}_{=\sqrt{n}} \le \sqrt{n} \underbrace{\|E\|_2}_{=1}.$$

**Exercise 6.4.** Let  $A \in \mathbb{R}^{2 \times 2}$  be symmetric and assume that

$$\det(A) = 1$$
 and  $\frac{\operatorname{trace}(A)}{2} = N$ ,

where N > 0. Show the following:

- (a) It is necessarily true that  $N \geq 1$ .
- **(b)** It is

$$\operatorname{cond}_2(A) \ge 4N^2 - 2N - 1 \quad \text{for } N \ge 1.$$
 (1)

**Suggested Solution.** A sym.  $\Rightarrow V^T A V = D = \operatorname{diag}(\lambda_1, \lambda_2)$  for some  $V \in \mathcal{O}_n$  and  $\lambda_{1/2} \in \mathbb{R}$ . Since  $|A| = \lambda_1 \lambda_2 = 1$  and  $\operatorname{trace}(A) = \lambda_1 + \lambda_2 > 0$  it follows (without loss of generality) that  $\lambda_1 \geq \lambda_2 > 0$ . First, we transform the quadratic characteristic polynomial into normal form

$$\chi_A(\lambda) = \chi_D(\lambda) 
= (\lambda_1 - \lambda)(\lambda_2 - \lambda) 
= \lambda^2 - (\lambda_1 + \lambda_2)\lambda + \lambda_1\lambda_2 
= (\lambda - N)^2 - (N^2 - 1).$$
(2)

(a) From (2) we deduce that

$$\lambda_{1/2} \in \mathbb{R} \iff N^2 - 1 \geq 0 \stackrel{N \geq 0}{\Leftrightarrow} N \geq 1.$$

(b) Additionally, from Equation 2 we can write the eigenvalues of A in terms of N:

$$\lambda_{1/2} = N \pm \sqrt{N^2 - 1}.$$

From this we can finally deduce

$$\operatorname{cond}_2(A) = \frac{\lambda_1}{\lambda_2} = \frac{N + \sqrt{N^2 - 1}}{N - \sqrt{N^2 - 1}} = N^2 + 2N\sqrt{N^2 - 1} + N^2 - 1.$$
 (3)

The claim follows when rhs of  $(3) \ge \text{rhs of } (1)$ :

$$2N^2 + 2N\sqrt{N^2 - 1} - 1 \ge 4N^2 - 2N - 1$$
  

$$\Leftrightarrow \qquad \sqrt{N^2 - 1} \ge N - 1$$
  

$$\Leftrightarrow \qquad N^2 - 1 \ge N^2 - 2N + 1$$
  

$$\Leftrightarrow \qquad N \ge 1,$$

which is true by (a).

**Exercise 6.5** (Octave). Write an Octave program that calculates the largest and smallest eigenvalues of a matrix  $A \in Gl(n)$ . Avoid the calculation of the inverse matrix  $A^{-1}$  and use the LU decomposition of A instead (without row permutations).

#### Suggested Solution.

```
function [lambdamin,lambdamax] = ex6_5_powermethod(A,nmax)
  % nmax = maximum number of iterations
  % Computation of the largest and smallest absolute value
  [n,n] = size(A);
  % Initialize start vector
  v = ones(n,1);
  % Computation of lambdamax
  for j=1:nmax
    v = A*v;
    v = v/norm(v);
    lambdamax = v'*A*v;
  % Computation of lambdamin
  [L,U] = LU_wo_permutations(A); v = ones(n,1);
  for j=1:nmax
    % Solve SLE Ax = v; x is then the next iterate
    y = L \ x = U \ x = x;
    v = v/norm(v);
    lambdamin = v'*A*v;
  end
endfunction
```