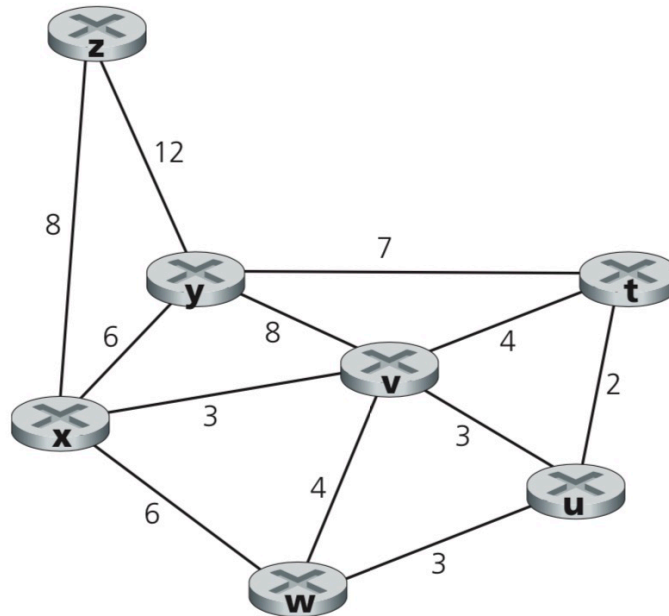


Exercise

P3

Consider the following network. With the indicated link costs, use Dijkstra's shortest-path algorithm to compute the shortest path from x to all network nodes. Show how the algorithm works by computing a table similar to Table 5.1.

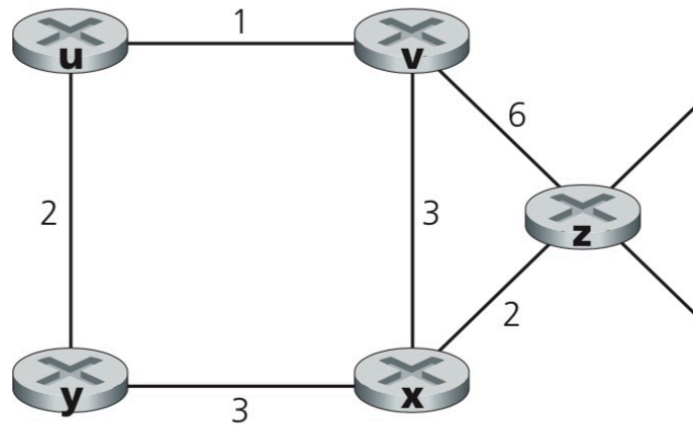


Answer

Step	N'	D(y),p(y)	D(z),p(z)	D(u),p(u)	D(v),p(v)	D(w),p(w)	D(t),p(t)
0	x	6,x	8,x	∞ ,x	3,x	6,x	∞ ,x
1	xv	6,x	8,x	6,v	3,x	6,x	7,v
2	xvy	6,x	8,x	6,v	3,x	6,x	7,v
3	xvyu	6,x	8,x	6,v	3,x	6,x	7,v
4	xvyuw	6,x	8,x	6,v	3,x	6,x	7,v
5	xvyuwt	6,x	8,x	6,v	3,x	6,x	7,v
6	xvyuwtz	6,x	8,x	6,v	3,x	6,x	7,v

P5

Consider the network shown below, and assume that each node initially knows the costs to each of its neighbors. Consider the distance-vector algorithm and show the distance table entries at node z.



Answer

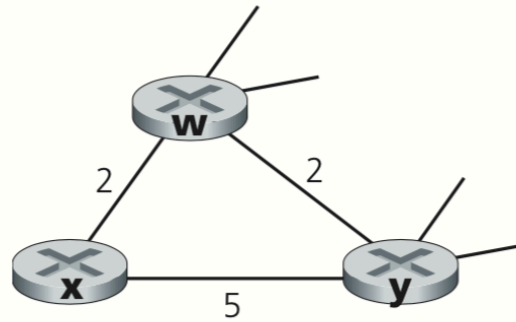
	u	v	x	y	z
v	∞	∞	∞	∞	∞
x	∞	∞	∞	∞	∞
z	∞	6	2	∞	0

	u	v	x	y	z
v	1	0	3	∞	6
x	∞	3	0	3	2
z	7	5	2	5	0

	u	v	x	y	z
v	1	0	3	3	5
x	4	3	0	3	2
z	6	5	2	5	0

P7

Consider the network fragment shown below. x has only two attached neighbors, w and y . w has a minimum-cost path to destination u (not shown) of 5, and y has a minimum-cost path to u of 6. The complete paths from w and y to u (and between w and y) are not shown. All link costs in the network have strictly positive integer values.



- Give x 's distance vector for destinations w , y , and u .
- Give a link-cost change for either $c(x,w)$ or $c(x,y)$ such that x will inform its neighbors of a new minimum-cost path to u as a result of executing the distance-vector algorithm.
- Give a link-cost change for either $c(x,w)$ or $c(x,y)$ such that x will *not* inform its neighbors of a new minimum-cost path to u as a result of executing the distance-vector algorithm.

Answer

a. $D_x(w)=2$, $D_x(y)=4$, $D_x(u)=7$

b. If $c(x,w)$ changes. If $c(x,w) > 6$, the least cost path will pass through y and cost 11. x will inform its neighbours. If $c(x,w) < 2$, the least cost path will be the same one but cost $5+c(x,w)$. x will inform its neighbours as well.

If $c(x,y)$ changes. If $c(x,y) < 1$, the least cost path will pass through y and cost $6+c(x,w)$. x will inform its neighbours.

c. If $c(x,y)$ changes. If $c(x,y) \geq 1$, x won't inform its neighbours.

If $c(x,w)$ changes. If $c(x,w) \geq 2$ and $c(x,w) < 6$, x won't inform its neighbours.

P12

Describe how to detect a loop in BGP.

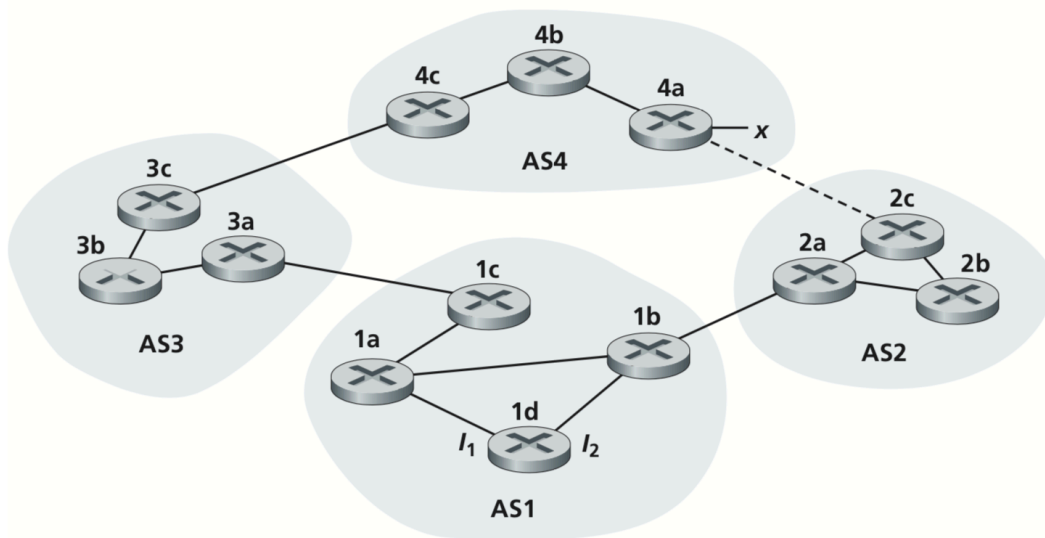
Answer

If a BGP peers receive a BGP message that the AS path contains its own AS number, then using that loop will result in a loop.

P14

Consider the network shown below. Suppose AS3 and AS2 are running OSPF for their intra-AS routing protocol. Suppose AS1 and AS4 are running RIP for their intra-AS routing protocol. Suppose eBGP and iBGP are used for the inter-AS routing protocol. Initially suppose there is *no* physical link between AS2 and AS4.

- Router 3c learns about prefix *x* from which routing protocol: OSPF, RIP, eBGP, or iBGP?
- Router 3a learns about *x* from which routing protocol?
- Router 1c learns about *x* from which routing protocol?
- Router 1d learns about *x* from which routing protocol?



Answer

- eBGP
- iBGP
- eBGP
- iBGP