

**Uge 46****Ex 1**

Which of the following formulas are satisfiable (give a satisfying assignment)? Which are not (give reasons)?

- a)  $A \wedge B$
- b)  $A \vee B$
- c)  $A \rightarrow B$
- d)  $A \wedge \neg A$
- e)  $A \vee \neg A$
- f)  $(A \rightarrow B) \wedge (B \rightarrow A)$
- g)  $(A \rightarrow B) \wedge (B \rightarrow A) \wedge A$
- h)  $(A \rightarrow B) \wedge (B \rightarrow A) \wedge \neg A$
- i)  $(A \rightarrow B) \wedge (B \rightarrow \neg A) \wedge (\neg A \rightarrow \neg B) \wedge (\neg B \rightarrow A)$

- a.  $A = 1, B = 1$
- b.  $A = 1, B = 0$
- c.  $A = 0, B = 1$  (if A then B, remember same as  $\neg A \vee B$ )
- d. Nope not satisfiable
- e. Yep, always true.  $A = 1$
- f.  $A = 1, B = 1$
- g.  $A = 1, B = 1$
- h.  $A = 0, B = 0$
- i. Nope, clearer when written out:

$$(\neg A \vee B) \wedge (\neg B \vee \neg A) \wedge (A \vee \neg B) \wedge (B \vee A)$$

led 2 og 4 indicates they cannot have same truth value.

led 1 og 3 indicates they cannot have different truth value.

No options left.

## Ex 2

Two formulas are equivalent, if the same assignments satisfy both of them.

Which of the following formulas are equivalent?

- a)  $\neg A \wedge B$
- b)  $\neg A \vee B$
- c)  $A \rightarrow B$
- d)  $(A \rightarrow B) \wedge (\neg B \rightarrow A)$
- e)  $(\neg A \rightarrow B) \wedge (\neg B \rightarrow \neg A)$

making truthtables:

$\neg A \wedge B$			$\neg A \vee B$			$A \rightarrow B$			$(A \rightarrow B) \wedge (\neg B \rightarrow A)$			$(\neg A \rightarrow B) \wedge (\neg B \rightarrow \neg A)$		
A	B	O	A	B	O	A	B	O	A	B	O	A	B	O
0	0	0	0	0	1	0	0	1	0	0	0	0	0	0
0	1	1	0	1	1	0	1	1	0	1	1	0	1	1
1	0	0	1	0	0	1	0	0	1	0	0	1	0	0
1	1	0	1	1	1	1	1	1	1	1	1	1	1	1

Also seen when writing out:

$$\neg A \wedge B$$

$$\neg A \vee B$$

$$\neg A \vee B$$

$$(\neg A \vee B) \wedge (B \vee A)$$

$$(A \vee B) \wedge (B \vee \neg A)$$

We find that the following pairs are equal:

- $b$  and  $c$
- $d$  and  $e$

**Ex 3**

We remember:

### Conversion to CNF

- Implications can be replaced by disjunction:
  - $A \rightarrow B$  converted to  $\neg A \vee B$
- DeMorgan's rules specify how to move negation "inwards":
  - $\neg(A \wedge B) = \neg A \vee \neg B$
  - $\neg(A \vee B) = \neg A \wedge \neg B$
- Double negations can be eliminated:
  - $\neg(\neg A) = A$
- Conjunction can be distributed over disjunction:
  - $A \vee (B \wedge C) = (A \vee B) \wedge (A \vee C)$

A and B are literals.

Convert the following formulas into CNF:

- a)  $\neg A \wedge B$
- b)  $\neg A \vee B$
- c)  $A \rightarrow B$
- d)  $(A \rightarrow B) \wedge (\neg B \rightarrow A)$
- e)  $(\neg A \rightarrow B) \wedge (\neg B \rightarrow \neg A)$
- f)  $A \rightarrow (\neg (B \wedge D))$
- g)  $A \rightarrow (\neg (B \vee D))$
- h)  $A \rightarrow (\neg (B \rightarrow (C \wedge D)))$

- a. is in CNF
- b. is in CNF
- c.  $\neg A \vee B$
- d.  $(\neg A \vee B) \wedge (B \vee A)$
- e.  $(A \vee B) \wedge (B \vee \neg A)$
- f.

$$\neg A \vee (\neg (B \wedge D))$$

$$\neg A \vee \neg B \vee \neg D$$

g.

$$\neg A \vee (\neg (B \vee D))$$

$$\neg A \vee (\neg B \wedge \neg D)$$

h.

$$(\neg A \vee \neg B) \wedge (\neg A \vee \neg D)$$

$$\neg A \vee (\neg(B \rightarrow (C \wedge D)))$$

$$\neg A \vee (\neg(\neg B \vee (C \wedge D)))$$

$$\neg A \vee ((B \wedge \neg(C \wedge D)))$$

$$\neg A \vee (B \wedge (\neg C \vee \neg D))$$

$$\neg A \vee ((B \wedge \neg C) \vee (B \wedge \neg D))$$

$$(\neg A \vee (B \wedge \neg C)) \wedge (\neg A \vee (B \wedge \neg D))$$

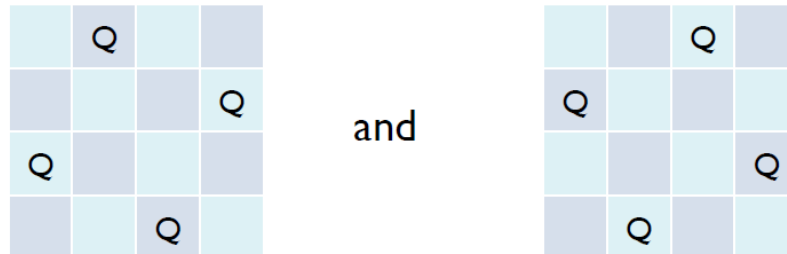
$$(\neg A \vee (B \wedge \neg C)) \wedge (\neg A \vee (B \wedge \neg D))$$

$$(\neg A \vee B) \wedge (\neg A \vee \neg C) \wedge (\neg A \vee B) \wedge (\neg A \vee \neg D)$$

$$(\neg A \vee \neg C) \wedge (\neg A \vee B) \wedge (\neg A \vee \neg D)$$

**Ex 4**

Solutions to N-Towers and N-Queens are symmetric:



- a) Write two clauses that forbid solutions where there is a queen in the right half of the first row.
- b) Instead of adding two clauses, change an existing clause.

$x_{1,1}$	$x_{1,2}$	$x_{1,3}$	$x_{1,4}$
$x_{2,1}$	$x_{2,2}$	$x_{2,3}$	$x_{2,4}$
$x_{3,1}$	$x_{3,2}$	$x_{3,3}$	$x_{3,4}$
$x_{4,1}$	$x_{4,2}$	$x_{4,3}$	$x_{4,4}$

a. clauses:

$$\neg x_{1,3}$$

$$\neg x_{1,4}$$

b. changing where the tower can be:

from:

$$x_{1,1} \vee x_{1,2} \vee x_{1,3} \vee x_{1,4}$$

to:

$$x_{1,1} \vee x_{1,2}$$

**Ex 5**

- Install **lingeling** or another compatible SAT solver
- Alternatively, use a Javascript SAT solver, e.g.:
  - <https://www.msoos.org/2013/09/minisat-in-your-browser/>
- Test it using the following input saved as `test.cnf`

```
p cnf 4 6
-1 -2 0
-1 -3 0
-2 -4 0
-3 -4 0
1 2 0
3 4 0
```

**Ex 6**

The formula from Slide 11 contains redundant information. For example,  $X_{1,1} \rightarrow \neg X_{1,2}$  and  $X_{1,2} \rightarrow \neg X_{1,1}$  are equivalent.

Understand and remove these redundancies:

- Why do these redundancies occur?
- Identify all such redundancies!
- Write down a simplified formula without redundancies!
- Convert the simplified formula into CNF!
- Write the formula in DIMACS format!
- Run the lingeling solver on it and interpret the result!

$X_{1,1} \rightarrow \neg X_{1,2}$	"Tower at (1,1) attacks to the right"
$X_{1,1} \rightarrow \neg X_{2,1}$	"Tower at (1,1) attacks downwards"
$X_{1,2} \rightarrow \neg X_{1,1}$	"Tower at (1,2) attacks to the left"
$X_{1,2} \rightarrow \neg X_{2,2}$	"Tower at (1,2) attacks downwards"
$X_{2,1} \rightarrow \neg X_{2,2}$	"Tower at (2,1) attacks to the right"
$X_{2,1} \rightarrow \neg X_{1,1}$	"Tower at (2,1) attacks upwards"
$X_{2,2} \rightarrow \neg X_{1,2}$	"Tower at (2,2) attacks to the left"
$X_{2,2} \rightarrow \neg X_{2,1}$	"Tower at (2,2) attacks upwards"
$X_{1,1} \vee X_{1,2}$	"Tower in first row"
$X_{2,1} \vee X_{2,2}$	"Tower in second row"

- Because of symmetry
- Redundant pairs:
  - 1 and 3
  - 2 and 6
  - 4 and 7
  - 5 and 8
- without redundancies:
  - $X_{1,1} \rightarrow \neg X_{2,1}$
  - $X_{1,2} \rightarrow \neg X_{1,1}$
  - $X_{1,2} \rightarrow \neg X_{2,2}$
  - $X_{2,1} \rightarrow \neg X_{2,2}$
  - $X_{1,1} \vee X_{1,2}$
  - $X_{2,1} \vee X_{1,2}$

d. Convert into CNF

$$\begin{aligned} & (X_{1,1} \rightarrow \neg X_{2,1}) \wedge (X_{1,2} \rightarrow \neg X_{1,1}) \wedge (X_{1,2} \rightarrow \neg X_{2,2}) \wedge (X_{2,1} \rightarrow \neg X_{2,2}) \\ & \quad \wedge (X_{1,1} \vee X_{1,2}) \wedge (X_{2,1} \vee X_{1,2}) \\ & (\neg X_{1,1} \vee \neg X_{2,1}) \wedge (\neg X_{1,2} \vee \neg X_{1,1}) \wedge (\neg X_{1,2} \vee \neg X_{2,2}) \wedge (\neg X_{2,1} \vee \neg X_{2,2}) \\ & \quad \wedge (X_{1,1} \vee X_{1,2}) \wedge (X_{2,1} \vee X_{2,2}) \end{aligned}$$

e. Write in DIMACS:

$X_{1,1}$	$X_{1,2}$
$X_{2,1}$	$X_{2,2}$

```
p cnf 4 8
-1 -3 0
-2 -1 0
-2 -4 0
-3 -4 0
1 2 0
3 4 0
```

f. Run and interpret:

```
Bash on Ubuntu on Windows
c SATISFIABLE
v -1 2 3 -4 0
c 0.000 0% simplifying
c 0.000 0% search
c =====
c 0.000 100% all
c
c Run and interpret:
c 0 conflicts, 0.0 confs/sec
c 0 ternaries, 0.0 confs/ternary
c 0 binaries, 0.0 confs/binary
c 0 iterations, 0.0 confs/iteration
c
c 0 reductions, 0.0 redus/sec, 0.0 confs/reduction
c 0 restarts, 0.0 rests/sec, 0.0 confs/restart
c 0 decisions, 0.0 decis/sec, 0.0 decis/conflict
c 0 propagations, 0.0 props/sec, 0.0 props/decision
c 0.0 seconds, 0.0 MB
```