Question 1: Functional Dependencies I [15 points]

Consider the following legal instance of a relational schema S with attributes \widetilde{XYZ} :

5	X	Y	Z
	m	20	T
	m	10	F
	О	30	T
	\overline{n}	30	T
	О	20	T

Table 1: Legal instance of schema S for question 2.1

(a) Which of the following dependencies are $violated$ by the instances of S in Table 1?						
i. [2 points] \boxtimes Yes \square No : $X \to Y$ is violated.						
ii. [2 points] \boxtimes Yes \square No : $Z \to X$ is violated.						
iii. [2 points] \square Yes \bowtie No : $Y \to Z$ is violated.						
iv. [2 points] \square Yes \boxtimes No : $XY \rightarrow Z$ is violated.						
v. [2 points] \boxtimes Yes \square No : $YZ \to X$ is violated.						
vi. [2 points] \boxtimes Yes \square No : $XZ \to Y$ is violated.						
(b) [3 points] By only observing the instance of S in Table 1, can you identify the functional						
dependencies that hold on schema S? Why?						
□ Yes ☑ No						
Because Table 1 is just a legal instance of schema S.						
There may be counterexamples in other instances of schema S.						

Question 2: Functional Dependencies II For the next set of questions consider the relational schema $\mathcal{R} = \{P, Q, R\}$ the set of functional dependencies FD:	
$Q \rightarrow U$	(1)
$egin{array}{ccc} oldsymbol{arphi} & oldsymbol{arphi} & oldsymbol{arphi} & oldsymbol{V} & oldsymbol{arphi} & oldsymbol{V} & oldsymbol{arphi} &$	(2)
$PQ \rightarrow WST$	(3)
$SU \rightarrow TR$	(4)
$VT \rightarrow RW$	(5)
$R \rightarrow W$	(6)
(a) [8 points] Which of the following is a minimum cover of the FD? if none, mark accordingly, and give your <i>own</i>. answer.i. The given FDs (Eq 1-6), is a minimum cover already.	Mark all that qualify;
ii. $\{Q \rightarrow U, U \rightarrow V, PQ \rightarrow S, SU \rightarrow T, SU \rightarrow R, VT \rightarrow R, VT\}$	$T \to W R \to W$
iii. $\{Q \rightarrow U, U \rightarrow V, PQ \rightarrow S, SU \rightarrow T, PQ \rightarrow W, VT \rightarrow R, PQ \rightarrow W, $	
iv. $\{Q \to U, U \to V, PQ \to S, SU \to T, VT \to R, R \to W\}$	*
v. $\{Q \rightarrow U, U \rightarrow V, PQ \rightarrow S, SU \rightarrow T, SU \rightarrow R, VT \rightarrow R, PQ\}$	$Q \to T, R \to W$
vi. none of the above - the cover is	
 (b) Yes/No: Which of the following functional dependencies can be ded set of functional dependencies (Eq. (1)-(6))? i. [3 points]	
(c) [3 points] True or False: The attribute closure $\{Q\}^+$ is $\{Q, U, V\}$.	
☐ True ☐ False	
(d) [3 points] True or False: The attribute closure $\{PQ\}^+$ is $\{P,Q,W\}$ \square True \bowtie False	$\{S,T\}.$

Question 3: Decompositions......[20 points]

For this set of questions, consider the relation with attributes, $\mathcal{X} = \{A, B, C, D, E, F\}$, Let the following functional dependencies FD be defined over the relation \mathcal{X} :

$$A \to B$$
$$B \to CD$$
$$E \to F$$

- (a) [2 points] Provide the attribute closure of $\{AB\}$. $\{AB\}^+ = \{ABCD\}$
- (b) Consider the decomposition AB, BCD, EF. Mark 'True' or 'False':
 - i. [3 points] □ True ☑ False: It is lossless
 - ii. [3 points] ☑ True □ False: It is dependency-preserving
- (c) Consider the decomposition AB, BCDF, EF. Mark 'True' or 'False':
 - i. [3 points] □ True ☑ False: It is lossless
 - ii. [3 points]

 ☐ True ☐ False: It is dependency-preserving
- (d) Consider the decomposition ABCEF, EBD. Mark 'True' or 'False':
 - i. [3 points] ☑ True ☐ False: It is lossless
 - ii. [3 points]

 ✓ True

 □ False: It is dependency-preserving

	ributes, $\mathcal{E} = \{P, Q, R, S\}$. Suppose				
	$PQ \rightarrow R$	(7)			
	$PQ \rightarrow S$	(8)			
	$R \rightarrow P$	(9)			
	$S \rightarrow Q$	(10)			
(a) [6 points] List <i>all</i> the c	candidate key(s) for \mathcal{E} . $\{PQ\}$	QR) {RS} {PS}			
(b) [2 points] Is the relation	[6 points] List all the candidate key(s) for \mathcal{E} . $\{PQ\}$, $\{QR\}$, $\{QS\}$, $\{PS\}$ [2 points] Is the relation \mathcal{E} in BCNF? \square Yes \bowtie No				
) From the list below, select all applicable choices to justify whether \mathcal{E} is (or is not) in BCNF.					
Note : when we refer to a super key.	the main requirement for BCNF, we	e mean: every determinant is			
i. [1 point] □ True	False: All FD's satisfy the ma	-			
ii. [1 point] \square True	☐ False: FD (7) violates the ma	•			
iii. [1 point] □ True	☐ False: FD (8) violates the ma	in requirement.			
iv. [1 point] True	□ False: FD (9) violates the ma	•			
v. [1 point] 🗹 True	□ False: FD (10) violates the m	ain requirement.			
(d) [2 points] Is the relation	on \mathcal{E} in 3NF? \boxtimes Yes \square No				
(e) From the list below, sele	ect all applicable choices to justify w	hether \mathcal{E} is (or is not) in 3NF.			
Note: when we refer to $X \to A$, A is part of a c	o the secondary requirement for 31 candidate key.	NF, we mean: for every FD			
i. [1 point] 🗹 True	☐ False: All FD's satisfy the sec	condary requirement.			
ii. [1 point] □ True	☐ False: FD (7) violates the sec				
iii. [1 point] 🗆 True	False: FD (8) violates the sec	ondary requirement.			
iv. [1 point] □ True	False: FD (9) violates the sec	ondary requirement.			
v. [1 point] True	☐ False: FD (10) violates the se	condary requirement.			
	decomposition of \mathcal{E} that is lossless, sible. $R_1 = \{P, Q, R, S\}$	dependency preserving, and			
(g) [8 points] Give a BCN	IF decomposition of $\mathcal E$ that is lossle	ess, and has as few tables as			
possible.	R1= {P.R}				
	$R_2 = \{Q,S\}$				
	$R_3 = \{R, S\}$				
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