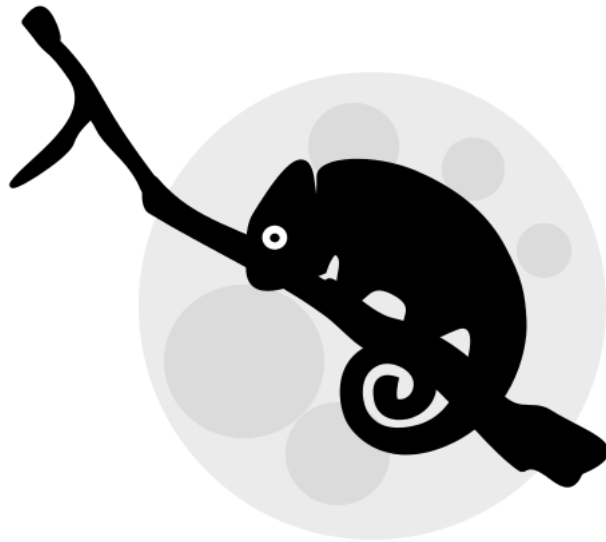


# The book of Revelation

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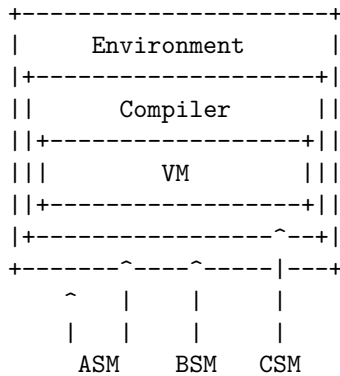
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# 1 ASVM

## 1.1 Architecture

### 1.1.1 Environment

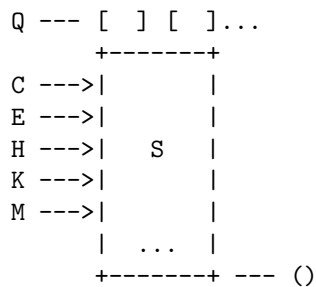


Source types:

- .asm - high level ASM.
- .bsm - human-readable barebones ASM.
- .csm - raw, compiled bytecode

### 1.1.2 Compiler

### 1.1.3 VM



Registers & values:

- Q - QValue register
- C - code pointer
- E - environment stack pointer
- H - handler stack pointer
- K - continuation stack pointer
- M - metacontinuation stack pointer
- S - store pointer (memory start)
- () - fnord value (memory end)

## 1.2 Memory model

### 1.2.1 Layout

- Tagged memory

```
15 [--|---|---|---|-----] 0
    ^   ^   ^   ^   ^
    |   |   |   |   |_ 8 type/operator bits
    |   |   |   |   |_ 1 immutability bit
    |   |   |   |   |_ 2 cons packing bits (optional)
    |   |   |   |   |_ 3 unused bits (reserved for future use)
    |_ 2 GC bits (may require more)
```

- TValue

Layout:

```
x86    - [-2-|---4--]      ---> 6
         [-2-|x2x|---4--] ---> 8, word-aligned
amd64  - [-2-|----8----] ---> 10
         [-2-|---6---]     ---> 8, word-aligned, memory-magic
```

Variants:

```
[TAG|--PTR--] ---> anything that uses pointers
[TAG|--VAL--] ---> anything that uses values
```

Cons:

- No big numbers without any special treatment.
- Interfacing arrays and native calls to D will be complicated.
- Tag might need additional cons-packing bits.
- Tag appears twice in a pair.

Pros:

- Doesn't need as much padding.
- No dangling null pointers.
- Cons packing is trivial.
- Allows for many different kinds of objects to be implemented.
- Allows other kinds of languages to be implemented.
- Maps to Lisp well (was used in Lisp machines).

- TBlobs

Consist of several consecutive QValues.

Pairs, triples, etc:

```
[pair|--PTR--] [fnord|--PTR--]
[triple|--PTR--] [fnord|--PTR--] [fnord|--PTR--]
```

Cons packing:

```
[00|--VAL--]          ---> there's no cdr (errors on read/write)
[01|--VAL--]          ---> car is (at) val, cdr is null
[10|--VAL1-][TAG|--VAL2-] ---> car is (at) val1, cdr is at val2
[11|--VAL1-][TAG|--VAL2-] ---> car is (at) val1, cdr is val2
^--- Cons encoding part of the tag field.
```

- Tuples and Lists

Primitives:

- %car - first pointer of a pair
- %cdr - second pointer of a pair

- Arrays and Vectors

Primitives:

- %ref - array/vector pointer and offset pair
- %slice - array/vector pointer, start and end pointer triple

- Compound types

Used to implement sealer/unsealer pattern.

```
# Could be a hash.
(var *compound-type* 0)

# Could use a separate Q type (Type?) and make use of unique references and is? predicate.
(function make-type ()
  (do (var t *compound-type*)
    (set! *compound-type* (+ 1 *compound-type*))
    (tuple t
      (lambda (o)
        (cons t o))
      (lambda (o)
        (if (and (tuple? o)
                  (equal? (car o) t))
            (cdr o)
            (error "Type mismatch."))))))
(function typeof (o)
  (when (pair? o)
    (car o)))

(var (T sealT unsealT) (make-type))

(var foo (sealT (tuple 1 2 3)))

# Might facilitate predicate-based type pattern matching.
(function baz (v)
  (case (typeof v)
    (T (unsealT v))
    (X (unsealX v))
    ...))
```

### 1.2.2 Allocator

### 1.2.3 Garbage collection

- GC bits

```
[00|--VAL--] ---> unmanaged (pinned)
[01|--VAL--] ---> undecided
[10|--VAL--] ---> undecided
[11|--VAL--] ---> undecided
```

### 1.2.4 OpCode encoding

Always pairs → type part of the tag can be used as the operator type.

Example:

(%foo bar baz) is...

```
[%foo|-car-|-cdr-]---->[tag|baz]
      |
      v
    [tag|bar]
```

...instead of...

```
[pair|-car-|-cdr-]---->[pair|-car-|-cdr-]---->[tag|baz]
      |                               |
      v                               v
    [tag|%foo]                     [tag|bar]
```

## 1.3 Threading

### 1.3.1 Actor model

- Threading
  - %tid - returns current threads ID.
  - %spawn - spawns a thread evaluating given bytecode.
  - %send - sends a bunch of immutable data to a thread.
  - %receive - receives a bunch of data.

### 1.3.2 Processes

### 1.3.3 Message passing

## 1.4 Combinators

### 1.4.1 Vau calculus

Basics:

```
((vau x e x) foo)      ----> foo
((vau x e e) foo)      ----> dynamic environment
((vau (x y z) e z) 1 2 3) ----> 3
```

Implementing lambda:

```
lambda => (vau (args body) env
              (wrap (eval (tuple 'vau args () body)
                             env)))
```

Primitives

- %vau - creates a lexically scoped operative combinator taking dynamic environment.
- %wrap - induces arg evaluation allowing for applicative combinators.

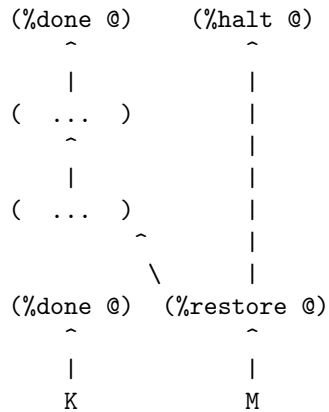
### 1.4.2 Argument evaluation

## 1.5 Formal operational semantics

### 1.5.1 Environments

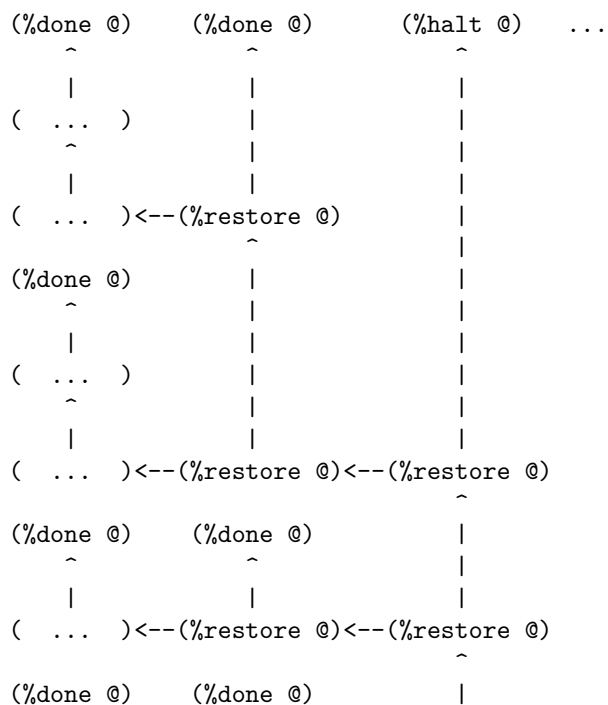
### 1.5.2 Continuations

- @ register  
Instead of value stores the return address where the value should be stored.
- Metacontinuations  
Additional M stack containing continuation segments.

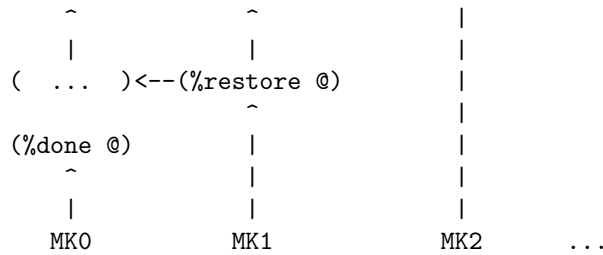


Primitives:

- %restore - sets K to the stored continuation stack segment and applies it to the continuation hole.
- %done - ends the current continuation segment and invokes the M register.
- %halt - ends the flow of the program.
- Generalized metacontinuations  
Multiple metacontinuation stacks with multiple segments each.







Primitives:

- %restore - pushes a stored continuation stack segment onto the MK register.
- %done - pops the MK register leaving the rest of the meta-stack.
- %halt - ends the flow of the program.

Possible primitives:

- %done-if - premature MK register popping (if @ != ()).
- %select - depending on @ pushes one of its children onto the MK stack.

Notes:

- Might be really cool. Especially because it doesn't require constant consing of the continuation stack.
- All the code can be pre-transformed into dataflow format and then executed with no further transformations.
- Used to implement delimited continuations.

### 1.5.3 Error handling

## 1.6 Interfacing with D

### 1.6.1 Native calls

Implementation:

```
[delegate|ptr|func]---> raw D function pointer
|
v
memory location containing the closure
```

```
[native|func]----> raw D function pointer
[nativec|func]---> raw C function pointer
```

Example usage:

```
ASM.define("foo", x => x);
ASM.foo = x => x;
```

### 1.6.2 Native types

Implementation:

```
[user-type|type|data]---> raw D data
|
v
D typeid
```

Example usage:

```

struct Test {
    int bar;
    string foo;
}

// ...

ASM.defineType!Test;

ASM.define("foobar", (scpe, args) {
    if(args.car.type == Type.UserDefined)
    if(args.car.userType == typeid(Type))
        // Do shit
    return ASM.fnord;
});

ASM.doString(q{
    (var baz (scope
        (var _inner (newTest))

        (function getFoo ()
            (getTestFoo _inner))
        (function setFoo (newVal)
            (setTestFoo _inner newVal))

        (function getBar ()
            (getTestBar _inner))
        (function setBar (newVal)
            (setTestBar _inner newVal))))
    ((baz setFoo) "Test")
    (foobar baz)
    (writeln (baz getBar))
});

```

### 1.6.3 Dynamic FFI

### 1.6.4 Loading any ASM version

## 2 ASM programming language

### 2.1 Phases of evaluation

#### 2.1.1 Lexical analysis

#### 2.1.2 Static analysis

#### 2.1.3 Code generation

#### 2.1.4 Optimisation

#### 2.1.5 Evaluation

### 2.2 Lexical

#### 2.2.1 Comments

- Metadata
- Opts
- Expression comments
- Shebang parameters problem
- Multiline comments

#### 2.2.2 Numbers

#### 2.2.3 Symbols

#### 2.2.4 Identifiers

#### 2.2.5 Tuples

#### 2.2.6 Vectors

#### 2.2.7 Strings

#### 2.2.8 Reserved keywords & special tokens

### 2.3 Semantics

#### 2.3.1 Immutability

#### 2.3.2 Atoms

- Numbers
- Symbols
- Booleans
- Unit Type

#### 2.3.3 Combinators

- Applicative combinators

```
(function (foo bar baz)
  (* bar baz))
```

- Named call parameters syntax

```
(foo (bar . 23) (baz . 5))
(foo bar=23 baz=5)
(foo --bar 23 --baz 5)
```

- Operative combinators  
According to Vau calculus as defined by John Shutt.

```
(var lambda (%vau (args body) env
  (%wrap (eval `($%vau $args ignored $body)
    env))))
(var wrap (%lambda (combinator)
  (%lambda args
    (eval `($combinator $@args)))))
```

### 2.3.4 Ranges

As defined by Andrei Alexandrescu in On iteration.

Input:

```
      InputRange ---> front, empty?, popFront!
      ^
      |
ForwardRange ---> save (deep copy)
      ^      ^
      /        \
BidirectionalRange -+--> back, popBack!
                    \
                    InfiniteRandomAccessRange ---> [] (indexing)
                    ^
                    |
                    FiniteRandomAccessRange ---> [] (indexing)
```

Output:

```
OutputRange ---> put
```

Output ranges could be dropped in favour of impure functions -

```
(put 23) ---> (foo 23 'bar '(1 2 3)).
```

- With macro

```
# Could be generalized to all scopes.
(with someRange      (do
  (pop!)              ==> ((someRange pop!))
  (put! 'foo))        ((someRange put!) 'foo))
```

- Example

```
(function circular (tpl)
  (scope (var offset tpl)
    (function empty? ()
      '())
    (function front ()
      (car offset))
    (function popFront! ()
      (set! offset (cdr offset))
      (when (not offset)
        (set! offset tpl)))))

(var foo (circular '(1 2 3)))
```

Together with type tagging/boxing/sealing might prove to be quite nice:

```
(function chained (r1 r2)
  (seal 'range
    (scope ...
      ...)))
```

- Tuples
- Vectors  
Vectores are heterogenous.

```
[vec|stored-type] [int|length] [ptr|data]
```

Vec is the vector tag. Stored-type contains ASM typeinfo of the types of stored elements or None for an empty vector or Any for a vector containing different typed-values. Length contains the length of the data in the vector. Ptr is the data pointer.

Pros:

- merges vectors and arrays
- O(1) best case typeinfo
- the type info will be used by D code facilitating data-packing

Cons:

- implies immutability
- 2-3 word header in addition to the actual data
- Vectors of multi-word values should be un-allowed (using vectors of pointers instead)

Optimisations: Data-pack same-typed-data vectors of basic types not to wrap them with an ASM typeinfo (requiring simple-type packing, unpacking, vector copy and vector slice).

Store vec and stored-type in length making it 32 bit long instead of 48 bits:

```
[vec|stored-type|length] [ptr|data]
```

Pros:

- 1-2 word header

Cons:

- 32 bits for length

Use additional cons-packing value to indicate a vectorized type (followed by length and data pointer).

```
[int|1] <-- int
[int*|4] [ptr|--->] ... \footnote{DEFINITION NOT FOUND: 1 }\textsuperscript{,}\,\footnote{DEFINITION NOT FOUND: 1 }
```

Pros:

- 1-2 word header
- 48 bit length
- cons-packed tuples will easily vectorize

Cons:



## 3 Code Cube

### 3.1 cs

#### 3.1.1 core

- Type predicates
- Type conversions
- Basic math
- import  
Imports symbols, loads modules, manages scopes:  
(import func from ModuleA as AFunc all from ModuleB)

- let  
Immutable let and mutable var:

```
(let ((foo bar))  
  # foo is immutable  
)
```

```
(var ((foo bar))  
  # foo is mutable  
)
```

- module/program/class/application etc  
Wraps a bunch of functions and state into a single, named unit:

```
(module Foo  
  private (function (foo bar baz)  
    (bar baz baz))  
  public (var bar)  
  export (function (herp derp)  
    (derp derp derp)))
```

Dependency injection:

```
(module Math (printer allocator)  
  (function (matrix m n)  
    (allocator.malloc (* bar baz)))  
  
  (function (printm matrix)  
    (foreach e in matrix  
      do (printer.print e))))  
  
(import (Math my-logger kewl-allocator))  
  
(Math.printm (Math.matrix 3 3))
```

- case/switch/match/type-dispatch  
Switch-like control structure, with fallthrough + case goto, case ranges etc:

```
(switch a  
  case b (foo bar baz)  
  case c (faz baz baz)  
  default foo)
```

### 3.1.2 memory

- GC
  - `collect!` - does a collection.
  - `minimise!` - minimises memory use.
  - `disable!` - stops GC.
  - `enable!` - resumes GCs work.
- Allocator

### 3.1.3 thread

### 3.1.4 error

- `(handle e handler)` - handler = (error-object handling-function)
- `(raise error-object)`
- `warn` - runtime warning
- `assert` - check condition and rise errors

### 3.1.5 reader

Based on dynamic PEG parser generator, because it doesn't need separate lexing phase. Reader macros will be grammar based.

Implementation:

```
(syntax (grammar-declarator)
  transform
  ...)

(grammar ((grammar-declarator) transform)
  ...)
```

grammar-declarator:

Name arrow Rules

transform:

Any code, really.

Name:

Rule name - used inside of it for transforms and outside for parsing.

Rules - implicitly wrap in a sequence:

- `(a b c ...)` - sequence
- `(/ a b c ...)` - ordered choice
- `(* Rules)` - zero or more repeats of the Rules
- `(+ Rules)` - one or more repeats of the Rules
- `(? Rules)` - optional Rules
- `(! Rules)` - not Rules
- `(& Rules)` - and Rules
- `(: Rules)` - consumes input and drops captures



- (`~` Rules) - concatenates captures

arrow:

- `<-` - basic
- `<` - spacing consuming
- `<~` - concatenative

### 3.1.6 writer

Using pattern matching and string embeds, possibly sewn together with the reader. Might be of use for the bytecode/crosscode compiler.

## 3.2 cc

### 3.2.1 docs

Used for documenting code, using... code in the comments. Something along these lines (needs more work):

```
#? (ASMDoc
#?   This function does some stuff and returns other stuff.
#?   --params
#?       bar - an integer,
#?   --returns - another integer,
#?   --example
#?       (var baz (foo 23))
#? )
(function foo (bar)
  (doStuff bar))
```

### 3.2.2 test

Automated unittest runner:

```
(unittest Foo
  assert (equal? bar baz)
  assert (foo bar baz)
  test Bar
    assert (foo bar baz)
    assert (foo bar baz)
    log "herp derp"
  test Baz
    assert (bar foo faz)
  finally (derp herp))
```

### 3.2.3 dbc

function macro - creates a function with all kinds of cool stuff:

```
(function (foo bar baz)
  in (equal? bar 23)
  in (> baz bar)
  out (< result bar)
  body (bar baz))
```

- `□` enforce - makes sure an operation will succeed.

### 3.2.4 ranges

- Collection manipulation

- $\boxminus$  join - if the second argument is a collection - prepends it the the first argument, if it's not a collection - joins both arguments into a pair. Creates a new collection. Examples:
  - \* (join 1 '[1 2 3]) -> [1 1 2 3]
  - \* (join '(a b) '[1 2 3]) -> [(a b) 1 2 3]
  - \* (join 'a 'b) -> (a b)
- $\boxplus$  append if argument types match - appends element or a collection to another collection, if types don't match - appends the second argument to the collection. Creates a new collection. Examples:
  - \* (append '[1 2 3] 4) -> [1 2 3 4]
  - \* (append '[1 2 3] '[4 5 6]) -> [1 2 3 4 5 6]
  - \* (append '(2 3) '[2 3]) -> (2 3 [2 3])
- $\boxtimes$  first - returns a reference to the first element of a mutable collection, or its value for an immutable collection.
- $\boxtimes$  rest - returns a new collection referencing the rest part of the old one.
- $\boxtimes$  second, third, fourth etc.
- $\boxtimes$  nth - returns nth element of a collection.
- $\boxtimes$  map - maps an operation to a collection collecting results.
- $\boxtimes$  reduce - maps an operation to a collection reducing it to a single value.
- $\square$  ? slice - slices a collection creating subcollection.
- $\boxminus$  ? push, push-back, pop, pop-back.
- ? etc

- Collection creation

- $\boxtimes$  list - returns a list consisting of the call args.
- $\boxtimes$  tuple - returns a tuple consisting of the call args.
- $\boxtimes$  set - returns a set consisting of the call args.
- $\boxtimes$  scope - reuturns a scope with call args defined in it.
- ? etc.

- APL-esque array processing

- + monadic conjugate, dyadic plus
- - monadic negate, dyadic minus
- $\div$  div m reciprocal, d divide
- $\times$  mul m sign of, d multiply
- upstil - m ceiling, d maximum
- downstil - m floor, d minimum
- \* exp m exponential, d power
- ! bang m factorial, d binomial
- | stile m magnitude, d residue
- log m natural logarithm, d logarithm
- circle m times pi, d circular function
- domino m matrix inverse of, d matrix division

- decode d decode
- encode, d encode
- ? roll m roll, d deal
- land d lang/lowest common multiple
- lor d logical or/greatest common divisor
- nand d nand
- nor d nor
- < less d less than
- > greater d greater than
- leq d less or equal
- geq d greater or equal
- = eq d equal to
- neq d not qual to
- equnderbar m depth of, d matches
- nequnderbar d not match
- rho m shape of, d reshape
- , comma m ravel, d cotenate/laminate
- commabar m columnize, d catenate along first axis
- circle stile m reverse, d rotate
- circle bar m reverse along first axis, d rotate along first axis
- transpose m transpose, d general transpose
- ↑ up m mix, d take
- ↓ down m split, d drop
- left shoe m enclose, d partitioned enclose
- epsilon m enlist, d member of
- index m array/default, d index
- right shoe m first, d pick
- / slash d replicate, o reduce
- slashbar d replicate along first axis, o reduce along first axis
- \ slope d expand, o scan
- slopebar d expand along first axis, o scan along first axis
- ~ tilde m not, d without
- union m unique, d union
- intersection d intersection
- left tack m same, d left
- right takc m same, d right
- iota m index generator, d index of
- epsilon underbar d find
- grade up m grade up, d collated grade up
- grade down m gnade down, d collated grade down
- ^ high minus - same as minus
- ‘ quote - string delimiter
- ← left d assignment
- zilde niladic - empty numeric vector (same as (iota 0))
- thorn m format, d format by specification

- diamond - statement separator
- comment - comment
- del - self reference
- alpha - left argument of a dyadic function
- omega - right argument of a dyadic function
- quad - system name prefix
- " dieresis o each
- dieresis tilde o commute
- dierosis star o power
- . dot o inner product ((dot jot) produces outer product)
- jot o compose ((dot jot) produces outer product)

- VLists

O(log n) indexing. If offset is 0, vlist preallocates additional chunk of data.

```
[offset|list]--->[vector|next]--->[vector|next]--->[vector|next]--->()
|               |               |               |
|               |               |               v
|               |               v               [length|data]--->|0| | |
|               v               [length|data]--->|1|2|
|               [length|data]--->|3|4|5|6|
|-----^-----
```

- vectorize

```
(macro vectorize (fun)
  (let ((old-fun (gensym)))
    `(let (($old-fun $fun))
      (function $fun (vec)
        (map fun vec))))))
```

- loop

Common Lisp like loop macro:

```
(loop for foo in bar
      for baz being each hash-key of goo
      when gaz
      do gar)
```

- for

```
(for x <- foo
  y <- x
  if (> y 23)
  yield y)
```

### 3.2.5 io

- Input

- □ readln - Unformatted (string) reads.
- □ read - Formatted reads.
- □ ? load/open - Loads a file for reading (as a Scope/Stream with read defined accordingly).
- □ ? close - closes an input stream.

- `eof?` - returns 'yup/the object if it has reached EOF.
- `? etc`
- Output

- `write` - writes string representation of the args.
- `? etc`

### 3.2.6 math

- `sqrt`
- `pow`
- `exp`
- `min/max/clamp`
- `etc`

### 3.2.7 random

- `Marsane Twister`
- `Gaussian distribution`

### 3.2.8 object

- `opElvis syntax`

```
foo ? bar == (if-non-fnord foo bar) == (if foo foo bar)
```

Implementation:

```
(macro if-non-fnord (foo bar)
  (let ((_foo (gensym)))
    `(let (($_foo $foo))
      (if $_foo
          $_foo
          $bar))))

(syntax "\?" (parsed parsing)
  `(if-non-fnord $(pop-front! parsed)
                  $(read-expression! parsing)))
```

Example:

```
# map - 1d or 2d list
(let ((leny (length map))
      (lenx (length (car map)) ? 1))
  (do-stuff lenx leny) ? 23)

(let ((leny (length map))
      (lenx (if-non-fnord (length (car map))
                           1)))
  (if-non-fnord (do-stuff lenx leny)
                 23))
```

```
(let ((leny (length map))
      (lenx (let ((__GENSYM0 (length (car map))))
                (if __GENSYM0
                    __GENSYM0
                    1)))))
  (let ((__GENSYM1 (do-stuff lenx leny)))
    (if __GENSYM1
        __GENSYM1
        23)))
```

- Dynamic dispatch

```
(defmethod foo (bar baz) body) (foo bar baz) <=> ((get bar 'foo) baz) # Dynamic dispatch
```

or

```
(with bar (foo baz)) # With macro
```

or

```
(function foo (baz) body) (connect bar foo) # Slots
```

### 3.2.9 babel

- JSON

- XML

- SVG

Returns a wellformed SVG string:

```
(SVG 100 100
  (circle 50 50
    '(255 255 100)))
```

- YAML

- $\text{\LaTeX}$

Returns a wellformed  $\text{\LaTeX}$  string:

```
(LaTeX
  "The following equation is herp derp derp:"
  (equation "a^2 + b^2 = c^2")
  (equation "\herp = \derp"))
```

- dot

– ASM AST/module dependancy -> graphviz utility.

- iexpr

```
(package foo
  (function (bar arg0 arg1)
    (if (and (atom? arg0)
              (atom? arg1))
```

```

      (* arg0 arg 1)
      (apply + (append arg0 arg1))))
(var gun (bar 2 3))

```

```

||
\\|/
\ /

```

```

package foo
function (bar arg0 arg1)
  if and atom? arg0
    atom? arg1
      (* arg0 arg1)
      apply +
        (append arg0 arg1)
var gun
  (bar 1 2)

```

## 4 References