

Substitution Method

The substitution method

1. Guess a solution
2. Use induction to prove that the solution works

Substitution method

- Guess a solution
 - $T(n) = O(g(n))$
 - Induction goal: **apply the definition of the asymptotic notation**
 - $T(n) \leq d g(n)$, for some $d > 0$ and $n \geq n_0$ (strong induction)
 - Induction hypothesis: $T(k) \leq d g(k)$ for all $k < n$
- Prove the induction goal
 - Use the **induction hypothesis** to **find some values of the constants d and n_0** for which the **induction goal** holds

Example 1

$$T(n) = c + T(n/2)$$

- Guess: $T(n) = O(\lg n)$
 - Induction goal: $T(n) \leq d \lg n$, for some d and $n \geq n_0$
 - Induction hypothesis: $T(n/2) \leq d \lg(n/2)$
- Proof of induction goal:

$$\begin{aligned} T(n) &= T(n/2) + c \leq d \lg(n/2) + c \\ &= d \lg n - d + c \leq d \lg n \end{aligned}$$

$$\text{if: } -d + c \leq 0, d \geq c$$

Example 2

$$T(n) = T(n-1) + n$$

- Guess: $T(n) = O(n^2)$
 - Induction goal: $T(n) \leq dn^2$, for some d and $n \geq n_0$
 - Induction hypothesis: $T(n-1) \leq d(n-1)^2$ for all $k < n$
- Proof of induction goal:

$$T(n) = T(n-1) + n \leq d(n-1)^2 + n$$

$$= dn^2 - (2dn - d - n) \leq dn^2$$

$$\text{if: } 2dn - d - n \geq 0 \Leftrightarrow d \geq n/(2n-1) \Leftrightarrow d \geq 1/(2 - 1/n)$$

- For $n \geq 1 \Rightarrow 2 - 1/n \geq 1 \Rightarrow$ any $d \geq 1$ will work

Example 3

$$T(n) = 3T(n/4) + cn^2$$

- Guess: $T(n) = O(n^2)$
 - Induction goal: $T(n) \leq dn^2$, for some d and $n \geq n_0$
 - Induction hypothesis: $T(n/4) \leq d(n/4)^2$
- Proof of induction goal:

$$\begin{aligned} T(n) &= 3T(n/4) + cn^2 \\ &\leq 3d(n/4)^2 + cn^2 \\ &= (3/16)d n^2 + cn^2 \\ &\leq d n^2 \quad \text{if: } d \geq (16/13)c \end{aligned}$$

- Therefore: $T(n) = O(n^2)$

Example 4

$$T(n) = 2T(n/2) + cn$$

- **Guess:** $T(n) = O(n \log n)$
- **Proof by Mathematical Induction:**

Prove that $T(n) \leq d n \log n$ for $d > 0$

$$T(n) \leq 2(d \cdot n/2 \cdot \log n/2) + cn$$

(where $T(n/2) \leq d \cdot n/2 (\log n/2)$ by induction hypothesis)

$$\leq dn \log n/2 + cn$$

$$= dn \log n - dn + cn$$

$$= dn \log n + (c-d)n$$

$$\leq dn \log n \quad \text{if } d \geq c$$

- Therefore, $T(n) = O(n \log n)$

Example 5

$$T(n) = 2T(n/2) + n$$

- Guess: $T(n) = O(n \lg n)$
 - Induction goal: $T(n) \leq dn \lg n$, for some d and $n \geq n_0$
 - Induction hypothesis: $T(n/2) \leq d(n/2) \lg(n/2)$
- Proof of induction goal:

$$\begin{aligned} T(n) &= 2T(n/2) + n \leq 2d(n/2)\lg(n/2) + n \\ &= dn \lg n - dn + n \leq dn \lg n \end{aligned}$$

$$\text{if: } -dn + n \leq 0 \Rightarrow d \geq 1$$