REGULAR EXPRESSION MODULE 2

- Regular Language:
 - · A language is regular if there exists a finite acceptor for it
 - Hence, every regular language can be described by using some NFA/DFA
- Regular Expression:
 - One of the way of describing regular languages which consists of strings, symbols and operators
 - This notation involves:
 - Combination of strings of symbols from some alphabet Σ
 - Parentheses
 - Operators + , · and *

Examples of Regular Expression:

- For language L1 = { a }
 Regular Expression R = a
- 2. Language $L2 = \{a, b, c\}$

R. E.
$$R = a + b + c$$

Here + represents union operation

Similarly . represents concatenation

* represents star - closure

Examples of Regular Expression:

3. Suppose R = (a + (b.c))*
It stands for star closure of { a } U { bc }
Then language will be
L = { €, a, bc, bca, abc, aa, aaa, bcbc, aabc, ... }

Formal definition:

Let Σ be a given alphabet then

- Φ, €, and a € Σ are all regular expressions.
 These are called Primitive Regular Expression
- 2. If r1 and r2 are regular expressions then r1 + r2, r1 . r2, r1* and (r1) are all regular expressions
- 3. A string is a regular expression if and only if it can be derived from the primitive regular expression by a finite number of applications of the rules defined in statement 2

Example:

If
$$\Sigma = \{ a, b, c \}$$

Then the string $(a + b \cdot c) * \cdot (c + a)$ is a regular expression

(a + b +) is not a regular expression

Language associated with Regular Expression

- Regular expressions can be used to describe some simple languages.
- If r is a regular expression then L (r) denote the language associated with r
- The language L (r) denoted by any regular expression r is defined by following rules:
 - 1. Φ is a regular expression denoting the empty set
 - 2. € is a regular expression denoting { € }
 - 3. For every $a \in \Sigma$, a is a regular expression denoting $\{a\}$

Language associated with Regular Expression

If r1 and r2 are regular expression then

1.
$$L(r1 + r2) = L(r1) U L(r2)$$

2.
$$L(r1.r2) = L(r1).L(r2)$$

3.
$$L(r1*) = (L(r1))*$$

4.
$$L((r1)) = L(r1)$$

Language associated with Regular Expression

- r* is known as kleen closure which indicate occurrences of r for ∞ number of times
- Examples:

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    If Σ = { a } and R = a*
    Then R = { €, a, aa, aaa, aaaa, ... }
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2. If
$$\Sigma = \{a\}$$
 and $R = a+$

Then
$$R = \{a, aa, aaa, aaaa, ...\}$$

This is known as Positive Closure

• Ex. 1 Exhibit the language in set notation

$$L(a*.(a+b))$$

Solution:

$$L(a*.(a+b)) = L(a*).L(a+b)$$

$$= (L(a))*.(L(a)UL(b))$$

$$= \{ \epsilon, a, aa, aaa, ... \}. \{ a, b \}$$

$$= \{ a, aa, aaa, ... , b, ab, aab, aaab, ... \}$$

Ex. 2 Represent the language for given regular expression

$$r = (a + b)^* . (a + bb)$$
 where $\Sigma = \{a, b\}$

Solution:

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L(r) = L((a + b)*.(a + bb))

= L((a + b)*).L(a + bb)

= (L(a + b))*.(L(a)UL(bb))

= (L(a)UL(b))*.({a, bb})

= ({a, b})*.{a, bb}

= {e, a, b, aa, ab, ba, ...}.{a, bb}

= {a, bb, aa, abb, ba, bbb, aaa, aabb, aba, abbb, ...
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L (r) is a language consists of all the strings terminated by either a or bb

Ex. 3 Exhibit the language for given regular expression

$$r = (aa)^* . (bb)^* . b$$

Solution:

 $L(r) = \{ b, aab, bbb, aabbb, aaaabbb, ... \}$

$$L(r) = \{ a^{2n} \cdot b^{2m+1} \mid n \ge 0, m \ge 0 \}$$

It denotes set of all strings with an even number of a's followed by odd number of b's

Ex. 4 Write the regular expression for the language accepting all combinations of a's over the set $\Sigma = \{a\}$

Solution:

$$\Sigma = \{a\}$$

Possible set of language $L = \{ \epsilon, a, aa, aaa, aaaa, ... \}$

$$R = a^*$$

This is known as Kleen Closure of a

Ex. 5 Write the regular expression for the language accepting all combinations of a's except the null string over the set $\Sigma = \{a\}$ Solution:

$$\Sigma = \{a\}$$

Possible set of language $L = \{ a, aa, aaa, aaaa, ... \}$

$$R = a +$$

This is known as Positive Closure of a

Ex. 6 Design the regular expression for the language containing all the strings having any number of a's and b's over the set $\Sigma = \{a, b\}$ Solution:

$$\Sigma = \{ a, b \}$$

Possible set of language $L = \{ \epsilon, a, b, aa, ab, ba, bb, aaa, ... \}$

$$R = (a + b)^*$$

Ex. 7 Construct the regular expression for the language containing all the strings having any number of a's and b's except the null string over the set $\Sigma = \{a, b\}$

Solution:

$$\Sigma = \{ a, b \}$$

Possible set of language $L = \{a, b, aa, ab, ba, bb, aaa, ...\}$

$$R = (a + b)^{+}$$

Ex. 8 Construct the regular expression for the language containing all the strings which are ended with 00 over the set $\Sigma = \{0, 1\}$ Solution:

$$\Sigma = \{ 0, 1 \}$$

R. E. = (Any combination of 0's and 1's) . 00 = (0 + 1)*00

Possible set of language $L = \{00, 000, 100, 0100, 1100, ...\}$

$$R = (0 + 1)*00$$

Ex. 9 Write regular expression for the language accepting the string which are starting with 1 and ending with 0 over the set $\Sigma = \{0, 1\}$ Solution:

$$\Sigma = \{ 0, 1 \}$$

R. E. = 1 . (Any combination of 0's and 1's) . 0
= 1 .
$$(0 + 1)^*$$
 . 0

Possible set of language L = { 10, 100, 110, 1000, 1010, ... }

$$R = 1 (0 + 1)* 0$$

Ex. 10 What is the regular expression for the language starting and ending with a and having any combination of b in between over the set $\Sigma = \{a, b\}$ Solution:

$$\begin{split} \Sigma &= \{ \ a, \ b \ \} \\ R. \ E. &= a \ . \ (\ Any \ combination \ of \ b) \ . \ a \\ &= a \ . \ (\ b \)^* \ . \ a \\ \end{split}$$
 Possible set of language L = { aa, aba, abba, abbba, ... }

R = ab*a

Ex. 11 For $\Sigma = \{ 0, 1 \}$ Give regular expression such that $L(r) = \{ w \in \Sigma^* \mid w \text{ has at least one pair of consecutive zeros } \}$ Solution:

$$\Sigma = \{ 0, 1 \}$$

R. E. =
$$(0 + 1)^*$$
. 00. $(0 + 1)^*$

Possible set of language $L = \{00, 000, 100, 001, 1001, 0001, ...\}$

$$R = (0 + 1)* 00 (0 + 1)*$$

$$= (0 + 1)^* (00)^+ (0 + 1)^*$$

Ex. 12 Describe in simple English language

Given regular expression is $r = (a + ab)^*$

Solution:

$$\Sigma = \{ a, b \}$$

R. E. =
$$(a + ab)^*$$

Possible set of language $L = \{ \epsilon, a, ab, aba, aab, abab, aaa, ... \}$

The language is beginning with zero or any number of a's but not having consecutive b's

Ex. 13 Write regular expression to denote the language L over Σ^* where

 $\Sigma = \{ \text{ a, b, c} \}$ in which every string will be such that any number of a's followed by any number of b's followed by any number of c's

Solution:

$$\Sigma = \{ a, b, c \}$$

R. E. = (any number of a's). (any number of b's). (any number of c's)

R. E. =
$$a*b*c*$$

Possible set of language $L = \{ \epsilon, a, ab, abc, aabbc, abbccc, aaa, ... \}$

Ex. 14 Write regular expression to denote the language L over Σ^* where

 $\Sigma = \{ a, b, c \}$ in which every string will be such that at least one a followed by at least one b followed by at least one c

Solution:

$$\Sigma = \{ a, b, c \}$$

R. E. = $a^+b^+c^+$

Possible set of language $L = \{ abc, aabc, abbc, aabbc, abbcc, aabbcc, ... \}$

REGULAR EXPRESSION MODULE 2

Ex. 15 Write regular expression to denote the language L over Σ^* where

 $\Sigma = \{~0,~1~\}$ in which every string will be such that at all strings which begins or end with 00~or~11

Solution:

$$R = L1 + L2$$

L1 – The strings which begin with 00 or 11

L2 – The strings which end with 00 or 11

$$L1 = (00 + 11)$$
. (Any combination of 0's and 1's)

$$= (00 + 11).(0 + 1)*$$

L2 = (Any combination of 0's and 1's).(00 + 11)

$$= (0 + 1)*.(00 + 11)$$

$$P = (00 \pm 11)(0 \pm 1)* \pm (0 \pm 1)*(00 \pm 11)$$

Ex. 16 Write regular expression to denote the language L over Σ^* where

 $\Sigma = \{ \text{ a, b} \}$ such that the third character from right end of the string is always a

Solution:

$$\Sigma = \{ a, b \}$$

R. E. = (Any number of character a's and b's). a. (a or b). (a or b)

$$R = (a + b)^* \cdot a \cdot (a + b) \cdot (a + b)$$

Possible set of language $L = \{ aaa, aba, abb, aaba, aabb, bababb, ... \}$

Ex. 17 Write regular expression to denote the language L over Σ^* where $\Sigma = \{a, b\}$ which accepts all the strings with at least two b's Solution:

$$\Sigma = \{ a, b \}$$

The two b's can be surrounded with any number of a's or b's in between

 $R.\ E.=$ (Any combination of a's and b's) . b . (Any combination of a's and b's) . b . (Any combination of a's and b's)

$$R = (a + b)^* . b . (a + b)^* . b . (a + b)^*$$

Possible set of language $L = \{ bb, bba, abb, bbb, aabb, abba, aabbaa, ... \}$

Ex. 18 Write regular expression to denote the language L having strings which should have at least one 0 and at least 1 where $\Sigma = \{0, 1\}$

Solution:

$$R = L1 + L2$$

L1 – The strings in which 0 precedes 1

L2 – The strings in which 1 precedes 0

$$L1 = (0 + 1)*.0.(0 + 1)*.1.(0 + 1)*$$

$$L2 = (0 + 1)^* . 1 . (0 + 1)^* . 0 . (0 + 1)^*$$

We can write resultant expression as

$$R = (0 + 1)*0(0 + 1)*1(0 + 1)*+(0 + 1)*1(0 + 1)*0(0$$

Ex. 19 Construct regular expression which contains even length of string over the set $\Sigma = \{\ 0\ \}$

Solution:

Given $\Sigma = \{ 0 \}$

The language with even number of 0's can be written as

$$L = \{ \epsilon, 00, 0000, 000000, \dots \}$$

We cann say that 00 always comes in pair.

$$R = (00)^*$$

Ex. 20 Write regular expression which denotes a language over the set $\Sigma = \{1\}$ having odd length of strings

Solution:

Given
$$\Sigma = \{ 1 \}$$

The language with odd number of 1's can be written as

$$L = \{ 1, 111, 111111, \dots \}$$

The regular expression can be written as

$$R = 1.(11)*$$

Ex. 21 Find a regular expression for the language

 $L = \{ w \in \{ 0, 1 \}^* \mid w \text{ has no pair of consecutive zeros } \}$

Solution:

The regular expression involves repetition of the strings of the form 1 ... 1

We can represent this as r1 = (1*011*)*

If the string can be terminated by 0 then we can modify the expression as

$$r1 = (1*011*)*.(0 + \epsilon)$$

The expression does not represent the strings with all 1's and strings with all 1's ending with 0

Ex. 21 Find a regular expression for the language

 $L = \{ w \in \{ 0, 1 \}^* \mid w \text{ has no pair of consecutive zeros } \}$

Solution:

It can be represented as

$$r2 = 1*.(0+\epsilon)$$

We can give resultant regular expression as

$$R = r1 + r2$$

$$R = (1*011*)*.(0+\epsilon) + 1*.(0+\epsilon)$$

Alternate Solution: $r = (1 + 01)^* \cdot (0 + \epsilon)$

$$r = (1*(01)*1*)*.(0 + \epsilon)$$

Ex. 22 Obtain regular expression such that

L(R) = { w | w \in { 0, 1}* and w has at least single occurrence of three consecutive zeros}

Solution:

$$\Sigma = \{ 0, 1 \}$$

R = (Any combination of 0's and 1's).000.(Any combination of 0's and 1's)

$$R = (0 + 1)^*.000.(0 + 1)^*$$

Possible set of language L = { 000, 1000, 0001, 10001, 00000011, ...

Ex. 23 Obtain regular expression such that

The set of all strings over $\Sigma = \{\ 0\ ,\ 1\ \}$ without length two Solution:

$$\Sigma = \{ 0, 1 \}$$

The regular expression can be splitted into three parts as

$$R = r1 + r2 + r3$$

Where r1 - Strings with length three or more

r2 - String with length one

r3 - String with length zero

Ex. 23 Obtain regular expression such that

The set of all strings over $\Sigma = \{\ 0\ ,\ 1\ \}$ without length two

Solution:

The expressions can be written as

$$R1 = (0+1)(0+1)(0+1)^{+}$$

$$R2 = (0 + 1)$$

$$R3 = \epsilon$$

Final Expression:

$$R = (0+1)(0+1)(0+1)^{+} + (0+1) + \epsilon$$

Ex. 24 Obtain regular expression over $\Sigma = \{\,0\,,\,1\,\}$ such that the set of all strings with number of zeros are odd Solution:

We need at least one zero to make count of zeros to odd. This zero can be surrounded by any number of 1's. Therefore we can write it as R1 = 1*.0.1*

Now this expression can be followed with pair of zeros. Those pair of zeros can be surrounded by any number of 1's. It can be written as R2 = (1*.0.1*.0.1*).

This R2 can be repeated any number of times.

Ex. 24 Obtain regular expression over $\Sigma = \{\,0\,,\,1\,\}$ such that the set of all strings with number of zeros are odd Solution (Contd.) :

When we concatenate R1 and R2 then no of 1's will remain odd

R = R1.R2

R = 1*.0.1*.(1*.0.1*.0.1*)*

This expression can be further reduced as R = 1*.0.1*.(0.1*0.1*)*

Final Answer: R = 1*.0.1*.(0.1*.0.1*)*

Ex. 25 Obtain regular expression over $\Sigma = \{\,0\,\,,\,1\,\}$ such that the set of all strings containing both 11 and 010 as substring

Solution:

$$\Sigma = \{ 0, 1 \}$$

We split regular expression in two parts:

R1 = The strings in which 11 precedes 010

$$= (0+1)^*.11.(0+1)^*.010.(0+1)^*$$

R2 = The strings in which 010 precedes 11

$$= (0+1)^*.010.(0+1)^*.11.(0+1)^*$$

Final Expression: R = R1 + R2

$$R = (0+1)*11(0+1)*010(0+1)* + (0+1)*010(0+1)*11(0+1)*$$

Ex. 26 Obtain regular expression over $\Sigma = \{0, 1\}$ such that the set of all strings begin or end with 00 or 11

Solution:

$$\Sigma = \{ 0, 1 \}$$

Consider six different cases:

R1 = Strings beginning with 00 and ending with 11

R2 = Strings beginning with 11 and ending with 00

R3 = Strings beginning with 00

R4 = Strings beginning with 11

R5 = Strings ending with 00

R6 = Strings ending with 11

Ex. 26 Obtain regular expression over $\Sigma = \{0, 1\}$ such that the set of all strings begin or end with 00 or 11

Solution:

Regular expression can be written as

$$R1 = 00 (0+1)*11$$

$$R2 = 11 (0+1)*00$$

$$R3 = 00 (0+1)*$$

$$R4 = 11 (0+1)*$$

$$R5 = (0+1)*00$$

$$R6 = (0+1)*11$$

Ex. 26 Obtain regular expression over $\Sigma = \{0, 1\}$ such that the set of all strings begin or end with 00 or 11

Solution:

Regular expression can be written as

$$R = R1 + R2 + R3 + R4 + R5 + R6$$

$$R = 00 (0+1)*11 + 11(0+1)*00 + 00 (0+1)* + 11 (0+1)* + (0+1)*00 + (0+1)*11$$