

Name: \_\_\_\_\_

Std No: \_\_\_\_\_

**Q.1 (1.5 \* 4 = 6) Circle exactly one best answer for each of the following four multiple-choice problems.**

1. The class  $NP$  is the set of all decision problems that:
  - a. Can be solved by polynomial-time algorithms.
  - b. Can definitely *not* be solved by polynomial-time algorithms.
  - c. Have polynomial-time algorithms that can verify potential solutions.
  - d. All of the above.
  - e. None of the above.
2. The class  $NP$ -complete is the set of all decision problems that:
  - a. Can be solved by polynomial-time algorithms.
  - b. Can definitely *not* be solved by polynomial-time algorithms.
  - c. Have polynomial-time algorithms that can verify potential solutions.
  - d. All of the above.
  - e. None of the above.
3. Suppose  $X \leq_p Y$ . Which must be true?
  - a. Problem  $X$  is polynomial-time reducible to problem  $Y$ .
  - b. Problem  $Y$  is polynomial-time reducible to problem  $X$ .
  - c. Problems  $X$  and  $Y$  are equivalent in terms of computational complexity.
  - d. Both (a) and (c).
  - e. None of the above.
4. Suppose problem  $X$  is in class  $P$ , problem  $Y$  is in class  $NP$ , and  $Y \leq_p X$ . Which must be true?
  - a. Problem  $Y$  is in class  $P$ .
  - b. Problem  $Y$  is  $NP$ -complete.
  - c.  $P = NP$
  - d. Both (a) and (c)
  - e. None of the above.

**Q.2 (4\*1\*5 = 6) Indicate whether the following statements are True or False (T/F)**

- i) If PRIME is in  $NP$ -complete, then  $P \neq NP$ .
- ii)  $NP=co-NP \Rightarrow P = NP$ .
- iii)  $NP$ -hard problems which are not  $NP$ -complete are not in  $NP$ .
- iv) The TAUTOLOGY problem asks if a given Boolean formula is true for all possible assignments to the Boolean variables. TAUTOLOGY is in  $co-NP$ .

**Q. 3 (4+4 = 8 pts)**

i) Suppose you have found a problem which is in  $NP$  but not in  $P$ . Why do you think your finding is important?

ii) Point out and briefly discuss the fallacy in the following “proof” that  $P \neq NP$ :  
“To see if a 3-SAT formula is satisfiable, we need to look at  $2^n$  possible truth assignments. This takes exponential time, so 3-SAT is not in  $P$ . But it is in  $NP$ , so  $P \neq NP$ .”