

Roll - 1705045

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Ans. to -1

Q) In computer, ~~total~~ total CPU clock cycle =  $\text{Multiplication Instruction count} \times \text{clock cycle for multiplication} + \text{Addition count} \times \text{Addition clock cycle}$

For A, CPU time,  $\text{CPU}_A = \text{cycle time} \times \text{CPU clock cycle}$

$$= \cancel{500 \text{ ps}} \times \text{Instruction count} \times \text{CPI}$$

$$= 500 \cancel{\text{ ps}} \times (3 \times 3 + 2 \times 2) \text{ ps}$$

$$= 500 \times (9 + 4) \text{ ps}$$

$$= 6500 \text{ ps}$$

CPU time in B,  $\text{CPU}_B = \text{cycle time} \times \text{CPU clock cycle}$

$$= 450 \times (3 \times 4 + 2 \times 1)$$

$$= 450 \times (12 + 2)$$

$$= 6300 \text{ ps}$$

So, CPU ~~ex~~ time B is less. Hence, regard the program, B's performance is better, as

$$\text{CPU}_A > \text{CPU}_B.$$

Q.

Ans. to q. - 2

Here, we will try to modify  $X_i$  as a function of  $A_i B_i$  and  $Y_i$  as a function of  $C_i$ .

$S_2$	$S_1$	$C_{in}$	$X_i$	$Y_i$	F (through Parallel Adder)
0	0	0	$A_i B_i$	0	$AB + C$
0	0	1	$A_i B_i$	$C_i$	$AB + C + 1$
0	1	0	$A_i B_i$	0	$AB$
0	1	1	$A_i B_i$	0	$AB + 1$
1	0	0 (forcefully)	$A_i B_i$	1	$(AB)'$
1	1	0 (forcefully)	$A_i B_i$	0	$AB$

In this two,  
 $C_{in} = 0$  forcefully,  
 (all step)  
 So,  $F = X_i \oplus Y_i$

So,

$$X_i = A_i B_i$$

For  $Y_i$ ,

$S_2 \backslash S_1$	$C_{in}$	00	01	11	10
00			1	1	
01					
11					
10		1	1		

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Q8

So,

$$Y_i = S'_2 S'_1 C_p + S_2 S'_1 C'_{in}$$

And,  $Z_i = \overline{S_2} C_{Adder,i}$

Here,  $C_{Adder}$  means parallel adder's carry outputs. Specifically,  $C_{Adder,0} = C_{in}$