

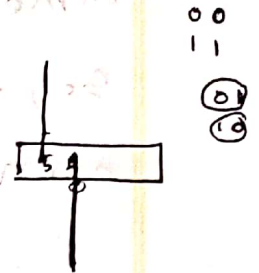
ID-1705045

① Yes, we can do it.

② Reason:-

Example matrix —

(1,1)	(1,2)	(1,3)	(1,4)
(2,1)	(2,2)	(2,3)	(2,4)
(3,1)	(3,2)	(3,3)	(3,4)
(4,1)	(4,2)	(4,3)	(4,4)



	c-bit	Parity of data bits
Parity set A →	1	3, 5, 7, 9, 11
Parity set B →	2	3, 6, 7, 10, 11
Parity set C →	:	

Errorred bit are boxed.

As burst of length at $k+1$ bits, only one row will have at most 2 erroneous bit. All other rows will have at most 1 erroneous bit.

When there will be 1 erroneous bit in a row, of course at least one 1 ~~check~~ parity-set will have even parity.

Now, when there will be at most 2 erroneous bit in a row, we need to proof that all parity-sets will not have odd parity. In case of one parity set to have odd parity even after having erroneous bits, there must be 2 erroneous bit. Otherwise, it would have even parity.

Now, we know, every number's binary representation are distinct. So, the row and 2 erroneous bit we are talking about, there must be some parity-set having not both of them. So, that parity-set must ~~have~~ have its parity even. And, we can detect this.

So, we will simply detect by checking if there are any even parity.