

Name

Student ID

There are **four** questions. Answer ALL.

1. **[6 points]** Consider the following solution for the producer-consumer problem. Give an example of a scenario where this code fails. Assume down and up operations have their standard definition.

<pre>#define N = 100 typedef int semaphore; semaphore mutex = 1; int count = 0; void producer (void) { int item; while (TRUE){ item = produce(); down(&mutex); if (count == N) sleep(); inset(item); count = count + 1; if (count == 1) wakeup(consumer); up(&mutex); } }</pre>	<pre>void consumer (void) { int item; while (TRUE){ down(&mutex); if (count == 0) sleep(); item = remove-item(); count = count - 1; if (count == N-1) wakeup(producer); up(&mutex); consume(item); } }</pre>	<p><u>Answer:</u></p> <p><u>When N ==1</u> Producer enters the critical region by down(mutex). Now since count = N, producer will call sleep();</p> <p>On the other hand, consumer will try to enter the critical region by calling down (mutex), but it cannot. So it will sleep too.</p> <p>Both will sleep forever.</p>
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2. **[4 points]** Now consider yet another solution for the producer-consumer problem that eliminates the failure scenario that is present in the solution in Question No 1. Give an example scenario to show the demerit of this solution.

<pre>#define N = 100 typedef int semaphore; semaphore mutex = 1 int count = 0; void producer (void) { int item; while (TRUE){ item = produce(); down(&mutex); if (count < N) { inset(item); count = count + 1;} up(&mutex); } }</pre>	<pre>void consumer (void) { int item; while (TRUE){ down(&mutex); if (count > 0) { item = remove(); count = count - 1;} up(&mutex); consume(item); } }</pre>	<p><u>Answer:</u></p> <p>Busy waiting problem. Wastage of resources.</p> <p>Suppose when count = N, the producer cannot insert an item. Suppose the scheduler does not schedule consumer for a long period of time.. In the meantime, producer will unnecessarily execute the loop utilizing the CPU.</p>
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3. **[6 points]** Consider you have two processes called Examiner and Scrutinizer. The Examiner process performs the following tasks: grade and update_marksheet. The Scrutinizer process performs the following tasks: verify_grade and verify_update_marksheet. The tasks of Examiner and Scrutinizer processes must be executed in the following order: grade → verify_grade → update_marksheet → verify_update_marksheet. In the skeleton code for Examiner and Scrutinizer given below, initialize and insert down and up operations of the given semaphores to ensure the above mentioned execution sequence.

<pre>semaphore e ← 1 semaphore s ← 0</pre>	<pre>// Initialize e and s</pre>
<pre>Process Examiner { while(TRUE){ Down e grade(); Up s Down e update_marksheet(): Up s } }</pre>	<pre>Process Scrutinizer { while(TRUE){ Down s verify_grade(); Up e Down s verify_update_marksheet(): Up e } }</pre>

4. **[4 points]** Write one advantage and one disadvantage of an User-level Thread over a Kernel-level Thread.

Advantage: Context switching is cheaper.
Disadvantage: Less control of Kernel. A buggy user level thread may occupy CPU for a long time.