

5)

$$\rightarrow a) f(n) = 5n^5 - 4n^4 + 3n^3 + \sqrt{2}n^2 + n - 5$$

- $5n^5$ is the dominant term because it grows faster than any other term, therefore the function is $O(n^5)$.

$$\rightarrow b) f(n) = (n^3 - (\log(n))^3)(n^2 - \log(n)) + 55n^3$$

$$= (n^3 \cdot n^2) + (n^3 \cdot -\log(n)) + (-(\log(n))^3 \cdot n^2) + (-(\log(n))^3 \cdot -\log(n)) + 55n^3$$

$$= n^5 - n^3 \log(n) - n^2 (\log(n))^3 + (\log(n))^4 + 55n^3$$

- The dominant term is n^5 because it grows the fastest. Therefore the function is $O(n^5)$.

$$\rightarrow c) f(n) = \frac{5n^4 + 10n^3 - 100n^2 - n - 5}{6n^2}$$

$$= \frac{5n^4}{6n^2} + \frac{10n^3}{6n^2} + \frac{100n^2}{6n^2} + \frac{-n}{6n^2} + \frac{-5}{6n^2}$$

$$= \frac{5}{6}n^2 + \frac{10}{6}n + \frac{-100}{6} + \frac{-1}{6n} + \frac{-5}{6n^2}$$

- The dominant term is $\frac{5}{6}n^2$, therefore the function is $O(n^2)$.

$$\rightarrow d) f(n) = \log(n^3 + n + 30) + n^2 \log(n + 4)$$

$$= \log(n^3) + \log(n)$$

- $= O(\log(n)) + O(n^2 \log(n))$ // BigO of both terms.

- Since the dominant term is $O(n^2 \log(n))$ that is the BigO of this function.

$$\rightarrow e) f(n) = (n \log(n) + 5)^2 + (\log(n) + 5)(n^2 + 5)$$

$$= [(n \log(n))^2 + 2(n \log(n))(5) + 5^2] + [(\log(n))(n^2) + (\log(n))(5) + (5)(n^2) + (5)(5)]$$

$$= n^2 \log^2(n) + 2n \log(n) + 25 + n^2 \log(n) + \log(n) + n^2 + 25$$

- The dominant term is $n^2 \log^2(n)$, therefore the function is $O(n^2 \log^2(n))$.

$$\rightarrow f) f(n) = \log[(5n^5 + 7n^3 + 10)^2 (3n^3 + 4n + 50)]$$

$$- = \log(5n^5 + 7n^3 + 10)^2 + \log(3n^3 + 4n + 50)$$

$$- 5n^5 + 7n^3 + 50 \approx n^5 ; \log((n^5)^2) = \log(n^{10}) = 10 \log(n)$$

$$- 3n^3 + 4n + 50 \approx n^3 ; \log(n^3) = 3 \log(n)$$

- In this case once we add the two sides, we are left with $53 \log(n)$, but since constant factors are ignored, the function is $\underline{O(\log n)}$.

2)

$$\rightarrow a) f(n) = n^2 + 50 \in O(n^3)$$

$$- n^2 + 50 \leq cn^3$$

$$- \frac{n^2}{n^3} + \frac{50}{n^3} \leq c$$

$$- \frac{1}{n} + \frac{50}{n^3} \leq c$$

$$- \text{For } n \geq 2, \frac{1}{n} \leq \frac{1}{2}$$

$$- \text{Thus } n_0 = 2, \frac{1}{2} + \frac{50}{8} = \frac{1}{2} + \frac{5}{4} = \frac{7}{4} \leq 2$$

$$- \underline{c=2, n_0=2} ; \text{ Thus } O(n^3) \checkmark$$

$$\rightarrow b) n! \in O(n^n)$$

$$- n! \leq n^n \text{ for all } n \geq 5.$$

$$- \underline{c=5, n_0=5} ; \text{ Thus } O(n^n) \checkmark$$

$$\rightarrow c) f(n) = 3n^3 + 7n + 50 \in O(n^3)$$

$$- 3n^3 + 7n + 50 \leq cn^3$$

$$- 3 + \frac{7}{n^2} + \frac{50}{n^3} \leq c$$

$$- \text{For } n \geq 5, \frac{7}{n^2} \leq 7, \frac{50}{n^3} \leq 50$$

$$- \text{Thus } n \geq 2 ; 3 + \frac{7}{n^2} + \frac{50}{n^3} \leq 3 + 7 = 10$$

$$- \underline{c=10, n_0=2} ; \text{ Thus } O(n^3) \checkmark$$

→ d) $1^3 + 2^3 + 3^3 + \dots + n^3 \in O(n^4)$ // Sum of cubes

$$- S(n) = \left(\frac{n(n+1)}{2}\right)^2$$

$$- S(n) = \frac{n^2(n+1)^2}{4}$$

$$- S(n) = \frac{n^2(n^2+2n+1)}{4}$$

$$- S(n) = \frac{n^4 + 2n^3 + n^2}{4}$$

$$- n^4 + 2n^3 + n^2/4 \leq cn^4$$

$$- \frac{n^4}{4n^4} + \frac{2n^3}{4n^4} + \frac{n^2}{4n^4} \leq c$$

$$- \frac{1}{4} + \frac{2}{4n} + \frac{1}{4n^2} \leq c$$

$$- \text{For } n \geq 2, \frac{1}{4} + \frac{2}{8} + \frac{1}{36} = \frac{53}{36}$$

$$- \text{Thus } c=5, \frac{1}{4} + \frac{2}{4n} + \frac{1}{4n^2} \leq n$$

$$- \underline{c=5}, \underline{n_0=2}; \text{ Thus } O(n^4) \checkmark$$

3)

→ - First function worst case: Either the first if or the second else if will trigger a loop that runs n^2 times ($n * n$). The else block (loop runs 30 times) is a constant.

$$- O(n) \cdot O(n^2) = \underline{O(n^3)}$$

→ Second function worst case: Both loops run n times.

$$- \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$$- \underline{O(n^2)}$$

→ Third function worst case: Outer loop runs \checkmark times and i is halved each time. Inner loop j doubles each time and runs $\log(n)$ times.

$$- O(\log(n)) \cdot O(\log(n)) = O(\log^2(n))$$

→ Forth function worst case: Outer loop, i is halved each time and runs $\log(n)$ times. Middle loop doubles each time so it also runs $\log(n)$ times.

Inner loop increments by 2, so it will run $\frac{n}{2}$ times, which is $O(n)$.

$$- O(\log(n)) \cdot O(\log(n)) \cdot O(n) = \underline{O(n \log^2(n))}$$

4)

→ a) Give the n^{th} term of the sequence.

- $a_n = a_1 + (n-1) \cdot d$; $a_n = 28 + (n-1) \cdot 3$

- $a_n = 28 + 3n - 3$

- $\underline{a_n = 25 + 3n}$

→ b) What term is 552 in this sequence?

- $552 = 25 + 3n$

- $87 = 3n$

- $\underline{n = 29}$, 552 is the 29th term in the sequence.

→ c) What are the first 50 numbers in the sequence?

- For $n=1$: $a(1) = 3(1) + 25 = \underline{28}$

- For $n=2$: $a(2) = 3(2) + 25 = \underline{31}$

- For $n=3$: $a(3) = 3(3) + 25 = \underline{34}$

- For $n=4$: $a(4) = 3(4) + 25 = \underline{37}$

- For $n=5$: $a(5) = 3(5) + 25 = \underline{40}$

- For $n=6$: $a(6) = 3(6) + 25 = \underline{43}$

- For $n=7$: $a(7) = 3(7) + 25 = \underline{46}$

- For $n=8$: $a(8) = 3(8) + 25 = \underline{49}$

- For $n=9$: $a(9) = 3(9) + 25 = \underline{52}$

- For $n=50$: $a(50) = 3(50) + 25 = \underline{55}$

→ d) Find the sum of first 50 numbers in the sequence.

- $S_n = \frac{n}{2} \cdot (a_1 + a_n)$

- $S_n = \frac{50}{2} \cdot (28 + 55) = 5 \cdot 83 = \underline{415}$