



Field-sensitive Points-to Analysis with CFL-Reachability

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17.12.2021

- Context-free grammar — $G = (\Sigma, N, P, S)$
 - ▶ Σ — set of terminal symbols
 - ▶ N — set of non-terminal symbols
 - ▶ P — set of production rules $\{A \rightarrow \alpha \mid A \in N, \alpha \in (\Sigma \cup N)^*\}$
 - ▶ $S \in N$ — start symbol
- Graph — $\mathcal{G} = (V, E, I)$
 - ▶ V — set of vertices
 - ▶ $E \subseteq V \times V$ — set of edges
 - ▶ $I : E \rightarrow L$
- $\{(v_1, v_n) : \exists p = (e_1, e_2, \dots, e_n) \in E^*, \text{src}(e_1) = v_1, \text{dst}(e_n) = v_n, I(e_1) \cdot \dots \cdot I(e_n) \in L(G)\}$

Field-sensitive Points-to Analysis

- Graph

Relation	Statement	Graph edge
$alloc \subseteq \text{Vars} \times \text{Objects}$	$x = \text{new Obj}();$	$x \xrightarrow{alloc} h$
$assign \subseteq \text{Vars} \times \text{Vars}$	$x = y;$	$x \xrightarrow{assign} y$
$load_f \subseteq \text{Vars} \times \text{Vars}$	$x = y.f;$	$x \xrightarrow{load_f} y$
$store_f \subseteq \text{Vars} \times \text{Vars}$	$x.f = y;$	$x \xrightarrow{store_f} y$

- Context-free grammar defining analysis

$PointsTo \rightarrow (assign \mid load_f \text{ Alias } store_f)^* alloc$

$Alias \rightarrow PointsTo \ FlowsTo$

$\forall f \in Fields$

$FlowsTo \rightarrow \overline{alloc} (\overline{assign} \mid \overline{store_f} \text{ Alias } \overline{load_f})^*$

Example Graph

v1 = new Obj(); // h1

v2 = new Obj(); // h2

v4 = new Obj(); // h3

v6 = new Obj(); // h4

v5 = v4;

v5 = v6;

v1.h = v1;

v2.g = v1;

v4.f = v2;

v7 = v5.f;

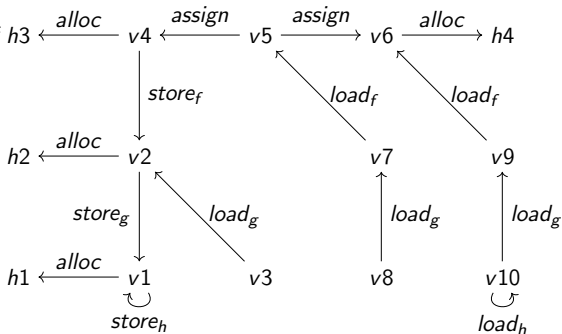
v9 = v6.f;

v3 = v2.g;

v8 = v7.g;

v10 = v9.g;

v10 = v10.h;



Example Graph with PointsTo Edges

```
v1 = new Obj(); // h1
```

```
v2 = new Obj(); // h2
```

```
v4 = new Obj(); // h3
```

```
v6 = new Obj(); // h4
```

```
v5 = v4;
```

```
v5 = v6;
```

```
v1.h = v1;
```

```
v2.g = v1;
```

```
v4.f = v2;
```

```
v7 = v5.f;
```

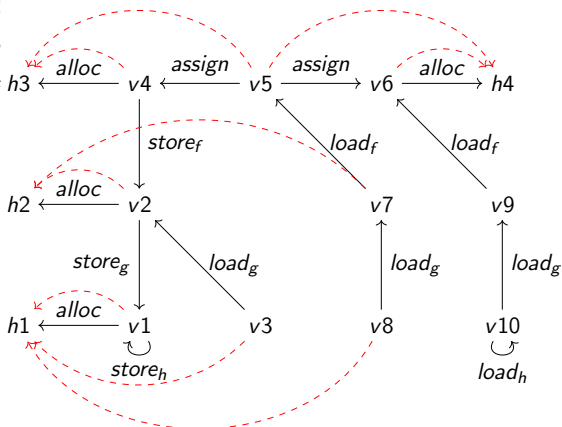
```
v9 = v6.f;
```

```
v3 = v2.g;
```

```
v8 = v7.g;
```

```
v10 = v9.g;
```

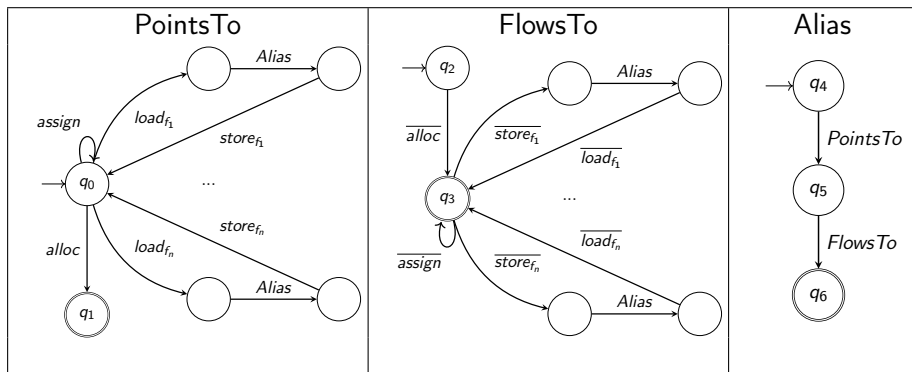
```
v10 = v10.h;
```



Algorithms Based on Linear Algebra Operations

- Matrix multiplication based algorithm
 - ▶ Grammar should be converted to a Weak Chomsky Normal Form which significantly increases its size
 - ▶ The performance of the algorithm depends on the grammar size
- Kronecker Product based algorithm
 - ▶ Grammar should be converted to a recursive state machine (RSM)
 - ▶ For the intersection of RSM and a directed labeled graph, they are represented as a composition of adjacency boolean matrices for each label

RSM for Field-sensitive Points-to Analysis



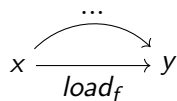
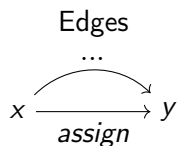
One Terminal, One Non-terminal Representation

- The graph and RSM are represented as adjacency matrices, whose elements are integers
- The high-order bits encode the non-terminal, the low-order bits encode the terminal



- Operation for Kronecker product
 - ▶ $times(x, y) = (x_{nonterm} = y_{nonterm} \neq 0 \text{ or } x_{term} = y_{term} \neq 0)$

Adapting Graph to this Representation



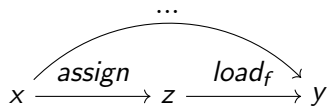
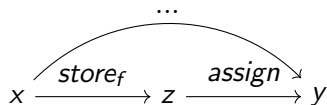
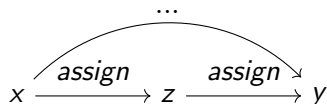
New variable

$z = y; x = z;$

$z = y; x.f = z;$

$z = y.f; x = z;$

Result



- The algorithm based on the Kronecker product has been optimized for the field-sensitive points-to analysis problem
- The developed algorithm was implemented within the CFPQ_PyAlgo¹ framework
- The correctness of the algorithm has been checked on the Dataset²

Future Work:

- Extend Benchmark Graphalytics³
 - ▶ Investigate the performance of the developed algorithm
 - ▶ Compare performance with other CFL-Reachability algorithms

¹https://github.com/JetBrains-Research/CFPQ_PyAlgo

²<https://bitbucket.org/jensdietrich/gigascale-pointsto-oopsla2015>

³<https://graphalytics.org/>