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EXCEL LONDON / UNITED KINGDOM

DIVIDE ET IMPERA:
MEMORY RANGER RUNS DRIVERS IN
ISOLATED KERNEL SPACES

Igor Korkin, Ph.D

WHOAMI

- MEPhI Alumni, PhD in Cyber Security, published 23 papers
- Area of interest is Windows Kernel security:
 - Memory Forensics
 - Rootkits Detection
 - Bare-Metal Hypervisors
- Fan of academic cross-disciplinary research igorkorkin.blogspot.com
- Love traveling and powerlifting (6) igor.korkin

AGENDA

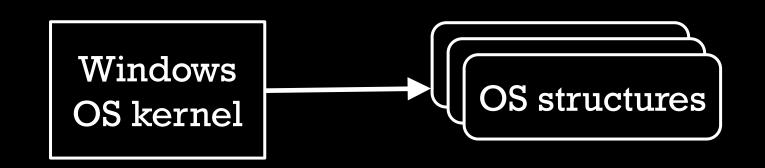
Attacking the kernel-mode memory

Existing protection: Windows built-in security and research projects

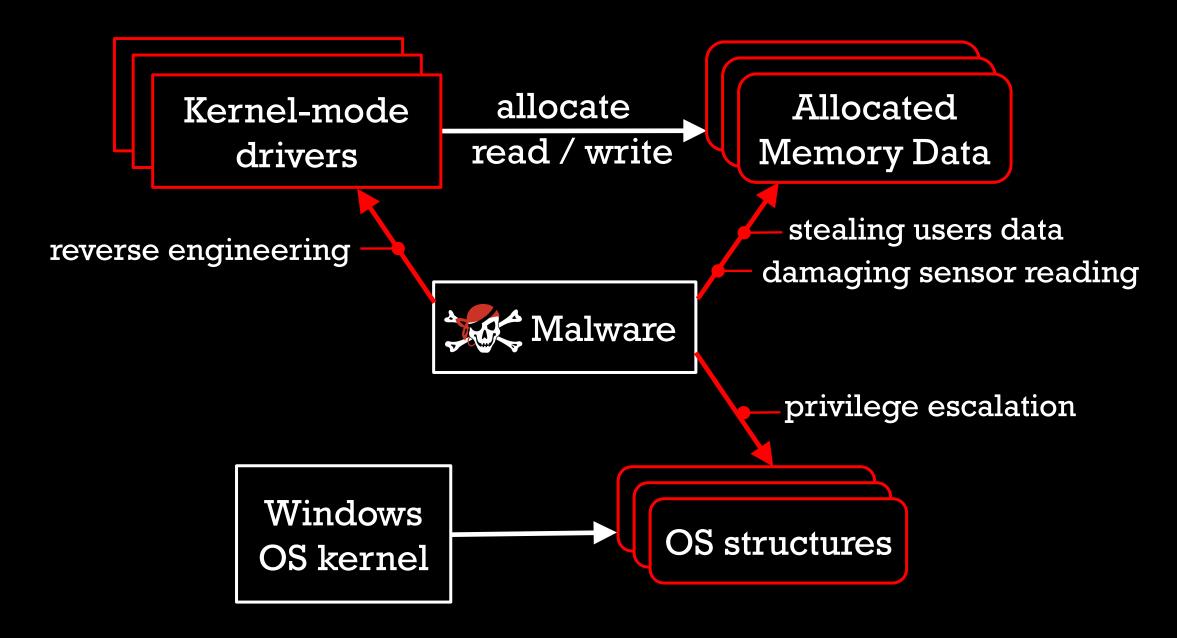
MemoryRanger hypervisor: idea, details, demos

ATTACKS ON KERNEL MODE MEMORY





ATTACKS ON KERNEL MODE MEMORY



TWO HOUSES WITH PRIVATE ART COLLECTIONS

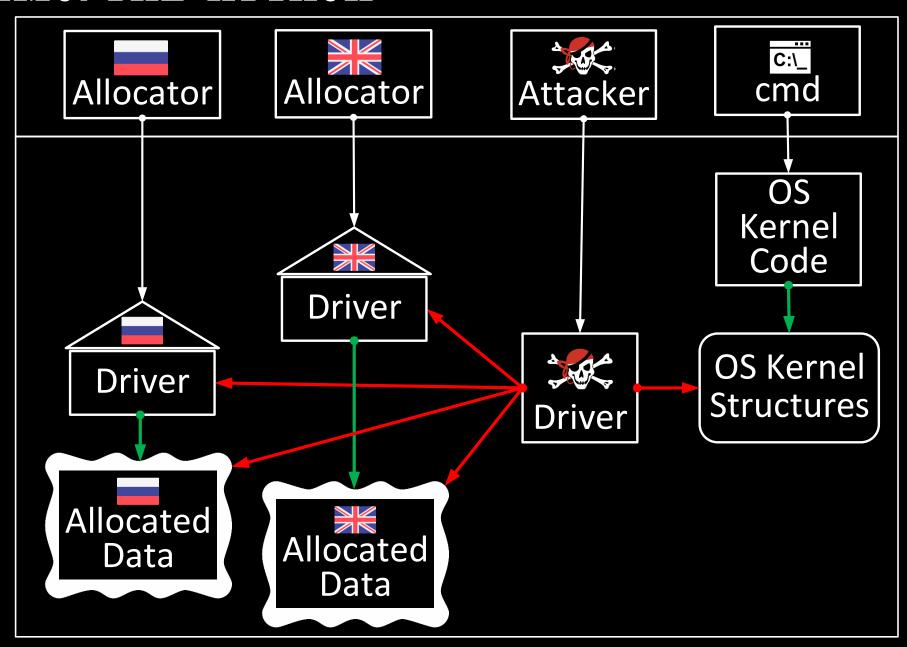






Hi all! I am a hacker-attacker!
I will inspect these houses to
steal and damage painting!

DEMO: THE ATTACK



DEMO: THE ATTACK

The online version is here -

https://www.youtube.com/embed/HNxc-tjy3QA?vq=hd1080

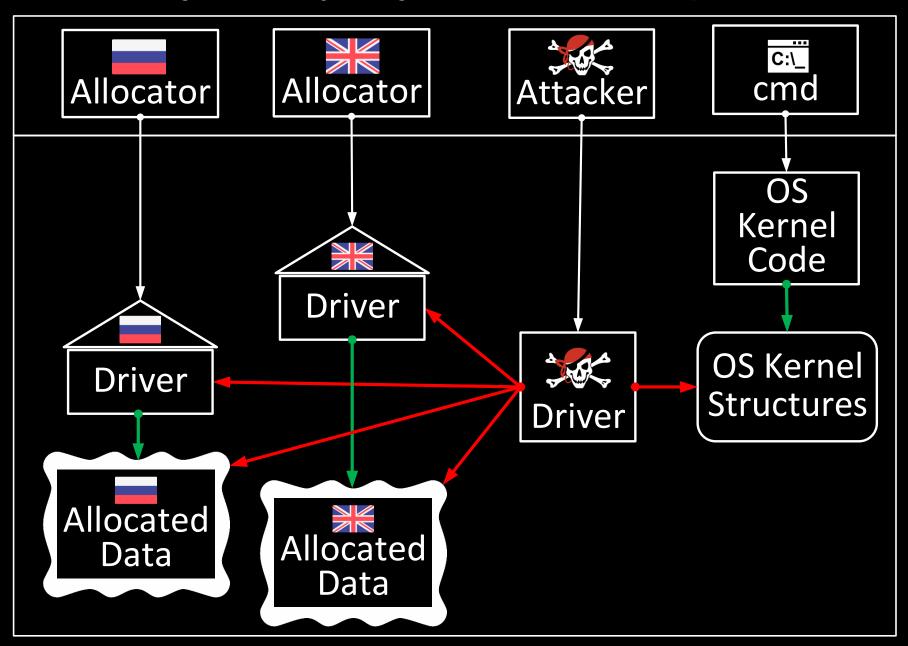
TWO HOUSES WITH PRIVATE ART COLLECTIONS







THE ATTACK HAS NOT BEEN PREVENTED

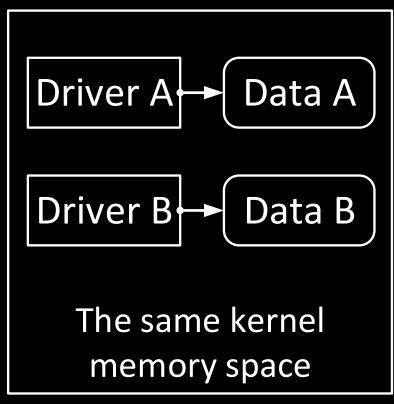


BACKGROUND ANALYSIS

Memory protection projects	Malware attacks on								
		e of OS &		OS data:	Data of third-				
	third-pa	arty drivers	interr	nal structures	party drivers				
	Read	Write	Read	Write	Read	Write			
Windows Security	-	BSOD 0xBE by Device Guard	-	BSOD 0x109 by PatchGuard	-	_			
PrivGuard	_	_	-	+	_	-			
LAKEED	+	+	+	+	_	-			
LKMG	_	+	+	+	+	+			
rR^X	+	+	-	_	_	-			
AllMemPro	-	_	+	+	+	+			
Memory Ranger	+	+	+	+	+	+			

IDEA OF DRIVERS EXECUTION ISOLATION

Now all drivers share the same memory space

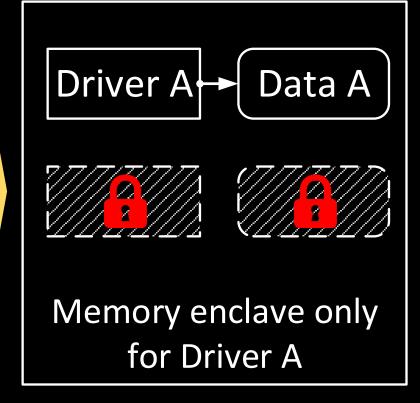


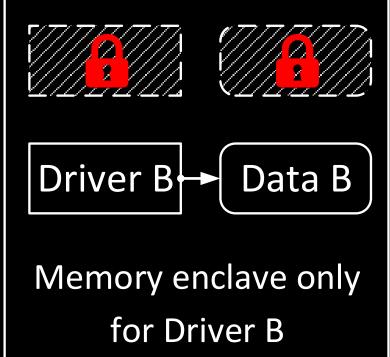
IDEA OF DRIVERS EXECUTION ISOLATION

Now all drivers share the same memory space

Driver A Data A Driver B Data B The same kernel memory space

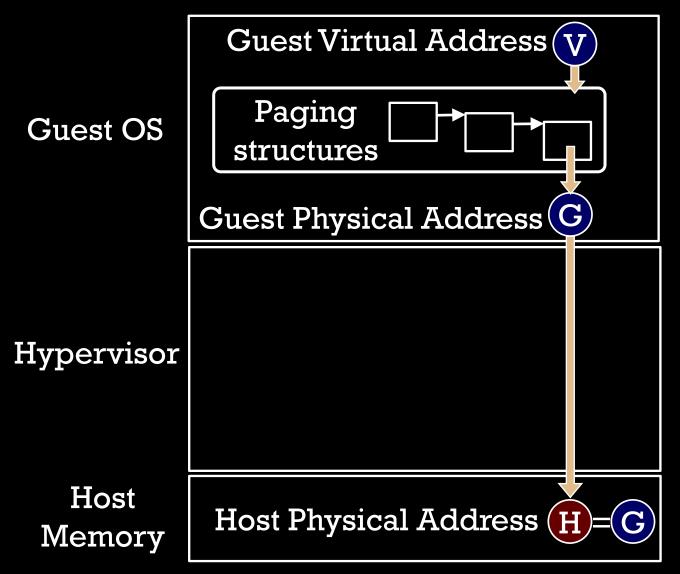
Let's execute these two drivers into separate memory enclosures





PROCESSING MEMORY ACCESS: EPT FEATURE

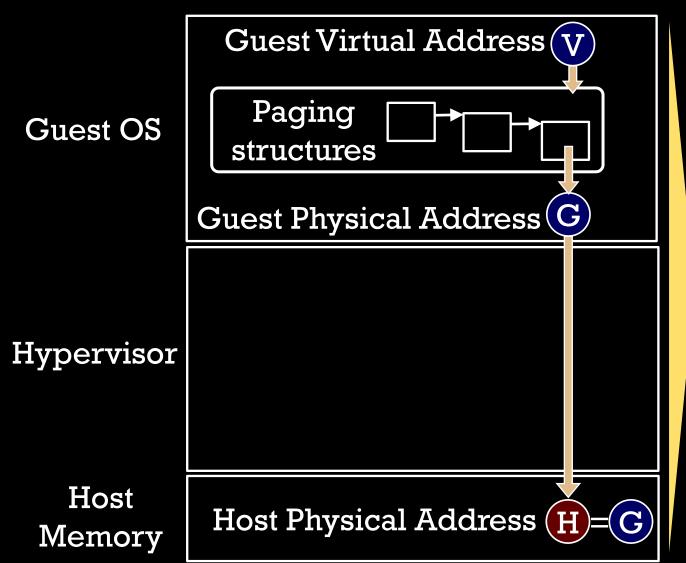
VT-x without EPT

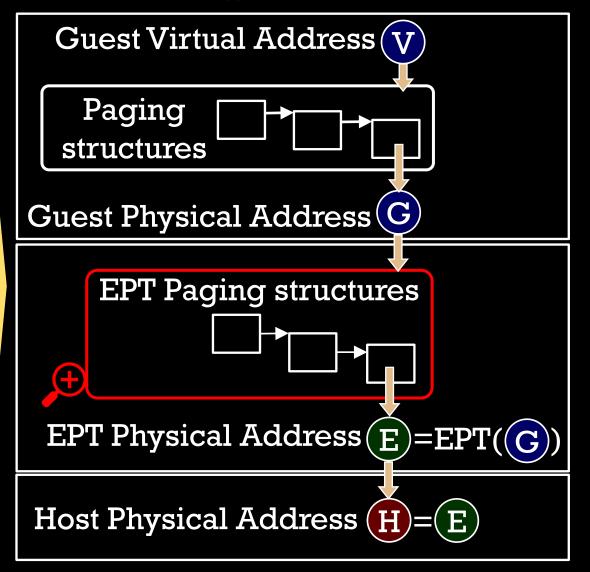


PROCESSING MEMORY ACCESS: EPT FEATURE

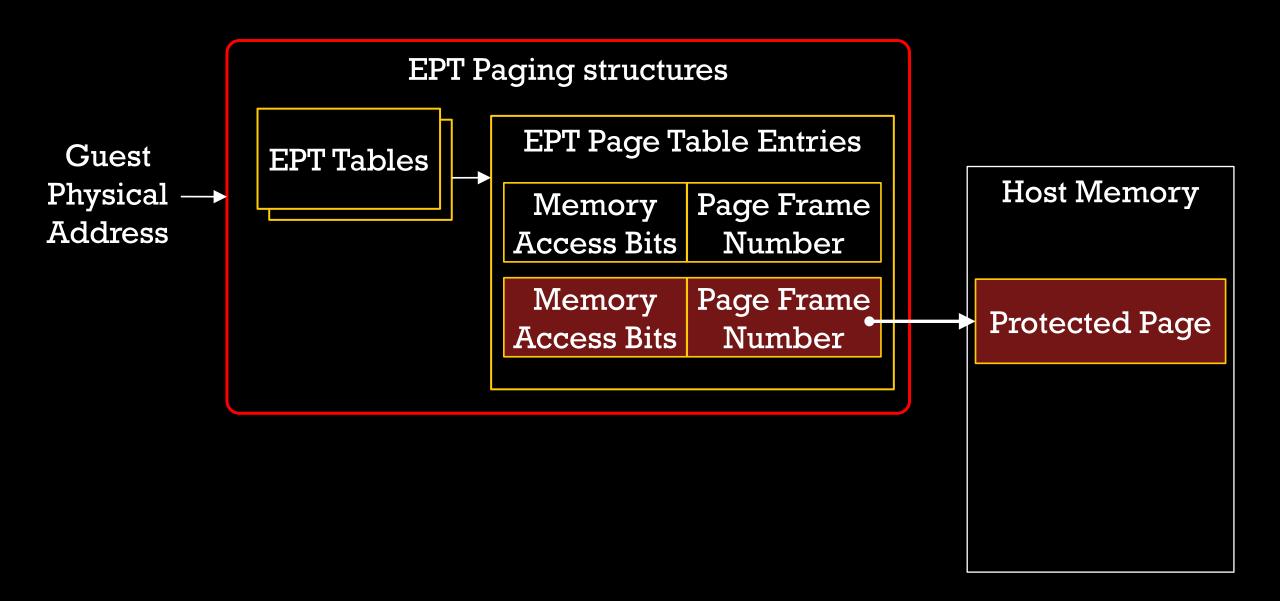
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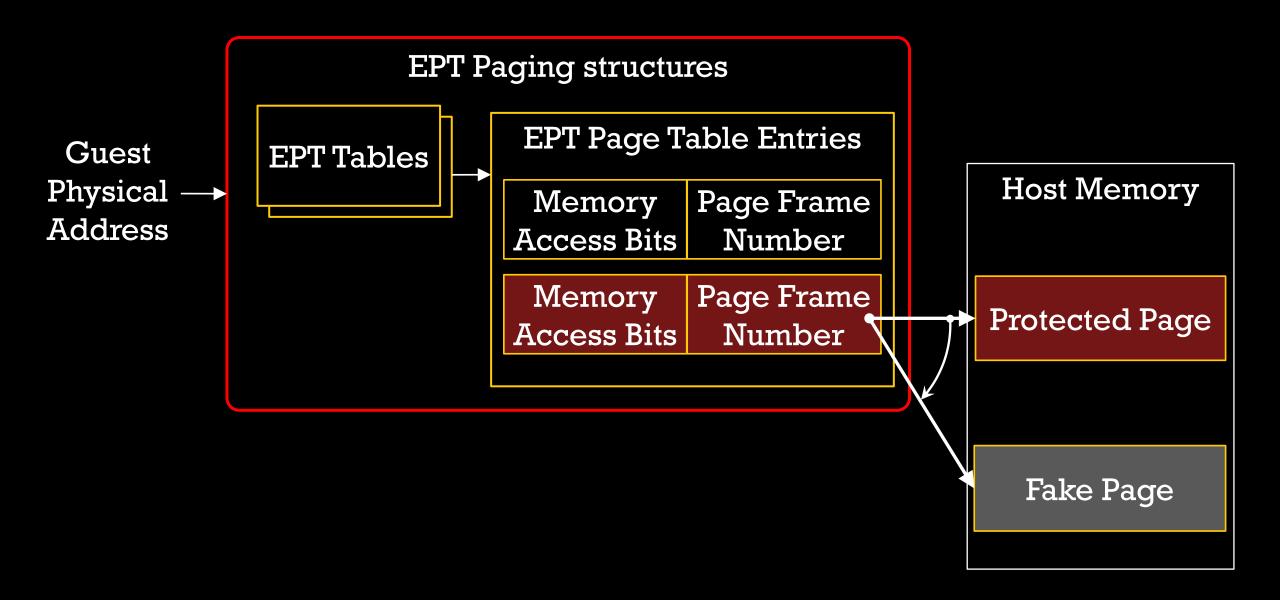




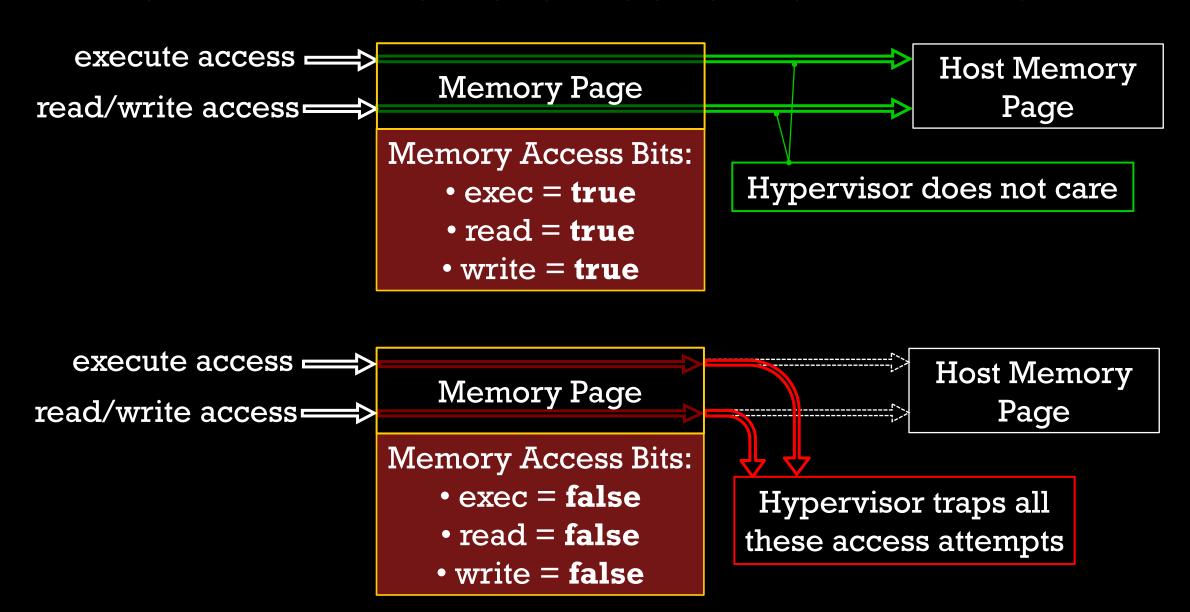
INSIDE EPT PAGING STRUCTURES. EPT PFN



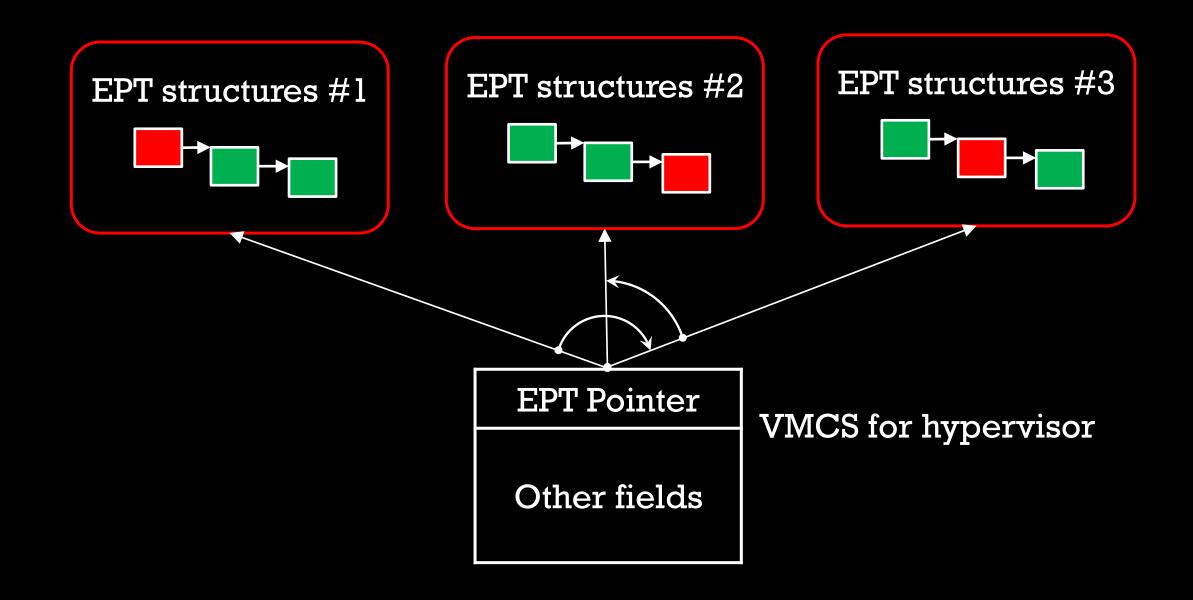
INSIDE EPT PAGING STRUCTURES. EPT PFN

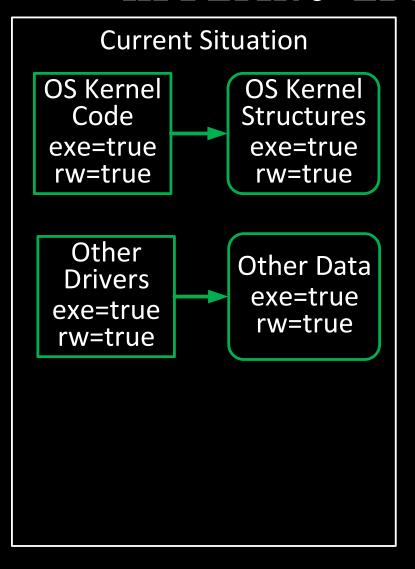


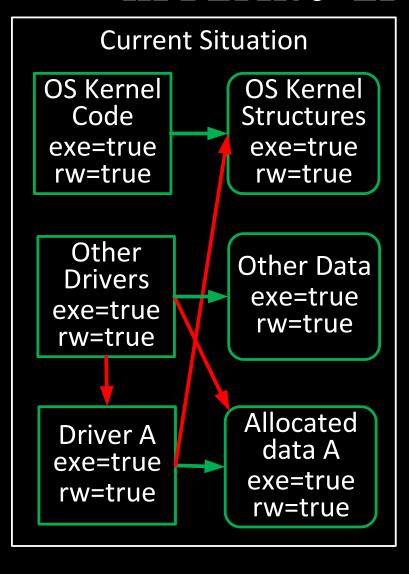
INSIDE EPT PAGING STRUCTURES. EPT BITS

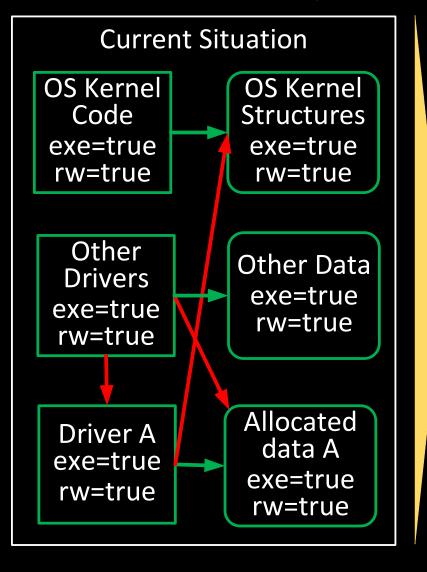


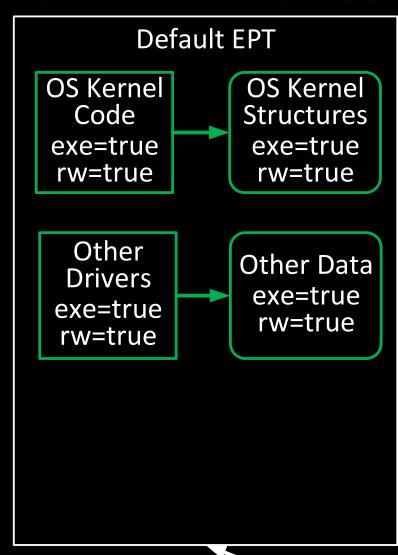
INSIDE EPT PAGING STRUCTURES

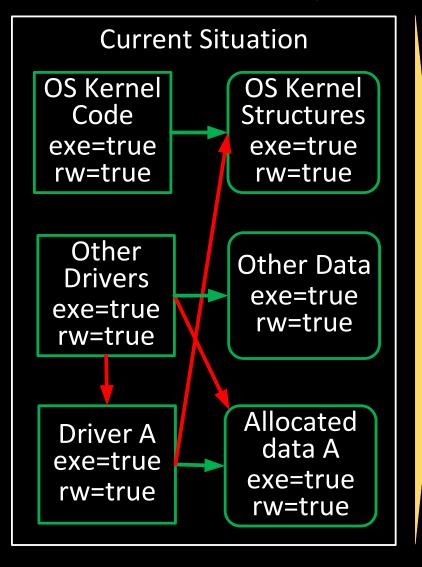


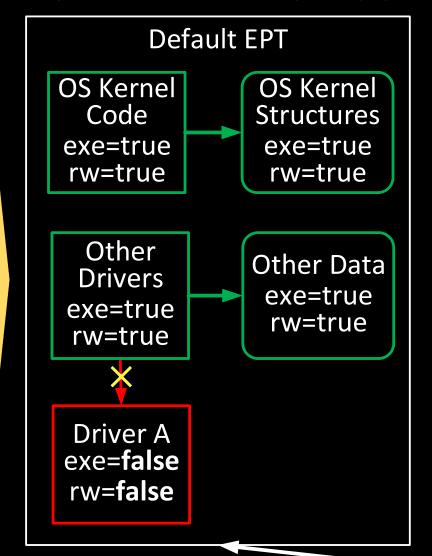


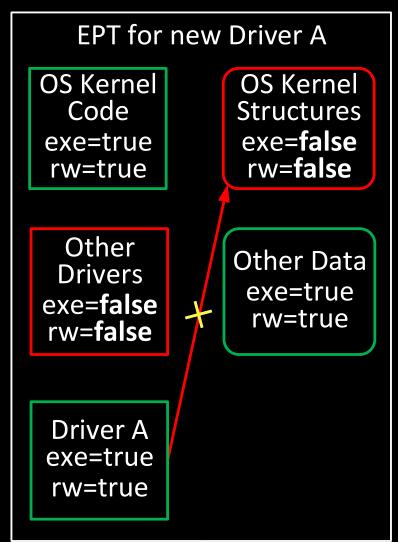


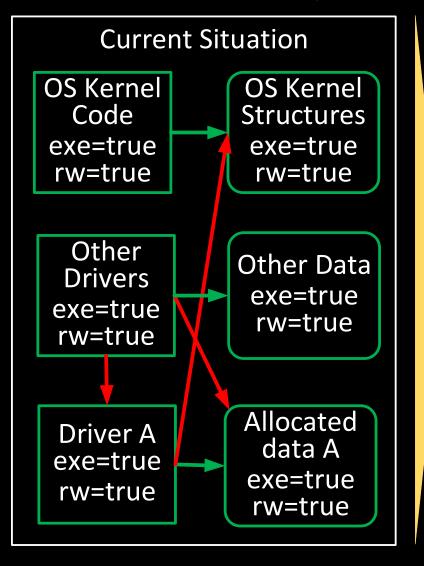


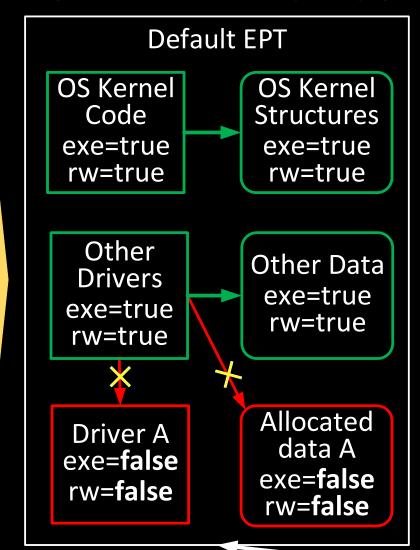


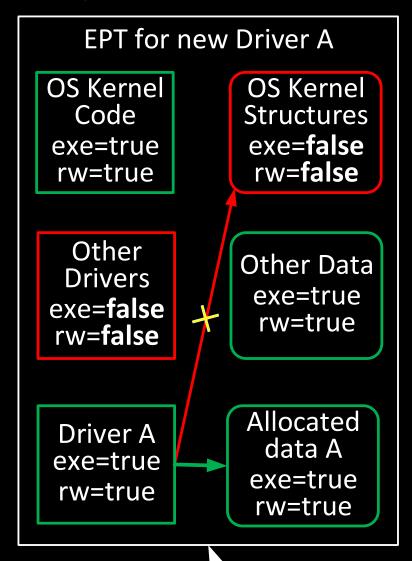






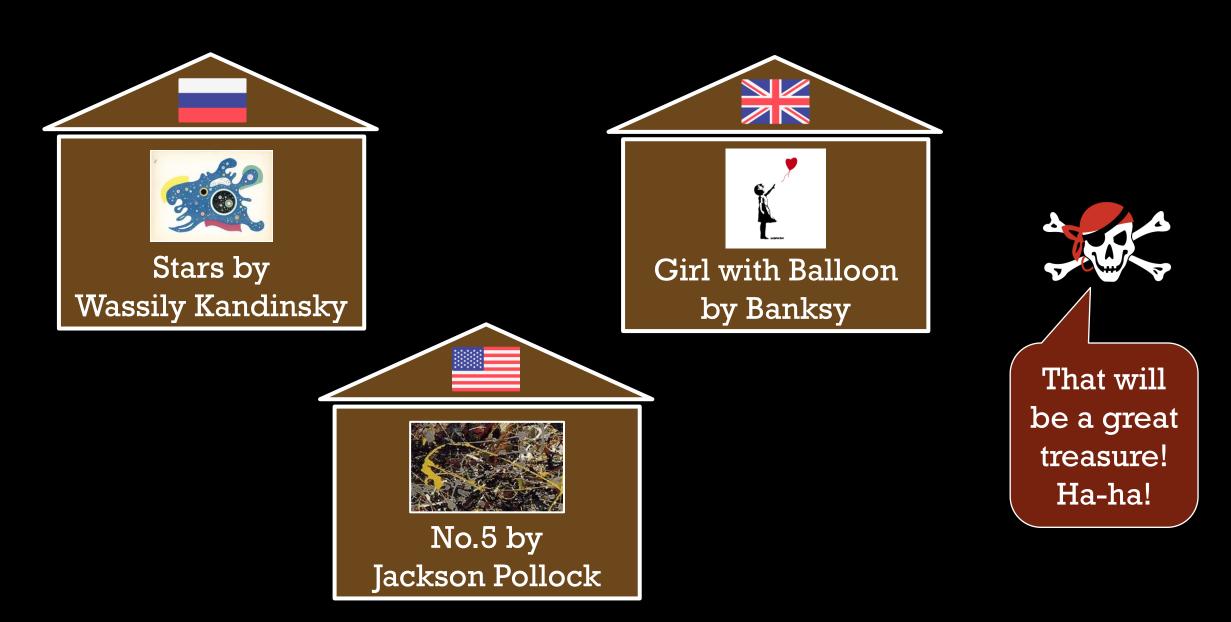




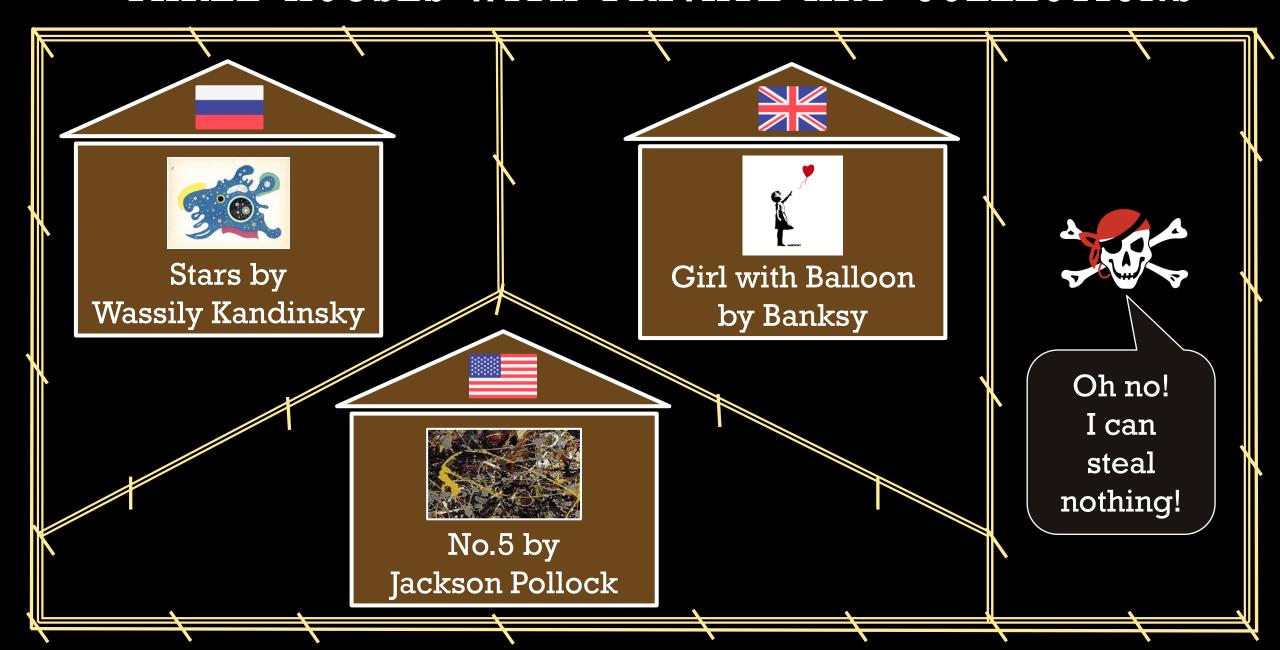


EPT pointer

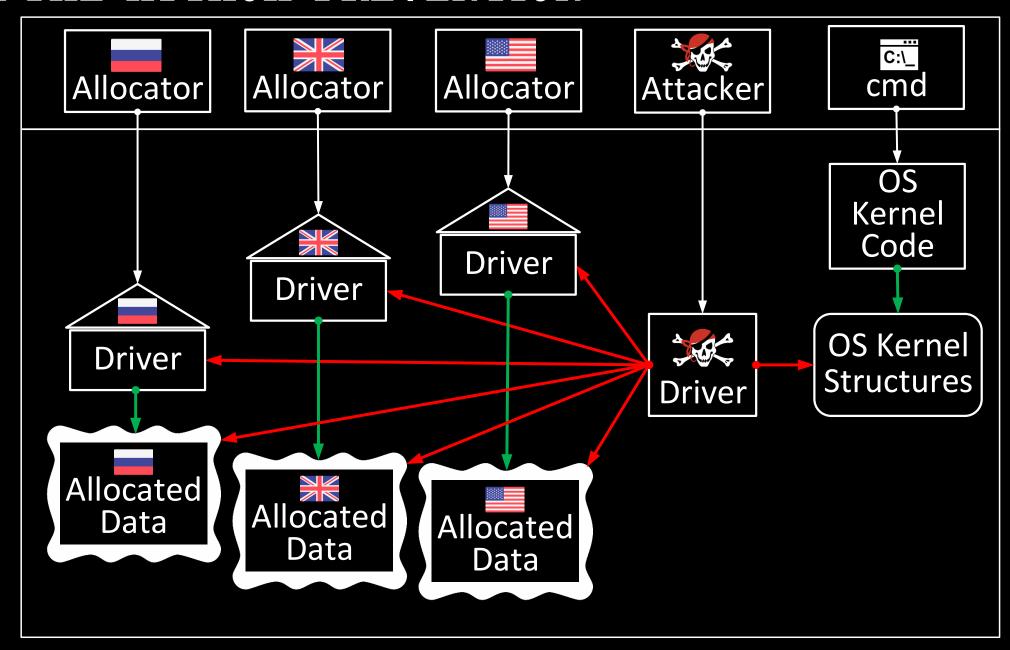
THREE HOUSES WITH PRIVATE ART COLLECTIONS



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DEMO: THE ATTACK PREVENTION

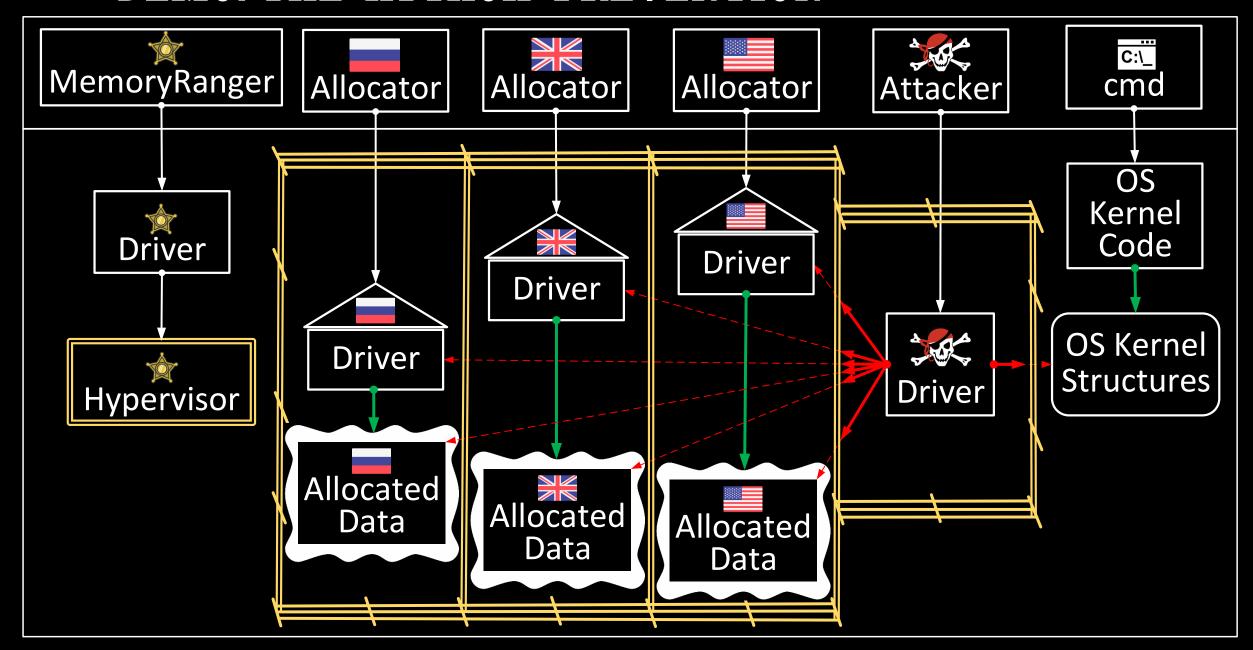


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DEMO: THE ATTACK PREVENTION



MEMORY RANGER: PRINCIPLE OF LEAST PRIVILEGE

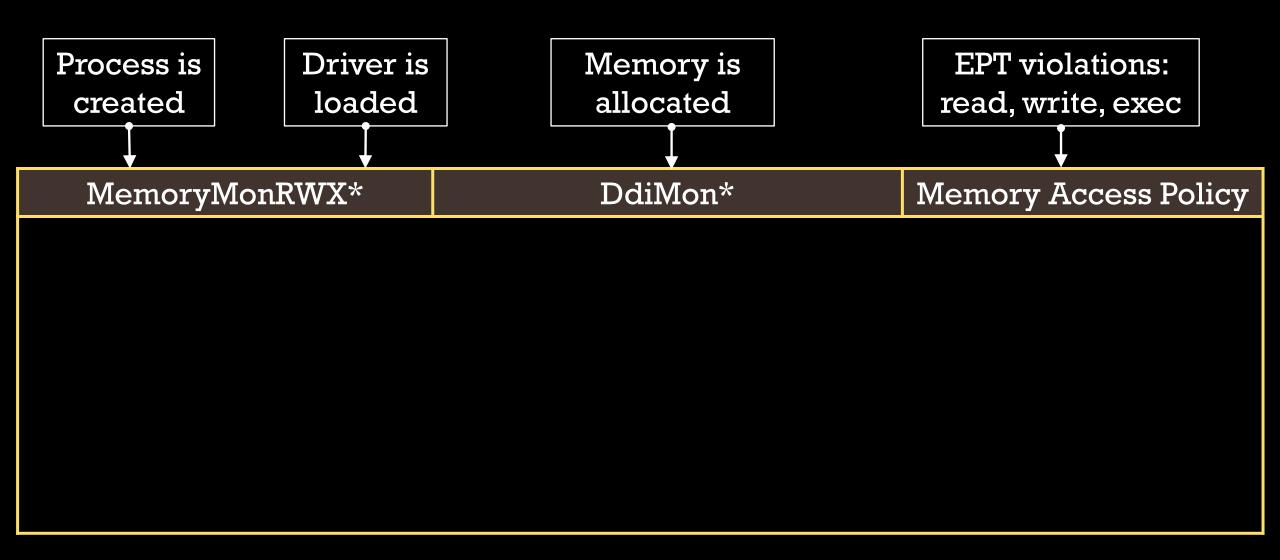
Kernel-	Drivers Code					
mode drivers						
	✓					
		✓				
			~			
				~		
OS kernel	~	✓	>	✓		

MEMORY RANGER: PRINCIPLE OF LEAST PRIVILEGE

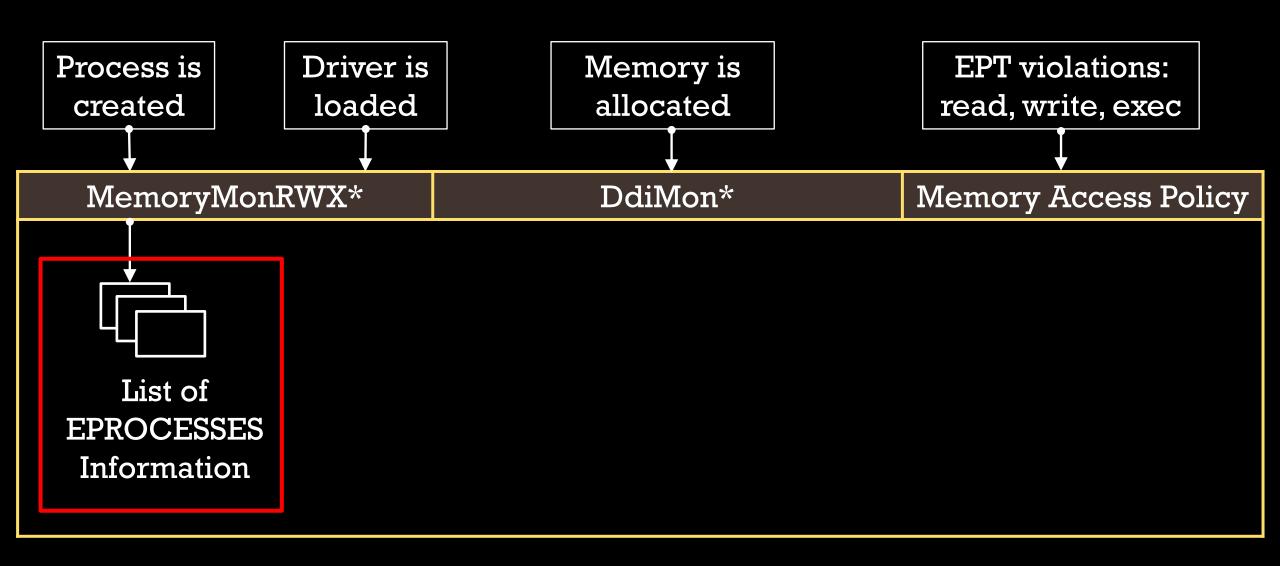
Kernel- mode drivers	Drivers Code			Allocated Memory Data				
								EPROCESS structures
	✓				✓			
		✓				~		
			~				✓	
				~				
OS kernel	✓	✓	~	>	~	✓	✓	✓

MEMORY RANGER: PRINCIPLE OF LEAST PRIVILEGE

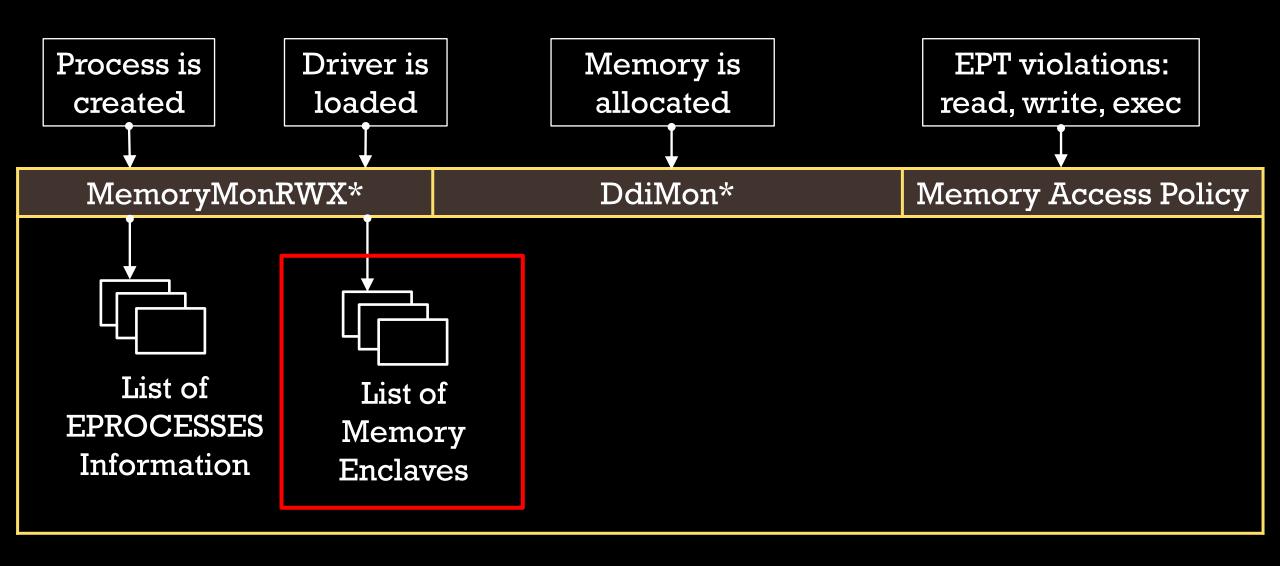
Kernel- mode drivers	Drivers Code				Allocated Memory Data			
								EPROCESS structures
	~	×	X	×	~	X	×	×
	×	✓	×	×	×	~	×	×
	×	×	~	×	×	X	~	×
	×	×	×	~	×	X	×	×
OS kernel	✓	✓	~	~	✓	✓	✓	✓



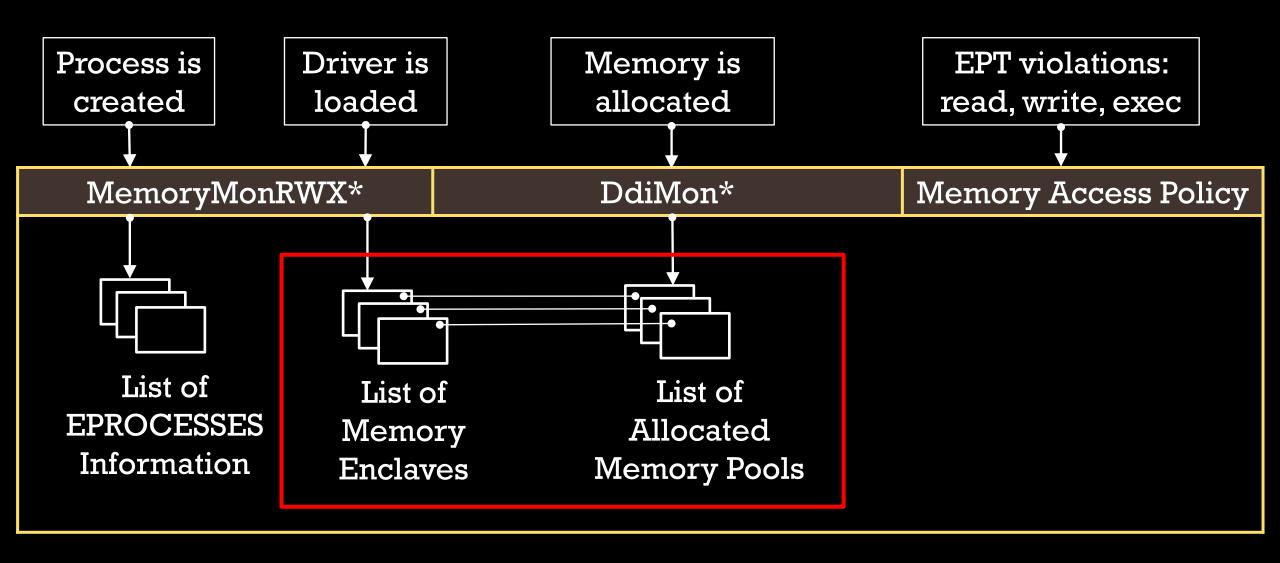
^{*} by Satoshi Tanda, @standa_t, https://github.com/tandasat



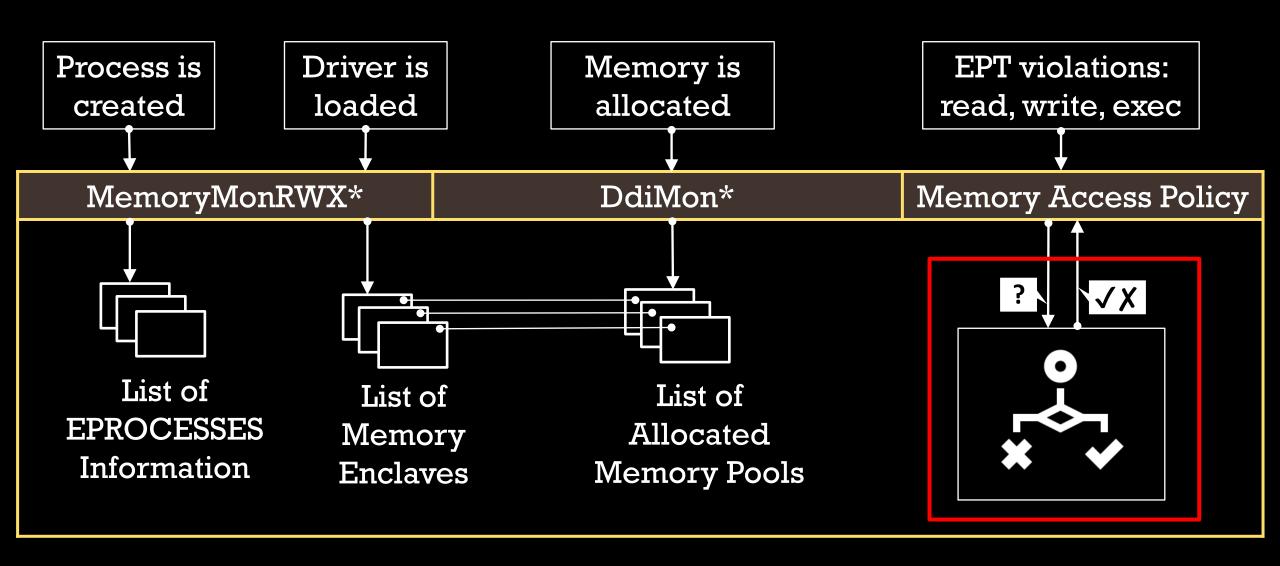
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```
switch (exit_reason){
      case (execute_violation):
            change_ept();
            break;
      case (read_violation|| write_violation):
            if (access_legal()==false){
                   set_pte(pfn, read | write, fake_page);
                   set_monitor_trap_flag();
                   break;
      case (monitor_trap_flag):
            set_pte(pfn, no_access, original_page);
            clear_monitor_trap_flag();
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```

HOW TO PROTECT YOUR DATA IN MEMORY?

- 1. Callback creating a list of protected objects
 - Add objects' addresses & sizes to the list
 - Restrict memory access for objects memory via EPT

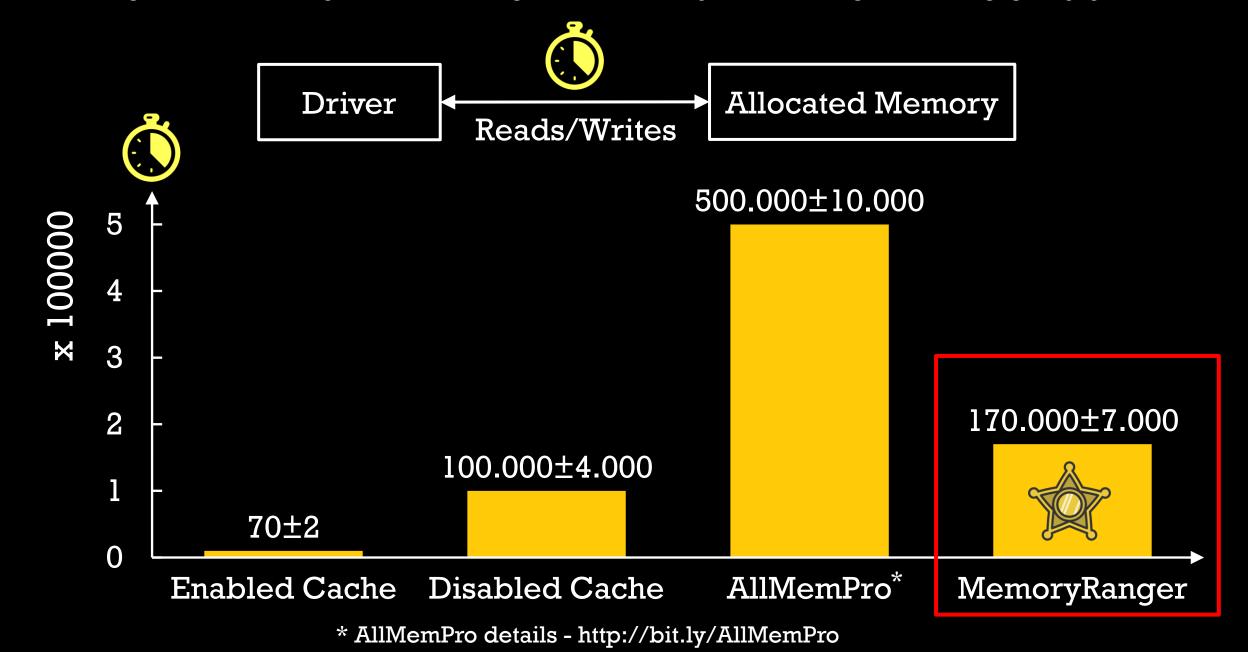
- 2. EPT dispatcher processing EPT violations for this data
 - type_of_access read or write
 - guest_ip is the 'source address'
 - fault_va is the 'destination address'
 - Temporary allow access to the data using MTF
 - Redirect access to the fake data using MTF and EPT.PFN

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MEMORY RANGER BENCHMARKS: MEMORY ACCESS TIME



BLACK HAT SOUND BYTES OR CONCLUSION

Kernel-mode memory is out of control

 MemoryRanger isolates drivers execution by using a specific EPT structure for each driver

MemoryRanger seems to prevent Spectre and Meltdown CPU attacks:

research is ongoing

Dīvide et Imperā* from Latin divide and rule

Thank you!

igor.korkin@gmail.com Igor Korkin

All the details & my CV are here igorkorkin.blogspot.com









