

APPLYING MEMORY FORENSICS TO ROOTKIT DETECTION

Igor Korkin Ivan Nesterov

CDFSL 2014

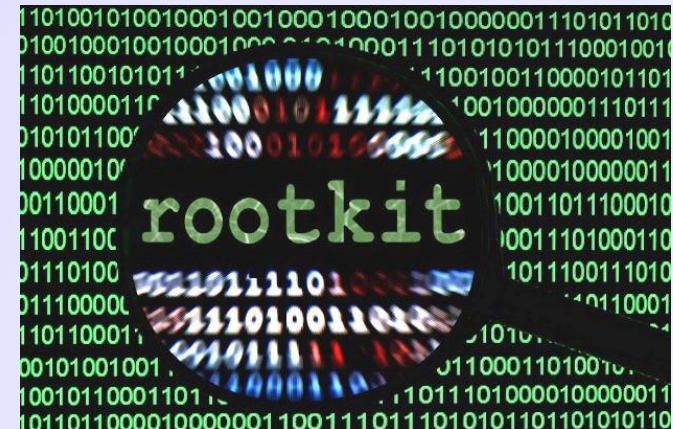
Goals of memory forensics

**Passwords, crypto keys and etc.
revealing software**



Software reverse engineering

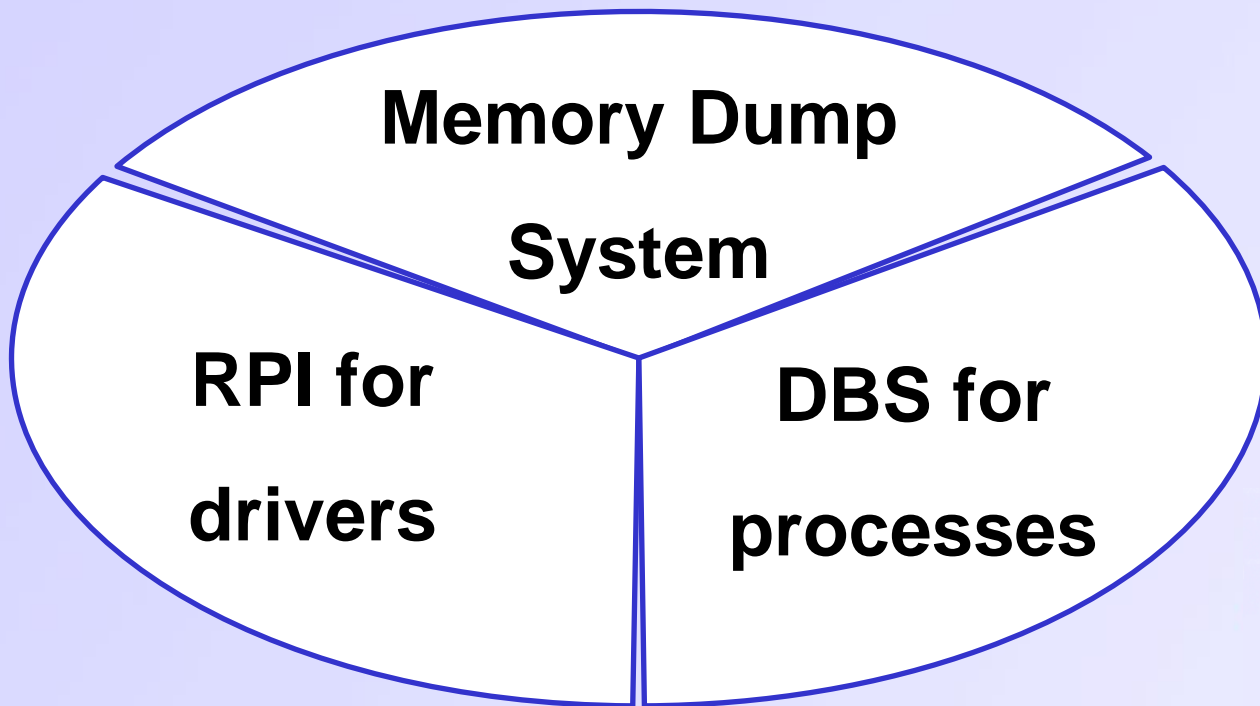
Rootkits analysis & detection



Agenda

1. Review of dump & analysis tools in rootkit conditions

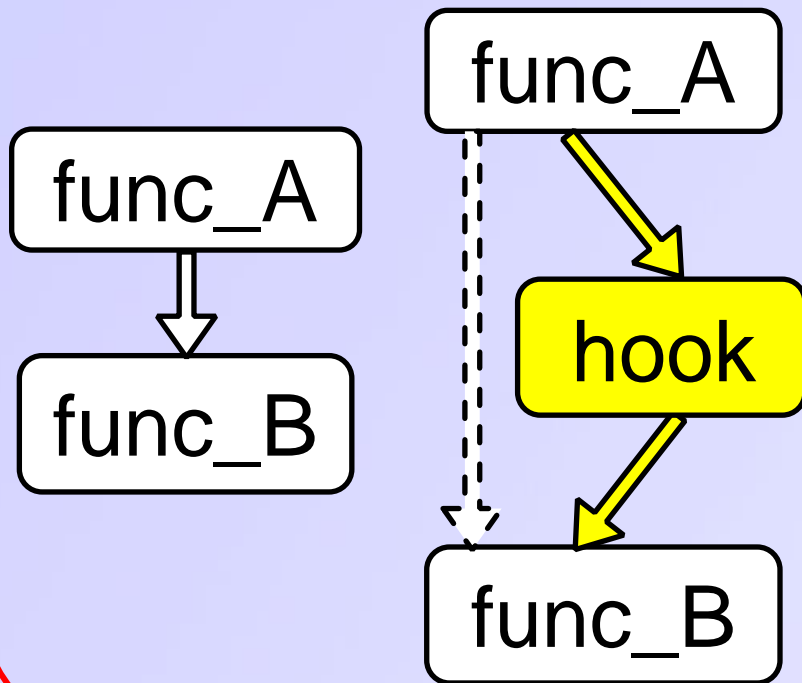
2-3. MASHKA – Malware Analysis System for Hidden Knotty Anomalies:



Review of rootkits techniques

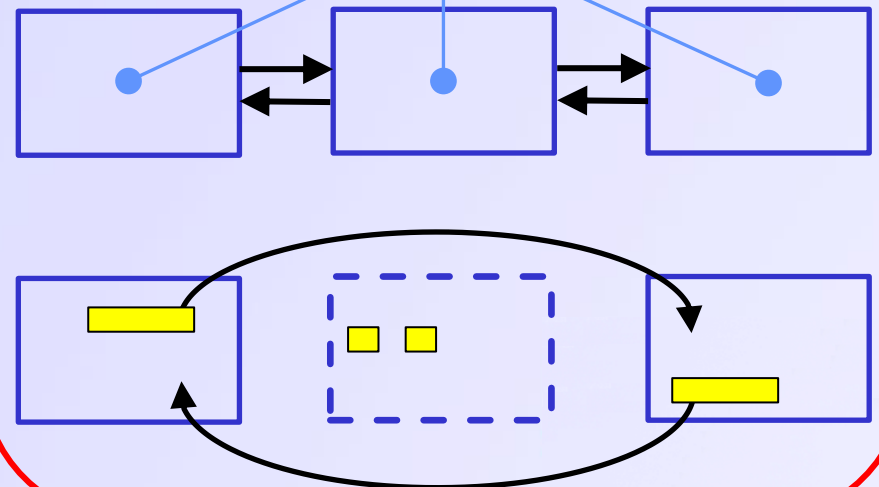
Rootkits techniques – malware hiding from OS & AV

function hooking

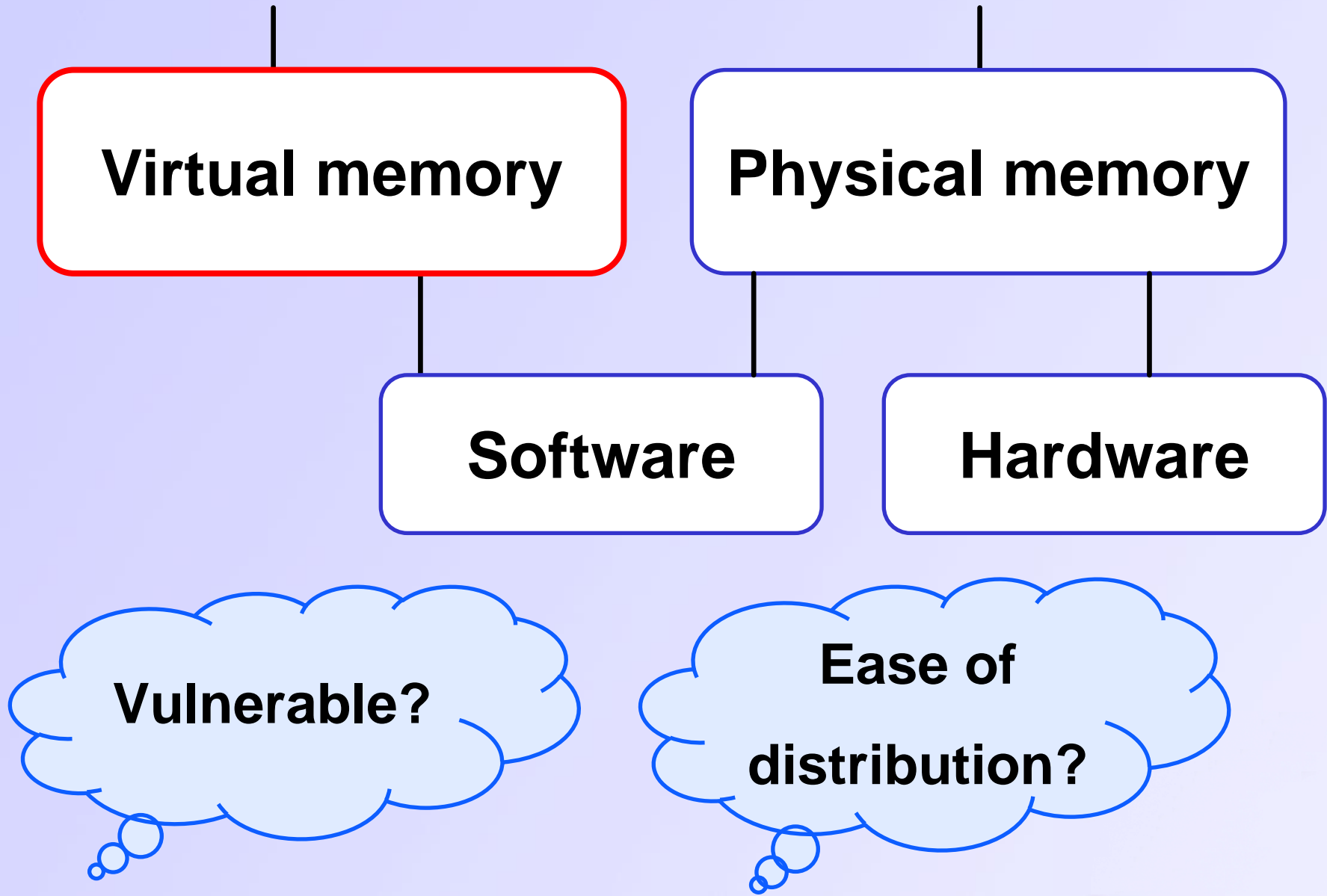


object manipulation (byte modification)

EPROCESS structures



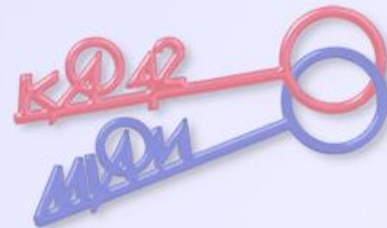
Dump approaches classification



Dump approaches are either vulnerable or non applicable in enterprises

	Hooking resilience	Ease of distribution
Software	-	+
Hardware	+	-

Why are software approaches vulnerable?



Details of dump & analysis tools

Typical dump & analysis tool

Memory mapping
routines

J.Stuttgen, M.Cohen (`13)

Hook

ZwWriteFile or analogue

L.Milkovic (`12)

Hook

Analysis of kernel OS
structures

T.Haruyama, H.Suzuki (`12)

Byte Modification



**What can we do under these
circumstances?**



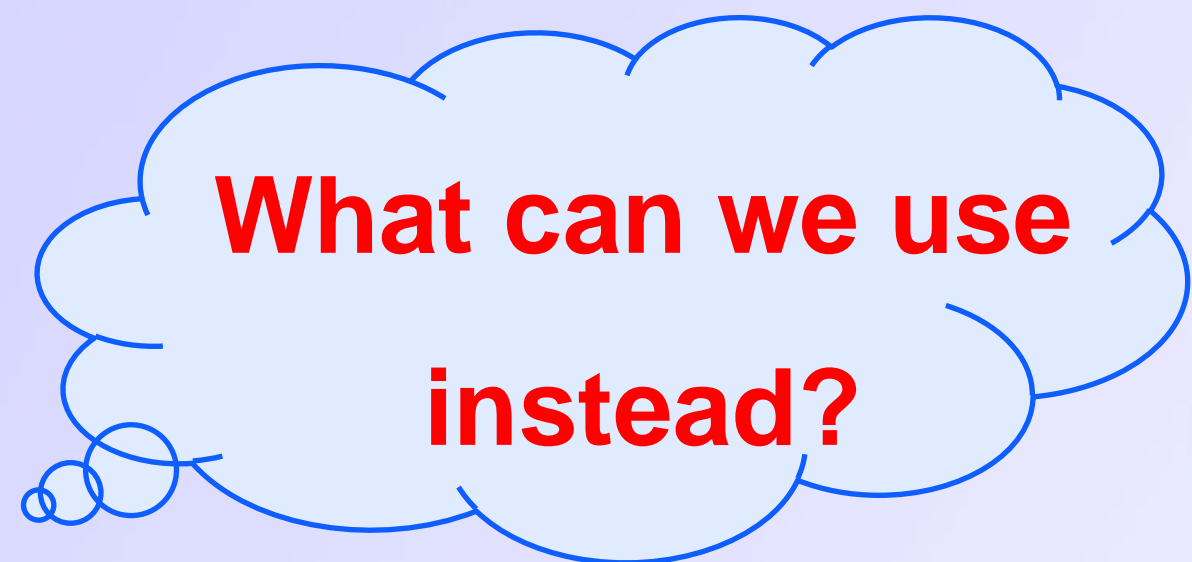
**What can we do under these
circumstances?**

Let's omit the functions!



**What can we do under these
circumstances?**

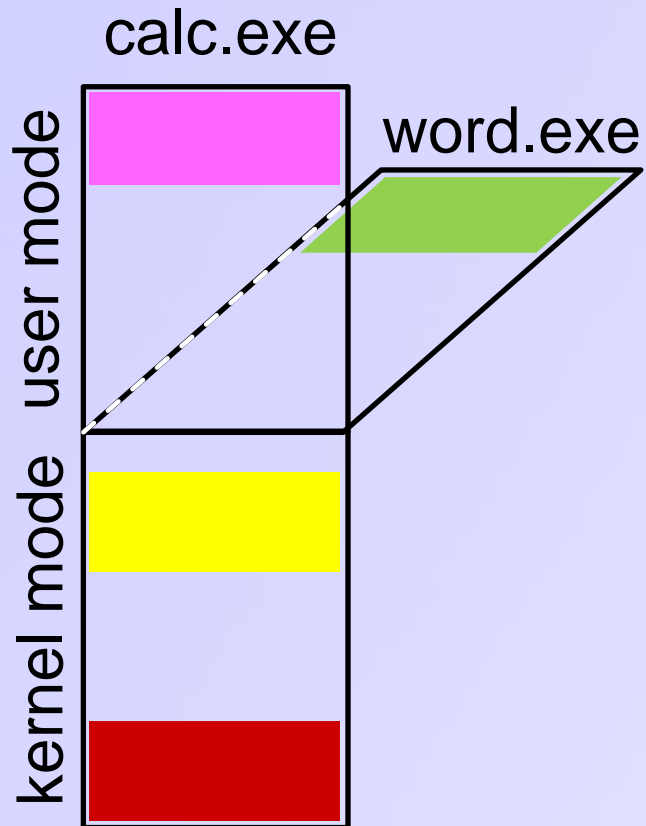
Let's omit the functions!



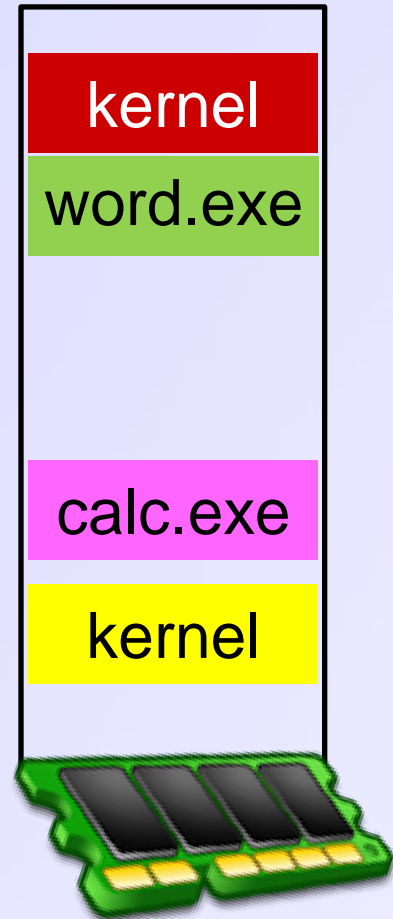
**What can we use
instead?**

Virtual and Physical memory

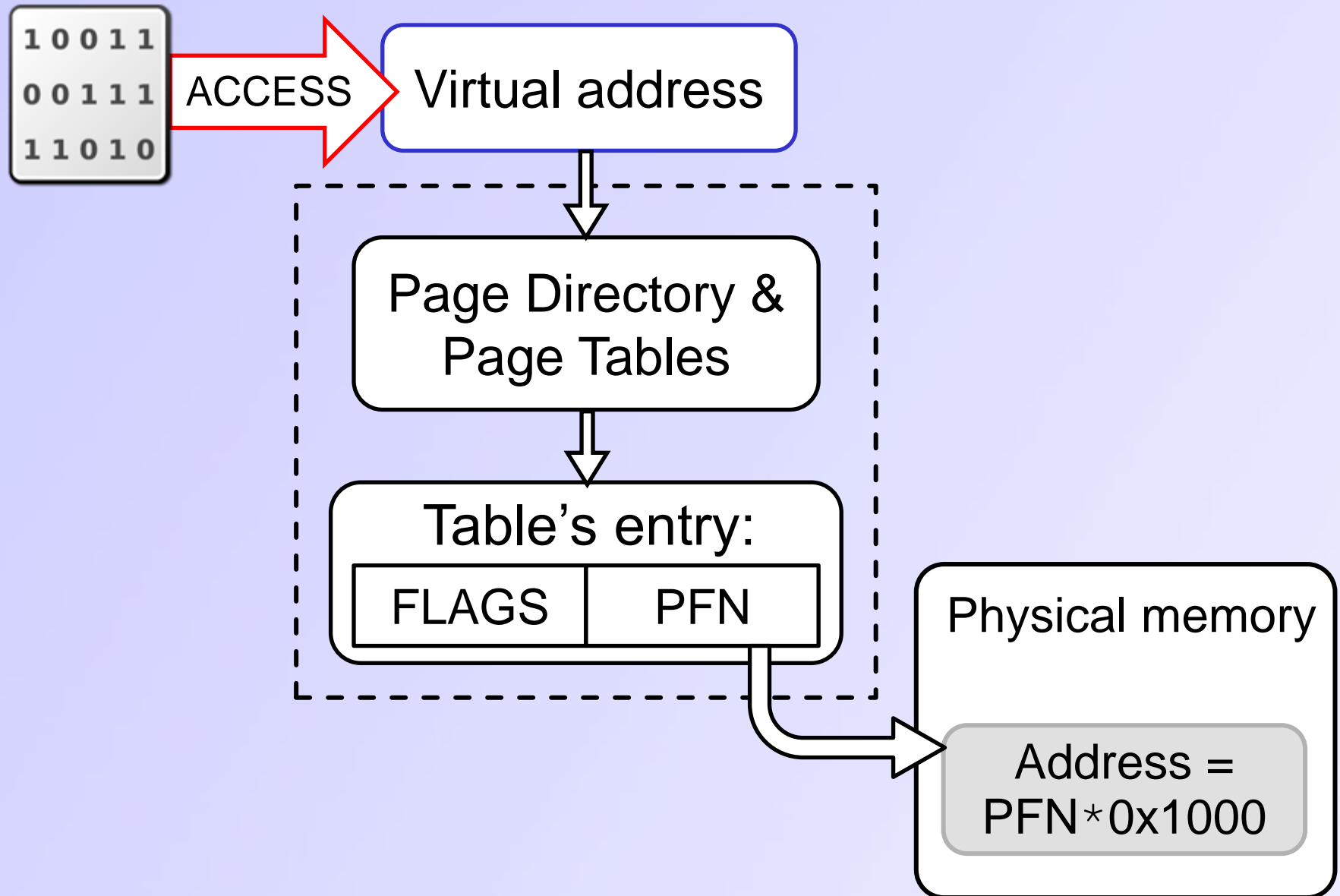
Virtual memory



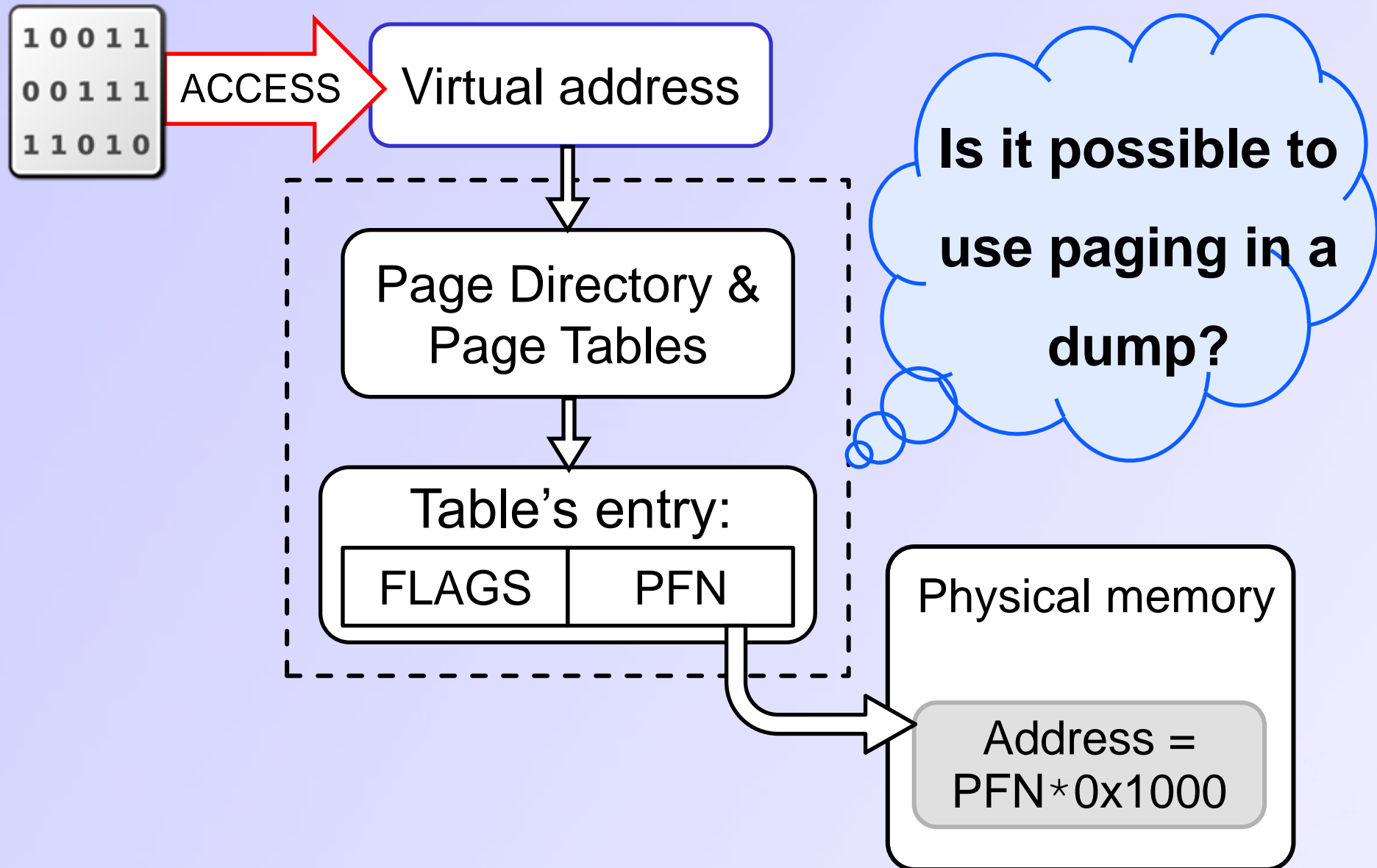
Physical memory



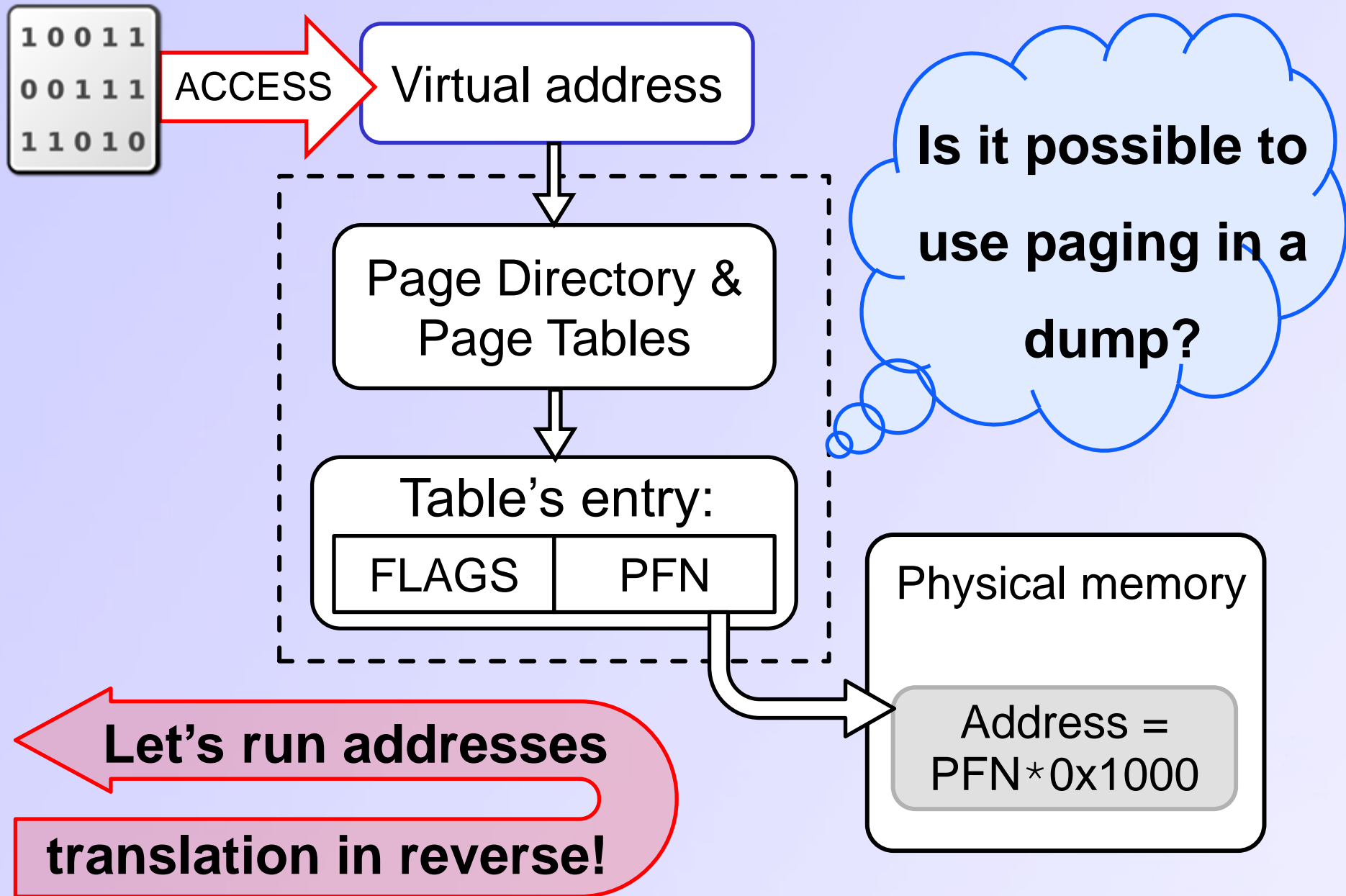
How does addresses translation work?



How does addresses translation work?



How does addresses translation work?



MASHKA's memory dump algorithm

Page Directory

<i>i</i>	others	P	PS
5	BE 3C	0	
6			
7			
...			

Go to next entry

MASHKA's memory dump algorithm

Page Directory

<i>i</i>	others	P	PS
5	BE 3C	0	
6	BE 4C	1	1
7			
...			

Go to next entry

Save memory page
(4 Mb or 2Mb) by *i*

MASHKA's memory dump algorithm

Page Directory

<i>i</i>	others	P	PS
5	BE 3C	0	
6	BE 4C	1	1
7	BE FF	1	0
...			

Go to next entry

Save memory page
(4 Mb or 2Mb) by *i*

Go to Page Table

Page Table

<i>j</i>	others	P
0		
1		
...		

MASHKA's memory dump algorithm

Page Directory

<i>i</i>	others	P	PS
5	BE 3C	0	
6	BE 4C	1	1
7	BE FF	1	0
...			

Go to next entry

Save memory page
(4 Mb or 2Mb) by *i*

Go to Page Table

Page Table

<i>j</i>	others	P
0	BF 00	0
1		
...		

Go to next entry

MASHKA's memory dump algorithm

Page Directory

<i>i</i>	others	P	PS
5	BE 3C	0	
6	BE 4C	1	1
7	BE FF	1	0
...			

Go to next entry

Save memory page
(4 Mb or 2Mb) by *i*

Go to Page Table

Page Table

<i>j</i>	others	P
0	BF 00	0
1	BF 01	1
...		

Go to next entry

Save memory page
(4 Kb) by *i* & *j*

MASHKA's memory dump algorithm

Page Directory

<i>i</i>	others	P	PS
5	BE 3C	0	
6	BE 4C	1	1
7	BE FF	1	0
...			

Go to next entry

Save memory page
(4 Mb or 2Mb) by *i*

Go to Page Table

Page Table

<i>j</i>	others	P
0	BF 00	0
1	BF 01	1
...		

Go to next entry

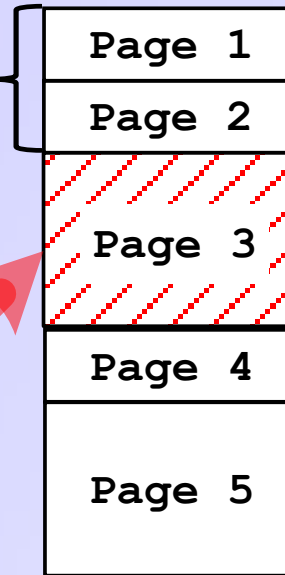
Save memory page
(4 Kb) by *i* & *j*

MASHKA's dump algorithm details

Dump File (300Mb)

DumpOffset_3

Virtual Memory (4GB)



Struct File

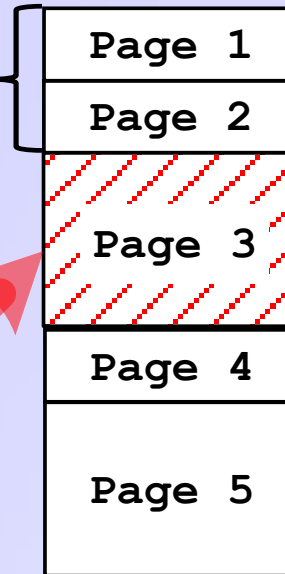
StartAddr_1	FinishAddr_1	DumpOffset_1
StartAddr_2	FinishAddr_2	DumpOffset_2
StartAddr_3	FinishAddr_3	DumpOffset_3
...		
StartAddr_5	FinishAddr_5	DumpOffset_5

MASHKA's dump algorithm details

Dump File (300Mb)

DumpOffset_3

Virtual Memory (4GB)

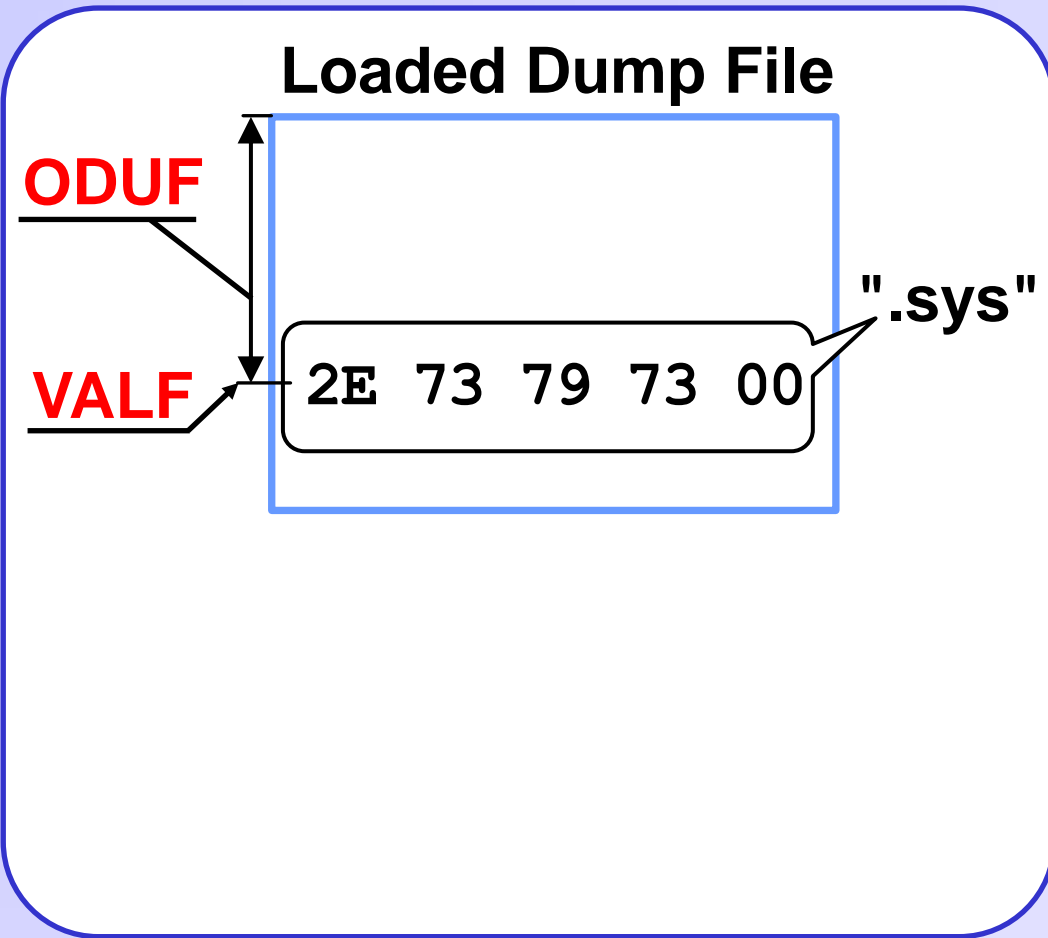


Struct File

StartAddr_1	FinishAddr_1	DumpOffset_1
StartAddr_2	FinishAddr_2	DumpOffset_2
StartAddr_3	FinishAddr_3	DumpOffset_3
...		
StartAddr_5	FinishAddr_5	DumpOffset_5

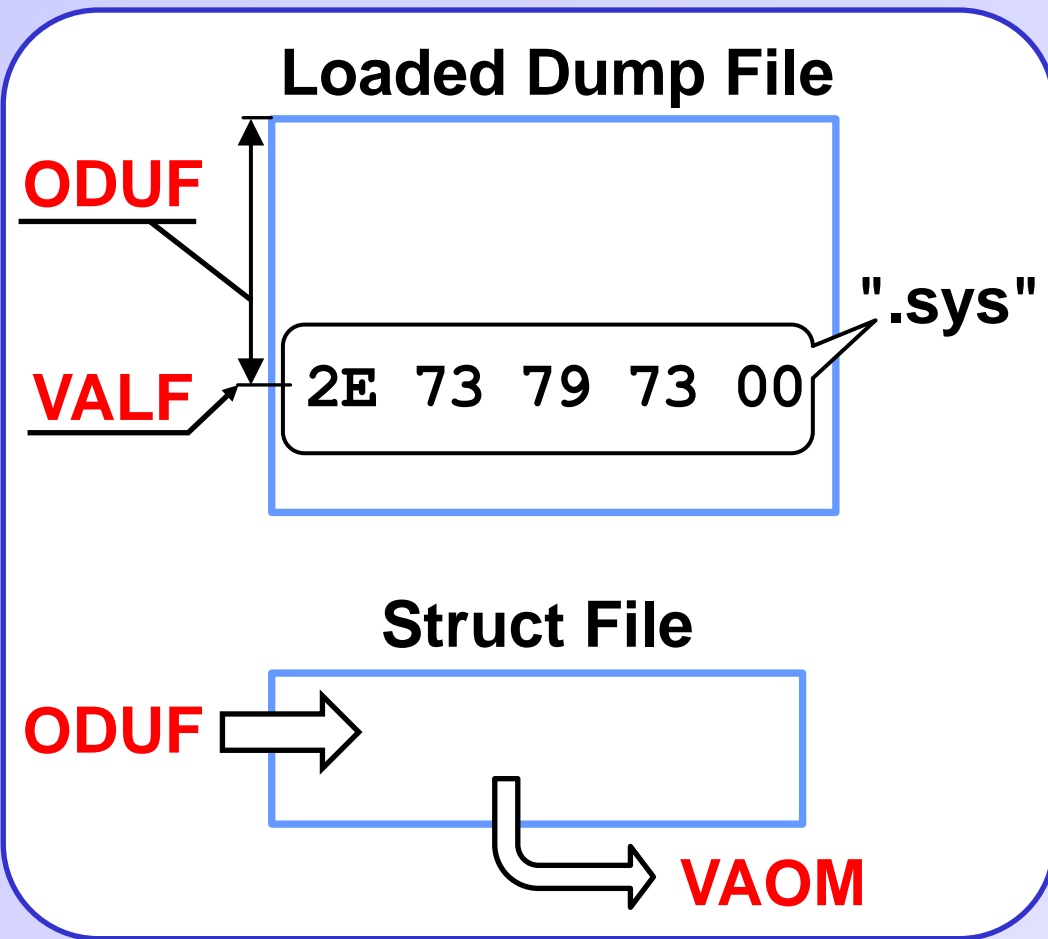
How should new files be used?

MASHKA in memory forensics tasks



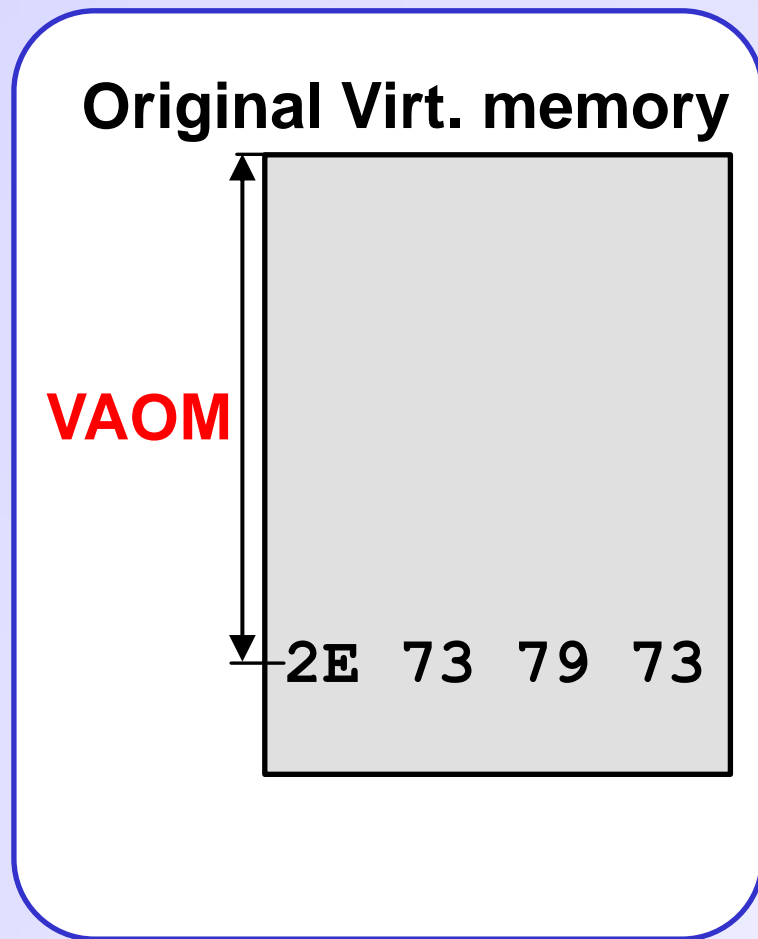
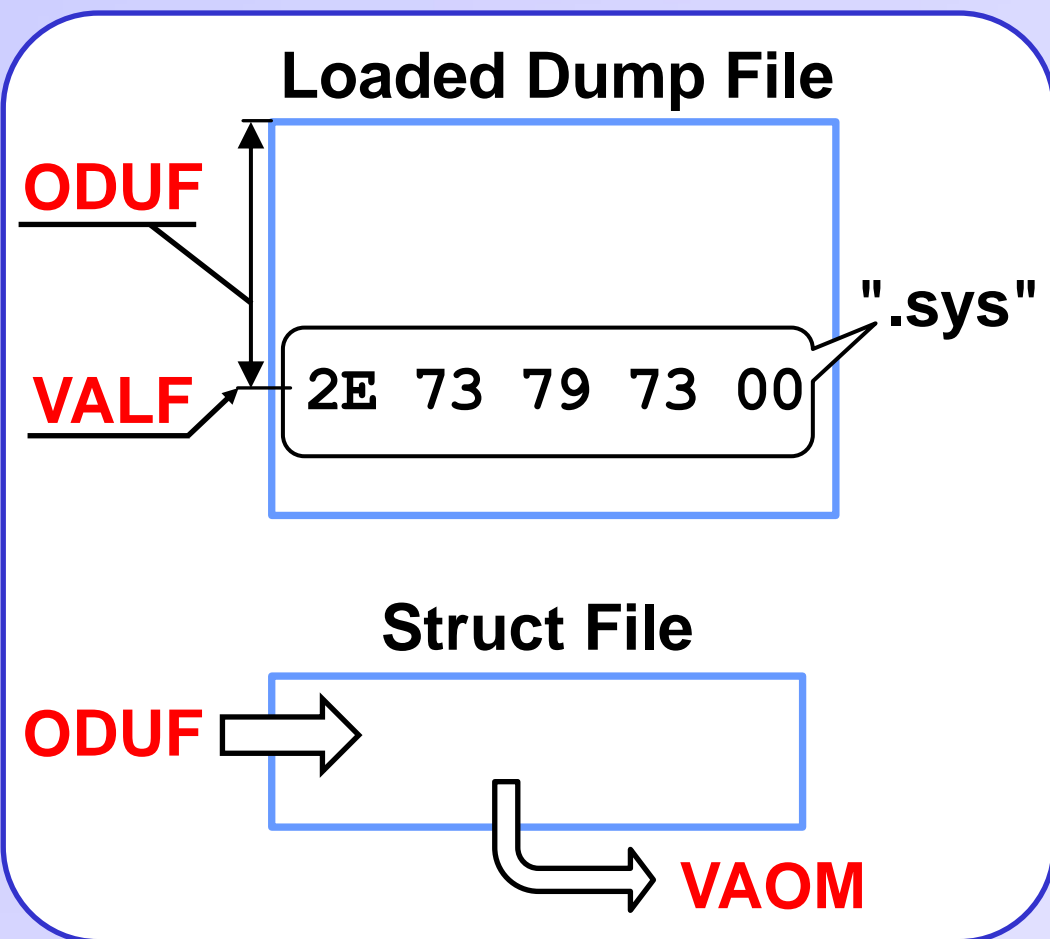
VALF	Virtual Address in the Loaded dump File
ODUF	Offset in DUmP File

MASHKA in memory forensics tasks



VAOM	Virtual Address in the Original virt. Memory
VALF	Virtual Address in the Loaded dump File
ODUF	Offset in DUmP File

MASHKA in memory forensics tasks

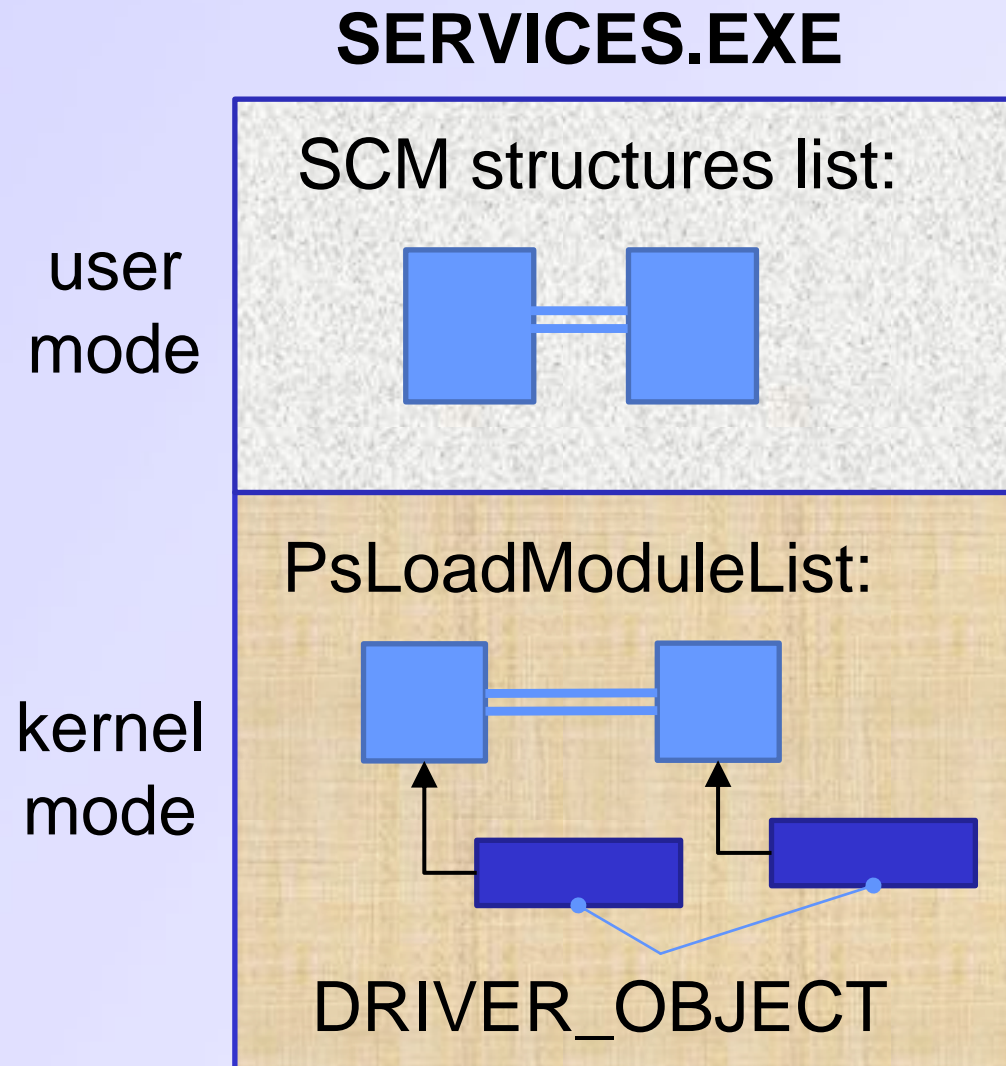


VAOM	Virtual Address in the Original virt. Memory
VALF	Virtual Address in the Loaded dump File
ODUF	Offset in DUmP File

How is VAOM etc used?



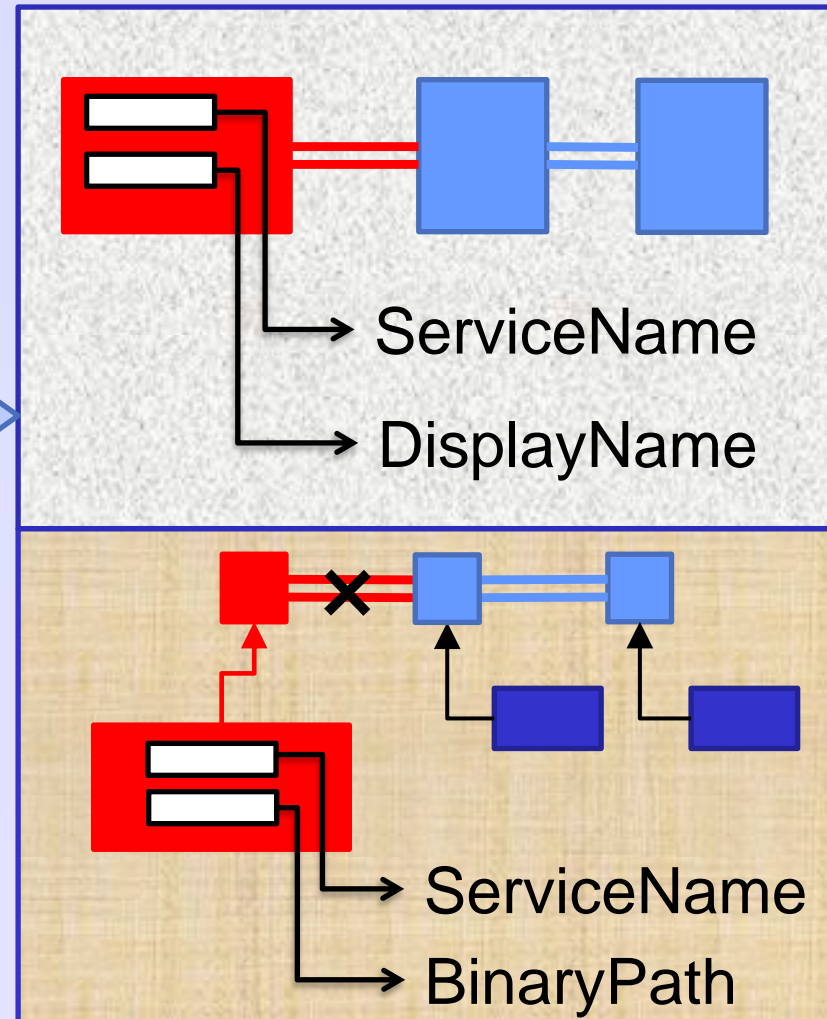
Use MASHKA in drivers forensics



Use MASHKA in drivers forensics

CreateService(ServiceName, DisplayName, BinaryPath,...)

SCM structure,
DRIVER_OBJECT and
others will be added

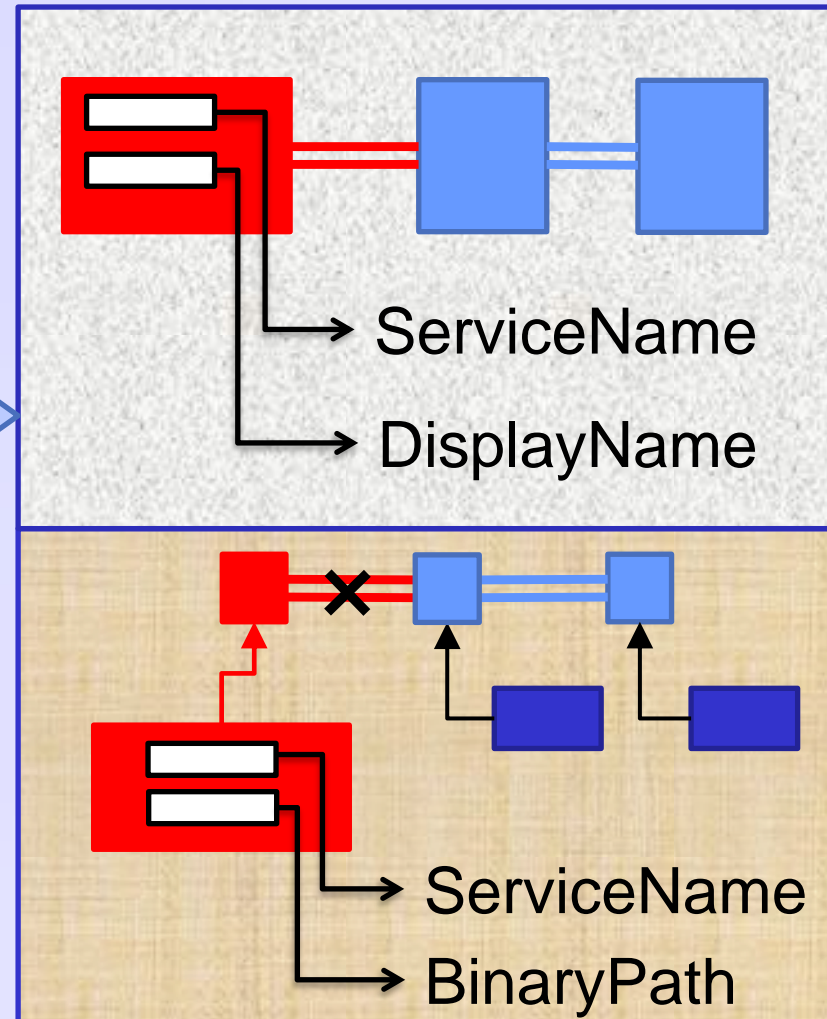


Use MASHKA in drivers forensics

CreateService(ServiceName, DisplayName, BinaryPath,...)

SCM structure,
DRIVER_OBJECT and
others will be added

ServiceName > VAOMs of 'SN'



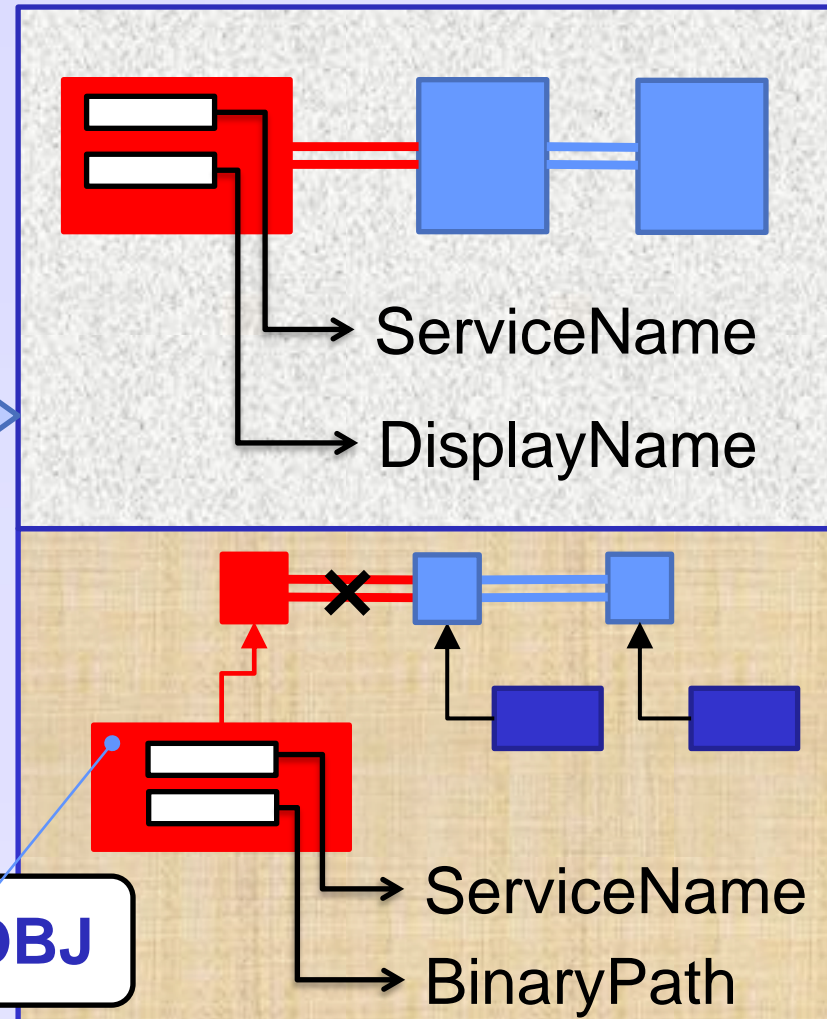
Use MASHKA in drivers forensics

CreateService(ServiceName, DisplayName, BinaryPath,...)

SCM structure,
DRIVER_OBJECT and
others will be added

ServiceName > VAOMs of 'SN'

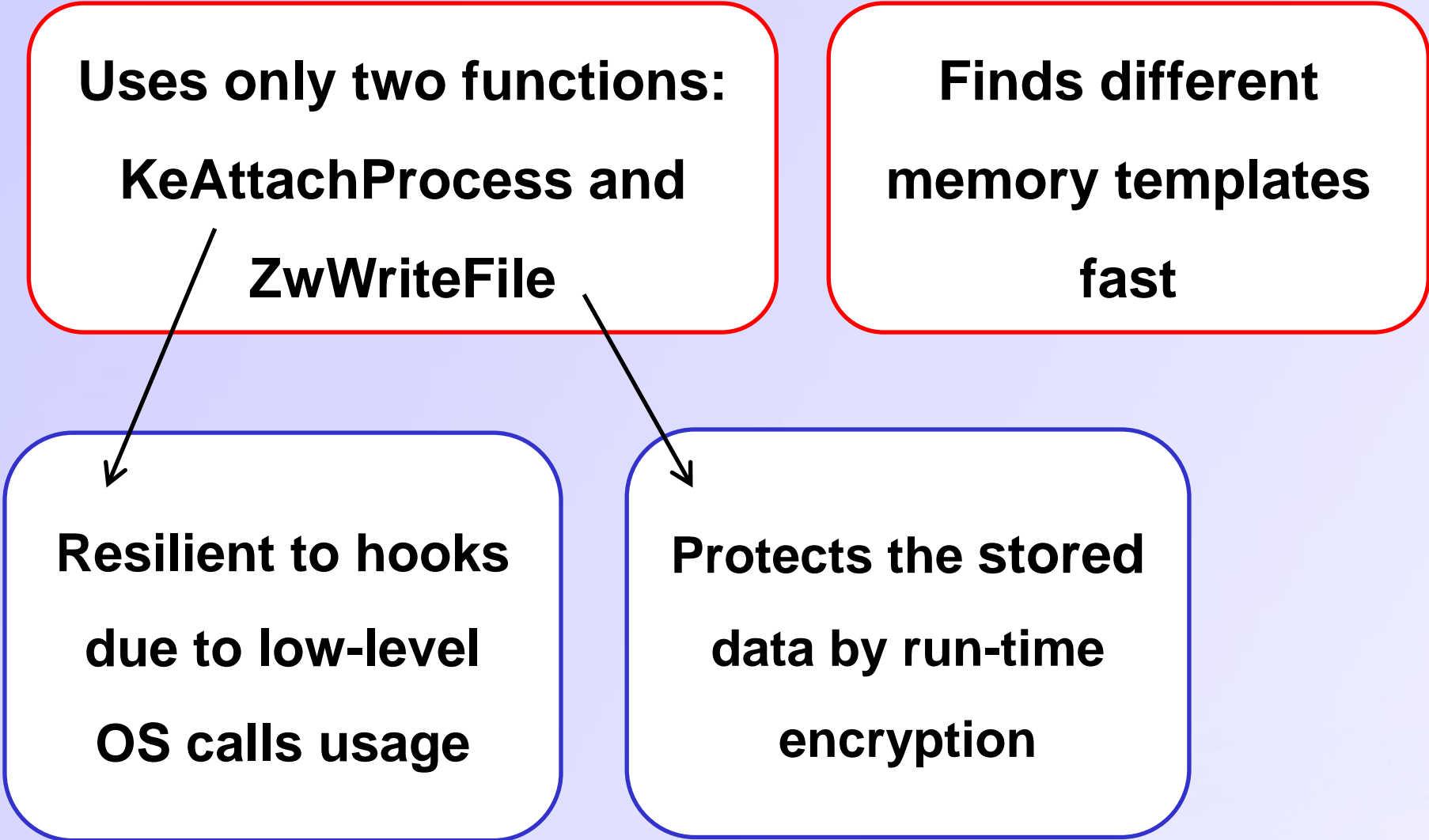
VAOMs of 'SN' > VAOM of DRV_OBJ



Advantages of MASHKA

**Uses only two functions:
KeAttachProcess and
ZwWriteFile**

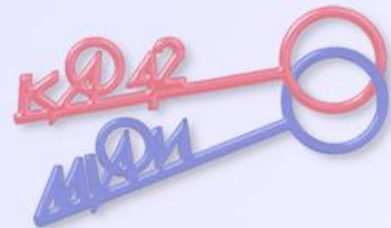
**Finds different
memory templates
fast**



**Resilient to hooks
due to low-level
OS calls usage**

**Protects the stored
data by run-time
encryption**

How to apply MASHKA to processes detection?



OS processes list handling

how can the
process be
hidden?

ZwQuerySystemInformation
hooking

or

PsActiveProcessList
modifying

How to detect a hidden process?

Process detection approaches review

Heuristic
analyzer

- hooking functions such as SwapContext or KiFastCallEntry

Object
structure
lists

- a processes' list from CSRSS.EXE
- a processes handle table list

Static
signature
scans

- static signatures by Schuster ('07)
- robust signatures by Dolan-Gavitt ('09)
- structures location by Grizzard ('10)

Process detection approaches review

Heuristic
analyzer

- hooking functions such as SwapContext or KiFastCallEntry

Object
structure
lists

- a processes' list from CSRSS.EXE
- a processes handle table list

Static
signature
scans

- static signatures by Schuster ('07)
- robust signatures by Dolan-Gavitt ('09)
- structures location by Grizzard ('10)

Analysis of static signature scan

GMER, PowerTool and XueTr use it

Scan is based on

some EPROCESS fields
values are either known or
exceed the constant,
e.g. 0x8000_0000

Disadvantages

vulnerable to field
modifications

difficult to achieve
portability

Analysis of static signature scan

GMER, PowerTool and XueTr use it

Scan is based on

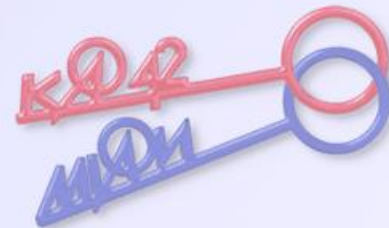
some EPROCESS fields
values are either known or
exceed the constant,
e.g. 0x8000_0000

Disadvantages

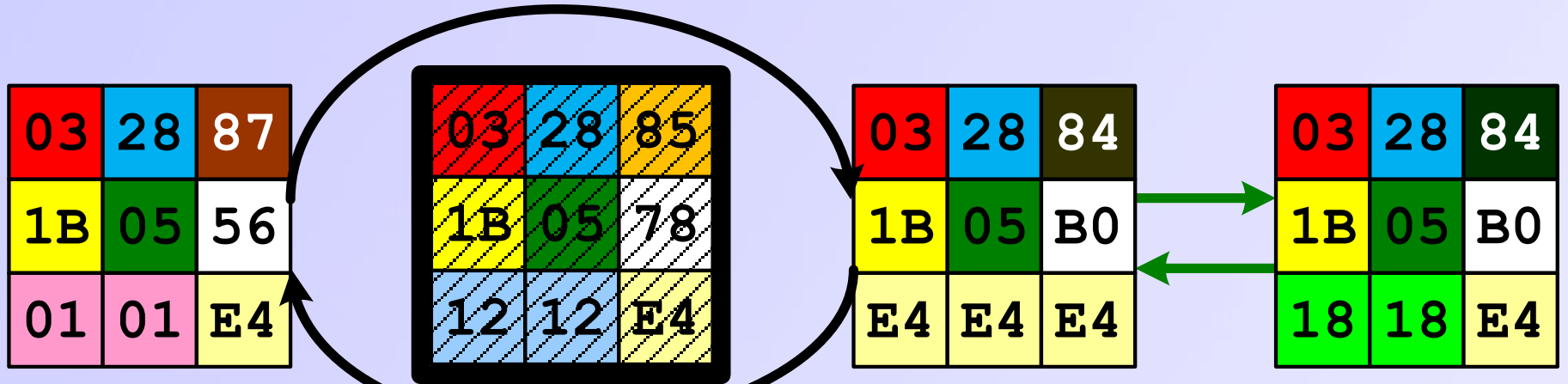
vulnerable to field
modifications

difficult to achieve
portability

**How can we improve
signature scans?**

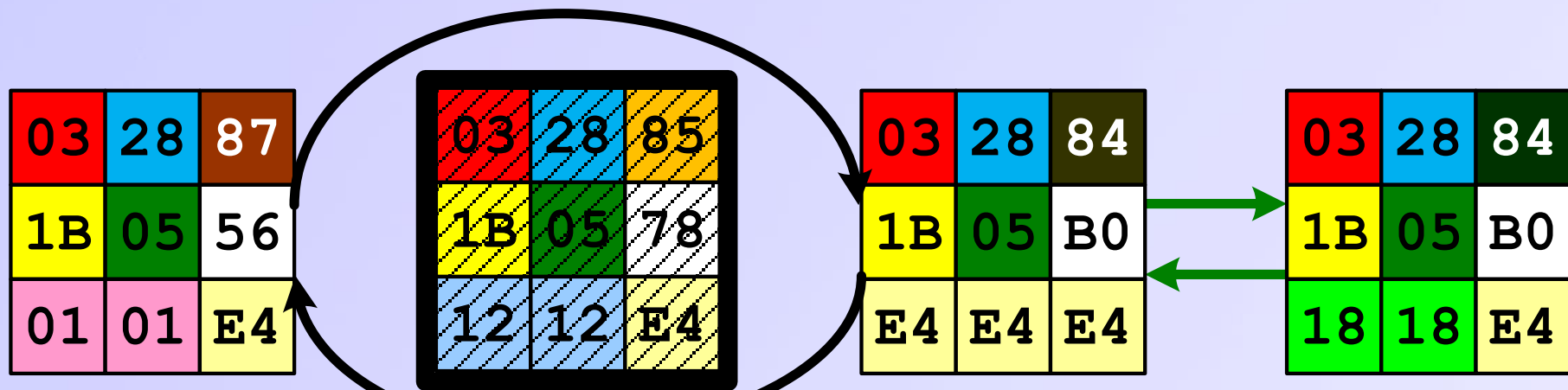


Objects structures typical design



Objects structures

Objects structures typical design



Objects structures

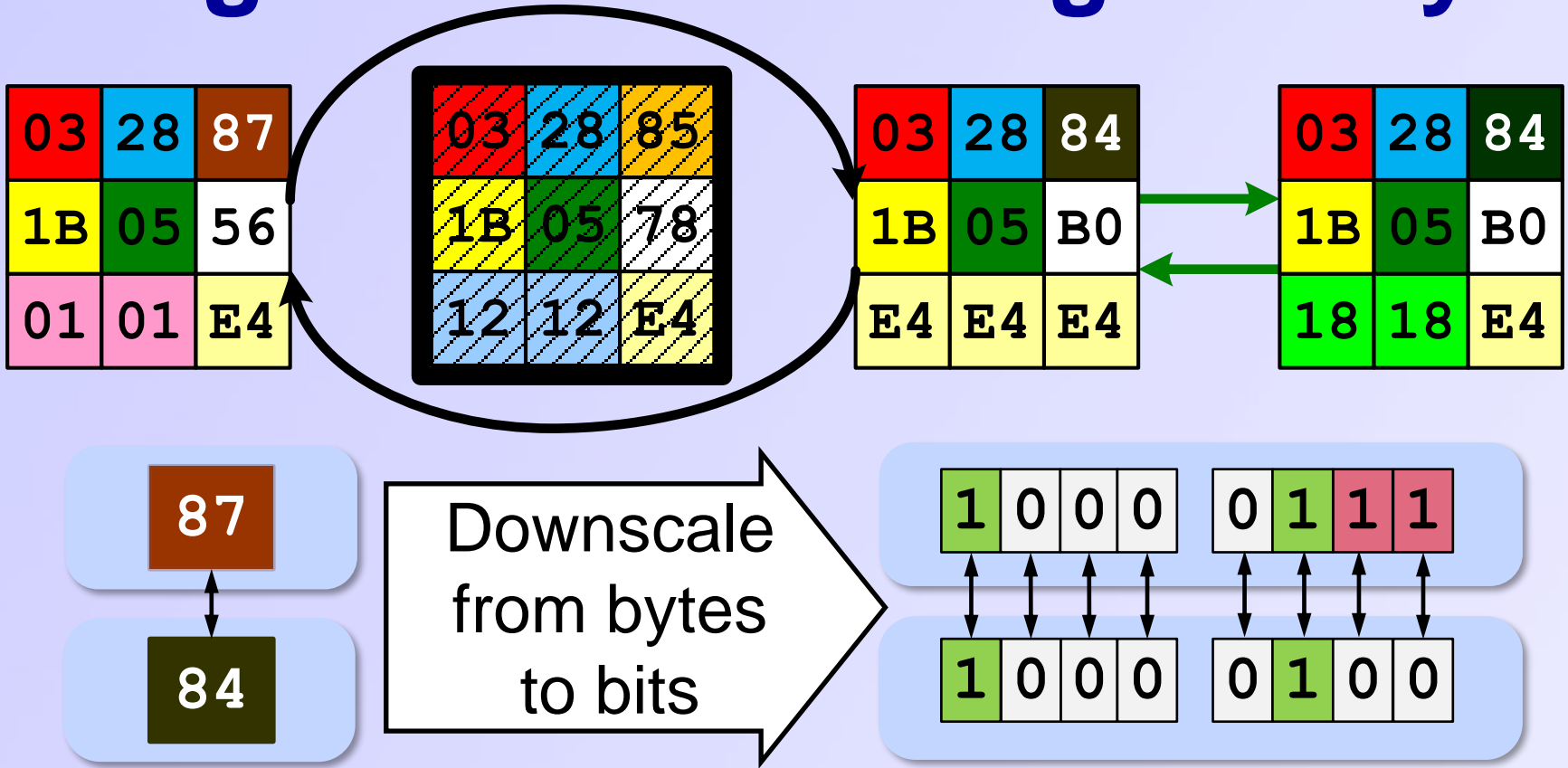
03	28	-
1B	05	-
-	-	E4

**Dynamic Byte
Signature
memory pattern**

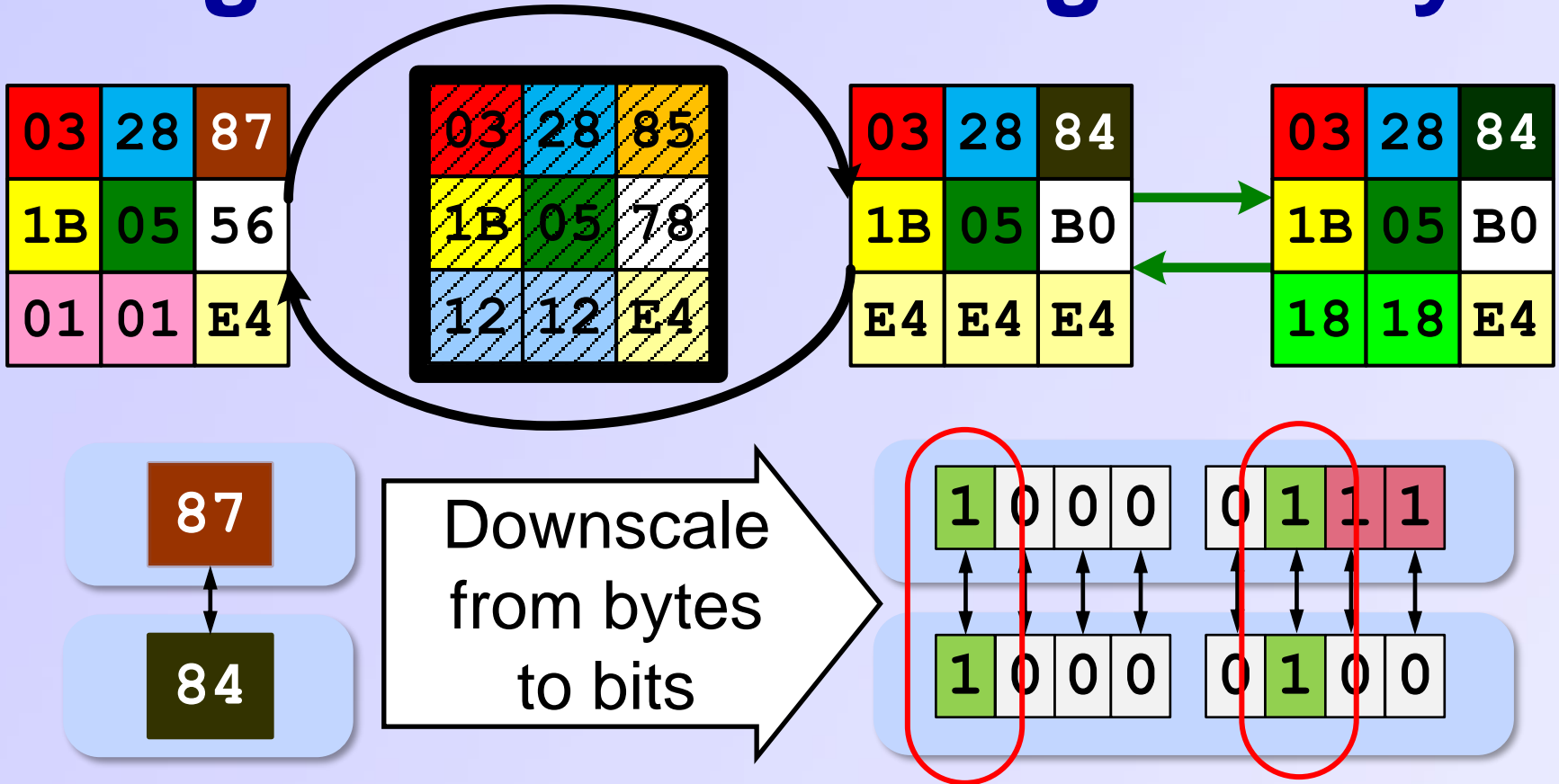
Process detection with Dynamic Byte Signature

- 1. Create Dynamic Byte Signature by using
EPROCESS structures in PsActiveProcessList**
- 2. Use byte to byte DBS search to find all
EPROCESS structures**
- 3. Compare a new list with
NtQuerySystemInformation list**

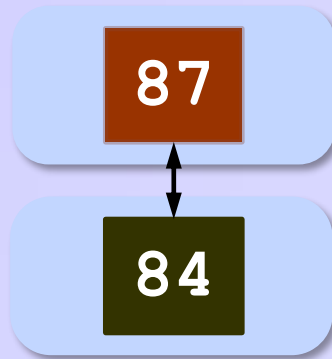
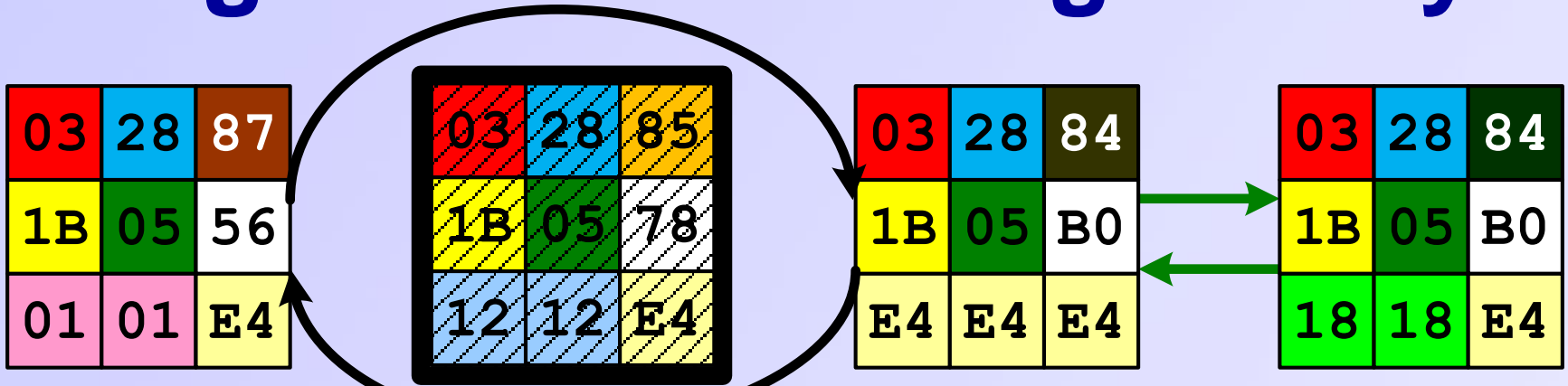
Bit signature = thorough analysis



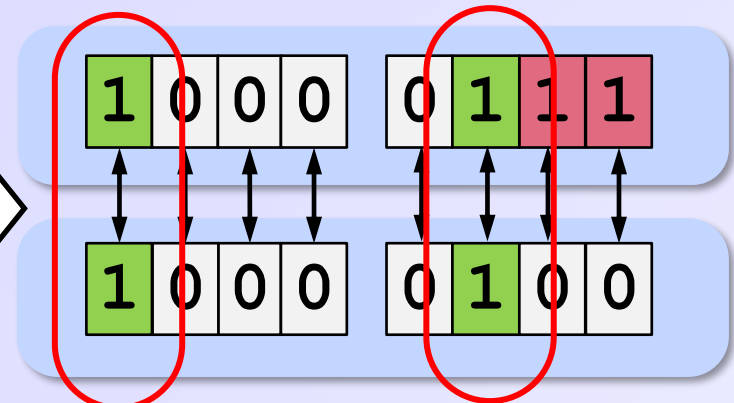
Bit signature = thorough analysis



Bit signature = thorough analysis



Downscale
from bytes
to bits



Dynamic bit signature:

03	28	-
1B	05	-
-	-	E4

Dynamic Bit Signature Analysis

DBS features	Advantages
Automatic learning	Easily portable
Bit based analysis	More thorough analysis
Probabilistic check	Able to recognize structures even without full pattern match

**What about hidden drivers
and their detection?**



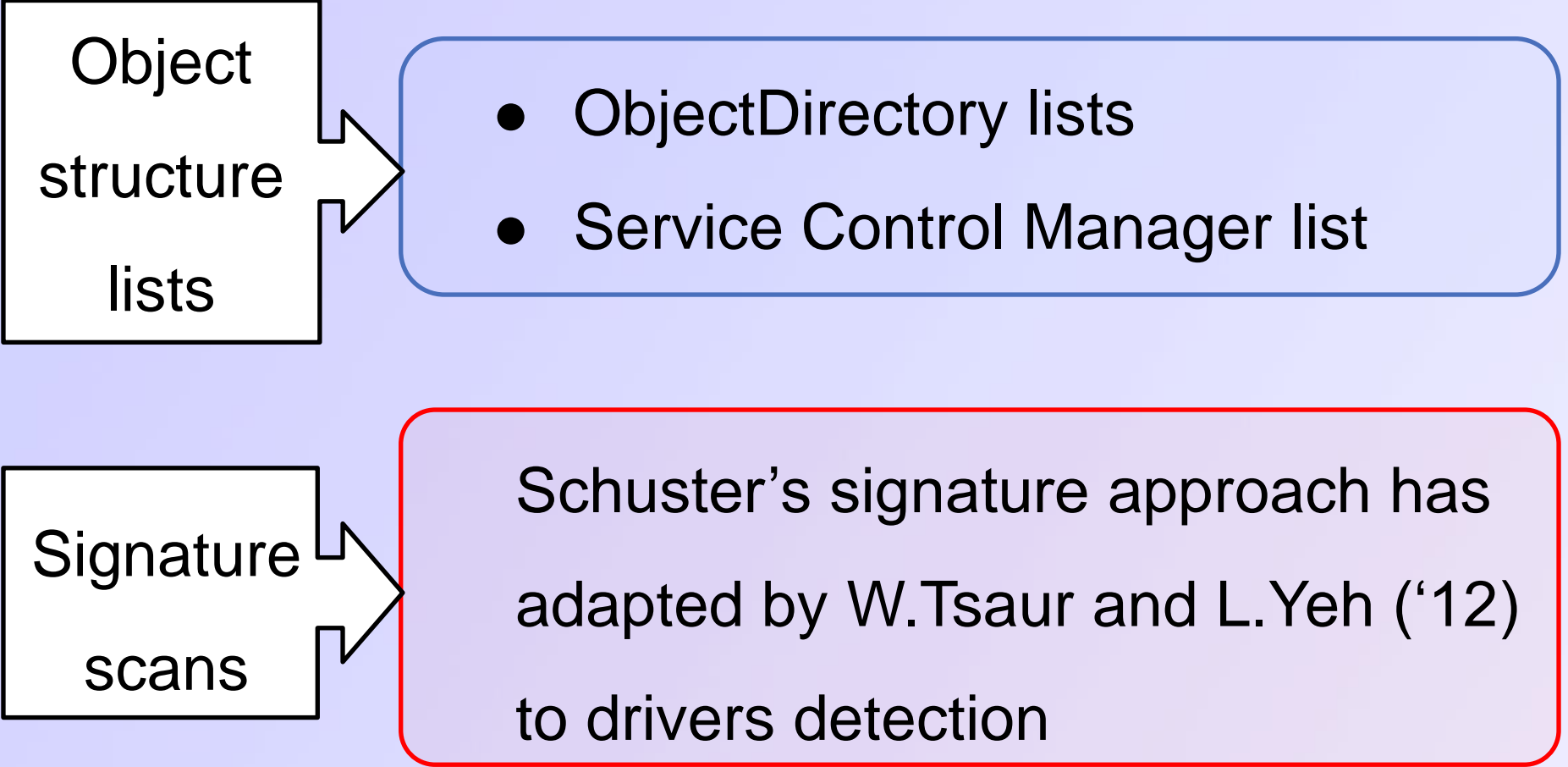
Hidden drivers have similar cases

	List view	Activity to hide
Processes	TaskMgr.exe	PsActiveProcessList modification
Drivers	DriverQuery.exe	PsLoadedModuleList modification

ZwQuerySystemInformation hooking leads to
processes & drivers hiding

Drivers detection approaches review

Object
structure
lists



- ObjectDirectory lists
- Service Control Manager list

Signature
scans

Schuster's signature approach has adapted by W.Tsaur and L.Yeh ('12) to drivers detection

**Is it possible to adapt DBS
for driver detection?**

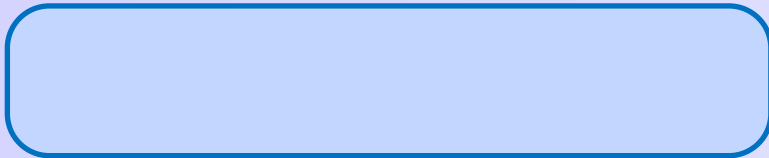
Is it possible to adapt DBS for driver detection?

DBS only can detect structures with a lot of fields

Is it possible to adapt DBS for driver detection?

DBS only can detect structures with a lot of fields

DRIVER_OBJECT



EPROCESS



Rating Point Inspection (RPI)

RPI improvements over DBS

- RPI utilizes additional weight matrix for precise pattern matching
- RPI use selective matching algorithm

If one of the checks is true	
DBS	RPI
add 1 point	1, 2 or etc. points are added to the final score

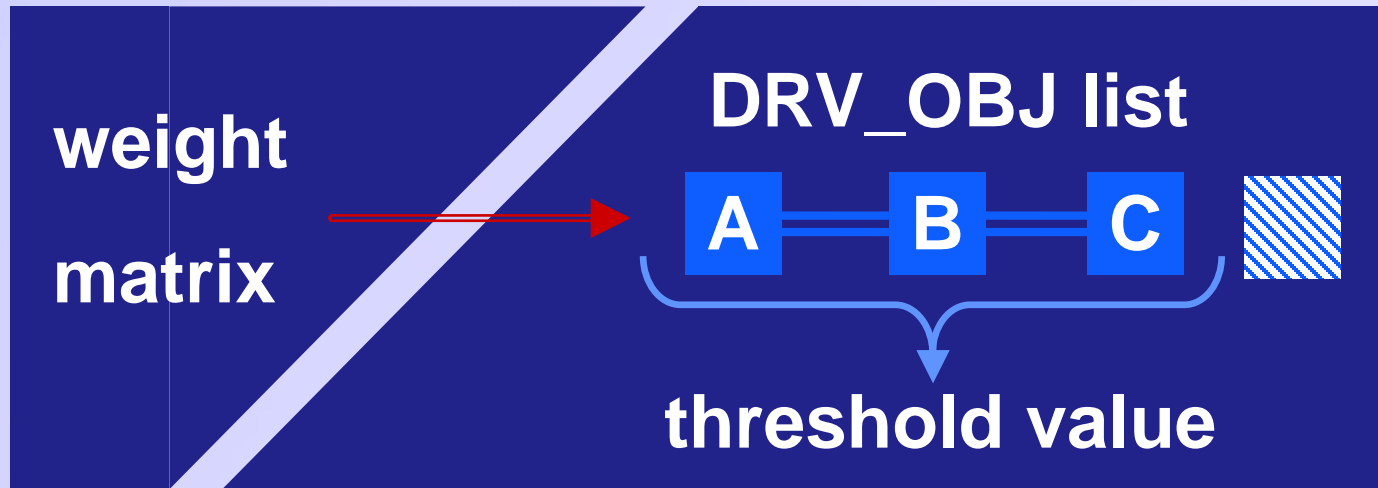
**Description of weight matrix for
DRIVER_OBJECT
is in the corresponding paper**



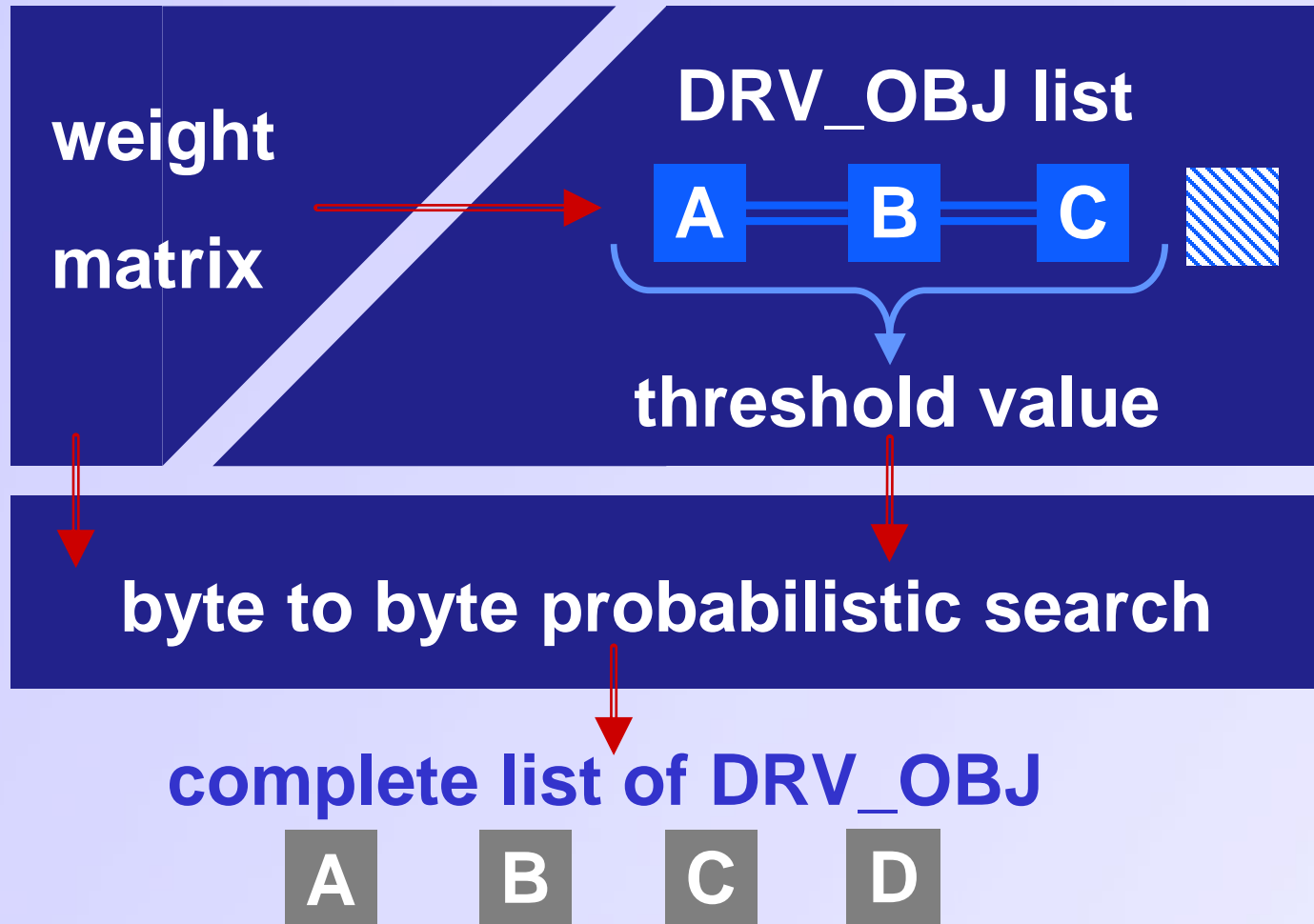
How does RPI detect drivers?



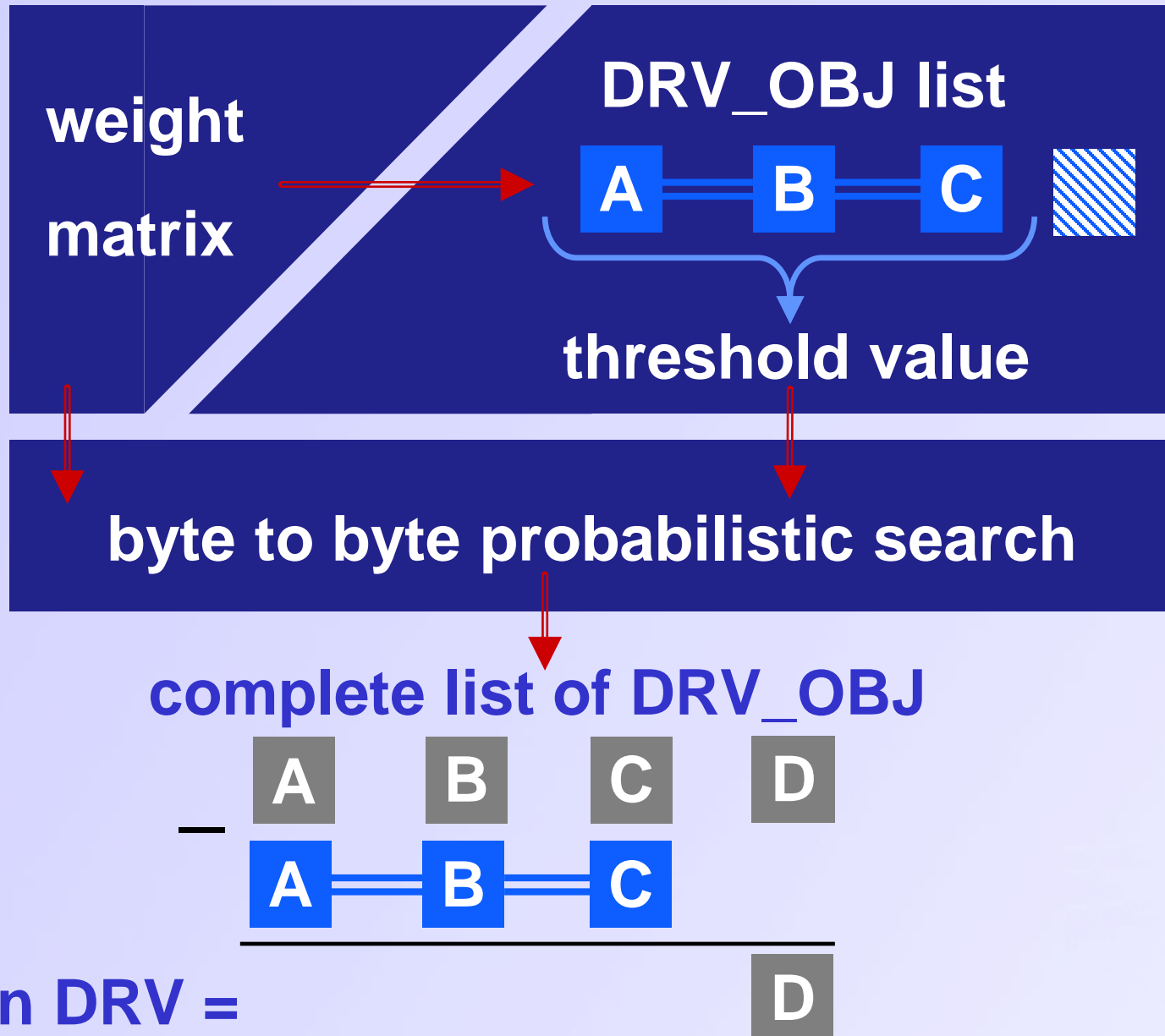
How does RPI detect drivers?



How does RPI detect drivers?



How does RPI detect drivers?



MASHKA's achievements

Reveals rootkits:

- Deliberately hidden processes and drivers
- Virus.Win32.Sality.q
- Trojan.Win32.VB.aqt
- Hidden drivers by ATSIIV

MASHKA's achievements

Reveals rootkits:

- Deliberately hidden processes and drivers
- Virus.Win32.Sality.q
- Trojan.Win32.VB.aqt
- Hidden drivers by ATSIV

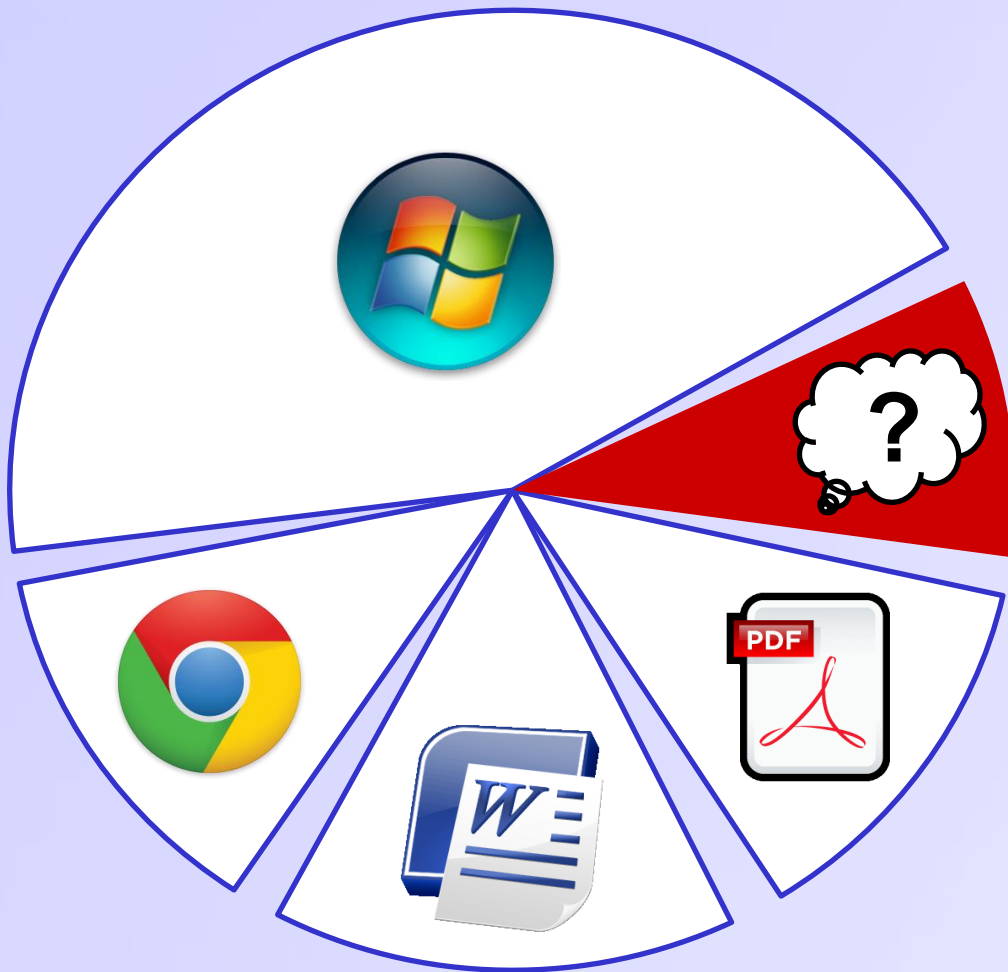
Existing anti-rootkits PowerTool, TDSSKiller, Xuetr fail,
but MASHKA can detect them

Demo: bit.ly/win8t6st

What is the pie filling?



What is the pie filling?



Igor Korkin, Ph.D

igor.korkin@gmail.com

sites.google.com/site/iykorkin



ADDITIONAL



WHAT IS IT MASHKA?

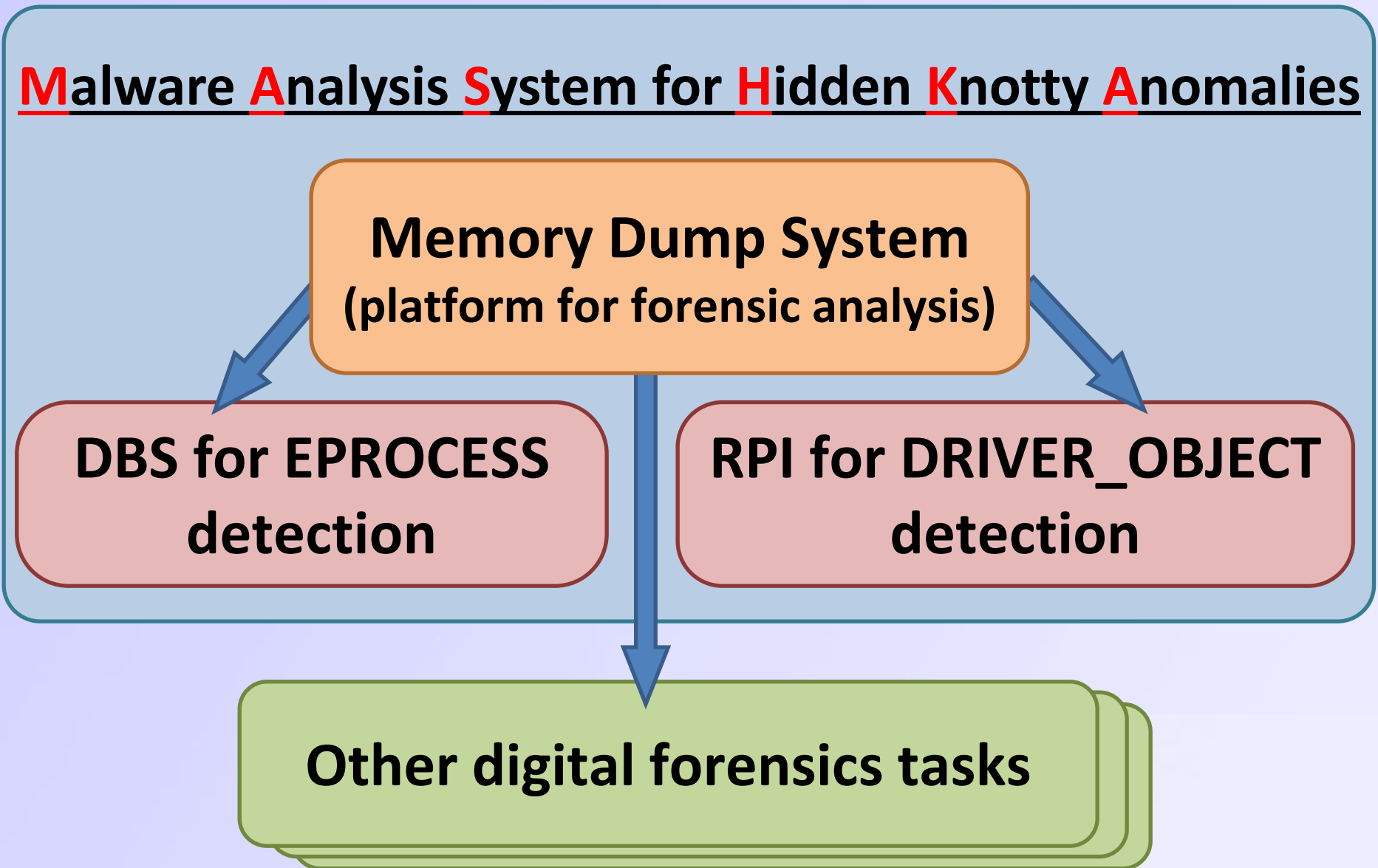
Malware **A**nalysis **S**ystem for **H**idden **K**notty **A**nomalies

Memory Dump System
(platform for forensic analysis)

**DBS for EPROCESS
detection**

**RPI for DRIVER_OBJECT
detection**

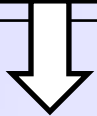
Other digital forensics tasks



MASHKA IN MEMORY FORENSICS TASKS

Various search signatures: char and wide char strings, byte fragments include addresses

As a result we receive: **What can we do with it?**

Name	Definition	
VALF	v irtual a ddress of the l oaded dump f ile	read data
ODUF	corresponding o ffset in d ump f ile	calculate offsets
VAOM	v irtual a ddress of the o riginal m emory	find value in dump

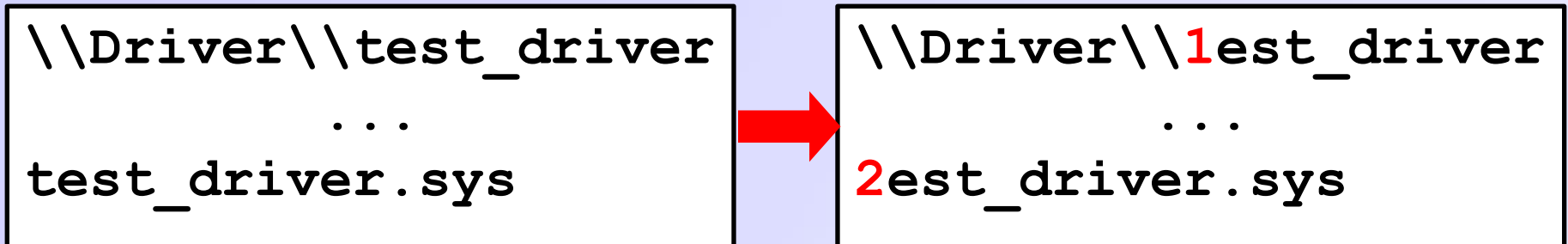
USE MASHKA TO RESEARCH DRIVERS

- 1. Run Windows under WinDbg control**
- 2. Install a test driver with 'ServiceName', 'DisplayName' and 'BinaryPath'**
- 3. Hide this driver structure by unlinking from PsLoadedModuleList**
- 4. Check the system with anti-rootkit tool**
- 5. Dump memory with the help of MASHKA**

USE MASHKA TO RESEARCH DRIVERS

6. Search strings from step 2 and save their 'VAOM'

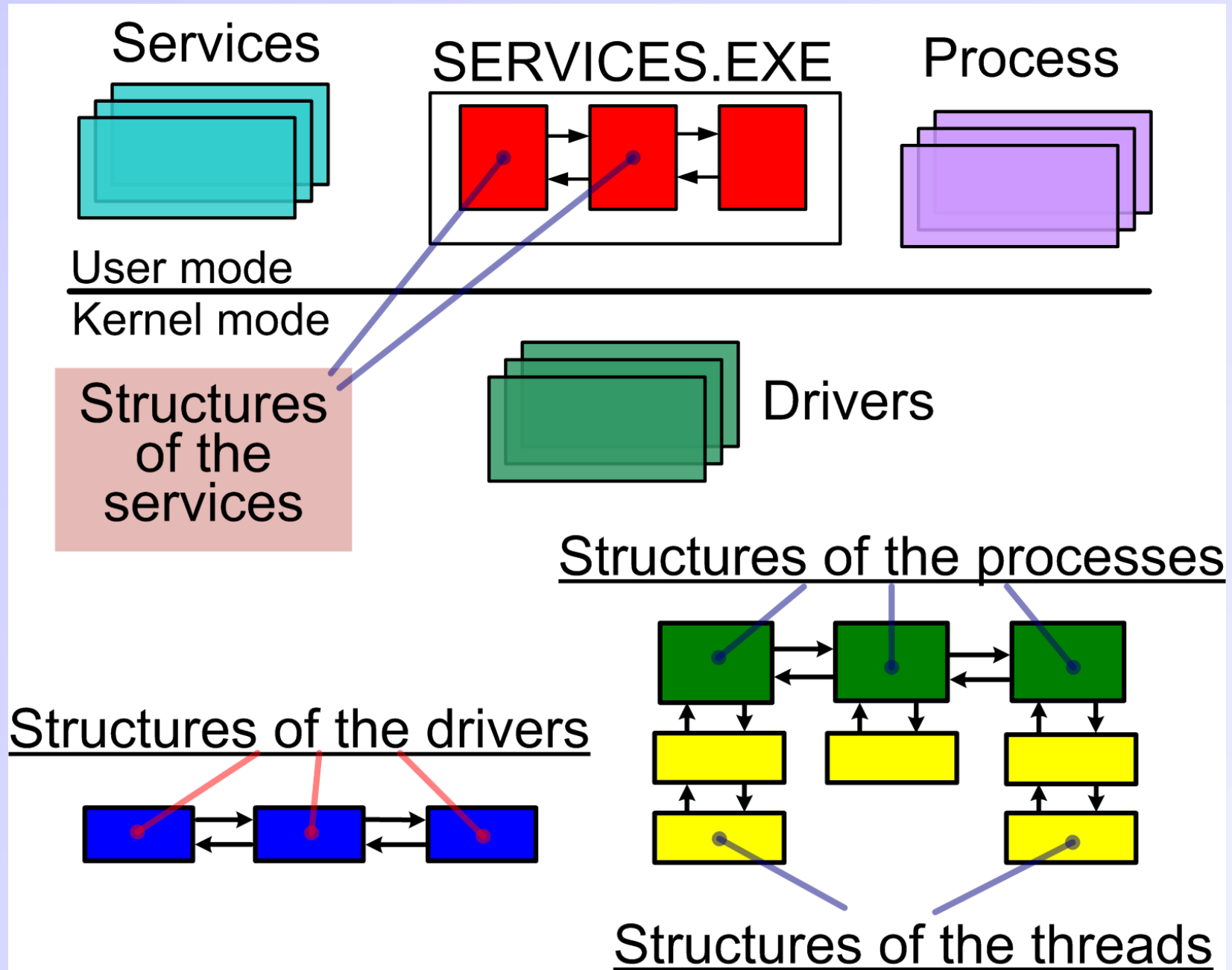
7. By WinDbg and strings VAOM change their content



8. Check the system repeatedly. Detection tools will give us a changed name.

By known 'VAOM' run further analysis

PROBLEM STATEMENT - ROOTKITS IN WINDOWS



ANALYSIS OF CURRENT APPROACHES TO DETECTION IN FACE OF OPPOSITIONS

**Cross-view detection is the main point
for all tools**

Low-level mechanisms:

- **Heuristic analyzer**
- **Additional object structure lists**
- **Signature scans are based on byte to
byte search of fragments of objects
structures in memory**

ANALYSIS OF SIGNATURE SCANS

- The fact that some fields' values are either known or exceed the constant, for example 0x8000_0000
- Parts of this method are implemented in the popular tools such as *GMER*, *PowerTool*, *XueTr*

Method's disadvantages:

- vulnerable to field modifications: If at least one byte does not match, the signature scan will miss the structure
- difficult to achieve portability on different versions of Windows OS, as it requires a lot of manual work

RPI FOR DRIVER_OBJECT DETECTION

Condition	Score
if (DRIVER_OBJECT_32.Type == 0x04)	1
if (DRIVER_OBJECT_32.Size == 0xa8)	1
if (chk_unicode_string(&DRIVER_OBJECT_32.DriverName))	2
if (chk_unicode_string(DRIVER_OBJECT_32.HardwareDatabase))	2
if ((DRIVER_OBJECT_32.MajorFunction[0]) >> 31)	2
if (max_same_major_functions(&DRIVER_OBJECT_32) >= min_major_function)	2
check_function_prologue(addr)	4

'global_scope' is a sum of points

THE 'CHECK_FUNCTION_PROLOGUE (ADDR)' FUNCTION

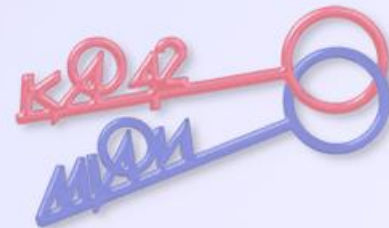
Condition	Result
<pre>If (((addr[i+0] == 0x55) && (addr[i+1] == 0x89) && (addr[i+2] == 0xe5)) ((addr[i+0] == 0x55) && (addr[i+1] == 0x8b) && (addr[i+2] == 0xec)) ((addr[i+0] == 0x53) && (addr[i+1] == 0x56)) ((addr[i+0] == 0x56) && (addr[i+1] == 0x57)) ((addr[i+0] == 0x56) && (addr[i+1] == 0x57)) ((addr[i+0] == 0x8b) && (addr[i+1] == 0xff)))</pre>	true or false

RPI APPLYING

1. Calculate all values, such as 'min_major_function' and 'global_scope'
2. Perform a byte-to-byte search by calculating the sum of points for each memory region
3. DRIVER_OBJECT structure is found if the probabilistic comparing of matching points with the 'global_scope' value is true
4. Compare the RPI-matching list with the drivers list, which has been obtained by ZwOpenDirectoryObject

FUTURE PLANS OF HOW TO USE & IMPROVE MASHKA

- Detection Shadow Walker-like Rootkits
- GPU Utilization in Memory Forensics
- The Idea of Cloud Anti-Rootkit or
Anti-Rootkit as a Service
- The Center of Mass of Kernel Mode Structures
- Digital Forensics in Education



TESTING RESULTS OF MASHKA

DBS approach has been successfully tested

deliberately
hidden
objects

real rootkits:

- Virus.Win32.Sality.q (Kaspersky Lab)
- Trojan.Win32.VB.aqt (Kaspersky Lab)

RPI approach has been successfully tested

deliberately
hidden
objects

real
rootkits

for hidden drivers which
were loaded by *ATSIV*
(Linchpin Labs)

In the latter case popular tools such as PowerTool, TDSSKiller, Xuetr cannot detect a hidden driver, but the RPI can

Demo - bit.ly/win8t6st

CONCLUSIONS

- Level of sophisticated malware increases
- Vulnerability of Windows OS
- Popular dump systems are vulnerable to intruder attacks
- Popular anti-rootkits are stopped by malware
- To prevent a possible attack, continue to maintain systems

CONCLUSIONS

- Use the page tables to memory dump
- Dynamic bit signatures can detect structures which have a typical design with a lot of members
- Rating point inspection can detect structure by detailed analysis of its members