# ICPC Lib

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| Contents |   |          | 7.7 dinic_with_scaling.cpp                |     |
|----------|---|----------|---|-----|
|          |   |          | 7.8 floyd_warshall.cpp                    | 11  |
| 1        | Bitmasks                                    | 1        | 7.9 kosaraju.cpp                          |     |
|          | 1.1 enumerate_submasks.cpp                  | 1        | 7.10 lexicographical_minimal_toposort.cpp | 12  |
|          | 1.2 sos_dp.cpp                              | 1        | 7.11 tarjan.cpp                           | 12  |
| 2        | Combinatorics                               | 2        | 8 Number Theory                           | 13  |
|          | 2.1 Fórmula de Legendre                     | 2        | 8.1 fastexp.cpp                           |     |
|          | 2.2 Teorema de Lucas                        | 2        | 8.2 $\operatorname{spf.cpp}$              | 13  |
|          | 2.3 Teorema de Sprague-Grundy e game-theory | 2        |   |     |
|          | 2.4 Stars and bars                          | 2        | 9 Range Queries                           | 13  |
|          |   |          | 9.1 bit.cpp                               |     |
| 3        | $count\_inversions\_merge\_sort.cpp$        | 2        | 9.2 segtree.cpp                           | 13  |
| 4        | dsu.cpp                                     | 2        | 10 Strategy                               | 14  |
|          |   |          | 10.1 Debugging:                           |     |
| 5        | fft.cpp                                     | 3        | 10.2 Início (primeira 1h30)               |     |
| c        | Geometry                                    | 9        | 10.3 Meio                                 |     |
| U        | 6.1 convex_hull.hpp                         | <b>ე</b> | 10.4 Final (última 1h30)                  | 14  |
|          | 6.2 point.hpp                               | ა<br>4   | 11 (1)                                    | 4 P |
|          | 6.3 polygon.hpp                             | 5        | 11 Strings                                | 15  |
|          |   | 6        | 11.1 rabin_karp.cpp                       |     |
|          | 6.4 triangle.hpp                            | O        | 11.2 suffix_array.cpp                     |     |
| 7        | Graphs                                      | 6        | 11.3 trie.cpp                             | 16  |
| •        | 7.1 2sat.cpp                                | 6        | 12 template.cpp                           | 17  |
|          | 7.2 bellman_ford.cpp                        | 7        | 12 template.cpp                           | 11  |
|          | 7.3 cycle_detection.cpp                     | 8        | 13 Trees                                  | 17  |
|          | 7.4 dfs_toposort.cpp                        | 8        | 13.1 binary_lifting.cpp                   | 17  |
|          | 7.5 dijkstra.cpp                            | 8        | 13.2 centroid.cpp                         |     |
|          | 7.6 dinic cpp                               | 9        | 13.3 hld no range update.cpp              |     |

| 13.4 | hld_with_range_update.cpp | 19 |
|------|---------------------------|----|
| 13.5 | lca.cpp                   | 21 |
| 13.6 | path_queries.cpp          | 21 |

### 1 Bitmasks

### 1.1 enumerate\_submasks.cpp

```
// Enumerate all submasks of mask in 2^(active bits of mask), which is
    optimal
// in DECREASING order
for (i64 submask = mask: submask > 0: submask = (submask - 1) & mask) {
    // do stuff
// NOTE: needs to treat 0 submask separately
1.2 \quad sos\_dp.cpp
#include <bits/stdc++.h>
using namespace std;
#define i64 int64 t
vector<i64> sosdp(const vector<i64> A, const i64 B, bool inverse =
    false){
    assert(A.size() == (111 << B));
    vector < i64 > S = A;
    for (i64 i = 0; i < B; i++){</pre>
         // introduce possible diff at i
         for (i64 mask = (111 << B) - 1; mask >= 0; mask--){
             if (mask & (111 << i)){</pre>
                 if (inverse){
                      S[mask] -= S[mask ^ (111 << i)];</pre>
                 } else {
                      S[mask] += S[mask ^ (111 << i)];
             }
         }
    return S;
vector < i64 > not_arr(const vector < i64 > A, const i64 B){
    vector < i64 > S(111 << B);</pre>
    for (i64 mask = 0; mask < (111 << B); mask++){</pre>
        S[mask] = A[((111 << B) - 1) ^ mask];
    return S;
// SOSP(A)[x] = sum(y supermask of x)A[x] = sum(y submask of \simx)A[\simy]
// = SOS(\sim A)[\sim x] = \sim SOS(\sim A)
// otherwise if SOSP(S) = A then
//\sim SOS(\sim S) = A ==> SOS(\sim S) = \sim A ==> \sim S = SOS^{-1}(\sim A) ==> S = \sim
    SOS^{-1}(\sim A)
vector < i64 > supersetdp (const vector < i64 > A, const i64 B, bool inverse
    return not_arr(sosdp(not_arr(A, B), B, inverse), B);
```

### 2 Combinatorics

### 2.1 Fórmula de Legendre

Para primo  $p, v_p(n!) = \sum_{k=1}^{\infty} \left\lfloor \frac{n}{p^k} \right\rfloor, e v_p(\binom{n}{k}) = v_p(n!) - v_p(k!) - v_p((n-k)!).$ 

Para 
$$m = \prod p_i^{e_i}$$
 composto,  $v_m(n!) = \min_i \left\lfloor \frac{v_{p_i}(n!)}{e_i} \right\rfloor$ .

### 2.2 Teorema de Lucas

Para p primo e  $n = n_k p^k + ... + n_1 p + n_0$  e  $m = m_k p^k + ... + m_1 p + m_0$ ,

$$\binom{n}{m} \equiv \prod_{i=0}^{k} \binom{n_i}{m_i} \mod p$$

### 2.3 Teorema de Sprague-Grundy e game-theory

Nim: ganha se e só se o XOR dos números da pilha é maior que 0:  $a_1 \oplus \ldots \oplus a_n \neq 0$ .

Estados terminais tem número de Grundy n(v) = 0 e não-terminais tem número de Grundy  $n(v) = \max_{v \to u} n(u)$ . Assim, um estado é **vencedor** (em um jogo em que quem não tem movimentos perde) se e só se  $n(v) \neq 0$ . Um único jogo é equivalente a uma **pilha de Nim com (limitados) acréscimos**.

Assim, se temos dois "sub-jogos" independentes  $S_1$  e  $S_2$  então seu número de Grundy é o XOR dos números de Grundy individuais,  $n(S_1 + S_2) = n(S_1) \oplus n(S_2)$ .

### 2.4 Stars and bars

Número de soluções inteiras para  $x_1+x_2+\ldots+x_k=n$  com  $x_i\geq 1$  para todo i é  $\binom{n-1}{k-1}$  (n-1) posições - aquelas entre j e j+1 para  $j=1\ldots n-1$  - para k-1 barras separadoras, sendo que cada posição só pode ser ocupada por no máximo uma barra)

Número de soluções inteiras para  $x_1 + x_2 + \ldots + x_k = n$  com  $x_i \ge 0$  para todo i é  $\binom{n+k-1}{k-1}$  (n unidades, k-1 barras separadoras, qualquer permutação é válida).

## ${f 3}$ count\_inversions\_merge\_sort.cpp

```
#include <bits/stdc++.h>
using namespace std;
// sorts and returns count of inversions. Modifies original array!
int inversions(vector<int>& a, int 1, int r){
    if (1 == r) return 0;
    const int m = 1 + (r - 1) / 2;
    int total = inversions(a, l, m) + inversions(a, m + 1, r);
    vector<int> left, right;
    for (int i = 1; i <= m; i++) left.emplace_back(a[i]);</pre>
    for (int i = m + 1; i <= r; i++) right.emplace_back(a[i]);</pre>
    int idx = 1;
    int pl = 0, pr = 0;
    while (pl < left.size() || pr < right.size()){</pre>
        // we only count the inversion pairs from the perspective of
            the smaller
        // element on the right
        // 1 3 5 | 2 4
        // if right[pr] < left[pl] then it is smaller than
            [left[pl]..[left[|left|-1]]]
        // so contributes for left-pl cross-inversions
        if (pr == right.size())
            a[idx++] = left[pl++];
        else if (pl == left.size())
            a[idx++] = right[pr++];
        else {
            if (left[pl] < right[pr])</pre>
                a[idx++] = left[pl++];
            else {
                a[idx++] = right[pr++];
                total += left.size() - pl;
        }
    return total;
    dsu.cpp
#include <bits/stdc++.h>
using namespace std;
struct DSU {
    vector<int> parent, sz;
    vector < vector < int >> elements;
    DSU(int n){
        parent.resize(n);
        iota(parent.begin(), parent.end(), 0);
        sz.assign(n, 1);
        elements.resize(n);
        for (int i = 0; i < n; i++){</pre>
            elements[i].emplace_back(i);
```

```
}
    int find(int x){
        if (parent[x] != x){
            parent[x] = find(parent[x]); // path-compression
        return parent[x];
    }
    void unite(int x, int y){
        x = find(x):
        y = find(y);
        if (x == y){
            return:
        } else if (sz[x] > sz[y]) {
            swap(x, y);
        sz[y] += sz[x];
        for (int z: elements[x]){
            elements[y].emplace_back(z);
        parent[x] = y;
    }
};
    fft.cpp
// CODE TAKEN FROM CP-ALGORITHMS
#include <bits/stdc++.h>
using namespace std;
using cd = complex < double >;
const double PI = acos(-1);
void fft(vector < cd> & a. bool invert) {
    int n = a.size():
    for (int i = 1, j = 0; i < n; i++) {</pre>
        int bit = n >> 1;
        for (; j & bit; bit >>= 1)
               `= bit;
        j ^= bit;
        if (i < j)
            swap(a[i], a[j]);
    }
    for (int len = 2; len <= n; len <<= 1) {
        double ang = 2 * PI / len * (invert ? -1 : 1);
        cd wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i += len) {</pre>
            cd w(1):
            for (int j = 0; j < len / 2; j++) {
                cd u = a[i+j], v = a[i+j+len/2] * w;
                a[i+j] = u + v;
```

```
a[i+j+len/2] = u - v;
                w *= wlen;
            }
        }
    }
    if (invert) {
        for (cd & x : a)
            x /= n;
}
vector<int> multiply(vector<int> const& a, vector<int> const& b) {
    vector < cd > fa(a.begin(), a.end()), fb(b.begin(), b.end());
    int n = 1:
    while (n < a.size() + b.size())</pre>
        n <<= 1:
    fa.resize(n);
    fb.resize(n);
    fft(fa, false);
    fft(fb, false);
    for (int i = 0; i < n; i++)</pre>
        fa[i] *= fb[i]:
    fft(fa, true);
    vector < int > result(n);
    for (int i = 0: i < n: i++)
        result[i] = round(fa[i].real());
    return result;
    Geometry
6.1 convex_hull.hpp
#include "./point.hpp"
#include "./polygon.hpp"
#define i64 int64 t
// counterclockwise convex hull, can include collinear points or not
vector < pt > convex_hull(vector < pt > poly, bool INCLUDE_COLLINEAR) {
    // choose bottommost point as pivot
    pt pivot = poly[0];
    const i64 n = poly.size();
    for (i64 i = 1; i < n; i++){}
        if (poly[i].y < pivot.y || (poly[i].y == pivot.y && poly[i].x</pre>
            < pivot.x)){
            pivot = poly[i];
        }
    sort(poly.begin(), poly.end(),
        [&pivot](pt p, pt q){return polarComp(pivot, p, q);}
    ); // pivot will be always poly[0] after sort
```

```
// since including collinear points, we want to privilege the
       farthest collinear points
    // when coming back (e.g (0,0), (1,0), (2,0), (2,1), (2,2), (1,1))
    if (INCLUDE COLLINEAR) {
        i64 i = n - 1;
        while (i > 0 && orientation(pivot, poly[i], poly.back()) == 0){
        reverse(poly.begin() + i + 1, poly.end());
    }
    vector<pt> hull = {pivot};
    for (i64 i = 1: i < n: i++){
        while (
            hull.size() >= 2
            && (
                orientation(hull[hull.size() - 2], hull[hull.size() -
                    1], poly[i]) == 1
                || (!INCLUDE_COLLINEAR && orientation(hull[hull.size()
                    - 2], hull[hull.size() - 1], poly[i]) == 0)
            )
        ) {
            // goes clockwise, break orientation
            hull.pop_back();
        hull.emplace_back(poly[i]);
    return hull:
}
6.2 point.hpp
#include <bits/stdc++.h>
using namespace std;
#define i64 int64 t
struct pt {
    i64 x, y;
    pt(i64 x, i64 y) : x(x), y(y) {}
    pt(const pt \& p) : x(p.x), y(p.y) \{\}
    i64 cross(pt p){
        return x * p.y - y * p.x;
    i64 dot(pt p){
        return x * p.x + y * p.y;
    void operator=(const pt q){
        x = q.x;
        y = q.y;
};
pt operator-(pt p1, pt p2){
    return pt(p1.x - p2.x, p1.y - p2.y);
```

```
pt operator+(pt p1, pt p2){
    return pt(p1.x + p2.x, p1.y + p2.y);
pt operator*(pt p1, i64 k){
    return pt(k * p1.x, k * p1.v);
pt operator*(i64 k, pt p1){
    return pt(k * p1.x, k * p1.y);
i64 orientation(pt p, pt q, pt r){
    i64 \ o = (q - p).cross(r - p);
    if (o > 0){
        return -1; // counterclockwise
    } else if (o < 0){
        return 1; // clockwise
    } else {
        return 0: // collinear
i64 ccw(pt p, pt q, pt r) { return orientation(p, q, r) < 0; }
bool is_collinear(pt p, pt q, pt r){
    return (q - p).cross(r - p) == 0;
// is p on segment q,r?
bool on_segment(pt p, pt q, pt r){
    // pqr is collinear
    // qp and pr have the same sign/orientation (as p in between q and
       r)
    return (
        is_collinear(p, q, r)
        && ((p - q).dot(r - p) >= 0)
    );
// counterclockwise polar sort
bool polarComp(pt pivot, pt p, pt q){
    auto d2 = [](pt a){}
        return a.x * a.x + a.y * a.y;
    i64 o = orientation(pivot, p, q);
    if (o == -1) return true:
    else if (o == 1) return false;
        // collinear => return closest point
        return d2(p - pivot) < d2(q - pivot);</pre>
// line intersection
```

```
bool line_line_intersect(pt a, pt b, pt c, pt d){
    return (
        (b - a).cross(d - c) != 0
        || (b - a).cross(c - a) == 0
    ); // not parallel, or equal lines (ab||cd and abc collinear)
// line ab, segment cd
bool line_segment_intersect(pt a, pt b, pt c, pt d){
    // c and d are in different half-planes
    // ==> different orientations
    // OR any triple abc, abd is collinear
    return orientation(c, a, b) * orientation(d, a, b) != 1;
// https://cp-algorithms.com/geometry/check-segments-intersection.html
bool segment_segment_intersect(pt a, pt b, pt c, pt d){
    auto inter1d = [](i64 11, i64 r1, i64 12, i64 r2){
        if (11 < r1) swap(11, r1):
        if (12 < r2) swap(12, r2);
        return max(11, 12) <= min(r1, r2);</pre>
    };
    if (orientation(a, b, c) == 0 && orientation(a, b, d) == 0){
        // all four are collinear
        // reduces to 1d interval intersection
        return inter1d(a.x, b.x, c.x, d.x) && inter1d(a.y, b.y, c.y,
           d.v);
    }
    return
        // different half-planes: line ab intersects segment cd
        (orientation(a, b, c) != orientation(a, b, d))
        // different half-planes: line cd intersects segment cd
        && (orientation(c, d, a) != orientation(c, d, b))
// horizontal ray from a, direction +-1, intersects segment cd?
// DOES include intersection at endpoints (cd) by default,
// verified by CSES "Point in Polygon"
bool horizontal_ray_segment_intersect(pt a, pt c, pt d, i64 direction
   = 1){}
    if (c.y \le a.y \&\& a.y \le d.y){ // different half-planes from ray,
       c below or at line, d above
        // now we need d in the same half-plane as the ray
           (considering line ac)
        // so the intersection happens on the ray and not on the other
           side
        if (direction == 1){
            return orientation(c, a, d) > 0;
        } else {
            return orientation(c, a, d) < 0;</pre>
    } else if (d.y <= a.y && a.y < c.y){</pre>
        return horizontal_ray_segment_intersect(a, d, c, direction);
```

```
return false;
6.3 polygon.hpp
#include "./point.hpp"
#include "./triangle.hpp"
// shoelace formula
double area(vector<pt>& p){
    i64 a = 0:
    i64 n = p.size();
    for (i64^{\circ}i = 0; i < n; i++){}
        a += p[i].cross(p[(i + 1) % n]); // area of OP[i]P[i+1]
    return (double)a / 2.0f;
}
// O(log n) check for boundary/inside convex polygon
// take the first pt (lexicographically) as reference (PO), and sort
   by polar order
// since polygon is convex, then P = U P[0]P[i]P[i+1] for i=1..n-2
// do binary search to find i, and then check for being inside triangle
// check for boundaries separately
const i64 INSIDE = 1;
const i64 BOUNDARY = 0;
const i64 OUTSIDE = -1;
i64 inside_polygon_convex(pt q, vector <pt> p){
    const i64 n = p.size();
    pt pivot = p[0];
    for (pt a: p){
        if (make_pair(a.x, a.y) < make_pair(pivot.x, pivot.y)){</pre>
            pivot = a;
    sort(p.begin(), p.end(), [&pivot](pt a, pt b){return
       polarComp(pivot, a, b);});
    i64 l = 1, r = n - 2, i = -1;
    while (1 \le r)
        i64 m = 1 + (r - 1) / 2;
        // is in P[0]P[m]P[m+1] iff P[m]P[0]Q is ccw but P[m+1]P[0]Q
           is cw
        if (ccw(p[m], p[0], q)){
            if (!ccw(p[m + 1], p[0], q)){
                i = m;
                break:
            } else {
                1 = m + 1;
        } else {
            r = m - 1:
   }
    // inside_triangle includes boudaries
    if (i == -1 || !inside_triangle(q, p[0], p[i], p[i + 1])) return
```

```
OUTSIDE:
    else if (
        on_segment(q, p[i], p[i + 1])
        || (i == 1 \&\& on_segment(q, p[0], p[1]))
        || (i == n - 2 \&\& on_segment(q, p[0], p[n - 1]))
    ) {
        return BOUNDARY:
    } else return INSIDE:
// Raycasting algorithm, works on all polygons, convex or not.
   Complexity: O(N)
// if the ray intersects a vertex, then it only counts if orientation
// (so it does not hit "tangentially", but cuts through)
i64 inside_polygon(pt q, vector<pt>& p){
    const i64 n = p.size();
    // draw horizontal rav
    i64 ray_intersections = 0;
    for (i64 i = 0; i < n; i++)
        // p[i]p[i+1]
        if (on_segment(q, p[i], p[(i + 1) % n])){
            return BOUNDARY;
        ray_intersections += horizontal_ray_segment_intersect(q, p[i],
           p[(i + 1) % n]);
    }
    // to be inside: odd exits
    if (ray intersections % 2 == 1){
        return INSIDE;
    } else {
        return OUTSIDE;
6.4 triangle.hpp
#include "./point.hpp"
// pgr oriented area
// can also be calculated by shoelace
double area(pt p, pt q, pt r){
    return (double)((q - p).cross(r - p)) / 2.0f;
bool inside_triangle(pt q, pt p1, pt p2, pt p3){
    i64 abs_area = 0;
    pt p[3] = \{p1, p2, p3\};
    for (i64 i = 0; i < 3; i++){
        abs_area += abs((p[i] - q).cross(p[(i + 1) % 3] - q));
    i64 \text{ total} = abs((p[1] - p[0]).cross(p[2] - p[0]));
    return abs_area == total;
```

### 7 Graphs

### 7.1 2sat.cpp

```
#include <bits/stdc++.h>
using namespace std;
namespace SAT2 {
    vector<vector<int>> build_graph(const vector<pair<int, int>>
        &clauses. int n) {
        vector < vector < int >> g(2 * n);
        for (auto[x, y]: clauses){
            // (x or y) clause = (\sim x => y) clause and (\sim y => x) clause
            g[(x + n) \% (2 * n)].emplace_back(y);
            g[(y + n) \% (2 * n)].emplace_back(x);
        return g;
    vector < vector < int >> transpose (vector < vector < int >> &g) {
        vector < vector < int >> gt(g.size());
        for (int u = 0; u < g.size(); u++){</pre>
            for (int v: g[u]){
                 gt[v].emplace_back(u);
        }
        return gt;
    void forward dfs(
        const int u, vector<vector<int>> &g, vector<bool> &vis,
            vector<int> &stack
        vis[u] = true:
        for (int v: g[u]){
            if (!vis[v]){
                forward_dfs(v, g, vis, stack);
        }
        stack.emplace_back(u);
    void backward dfs(
        const int u, vector < vector < int >> &g, vector < bool > &vis,
        vector < int > & comp, int c = 0
   ) {
        vis[u] = true;
        comp[u] = c;
        for (int v: g[u]){
            if (!vis[v]){
                backward_dfs(v, g, vis, comp, c);
        }
    }
```

```
pair < vector < int >, int > scc(const vector < pair < int , int >> & clauses,
    vector < vector < int >> g = build_graph(clauses, n);
    vector < vector < int >> gt = transpose(g):
    const int s = g.size();
    vector < int > stack, comp(s);
    vector <bool> vis(s. false):
    for (int u = 0; u < s; u++){
        if (!vis[u]){
            forward_dfs(u, g, vis, stack);
    }
    reverse(stack.begin(), stack.end());
    fill(vis.begin(), vis.end(), false);
    int c = 0;
    for (int u: stack){
        if (!vis[n]){
            backward_dfs(u, gt, vis, comp, c);
    return {comp, c};
pair < vector < bool > , bool > solve(const vector < pair < int , int >> &
   clauses. int n){
    // run scc (Kosaraju)
    auto[comp, num_comps] = scc(clauses, n);
    // build quotient graph
    vector < vector < int >> g_scc(num_comps);
    for (auto[x, y]: clauses){
        // (x or y) = (\simx => y) = (\simy => x)
        g_scc[comp[(x + n) \% (2 * n)]].emplace_back(comp[y]);
        g_scc[comp[(y + n) \% (2 * n)]].emplace_back(comp[x]);
    // build valuation: mark as true the first one that appears
    vector < bool > valuation(n):
    for (int u = 0; u < n; u++) {
        if (comp[u] < comp[u + n]){
            valuation[u] = false;
        } else {
            valuation[u] = true;
    // check valuation
    bool ok = true:
    for (auto[x, y]: clauses){
        bool left = ((x < n) ? valuation[x] : !valuation[x - n]);</pre>
        bool right = ((y < n) ? valuation[y] : !valuation[y - n]);</pre>
        ok &= left | right;
    }
```

```
return {valuation, ok};
};
7.2 bellman_ford.cpp
#include <bits/stdc++.h>
using namespace std;
// edges: (a, b, c) = directed edge from a to b with weight c
vector<int> negative_cycle(const vector<tuple<int, int, int>> edges,
   int n){
   // virtual vertex connects to all others with weight 0
    // so ALL negative cycles are reachable!
    // for that effect we just need to set d[u] = 0 for all u
    // (we don't need an actual new vertex)
    vector<int> d(n, 0);
    vector<int> source_list;
    vector<int> prev(n, -1);
    // all paths of size s
    for (int s = 0: s \le n: s++){
        bool any = false;
        for (auto[a, b, c]: edges){
            if (d[a] + c < d[b])
                any = true;
                prev[b] = a;
                d[b] = d[a] + c;
        if (!any) break;
    // run a (single) new iteration of relaxation
    int target = -1;
    for (auto[a, b, c]: edges){
        if (d[a] + c < d[b])
            target = b;
            break;
        }
    }
    if (target != -1){
        vector < int > cyc;
        // might be a cycle end or a vertex reachable from a negative
           cvcle
        int curr = target;
        vector < bool > vis(n);
        while (!vis[curr]) {
            cvc.emplace_back(curr);
            vis[curr] = true;
            // we might have prev[curr] == curr
            // (in case of self negative edges!)
            curr = prev[curr];
        // now curr is a visited vertex (closed the cycle!)
        cvc.emplace_back(curr);
        reverse(cyc.begin(), cyc.end());
```

```
while (cyc.back() != cyc[0]){
            cyc.pop_back(); // remove vertices reachable from the
               cycle but not in it
        return cyc;
    } else {
        return vector < int > ();
7.3 cycle_detection.cpp
#include <bits/stdc++.h>
using namespace std;
#define WHITE 0
#define GRAY 1
#define BLACK 2
bool has_cycle = false;
void dfs(vector<vector<int>> &g, int u, vector<int> &color){
    if (color[u] != WHITE){
        return:
    color[u] = GRAY:
    for (int v: g[u]){
        if (color[v] == GRAY){
            has_cycle = true;
        } else if (color[v] == WHITE){
            dfs(g, u, color);
    color[u] = BLACK;
7.4 dfs_toposort.cpp
#include <bits/stdc++.h>
using namespace std;
namespace Graph {
    int c;
    void dfs(
        int u, vector<vector<int>> &G, vector<bool> &vis, vector<int>
           &order
    ) {
        vis[u] = true;
        for (int v: G[u]){
            if (!vis[v]){
                dfs(v, G, vis, order);
        }
        order[u] = c;
        --c;
    }
```

```
vector<int> toposort(vector<vector<int>> &G){
        c = G.size() - 1:
        const int n = G.size():
        vector < bool > vis(n, false);
        vector < int > order(n);
        for (int u = 0; u < n; u++) {
            if (!vis[u]){
                dfs(u, G, vis, order);
        }
        // validate order
        for (int u = 0; u < n; u++) {
            for (int v: G[u]){
                if (order[u] >= order[v]) {
                     return vector < int > ();
            }
        }
        return order;
     dijkstra.cpp
#include <bits/stdc++.h>
using namespace std;
const int INF = (int)2e9;
// directed edges are pairs (weight, endpoint)
void dijkstra(vector<vector<pair<int, int>>> &adj, int dis[], int s)
    for (int i = 0; i < (int)adj.size(); i++) dis[i] = INF;</pre>
    priority_queue < pair < int , int > , vector < pair < int , int >> ,
       greater < pair < int , int >> > pq;
    pq.emplace(0, s);
    dis[s] = 0;
    while (!pq.empty())
        auto [dis_u, u] = pq.top();
        pq.pop();
        if (dis_u > dis[u])
            continue; // skip suboptimal
        // since dis[u] is always updated when better path is found,
        // dis[u] <= d for each [d, u] state on the priority queue
        for (auto [w, v] : adj[u])
            if (dis[u] + w < dis[v])</pre>
                dis[v] = dis[u] + w;
                pq.emplace(dis[v], v);
                // sorting by dis[v] is the same as sorting by the
```

```
weights w[u, v]
                // since dis[v] - w(u, v) = dis[u] for all v (constant
            }
        }
    }
7.6 dinic.cpp
// Tested on CSES - Police Chases
#include <bits/stdc++.h>
using namespace std;
#define i64 int64_t
Dinic (overview):
// make residual graph (edges and backedges)
// then on each iteration
// do bfs to make layer graph
// do dfs to find blocking flow
*/
const i64 INF = (i64)2e18;
struct Edge {
    i64 u, v, c, f;
    i64 cap() const { return c - f; }
    Edge (i64 from, i64 to, i64 cap, i64 flow) : u(from), v(to),
       c(cap), f(flow) {}
}:
// for edge indices 0..m-1 in original network
// we represent forward edge i as 2 * i and backedge as 2*i + 1
// thus forward and backedge of indices i1, i2 relate in this way:
// i1 ^ 1 = i2, i2 ^ 1 = i1
struct Dinic {
    i64 n, s, t;
    vector < vector < i64>> adj; // maps vertex to edge indices
    vector < Edge > edges;
    // performance optimization to push flow possibly more than once
    // after each bfs/layer graph build
    vector < i64 > cur_edge_ptr;
    // performance optimization to avoid looking at blocked nodes
    vector < i64 > blocked:
    vector < i64 > level;
    queue < i64 > q;
    void add_edge(i64 u, i64 v, i64 c){
        // set flow at c so effective cap of back edge starts at 0
        Edge forward(u, v, c, 0), backward(v, u, c, c);
        adj[u].emplace_back(edges.size());
        edges.emplace_back(forward);
        adj[v].emplace_back(edges.size());
        edges.emplace_back(backward);
    }
```

```
Dinic(const i64 n, i64 s, i64 t) : n(n), s(s), t(t){
    adi.resize(n):
    level.resize(n):
    cur_edge_ptr.assign(n, 0);
    blocked.assign(n, false);
    edges.reserve(2 * n);
// level vector --> implicit layer graph (by takings edges where d
   increases by 1)
// pass by ref to optimize copies
void layer_graph(){
    fill(level.begin(), level.end(), INF);
    level[s] = 0;
    q.emplace(s);
    while (!q.empty()){
        i64 u = q.front();
        q.pop();
        for (i64 id: adj[u]){
            const Edge& e = edges[id]; // u -> v
            if (e.cap() == 0) continue; // edge does not exist
            i64 v = e.v:
            if (level[v] == INF){
                level[v] = level[u] + 1;
                q.emplace(v);
        }
    while (!q.empty()) q.pop();
// push_flow on layer dag of residual graph
i64 push flow(i64 u. vector<i64>& blocked. vector<i64>& level. i64
   dF){
    if (u == t){
        return dF; // already at sink
    bool all_neighbors_blocked = true; // performance optimization
    i64 total_pushed = 0;
    for (; cur_edge_ptr[u] < (i64)adj[u].size() && dF > 0;
       cur_edge_ptr[u]++){
        i64 id = adj[u][cur_edge_ptr[u]];
        // edges[id] is forward edge (relative to u): u -> v
        i64 v = edges[id].v;
        if (level[v] != level[u] + 1) continue; // not dag edge
        if (edges[id].cap() == 0 || blocked[v]) continue; // edge
           does not exist
        i64 pushed = push_flow(v, blocked, level,
           min(edges[id].cap(), dF));
        edges[id].f += pushed;
        edges[id ^ 1].f -= pushed;
        total_pushed += pushed;
        dF -= pushed;
```

```
if (edges[id].cap() > 0 && !blocked[v]) { // edge still
                 all_neighbors_blocked = false;
            }
        }
        if (all_neighbors_blocked) blocked[u] = true;
        return total_pushed;
    }
    i64 maxflow(){
        while (true) {
            laver_graph();
            if (level[t] == INF) {
                 break;
            } // no more augmenting paths
            // can't assume anyone is blocked (reuse block vector),
                since we might try to unblock by undoing flow
            fill(cur_edge_ptr.begin(), cur_edge_ptr.end(), 0); //
                reset edge ptr
            fill(blocked.begin(), blocked.end(), false); // reset
                blocked
            while (push_flow(s, blocked, level, INF)) {}
        i64 \text{ maxFlow} = 0;
        for (i64 id: adj[s]) {
            maxFlow += edges[id].f;
        return maxFlow;
    }
    // assumes dinic was called first
    vector < i64 > mincut() {
        layer_graph();
        vector < i64 > partition_num(n);
        for (i64 u = 0; u < n; u++){}
            if (level[u] < INF){</pre>
                 partition num[u] = 0:
            } else {
                 partition_num[u] = 1;
            }
        return partition_num;
    }
    dinic_with_scaling.cpp
// Tested on CSES - Police Chases
#include <bits/stdc++.h>
using namespace std;
#define i64 int64_t
// Same as dinic but with edge scaling
// Complexity is O(VE \log U) where U = \max \text{ edge capacity}
// WE ONLY INCLUDE WHAT CHANGED
```

const i64 INF = (i64)(111 << 60); // best to use a power of 2</pre>

};

```
// level vector --> implicit layer graph (by takings edges where d
   increases by 1)
// NOTE: when using dinic with scaling for mincut, do layer graph with
   default lim of 1
void layer_graph(i64 lim = 1){
   fill(level.begin(), level.end(), INF);
    level[s] = 0;
    q.emplace(s);
    while (!q.empty()){
        i64 u = q.front();
        q.pop();
        for (i64 id: adj[u]){
            const Edge& e = edges[id]; // u -> v
            if (e.cap() < lim) continue; // edge does not exist</pre>
            i64 v = e.v:
            if (level[v] == INF){
                level[v] = level[u] + 1;
                q.emplace(v);
       }
    while (!q.empty()) q.pop();
// push_flow on layer dag of residual graph
bool can_push_flow(i64 u, vector<i64>& level, i64 flow){
   if (u == t){
        return true; // already at sink
    for (; cur_edge_ptr[u] < (i64)adj[u].size(); cur_edge_ptr[u]++){</pre>
        i64 id = adj[u][cur_edge_ptr[u]];
        // edges[id] is forward edge (relative to u): u -> v
        i64 v = edges[id].v;
        if (level[v] != level[u] + 1) continue; // not dag edge
        if (edges[id].cap() < flow) continue; // edge does not exist</pre>
        i64 pushed = can push flow(v. level. flow):
        if (pushed){
            edges[id].f += flow;
            edges[id ^ 1].f -= flow;
            return true;
        }
   return false:
i64 maxflow(){
    for (i64 lim = INF; lim >= 1; ){
        layer_graph(lim);
        if (level[t] == INF) {
            lim >>= 1;
            continue;
        } // no more augmenting paths
```

```
// can't assume anyone is blocked (reuse block vector), since
            we might try to unblock by undoing flow
        fill(cur_edge_ptr.begin(), cur_edge_ptr.end(), 0); // reset
            edge ptr
        while (can_push_flow(s, level, lim)) {}
    i64 \text{ maxFlow} = 0:
    for (i64 id: adj[s]) {
        maxFlow += edges[id].f;
    return maxFlow;
    floyd_warshall.cpp
#include <bits/stdc++.h>
using namespace std;
const int INF = (int)2e9; // might change depending on the problem
vector < vector < int >> floydWarshall (const vector < tuple < int, int >> &
   edges, int n){
    vector < vector < int >> d(n, vector < int > (n, INF));
    for (int i = 0; i < n; i++){</pre>
        d[i][i] = 0:
    }
    for (auto[a, b, w]: edges){
        d[a][b] = min(d[a][b], w);
    for (int k = 0; k < n; ++k) {
        for (int i = 0; i < n; ++i) {</pre>
            for (int j = 0; j < n; ++ j) {
                 if (d[i][k] < INF && d[k][j] < INF)</pre>
                     d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
        }
    return d;
    kosaraju.cpp
#include <bits/stdc++.h>
using namespace std;
vector < vector < int >> transpose (const vector < vector < int >> &g) {
    vector < vector < int >> gt(g.size());
    for (int u = 0; u < g.size(); u++){</pre>
        for (int v: g[u]){
            gt[v].emplace_back(u);
    return gt;
```

```
void forward dfs(
     const int u, const vector < vector < int >> &g, vector < bool > &vis,
        vector<int> &stack
) {
    vis[u] = true:
    for (int v: g[u]){
         if (!vis[v]){
             forward_dfs(v, g, vis, stack);
     stack.emplace_back(u);
void backward_dfs(
     const int u, const vector<vector<int>> &g, vector<bool> &vis,
     vector < int > & comp, int c = 0
) {
     vis[u] = true;
     comp[u] = c;
     for (int v: g[u]){
         if (!vis[v]){
             backward_dfs(v, g, vis, comp, c);
    }
pair < vector < int > , int > scc(const vector < vector < int > > &g , int n) {
     vector < vector < int >> gt = transpose(g);
     const int s = g.size();
     vector < int > stack, comp(s);
     vector < bool > vis(s, false);
     for (int u = 0; u < s; u++) {
         if (!vis[u]){
             forward_dfs(u, g, vis, stack);
    }
    reverse(stack.begin(), stack.end());
     fill(vis.begin(), vis.end(), false);
    int c = 0:
    for (int u: stack){
         if (!vis[u]){
             backward_dfs(u, gt, vis, comp, c);
    return {comp, c};
7.10 lexicographical_minimal_toposort.cpp
#include <bits/stdc++.h>
using namespace std;
#define i64 int64_t
// toposort with bfs, returns the lexicographically smallest
```

```
topological order
// Runs in O(n log n) because of sorting
// Returns boolean indicating if topo order exists, and the order in
   case it does
pair < bool, vector < i64>> lexicographical_toposort(const
   vector < vector < i64 >> & dag) {
    const i64 n = dag.size();
    // dependencies_met[i]: how many j such that j->i
    // have received their order number
    vector < i64 > dependencies_met(n, 0),
                in_degree(n, 0),
                order(n, -1);
    for (i64 u = 0; u < n; u++){}
        for (i64 v: dag[u]){
            ++in_degree[v];
        }
    priority_queue < i64, vector < i64>, greater < i64>> q;
    // min pq guarentees lexicographical order and topo order
    // by only pushing when all dependencies are met
    for (i64 u = 0; u < n; u++){}
        if (!in_degree[u]) q.emplace(u);
    i64 c = 0;
    while (!q.empty()){
        i64 u = q.top();
        order[u] = c++;
        q.pop();
        for (i64 v: dag[u]){
            ++dependencies_met[v];
            if (dependencies_met[v] == in_degree[v]){
                q.emplace(v);
        }
    }
    for (i64 u = 0; u < n; u++){}
        // circular dependencies exist
        if (order[u] == -1) return {false, {}};
    return {true, order};
7.11 tarjan.cpp
#include <bits/stdc++.h>
using namespace std;
const int INF = (int)2e9;
namespace Tarjan {
    vector < int > lowlink, t_in;
    vector < pair < int , int >> dfs_tree_edges;
    int timer;
```

void dfs(int u, const vector<vector<int>> &g, vector<bool> &vis,

```
int parent = -1){
    vis[u] = true:
    t_in[u] = lowlink[u] = timer++;
    for (int to: g[u]){
        if (to == parent) continue;
        else if (vis[to]){
            // back edge on DFS tree
            lowlink[u] = min(t_in[to], lowlink[u]);
        } else {
            dfs(to, g, vis, u);
            dfs_tree_edges.emplace_back(u, to);
            lowlink[u] = min(lowlink[to], lowlink[u]);
        }
    }
}
// lowlink(u) = min(t_in(u), min_{p>u}t_in(u), min_{u->to} em
   tree lowlink (to))
vector < vector < pair < int , int >>> dfs_tree_bridges (const
   vector < vector < int >> &g) {
    const int n = g.size();
    vector < bool > vis(n, false);
    lowlink = vector < int > (n, INF);
    t_in = vector < int > (n, INF);
    dfs_tree_edges.clear();
    timer = 0;
    for (int u = 0; u < n; u++) {
        if (!vis[u]){
            dfs(u, g, vis);
        }
    vector < vector < pair < int , int >>> h(n);
    for (auto[v, to]: dfs_tree_edges){
        if (lowlink[to] > t_in[v]){
            h[v].emplace_back(to, 1);
        } else {
            h[v].emplace_back(to, 0);
    return h:
Number Theory
```

### 8.1 fastexp.cpp

```
#include <bits/stdc++.h>
using namespace std;
const int MOD = (int)1e9 + 7; // NOTE: MIGHT CHANGE
```

```
int bexp(int a, int p){
    if (p == 0){
        return 1;
    } else {
        int m = bexp(a, p / 2);
        if (p % 2 == 1){
            return (((m * m) % MOD) * a) % MOD;
        } else {
            return (m * m) % MOD;
        }
    }
int inv(int a){
    return bexp(a, MOD - 2);
8.2 spf.cpp
Given N, calculates the smallest prime factor (spf) that divides K for
   each 1 \leftarrow K \leftarrow N
If K is prime, that prime factor will be K itself
#include <bits/stdc++.h>
using namespace std;
void spf(int N, int ans[]){
  // ans is a N+1 sized array
  for (int i = 0; i <= N; i++){</pre>
    ans[i] = i:
  int is_prime[N+1]; // O means is prime, 1 means is composite
  memset(is_prime, N + 1, 0);
  is_prime[0] = 1; is_prime[1] = 1;
  for (int i = 2; i <= N; i++){</pre>
    if (is_prime[i] == 0){
      for (int mul = i; mul <= N; mul += i){</pre>
        ans[mul] = min(ans[mul], i);
        // last value will be the smallest divisor of mul greater than
            1 (which is the spf)
      }
    Range Queries
```

### 9.1 bit.cpp

```
#include <bits/stdc++.h>
using namespace std;
vector<int> BIT(const vector<int> &a){
    int n = a.size():
    vector < int > B(n + 1, 0);
    for (int i = 0; i < n; i++){
        update(B, i, a[i]); // sum a[i] to position i+1 of the BIT
   return B;
int get(const vector<int> &BIT, int i){
   int s = 0;
    while (i \ge 0)
        s += BIT[i];
       // i & -i == 2^lsb(i)
       i = i ^ (i & (-i));
   return s;
void update(vector<int> &BIT, int i, int x){
    const int n = BIT.size() - 1;
    while (i \le n)
        BIT[i] += x;
        i += i & (-i);
9.2 segtree.cpp
#include <bits/stdc++.h>
using namespace std;
// default seg tree (in this example, a Max Seg)
struct SegTree {
    vector<int> seg;
    vector<int> a:
    int n;
    SegTree(const vector<int>& a) : a(a), n(a.size()) {
        seg.resize(4 * n);
        build(0, n - 1, 0);
    inline int merge(int a, int b){
        return max(a, b); // <--- CHANGE HERE</pre>
   }
    void build(int 1, int r, int node){
        if (1 == r) seg[node] = a[1];
        else {
            const int m = 1 + (r - 1) / 2;
            build(1, m, 2 * node + 1);
```

```
build(m + 1, r, 2 * node + 2);
        seg[node] = merge(seg[2 * node + 1], seg[2 * node + 2]);
    }
}
int query(int 1, int r){
    return _query(1, r, 0, n - 1, 0);
void update(int p, int x){
    _{update(p, x, 0, n - 1, 0)};
int _query(int 1, int r, int t1, int tr, int node){
    if (1 == t1 && r == tr){
        return seg[node];
    } else {
        int tm = tl + (tr - tl) / 2;
        if (r <= tm){
            return _query(1, r, tl, tm, 2 * node + 1);
        } else if (1 > tm){
            return _{query}(1, r, tm + 1, tr, 2 * node + 2);
        } else {
            return merge (
                _{query(1, tm, tl, tm, 2 * node + 1),}
                _{query(tm + 1, r, tm + 1, tr, 2 * node + 2)}
            );
        }
void _update(int p, int x, int tl, int tr, int node){
    if (tl == p && tr == p){
        seg[node] = x;
        a[p] = x;
    } else {
        int tm = tl + (tr - tl) / 2;
        if (p <= tm){</pre>
            _{update(p, x, tl, tm, 2 * node + 1)};
            _update(p, x, tm + 1, tr, 2 * node + 2);
        seg[node] = merge(seg[2 * node + 1], seg[2 * node + 2]);
    }
}
```

};

### 10 Strategy

### 10.1 Debugging:

- Nem sempre o bug está na parte mais complexa do código: leia tudo, até o template
- Verifique o valor do MOD, do INF e das constantes, dos tamanhos dos arrays.
- Criar corner cases (n = 0, 1, casos especiais que forçam alguma coisa, etc) e casos pequenos
- Colocar asserts para detectar índices out-of-bounds ou outros problemas
- Minimizar estado global, separar lógica em funções menores

### 10.2 Início (primeira 1h30)

- Ler questões
- Passar no máximo 10-15 min por questão
- Implementar apenas as fáceis
  - Priorizar primeiro aquelas que são simples e rápidas de implementar (em detrimento por ex. de uma aplicação trivial de um algoritmo longo, como uma SegTree ou fluxo)
- Fase mais individual
- Priorizar para implementação aquela pessoa que tem a ideia mais clara da solução

### 10.3 Meio

- Ler o restante da prova
- Olhar os standings para selecionar a próxima questão
- Analisar mais profundamente no papel para conseguir passar as médias/difíceis
- Evitar ficar três pessoas em uma mesma questão

### 10.4 Final (última 1h30)

• Foco do time todo em uma ou duas questões mais prováveis de passarem

### 11 Strings

### 11.1 rabin\_karp.cpp

```
#include <bits/stdc++.h>
using namespace std;
#define i64 int64 t
vector < i64 > rollingHash(const vector < int > & v, i64 p, i64 m) {
    const i64 n = v.size();
    vector < i64 > prefH(n):
    vector < i64 > p_pow(n);
    // prefH = hash(v[0...i]) = sum(j <= i) v[j] * p^j mod m
    p_pow[0] = 1;
    prefH[0] = v[0] % m;
    for (i64 i = 1; i < n; i++){
        p_pow[i] = (p_pow[i - 1] * p) % m;
        prefH[i] = (v[i] * p_pow[i] + prefH[i - 1]) % m;
    return prefH;
}
// lowercase
vector<i64> rollingHash(const string& s){
    vector<int> s_repr(s.size());
    for (i64 i = 0; i < s.size(); i++){</pre>
        s_repr[i] = s[i] - 'a';
    return rollingHash(s_repr, 31, (i64)1e9 + 7);
}
i64 bexp(const i64 a, const i64 p, const i64 mod){
    if (p == 0) {
        return 1:
    } else {
        const i64 b = bexp(a, p / 2, mod);
        if (p \% 2 == 0){
            return (b * b) % mod;
            return (a * ((b * b) % mod)) % mod;
    }
}
vector<i64> prime_powers(const i64 n, const i64 p = 31, const i64 m =
   (i64)1e9 + 7){}
    vector < i64 > p_pow(2 * n + 1);
    p_pow[n] = 1;
    p_pow[n-1] = bexp(p, m-2, m); // fermat little theorem
    for (i64 i = 1: i <= n: i++){
        p_pow[n + i] = (p_pow[n + i - 1] * p) % m;
    for (i64 i = 2; i <= n; i++){</pre>
        p_pow[n - i] = (p_pow[n - i + 1] * p_pow[n - 1]) % m;
```

```
return p_pow;
vector<i64> rabinKarpMatch(const string& s, const string& t){
    vector<i64> prefix_s = rollingHash(s), prefix_t = rollingHash(t);
    const i64 mod = (i64)1e9 + 7, p = 31;
    const i64 HS = prefix_s.back(); // complete hash of s
    // precompute prime powers mod M
    const vector<i64> p_pow = prime_powers(t.size(), p, mod);
    vector < i64 > matches; // starting points of matches
    const i64 offset = t.size():
    for (i64 i = 0; i + s.size() <= t.size(); i++){</pre>
        // hash of t[i...i + |s| - 1] == (hash(t[0...i + |s| - 1]) -
           hash(t[0...i-1])) * p^-i
        const i64 h = ((prefix_t[i + s.size() - 1] - (i > 0 ?
            prefix_t[i - 1] : 0) + mod) * p_pow[offset - i]) % mod;
        if (h != HS){
            continue:
        // compare
        bool ok = true;
        for (i64 j = 0; (j < s.size()) && ok; j++){
            ok = ok && (s[i] == t[i + i]);
        if (ok) matches.emplace_back(i);
   return matches;
11.2 suffix_array.cpp
// TESTED on first suffix array problem of Codeforces EDU
#include <bits/stdc++.h>
using namespace std;
#define ALPHABET 256
void csort(vector<int>& p, const vector<int>& classes, int offset){
    const int n = p.size();
    vector < int > cnt(max(ALPHABET, n), 0), ptemp(n);
    for (int i = 0; i < n; i++){</pre>
        ++cnt[classes[i]];
    for (int i = 1; i < max(ALPHABET, n); i++){
        cnt[i] += cnt[i - 1];
    // distribute on blocks in reverse order to make sorting stable
    // distribute according to order of k-th element (k=offset)
    for (int i = n - 1; i \ge 0; i--){
        int offseted_pos = (p[i] + offset) % n;
        ptemp[--cnt[classes[offseted_pos]]] = p[i];
```

```
swap(ptemp, p);
pair < vector < int >, vector < vector < int >>> suffix_array(const string& s) {
    int lg = 0;
    string t = s;
    t.push_back('$');
    int n = t.size();
    while ((111 << lg) < n) ++ lg; // now 2^lg >= n
    vector < int > p(n);
    vector < vector < int >> c(lg + 1, vector < int > (n));
    for (int i = 0: i < n: i++){
        p[i] = i;
        c[0][i] = (int)t[i];
    }
    csort(p, c[0], 0);
    // check offset of size 2^k
    for (int k = 0; k \le lg; k++) {
        csort(p, c[k], (111 << k)); // sort by second, stably
        csort(p, c[k], 0); // sort by first, stably
        if (k == lg) break;
        // redo equivalence classes
        c[k + 1][p[0]] = 0: // first element receives class=0
        int cls = 0;
        for (int i = 1; i < n; i++) {
            if (
                (c[k][p[i-1]] < c[k][p[i]])
                || (c[k][(p[i-1] + (111 << k)) % n] < c[k][(p[i] +
                    (111 << k)) % n])
            ) {
                ++cls;
            c[k + 1][p[i]] = cls;
        }
    }
    return {p, c};
int lcp(int i, int j, const vector<vector<int>>& cls){
    const int n = cls[0].size():
    int lg = 0;
    while ((111 << lg) < n) ++ lg; // now 2^lg >= n
    int ans = 0;
    for (int k = lg; k >= 0; k--){
        if (cls[k][i] == cls[k][j]){
            i = (i + (111 << k)) \% n;
            j = (j + (111 << k)) \% n;
            ans += (111 << k);
       }
    return ans;
// LCS in O(N log N)
// TESTED on second suffix array problem of Step 2 of Codeforces EDU
tuple<int, int, int> lcs(string s, string t){
```

```
auto [p, c] = suffix_array(s + "#" + t);
    // NOTE: the sentinel must be different to avoid having cyclic
       shifts breaking stuff
    const int n = s.size(), m = t.size();
    const int sz = n + 1 + m + 1; // middle # and final $
    vector<int> lcp_array(sz - 1);
    for (int i = 0; i + 1 < sz; i++) lcp_array[i] = lcp(p[i], p[i +
       1]. c):
    int lcs_sz = 0, s_idx = -1, t_idx = -1;
    for (int i = 0; i + 1 < sz; i++){
        int left = p[i], right = p[i + 1];
        if (lcp_array[i] > lcs_sz && min(left, right) < n && max(left,</pre>
           right) > n){
            // p[i] in s, p[i + 1] in t or vice-versa
            lcs_sz = lcp_array[i];
            s_idx = min(left, right);
            t_idx = max(left, right);
        }
   }
    return {lcs_sz, s_idx, t_idx};
11.3 trie.cpp
#include <bits/stdc++.h>
using namespace std;
template < int K > // sum of characters in the set
struct Trie {
    int children[K][26];
    int cnt[K]; // cnt of each word
    int ptr = 1;
   Trie(){
        for (int i = 0; i < K; i++){</pre>
            for (int j = 0; j < 26; j++) {
                children[i][j] = -1;
            cnt[i] = 0;
   }
    void insert(string& u, int i=0, int node=0){
        int idx = u[i] - 'a';
        if (children[node][idx] == -1){
            // create node
            children[node][idx] = ptr++;
        if (i + 1 < u.size()){</pre>
            insert(u, i + 1, children[node][idx]);
        } else {
            ++cnt[children[node][idx]];
   }
```

```
int search(string &u, int i=0, int node=0){
    if (i == u.size()){
        return cnt[node]:
    } else {
        int idx = u[i] - 'a';
        if (children[node][idx] == -1){
            return 0:
        } else {
            return search(u, i + 1, children[node][idx]);
        }
    }
}
bool isleaf(int node){
    bool ok = true;
    for (int i = 0; i < 26; i++){
        ok = ok && children[node][i] == -1;
    return ok;
}
void remove(string &u, int i=0, int node=0){
    int idx = u[i] - 'a';
    if (children[node][idx] != -1){
        remove(u, i + 1, children[node][idx]);
        if (isleaf(children[node][idx]) &&
           cnt[children[node][idx]] == 0){
            children[node][idx] = -1;
    } else {
        // reached leaf
        cnt[node] = max(cnt[node] - 1, 0);
    }
```

### 12 template.cpp

};

```
#define TESTCASES

#include <bits/stdc++.h>
using namespace std;
#define fastio
    ios_base::sync_with_stdio(false);cin.tie(NULL);cout.tie(NULL)
#define endl '\n'

#define i64 int64_t
#define i128 __int128
#define all(x) begin(x),end(x)
#define print(msg, v) cerr << msg; for(auto it = v.begin(); it != v.end(); it++){cerr << *it << " ";} cerr << endl;
#define printgraph(msg, G) cerr << msg << endl; \
    for (i64 u = 0; u < G.size(); u++) {
        cerr << "G[" << u << "]="; \
        for (i64 v: G[u]) {</pre>
```

```
cerr << v << " "; \
        } \
        cerr << endl: \
    }
template <typename T, typename U>
ostream& operator << (ostream& out, pair <T, U> x)
    out << x.first << " " << x.second;
    return out;
template <typename T>
ostream& operator << (ostream& out, vector <T> v)
    for (const auto& x: v){
        out << x << " ":
    return out;
template <typename T>
using min_pq' = priority_queue<T, vector<T>, greater<T>>;
template <typename T>
using max_pq = priority_queue <T>;
void solve(){}
signed main(){
    fastio;
    int t = 1
#ifdef TESTCASES
    cin >> t:
#endif
    while (t--) {
        solve();
```

### 13 Trees

### 13.1 binary\_lifting.cpp

```
#include <bits/stdc++.h>
using namespace std;

const int N = 2e5, LOGN = 20;
int lift[LOGN + 1][N];

void binary_lifting(const vector<vector<int>> &forest){
    const int n = forest.size();
    vector<int> parent(n, -1);
    for (int u = 0; u < n; u++){
        for (int v: forest[u]){
            parent[v] = u;
        }
}</pre>
```

```
for (int u = 0; u < n; u++) {
        lift[0][u] = parent[u];
    for (int k = 1: k \le LOGN: k++) {
        for (int u = 0; u < n; u++) {
            lift[k][u] = lift[k - 1][lift[k - 1][u]];
int jump(int u, int k){
    for (int i = 0; i <= LOGN; i++){</pre>
        if ((k >> i) & 1) {
            u = lift[i][u]:
    return u;
13.2 centroid.cpp
#include <bits/stdc++.h>
using namespace std;
namespace CentroidDecomposition {
    int subtree_size(int u, const vector<vector<int>>& adj,
       vector<int>& s, vector<int>& blocked, int p = -1){
        s[u] = 1:
        for (int v: adj[u]){
            if (v == p) continue;
            if (blocked[v]) continue;
            s[u] += subtree_size(v, adj, s, blocked, u);
        return s[u]:
    }
    // find centroid of subtree
    int find_centroid(int u, const vector < vector < int >> & adj,
       vector<int>& s, vector<int>& blocked, int total_size, int p =
       -1){
        for (int v: adj[u]){
            if (v == p) continue;
            if (blocked[v]) continue:
            if (2 * s[v] >= total_size){
                return find_centroid(v, adj, s, blocked, total_size,
                    u); // there can be only one such vertex
            }
        return u;
    // total_size = size of subtree rooted at root (considering
       deletions)
    // @return centroid subtree root
    int _build_centroid_tree(
```

```
const vector < vector < int >> & adj,
        vector < int > & blocked,
        vector < vector < int >>& centroid_adj,
        vector < int > % sz,
        int total size.
        int root = 0
    ) {
        subtree_size(root, adj, sz, blocked);
        int c = find centroid(root, adi, sz. blocked, total size);
        blocked[c] = true:
        for (int v: adj[c]){
            if (blocked[v]) continue;
            // build centroid subtree on partition and get its root
            int c2 = _build_centroid_tree(adj, blocked, centroid_adj,
                sz, sz[v], v);
            centroid_adj[c].emplace_back(c2);
        return c;
    }
    // @return pair (centroid tree root, directed centroid tree)
    pair < int , vector < vector < int >>> centroid_tree (const
        vector < vector < int >> & adj) {
        const int n = adj.size();
        vector < int > blocked(n, false), sz(n);
        vector < vector < int >> centroid_adj(n);
        int master_centroid = _build_centroid_tree(
            adj, blocked, centroid_adj, sz, n
        return {master_centroid, centroid_adj};
13.3 hld_no_range_update.cpp
// HLD implementation: no range updates
// Tested on CSES Path Queries (https://cses.fi/problemset/task/1138)
// Query: sum of vertices from root to v
// Update: set value of v to x
// uses multiple seg trees. Marginally faster than a single segtree
// see hld_with_range_update for better descriptions
#include <bits/stdc++.h>
using namespace std;
#define i64 int64_t
using tree = vector<vector<i64>>;
struct HLD {
    vector < i64 > heavy;
    vector < i64 > sz;
    vector < i64 > head:
    vector < i64 > pos;
    vector < i64 > parent;
    // NOTE: SegTree is the segment tree implementation
    // NOTE: SegTree needs default constructor
```

vector < SegTree > segs;

```
vector < i64 > a:
vector < i64 > depth;
HLD(const tree& adj, const vector < i64 > & _a, i64 root = 0) : a(a) {
    const i64 n = adj.size();
    heavy.resize(n);
    sz.resize(n):
    head.resize(n):
    pos.resize(n);
    parent.resize(n);
    segs.resize(n);
    subtree(root, adj);
    dfs(root, adi);
    // build segment trees
    for (i64 h = 0; h < n; h++){}
        if (h == head[h]){
            vector \langle i64 \rangle path;
            i64 \text{ cur} = h:
            while (cur != -1 && head[cur] == h){
                 path.emplace_back(a[cur]); // add value
                 cur = heavy[cur];
            segs[h] = SegTree(path);
    }
void subtree (i64 u, const tree & adj, i64 p = -1, i64 d = 0) {
    sz[u] = 1;
    depth[u] = d;
    for (i64 v: adj[u]){
        if (v == p) continue;
        subtree(v, adj, u, d + 1);
        sz[u] += sz[v];
    }
}
void dfs(i64 u. const tree& adi. i64 p = -1. i64 h = 0. i64 d = 0){
    i64 cur_sz = 0;
    pos[u] = d;
    head[u] = h;
    parent[u] = p:
    heavy[u] = -1;
    for (i64 v: adj[u]){
        if (v == p) continue;
        if (sz[v] > cur sz)
            cur_sz = sz[v];
            heavy[u] = v:
        }
    for (i64 v: adj[u]){
        if (v == p) continue;
        if (v == heavy[u]){
```

```
dfs(v, adj, u, h, d + 1);
                 dfs(v, adj, u, v, 0);
        }
    }
    // query from u to p^min(k, depth[u])(u), inclusive
     // NOTE: SUM QUERY HERE, other variations are possible
    i64 query_up(i64 u, i64 k){
        i64 s = 0;
         while (u != -1 \&\& k >= pos[u]) {
             s += segs[head[u]].query(0, pos[u]);
             u = parent[head[u]];
        if (u != -1){ // k < pos[u]</pre>
             s += segs[head[u]].query(pos[u] - k, pos[u]);
        return s;
     i64 query_path(i64 u, i64 v){
        // NOTE: need to implement LCA separetely!
         i64 l = lca(u, v);
         return query_up(u, depth[u] - depth[l]) + query_up(v, depth[v]
            - depth[1]) - a[1];
    }
     void update(i64 u, i64 x){ // set value to x
         segs[head[u]].update(pos[u], x);
};
13.4 hld_with_range_update.cpp
// HLD implementation
// Tested on CSES Path Queries (https://cses.fi/problemset/task/1138)
// and range update tested on problem L from ICPC Gran Premio de
    Mexico 2025 3ra fecha
 // Query: sum of vertices from root to v
// Update: set value of v to x
// Call init_seg before using
// Implementation note: this is an implementation where
// you use a single seg-tree and the node positions are their
// time of visit in the DFS, where you **visit heavy edges first**
// (so heavy paths make contiguous ranges)
// However, this leads to a larger segtree and can be slower!
// An alternative is to use an array of segtrees (one for each path)
#include <bits/stdc++.h>
using namespace std;
#define i64 int64_t
// heavy edge is the edge with the largest subtree
\mid // going through light (i.e not heavy) edge reduces size by half
```

```
// heavy path goes up to just the vertex before light edge
template < size t N> struct HLD {
    const vector < vector < i64 >> adi:
    vector < i64 > a;
    i64 parent[N], sz[N], depth[N],
        heavy[N], // heavy[u] = v <==> uv is heavy edge
        head[N]. // head[u] = start of heavy path that goes through u
        pos[N]: // visit pos of vertex in dfs
    i64 global_pos = 0;
    i64 n, root;
    HLD(const vector < vector < i64 >> & adj, i64 root = 0) :
        adj(adj), n(adj.size()), root(root)
    {
        subtree(root);
        dfs(root);
    }
    void init_seg(const vector < i64 > & vec) {
        a = vec;
        vector < i64 > b(n);
        for (i64 i = 0; i < n; i++) b[pos[i]] = a[i];</pre>
        seg::init(b):
    }
    void subtree (i64 u, i64 p = -1, i64 d = 0) {
        sz[u] = 1; depth[u] = d; parent[u] = p;
        for (i64 v: adj[u]){
            if (v == p) continue;
            subtree(v, u, d + 1);
            sz[u] += sz[v];
        }
    }
    void dfs(i64 u, i64 p = -1, i64 h = 0){
        i64 cur_sz = 0;
        pos[u] = global_pos++;
        head[u] = h;
        heavy[u] = -1;
        for (i64 v: adj[u]){
            if (v == p) continue;
            if (sz[v] > cur_sz){
                cur_sz = sz[v];
                heavv[u] = v;
            }
        }
        if (heavy[u] != -1) dfs(heavy[u], u, h);
        for (i64 v: adi[u]){
            if (v == p || v == heavy[u]) continue;
            dfs(v, u, v);
        }
    }
    // query from u to p^{k} - 1(u)
    // NOTE: sum query here, other variations are possible
    i64 query_up(i64 u, i64 k){
```

```
i64 s = 0;
    while (u != -1 \&\& k >= pos[u] - pos[head[u]] + 1){ // full}
        s += seg::query(pos[head[u]], pos[u]);
        k -= pos[u] - pos[head[u]] + 1;
        u = parent[head[u]];
    if (u != -1 \&\& k >= 1){
        s += seg::query(pos[u] - (k - 1), pos[u]);
    return s;
i64 lca(i64 u, i64 v){
    if (pos[v] < pos[u]) swap(u, v);
    // now u is above or in another branch
    if (head[u] == head[v]) return u;
    else return lca(u, parent[head[v]]);
// query when values are on vertices
i64 query_path_vertex(i64 u, i64 v){
    i64 l = lca(u, v);
    return (query_up(u, depth[u] - depth[l]) + query_up(v,
        depth[v] - depth[1]) + query_up(1, 1));
}
// query when a[i] = value for edge i->parent[i]
i64 query_path_edge(i64 u, i64 v){
    i64 l = lca(u, v);
    return query_up(u, depth[u] - depth[l]) + query_up(v, depth[v]
        - depth[1]);
}
void update_pt(i64 u, i64 x){ // set value to x
    seg::update_pt(pos[u], x);
    a[u] = x:
void update_path_vertices(i64 u, i64 v, i64 x){ // same stuff as
    if (pos[v] < pos[u]) swap(u, v);</pre>
    // now u is above or in another branch
    if (head[u] == head[v]) seg::update_range(pos[u], pos[v], x);
    else {
        seg::update_range(pos[head[v]], pos[v], x);
        update_path_vertices(u, parent[head[v]], x);
    }
}
void update_path_edges(i64 u, i64 v, i64 x){ // same stuff as lca
    if (pos[v] < pos[u]) swap(u, v);
    // now u is above or in another branch
    if (head[u] == head[v]) {
        if (pos[u] < pos[v]) seg::update_range(pos[u] + 1, pos[v],</pre>
```

```
x);
        }
        else {
            seg::update_range(pos[head[v]], pos[v], x);
            update_path_edges(u, parent[head[v]], x);
        }
    }
    // subtree of v is entered at pos[v] and visited in contiguous
    i64 query_subtree(i64 u){
        return seg::query(pos[u], pos[u] + sz[u] - 1);
    }
    void update_subtree(i64 u, i64 x){
        seg::update_range(pos[u], pos[u] + sz[u] - 1, x);
    }
};
13.5 lca.cpp
// tested on CSES company queries II
#include <bits/stdc++.h>
using namespace std;
#define i64 int64_t
namespace lca {
    const i64 N = 2e5 + 100, LOGN = 20;
    i64 parent[N], level[N];
    i64 blift[LOGN + 1][N];
    void dfs(i64 u, const vector < vector < i64 >> & adj, i64 p = -1, i64 d
       = 0){}
        parent[u] = p;
        level[u] = d;
        for (i64 v: adj[u]){
            if (v == p) continue;
            dfs(v, adj, u, d + 1);
        }
    void init(const vector < vector < i64 >> & adj, i64 root = 0) {
        const i64 n = adj.size();
        dfs(root, adj);
        parent[root] = root; // fix root
        for (i64 u = 0; u < n; u++) blift[0][u] = parent[u];</pre>
        for (i64 k = 1; k <= LOGN; k++) {
            for (i64 u = 0; u < n; u++){}
                blift[k][u] = blift[k - 1][blift[k - 1][u]];
        }
    }
    i64 up(i64 u, i64 k){
        for (i64 i = 0; i <= LOGN; i++){</pre>
            if (k & (111 << i)){
                u = blift[i][u];
            }
```

```
return u;
    i64 lca(i64 u, i64 v){
        if (level[v] < level[u]) swap(u, v);</pre>
        v = up(v, level[v] - level[u]);
        if (u == v) return u;
        // loop invariant: u and v are distinct
        // (have not reached a common ancestor yet)
        for (i64 i = LOGN; i >= 0; i--){}
            if (blift[i][u] != blift[i][v]){
                // keeps loop invariant -> greedily take it!
                u = blift[i][u];
                v = blift[i][v]:
        // at the end of the invariant we must be just one level below
        // otherwise we could have jumped more
        return parent[u];
13.6 path_queries.cpp
// Tested on CSES Distance Queries
// sum queries on paths. Values are on vertices
#include <bits/stdc++.h>
using namespace std;
const int N = 2e5, LOGN = 20;
int lift[LOGN + 1][N];
int sums[LOGN + 1][N];
void binary_lifting(const vector<int>& parent, const vector<int>&
   values){
    const int n = parent.size();
    for (int u = 0; u < n; u++) {
        lift[0][u] = parent[u];
        sums[0][u] = values[u]:
    for (int k = 1; k \le LOGN; k++) {
        for (int u = 0; u < n; u++) {
            lift[k][u] = lift[k - 1][lift[k - 1][u]];
            // u..p^(2^{k - 1})(u) and then p^(2^{k - 1}+1)(u) to
                p^{(2k(u))}
            sums[k][u] = sums[k - 1][u] + sums[k - 1][lift[k - 1][u]];
        }
   }
int jump(int u, int k){
    for (int i = 0; i <= LOGN; i++){</pre>
```

```
if ((k >> i) & 1){
            u = lift[i][u];
        }
    return u;
// v[u] + v[p[u]] + ... + v[p^{2^k - 1}(u)]
int query_sum(int u, int k){
    int s = 0;
    for (int i = 0; i <= LOGN; i++){</pre>
        if ((k >> i) & 1) {
            s += sums[i][u];
            u = lift[i][u];
        }
    return s;
int LCA(
    int u, int v, const vector < int > &level,
    const vector <int > &parent
) {
    // invariant: level[u] <= level[v]</pre>
    if (level[v] < level[u]) swap(u, v);</pre>
    v = jump(v, level[v] - level[u]);
    if (u == v){
        return u;
    // loop invariant: u and v are distinct
    // (we have not reached a common ancestor yet)
    for (int i = LOGN; i >= 0; i--){
        if (lift[i][u] != lift[i][v]){
            // keeps loop invariant -> greedily take it!
            u = lift[i][u];
            v = lift[i][v];
        }
    }
    // at the end of the invariant we must be just one level below the
    // otherwise we could have jumped more
    return parent[u];
}
int path_sum(int u, int v, const vector<int> &level, const vector<int>
   &parent){
    int lca = LCA(u, v, level, parent);
    int left = query_sum(u, level[u] - level[lca]), right =
        query_sum(v, level[v] - level[lca]);
    return left + right + sums[0][lca];
```