ICPC Lib

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1 binary_lifting.cpp

```
#include <bits/stdc++.h>
using namespace std;
const int N = 2e5, LOGN = 20;
int lift[LOGN + 1][N];
void binary_lifting(const vector<vector<int>> &forest){
    const int n = forest.size();
    vector < int > parent(n, -1);
    for (int u = 0; u < n; u++) {
        for (int v: forest[u]){
            parent[v] = u;
    }
    for (int u = 0; u < n; u++){
        lift[0][u] = parent[u];
    for (int k = 1; k <= LOGN; k++){</pre>
        for (int u = 0; u < n; u++) {
            lift[k][u] = lift[k - 1][lift[k - 1][u]];
int jump(int u, int k){
    for (int i = 0; i <= LOGN; i++){</pre>
        if ((k >> i) & 1) {
            u = lift[i][u]:
    return u;
```

2 Combinatorics

2.1 Fórmula de Legendre

Para primo p, $v_p(n!) = \sum_{k=1}^{\infty} \left\lfloor \frac{n}{p^k} \right\rfloor$, e $v_p\left(\binom{n}{k}\right) = v_p(n!) - v_p(k!) - v_p((n-k)!)$.

Para
$$m = \prod p_i^{e_i}$$
 composto, $v_m(n!) = \min_i \left\lfloor \frac{v_{p_i}(n!)}{e_i} \right\rfloor$.

2.2 Teorema de Lucas

Para p primo e $n = n_k p^k + ... + n_1 p + n_0$ e $m = m_k p^k + ... + m_1 p + m_0$,

$$\binom{n}{m} \equiv \prod_{i=0}^{k} \binom{n_i}{m_i} \mod p$$

2.3 Teorema de Sprague-Grundy e game-theory

Nim: ganha se e só se o XOR dos números da pilha é maior que 0: $a_1 \oplus \ldots \oplus a_n \neq 0$.

Estados terminais tem número de Grundy n(v) = 0 e não-terminais tem número de Grundy $n(v) = \max_{v \to u} n(u)$. Assim, um estado é **vencedor** (em um jogo em que quem não tem movimentos perde) se e só se $n(v) \neq 0$. Um único jogo é equivalente a uma **pilha de Nim com (limitados) acréscimos**.

Assim, se temos dois "sub-jogos" independentes S_1 e S_2 então seu número de Grundy é o XOR dos números de Grundy individuais, $n(S_1 + S_2) = n(S_1) \oplus n(S_2)$.

3 count_inversions_merge_sort.cpp

```
#include <bits/stdc++.h>
using namespace std;
// sorts and returns count of inversions. Modifies original array!
int inversions(vector<int>& a, int 1, int r){
    if (1 == r) return 0;
    const int m = 1 + (r - 1) / 2;
    int total = inversions(a, 1, m) + inversions(a, m + 1, r);
    vector<int> left, right;
    for (int i = 1; i <= m; i++) left.emplace_back(a[i]);</pre>
    for (int i = m + 1; i <= r; i++) right.emplace_back(a[i]);</pre>
    int idx = 1:
    int pl = 0, pr = 0;
    while (pl < left.size() || pr < right.size()){</pre>
        // we only count the inversion pairs from the perspective of
        // element on the right
        // 1 3 5 | 2 4
        // if right[pr] < left[pl] then it is smaller than
            [left[pl]..[left[|left|-1]]]
        // so contributes for left-pl cross-inversions
        if (pr == right.size())
            a[idx++] = left[pl++];
        else if (pl == left.size())
            a[idx++] = right[pr++];
        else {
```

};

```
struct DSU {
    vector < int > parent, sz;
    vector < vector < int >> elements;
    DSU(int n){
        parent.resize(n);
        iota(parent.begin(), parent.end(), 0);
        sz.assign(n, 1);
        elements.resize(n);
        for (int i = 0; i < n; i++){
            elements[i].emplace_back(i);
    }
    int find(int x){
        if (parent[x] != x){
            parent[x] = find(parent[x]); // path-compression
        return parent[x];
    }
    void unite(int x, int y){
        x = find(x):
        y = find(y);
        if (x == y){
            return;
        } else if (sz[x] > sz[y]) {
            swap(x, y);
        }
        sz[y] += sz[x];
        for (int z: elements[x]){
            elements[y].emplace_back(z);
        parent[x] = y;
```

5 fft.cpp

```
// CODE TAKEN FROM CP-ALGORITHMS
#include <bits/stdc++.h>
using namespace std;
using cd = complex < double >;
const double PI = acos(-1):
void fft(vector < cd> & a, bool invert) {
    int n = a.size();
    for (int i = 1, j = 0; i < n; i++) {
        int bit = n >> 1;
        for (; j & bit; bit >>= 1)
               ^= bit;
        j ^= bit;
        if (i < j)
            swap(a[i], a[j]);
    }
    for (int len = 2; len <= n; len <<= 1) {</pre>
        double ang = 2 * PI / len * (invert ? -1 : 1);
        cd wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i += len) {</pre>
            cd w(1);
            for (int j = 0; j < len / 2; j++) {
                 cd u = a[i+j], v = a[i+j+len/2] * w;
                 a[i+j] = u + v;
                 a[i+j+len/2] = u - v;
                w *= wlen;
            }
        }
    }
    if (invert) {
        for (cd & x : a)
            x /= n;
}
vector<int> multiply(vector<int> const& a, vector<int> const& b) {
    vector < cd > fa(a.begin(), a.end()), fb(b.begin(), b.end());
    int n = 1:
    while (n < a.size() + b.size())</pre>
        n <<= 1;
    fa.resize(n);
    fb.resize(n);
    fft(fa, false);
    fft(fb, false);
    for (int i = 0; i < n; i++)</pre>
        fa[i] *= fb[i];
    fft(fa, true);
    vector<int> result(n);
```

```
for (int i = 0; i < n; i++)
    result[i] = round(fa[i].real());
return result;</pre>
```

6 Geometry

6.1 convex_hull.hpp

```
#include "./point.hpp"
#include "./polygon.hpp"
#define i64 int64_t
// counterclockwise convex hull, can include collinear points or not
vector < pt > convex_hull (vector < pt > poly, bool INCLUDE_COLLINEAR) {
    // choose bottommost point as pivot
    pt pivot = poly[0];
   const i64 n = poly.size();
   for (i64 i = 1; i < n; i++){}
        if (poly[i].y < pivot.y || (poly[i].y == pivot.y && poly[i].x</pre>
           < pivot.x)){
            pivot = poly[i];
        }
   }
    sort(poly.begin(), poly.end(),
        [&pivot](pt p, pt q){return polarComp(pivot, p, q);}
    ): // pivot will be always poly[0] after sort
    // since including collinear points, we want to privilege the
       farthest collinear points
    // when coming back (e.g (0,0), (1,0), (2,0), (2,1), (2,2), (1,1))
    if (INCLUDE COLLINEAR) {
        i64 i = n - 1;
        while (i > 0 && orientation(pivot, poly[i], poly.back()) == 0){
        reverse(poly.begin() + i + 1, poly.end());
   }
    vector < pt > hull = {pivot};
    for (i64 i = 1: i < n: i++)
        while (
            hull.size() >= 2
            && (
                orientation(hull[hull.size() - 2], hull[hull.size() -
                    1], poly[i]) == 1
                || (!INCLUDE_COLLINEAR && orientation(hull[hull.size()
                    - 2], hull[hull.size() - 1], poly[i]) == 0)
            )
        ) {
            // goes clockwise, break orientation
            hull.pop_back();
        }
        hull.emplace_back(poly[i]);
    return hull;
```

```
6.2 point.hpp
#include <bits/stdc++.h>
using namespace std;
#define i64 int64_t
struct pt {
    i64 x, v;
    pt(i64 x, i64 y) : x(x), y(y) {}
    pt(const pt \& p) : x(p.x), y(p.y) \{\}
    i64 cross(pt p){
        <u>return</u> x * p.y - y * p.x;
    i64 dot(pt p){
        return \bar{x} * p.x + y * p.y;
    void operator=(const pt q){
        x = q.x;
        y = q.y;
    }
};
pt operator-(pt p1, pt p2){
    return pt(p1.x - p2.x, p1.y - p2.y);
pt operator+(pt p1, pt p2){
    return pt(p1.x + p2.x, p1.y + p2.y);
pt operator*(pt p1, i64 k){
    return pt(k * p1.x, k * p1.y);
pt operator*(i64 k, pt p1){
    return pt(k * p1.x, k * p1.y);
i64 orientation(pt p, pt q, pt r){
    i64 o = (q - p).cross(r - p);
    if (o > 0){
        return -1: // counterclockwise
    } else if (o < 0){</pre>
        return 1; // clockwise
        return 0; // collinear
i64 ccw(pt p, pt q, pt r) { return orientation(p, q, r) < 0; }
bool is_collinear(pt p, pt q, pt r){
```

return (q - p).cross(r - p) == 0;

```
// is p on segment q,r?
bool on_segment(pt p, pt q, pt r){
    // pgr is collinear
    // qp and pr have the same sign/orientation (as p in between q and
    return (
        is_collinear(p, q, r)
        && ((p - q).dot(r - p) >= 0)
    );
// counterclockwise polar sort
bool polarComp(pt pivot, pt p, pt q){
    auto d2 = [](pt a){}
        return a.x * a.x + a.y * a.y;
    i64 o = orientation(pivot, p, q);
    if (o == -1) return true;
    else if (o == 1) return false;
    else {
        // collinear => return closest point
        return d2(p - pivot) < d2(q - pivot);</pre>
// line intersection
bool line_line_intersect(pt a, pt b, pt c, pt d){
    return (
        (b - a).cross(d - c) != 0
        || (b - a).cross(c - a) == 0
    ); // not parallel, or equal lines (ab||cd and abc collinear)
// line ab, segment cd
bool line_segment_intersect(pt a, pt b, pt c, pt d){
    // c and d are in different half-planes
    // ==> different orientations
    // OR any triple abc, abd is collinear
    return orientation(c, a, b) * orientation(d, a, b) != 1;
// https://cp-algorithms.com/geometry/check-segments-intersection.html
bool segment_segment_intersect(pt a, pt b, pt c, pt d){
    auto inter1d = [](i64 11, i64 r1, i64 12, i64 r2){
        if (l1 < r1) swap(l1, r1);</pre>
        if (12 < r2) swap(12, r2);</pre>
        return max(11, 12) \le min(r1, r2);
    };
    if (orientation(a, b, c) == 0 \&\& orientation(a, b, d) == 0){
        // all four are collinear
        // reduces to 1d interval intersection
        return inter1d(a.x, b.x, c.x, d.x) && inter1d(a.y, b.y, c.y,
           d.y);
```

```
}
    return
        // different half-planes: line ab intersects segment cd
        (orientation(a, b, c) != orientation(a, b, d))
        // different half-planes: line cd intersects segment cd
        && (orientation(c, d, a) != orientation(c, d, b))
// horizontal ray from a, direction +-1, intersects segment cd?
// DOES include intersection at endpoints (cd) by default,
// verified by CSES "Point in Polygon"
bool horizontal_ray_segment_intersect(pt a, pt c, pt d, i64 direction
    if (c.y <= a.y && a.y < d.y){ // different half-planes from ray,</pre>
       c below or at line, d above
        // now we need d in the same half-plane as the ray
            (considering line ac)
        // so the intersection happens on the ray and not on the other
            side
        if (direction == 1){
            return orientation(c, a, d) > 0;
        } else {
            return orientation(c, a, d) < 0;</pre>
    } else if (d.y \le a.y \&\& a.y \le c.y){
        return horizontal_ray_segment_intersect(a, d, c, direction);
    return false:
6.3 polygon.hpp
#include "./point.hpp"
#include "./triangle.hpp"
// shoelace formula
double area(vector<pt>& p){
    i64 a = 0;
    i64 n = p.size();
    for (i64 i = 0; i < n; i++){</pre>
        a += p[i].cross(p[(i + 1) % n]); // area of OP[i]P[i+1]
    return (double)a / 2.0f;
// O(log n) check for boundary/inside convex polygon
// take the first pt (lexicographically) as reference (PO), and sort
   by polar order
// since polygon is convex, then P = U P[0]P[i]P[i+1] for i=1..n-2
// do binary search to find i, and then check for being inside triangle
// check for boundaries separately
const i64 INSIDE = 1;
const i64 BOUNDARY = 0;
const i64 OUTSIDE = -1;
i64 inside_polygon_convex(pt q, vector<pt> p){
```

```
const i64 n = p.size();
   pt pivot = p[0];
   for (pt a: p){
        if (make_pair(a.x, a.y) < make_pair(pivot.x, pivot.y)){</pre>
            pivot = a:
        }
   }
    sort(p.begin(), p.end(), [&pivot](pt a, pt b){return
       polarComp(pivot, a, b);});
    i64 \ 1 = 1, r = n - 2, i = -1;
    while (1 \le r)
        i64 m = 1 + (r - 1) / 2;
        // is in P[0]P[m]P[m+1] iff P[m]P[0]Q is ccw but P[m+1]P[0]Q
        if (ccw(p[m], p[0], q)){
            if (!ccw(p[m + 1], p[0], q)){
                i = m;
                break;
            } else {
                1 = m + 1;
        } else {
           r = m - 1;
    // inside_triangle includes boudaries
    if (i == -1 \mid | !inside_triangle(q, p[0], p[i], p[i + 1])) return
       OUTSIDE;
    else if (
        on_segment(q, p[i], p[i + 1])
        || (i == 1 && on_segment(q, p[0], p[1]))
        || (i == n - 2 \&\& on_segment(q, p[0], p[n - 1]))
   ) {
        return BOUNDARY;
   } else return INSIDE;
// Raycasting algorithm, works on all polygons, convex or not.
   Complexity: O(N)
// if the ray intersects a vertex, then it only counts if orientation
// (so it does not hit "tangentially", but cuts through)
i64 inside_polygon(pt q, vector <pt>& p){
    const i64 n = p.size();
    // draw horizontal ray
    i64 ray_intersections = 0;
   for (i64 i = 0; i < n; i++){}
        // p[i]p[i+1]
        if (on_segment(q, p[i], p[(i + 1) % n])){
            return BOUNDARY;
        ray_intersections += horizontal_ray_segment_intersect(q, p[i],
           p[(i + 1) % n]);
    // to be inside: odd exits
```

}

```
if (rav_intersections % 2 == 1){
        return INSIDE;
    } else {
        return OUTSIDE;
}
6.4 triangle.hpp
#include "./point.hpp"
// pgr oriented area
// can also be calculated by shoelace
double area(pt p, pt q, pt r){
    return (double)((q - p).cross(r - p)) / 2.0f;
bool inside_triangle(pt q, pt p1, pt p2, pt p3){
    i64 abs_area = 0;
    pt p[3] = \{p1, p2, p3\};
    for (i64 i = 0; i < 3; i++){}
        abs_area += abs((p[i] - q).cross(p[(i + 1) % 3] - q));
    i64 total = abs((p[1] - p[0]).cross(p[2] - p[0]));
    return abs_area == total;
    Graphs
7.1 2sat.cpp
#include <bits/stdc++.h>
using namespace std;
namespace SAT2 {
    vector < vector < int >> build_graph(const vector < pair < int , int >>
        &clauses, int n){
        vector < vector < int >> g(2 * n);
        for (auto[x, y]: clauses){
            // (x or y) clause = (\simx => y) clause and (\simy => x) clause
            g[(x + n) \% (2 * n)].emplace_back(y);
            g[(y + n) \% (2 * n)].emplace_back(x);
        return g;
    vector < vector < int >> transpose (vector < vector < int >> &g) {
        vector < vector < int >> gt(g.size());
        for (int u = 0; u < g.size(); u++){}
            for (int v: g[u]){
                gt[v].emplace_back(u);
            }
        return gt;
```

```
}
void forward dfs(
    const int u, vector<vector<int>> &g, vector<bool> &vis,
        vector<int> &stack
) {
    vis[u] = true:
    for (int v: g[u]){
        if (!vis[v]){
            forward_dfs(v, g, vis, stack);
    stack.emplace_back(u);
}
void backward dfs(
    const int u, vector < vector < int >> &g, vector < bool > &vis,
    vector<int> &comp, int c = 0
) {
    vis[u] = true;
    comp[u] = c;
    for (int v: g[u]){
        if (!vis[v]){
            backward_dfs(v, g, vis, comp, c);
    }
}
pair < vector < int >, int > scc(const vector < pair < int , int >> & clauses ,
   int n){
    vector < vector < int >> g = build_graph(clauses, n);
    vector < vector < int >> gt = transpose(g);
    const int s = g.size();
    vector < int > stack, comp(s);
    vector < bool > vis(s, false);
    for (int u = 0; u < s; u++){</pre>
        if (!vis[u]){
            forward_dfs(u, g, vis, stack);
    reverse(stack.begin(), stack.end());
    fill(vis.begin(), vis.end(), false);
    int c = 0:
    for (int u: stack){
        if (!vis[u]){
             backward_dfs(u, gt, vis, comp, c);
             ++c:
        }
    return {comp, c};
}
pair < vector < bool > , bool > solve(const vector < pair < int , int >> &
   clauses, int n){
    // run scc (Kosaraju)
```

```
auto[comp, num_comps] = scc(clauses, n);
        // build quotient graph
        vector < vector < int >> g_scc(num_comps);
        for (auto[x, y]: clauses){
            // (x or y) = (\simx => y) = (\simy => x)
            g_scc[comp[(x + n) \% (2 * n)]].emplace_back(comp[y]);
            g_scc[comp[(y + n) \% (2 * n)]].emplace_back(comp[x]);
        // build valuation: mark as true the first one that appears
        vector < bool > valuation(n):
        for (int u = 0: u < n: u++) {
            if (comp[u] < comp[u + n]){
                valuation[u] = false;
            } else {
                valuation[u] = true;
        }
        // check valuation
        bool ok = true:
        for (auto[x, y]: clauses){
            bool left = ((x < n) ? valuation[x] : !valuation[x - n]);</pre>
            bool right = ((y < n) ? valuation[y] : !valuation[y - n]);</pre>
            ok &= left | right:
        return {valuation, ok};
};
7.2 bellman_ford.cpp
#include <bits/stdc++.h>
using namespace std;
// edges: (a, b, c) = directed edge from a to b with weight c
vector<int> negative_cycle(const vector<tuple<int, int, int>> edges,
   int n){
   // virtual vertex connects to all others with weight 0
   // so ALL negative cycles are reachable!
    // for that effect we just need to set d[u] = 0 for all u
    // (we don't need an actual new vertex)
    vector<int> d(n, 0);
    vector<int> source_list;
    vector<int> prev(n, -1);
    // all paths of size s
    for (int s = 0; s \le n; s++){
        bool any = false;
        for (auto[a, b, c]: edges){
            if (d[a] + c < d[b]){
                any = true;
                prev[b] = a;
                d[b] = d[a] + c;
```

```
if (!any) break;
    }
    // run a (single) new iteration of relaxation
    int target = -1;
    for (auto[a, b, c]: edges){
        if (d[a] + c < d[b])
            target = b;
            break;
        }
    }
    if (target != -1){
        vector<int> cyc;
        // might be a cycle end or a vertex reachable from a negative
        int curr = target;
        vector < bool > vis(n):
        while (!vis[curr]) {
            cyc.emplace_back(curr);
            vis[curr] = true;
            // we might have prev[curr] == curr
            // (in case of self negative edges!)
            curr = prev[curr]:
        // now curr is a visited vertex (closed the cycle!)
        cyc.emplace_back(curr);
        reverse(cyc.begin(), cyc.end());
        while (cyc.back() != cyc[0]){
            cyc.pop_back(); // remove vertices reachable from the
               cycle but not in it
        return cyc;
    } else {
        return vector < int > ();
7.3 cycle_detection.cpp
#include <bits/stdc++.h>
using namespace std;
#define WHITE O
#define GRAY 1
#define BLACK 2
bool has_cycle = false;
void dfs(vector<vector<int>> &g, int u, vector<int> &color){
    if (color[u] != WHITE){
        return;
    color[u] = GRAY:
    for (int v: g[u]){
        if (color[v] == GRAY){
            has_cycle = true;
        } else if (color[v] == WHITE){
```

}

```
dfs(g, u, color);
    }
    color[u] = BLACK:
7.4 dfs_toposort.cpp
#include <bits/stdc++.h>
using namespace std;
namespace Graph {
    int c;
    void dfs(
        int u, vector<vector<int>> &G, vector<bool> &vis, vector<int>
    ) {
        vis[u] = true;
        for (int v: G[u]){
            if (!vis[v]){
                dfs(v, G, vis, order);
        order[u] = c;
        --c;
    vector<int> toposort(vector<vector<int>> &G){
        c = G.size() - 1;
        const int n = G.size();
        vector < bool > vis(n, false);
        vector < int > order(n);
        for (int u = 0; u < n; u++) {
            if (!vis[u]){
                dfs(u, G, vis, order);
            }
        // validate order
        for (int u = 0: u < n: u++) {
            for (int v: G[u]){
                if (order[u] >= order[v]) {
                    return vector<int>();
            }
        return order;
7.5 dijkstra.cpp
#include <bits/stdc++.h>
using namespace std;
```

```
const int INF = (int)2e9;
// directed edges are pairs (weight, endpoint)
void dijkstra(vector<vector<pair<int, int>>> &adj, int dis[], int s)
    for (int i = 0; i < (int)adj.size(); i++) dis[i] = INF;</pre>
    priority_queue <pair <int , int > , vector <pair <int , int >> ,
        greater < pair < int , int >> > pq;
    pq.emplace(0, s);
    dis[s] = 0:
    while (!pq.empty())
        auto [dis_u, u] = pq.top();
        pq.pop();
        if (dis u > dis[u])
            continue; // skip suboptimal
        // since dis[u] is always updated when better path is found,
        // dis[u] <= d for each [d, u] state on the priority queue
        for (auto [w, v] : adj[u])
            if (dis[u] + w < dis[v])</pre>
            {
                dis[v] = dis[u] + w;
                pq.emplace(dis[v], v);
                // sorting by dis[v] is the same as sorting by the
                    weights w[u, v]
                // since dis[v] - w(u, v) = dis[u] for all v (constant
                    difference)
            }
        }
    }
7.6 dinic.cpp
// Tested on CSES - Police Chases
#include <bits/stdc++.h>
using namespace std:
#define i64 int64 t
Dinic (overview):
// make residual graph (edges and backedges)
// then on each iteration
// do bfs to make layer graph
// do dfs to find blocking flow
const i64 INF = (i64)2e18;
struct Edge {
    i64 u, v;
    i64 c, f;
    i64 other(i64 x){
        return (x == u ? v : u);
```

```
i64 cap_from(i64 x){
        if (x == u){
            return c - f;
        } else {
            return f;
   }
    void add_flow(i64 from, i64 dF){
        if (from == u) { // consume flow
            f += dF;
        } else { // undo flow
            f -= dF:
        }
    Edge(i64 from, i64 to, i64 cap) : u(from), v(to), c(cap), f(0) {}
};
struct Dinic {
    i64 n, s, t;
    vector < vector < Edge * >> adj; // [cap, to] pairs
    void add_edge(i64 u, i64 v, i64 c){
        // must be heap-allocated, or else will disappear when leaving
           function
        Edge* e = new Edge(u, v, c);
        adi[u].emplace back(e):
        adj[v].emplace_back(e);
    Dinic(const vector < vector < pair < i64, i64 >>> g, i64 s, i64 t)
        : n(g.size()), s(s), t(t) {
        adj.resize(n);
        for (i64 u = 0; u < n; u++){}
            for (auto[c, v]: g[u]){
                add edge(u. v. c):
            }
        }
   }
    // dag adj list, level vector
    pair < vector < vector < Edge * >> , vector < i64 >> make_dag() {
        vector < vector < Edge * >> dag(n);
        vector < i64 > level(n, INF);
        queue < i64 > q;
        q.emplace(s);
        level[s] = 0;
        vector < i64 > vis(n, false);
        while (!q.empty()){
            i64 u = q.front();
            q.pop();
            if (vis[u]) continue;
            vis[u] = true;
            for (Edge* e: adj[u]){
                if (e->cap_from(u) == 0) continue; // edge does not
```

```
exist
            i64 v = e \rightarrow other(u):
            if (!vis[v]){
                level[v] = level[u] + 1:
                dag[u].emplace_back(e);
                q.emplace(v);
        }
    return {dag, level};
}
// push_flow on layer dag of residual graph
i64 push_flow(i64 u, vector<i64>& blocked, vector<i64>& level, i64
   dF = INF)
    if (u == t){
        return dF; // already at sink
    bool all neighbors blocked = true:
    i64 pushedF = 0;
    for (Edge* e: adj[u]){
        if (dF == 0) break;
        i64 v = e \rightarrow other(u):
        if (level[v] != level[u] + 1) continue; // not dag edge
        if (e->cap_from(u) == 0 || blocked[v]) continue;
        // NOTE: this push flow might block v!
        i64 delta = push_flow(v, blocked, level,
            min(e->cap_from(u), dF));
        e->add_flow(u, delta);
        dF -= delta;
        pushedF += delta;
        if (e->cap_from(u) > 0 && !blocked[v]){  // still valid
            after pushing
            all_neighbors_blocked = false;
        }
    if (all_neighbors_blocked) blocked[u] = true;
    return pushedF;
}
i64 maxflow(){
    while (true) {
        auto[dag, level] = make_dag();
        if (level[t] == INF) break; // no more augmenting paths
        // can't assume anvone is blocked (reuse block vector).
            since we might try to unblock by undoing flow
        vector < i64 > blocked(n, false); // indicates that vertex
            is blocked: can't push flow from it
        push_flow(s, blocked, level);
    i64 \text{ maxFlow} = 0;
    for (Edge* e: adj[s]){
        if (s == e \rightarrow u){ // is origin vertex: forward edge
            maxFlow += e->f;
```

```
} else { // backward edge
                 maxFlow -= e->f;
        return maxFlow;
    // assumes dinic was called first
    vector<i64> mincut(){
        auto[_, level] = make_dag();
        vector < i64 > partition_num(n);
        for (i64 u = 0: u < n: u++)
            if (level[u] < INF){</pre>
                 partition_num[u] = 0;
            } else {
                 partition_num[u] = 1;
        return partition_num;
};
7.7 floyd_warshall.cpp
#include <bits/stdc++.h>
using namespace std;
const int INF = (int)2e9; // might change depending on the problem
vector < vector < int >> floydWarshall(const vector < tuple < int, int >> &
    edges, int n){
    vector < vector < int >> d(n, vector < int > (n, INF));
    for (int i = 0; i < n; i++){
        d[i][i] = 0:
    for (auto[a, b, w]: edges){
        d[a][b] = min(d[a][b], w):
    for (int k = 0; k < n; ++k) {
        for (int i = 0; i < n; ++i) {
            for (int j = 0; j < n; ++ j) {
                 if (d[i][k] < INF && d[k][j] < INF)</pre>
                     d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
        }
    }
    return d;
     kosaraju.cpp
#include <bits/stdc++.h>
using namespace std;
vector<vector<int>> transpose(const vector<vector<int>> &g){
```

```
vector < vector < int >> gt(g.size());
    for (int u = 0; u < g.size(); u++){</pre>
        for (int v: g[u]){
            gt[v].emplace_back(u);
    return gt;
void forward_dfs(
    const int u, const vector<vector<int>> &g, vector<bool> &vis,
        vector<int> &stack
) {
    vis[u] = true;
    for (int v: g[u]){
        if (!vis[v]){
            forward_dfs(v, g, vis, stack);
    stack.emplace_back(u);
void backward_dfs(
    const int u, const vector < vector < int >> &g, vector < bool > &vis,
    vector<int> &comp, int c = 0
    vis[u] = true;
    comp[u] = c;
    for (int v: g[u]){
        if (!vis[v]){
            backward_dfs(v, g, vis, comp, c);
pair < vector < int > , int > scc(const vector < vector < int >> &g , int n) {
    vector < vector < int >> gt = transpose(g);
    const int s = g.size();
    vector < int > stack, comp(s);
    vector < bool > vis(s, false);
    for (int u = 0; u < s; u++){
        if (!vis[u]){
            forward_dfs(u, g, vis, stack);
    reverse(stack.begin(), stack.end());
    fill(vis.begin(), vis.end(), false);
    int c = 0:
    for (int u: stack){
        if (!vis[u]){
            backward_dfs(u, gt, vis, comp, c);
            ++c;
        }
    return {comp, c};
```

7.9 lexicographical_minimal_toposort.cpp

```
#include <bits/stdc++.h>
using namespace std;
#define i64 int64_t
// toposort with bfs, returns the lexicographically smallest
   topological order
// Runs in O(n log n) because of sorting
// Returns boolean indicating if topo order exists, and the order in
pair < bool, vector < i64 >> lexicographical_toposort(const
   vector < vector < i64 >> & dag) {
    const i64 n = dag.size():
    // dependencies_met[i]: how many j such that j->i
    // have received their order number
    vector < i64 > dependencies_met(n, 0),
                in_degree(n, 0),
                order(n, -1);
    for (i64 u = 0: u < n: u++){}
        for (i64 v: dag[u]){
            ++in_degree[v];
    priority_queue < i64, vector < i64>, greater < i64>> q;
    // min pq guarentees lexicographical order and topo order
    // by only pushing when all dependencies are met
    for (i64 u = 0; u < n; u++){}
        if (!in_degree[u]) q.emplace(u);
    i64 c = 0;
    while (!q.empty()){
        i64 u = q.top();
        order[u] = c++:
        q.pop();
        for (i64 v: dag[u]){
            ++dependencies_met[v];
            if (dependencies_met[v] == in_degree[v]){
                q.emplace(v);
        }
    for (i64 u = 0; u < n; u++){}
        // circular dependencies exist
        if (order[u] == -1) return {false, {}};
    return {true, order};
```

7.10 tarjan.cpp

#include <bits/stdc++.h>

```
using namespace std;
const int INF = (int)2e9;
namespace Tarjan {
    vector < int > lowlink, t_in;
    vector < pair < int , int >> dfs_tree_edges;
    int timer:
    void dfs(int u, const vector<vector<int>> &g, vector<bool> &vis,
            int parent = -1){
        vis[u] = true;
        t in[u] = lowlink[u] = timer++:
        for (int to: g[u]){
            if (to == parent) continue;
            else if (vis[to]){
                // back edge on DFS tree
                lowlink[u] = min(t_in[to], lowlink[u]);
            } else {
                dfs(to, g, vis, u);
                dfs_tree_edges.emplace_back(u, to);
                lowlink[u] = min(lowlink[to], lowlink[u]);
            }
        }
    }
    // lowlink(u) = min(t_in(u), min_{p>u}t_in(u), min_{u>to} em
       tree}lowlink(to))
    vector < vector < pair < int , int >>> dfs_tree_bridges (const
       vector<vector<int>> &g) {
        const int n = g.size();
        vector < bool > vis(n, false);
        lowlink = vector < int > (n, INF);
        t_in = vector < int > (n, INF);
        dfs_tree_edges.clear();
        timer = 0;
        for (int u = 0; u < n; u++) {
            if (!vis[u]){
                dfs(u, g, vis);
        }
        vector < vector < pair < int , int >>> h(n);
        for (auto[v, to]: dfs_tree_edges){
            if (lowlink[to] > t_in[v]){
                h[v].emplace_back(to, 1);
            } else {
                h[v].emplace_back(to, 0);
        return h:
```

}

Number Theory

#include <bits/stdc++.h>

8.1 fastexp.cpp

```
using namespace std;
const int MOD = (int)1e9 + 7; // NOTE: MIGHT CHANGE
int bexp(int a, int p){
   if (p == 0){
        return 1:
   } else {
        int m = bexp(a, p / 2);
        if (p % 2 == 1){
            return (((m * m) % MOD) * a) % MOD;
        } else {
            return (m * m) % MOD;
   }
int inv(int a){
    return bexp(a, MOD - 2);
8.2 spf.cpp
Given N, calculates the smallest prime factor (spf) that divides K for
   each 1 \le K \le N
If K is prime, that prime factor will be K itself
#include <bits/stdc++.h>
using namespace std;
void spf(int N, int ans[]){
// ans is a N+1 sized array
 for (int i = 0; i <= N; i++){</pre>
    ans[i] = i:
  int is_prime[N+1]; // O means is prime, 1 means is composite
  memset(is_prime, N + 1, 0);
  is_prime[0] = 1; is_prime[1] = 1;
  for (int i = 2: i <= N: i++){
   if (is_prime[i] == 0){
     for (int mul = i; mul <= N; mul += i){</pre>
        ans[mul] = min(ans[mul], i);
        // last value will be the smallest divisor of mul greater than
           1 (which is the spf)
     }
   }
```

```
}
}
```

9 Range Queries

a.assign(n, val);

9.1 bit.cpp

```
#include <bits/stdc++.h>
using namespace std;
vector < int > BIT (const vector < int > &a) {
    int n = a.size():
    vector < int > B(n + 1, 0):
    for (int i = 0; i < n; i++){</pre>
        update(B, i, a[i]); // sum a[i] to position i+1 of the BIT
    return B;
}
int get(const vector<int> &BIT, int i){
    int s = 0;
    while (i >= 0){
        s += BIT[i];
        // i & -i == 2^lsb(i)
        i = i ^ (i & (-i));
    return s;
}
void update(vector<int> &BIT, int i, int x){
    const int n = BIT.size() - 1;
    while (i \le n) {
        BIT[i] += x:
        i += i & (-i);
    }
9.2 segtree.cpp
#include <bits/stdc++.h>
using namespace std;
const int INF = (int)2e9;
// default seg tree (in this example, a Max Seg)
struct MaxSegTree {
    vector < int > seg;
    vector < int > a;
    int n:
    MaxSegTree(int n, int val = -INF) : n(n) {
        seg.assign(4 * n, val);
```

```
}
    int query(int 1, int r){
        return _query(1, r, 0, n - 1, 0);
    void update(int p, int x){
        _{update(p, x, 0, n - 1, 0)};
private:
    int _query(int 1, int r, int t1, int tr, int node){
        if (1 == t1 && r == tr){
            return seg[node];
        } else {
            int tm = t1 + (tr - t1) / 2:
            if (r <= tm){
                return _query(1, r, tl, tm, 2 * node + 1);
            } else if (1 > tm){
                return _query(1, r, tm + 1, tr, 2 * node + 2);
            } else {
                return max (
                    _{query}(1, tm, tl, tm, 2 * node + 1),
                    _{query(tm + 1, r, tm + 1, tr, 2 * node + 2)}
                );
            }
        }
    }
    void _update(int p, int x, int tl, int tr, int node){
        if (tl == p && tr == p){
            seg[node] = x;
            a[p] = x;
        } else {
            int tm = t1 + (tr - t1) / 2;
            if (p <= tm){
                _{update(p, x, tl, tm, 2 * node + 1)};
            } else {
                _update(p, x, tm + 1, tr, 2 * node + 2);
            seg[node] = max(seg[2 * node + 1], seg[2 * node + 2]);
        }
   }
};
```

10 Strategy

10.1 Início (primeira 1h30)

- Ler questões
- Passar no máximo 10-15 min por questão
- Implementar apenas as fáceis
 - Priorizar primeiro aquelas que são simples e rápidas de implementar (em detrimento por ex. de uma aplicação trivial de um algoritmo longo, como uma SegTree ou fluxo)
- Fase mais individual
- Priorizar para implementação aquela pessoa que tem a ideia mais clara da solução

10.2 Meio

- Ler o restante da prova
- Olhar os standings para selecionar a próxima questão
- Analisar mais profundamente no papel para conseguir passar as médias/difíceis
- Evitar ficar três pessoas em uma mesma questão

10.3 Final (última 1h30)

• Foco do time todo em uma ou duas questões mais prováveis de passarem

11 Strings

11.1 rabin_karp.cpp

```
#include <bits/stdc++.h>
using namespace std;
#define i64 int64 t
vector<i64> rollingHash(const vector<int>& v, i64 p, i64 m){
    const i64 n = v.size();
    vector < i64 > prefH(n);
    vector < i64 > p_pow(n);
    // prefH = hash(v[0...i]) = sum(j<=i) v[j] * p^j mod m
    p_pow[0] = 1;
    prefH[0] = v[0] \% m;
    for (i64 i = 1; i < n; i++){
        p_pow[i] = (p_pow[i - 1] * p) % m;
        prefH[i] = (v[i] * p_pow[i] + prefH[i - 1]) % m;
    return prefH:
// lowercase
vector<i64> rollingHash(const string& s){
    vector<int> s_repr(s.size());
    for (i64 i = 0: i < s.size(): i++){}
        s_repr[i] = s[i] - 'a';
    return rollingHash(s_repr, 31, (i64)1e9 + 7);
i64 bexp(const i64 a, const i64 p, const i64 mod){
    if (p == 0){
        return 1;
    } else {
        const i64 b = bexp(a, p / 2, mod);
        if (p \% 2 == 0){
            return (b * b) % mod;
            return (a * ((b * b) % mod)) % mod;
vector < i64 > prime_powers(const i64 n, const i64 p = 31, const i64 m =
   (i64)1e9 + 7){}
    vector < i64 > p_pow(2 * n + 1);
    p_pow[n] = 1;
    p_pow[n-1] = bexp(p, m-2, m); // fermat little theorem
    for (i64 i = 1: i <= n: i++){
        p_pow[n + i] = (p_pow[n + i - 1] * p) % m;
   for (i64 i = 2; i <= n; i++){</pre>
        p_pow[n - i] = (p_pow[n - i + 1] * p_pow[n - 1]) % m;
```

```
return p_pow;
vector<i64> rabinKarpMatch(const string& s, const string& t){
    vector<i64> prefix_s = rollingHash(s), prefix_t = rollingHash(t);
    const i64 mod = (i64)1e9 + 7, p = 31;
    const i64 HS = prefix_s.back(); // complete hash of s
    // precompute prime powers mod M
    const vector<i64> p_pow = prime_powers(t.size(), p, mod);
    vector < i64 > matches; // starting points of matches
    const i64 offset = t.size():
    for (i64 i = 0; i + s.size() <= t.size(); i++){</pre>
        // hash of t[i...i + |s| - 1] == (hash(t[0...i + |s| - 1]) -
           hash(t[0...i - 1])) * p^-i
        const i64 h = ((prefix_t[i + s.size() - 1] - (i > 0)?
           prefix_t[i - 1] : 0) + mod) * p_pow[offset - i]) % mod;
        if (h != HS) {
            continue;
        }
        // compare
        bool ok = true;
        for (i64 j = 0; (j < s.size()) && ok; j++){
            ok = ok && (s[i] == t[i + i]);
        if (ok) matches.emplace_back(i);
    return matches;
11.2 suffix_array.cpp
// TESTED on first suffix array problem of Codeforces EDU
#include <bits/stdc++.h>
using namespace std;
#define ALPHABET 256
void csort(vector<int>& p, const vector<int>& classes, int offset){
    const int n = p.size();
    vector < int > cnt(max(ALPHABET, n), 0), ptemp(n);
    for (int i = 0; i < n; i++){</pre>
        ++cnt[classes[i]];
    }
    for (int i = 1; i < max(ALPHABET, n); i++){
        cnt[i] += cnt[i - 1]:
    }
    // distribute on blocks in reverse order to make sorting stable
    // distribute according to order of k-th element (k=offset)
    for (int i = n - 1; i \ge 0; i--){
        int offseted_pos = (p[i] + offset) % n;
        ptemp[--cnt[classes[offseted_pos]]] = p[i];
    }
```

```
swap(ptemp, p);
pair < vector < int >, vector < vector < int >>> suffix_array(const string& s) {
    int lg = 0;
    string t = s;
    t.push_back('$');
    int n = t.size();
    while ((111 << lg) < n) ++ lg; // now 2^lg >= n
    vector < int > p(n);
    vector<vector<int>> c(lg + 1, vector<int>(n));
    for (int i = 0: i < n: i++){
        p[i] = i:
        c[0][i] = (int)t[i];
    csort(p, c[0], 0);
    // check offset of size 2^k
    for (int k = 0; k \le lg; k++) {
        csort(p, c[k], (111 << k)); // sort by second, stably</pre>
        csort(p, c[k], 0); // sort by first, stably
        if (k == lg) break;
        // redo equivalence classes
        c[k + 1][p[0]] = 0: // first element receives class=0
        int cls = 0;
        for (int i = 1; i < n; i++) {
            if (
                (c[k][p[i-1]] < c[k][p[i]])
                || (c[k][(p[i - 1] + (111 << k)) \% n] < c[k][(p[i] +
                    (111 << k)) % n])
            ) {
                ++cls:
            c[k + 1][p[i]] = cls;
    }
    return {p, c};
int lcp(int i, int j, const vector<vector<int>>& cls){
    const int n = cls[0].size():
    int lg = 0;
    while ((111 << lg) < n) ++ lg; // now 2^lg >= n
    int ans = 0;
    for (int k = lg; k \ge 0; k--){
        if (cls[k][i] == cls[k][j]){
            i = (i + (111 << k)) \% n;
            j = (j + (111 << k)) \% n;
            ans += (111 << k);
        }
    }
    return ans;
// LCS in O(N log N)
// TESTED on second suffix array problem of Step 2 of Codeforces EDU
tuple < int, int, int > lcs(string s, string t){
```

```
auto [p, c] = suffix_array(s + "#" + t);
    // NOTE: the sentinel must be different to avoid having cyclic
       shifts breaking stuff
    const int n = s.size(), m = t.size();
    const int sz = n + 1 + m + 1; // middle # and final $
    vector < int > lcp_array(sz - 1);
    for (int i = 0; i + 1 < sz; i++) lcp_array[i] = lcp(p[i], p[i +
       1]. c):
    int lcs_sz = 0, s_idx = -1, t_idx = -1;
    for (int i = 0; i + 1 < sz; i++){
        int left = p[i], right = p[i + 1];
        if (lcp_array[i] > lcs_sz && min(left, right) < n && max(left,</pre>
            right) > n){
            // p[i] in s, p[i + 1] in t or vice-versa
            lcs_sz = lcp_array[i];
            s_idx = min(left, right);
            t_idx = max(left, right);
        }
    }
    return {lcs_sz, s_idx, t_idx};
11.3 trie.cpp
#include <bits/stdc++.h>
using namespace std;
template <int K> // sum of characters in the set
struct Trie {
    int children[K][26];
    int cnt[K]; // cnt of each word
    int ptr = 1;
    Trie(){
        for (int i = 0; i < K; i++){
            for (int j = 0; j < 26; j++){
                children[i][j] = -1;
            cnt[i] = 0;
        }
    }
    void insert(string& u, int i=0, int node=0){
        int idx = u[i] - 'a';
        if (children[node][idx] == -1){
            // create node
            children[node][idx] = ptr++;
        if (i + 1 < u.size()){</pre>
            insert(u, i + 1, children[node][idx]);
        } else {
            ++cnt[children[node][idx]];
    }
```

```
int search(string &u, int i=0, int node=0){
        if (i == u.size()){
            return cnt[node]:
        } else {
            int idx = u[i] - 'a';
            if (children[node][idx] == -1){
                return 0:
            } else {
                return search(u, i + 1, children[node][idx]);
        }
    bool isleaf(int node){
        bool ok = true;
        for (int i = 0; i < 26; i++) {
            ok = ok && children[node][i] == -1;
        return ok;
    }
    void remove(string &u, int i=0, int node=0){
        int idx = u[i] - 'a';
        if (children[node][idx] != -1){
            remove(u, i + 1, children[node][idx]);
            if (isleaf(children[node][idx]) &&
                cnt[children[node][idx]] == 0){
                children[node][idx] = -1;
       } else {
            // reached leaf
            cnt[node] = max(cnt[node] - 1, 0);
        }
};
     template.cpp
#define TESTCASES
#define debug cerr
#include <bits/stdc++.h>
using namespace std;
#define fastio
   ios_base::sync_with_stdio(false); cin.tie(NULL); cout.tie(NULL)
#define endl '\n'
#define i64 int64 t
#define u64 uint64 t
#define i128 __int128
#define all(x) begin(x),end(x)
#define printmsg(msg, first, last) debug << msg; for(auto it = first;</pre>
   it != last; it++) {debug << *it << " ";} debug << endl;
```

#define printgraph(msg, G) debug << msg << endl; \</pre>

for (u64 u = 0; u < G.size(); u++) { \</pre>

```
debug << "G[" << u << "]="; \
        for (u64 v: G[u]) { \
            debug << v << " ": \
        } \
        debug << endl; \
    }
template <typename T, typename U>
ostream& operator << (ostream& out, pair <T, U> x)
    out << x.first << " " << x.second;
    return out;
template <typename T>
ostream& operator << (ostream& out, vector <T> v)
    for (const auto& x: v){
        out << x << " ":
    return out;
template <typename T>
using min_pq = priority_queue<T, vector<T>, greater<T>>;
template <typename T>
using max_pq = priority_queue <T>;
void solve(){
signed main(){
    fastio;
    int t:
#ifdef TESTCASES
    cin >> t;
#else
    t = 1:
#endif
    while (t--)
        solve():
     Trees
13.1 centroid.cpp
#include <bits/stdc++.h>
```

```
#include <bits/stdc++.h>
using namespace std;

namespace CentroidDecomposition {
   int subtree_size(int u, const vector<vector<int>& adj,
      vector<int>& s, vector<int>& blocked, int p = -1){
      s[u] = 1;
      for (int v: adj[u]){
        if (v == p) continue;
   }
}
```

```
if (blocked[v]) continue;
        s[u] += subtree_size(v, adj, s, blocked, u);
    return s[u]:
// find centroid of subtree
int find centroid(int u. const vector < vector < int >> & adi.
    vector<int>& s, vector<int>& blocked, int total_size, int p =
    -1){
    for (int v: adj[u]){
        if (v == p) continue;
        if (blocked[v]) continue;
        if (2 * s[v] >= total_size){
            return find_centroid(v, adj, s, blocked, total_size,
                u); // there can be only one such vertex
    }
    return u;
// total_size = size of subtree rooted at root (considering
    deletions)
// @return centroid subtree root
int _build_centroid_tree(
    const vector < vector < int >> & adj,
    vector < int > & blocked,
    vector < vector < int >> & centroid_adj,
    vector < int > & sz,
    int total_size,
    int root = 0
) {
    subtree_size(root, adj, sz, blocked);
    int c = find_centroid(root, adj, sz, blocked, total_size);
    blocked[c] = true;
    for (int v: adj[c]){
        if (blocked[v]) continue;
        // build centroid subtree on partition and get its root
        int c2 = _build_centroid_tree(adj, blocked, centroid_adj,
            sz, sz[v], v);
        centroid_adj[c].emplace_back(c2);
    return c;
}
// @return pair (centroid tree root, directed centroid tree)
pair < int , vector < vector < int >>> centroid_tree (const
    vector < vector < int >> & adj) {
    const int n = adj.size();
    vector < int > blocked(n, false), sz(n);
    vector < vector < int >> centroid_adj(n);
    int master centroid = build centroid tree(
        adj, blocked, centroid_adj, sz, n
    return {master_centroid, centroid_adj};
```

```
}
```

13.2 hld.cpp

```
// HLD implementation
// Tested on CSES Path Queries (https://cses.fi/problemset/task/1138)
// Query: sum of vertices from root to v
// Update: set value of v to x
// Implementation note: there is an implementation where
// you use a single seg-tree and the node positions are their
// time of visit in the DFS, where you **visit heavy edges first**
// (so heavy paths make contiguous ranges)
// However, this leads to a larger segtree and can be slower!
#include <bits/stdc++.h>
using namespace std;
#define i64 int64 t
// uses some modifications: now the heavy edge is the edge with the
   largest subtree
// instead of the restriction s(c) >= s(v) / 2
// Still, going through light (i.e not heavy) edge reduces size by half
// Heavy path goes up to just the vertex before light edge
using tree = vector <vector <i64>>;
struct HLD {
    vector < i64 > heavy; // heavy[u] = v <==> uv is heavy edge
    vector < i64 > sz;
    vector < i64 > head; // head[u] = start of heavy path that goes
       through u
    vector < i64 > pos; // depth of vertex in path. Head has depth 0
    vector < i64 > parent;
    // NOTE: SegTree is the segment tree implementation
    // NOTE: SegTree needs default constructor
    vector < SegTree > segs;
    vector < i64 > a;
    vector < i64 > depth;
    HLD(const tree& adj, const vector<i64>& _a, i64 root = 0) : a(a) {
        const i64 n = adj.size();
        heavv.resize(n):
        sz.resize(n):
        head.resize(n);
        pos.resize(n);
        parent.resize(n);
        segs.resize(n);
        subtree(root, adj);
        dfs(root, adj);
        // build segment trees
        for (i64 h = 0; h < n; h++){}
            if (h == head[h]){
                vector < i64 > path;
                i64 cur = h;
                while (cur != -1 \&\& head[cur] == h){
                    path.emplace_back(a[cur]); // add value
```

```
cur = heavv[cur];
            segs[h] = SegTree(path);
        }
   }
void subtree (i64 u, const tree & adj, i64 p = -1, i64 d = 0) {
    sz[u] = 1:
    depth[u] = d;
    for (i64 v: adj[u]){
        if (v == p) continue;
        subtree(v, adj, u, d + 1);
        sz[u] += sz[v]:
}
void dfs(i64 u, const tree& adj, i64 p = -1, i64 h = 0, i64 d = 0){
    i64 cur_sz = 0;
    pos[u] = d;
    head[u] = h;
    parent[u] = p;
    heavy[u] = -1;
    for (i64 v: adj[u]){
        if (v == p) continue;
        if (sz[v] > cur_sz){
            cur sz = sz[v]:
            heavy[u] = v;
        }
    }
    for (i64 v: adj[u]){
        if (v == p) continue;
        if (v == heavy[u]){
            dfs(v, adj, u, h, d + 1);
        } else {
            dfs(v, adj, u, v, 0);
    }
}
// query from u to p^min(k, depth[u])(u), inclusive
// NOTE: SUM QUERY HERE, other variations are possible
i64 query_up(i64 u, i64 k){
    i64 s = 0;
    while (u != -1 \&\& k >= pos[u]) {
        s += segs[head[u]].query(0, pos[u]);
        u = parent[head[u]];
    if (u != -1) \{ // k < pos[u] \}
        s += segs[head[u]].query(pos[u] - k, pos[u]);
    return s;
}
i64 query_path(i64 u, i64 v){
```

```
// TODO implement LCA!
        i64 l = lca(u, v);
        return query_up(u, depth[u] - depth[l]) + query_up(v, depth[v]
            - depth[1]) - a[1]:
    }
    void update(i64 u, i64 x){ // set value to x
        segs[head[u]].update(pos[u], x);
};
13.3 lca.cpp
#include <bits/stdc++.h>
using namespace std;
const int N = 2e5, LOGN = 20;
int lift[LOGN + 1][N];
void binary_lifting(const vector<vector<int>> &forest){
    const int n = forest.size();
    vector < int > parent(n, -1);
    for (int u = 0; u < n; u++) {
        for (int v: forest[u]){
            parent[v] = u;
        }
    }
    for (int u = 0; u < n; u++){
        lift[0][u] = parent[u];
    for (int k = 1: k \le LOGN: k++) {
        for (int u = 0: u < n: u++) {
            lift[k][u] = lift[k - 1][lift[k - 1][u]];
    }
}
int jump(int u, int k){
    for (int i = 0; i <= LOGN; i++){</pre>
        if ((k >> i) & 1){
            u = lift[i][u];
    return u;
int LCA(
    int u, int v, const vector <int> &level,
    const vector <int > &parent
    // invariant: level[u] <= level[v]</pre>
    if (level[v] < level[u]) swap(u, v);</pre>
    v = jump(v, level[v] - level[u]);
    if (u == v){
```

return u;

```
}
    // loop invariant: u and v are distinct
    // (we have not reached a common ancestor yet)
    for (int i = LOGN; i >= 0; i--){
        if (lift[i][u] != lift[i][v]){
            // keeps loop invariant -> greedily take it!
            u = lift[i][u];
            v = lift[i][v]:
    }
    // at the end of the invariant we must be just one level below the
    // otherwise we could have jumped more
    return parent[u];
13.4 path_queries.cpp
// Tested on CSES Distance Queries
// sum queries on paths. Values are on vertices
#include <bits/stdc++.h>
using namespace std;
const int N = 2e5, LOGN = 20;
int lift[LOGN + 1][N];
int sums[LOGN + 1][N];
void binary_lifting(const vector<int>& parent, const vector<int>&
   values){
    const int n = parent.size();
    for (int u = 0; u < n; u++) {
        lift[0][u] = parent[u];
        sums[0][u] = values[u];
    for (int k = 1; k \le LOGN; k++) {
        for (int u = 0; u < n; u++) {
            lift[k][u] = lift[k - 1][lift[k - 1][u]];
            // u..p^(2^{k - 1})(u) and then p^(2^{k - 1}+1)(u) to
                p^{(2k(u))}
            sums[k][u] = sums[k - 1][u] + sums[k - 1][lift[k - 1][u]]:
    }
int jump(int u, int k){
    for (int i = 0; i <= LOGN; i++){</pre>
        if ((k >> i) & 1){
            u = lift[i][u]:
    return u;
```

```
// v[u] + v[p[u]] + ... + v[p^{2^k - 1}(u)]
int query_sum(int u, int k){
    int s = 0;
    for (int i = 0; i <= LOGN; i++){</pre>
        if ((k >> i) & 1){
            s += sums[i][u];
            u = lift[i][u];
        }
    return s;
    int u, int v, const vector <int> &level,
    const vector <int > &parent
) {
    // invariant: level[u] <= level[v]</pre>
    if (level[v] < level[u]) swap(u, v);</pre>
    v = jump(v, level[v] - level[u]);
    if (u == v){
        return u;
    // loop invariant: u and v are distinct
    // (we have not reached a common ancestor yet)
    for (int i = LOGN; i >= 0; i--){
        if (lift[i][u] != lift[i][v]){
            // keeps loop invariant -> greedily take it!
            u = lift[i][u];
            v = lift[i][v];
        }
    }
    // at the end of the invariant we must be just one level below the
       LCA
    // otherwise we could have jumped more
    return parent[u];
int path_sum(int u, int v, const vector<int> &level, const vector<int>
   &parent){
    int lca = LCA(u, v, level, parent);
    int left = query_sum(u, level[u] - level[lca]), right =
       query_sum(v, level[v] - level[lca]);
    return left + right + sums[0][lca];
}
```