

Question No. 1

9(b)

5

One mark for each correct calculation as follows (Max 4)

- Node B (from Home) (Line 3 in table)
- Node C (from Home) (Line 4 in table)
- Node B and Node E (from A) (Lines 5 and 6 in table)
- Node F and Node School (from E) (Lines 7 and 8 in table)
- Node School (from F) (Line 9 in table)

One mark for correct path (Max 1):

- Home \Rightarrow A \Rightarrow E \Rightarrow F \Rightarrow School

	Node	Cost from Home Node (g)	Heuristic (h)	Total (f = g + h)
1	Home	0	14	14
2	A	1	10	11
3	B	5	7	12
4	C	4	9	13
5	B	1 + 3 = 4	7	11
6	E	1 + 6 = 7	3	10
7	F	7 + 1 = 8	3	11
8	School	7 + 5 = 12	0	12
9	School	8 + 3 = 11	0	11

Final Path	Home \Rightarrow A \Rightarrow E \Rightarrow F \Rightarrow School
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Question No. 2

5(a)

Working (**Max 3**)

5

May be seen on diagram

- Initialisation: setting Base to 0
- ... and the rest of the towns to ∞
- Evidence to show values at nodes being updated
- Evidence to show 'visited node(s)'

May be seen in working section of paper

- Evidence to show calculation of at least one route
- Evidence to show more than one route has been calculated for at least one town

Correct Answer (Max 2)

One mark for four correct values...

... One mark for all values correct

Town 1	Town 2	Town 3	Town 4	Town 5	Town 6
3	5	2	9	3	8

Question No. 3

8(c)	MP1	Linear search $O(n)$ and Binary search $O(\log_2 n) / O(\log n)$	3
	MP2	time to search increases linearly in relation to the number of items in the list for a linear search and logarithmically for a Binary search	
	MP3	time to search increases less rapidly for a binary search and time to search increases more rapidly for a linear search	

Graphs, Dictionary and Big O Notation

Question No. 4

(a) FUNCTION Hash(**Key** : STRING) RETURNS INTEGER
 DECLARE Number : INTEGER
 Number ← ASCII(LEFTSTRING(**Key**,1))
 // Number ← ASCII(**Key**[1])
 Number ← Number - 64
 RETURN Number
 // Result ← Number // Hash ← Number
 ENDFUNCTION

Accept ASC instead of ASCII

Accept LEFT instead of LEFTSTRING

Key can be a different identifier but must be the same in both places

[5]

(b) (i)

Index	Dictionary Key	Value
1		
2		
3	Computer	Rechner
4	Disk	Platte
5	Error	Fehler
6	File	Datei
7		
8		
:	:	:
:	:	:
1999		
2000		

Ignore spelling mistakes

1 mark for 2 correct pairs entered in correct slots

[2]

(ii) Collision / synonym / space already occupied / same index in array
 Overwrites previous key-value pair

reject error

[Max 2]

(iii) Create an overflow area

The 'home' record has a pointer to others with the same key // linked list
 OR

Store the overflow record at the next available address ...
 in sequence (= next available)

OR

Re-design the hash function // write a different/another algorithm
 to generate a wider range of indexes // enlarging storage space // to create fewer
 collisions

[2]

(iv) Mark as follows:

Check whether slot is empty:

```
IF Dictionary[Index,1] <>"" // != '' // > NULL // >
NONE
```

If not: update index: THEN Index \leftarrow <some value>

...to find an empty slot (loop / follow pointer / go to overflow area) reject FOR loop

Insert code between lines 20 and 30

```
21 WHILE Dictionary[Index,1] > ""
22   Index  $\leftarrow$  Index + 1
23   IF Index > 2000
24     THEN
25       Index  $\leftarrow$  1
26   ENDIF
27 ENDWHILE
```

[4]