a)	X=A.(\(\overline{A} + \overline{B}\)) A. \(\overline{A} \overline{B} + \overline{A}	[1] [1] [1]	3
b)	$X = A \cdot \overline{AB}$ $X = A \cdot \overline{A} + A \cdot \overline{B}$ = $0 + A \cdot \overline{B}$ = $A \cdot \overline{B}$	[1] [1] [1]	3
c)	logic circuit has: 1 AND gate and 1 NOT gate inputs to one NOT gate is B inputs to AND gate are A and output fro	[1] m NOT gate [1]	3
d)	A B X Y 0 0 0 0 0 0 1 0 1 1 0 1 1 1 1 0	[1] [1] [1] [1]	4
e)	half adder	[1]	1

	Inj	out	Out	put								
	X	Y	Α	В								
a) i)	0	0	0	0	1 mark for each correct column	2						
α, ι,	0	1	0	1	(A and B)							
	1	0	0	1								
	1	1	1	0								
a) ii)	Half ad	dder				1						
a) iii)	C // Carry S // Sum represents the carry part of the addition of two bits represents the sum part of the addition of two bits											
b) i)	A. (A	B + C))			2						
b) ii)	Allow follow through from (b)(i) A.(A.B+C) = A.A.B + A.C = A.B + A.C = A.(B+C) 1 mark for each correct simplification line – max 2 1 mark for A.(B+C) if correct answer to part (b)(i)											

	Input	t	Warking an an	Out	tput							
	PQ	R	Working space	J	K							
	0 0	0		0	0							
	0 0	1		0	1	1 mark each column						
a) i)	0 1	0		0	1		2					
,	0 1	1		1	0	If zero marks then 6 or 7 pairs correct - 1 mark						
	1 0	0		0	1	o or 7 pairs correct — I mark						
	1 0	1		1	0]						
	1 1	0		1	0							
	1 1	1		1	1							
a) ii)	Full adder											
a) iii)	C / Carry S / Sum represents the carry part of the addition of three bits represents the sum part of the addition of three bits											
b) i)	A. (A+	·B).	С				2					
b) ii)	(A+B).C Allow follow through from (b)(i) A. ((A+B).C) = A.(A.C + B.C) = A.A.C + A.B.C = A.C + A.B.C = A.C (1 + B) = A.C.1 = A.C 1 mark for each correct simplification line – max 3 [3] 1 mark for A.C if correct answer to part (b)(i) [1]											

a)	NOR								1		
h) i)	1 mark for X column, 1 mark for Y column										
		A E	3	Workin	g Space	х	Y				
	() ()			0	0		2		
b) i)	() 1				0	1		2		
		ı)			0	1				
		1 1				1	0				
b) ii)	Half a	dder							1		
b) iii)	□ X		oullet ed for) <u>ca</u> ed for) <u>su</u>						2		

	1 mark per bullet for working (max 4)							
	\$\overline{A}.\overline{B}.\overline{C}.\overline{D}\$ \square \overline{A}.\overline{B}.\overline{C}.\overline{D}\$ \overline{A}.\overline{B}.\overline{C}.\overline{D}\$ \overline{A}.\overline{B}.\overline{C}.\overline{D}\$ \overline{A}.\overline{B}.\overline{C}.\overline{D}\$ \overline{A}.\overline{B}.\overl							
	$= \overline{A}.\overline{B}.\overline{C}.\overline{D} \square \overline{C}.D \square C.D \square C.\overline{D} [+ \overline{A}.\overline{C}.\overline{D}.[B \square \overline{B}]$ $\square \text{Taking } \overline{A}.\overline{B} \text{ and } \overline{A}.\overline{C}.\overline{D} \text{ outside brackets (Associative Law)}$							
	$= \overline{A}.\overline{B}. \boxed{\overline{C}}. \boxed{\overline{D}} \square D[\square C. \boxed{D} \square \overline{\overline{D}}[] + \overline{A}.\overline{C}.\overline{D}. \boxed{B} \square \overline{\overline{B}}[$ $\square \text{Grouping } \overline{C}. \boxed{\overline{D}} \square D[\square C. \boxed{D} \square \overline{\overline{D}}[\text{ (Associative Law and Commutative}]}$							
c)	Law)	5						
0 ,	$= \overline{A}.\overline{B}.\overline{C}.\overline{C}.\overline{C} = C.\overline{C} = C.\overline{C}.\overline{C}.\overline{C}.\overline{C}.\overline{C}.\overline{C}.\overline{C}.C$							
	$= \overline{A}.\overline{B}.\overline{C} \square C[+ \overline{A}.\overline{C}.\overline{D}.\underline{C}]$							
	$=\overline{A}.\overline{B}.$ $\boxed{1}$							
	□ Replacing [D□D̄[with 1 and replacing [C̄□ C[with 1 (Use of Complement Law)							
	= Ā.B □ Ā.C.D □ Reducing first four terms to Ā.B and reducing last two terms to Ā.C.D (Use of Identity Law)							
	1 mark for correct answer = $\overline{A}.(\overline{B} \Box \overline{C}.\overline{D})$							

a)	$X = ((P \ XOR \ Q) \ XOR \ R)$ $Y = ((P \ XOR \ Q) \ AND \ R) \ OR \ (P \ AND \ Q)$ or $X = (\overline{P}.Q + P.\overline{Q}).R + (\overline{P}.Q + P.\overline{Q}).\overline{R}$ $Y = (\overline{P}.Q + P.\overline{Q}).R + P.Q$ One mark for correct use of XOR One mark for correct use of AND One mark for correct use of OR One mark for X correct One mark for Y correct	5
b) i)	X: Sum Y: Carry (out)	2
b) ii)	Carry (in)	1

	One mar					Q and R)				
	Α	В	С	Р	Q	R	Υ	Z			
	0	0	0	0	0	0	0	0			
	0	0	1	0	0	0	1	0			
a)	0	1	0	1	0	0	1	0		3	
	0	1	1	1	0	1	0	1			
	1	0	0	1	0	0	1	0			
	1	0	1	1	0	1	0	1			
	1	1	0	0	1	0	0	1			
	1	1	1	0	1	0	1	1			
b)	Full add	er							_	1	
	One mar	k for eac	h point								
c)	$Y = \overline{A} \overline{B} C + \overline{A} B \overline{C} + A \overline{B} \overline{C} + A B C$ Purpose: Sum bit										
	$Z = \overline{A} E$ Purpose:			A B C	+ A B (C					

a)	One mark per bullet point to max 3 Circuit / electronic components (construction) With two states Used for data storage elements // memory to store 1 bit of data	3
b)	One mark per bullet point to max 2 SR flip-flop has undefined / invalid / indeterminate state // JK flip-flop is stable Description of undefined / invalid / indeterminate state for SR // Description of why JK flip flop is stable JK flip flop has a clock pulse	2

		Α	В		X					
		0	0		1					
a)		0	1		1					
		1	0		1					
		1	1		0					
b) i)		S	R	Q	1	Q				
		1	0	0		1				
		1	1	0		1				
		0	1	1		0				
		1	1	1		0				
		0	0	1		1				
	S = 0 R = 0								_	
b) ii)	Produces $Q = 1$, $\overline{Q} = 0$ But Q and \overline{Q} should Becomes unstable	=1 // Q an be compl	nd Q have ements of	same each	e valu othe	ie er				3
c) i)	Clock (pulse)									1
c) ii)	All four possibilities are valid The 1-1 combination changes output to logical complement Unstable state avoided Invalid state cannot occur // the flip-flop is stable									1
d)	Memory // data stora Stores a single bit	ge								2

	Α	В		X			_			
	0	0		1						
a)	0	1		0				1		
	1	0		0						
	1	1		0						
		1		l _	l _	_	1			
			S	R	Q	Q				
	Initia	_	1	0	1	0				
b)	S change		0	0	1	0	(1)	4		
	R change	0	1	0	1	(1)				
	R change		0	0	0	1	(1)			
	S and R cha	anged to 1	1	1	0	0	(1)			
-) :)										
c) i)	Clock (pulse)							1		
	Max 2 marks	per proble	m – m a	ax 4 ma	ırks		<u>-</u>			
c) ii)	 Max 2 marks per problem – max 4 marks Problem 1 One combination of S and R gives NOT valid / indeterminate output // Q and Q have the same value The JK flip-flop does not allow for Q and Q to have the same value for any combination of inputs // Q and Q have to be complementary Problem 2 Inputs may not arrive at the same time The JK flip-flop has a clock pulse to synchronise inputs 									

		Α		В	Х								
		0		0	1								
a) i)		0		1	1								1
		1		0	1								
		1		1	0								
		Α	В		С	Х							
		0	0		0	1	1						
		0	0		1	1							
		0	1		0	1							
a) ii)	0 1			1	1							1	
		1	0		0	1							
		1	0		1	1							
		1	1		0	1							
	1 1				1	0							
					s	R	Q	Q					
			Initially		1	0	0	1	\dashv				
	R changed to 1				1	1	0	1		1			
b) i)	S changed to 0				0	1	1	0		1			3
			nanged		1	1	1	0		1			
	S		R chang		0	0	1	1					
	_		Q and	Q have	e same v	value	•	<u>. </u>					
b) ii)			Q and	Q shou	ıld be co	ompleme	ents of	each d	other				2
			Flip-flop	becor	nes uns	table		1 r	nark f	or eac	h point	, max 2	
									Init	ial	Fir	nal	
					Wor	king spa	ace	-	valı	ies	valı	ues	
	J	K	Clock			٠.			Q	Q	Q	Q	
	0	0	1						1	0	1	0	
c) i)	0	0	1	_				+	<u>0</u>	0	0	1	4
-,,,	0	1	1						0	1	0	1	
	1	0	1						1	0	1	0	
	1	0	1						0	1	1	0	
	1	1	1	1				-	0	0	0	0	
	<u> </u>		1							_		ded row	,

c) ii)	 □ S-R flip-flop has an invalid combination of S and R // The S_R flip flop allows both Q and Q to have the same value // S-R flip-flop inputs may arrive at different times □ The J-K flip-flop does not allow for Q and Q to have the same value // All four combination of values for J and K are valid // J-K flip-flop incorporates a clock pulse for synchronisation 	2
d)	□ A flip-flop can store either a 0 or a 1 □ Computers use bits to store data □ Flip-flops can therefore be used to store bits (of data) □ Memory can be created from flip-flops 1 mark for valid point, max 2	2

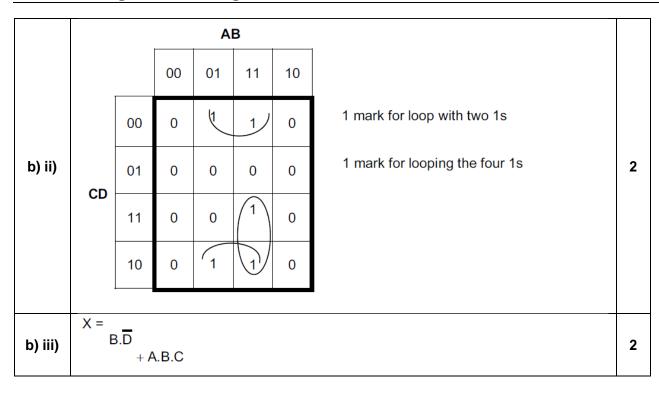
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-	(logic) Circuit // bi-stable Two Memory // data storage // registers // storing one bit of data JK/SR/D/T SR/JK/T/D	5
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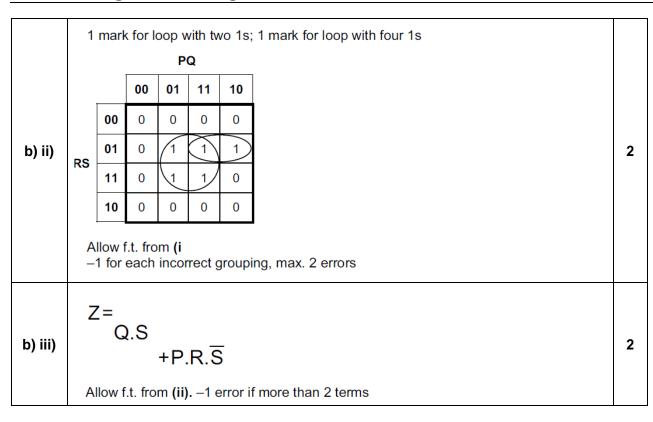
	•		Circuit 1		
		Α	В	Х	
a) ;)		0	0	1	
a) i)		0	1	1	
		1	0	1	
		1	1	0	
			Circuit 2		
		Α	В	X	
a) ii)		0	0	1	
		0	1	1	
		1	0	1	
		1	1	0	
b) i)	• circuit 1: $\overline{A.B}$ • circuit 2: $\overline{A} + \overline{B}$				
b) ii)	$\overline{A.B} \equiv \overline{A} + \overline{B}$				
c)	$ \frac{\overline{(A+B).B}}{\overline{(A+B)}} $ Mark as follows: $ \overline{(A+B)} $.B bar over whole express	sion			;
d)	$\overline{(A+B).B}$ $= \overline{(A+B)+B}$ $= (A+B)+\overline{B}$ $= A+(B+\overline{B})$ $= A+1$ $= 1$ allow f.t. from (c)				

		C	Circuit 1	
		А	В	X
a) i)		0	0	1
a) i)		0	1	0
		1	0	0
		1	1	0
		C	Circuit 2	
		А	В	X
2) ii)		0	0	1
a) ii)		0	1	0
		1	0	0
		1	1	0
b) i)	circuit 1 $\overline{\underline{A} + B}$ circuit 2 $\overline{\underline{A}}$ $\overline{\underline{B}}$			
b) ii)	$\overline{A+B} \equiv \overline{A}.\overline{B}$			
c)	$ \frac{\overline{(A.B)} + B}{\text{mark as :}} $ $ \frac{\overline{(A.B)}}{(A.B)} $ + B bar over whole express	ion		
d)	$\overline{(A.B) + B}$ $= \overline{(A.B) \cdot B}$ $= (A.B) \cdot \overline{B}$ $= (A.B) \cdot \overline{B}$ $= A.(B \cdot \overline{B})$ $= A.0$ $= 0$ allow f.t. from (c)			

a) i)	Ā .B.0	A.B	. C	A.B.C				3			
		00	01	AB 11	10						
a) ii)	c	_	0	1	0			1			
	1	0	1	1	0						
		00	01	AB 11	10						
a) iii)	c	0	0	1	0		1 mark for each loop				
	Allow f				0						
a) iv)	X = A.	+ B.						2			
				Al	В						
	ţ		00	01	11	10	1 mark row headings				
		00	0	1	1	0	1 mark column headings				
b) i)	CD	01	0	0	0	0	1 mark per 2 correct	4			
		11	0	0	1	0	rows (based on headings)				
		10	0	1	1	0					



a) i)	$Z=P.\overline{Q}.\overline{R} + P.\overline{Q}.R + P.Q.R$	3							
a) ii)	PQ 00 01 11 10 R 0 0 0 0 1 1 0 0 1 1	1							
a) iii)	1 mark each loop PQ 00 01 11 10 R 0 0 0 0 1 1 0 0 1 1 Allow f.t. from (ii)								
a) iv)	$Z=$ $P.\overline{Q}$ +P.R Allow f.t. from (iii)	1							
b) i)) 1 mark row headings. 1 mark column headings. 1 mark per 2 correct rows (based on headings) PQ 00 01 11 10 00 0 0 0 0 01 1 1 1 11 0 1 1 0 10 0 0 0	4							



a)	$S = (\overline{P} + (\overline{Q} + \overline{R})) \cdot R$ \overline{P} $(\overline{Q} + \overline{R})$ $(\overline{P} + (\overline{Q} + \overline{R}))$ R $(must be outside final brackets)$ 1 Or \overline{P} $(\overline{Q} + \overline{R})$ $\overline{P} + (\overline{Q} + \overline{R})$ $\overline{P} + (\overline{Q} + \overline{R})$ 1 $\overline{P} + (\overline{Q} + \overline{R})$ 1 1									
	_	Р	C)	R		Working space	S	_	
		0	0		0			0		
		0	0		1			1		
		0	1		0			0		
b)		0	1		1			1		2
		1	0		0			0		_
		1	0		1			0		
		1	1		0			0		
		1	1		1			0		
	2 mar	ks all	correct	, 1 mar	k seve	n correc	ct, 0 marks six or fewer	correct	_	
					Q	I	ſ			
c) i)		0	00	01	11	10				1
	R	1	1	1	0	0				
		1	1	P	Q	1	ı			
,			00	01	11	10				
c) ii)	R	0	0	0	0	0				1
	K	1	1	_1)	0	0				

c) iii)	$S = \overline{P} \cdot R$	1
d)	$S = (\overline{P} + (\overline{Q+R})) \cdot R$ $S = (\overline{P} + (\overline{Q} \cdot \overline{R})) \cdot R / / \overline{P} \cdot R + (\overline{Q+R}) \cdot R$ 1 $S = (\overline{P} \cdot R) + (\overline{Q} \cdot \overline{R} \cdot R)$ 1 $S = \overline{P} \cdot R + \overline{Q} \cdot 0$ $S = \overline{P} \cdot R + 0$ $S = \overline{P} \cdot R$ 1 1 1 1 1 1 1 1 1 1	3

a)	X = A B.C B + A.	A.(B + (E B.C	3 . C))								1 1 1	3
		Α	В	С		,	Workin	ıg Spa	ce	Х	1	
		0	0	0						0		
		0	0	1						0		
		0	1	0						0		
		0	1	1						0		
b)		1	0	0						1		2
		1	0	1						1		
		1	1	0						0		
		1	1	1						1		
	1 ma	ark first f	our entri	es, 1	mark	for the	last fo	our entr	ies			
c) i)				С	0	00 0	01 0 0	11 0 1	10 1			1
c) ii)				С	0	00 0	01 0 0	B 11 0	10			2
c) iii)	X = /	A.B + A.										2
d)	X = I	A.(<u>B</u> + (E A. <u>(</u> B + C A.B + A.)				1 (0			ust be co previous		2

a) i)	2 marks all products correct, 1 mark 2 or 3 products correct $X = \overline{A}.B.\overline{C} + \overline{A}.B.C + A.\overline{B}.\overline{C} + A.\overline{B}.C$	2							
a) ii)	1 mark for all correct bits AB 00 01 11 10 C 0 0 1 0 1 1 0 1	1							
a) iii)	1 mark for each correct loop AB 00 01 11 10 C 0 0 0 1 0 1	2							
a) iv)	1 mark per bullet – allow follow through from 4(a)(iii) • $\overline{A}.B$ • $+A.\overline{B}$ $X = \overline{A}.B + A.\overline{B}$	2							
b) i)									

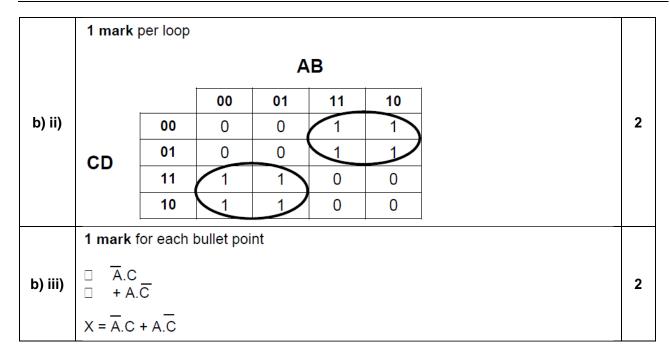
	1 mark for each correct loop							
b) ii)	CD	00 01 11 10	00 0 0 0		0 0 0 0		2	
b) iii)	1 mark per bullet $\overline{A}.B$ $+B.\overline{C}$ $X = \overline{A}.B + B.\overline{C}$						2	

a) ;)	1 mark for 2 or 3 cor	rect, 2 m	arks for	4 correct			-	2	
a) i)	$X = \overline{A}.B.C + A.\overline{B}.C + A.B.\overline{C} + A.B.C$								
	1 mark for the correc	t K-map							
a) ii)				A	AΒ				
			00	01	11	10		1	
	С	0	0	0	1	0			
	C	1	0	1	1	1			
	1 mark for each loop	max 3					-		
	АВ								
a) iii)			00	01	11	10		2	
	С	0	0	0		0			
		1	0 (\supset			
	1 mark for each pair.	Allow fo	llow thro	ugh from	(iii)				
a) iv)	• A.B							2	
4,117	• +B.C • +A.C							_	
	X = A.I	3+B.C+	A.C						

	1 mark per bullet poi			, booding	re volu	os only			
	Correct column I Correct column I mark for 2 correct I (based on headings)	headings rows or o	s and row	/ heading	gs – orde	r	or columns		
b) i)				A	\ Β			4	
b) i)			00	01	11	10		4	
		00	0	1	1	0			
	c.p.	01	0	0	1	0			
	CD	11	0	0	1	0			
		10	0	0	1	0			
	1 mark per loop						-		
	АВ								
			00	01	11	10			
b) ii)		00	0 <			0		2	
	c n	01	0	0	1	0			
	CD	11	0	0	1	0			
		10	0	0	\bigcup	0			
	1 mark per bullet point:								
b) iii)	• A.B • +B.C.D X = A.I	3+B.C.Ī	5					2	

a)	1 mark per bullet point to max 3: • Correct use of Idempotent law Y = Y.Y Y = Y + Y • Correct use of Complement law $0 = Y.\overline{Y}$ $1 = Y + \overline{Y}$ • Correct use of Distributive law $X(Y + Z) = X.Y + X.Z$ • Correct use of Redundancy law $X.\overline{Y} + Y = X + Y$ • Correct use of identity law $X.1 = X$ 1 mark for the correct answer For example: $X = A.\overline{B}.\overline{C} + A.B.\overline{C} + A.B.C \qquad Idempotent law$ $X = A.\overline{B}.\overline{C} + A.B.\overline{C} + A.B.C \qquad Distributive law$ $X = A.\overline{B}.\overline{C} + A.B.\overline{C} + A.B.C \qquad Complement/Inverse law$ $X = A.\overline{C}.(B + B) + A.B.(C + C) \qquad Complement/Inverse law$ $X = A.\overline{C}.(B + B) \qquad Correct answer$ $X = A.\overline{B}.\overline{C} + A.B.C \qquad Distributive law$ $X = A.\overline{C}.(B + B) + A.B.C \qquad Complement/Inverse law$ $X = A.\overline{C}.(B + B) + A.B.C \qquad Complement/Inverse law$ $X = A.\overline{C}.(B + B) + A.B.C \qquad Complement/Inverse law$ $X = A.\overline{C}.(B + B) + A.B.C \qquad Redundancy Law$ $X = A.(C + B.C) \qquad Redundancy Law$	4
b) i)	1 mark for first four as 0, 1 mark for 1011 A	2
b) ii)	1 mark for correct K-map AB 00 01 11 10 C 1 0 0 1 1 0	1
b) iii)	1 mark for each correct loop to max 2 AB 00 01 11 10 C 1 0 0 1 0	2
b) iv)	1 mark per bullet point: • A.C • + A.B X = A.C + A.B	2

					Δ	В				
a) i)				00	01	1	1	10		1
a) 1)		C)	1	1	0		1		•
			ı	1	1	0		1		
	1 mark f	or each	correc	t loop		۸D				
				00	01	AB	11	10		
a) ii)	С		0	1	1		0	/1	+	2
			1	1	1	7	0	1		
	1 mark p	er bullet	point							_
a) iii)	$\begin{array}{c c} \hline & \overline{A} \\ \hline & \overline{B} \end{array}$									2
	$X = \overline{A} \square \overline{E}$	3								
	1 mark c	orrect va	alues an	nd orde	er of row	and	colu	mn head	lings	
	3 marks 2 marks 1 mark ta	table en	tries co	ntain c	ne erro	r (bas	ed o	n headir	ngs) or	
F) ;)				A	λB					
b) i)			00	01	1 1	1	10)		4
		00	0	0		1	1			
		01	0	0		1	1			
	CD	11	1	1	(0	0			
		10	1	1		0	0			



a)	1 mark for 3 or 2 marks for all 8 $X = \overline{A}.\overline{B}.\overline{C} + \overline{A}.\overline{B}.$	5 corre	ect prod	ucts	4. <i>B</i> .C			
	1 mark for corre	ect ans	swer	А	В		_	
			00	01	11	10		
b)		0	1	0	0	0		
	C	1	1	1	1	1		
	1 mark per corre	ect loo	p	A	В	-		
			00	01	11	10		
c)		0	1	0	0	0		
	C -	1	1	1	1	1		
	1 mark per bulle	et poin	t.					
d)	•							
	$X = \overline{A}.\overline{B} + C // 2$	X = C -	+ Ā . B					

a) i)	1 mark for each 2 correct products, i.e. 3 marks for 6, 2 marks for 4 or 5, 1 mark for 2 or 3 $\mathbf{X} = \overline{\mathbf{A}}, \overline{\mathbf{B}}, \overline{\mathbf{C}} + \overline{\mathbf{A}}, \overline{\mathbf{B}}, \mathbf{C} + \overline{\mathbf{A}}, \overline{\mathbf{B}}, \overline{\mathbf{C}} + \overline{\mathbf{A}}, \overline{\mathbf{B}}, \overline{\mathbf{C}} + \overline{\mathbf{A}}, \overline{\mathbf{B}}, \overline{\mathbf{C}} + \overline{\mathbf{A}}, \overline{\mathbf{B}}, \overline{\mathbf{C}} + \overline{\mathbf{A}}, \overline{\mathbf{B}}, \overline{\mathbf{C}}$	3
	X = A.B.C + A.B.C + A.B.C + A.B.C + A.B.C 1 mark for the correct K-map	
a) ii)	AB 00	1
a) iil)	1 mark for each correct loop AB C 00 1 1 00 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1	2
a) iv)	 1 mark per bullet point: A + B X = A + B // X = B + A 	2
b)	$X = \left(\overline{(W} + X).(Y + \overline{Z})\right)$ One mark for correct use of <u>De Morgan's law</u> to + • $X = \overline{(W} + X) + \overline{(Y + \overline{Z})}$ One mark for correct use of <u>De Morgan's law</u> + to • $X = \overline{W}.\overline{X} + \overline{Y}.\overline{Z}$ One mark for correct answer • $X = W.\overline{X} + \overline{Y}.Z$	3

a)	1 m X = Y =	ark 2 A. B. A. B.	$C + \overline{A}$. C + A.	oroducts $\frac{B. \overline{C} + A}{\overline{B}. C + A}$	s correct a. B. C + a. B. C +	A.B.C A.B.C		rect no	incorre	ct produ	icts seen,	4
	On	e ma		each col		-map ma	x 2	0	UTPUT Y AB			
b)			00	01	11	10	00	01	11	10		2
,		0	0	1	0	1	0	0	1	0		
	С	1	1	0	1	0	0	1	1	1		
c) i)				OUTPUT e mark OUTPUT AB 01	vertica		rrect one		horizon OUTPUT AB 11		s correct	3
		0	0	1	0	1	0	0		0		
	С	1	1	0	1	0	0	1	Q			
c) ii)			rk for e		rrect pr	oduct an	nd no inc	orrect p	oroducts	s max 3		3
d)	Log X: Y:	jic cii	rcuit:	Full Ad Sum Carry	dder							3

a)	For X =	X 1 m Y 2 m · Ā.B.C · Ā.B.C	arks f + A.B	or 3 p .C	roduc				ther p	roduc	ts see	n	3
	One	e mark	for e	ach co	orrect	K-mar	o m	ax 2					
				OUTF A	PUT X .B					PUT Y .B	,		
b)			00	01	11	10		00	01	11	10		2
	С	0	1	0	0	0		0	0	0	0		
		1	0	0	1	0		0	1	1	1		
		mark mark						orrect	and r	no oth	ers ma	ax 2	
-> :>				OUTF A						PUT Y	•		•
c) i)			00	01	11	10		00	01	11	10		2
	С	0	1	0	0	0		0	0	0	0		
		1	0	0	1	0		0	1	1	1		
c) ii)		mark + B.C		ach co	rrect	produ	ct						2

a) i)		rk for 2 c				rect and	no other terms	
	One ma	arks for fu ark for a h arks for a	K-map wi	th one er		errors		
				PC	Q			
a) ii)			00	01	11	10		
٠,,		00	1	0	0	0		
	RS	01	0	0	1	0		
		11	0	0	1	0		
		10	1	0	0	0		
	One ma	rk for ead	ch correc	t loop m a	ax two			
				Р	Q			
			00	01	11	10		
a) iii)		00	(1)	0	0	0		
	RS	01	0	0	1	0		
	110	11	0	0	1	0		
		10	1	0	0	0		
a) iv)	 P.Q +P.0 							
b)	One ma One ma	rk for cor rk for cor rk for cor	rect use	of complor	utive law ementary dancy law otent law	V		
		R.S+P.G			Q.R.S			
		$\overline{S}(\overline{R}+R)$	`	,	20(4)			
	X = P.Q.	S + P.Q.	5 // P.Q.S	5(1)+P.0	J.S(1)			

a) i)	terms	rk for 2 co				for 4 cor	rect terms and no othe	r	
		rks for fu rk for a K irks for a	(-map wit	h one en		errors			
				P	Q				
a) ::)			00	01	11	10			
a) ii)		00	1	1	0	0			
	DO	01	0	0	0	0			
	RS	11	0	0	1	1			
		10	0	0	0	0			
	One mai	rk for eac	h correct	loop m a	x 2	1	·		-
				P	Q				
			00	01	11	10			
a) iii)		00	(1	1)	0	0			
~,,		01	0	0	0	0			
	RS	11	0	0	1	1)			
		10	0	0	0	0			
	0			•	•	0			_
a) iv)	• PR.		ITIL						
a) iv)	• +P.F X = PR.5	R.S 5 + P.R.S	3						
		rk for cor		of distribu	utive law				_
	One mai One mai Max two	rk for cor							
b)	e.g.								
	$X = \overline{P}.\overline{Q}.$	R.S + P.C			.Q.R.S				
		$\overline{S}(\overline{Q}+Q)$	`	,					
	X = PR.S	5 + P.R.S	5 // P.R.S	6 (1) + P.R	.S(1)				

a) i)	One mar $X = \bar{P}.\bar{Q}$						t	2
	One mar	k for eac	ch correct	pair of r	ows/colu	ımns to r	max 2	
				Р	Q			
-> ::>			00	01	11	10		2
a) ii)		00	1	0	0	0		2
	RS	01	1	0	0	0		
	1.0	11	0	0	1	0		
		10	0	0	1	0		
	One mar	k for eac	h correct	loop m a	ax 2			
				P	Q			
			00	01	11	10		_
a) iii)		00	<u>(1)</u>	0	0	0		2
	RS	01	(1)	0	0	0		
	N3	11	0	0	(1)	0		
		10	0	0	1/	0		
	One mar		llet point					
a) iv)	• \overline{P} . \overline{Q}							2
	• +P.	•						_
	$X = \overline{P}.\overline{Q}$							
		rk for cor						
		rk for cor $\bar{R}(\bar{S} + S)$				4W		
	_	$\bar{R}(3+3)$			1			
		, . (+ <i>)</i>	= - 6 (-)	,				
b)	or							2
	Two ma	rks for co	orrect use	e of redu	ndancy l	aw		
		$ar{Q}.ar{R}).ar{S}+Q.ar{R}+P.ar{Q}$		S + (P.	$Q.R).\bar{S}$ -	+ (P. Q. R	?).\$	

a) i)		ark for 2 (rks for 4 P.Q.R.S	correct		2
		ark two co arks for fo						
			00		Q	40	7	
a) ii)		00	00 1	01	11	10	1	2
	RS	01	0	0	0	0		
	RS	11	0	1	1	0		
		10	0	0	0	0		
	One ma	ark for ea	ch correc	ct loop m	ax 2			
				Р	Q			
\			00	01	11	10		
a) iii)		00	1)	0	0	(1)		2
	RS	01	0	0	0	0	-	
		11 10	0	0	1	0		
	One m	ark per l	hullet no	int				
	• QF	•	builet po					
a) iv)		ì.R.S						2
	$X = \overline{Q}R$		R.S or X =	= Q.R.S	$+\overline{Q}\overline{R}.\overline{S}$			
	One ma	ark for co	rrect use	of distrib	outive law	V		
	One ma	ark for co	rrect use	of comp	lement la	aw		
	$X = \overline{Q}.\overline{F}$	Ŗ.Ē(<u>P</u> +P) + Q.R.S	$(\overline{P} + P)$				
	$X = \overline{Q}.\overline{F}$	₹. ऽ (1) + ऽ	Q.R.S (1)					
b)	Or							2
	Two ma	arks for c	orrect us	e of redu	ındancy la	aw		
	$X = (\overline{Q}.$	$\overline{R}.\overline{S})\overline{P}+($	(Q.R.S)P	+ (Q.R.S	$\bar{P} + (Q.F)$	R.S)P)		
	$X = \overline{Q}.\overline{F}$	₹.\$ + Q.R.	.S					

	One mark per	two cor	rect prod	ucts (Ma	x 3)		
Α	(Z =) A B C D		D + A E	 B C D + A	_ A B C D +		
	A B C D + A		o correct	t rows or	columns	(Max 2)	
		ovory th	0 001100		В	(max 2)	
	`		00	01	11	10	
b) i)		00	0	0	1	0	
D) 1)	CD	01	0	0	1	1	
	CD	11	0	0	1	1	
		10	0	0	1	0	
	One mark for	correct le	оор (Мах	(2)			
				Α	В		
			00	01	11	10	
		00	0	0		0	
b) ii)	CD	01	0	0	1	1	
	CD	11	0	0	1	1	
		01	0	0	1	0	
	One mark per	correct	marking	point (Ma	ax 2)		
b) iii)	 A B // A + A D // + 	D	ai Milg	Pour (Me			
	(Z =) A B + A	D // A D	+ A B				
b) iv)	(Z =) A (B + D	-	· · · · · · · · · · · · · · · · · · ·				