PROJECT MANAGEMENT

18ME653 -PE3

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BOOKS

1. S Choudhary, "Project Management" Tata McGraw Hill Education Private Limited New Delhi.

2001

- 2. Harold Kerzner, "Project Management: A Systems Approach To Planning, Scheduling And Controlling", CBS publisher and distributers
- 3. L S Srinath "PERT and CPM Principles and Applications" Third Eddition
- 4. Dr. P N Modi "PERT and CPM"

TOOLS AND TECHNIQUES OF PROJECT MANAGEMENT

- Manager or admittatur constantly for looks forward to those techniques or methods which helps him in
- Planning
- Scheduling
- Controlling such activates
- The concept of network planning and critical path analysis
- ➤ It is technique through which largerer projects are broken down to individual jobs or events and arranged in logical Network
- > PERT and CPM are two management techniques.

Unit-4. Sheduling Techniques

- > 1. Gantt charts
- > 2. Milestone charts
- > 3. Work break down structure
- > 4.PERT and CPM Networks

GANTT CHART

Preparing a pattern for casting 4 weeks Preparing a mould 2 weeks Casting and cleaning operation of A 1 week Heat-treatment of A 2 weeks Obtaining and installing machine M 7 weeks Machining part B 5 weeks Assembling parts A and B3 weeks Preparing the test rig 4 weeks

Prepare pattern Prepare mould Cast & clean A Heat-treat Instal M/C M Machine part B Assemble A & B Prepare test rig Test assembly Pack & dispatch

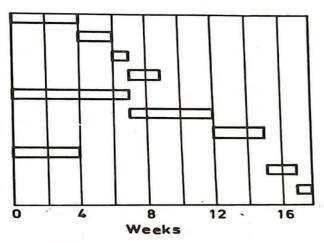


FIGURE 1-2

MILESTONECHARTS AND WORK BREAKDOWN STRUCTURER

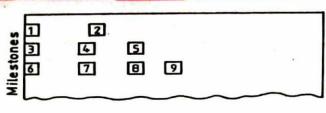


FIGURE 1-6

interdependencies between events. In a milestone chart, the events are in chronological, but not in a logical, sequence. A natural extension of the milestone chart was the network, where the events are connected by arrows in a logical sequence. This is shown in Fig. 1-7.

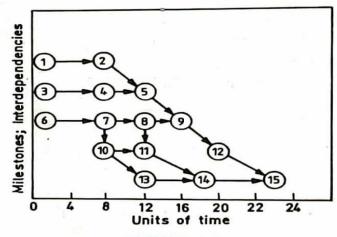
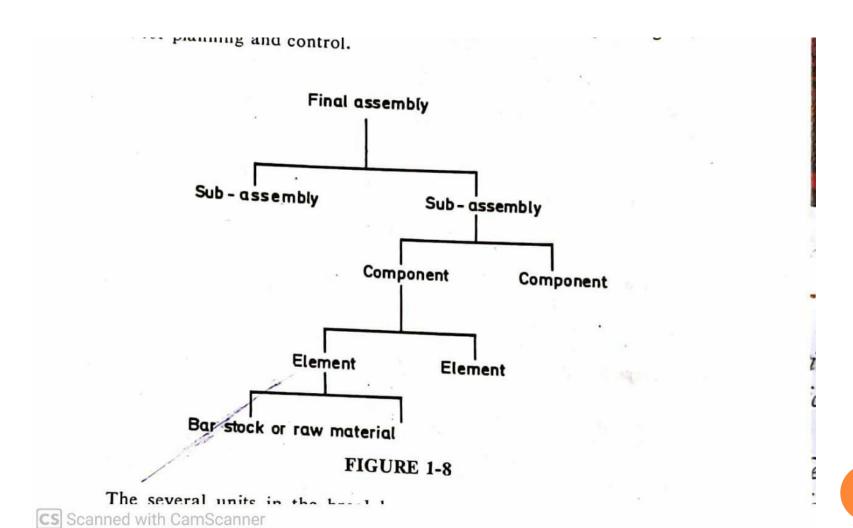
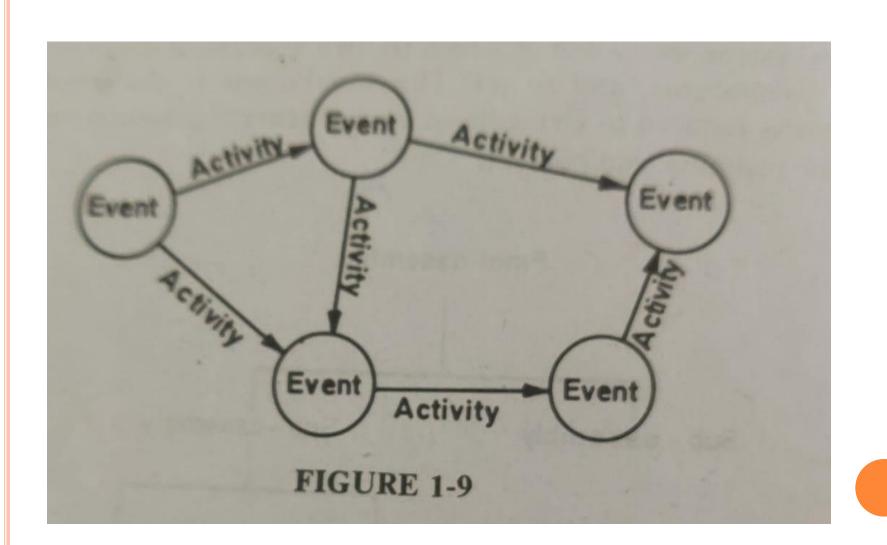


FIGURE 1-7





CPM-AND-PERT

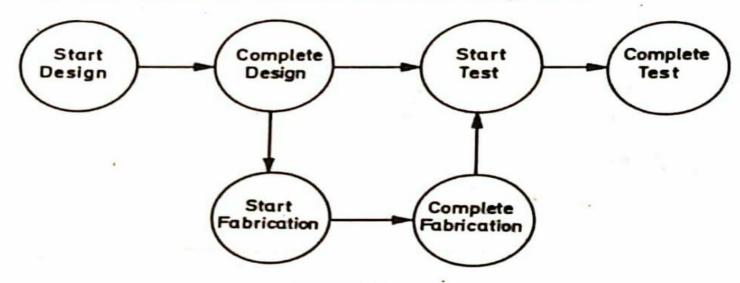


FIGURE 1-11

NETWORK TECHNIQUES

2.4 NUMBERING THE EVENTS

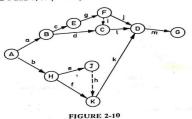
The event numbers should in some respect reflect their logical sequences. When a complicated network has been formed after numerous additions and deletions (which are unavoidable before a final acceptable network is obtained), the problem of assigning numbers to the events has to be considered. A rule devised by D. R. Fulkerson reduces this sequential numbering to the following steps:

- (i) An "initial" event is one which has arrows coming out of it and none entering it. In any network, there will be one such such event.
- (ii) Delete all arrows emerging from event 1. This will create at least one more "initial event".
 - (iii) Number these new initial events as "2, 3, ...".
- (iv) Delete all emerging arrows from these numbered events which will create new initial events.
 - (v) Follow step (iii).
- (vi) Continue until the last event which has no arrows emerging from it is obtained.

Let us consider the network shows !

emerge. Deleting these, we get events E and J which have only emerging arrows. Notice carefully that event K will have an arrow entering after deleting e and f, and hence cannot form initial events. Number events E and J as 4 and 5, respectively.

(v) Delete arrows g, d, and h. These give events F and K which will be numbered 6 and 7, respectively.



(vi) Delete arrows i, j, and k. Events C and D will get numbers

(vii) Finally, Event G which has no emerging arrows will form the event and is numbered 10.

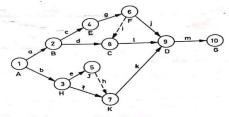


FIGURE 2-11

NETWORK TECHNIQUES

2.4 NUMBERING THE EVENTS

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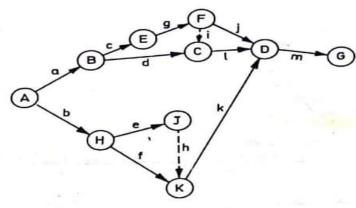
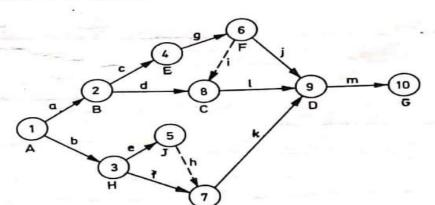


FIGURE 2-10

(vi) Delete arrows i, j, and k. Events C and D will get numbers 8 and 9, respectively.

(vii) Finally, Event G which has no emerging arrows will form the end event and is numbered 10.

The network with events numbered is shown in Fig. 2-11.



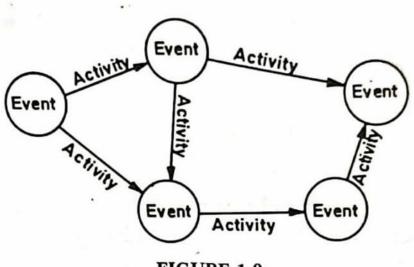
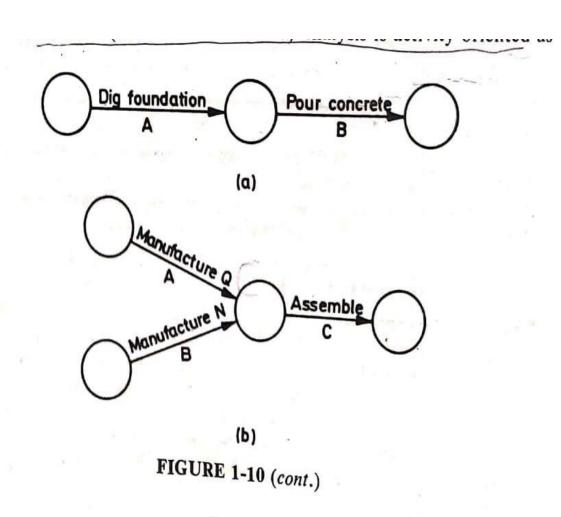


FIGURE 1-9

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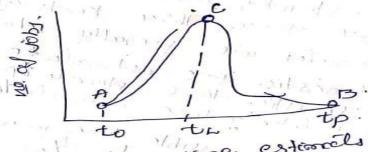


+ 6KI - event oxeculed Hireed Gepta there are ceclivities for which the associated time ducaration can be accurately estimated Activity - Each ceclevily - kekes some took doction tor its completion - This direction depends upon the necture of cectivity - there are some activities which are performed really no no data of trone existis dor such cultivities - Their time duration is un certain. - such cultavities use coolleal vousiceble cultavities ex: concedive activities such as sugarel design & development work by there are some activities dor which associated time can be estimated acadaly cere colleal deterministic type. these care Repetative in nattace PERT > Projects which consists i) the prejects which compaise of valuable type activities associated of deleanmestic type with phobabilistic time estimale heandled by CPM. empley PERT of network 2) CPM is cretivity >> PERT is event of enterol odiented 3) It is loased on 3-dime 3) it is based on lingle ine

estimates they care everage time requester? optimestic, pessimestitto execute the activities. a most likely time - PETR Takes in account 4) To det single -kine some uncertainty while hes-toliccel deter roxig be estimaling autivity available or best way of -tomes paedicting decta is by intelligence guess. The experie nce es pecison his techniqueal 5) Knowder not sake in to a/c 15 Bucquency distribution cult while extince we know three time estimates for PERT activity - optimestic, pessionxestic + most likely. - 90 the sange and optimestic to pessionestic there can be number of true estimates A the dequency distailabilion dos these com this is known as B-distainelion plotted we get to 1 to occurs oil the end. to - need not to be cel the michpt the Isaqueency destabletion rever. Mean (le) of standard devication Scanned with CamScanner

Time estimales - celter the plenning is delided: next is time suguished to execution of each activity - there is always plessure to do the job in short time. - many uncertantis are there to account all these there kinds as time estimates are generally obtained. i) The optimestic time estimate to 2) The pessionistic time estimate to 3) The most likely time estimate to 1) this is shootest possible time in which activity can be completed. - Ito 2) mæximum possible time to complete - considering abnormal situations - considerines abno 3) -time estimale which lies between LL optimestic a pessimistic time estimate - it assumes that things go in notional - things are usual. - All there cere selected in deces week or nosthy

- these 3 time estimates are given by peason based on his expedience 4 in formation - There is relationship between these time estimates a this is given by Arequeorcy distablelion cerve



time estimals to - every thing goes according to plan a solith minimum amount of difficulties.

tp- this time cessiernes their everything will not go according to plan 4 max portential disticulties will develop.

the the time that is there in The functional rounages

Transfer ...

to + 4-lm + 1p V-is variance (un centerainty) 3) V= 52 = (tp-to)2 Expertent time or everage-time at an activity is taken equal to mean. te-le to+4tm+tp wheel is the partiality this excepted excepted excepted in this excepted in this excepted in their excepted in the excepted in their excepted in the excepted in their excepted in the excepted in their excepted - Greater the varience the larger is uncertainty - vosiènce is unceeltaine 14. ex to = 5 days to = 9 days. LE = 5+4x7+9 = 14+28 = 42 = 7 decys. V = (9-5)2 - 0.444. tu-4 tm-7 tp-10 te = 4+4x7x10 - 7 deceys Compeding these two tt general scene V is more in and case - uncertaining and Completion of activity is more

different, and for each path we can get three time estimates based the optimistic, the most likely, or the pessimistic time estimate. These

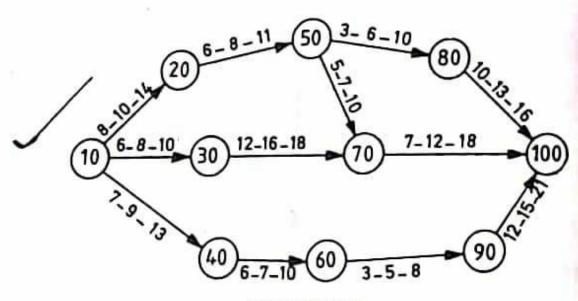


FIGURE 3-5

shown in Table 3-2.

TABLE 3-2



20-50. Proceeding this way, we arrive at events 70 and 80, the latter being the end event. The sequentially numbered network is shown in Fig. 3-9.

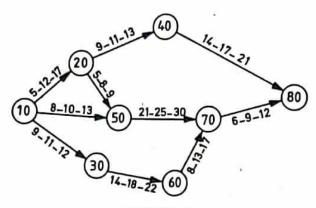


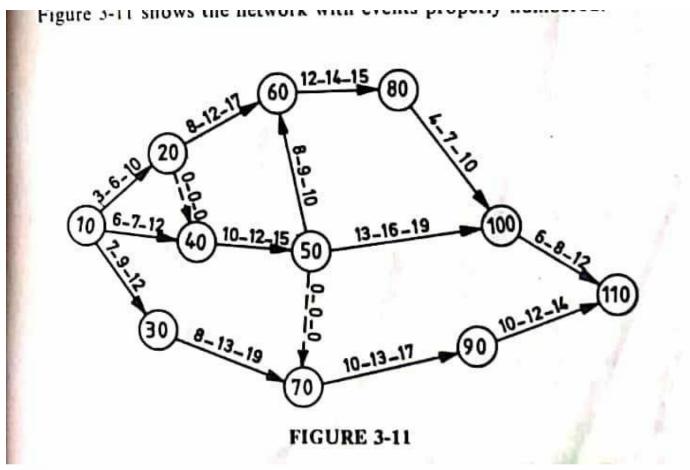
FIGURE 3-9

We shall now enter the data as shown in Table 3-4.

TABLE 3-4

Predecessor event	Successor event	10	t_L	t _P	σ^2	1 _E
10	20	5	12	17 -	4.0	11.67
10	30	9	11	12	0.25	10.83
10	50	8	10	13	0.69	10.17
20	40	9	11	13	0.44	11.00
20	50	5	8	9	0.44	7.67
30	60	14	18	22	1.78	18.00
40	80	14	17	21	1.36	17.17
50	70	21	25	30	2.25	25.18
60	70	8	13	17	2.25	12.83
70	80	6	9	12	1.0	9.00

In entering the event numbers, first the number of the start event is





			-	-5.01		
Predecessor event	Successor event	10	I_L	t _r	σ^2	I.E
10	20	3	6	10	1.36	6.17
10	30	7	9	12	0.70	9.17
10	40	6	7	12	1.0	7.67
20	40	0	0	0	0.0	0.0
20	60	8	12	.17	2.25	12.17
30	70	8	13	19	3.36	13.17
40	50	10	.12	15	0.70	12.17
50	60	8	9	10	0.11	9.00
50	100	13	16	19	1.00	16.00
60	80	12	14	15	0.25	13.83
70	90	10	13	17	1.36	13.17
80	100	4	7	10	1.00	7.00
90	110	10	12	- 14	0.44	12.00
100	110	6	8	12	1.00	8.33

The predecessor and successor events are entered in a manner similar to that explained in Example 1. The other entries are straightforward.

OUESTIONS

Computation 1

4.1 EARLIEST EXPECTED TIME

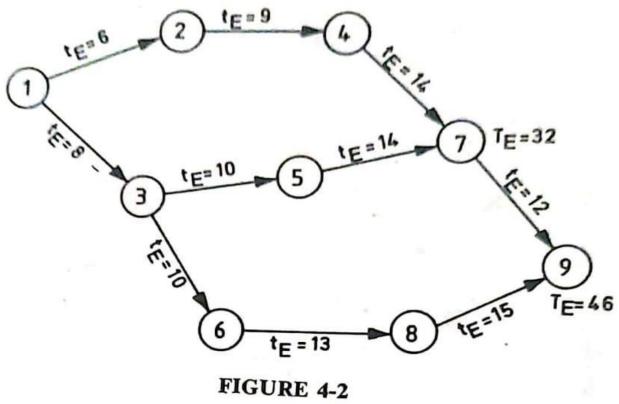
4.1 EARLIEST EXPECTED.

We are able to calculate the expected time estimates, we are able to calculate the expected time estimates, we are able to calculate the expected time. From the three time estimates, we are networks, enables us to find the or the average time which, in simple networks, a systematic method of determine or the average time which, in simple or the average time average time which, in simple or the average time avera critical path. But in large networks, a necessity. To achieve this, we need the critical path or paths becomes a necessity. The earliest expected the critical path or paths becomes—the earliest expected time to calculate, for each event, two time estimates—the earliest expected time which we shall discuss in this section, and the latest start time.

we should notice one important point here. The optimistic, most We should notice one important all refer to an activity or a job. The likely, and pessimistic time estimates, all refers to an activity or likely, and pessimistic tune estimates, and the least expected time or the average time te also refers to an activity connecting two events. However, both the earliest expected time and the latest start time refer to events.

The earliest expected time—denoted by TE—refers to the time when an event can be expected to be completed. It is computed by adding the $t_{E's}$ of the activity paths leading to that event. Let us consider the network shown in Fig. 4-1.

FIGURE 4-1



the predecessor and successor

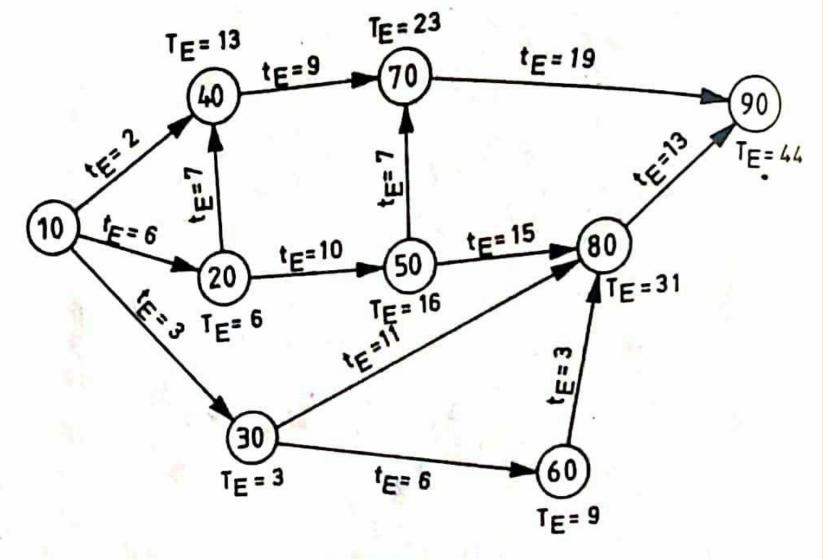


FIGURE 4-3

4.3 LATEST ALLOWABLE OCCURRENCE TIME

Here, we shall discuss the second time estimate regarding the event. The latest time by which an event must occur to keep the project on schedule is known as the latest allowable occurrence time, and is denoted by T_L . To explain this, let us assume that it has been agreed to complete the project within a certain allotted time called the contractual obligation time, denoted by T_S . This time refers to the occurrence of the end event. Let us consider once again the simple network already referred to in Fig. 4-1 and repeated in Fig. 4-4. Let us say that the contractual obligation time T_S for the project is 27. This means that the end event 4 must occur 27 units of time after the project is initiated.

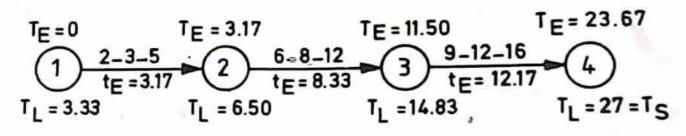


FIGURE 4-4

PERT AND CPM 48

When an event has more than one successor event, more than one T, will be available. Consider the network shown in Fig. 4-5.

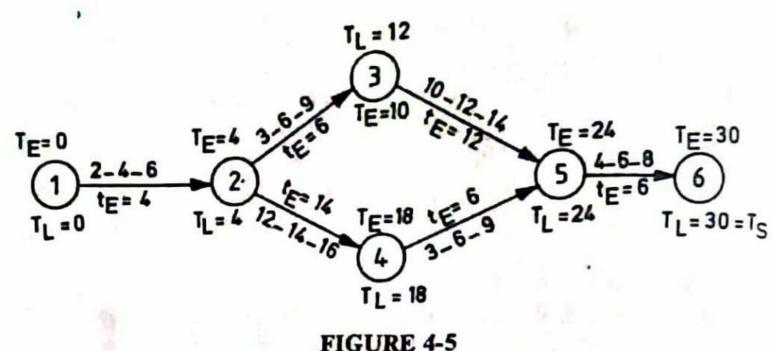


FIGURE 4-5

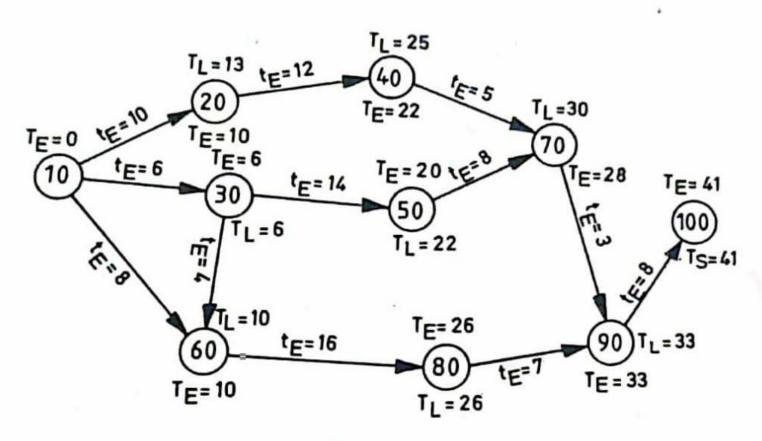


FIGURE 4-6

PROBLEM 1

Design the Network for fallowing set of activities of project. Determine

- i) Earliest Start
- ii) Earliest Finish
- iii) Latest start
- iv) Lates Finish and Identify the critical path and project duration.

Jo	b i-j
Successor event j	Predecessor event i
100	90'
90	80
90	70
80	60
70	50
70	40
60	30
60	10-
50	30-
40	20-
30	10-
20	10.

3 Je								
Successor Sevent j	Predecessor event i	$t_{\mathcal{O}}^{li}$	1 ^P	10	tE	T_E^j	T_L^i	T_L^+
100	90'	6	8	10	8	41	33	41
90	80	5	7	9	7	33	26	33
90	70	2	3	4	3	31	30	33
80	60	14	16	18	16	26	10	26
70	50	7	8	9	8	28	22	30
70	40	4	5	6	5	27	25	30
60	30	2	4	6	4	10	6	10
60	10-	6	8	10	8	8	2	10
50	30-	10	14	18	14	20	8	22
40	20-	8	12	16	12	22	13	25
30	10-	4	6	10	6	6	0	6
20	10-	4	10	16	10	10	. 3	13

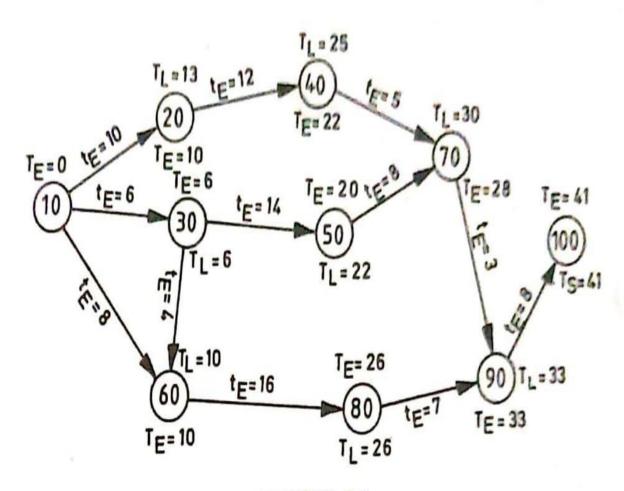


FIGURE 4-6

Jo Jo					-				
Successor event j	Predecessor event i		10	tĽ	$t_P^{(j)}$	t_E^{ij}	T_E^i	T_L^i	\tilde{I}_L^{\dagger}
100	90'	6	8	10	8	41	33	41	
90	80	5	7	9	7	33	26	33	
90	70	2	3	4	3	31	30	33	
80	60	14	16	18	16	26	10	26	
70	50	7	8	9	8	28	22	30	
70	40	4	5	6	5	27	25	30	
60	30	2	4	6	4	10	6	10	
60	10-	6	8	10	8	8	2	10	
50	30-	10	14	18	14	20	8	22	
40	20/	8	12	16	12	22	13	25	
30	10-	4	6	10	6	6	0	6	
20	10-	4	10	16	10	10	. 3	13	

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Jo		n		4:		-		
Successor event j	Predecessor event i	$t_{\mathcal{O}}^{li}$	1 ¹ L	t p	tE	T_E^i	T_L^i	T_L^+
100	90'	6	8	10	8	41	33	41
90	80	5	7	9	7	33	26	33
90	70	2	3	4	3	31	30	33
80	60	14	16	18	16	26	10	26
70	50	7	8	9	8	28	22	30
70	40	4	5	6	5	27	25	30
60	30	2	4	6	4	10	6	10
60	10-	6	8	10	8	8	2	10
50	30-	10	14	18	14	20	8	22
40	20-	8	12	16	12	22	13	2:
30	10-	4	6	10	6	6	0	
20	10-	4	10	16	10	10	- 3	1.

6.12 FLOAT

In the previous few sections, we considered the time elements associated with events. The next point of interest is the activities and their start and finish times. We can define the following for a given activity i-j:

Earliest Start Time This is the earliest occurrence time for the event from which the activity arrow originates, i.e., T_E^i .

Earliest Finish Time This is the earliest occurrence time for the event from which the activity arrow originates plus the duration for the activity, i.e., $= (T_E^i + t^{ij})$.

The total duration of time available for any job is the difference between its earliest start time and latest finish time. If *i-j* is the job under consideration, then

maximum time available = $T_L^i - T_E^i$.

If job i-j requires only t^{ij} units of time for its execution, the <u>total float</u> for job i-j is the difference between the maximum time available for the job and the actual time it takes, that is,

total float for
$$i-j = (T_L^J - T_E^i) - t^{ij}$$

= $(T_L^J - t^{ij}) - T_E^i$.

We notice this is equal to the latest start time for the activity minus its earliest start time. This value is obtained by taking the difference between

the values given in the sixth and fourth columns.

The second type of float defined is the "free float" for an activity. This is based on the possibility that all events occur at their earliest times, i.e., all activities start as early as possible. Consider two activities i-j and jk where the second activity j-k is a successor activity to i-j. Let the earliest occurrence time for event i be T_E^i and for event j, T_E^i . This means that the earliest possible start time for activity i-j is T'_E , and for activity j-k, it is Tk. Let the duration for activity i-j be til. Assume that i-j starts at $T_{\mathcal{S}}$ and takes t^{ij} units of time, and that the next activity j-k cannot start because its earliest possible start time T_E^i is greater than $(T_E^i + t^{ij})$. Then,

$$T_E^I - (T_E^I + t^{IJ})$$

is called the free float for activity i-j, i.e.,

free float for
$$i-j = T_E^j - (T_E^i + t^{ij})$$

 $=T_E^i$ — earliest finish time for *i-j*.

- 24-24 =

 $= T_E - \text{earliest finish time for } i-j$.

We can restate this as follows:

The free float for activity i-j is the difference between its earliest finish time and the earliest start time for its successor activity.

For example, in Table 6-2, the earliest finish time for activity 5-7 is 20 and the earliest start time for its successor activity 6-7 or 6-8 is 22. Hence, the free float for activity 5-7 is 2. The values for free float have been entered in column nine.

Another type of float, termed the "independent float", is also defined. Its basis is as follows. Let i-j be the activity of interest and h-i and j-k, respectively, be its predecessor and successor activities (Fig. 6-14). Let the preceding job h-i finish at its latest possible moment, which is T_L^i , and the succeeding job j-k start at its earliest possible moment, which is

FIGURE 6-14

 T_E^i . Then, activity *i-j* can take up any duration from t^{ij} to $(T_E^i - T_L^i)$ without in any way affecting the network. The difference between $(T_E^i - T_L^i)$ and t^{ij} is called the *independent float*, i.e.,

independent float for $i-j = (T_E^j - T_L^i) - t^{ij}$.

Consider activity 6-8 shown in the network of Fig. 6-13 redrawn partially in Fig. 6-15.

The latest finish time for the job preceding 6-8 is 22 and the earliest

(1) (2) y Job	(2)	(3)	(4) . Earl	(5) iest	(6) (7) Latest		(8) Total	(9) Free	(10) Indepen-
i	- j	Duration	Start	Finish,	Start.	Finish 3	float	float	dent float
1	2	4 .	0	4	0	4	0	0	0
1	3	12	70-54	12	2	14	2	0	0
1	4	10	0	10	2	12	2	(2)	2
2	4	. 8	4	12	4	12	0	0	0
2	5	6	4	10	6	12	2	0	0
3	6	8	12	20	IF 14	22	2	2	0
4	6	10	12	22	12	22	0	0	0
5	7	10	10	20	12	22	2	2	0
6	7	0	22	22	22	22	0	0	0
6	8	8	22	30	24	32	2	2	2
7	8	10	22	32	22	32	0	ō	0
8	9	6	32	38	32	38	0	0	0

III I HOIL . . ,