

Algorithmic Methods for Mathematical Models

Course Project

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May 26, 2024

Formal Problem Definition

- ▶ n : number of products
- ▶ x : height of the suitcase in millimeters
- ▶ y : width of the suitcase in millimeters
- ▶ c : limit to the total weight of the suitcase in grams
- ▶ p_i : price of the i -th product in euros
- ▶ w_i : weight of the i -th product in grams
- ▶ s_i : side length of the i -th product's (square) box in millimeters

Decision Variables

- ▶ Chosen_i : binary variable that is 1 if object i is chosen, and 0 otherwise.
- ▶ PointsX_i : the x-coordinate of the bottom-left corner of object i .
- ▶ PointsY_i : the y-coordinate of the bottom-left corner of object i .
- ▶ $\text{Overlap}_{i,j,d}$: binary variable indicating if objects i and j do not overlap in direction d , where $d \in \{1, 2, 3, 4\}$.

Objective Function

Maximize the total price of the chosen objects:

$$\text{maximize } \sum_{i=1}^n p_i \cdot \text{Chosen}_i$$

Max Weight Constraint

Ensure the total weight of the chosen objects does not exceed the suitcase's capacity:

$$\sum_{i=1}^n w_i \cdot \text{Chosen}_i \leq c$$

Coordinate Bounds Constraints

Ensure each object lies entirely within the suitcase's boundaries:

$$\forall i \in \{1, \dots, n\}, \quad \text{Points}X_i \geq 1$$

$$\forall i \in \{1, \dots, n\}, \quad \text{Points}Y_i \geq 1$$

$$\forall i \in \{1, \dots, n\}, \quad \text{Points}X_i + s_i - 1 \leq x$$

$$\forall i \in \{1, \dots, n\}, \quad \text{Points}Y_i + s_i - 1 \leq y$$

Non-Overlapping Constraints

Left/Right/Up/Down non-overlapping:

$$\forall i, j \in \{1, \dots, n\}, i \neq j, \quad \text{Points}X_i - \text{Points}X_j + s_i \leq \\ -M \cdot (\text{Chosen}_i + \text{Chosen}_j + \text{Overlap}_{i,j,1} - 3)$$

$$\forall i, j \in \{1, \dots, n\}, i \neq j, \quad \text{Points}X_i - \text{Points}X_j + s_i \leq \\ -M \cdot (\text{Chosen}_i + \text{Chosen}_j + \text{Overlap}_{i,j,2} - 3)$$

$$\forall i, j \in \{1, \dots, n\}, i \neq j, \quad \text{Points}X_i - \text{Points}X_j + s_i \leq \\ -M \cdot (\text{Chosen}_i + \text{Chosen}_j + \text{Overlap}_{i,j,3} - 3)$$

$$\forall i, j \in \{1, \dots, n\}, i \neq j, \quad \text{Points}X_i - \text{Points}X_j + s_i \leq \\ -M \cdot (\text{Chosen}_i + \text{Chosen}_j + \text{Overlap}_{i,j,4} - 3)$$

Non-Overlapping Constraints

At least one of the non-overlapping conditions is satisfied:

$$\forall i, j \in \{1, \dots, n\}, i \neq j, \quad \sum_{d=1}^4 \text{Overlap}_{i,j,d} \geq 1$$

Greedy heuristic

Input: A problem instance with n products, suitcase dimensions (x, y) , and weight capacity c .

Output: A feasible solution with selected products that maximize the total price.

Initialize solution as empty

Sort products by $(price/side/weight)$ in descending order

for each product in sorted products **do**

if product can fit in the suitcase and does not exceed weight limit **then**

 Add product to solution

 Update suitcase dimensions and weight capacity

end if

end for

return solution

Local search

Input: An initial solution, mode and strategy.

Output: An improved solution

bestSolution = initialSolution

improved = true

while improved **do**

 improved = false

for each product in bestSolution **do**

if strategy == EXCHANGE **then**

 newSolution = exchangeProduct(bestSolution, product)

else if strategy == REASSIGNMENT **then**

 newSolution = reassignProduct(bestSolution, product)

end if

if newSolution is better than bestSolution **then**

 bestSolution = newSolution

 improved = true

if mode == FIRST_IMPROVEMENT **then**

return bestSolution

end if

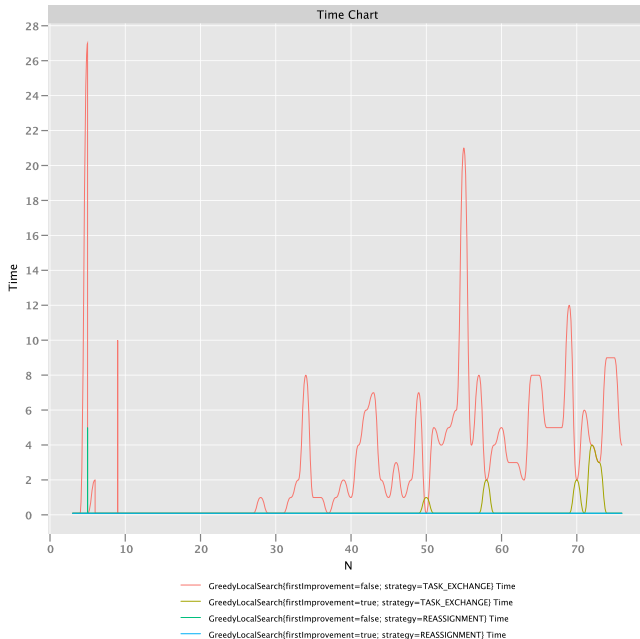
end if

end for

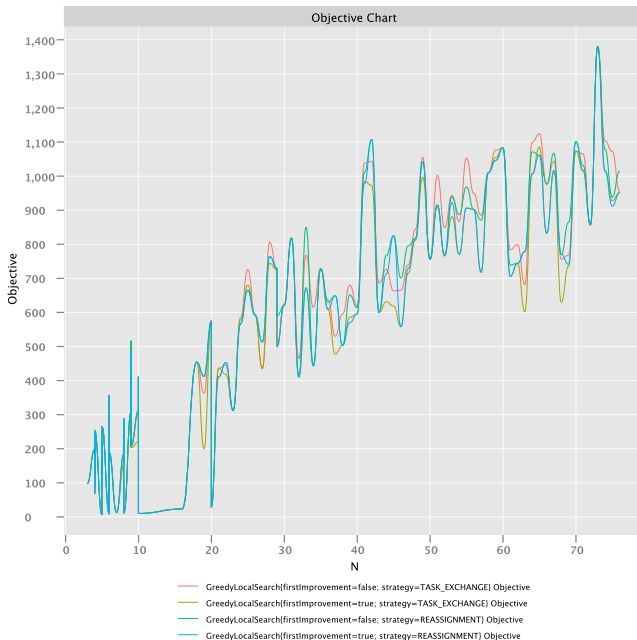
end while

return bestSolution

Local Search Time chart



Local Search Objective chart



GRASP

Input: A problem instance, maxIterations, alpha (for RCL threshold)

Output: The best solution found

bestSolution = null

for iteration = 1 to maxIterations **do**

 greedySolution = constructGreedyRandomizedSolution(problem, alpha)

 localOptimalSolution = LOCAL_SEARCH(greedySolution,

FIRST_IMPROVEMENT)

if bestSolution is null or localOptimalSolution is better than bestSolution

then

 bestSolution = localOptimalSolution

end if

end for

return bestSolution

GRASP (2)

Input: A problem instance, α (for RCL threshold)

Output: A feasible solution

Initialize solution as empty

Sort products by ($price/side/weight$) in descending order

while there are remaining products **do**

$qMax$ = maximum Q value in remaining products

$qMin$ = minimum Q value in remaining products

 threshold = $qMax - \alpha \times (qMax - qMin)$

 RCL = {products with Q value \geq threshold}

 selectedProduct = randomly select a product from RCL

if selectedProduct fits in the suitcase and does not exceed weight limit

then

 Add selectedProduct to solution

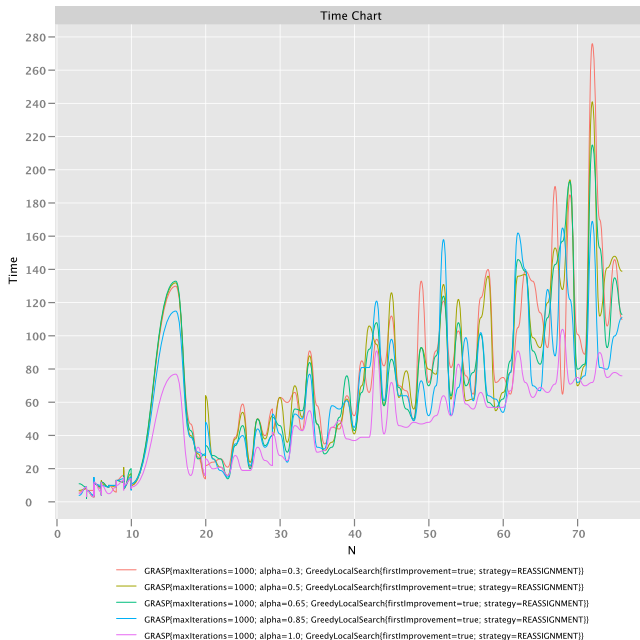
 Update suitcase dimensions and weight capacity

end if

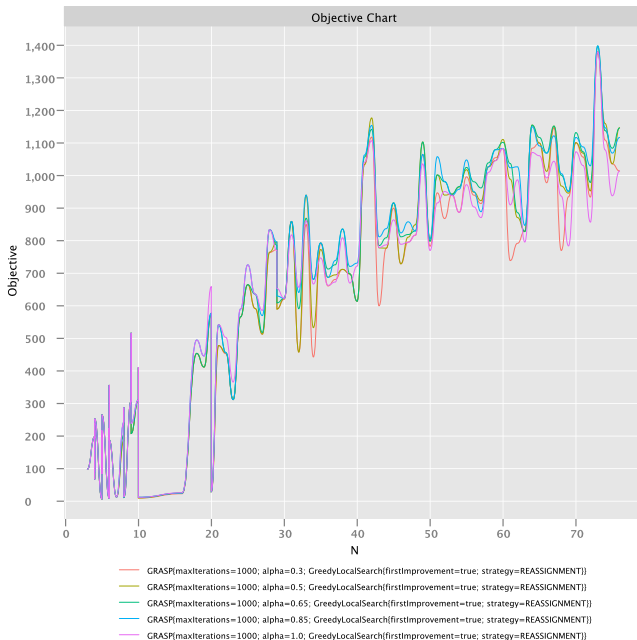
end while

return solution

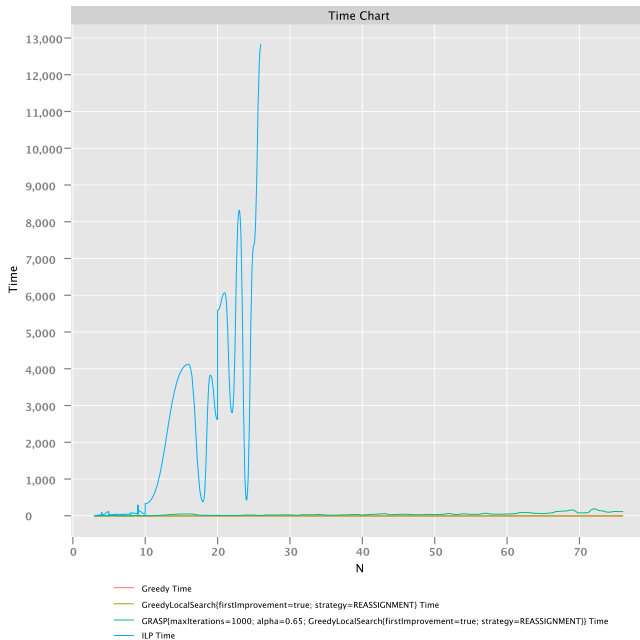
GRASP Time chart



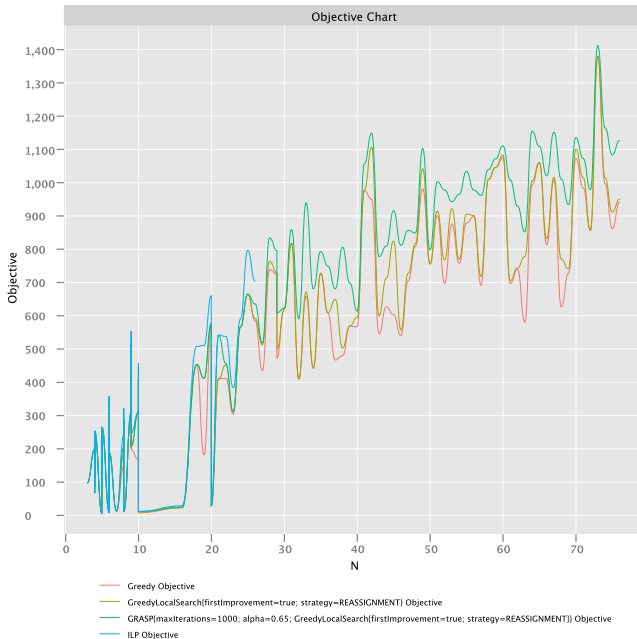
GRASP Objective chart



All Time chart



All Objective chart



Improvements

- ▶ Parallelize algorithm
- ▶ Better strategies
- ▶ Implement the state of the art for heuristic