Algorithmic Methods for Mathematical Models Course Project

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1 Formal Problem Definition

1.1 Inputs

- \bullet n: number of products
- x: height of the suitcase in millimeters
- y: width of the suitcase in millimeters
- c: limit to the total weight of the suitcase in grams
- p_i : price of the *i*-th product in euros
- w_i : weight of the *i*-th product in grams
- s_i : side length of the *i*-th product's (square) box in millimeters

1.2 Outputs

- Indices of the products chosen to maximize the accumulated price
- Arrangement of the products in the suitcase

2 Mathematical Formulation

2.1 Decision Variables

- Chosen,: binary variable that is 1 if object i is chosen, and 0 otherwise.
- Points X_i : the x-coordinate of the bottom-left corner of object i.
- Points Y_i : the y-coordinate of the bottom-left corner of object i.
- Overlap_{i,j,d}: binary variable indicating if objects i and j do not overlap in direction d, where $d \in \{1,2,3,4\}$.

2.2 Objective Function

Maximize the total price of the chosen objects:

$$\text{maximize } \sum_{i=1}^{n} p_i \cdot \text{Chosen}_i$$

2.3 Constraints

2.3.1 Max Weight Constraint

Ensure the total weight of the chosen objects does not exceed the suitcase's capacity:

$$\sum_{i=1}^{n} w_i \cdot \text{Chosen}_i \le c$$

2.3.2 Coordinate Bounds Constraints

Ensure each object lies entirely within the suitcase's boundaries:

$$\forall i \in \{1, \dots, n\}, \quad \text{Points} \mathbf{X}_i \geq 1$$

$$\forall i \in \{1, \dots, n\}, \quad \text{Points} \mathbf{Y}_i \geq 1$$

$$\forall i \in \{1, \dots, n\}, \quad \text{Points} \mathbf{X}_i + s_i - 1 \leq x$$

$$\forall i \in \{1, \dots, n\}, \quad \text{Points} \mathbf{Y}_i + s_i - 1 \leq y$$

2.3.3 Non-Overlapping Constraints

Ensure no two chosen objects overlap within the suitcase using the big-M method:

1. Horizontal non-overlapping to the right:

$$\forall i, j \in \{1, \dots, n\}, i \neq j, \quad \text{PointsX}_i - \text{PointsX}_j + s_i \leq -M \cdot (\text{Chosen}_i + \text{Chosen}_j + \text{Overlap}_{i,j,1} - 3)$$

2. Horizontal non-overlapping to the left:

$$\forall i,j \in \{1,\ldots,n\}, i \neq j, \quad \text{PointsX}_j - \text{PointsX}_i + s_j \leq -M \cdot (\text{Chosen}_i + \text{Chosen}_j + \text{Overlap}_{i,j,2} - 3)$$

3. Vertical non-overlapping upwards:

$$\forall i,j \in \{1,\ldots,n\}, i \neq j, \quad \text{PointsY}_i - \text{PointsY}_j + s_i \leq -M \cdot (\text{Chosen}_i + \text{Chosen}_j + \text{Overlap}_{i,j,3} - 3)$$

4. Vertical non-overlapping downwards:

$$\forall i, j \in \{1, \dots, n\}, i \neq j$$
, Points Y_j - Points $Y_i + s_j \leq -M \cdot (\text{Chosen}_i + \text{Chosen}_j + \text{Overlap}_{i,j,4} - 3)$

2.3.4 At Least One Not Overlapping Constraint

Ensure that for any two objects, at least one of the non-overlapping conditions is satisfied:

$$\forall i, j \in \{1, \dots, n\}, i \neq j, \quad \sum_{d=1}^{4} \text{Overlap}_{i,j,d} \geq 1$$

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3 Graphs

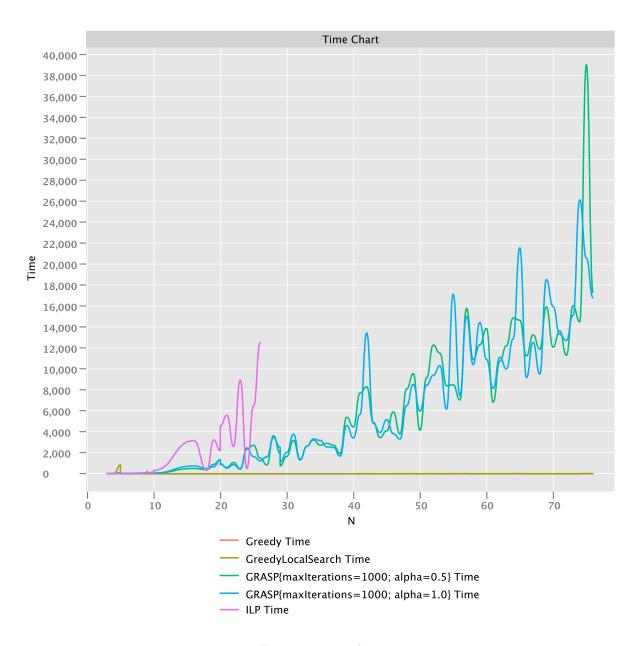


Figure 1: Time chart

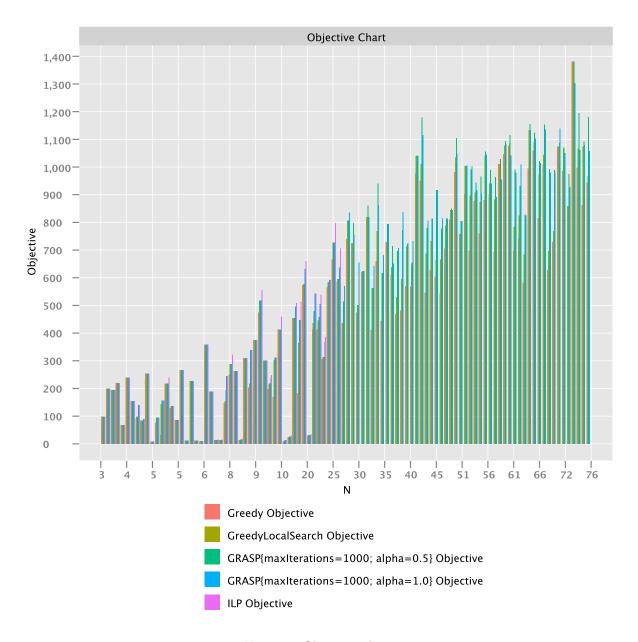


Figure 2: Objective chart