

# Exercise Session 24/05/2023

## Algorithms and Data Structures

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### Exercise 179 (f16)

Write an algorithm `BTreePrintRange(T, a, b)` that, given a B-tree  $T$  and two values  $a < b$ , prints all the keys  $k$  in  $T$  that are between  $a$  and  $b$ , that is,  $a < k < b$ . Recall that a node  $x$  in a B-tree has the following properties:

- $x.n$  is the number of keys
- $x.key[1] \leq x.key[2] \leq \dots \leq x.key[x.n]$  are the keys
- $x.leaf$  tells whether  $x$  is a leaf
- $x.c[1], x.c[2], \dots, x.c[x.n + 1]$  are the pointers to  $x$ 's children

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### Exercise 161 (f15)

A breadth-first search over a graph  $G$  returns a vector  $\pi$  that represents the resulting breadth-first tree, where the parent  $\pi[v]$  of a vertex  $v$  is the next-hop from  $v$  on the tree towards the source of the breadth-first search.

Question 1: Write an algorithm `BFS-First-Common-Ancestor( $\pi, u, v$ )` that finds the first common ancestor of two given nodes in the breadth-first tree, or null if  $u$  and  $v$  are not connected in  $G$ . The complexity of `BFS-First-Common-Ancestor` must be  $O(n)$ . Briefly analyze the space complexity of your solution

Question 2: Write an algorithm `BFS-First-Common-Ancestor-2( $\pi, D, u, v$ )` that is also given the distance vector  $D$  resulting from the same breadth first search. `BFS-First-Common-Ancestor-2` must be functionally equivalent to `BFS-First-Common-Ancestor` (as defined in Exercise 1) but with space complexity  $O(1)$ .