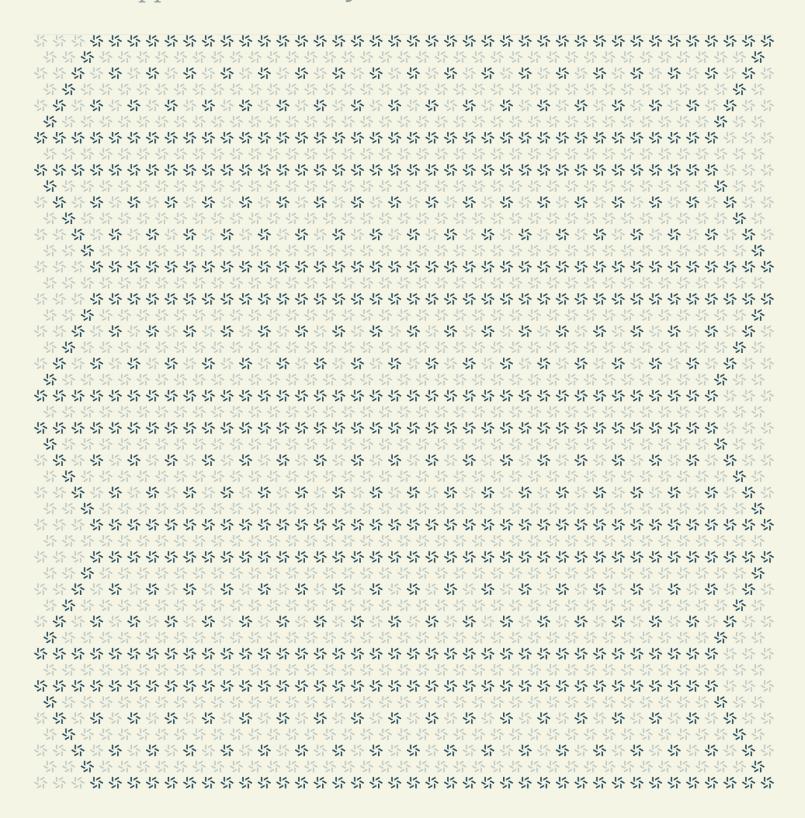


April 9, 2025

N1 Bridge

Solana Application Security Assessment





Contents

Abo	ut Zelli		4
1.	Over	view	4
	1.1.	Executive Summary	5
	1.2.	Goals of the Assessment	5
	1.3.	Non-goals and Limitations	5
	1.4.	Results	5
2.	Intro	duction	6
	2.1.	About N1 Bridge	7
	2.2.	Methodology	7
	2.3.	Scope	9
	2.4.	Project Overview	9
	2.5.	Project Timeline	10
3.	Detailed Findings		10
	3.1.	Insufficient block validation	11
	3.2.	Associated Token Account check not enforced on withdrawal account	14
	3.3.	Usage of Pubkey rather than Mint	16
4.	Disc	ussion	17
	4.1.	Block validation	18
	4.2.	Unused accounts	18
	4.3.	Test suite	18



5.	Threat Model		
	5.1. Program: bridge	19	
6.	Assessment Results	31	
	61 Disclaimer	32	



About Zellic

Zellic is a vulnerability research firm with deep expertise in blockchain security. We specialize in EVM, Move (Aptos and Sui), and Solana as well as Cairo, NEAR, and Cosmos. We review L1s and L2s, cross-chain protocols, wallets and applied cryptography, zero-knowledge circuits, web applications, and more.

Prior to Zellic, we founded the #1 CTF (competitive hacking) team > worldwide in 2020, 2021, and 2023. Our engineers bring a rich set of skills and backgrounds, including cryptography, web security, mobile security, low-level exploitation, and finance. Our background in traditional information security and competitive hacking has enabled us to consistently discover hidden vulnerabilities and develop novel security research, earning us the reputation as the go-to security firm for teams whose rate of innovation outpaces the existing security landscape.

For more on Zellic's ongoing security research initiatives, check out our website $\underline{\text{zellic.io}} \, \underline{\text{z}}$ and follow @zellic_io $\underline{\text{z}}$ on Twitter. If you are interested in partnering with Zellic, contact us at hello@zellic.io $\underline{\text{z}}$.



Zellic © 2025 ← Back to Contents Page 4 of 32



Overview

1.1. Executive Summary

Zellic conducted a security assessment for Layer N from March 20th, 2025 to March 25th, 2025. During this engagement, Zellic reviewed N1 Bridge's code for security vulnerabilities, design issues, and general weaknesses in security posture.

1.2. Goals of the Assessment

In a security assessment, goals are framed in terms of questions that we wish to answer. These questions are agreed upon through close communication between Zellic and the client. In this assessment, we sought to answer the following questions:

- Are there any vulnerabilities that could result in the loss of user funds?
- · Are access controls implemented correctly to prevent unauthorized operations?
- · Are there potential issues with the validation for proposed blocks?
- Does the custom Merkle implementation contain any bugs?

1.3. Non-goals and Limitations

We did not assess the following areas that were outside the scope of this engagement:

- · Front-end components
- · Infrastructure relating to the project
- · Key custody

Due to the time-boxed nature of security assessments in general, there are limitations in the coverage an assessment can provide.

1.4. Results

During our assessment on the scoped N1 Bridge programs, we discovered three findings. No critical issues were found. One finding was of medium impact and the other findings were informational in nature.

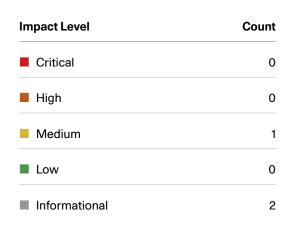
Additionally, Zellic recorded its notes and observations from the assessment for the benefit of Layer N in the Discussion section $(4, \pi)$.

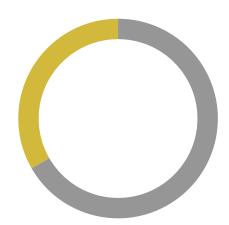
Zellic © 2025

← Back to Contents Page 5 of 32



Breakdown of Finding Impacts







2. Introduction

2.1. About N1 Bridge

Layer N contributed the following description of N1 Bridge:

A layer 1 blockchain designed for unlimited scale, featuring horizontal scalability, sub-ms latency, and congestion-free throughput.

2.2. Methodology

During a security assessment, Zellic works through standard phases of security auditing, including both automated testing and manual review. These processes can vary significantly per engagement, but the majority of the time is spent on a thorough manual review of the entire scope.

Alongside a variety of tools and analyzers used on an as-needed basis, Zellic focuses primarily on the following classes of security and reliability issues:

Basic coding mistakes. Many critical vulnerabilities in the past have been caused by simple, surface-level mistakes that could have easily been caught ahead of time by code review. Depending on the engagement, we may also employ sophisticated analyzers such as model checkers, theorem provers, fuzzers, and so on as necessary. We also perform a cursory review of the code to familiarize ourselves with the programs.

Business logic errors. Business logic is the heart of any smart contract application. We examine the specifications and designs for inconsistencies, flaws, and weaknesses that create opportunities for abuse. For example, these include problems like unrealistic tokenomics or dangerous arbitrage opportunities. To the best of our abilities, time permitting, we also review the contract logic to ensure that the code implements the expected functionality as specified in the platform's design documents.

Integration risks. Several well-known exploits have not been the result of any bug within the contract itself; rather, they are an unintended consequence of the contract's interaction with the broader DeFi ecosystem. Time permitting, we review external interactions and summarize the associated risks: for example, flash loan attacks, oracle price manipulation, MEV/sandwich attacks, and so on.

Code maturity. We look for potential improvements in the codebase in general. We look for violations of industry best practices and guidelines and code quality standards. We also provide suggestions for possible optimizations, such as gas optimization, upgradability weaknesses, centralization risks, and so on.

For each finding, Zellic assigns it an impact rating based on its severity and likelihood. There is no hard-and-fast formula for calculating a finding's impact. Instead, we assign it on a case-by-case basis based on our judgment and experience. Both the severity and likelihood of an issue affect its impact. For instance, a highly severe issue's impact may be attenuated by a low likelihood.

Zellic © 2025 ← Back to Contents Page 7 of 32



We assign the following impact ratings (ordered by importance): Critical, High, Medium, Low, and Informational.

Zellic organizes its reports such that the most important findings come first in the document, rather than being strictly ordered on impact alone. Thus, we may sometimes emphasize an "Informational" finding higher than a "Low" finding. The key distinction is that although certain findings may have the same impact rating, their *importance* may differ. This varies based on various soft factors, like our clients' threat models, their business needs, and so on. We aim to provide useful and actionable advice to our partners considering their long-term goals, rather than a simple list of security issues at present.

Finally, Zellic provides a list of miscellaneous observations that do not have security impact or are not directly related to the scoped programs itself. These observations — found in the Discussion $(\underline{4}, \pi)$ section of the document — may include suggestions for improving the codebase, or general recommendations, but do not necessarily convey that we suggest a code change.



2.3. Scope

The engagement involved a review of the following targets:

N1 Bridge Programs

Туре	Rust
Platform	Solana
Target	nord
Repository	https://github.com/n1xyz/nord >
Version	ed12f5dadaa555f401e7122085b8fecef90eddec
Programs	bridge

2.4. Project Overview

Zellic was contracted to perform a security assessment for a total of four person-days. The assessment was conducted by two consultants over the course of two calendar weeks.

Zellic © 2025 \leftarrow Back to Contents Page 9 of 32



Contact Information

The following project managers were associated with the engagement:

The following consultants were engaged to conduct the assessment:

Jacob Goreski

Frank Bachman

☆ Engineer frank@zellic.io
オ

Chad McDonald

片 Engagement Manager chad@zellic.io 제

Bryce Casaje

☆ Engineer

bryce@zellic.io

¬

2.5. Project Timeline

The key dates of the engagement are detailed below.

March 20, 2025	Kick-off call
March 20, 2025	Start of primary review period
March 25, 2025	End of primary review period

Zellic © 2025

← Back to Contents Page 10 of 32



3. Detailed Findings

3.1. Insufficient block validation

Target	bridge			
Category	Coding Mistakes	Severity	High	
Likelihood	Low	Impact	Medium	

Description

A block consists of the following fields:

```
pub struct Block {
    pub facts: BlockFacts,
    pub finalized: bool,
    pub slot_proposed: u64,
}

pub struct BlockFacts {
    pub prev_state_facts: StateFacts,
    pub next_state_facts: StateFacts,
    pub da_commitment: [u8; 32],
    pub withdrawal_root: [u8; 32],
}

pub struct StateFacts {
    pub app_state_commitment: [u8; 32],
    pub deposit_root: [u8; 32],
    pub last_deposit_index: u64,
    pub last_action_id: u64,
}
```

The propose_block instruction takes a BlockFacts struct as an argument to propose a new block to the bridge program, and it can only be run by the account set as the operator.

The operator is meant to be untrusted, so the implementation performs some validation on the block to assert that the block is correct.

It currently validates the following:

- 1. The block is proposed by the operator.
- 2. The data-availability (DA) fact is finalized.
- 3. The last deposit (if supplied) hashes to the provided deposit root in the next state facts.

Zellic © 2025 ← Back to Contents Page 11 of 32



And when a block is finalized via the finalize_block instruction, it performs some additional validation:

- 1. The previous state facts should match the current state facts.
- 2. The block has not been previously finalized.
- 3. The block has existed for enough time (based on challenge_period_slots).

However, the last_deposit account in the propose_block instruction does not have sufficient validation:

```
#[account(
    seeds = [DEPOSIT_SEED,
    &facts.next_state_facts.last_deposit_index.to_le_bytes()],
    bump,
)]
pub last_deposit: Option<Account<'info, Deposit>>,
```

Even if a last deposit does exist, it can still be set to None since the account is wrapped in an Option.

This messes up the validation for last_deposit_index and deposit_root for the next state facts:

```
match (
    facts.next_state_facts.last_deposit_index,
    facts.next_state_facts.deposit_root,
    &ctx.accounts.last_deposit,
) {
    (0, root, None) if root == [0; 32] => (),
    (_, root, Some(d)) if root == d.hash() => (),
    _ => return err!(BridgeError::InvalidDepositRoot),
};
```

A block can be proposed with a last_deposit_index set to zero and an empty deposit_root if last_deposit is set to None, even if a last deposit does exist.

In addition, there is no validation for the withdrawal_root field as well as the app_state_commitment, last_action_id, and last_deposit_index fields in the next state facts.

Impact

Invalid values for app_state_commitment, last_action_id, last_deposit_index, and deposit_root can cause errors in off-chain bridge infrastructure that trusts these values to be correct.

An invalid Merkle root for withdrawal_root would allow an attacker to verify any withdrawal proof

Zellic © 2025 ← Back to Contents Page 12 of 32



and drain funds from the bridge. However, this would require a compromised bridge operator to propose and finalize a malicious block.

Recommendations

Implement validation for fields that can be validated on chain.

For example, the last_deposit account should never be None if there has been a previous deposit.

Remediation

The issue has been acknowledged by Layer N, and the team intends for most block validation to happen off chain via manual intervention or a challenge mechanism, as mentioned in Discussion point 4.1. π .

Some extra checks were added in 716bb92c 7, and the following comment was provided:

A block is allowed to be proposed covering a smaller range of deposits than available at proposal time.

Off-chain bridge infrastructure that should NOT trust these values to be correct as they belong to an unfinalized block.

Furthermore, an invalid Merkle root for withdrawal_root would NOT allow an attacker to verify any withdrawal proof and drain funds from the bridge as that would require the block to be finalized.

While finalization is intended for the operator, there is no trust assumption on the signer selecting valid blocks to finalize. There is however a (planned) monetary incentive on the proposer. The validation of a block prior to finalization is done by the contract through proposal and finalization-time checks AND our fraud-proof challenge mechanism (planned) / manual review during longer challenge-period + manual intervention (interim).

Zellic © 2025 ← Back to Contents Page 13 of 32



3.2. Associated Token Account check not enforced on withdrawal account

Target	bridge			
Category	Coding Mistakes	Severity	Informational	
Likelihood	Low	Impact	Informational	

Description

It was found that in Withdraw, the $to_account$ is not enforced to be an Associated Token Account (ATA).

```
#[derive(Accounts)]
#[instruction(claim: WithdrawalClaim)]
pub struct Withdraw<'info> {
[...]
    #[account(mut, constraint = to_account.owner == claim.user)]
    pub to_account: InterfaceAccount<'info, TokenAccount>,
```

The Solana Token Program allows users to arbitrarily create many token accounts belonging to the same mint. Note that anyone can create these accounts for a specified owner.

It is therefore possible for a caller to create a separate account (still controlled by the owner) to withdraw the funds.

Impact

The caller can withdraw funds to a non-ATA token account. This is not a security issue as the token account would still be controlled and managed by the owner.

Recommendations

Enforce to_account in Withdraw to be an ATA.

Remediation

Layer N acknowledged the issue, and decided not to apply a remediation with the following explanation:



This is intended behaviour as a flexibility of the api that we get "for free" by not enforcing ATA. The signer of the withdraw transaction is responsible for the destination account and this should be harmless as the owner is validated.



3.3. Usage of Pubkey rather than Mint

Target	bridge			
Category	Coding Mistakes	Severity	Informational	
Likelihood	N/A	Impact	Informational	

Description

In the whitelist_asset and set_min_deposit instructions, the asset to configure is supplied by the operator as a Pubkey in the instruction arguments:

```
#[instruction(asset: Pubkey, min_deposit: u64)]
pub struct SetMinDeposit<'info> {
    // ...
}
```

The asset is intended to be a Solana Program Library (SPL) token mint, but this requirement is not set by the implementation.

Impact

The operator can whitelist an asset that is not an SPL token mint.

This does not cause any direct issues since the deposit_spl instruction derives the asset from the mint of a TokenAccount, but this should be fixed to prevent future security regressions.

Recommendations

The asset to configure should not be an instruction argument but instead a Mint account from the anchor-spl crate.

Zellic © 2025 \leftarrow Back to Contents Page 16 of 32



Remediation

This issue has been acknowledged by Layer N, and a fix was implemented in commit $\underline{ed392f7c} \overline{\nearrow}$.



Discussion

The purpose of this section is to document miscellaneous observations that we made during the assessment. These discussion notes are not necessarily security related and do not convey that we are suggesting a code change.

4.1. Block validation

The provided version of the code has no implementation to challenge proposed blocks or manually intervene and stop malicious blocks from being finalized.

However, Layer N intends to implement these features in future editions of the code.

As such, the propose_block and finalize_block instructions do not need to completely validate all fields of proposed blocks. Nevertheless, the fields that can be validated on chain should be validated to help prevent malicious blocks.

4.2. Unused accounts

The finalize_block instruction includes a redundant payer account that

- 1. incurs no fees (since there is no account creation in the instruction), and
- 2. serves no access-control purpose (since the instruction is permissionless).

As such, the payer account can be removed to simplify the interface and reduce potential confusion about its purpose.

This issue has been acknowledged by Layer N, and a fix was implemented in commit 72f1cbed 7.

4.3. Test suite

The test suite provided covered most of the instructions implemented by the program.

However, some improvements could still be made, especially in regards to edge cases and negative scenarios, as the test suite mostly focuses on verifying that each instruction works as expected with the correct signers and arguments.

In addition, none of the tests in the test suite simulate that the program functions correctly with more than one block.

Zellic © 2025 ← Back to Contents Page 18 of 32



Threat Model

As time permitted, we analyzed each instruction in the program and created a written threat model for the most critical instructions. A threat model documents the high-level functionality of a given instruction, the inputs it receives, and the accounts it operates on as well as the main checks performed on them; it gives an overview of the attack surface of the programs and of the level of control an attacker has over the inputs of critical instructions.

For brevity, system accounts and well-known program accounts have not been included in the list of accounts received by an instruction; the instructions that receive these accounts make use of Anchor types, which automatically ensure that the public key of the account is correct.

Discriminant checks, ownership checks, and rent checks are not discussed for each individual account; unless otherwise stated, the program uses Anchor types, which perform the necessary checks automatically.

Not all instructions in the audit scope may have been modeled. The absence of a threat model in this section does not necessarily suggest that an instruction is safe.

5.1. Program: bridge

Instruction: deposit_spl

This instruction allows the depositor to deposit SPL tokens to the bridge.

Input parameters

• amount: u64: The amount of SPL tokens to deposit.

Accounts

- depositor: The depositor account that is the authority for the deposited tokens pays for the new initialized accounts.
 - · Signer: Yes.
 - · Init: No.
 - PDA: No.
 - · Mutable: Yes.
 - Constraints: Must have sufficient lamports to cover the account-creation cost.
- deposit: The account to hold the new deposit data.
 - · Signer: No.
 - · Init: Yes.
 - PDA: Yes (derived from DEPOSIT_SEED and contract_storage.last_deposit_index + 1 with bump).

Zellic © 2025 ← Back to Contents Page 19 of 32



- · Mutable: Yes.
- Constraints: Must be a new account initialized with space 8 + 104 (discriminator + Deposit size).
- prev_deposit: The previous deposit data.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from DEPOSIT_SEED and contract_storage.last_deposit_index with bump).
 - · Mutable: No.
 - · Optional: Yes.
 - Constraints: Must be of type Deposit.
- asset_config: The configuration data for this asset.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from ASSET_CONFIG_SEED and from_account.mint with bump).
 - · Mutable: No.
 - Constraints: Must be an AssetConfig where min_deposit is less than or equal to amount.
- contract_storage: The global contract storage.
 - Signer: No.
 - Init: No.
 - PDA: Yes (derived from CONTRACT_STORAGE_SEED with bump).
 - · Mutable: Yes.
 - Constraints: Must be a ContractStorage account. If last_deposit_index is not 0, a prev_deposit must be provided.
- from_account: The SPL token account where tokens are taken from.
 - · Signer: No.
 - Init: No.
 - · PDA: No.
 - · Mutable: Yes.
 - · Constraints: Must be a valid SPL token account.
- to_account: The ATA where tokens will be sent.
 - · Signer: No.
 - Init: No.
 - · PDA: No.
 - · Mutable: Yes.
 - Constraints: Must be an ATA for the authority account, with the same mint as from_account.

Zellic © 2025 \leftarrow Back to Contents Page 20 of 32



- Transfers amount of the SPL token from from_account to to_account with the authority depositor.
- Creates and stores a Deposit where transfer holds the deposit information and prev_deposit_root is the hash of the previous deposit, if supplied.
- Increments contract_storage.last_deposit_index.
- Logs a Deposit event.

CPI

 transfer: Transfers the specified tokens from from_account to to_account, signed by depositor.

Instruction: finalize_block

This instruction finalizes a proposed block.

Input parameters

• block_id: u64: The ID for the block to finalize.

Accounts

- payer: The signer for this transaction.
 - · Signer: Yes.
 - Init: No.
 - PDA: No.
 - · Mutable: Yes.
 - · Constraints: None.
- block: The block to be finalized.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from BLOCK_STORAGE_SEED and block_id with bump).
 - · Mutable: Yes.
 - Constraints: The block's prev_state_facts must match contract_storage.fina_state_facts, the block must not be finalized, and the current time slot should be greater than block.slot_proposed + contract_storage.challenge_period_slots.

Zellic © 2025 ← Back to Contents Page 21 of 32



- **contract_storage**: The global contract storage.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from CONTRACT_STORAGE_SEED with bump).
 - · Mutable: Yes.
 - Constraints: Must be a ContractStorage account.

- Sets block.finalized to true.
- Sets contract_storage.fina_block_id to block_id.
- Sets contract_storage.fina_state_facts to block.facts.next_state_facts.
- Logs a FinalizeBlock event.

CPI

N/A.

Instruction: deposit_spl

This instruction finalizes a DA fact.

Input parameters

• fact: [u8; 32]: The DA fact to finalize.

Accounts

- payer: The signer for this transaction who pays account-creation fees.
 - · Signer: Yes.
 - Init: No.
 - PDA: No.
 - Mutable: Yes.
 - Constraints: Must have sufficient lamports to cover the account-creation cost.
- ${\bf fact_state_storage} : The \ account \ to \ store \ the \ DA-fact \ finalization.$
 - · Signer: No.
 - Init: Yes.

Zellic © 2025 \leftarrow Back to Contents Page 22 of 32



- PDA: Yes (derived from DA_FACT_STORAGE_SEED and fact with bump).
- · Mutable: Yes.
- Constraints: Must be a new account initialized with space 8 + 1 (discriminator + FactStateStorage size).

- Sets fact_state_storage to a FactStateStorage struct with the state FactState::Finalized.
- Logs a FinalizeDaFact event.

CPI

N/A.

Instruction: initialize

This instruction initializes the global contract storage.

Input parameters

- operator: Pubkey: The public key for the bridge operator role.
- challenge_period_slots: u64: The number of slots where a block can be challenged before it is finalized.
- app_state_commitment: [u8; 32]: The initial app state commitment value.

Accounts

- payer: The signer for this transaction who pays account-creation fees.
 - · Signer: Yes.
 - Init: No.
 - PDA: No.
 - Mutable: Yes.
 - Constraints: Must have sufficient lamports to cover the account-creation cost.
- program: The bridge program.
 - · Signer: Yes.
 - Init: No.
 - PDA: No.

Zellic © 2025

← Back to Contents Page 23 of 32



- Mutable: No.
- Constraints: Must be an account with the address crate::ID (which should be the bridge program).
- contract_storage: The global contract storage.
 - · Signer: No.
 - · Init: Yes.
 - PDA: Yes (derived from CONTRACT_STORAGE_SEED with bump).
 - · Mutable: Yes.
 - Constraints: Must be a new account initialized with space 8 + 144 (discriminator + ContractStorage size).

- Checks that crate::ID is not [0; 32].
- Creates and stores a ContractStorage account with the values from the input parameters.

CPI

N/A.

Instruction: propose_block

This instruction allows the operator to propose a new block to the bridge.

Input parameters

• facts: BlockFacts: Information about the new block to be proposed.

Accounts

- operator: The operator account who pays for the new initialized accounts.
 - · Signer: Yes.
 - Init: No.
 - PDA: No.
 - · Mutable: Yes.
 - Constraints: Must have the same public key as contract_storage.operator and have sufficient lamports to cover the account-creation cost.
- block: The account to hold the new proposed block data.

Zellic © 2025 \leftarrow Back to Contents Page 24 of 32



- · Signer: No.
- · Init: Yes.
- PDA: Yes (derived from BLOCK_STORAGE_SEED and contract_storage.last_block_id + 1 with bump).
- · Mutable: Yes.
- Constraints: Must be a new account initialized with space 8 + 233 (discriminator + Block size)
- last_deposit: The previous deposit data.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from DEPOSIT_SEED and facts.next_state_facts.last_deposit_index with bump).
 - · Mutable: No.
 - · Optional: Yes.
 - · Constraints: Must be of type Deposit.
- da_fact_state: The DA commitment account data.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from DA_FACT_STORAGE_SEED and facts.da_commitment with bump).
 - · Mutable: No.
 - Constraints: Must be FactStateStorage where da_fact_state.state is FactState::Finalized.
- contract_storage: The global contract storage.
 - · Signer: No.
 - · Init: No.
 - PDA: Yes (derived from CONTRACT_STORAGE_SEED with bump).
 - · Mutable: Yes.
 - Constraints: Must be a ContractStorage account.

- Checks that facts.next_state_facts.deposit_root is valid by ensuring it is either all zeros when facts.next_state_facts.last_deposit_index is 0 and last_deposit is None or matches the hash of last_deposit when a deposit exists.
- Creates and stores a Block with the facts from the input parameter.
- Increments contract_storage.last_block_id.
- Logs a ProposeBlock event.

Zellic © 2025 \leftarrow Back to Contents Page 25 of 32



CPI

N/A.

Instruction: set_min_deposit

This instruction allows the operator to set the minimum deposit for an asset.

Input parameters

- asset: Pubkey: The public key for the asset to configure.
- min_deposit: u64: The minimum deposit required for this asset.

Accounts

- operator: The operator account that is setting the minimum deposit.
 - · Signer: Yes.
 - Init: No.
 - PDA: No.
 - Mutable: No.
 - Constraints: Must have the same public key as contract_storage.operator.
- contract_storage: The global contract storage.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from CONTRACT_STORAGE_SEED with bump).
 - Mutable: No.
 - Constraints: Must be a ContractStorage account.
- asset_config: The asset-configuration account matching the public key.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from ASSET_CONFIG_SEED and asset with bump).
 - Mutable: Yes.
 - Constraints: The min_deposit input parameter must be greater than zero.

Additional checks and behavior

- Sets asset_config.min_deposit to the min_deposit from the input parameters.
- Logs a SetMinDeposit event.

Zellic © 2025 ← Back to Contents Page 26 of 32



CPI

N/A.

Instruction: whitelist asset

This instruction allows the operator to whitelist an asset for depositing and withdrawing.

Input parameters

- asset: Pubkey: The public key for the asset to configure.
- min_deposit: u64: The minimum deposit required for this asset.

Accounts

- operator: The operator account that is whitelisting the asset.
 - · Signer: Yes.
 - Init: No.
 - PDA: No.
 - Mutable: No.
 - Constraints: Must have the same public key as contract_storage.operator and enough lamports to pay for account-creation costs.
- $\bullet \ \ contract_storage: The \ global \ contract \ storage.$
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from CONTRACT_STORAGE_SEED with bump).
 - · Mutable: No.
 - Constraints: Must be a ContractStorage account.
- ${\tt asset_config:}$ The asset-configuration account matching the public key.
 - · Signer: No.
 - Init: Yes.
 - PDA: Yes (derived from ASSET_CONFIG_SEED and asset with bump).
 - · Mutable: Yes.
 - Constraints: Must be a new account initialized with space 8 + 8 (discriminator + AssetConfig size), and the min_deposit input parameter must be greater than zero.

Zellic © 2025 \leftarrow Back to Contents Page 27 of 32



- $\bullet \ \ Creates \ and \ stores \ a \ new \ AssetConfig \ with \ the \ min_deposit \ from \ the \ input \ parameters.$
- Logs a WhitelistAsset event.

CPI

N/A.

Instruction: withdraw

This instruction allows a user to withdraw tokens from the bridge.

Input parameters

Accounts

- payer: The account that will pay for the creation of the nullifier.
 - · Signer: Yes.
 - Init: No.
 - PDA: No.
 - · Mutable: Yes.
 - Constraints: Must have enough lamports to pay for account-creation costs.
- state_update: The finalized block that contains the withdrawal event.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from BLOCK_STORAGE_SEED and claim.block_id with bump).

Zellic © 2025 ← Back to Contents Page 28 of 32



- Mutable: No.
- Constraints: Must be a Block that is finalized.
- withdrawal_nullifier: The account to store the nullifier to ensure that a withdrawal cannot happen twice.
 - · Signer: No.
 - · Init: Yes.
 - PDA: Yes (derived from WITHDRAWAL_NULLIFIER_SEED, claim.block_id, and claim.leaf_index with bump).
 - · Mutable: Yes.
 - Constraints: Must be a new account initialized with space 8 (discriminator).
- from_account: The bridge SPL token account where tokens are taken from.
 - · Signer: No.
 - Init: No.
 - PDA: No.
 - · Mutable: Yes.
 - Constraints: Must be an ATA for the authority account, with the same mint as to_account.
- to_account: The user SPL token account where tokens are sent to.
 - · Signer: No.
 - Init: No.
 - PDA: No.
 - · Mutable: Yes.
 - Constraints: Must be a valid SPL token account owned by claim.user.
- authority: The authority account that owns tokens held by the bridge.
 - · Signer: No.
 - Init: No.
 - PDA: Yes (derived from AUTHORITY_SEED).
 - Mutable: No.
 - · Constraints: None.

- Takes the withdrawal_root from the state_update block provided and verifies that the user's withdrawal exists in the root.
- Transfers claim. amount tokens from from_account to to_account.
- Logs a Withdraw event.

Zellic © 2025 ← Back to Contents Page 29 of 32



CPI

 transfer: Transfers the specified tokens from from_account to to_account, signed by authority.

Module: merkle

This module contains functions for constructing Merkle roots and verifying Merkle paths. It is a custom implementation based on IETF RFC 6962. It creates a skewed Merkle tree. The tree is constructed by first selecting the largest subset of leaves that can form a complete binary Merkle tree and obtaining their root. To compute the overall Merkle root, the root of this subset is treated as a left sibling. The right sibling is computed by repeating this process recursively on the remaining leaves. It is implemented with in-place transforms.

While hashing leaves and intermediate nodes, the preimage is prefixed with LEAF_TAG and PAIR_TAG respectively to prevent second preimage attacks.

Function: proof_path

This function computes a minimal Merkle proof of inclusion of a leaf at index leaf_idx.

Inputs

- leaf_idx: The index of the leaf for which the path is to be derived.
- · leaves: The list of all leaves in the Merkle tree.
- digest_cb: The function encapsulating the logic for hashing a leaf struct.

Function: root_from_proof

This function computes the Merkle root from a proof of inclusion of a leaf at index $leaf_idx$.

Inputs

- leaf: The leaf for which the proof of inclusion is provided.
- leaf_idx: The index of the proven leaf.
- leaves_len: The total number of leaves used to calculate parity for all steps in Merkle proof verification.
- proof: The Merkle proof of inclusion of the leaf.
- digest_cb: The function encapsulating the logic for hashing a leaf struct.

Function: root_from_leaves

This function computes the Merkle root for a given list of leaves.

Zellic © 2025 ← Back to Contents Page 30 of 32



Inputs

- leaves: The list of all leaves in the Merkle tree.
- digest_cb: The function encapsulating the logic for hashing a leaf struct.



Assessment Results

At the time of our assessment, the reviewed code was not deployed to Solana Mainnet.

During our assessment on the scoped N1 Bridge programs, we discovered three findings. No critical issues were found. One finding was of medium impact and the other findings were informational in nature.

6.1. Disclaimer

This assessment does not provide any warranties about finding all possible issues within its scope; in other words, the evaluation results do not guarantee the absence of any subsequent issues. Zellic, of course, also cannot make guarantees about any code added to the project after the version reviewed during our assessment. Furthermore, because a single assessment can never be considered comprehensive, we always recommend multiple independent assessments paired with a bug bounty program.

For each finding, Zellic provides a recommended solution. All code samples in these recommendations are intended to convey how an issue may be resolved (i.e., the idea), but they may not be tested or functional code. These recommendations are not exhaustive, and we encourage our partners to consider them as a starting point for further discussion. We are happy to provide additional guidance and advice as needed.

Finally, the contents of this assessment report are for informational purposes only; do not construe any information in this report as legal, tax, investment, or financial advice. Nothing contained in this report constitutes a solicitation or endorsement of a project by Zellic.

Zellic © 2025 ← Back to Contents Page 32 of 32